

CSc 360

Operating Systems

Processes

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P1 is already out!

- Due Monday, Oct 2, thru Brightspace
 - please make sure to code your own assignment
 - if you use any helper function, do cite it clearly
 - don't waste your time/money to copy&paste
 - we have access to github and old files too
- Please attend tutorials to follow its schedule
- A simple shell interpreter (SSI), able to
 - execute external programs (e.g., `ls -l`)
 - change directories (i.e., internal commands)
 - execute programs in background

Too challenging?

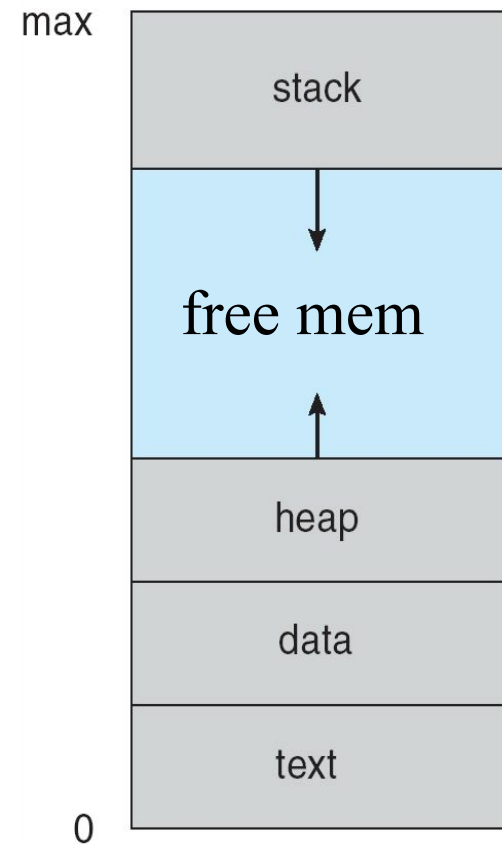
- We are here to help you
 - follow the suggested approach
 - discussed in the first lecture---it's effective!
 - attend lectures and tutorials: both!
 - get started earlier!
 - P1 discussion thread on Teams
 - get help and help others: TA standby in a day
 - *CSC consultant clinic/office: ECS 2nd floor*
 - office hours (tutorial and lecture instructors)

Responsible use of computers

- Through this course, we will know better about operating systems, how they work, and tricks and tips
- You can practice these things and skills on your own computers **with a VM**
- Do not attempt to trick or compromise computers ***also used by others***
- See UVic policies: IT Policy 6030

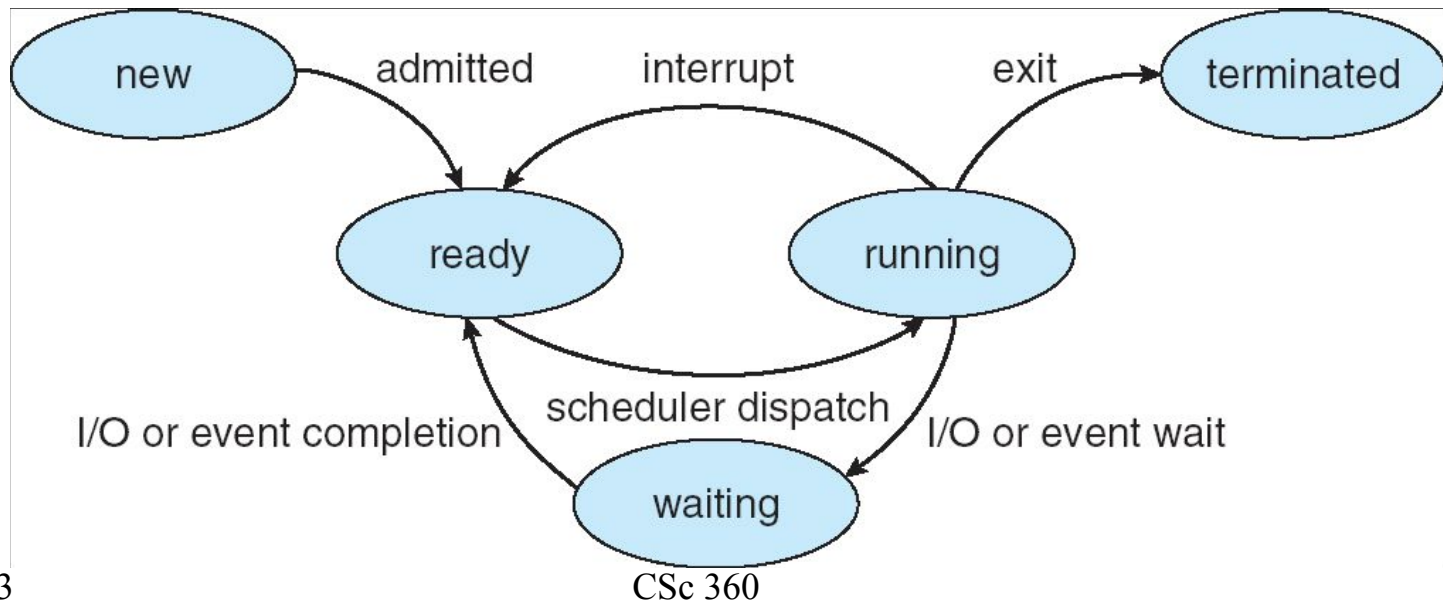
Processes

- Process: a program in execution
- Program: passive entity
 - static binary file on storage
 - e.g., `gcc -o hello hello.c; ls -l hello`
 - `-rwxrwxr-x 1 user group size date/time hello`
- Process: active entity; resource allocated!
 - `./hello`
 - text (code); data (static), stack, heap
 - process control block (PCB)



Process states

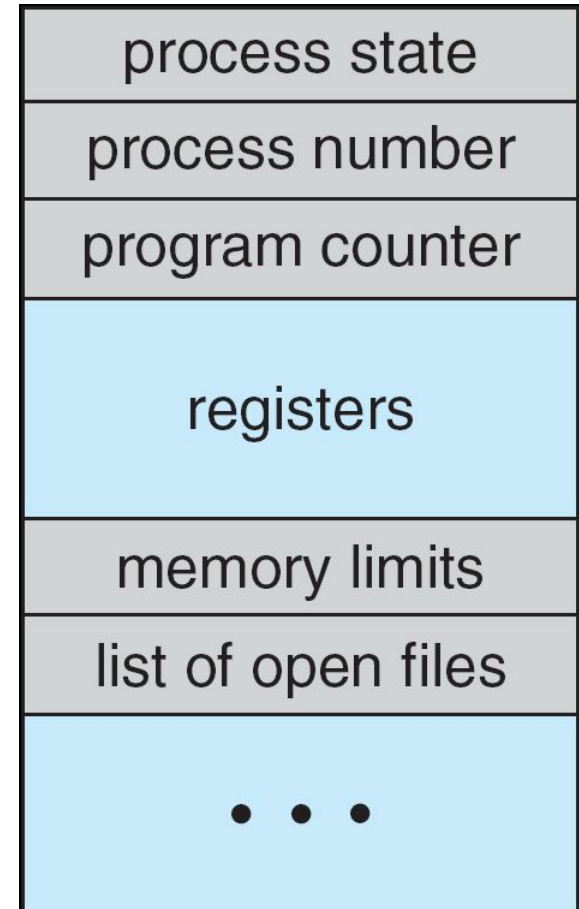
- E.g., one CPU (core)
 - one running process at any time
 - maybe many ready/waiting processes



why interrupt?

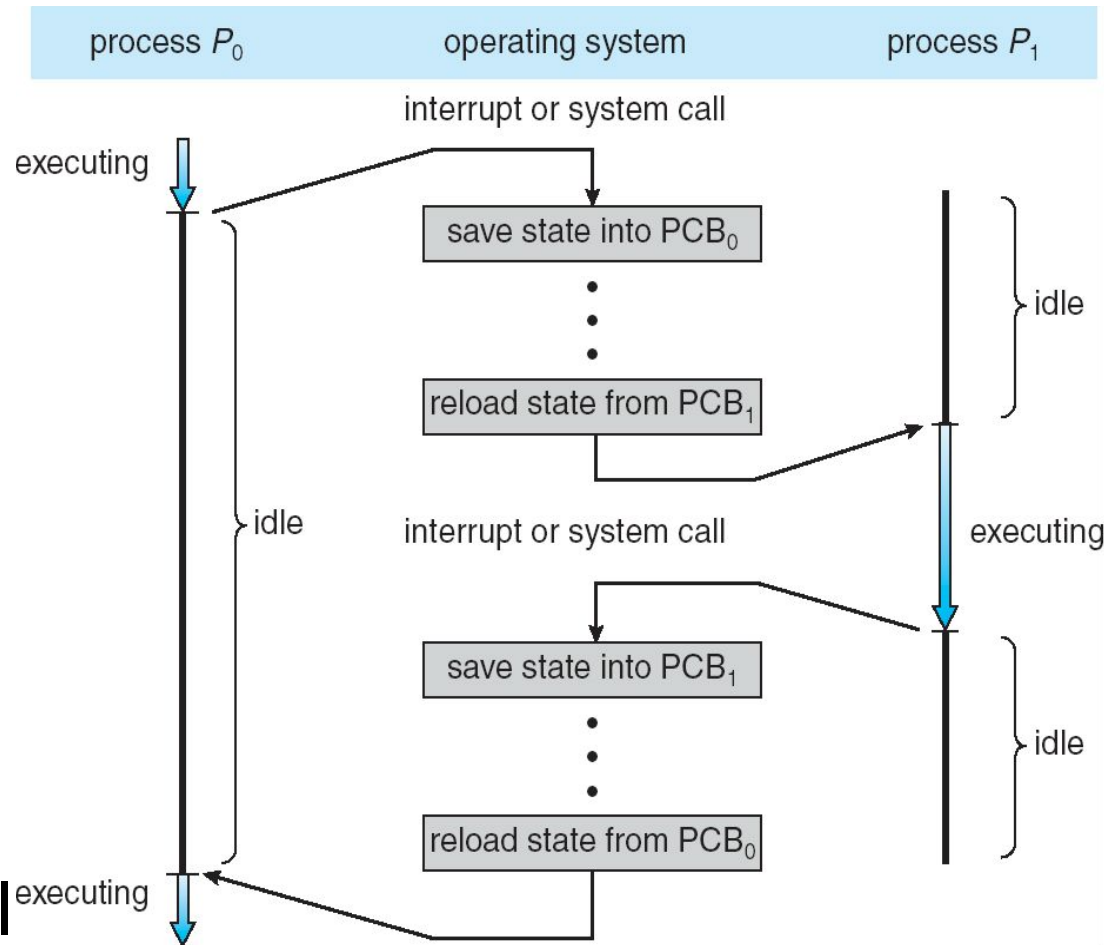
Process control blocks

- PCB: keep track processes
 - state: ready/running, etc
 - CPU
 - PC, registers, priority, etc
 - memory
 - memory control information
 - I/O
 - e.g., list of opened files
 - accounting



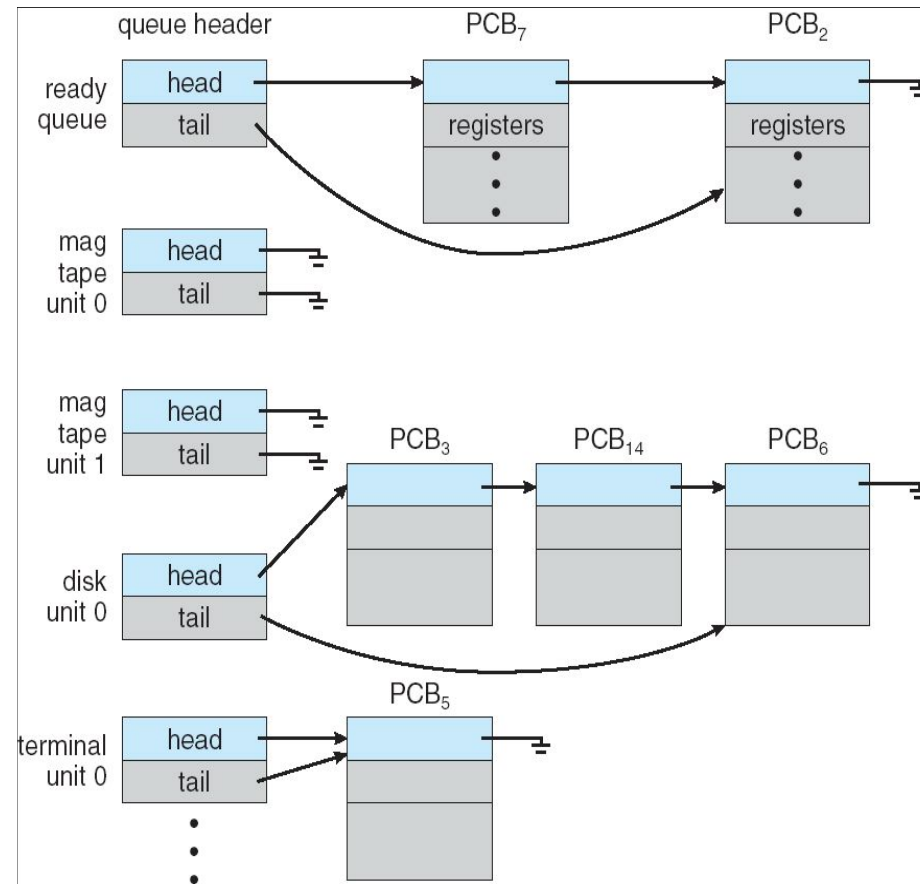
Context switching

- Context switch
 - save states
 - restore states
- When
 - timer
 - I/O, memory
 - trap
 - waiting sys call

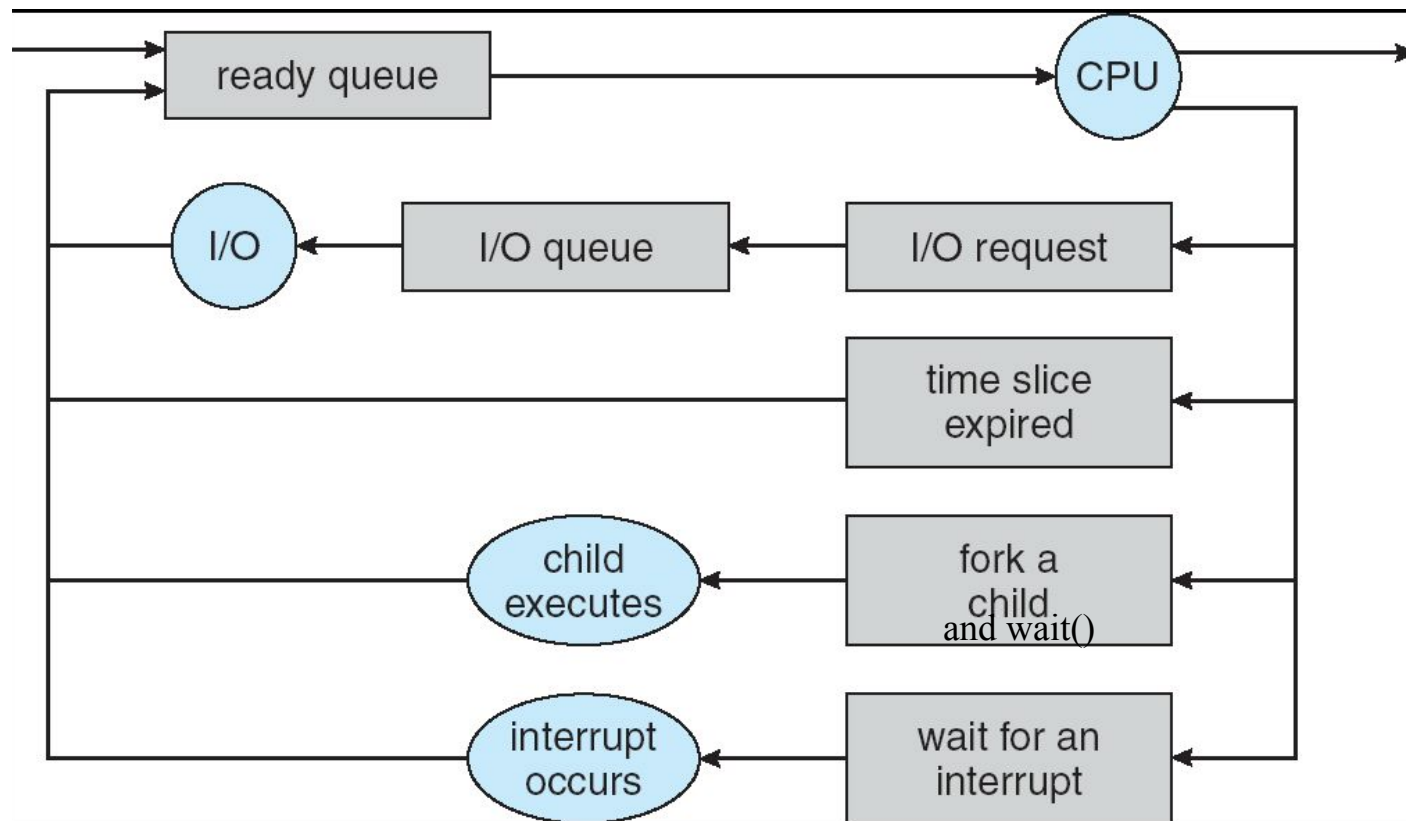


Process scheduling

- Multiprogramming
 - utilization
- Timesharing
 - interactive
- Scheduling queues
 - linked list
 - ready queue
 - I/O queue



Queuing system



Queuing scheduler

- Who's the next?
- Long-term scheduler
 - job scheduler (spooling)
 - get to the ready queue
 - CPU-intensive vs I/O intensive
- Short-term scheduler
 - CPU scheduler
 - frequency vs overhead

where's the long-term scheduler/gatekeeper?

More on scheduling

- Medium-term scheduler
 - who is NOT the next
 - reduce the degree of multiprogramming
 - swap-in/out
- Scheduling algorithms
 - first-come-first-server, shortest-job-first, priority, round-robin, fair and weighted fair, ...
 - more in Chapter 5

This lecture so far

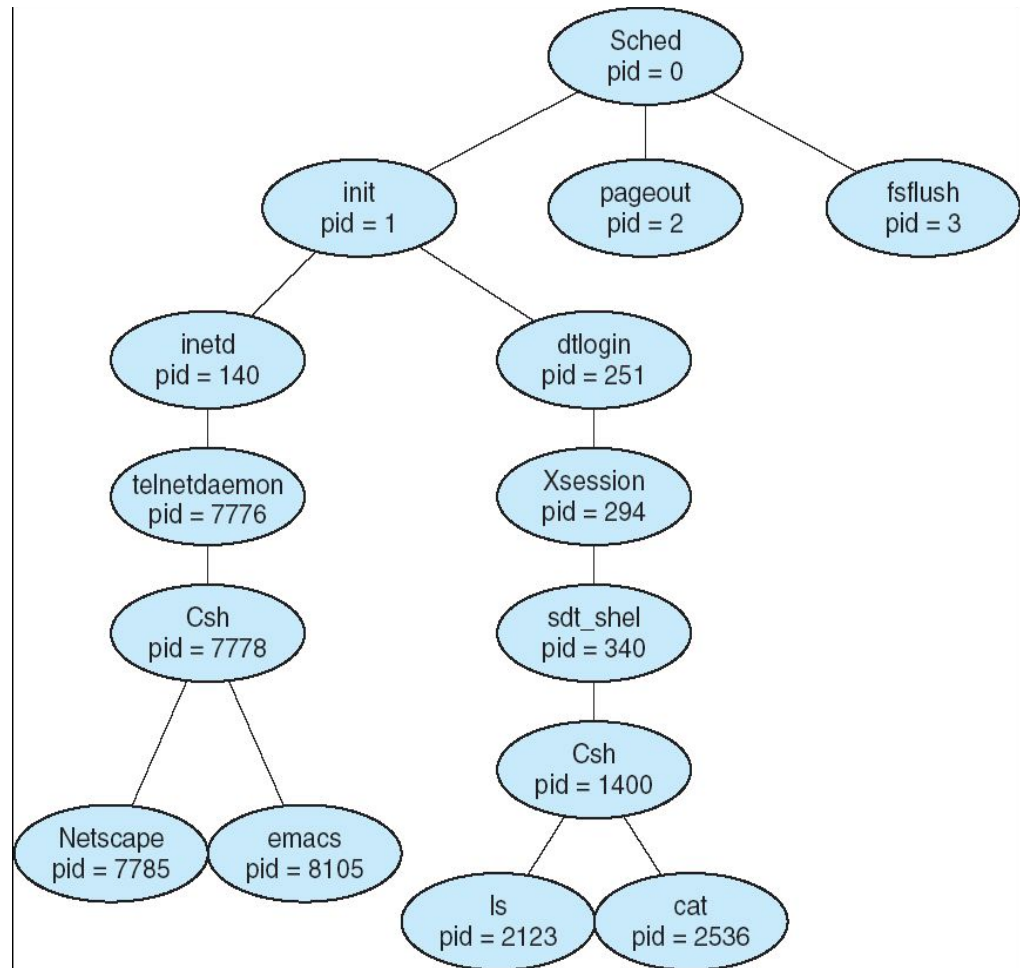
- Process and process scheduling
 - process vs program
 - process control block
 - context switch: what to save/restore
 - process scheduling
- Explore further
 - process status: `/bin/ps`
 - top CPU processes: `/usr/bin/top`

Process creation

- Creating processes
 - parent process: create child processes
 - child process: created by its parent process
- Process tree
 - recursive parent-child relationship; why tree?
 - `/usr/bin/pstree`
- Process ID (PID) and Parent PID (PPID)
 - usually nonnegative integer

Process tree

- sched (0)
 - init (1)
 - all user processes
 - pageout
 - memory
 - fsflush
 - file system



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pstree on linux.csc.uvic.ca

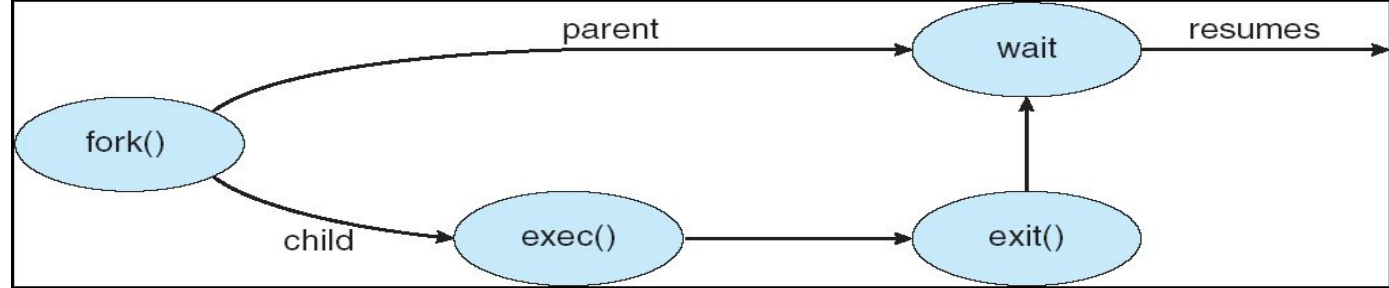
```
snmpd
sshd
  5*[sshd—sshd—sftp-server]
  sshd—sshd—bash—ssi
  sshd—sshd—bash—more
  sshd—sshd—bash—pstree
  sshd—sshd—csh—sftp-server
  sshd—sshd—bash
  sshd—sshd—bash—mutt
  sshd—sshd—tcsh—nano
  sshd—sshd—tcsh
syslogd
```

Parent vs child processes

- Process: running program + resources
- Resource sharing: possible approaches
 - all shared, or
 - some shared (e.g., read-only code), or
 - nothing shared*
- Process execution: possible approaches
 - parent waits until child finishes, or
 - parent and child run **concurrently***

fork(), exec*(), wait()

- Create a child process: fork()
 - return code < 0: error (in “parent” process)
 - return code = 0: you’re in child process
 - return code > 0: you’re in parent process
 - return code = child’s PID
- Child process: load a new program
 - exec*(): front-end for execve(file, arg, environ)
- Parent process: wait() and waitpid()
 - * details during tutorials



Example

```
int main()
{
    Pid_t pid;
    /* fork another process */
    pid = fork();
    if (pid < 0) { /* error occurred */
        fprintf(stderr, "Fork Failed");
        exit(-1);
    }
    else if (pid == 0) { /* child process */
        execvp("/bin/ls", "ls", NULL);
    }
    else { /* parent process */
        /*parent will wait for the child to complete*/
        wait (NULL);
        printf ("Child Complete");
        exit(0);
    }
}
```

Process termination

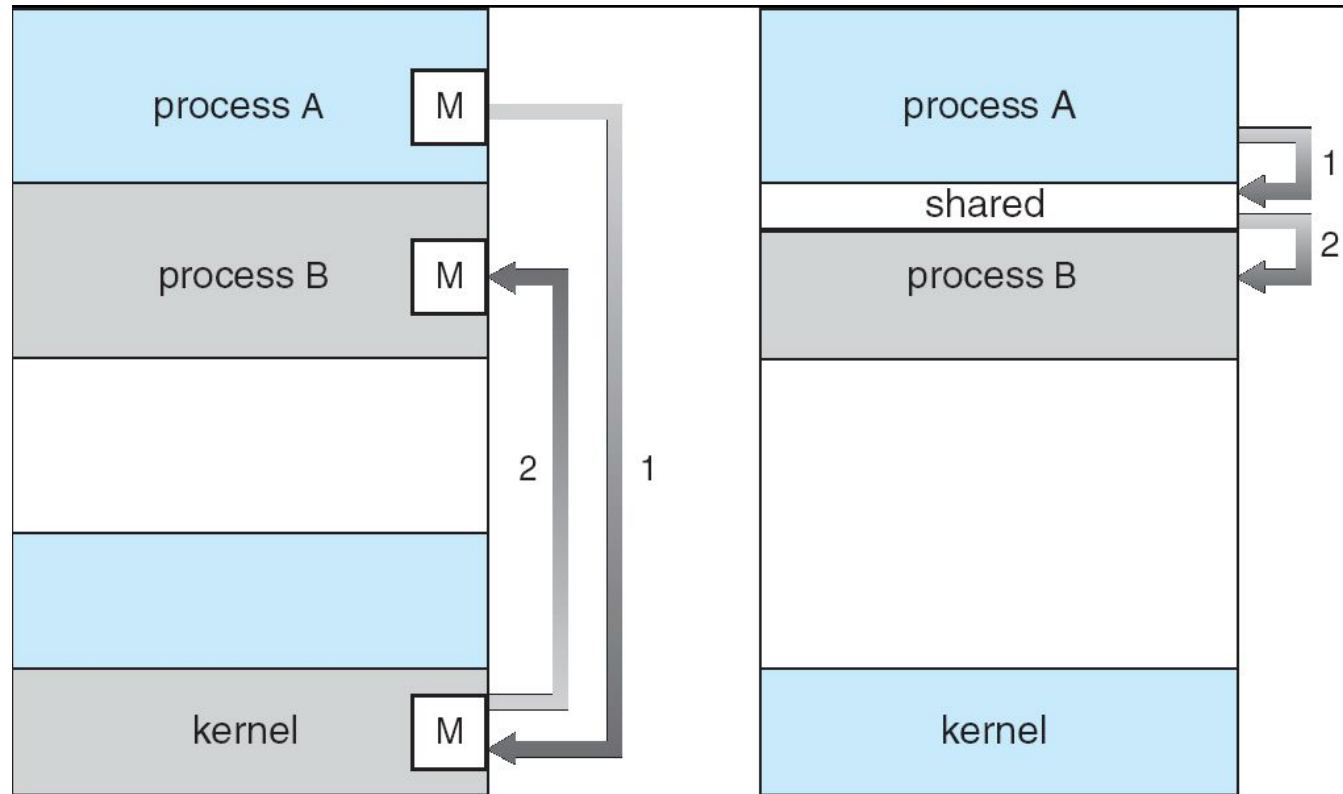
- Terminate itself: `exit()`
 - report status to parent process
 - release allocated resources
- Terminate child processes: `kill(pid, signal)`
 - actually send a signal to the child
 - child resource exceeded, child process no longer needed, and so on
 - parent is exiting
 - cascading termination, or find another parent

Process communication

- Independent process
 - standalone process
- Cooperating process
 - affected by or affecting other processes
 - sharing, parallel, modularity, convenience
- Process communication
 - shared memory
 - message passing

Message passing vs shared memory

- Overhead
- Protection



(a)

(b)

The 2nd half of this lecture

- Process operations
 - process creation
 - process tree
 - process termination
 - the need for inter-process communication
- Explore further
 - `/bin/ps`, `/usr/bin/top`, `/usr/bin/pstree`
 - how does a child process find its parent's PID?

Next lecture

- Inter-process communication
 - read OSC7 Chapter 3 (or OSC6 Chapter 4)