

i)

CSCI 404 Written HW #1

9)

```
void f(int n)
{
    int i = 2;
    while (i < n) {
        /* something that take O(1) time */
        i = i * i;
    }
}
```

Iteration 0	1	2	3	
$i = 2$	4	16	$16^2 = 256$	$i_k = 2^{2^k}$

• loop ends when $i_k \geq n$

$$\downarrow$$

$$2^{2^k} \geq n$$

$$\downarrow$$

$$2^k \geq \log_2 n$$

$$\downarrow$$

$$k \geq \log_2 (\log_2 n)$$

$$\downarrow$$

$$k \geq \log \log n$$

The function runs in $O(\log \log n)$ runtime complexity.

b.

```
void f2(int n)
{
```

```
  for(int i=1; i<=n; i++) { ← O(n) complexity
```

```
    if((i%(int)sqrt(n))==0) { ← happens  $\sqrt{n}$  times
```

```
      for(int k=0; k<pow(i,3); k++) { ← when true, runs  $i^3$  times
```

```
        /* something that takes O(k) time */
```

```
      }
```

```
    }
```

```
  }
```

```
}
```

if condition passes for \sqrt{n} multiples of \sqrt{n}

$i = j\sqrt{n}$ for $j=1, 2, \dots, k$

$\text{pow}(j\sqrt{n}, 3) = (j\sqrt{n})^3 = j^3 n^{3/2}$

$O(n^{3/2})$

complexity: $O(n^{5/2})$

c. for(int i=1; i<=n; i++) { ← O(n) complexity

```
  for(int k=1; k<=n; k++) { ← O(n) complexity
```

```
    if(A[k]==i) { ← worst case check every value, O(n) complexity
```

```
      for(int m=1; m<=n; m=m+m) {
```

```
        // something takes O(1) time //
```

```
        // contents of A[] are not changed //
```

```
      }
```

```
    }
```

```
  }
```

```
}
```

$m=1, 2, 4, 8, \dots$

$m=2^k$ where $k=0, 1, 2, \dots$

$2^k > n$

$k > \log_2 n \rightarrow k > \log n$

complexity: $O(n) \cdot O(n) \cdot \log n$
 $= O(n^2 \log n)$

d) int f(int n)
{

int *a = new int[10];

← O(1) complexity

int size = 10;

← O(1) complexity

for (int i = 0; i < n; i++) {

← O(n) complexity

if (i == size)

← has O(1) worst case: i = size every time

{

int newSize = 3 * size / 2;

int *b = new int[newSize];

for (int j = 0; j < size; j++) b[j] = a[j];

delete a;

a = b;

size = newSize;

}

a[i] = i * i;

}

}

resizes array

to $3/2$ of original
size

size = 10 → 15 → 22 → 33 → ...

size = $10 * 3/2^k \Rightarrow O(n)$

loop

total complexity: O(n)

Problem 2

```
struct Node {  
    int val;  
    Node* next;  
};
```

```
Node* llrec(Node* in1, Node* in2)  
{  
    if (in1 == nullptr) {  
        return in2;  
    }  
    else if (in2 == nullptr) {  
        return in1;  
    }  
    else {  
        in1->next = llrec(in2, in1->next);  
        return in1;  
    }  
}
```

a) in1 = 1, 2, 3, 4 in2 = 5, 6

1 5 2
in1->next = llrec(in2, in1->next)

6 2 1 4
in1->next = llrec(in2, in1->next)

3 6 2 1
in1->next = llrec(in2, in1->next)

↓ nullptr
6->next
in1->next = llrec(3, nullptr)
6->next = 3;

a. $in1 = 1, 2, 3, 4$ $in2 = 5, 6$

$1 \rightarrow 2 = (5 \rightarrow 6, 2 \rightarrow 3 \rightarrow 4);$

$6 \rightarrow nullptr = (2 \rightarrow 3 \rightarrow 4, 6 \rightarrow nullptr)$

$3 \rightarrow 4 = (6 \rightarrow nullptr, 3 \rightarrow 4)$

$nullptr = (3 \rightarrow 4, nullptr)$

$nullptr = 3 \rightarrow 4;$

↓

$6 \rightarrow 3 \rightarrow 4$

↓

$2 \rightarrow 6 \rightarrow 3 \rightarrow 4$

↓

$5 \rightarrow 2 \rightarrow 6 \rightarrow 3 \rightarrow 4$

↓

output: $1 \rightarrow 5 \rightarrow 2 \rightarrow 6 \rightarrow 3 \rightarrow 4$

b) input:

$in1 = nullptr$ $in2 = 2$

output: 2