

**Gebze Technical University  
Computer Engineering**

**CSE 222 - 2019 Spring**

**HOMEWORK 4 REPORT**

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# 1 QUESTION 1

## 1.1 Part A

```
public static LinkedListO<Integer> q1_iterative(LinkedListO<Integer> L) {
    if(null == L) return null;
    NodeO<Integer> head = L.head;
    if(null == head || null == head.data) return null;
    int maxLen = 0, len = 1, count = 1, maxFirstIndex = 0, firstIndex = 0;
    int last = head.data;
    head = head.next;
    while(null != head) {
        if(last <= head.data) {
            len++;
        }
        else {
            if(len > maxLen) {
                maxLen = len;
                maxFirstIndex = firstIndex;
            }
            firstIndex = count;
            len = 1;
        }
        count++;
        last = head.data;
        head = head.next;
    }
    if(len > maxLen) {
        maxLen = len;
        maxFirstIndex = firstIndex;
    }
    LinkedListO<Integer> S = new LinkedListO<Integer>();
    head = L.head;
    for(int i = 0; i < maxFirstIndex; i++) {
        head = head.next;
    }
    for(int i = 0; i < maxLen; i++) {
        S.add(head.data);
        head = head.next;
    }
    return S;
}
```

There is one while-loop to traverse all the numbers which are  $\theta(n)$  and then there are two for-loops to iterate to the first item of the sublist and one loop to copy sublist. All the for-loop is  $\theta(n)$ . But in the last for-loop, the add method is  $\theta(n)$  for LinkedLists. All the others are constant.

$$T(n) = n^2 + 2n + f(n) \rightarrow \theta(n^2)$$

## 1.2 Part B

Recursive function with LinkedList parameter is not possible. We should give as parameter the head node or use a helper function. I have chosen the second one.

```
public static LinkedListO<Integer> q1_recursive(LinkedListO<Integer> L) {
    if(null == L) return null;
    NodeO<Integer> head = L.head;
    int index = recursion(head, 0, 0, 0);
    LinkedListO<Integer> S = new LinkedListO<Integer>();
    for(int i = 0; i < index; i++) {
        head = head.next;
    }
    int last = head.data;
    while(null != head && last <= head.data) {
        last = head.data;
        S.add(last);
        head = head.next;
    }
    return S;
}

public static int recursion(NodeO<Integer> N, int maxFirstIndex, int maxLen, int index) {
    NodeO<Integer> head = N;
    if(null == head || null == head.data) return -1;
    int len = 1, count = 1, last = head.data;
    head = head.next;
    while(null != head) {
        if(last <= head.data) {
            len++;
        }
        else {
            if(len > maxLen) {
                maxLen = len;
                maxFirstIndex = index;
            }
            maxFirstIndex = recursion(head, maxFirstIndex, maxLen, index + count);
            break;
        }
        count++;
        last = head.data;
        head = head.next;
    }
    if(len > maxLen) {
        maxFirstIndex += index;
    }
    return maxFirstIndex;
}
```

For *q1\_recursive* method, the complexity is Trecursion(n) and two separate loops which are  $\theta(n)$  and  $\theta(n^2)$  so:

$$T(n) = Trecursion(n) + n^2 + n + \text{constants}$$

For *recursion* method,  $a \Rightarrow 1$  and  $b > 1$ ,  $T(n) = aT(n/b) + f(n)$

$n$  = size of numbers

$a$  = sorted sublists number

$n/a$  = size of each subproblem. (To apply Master Theorem, it is assumed that all subproblems are essentially the same size.  $a = b$ )

$f(n)$  = the cost of dividing the problem and merging solutions. ( $n/a$ )

$c = T(f(n)) = n/a$

$$T(n) = aT(n/a) + n/a$$

$a$  is not constant number. Master Theorem cannot be apply to this equation. But if we try:

$$T(n) = aT(n/a) + n/a \quad \log_b(a) = \log_a(a) \quad \log_b(a) = 1$$

$$f(n) = \theta(n^c * \log^k(n)) \quad \rightarrow \quad T(n) = \theta(n^c * \log^{(k+1)}(n))$$

$$c = 1, k = 0$$

$$c = \log_b(a)$$

$$T(n) = \theta(n * \log(n))$$

If we use induction,  $T(n) = aT(n/a) + n/a$ , and we guess  $T(n) \leq cn \log n$

$$T(n) \leq a * (c * (n/a) * \log(n/a)) + n/a$$

$$T(n) \leq cn * \log n - cn * \log a + n/a$$

$$T(n) \leq cn * \log n$$

$T(n)$  for *q1\_recursive* method,  $T(n) = T_{recursion}(n) + n^2 + n + \text{constants}$

$$T(n) = n \log n + n^2 + n + \text{constants} \quad \rightarrow \quad \theta(n^2)$$

## 2 QUESTION 2

```
int[] arr = DEFINED;
```

```
int x = DEFINED;
```

```
int idxL = 0, idxR = arr.length;
```

```
while(idxR >= idxL) {  
    int sum = arr[idxL] + arr[idxR];  
    if(sum > x) idxR--;  
    else if(sum < x) idxL++;  
    else break;  
}
```

**idxL and idxR are the special indexes.  $T(n) = n \rightarrow \theta(n)$**

## 3 QUESTION 3

This code has three inner for-loops.

$$T(n) = (2 * n) * n * (n / 3) = 2 * n * n * n / 3$$

$$T(n) = n^3 / 3 \text{ and } T(n) = n / 3 = O(\log_3(n))$$

$$O(n^2 * \log(n))$$