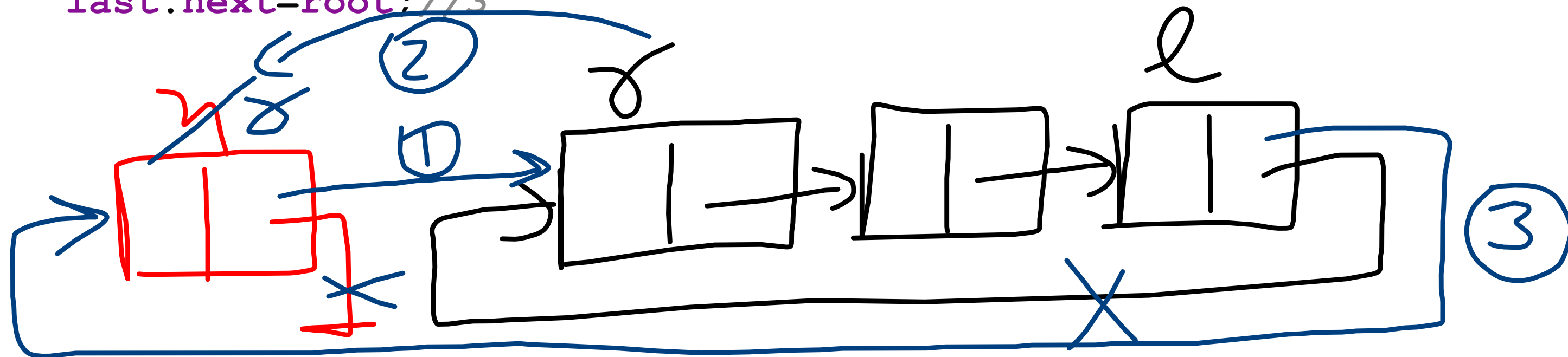
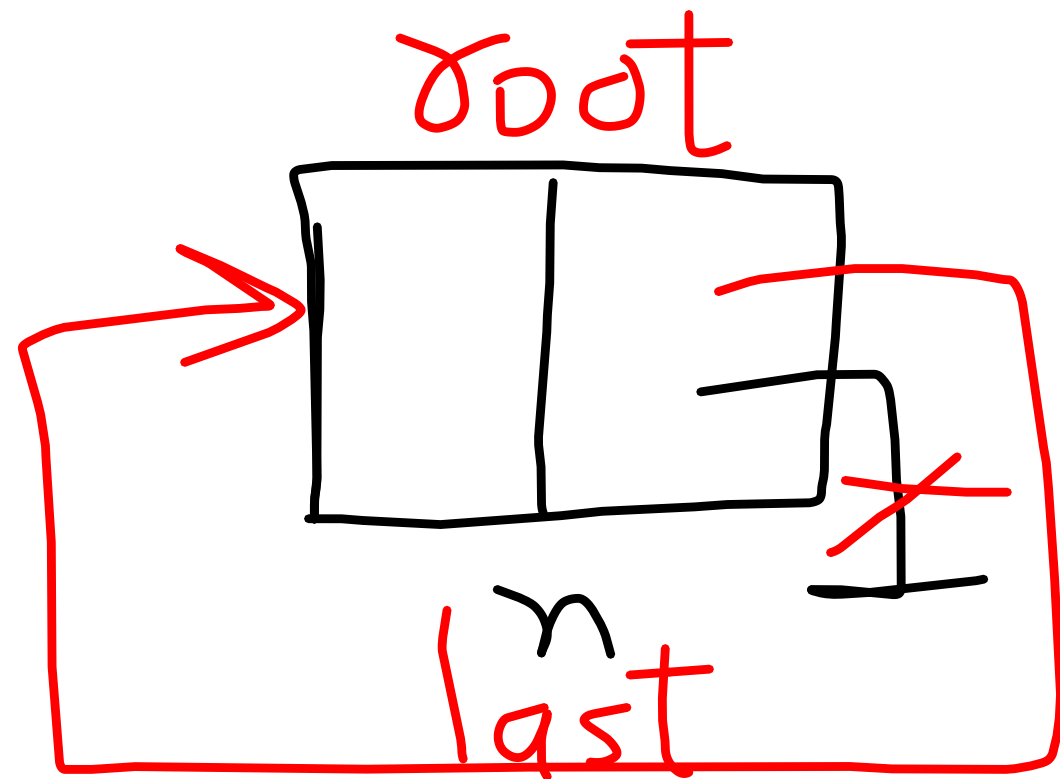


```

void insert_left(int data) {
    Node n = new Node(data); // created node
    if (root == null) // no root
    {
        root = last = n; // 1st becomes root
        last.next = root;
    }
    else {
        n.next = root; // 1 ✓
        root = n; // 2 ✓
        last.next = root; // 3
    }
}

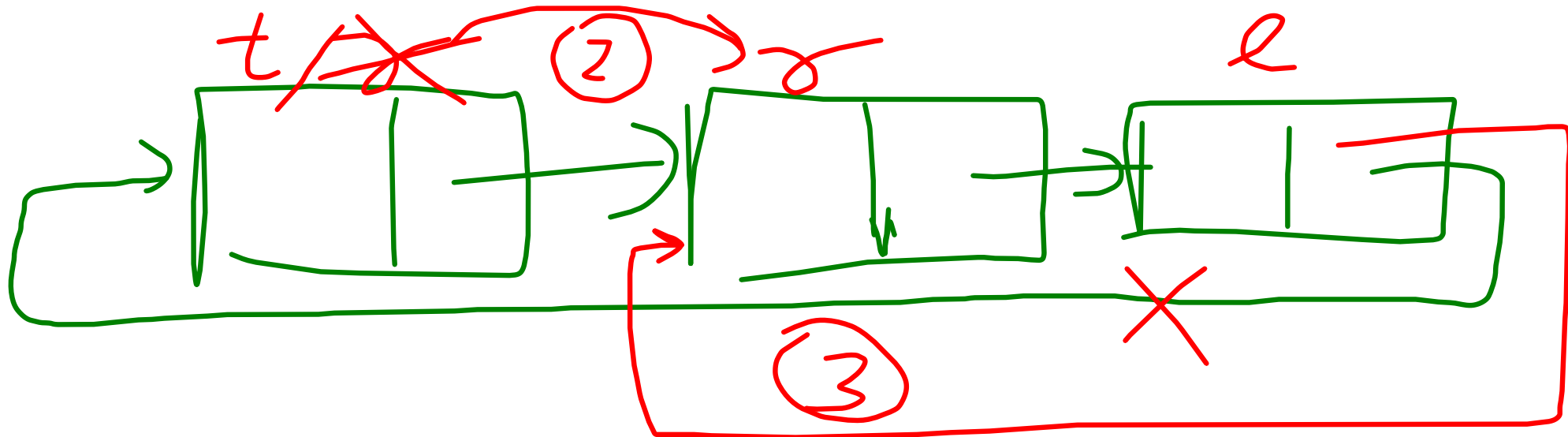
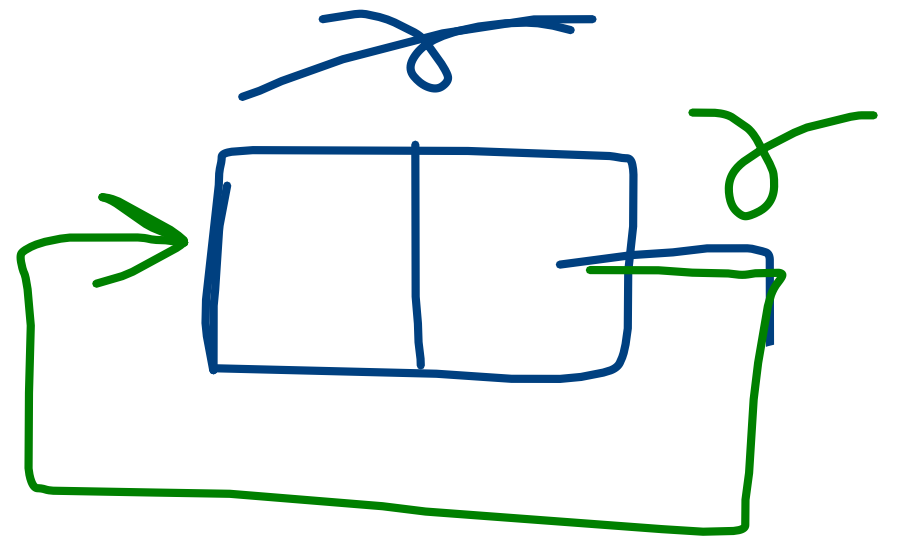
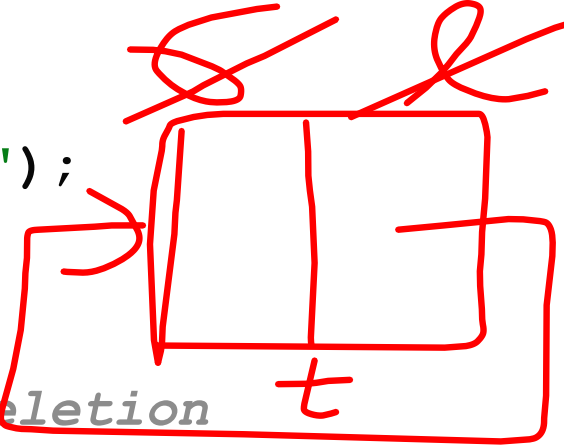
```




```

void delete_left() {
    if (root == null) //no root
        System.out.println("List is empty");
    else {
        Node t = root; //1
        if (root == last) //single node
        root = last = null; //manual deletion
        else {
            root = root.next; //2
            last.next = root; //3
        }
        System.out.println("Deleted:" + t.data); //3 response message of
deletion
    }
}

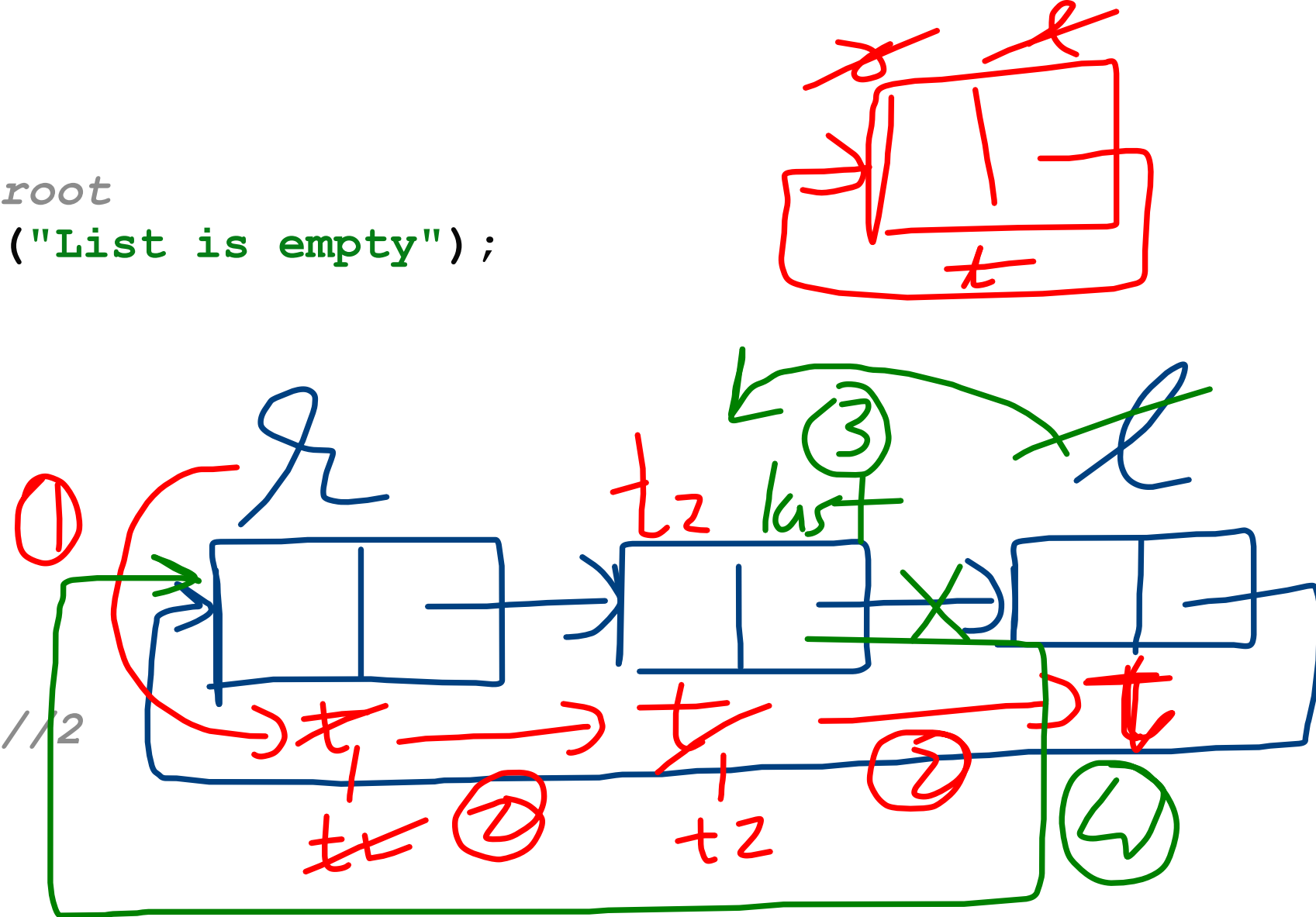
```



```

void delete_right() {
    if (root == null) //no root
        System.out.println("List is empty");
    else {
        Node t, t2;
        t = t2 = root; //1
        if (root == last)
            root = last = null;
        else
        {
            while (t != last) //2
            {
                t2 = t;
                t = t.next;
            }
            last = t2; //3
            last.next = root; //4
        }
        System.out.println("Deleted:" + t.data); //3 response
        message of deletion
            (t.next).data
    }
}

```

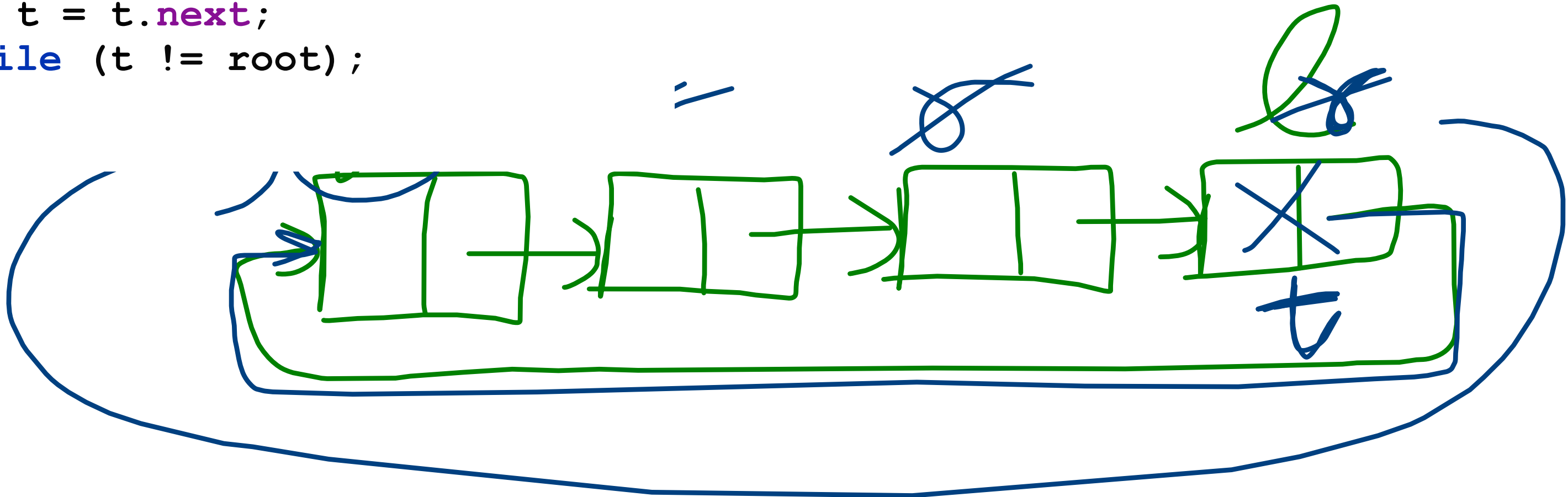


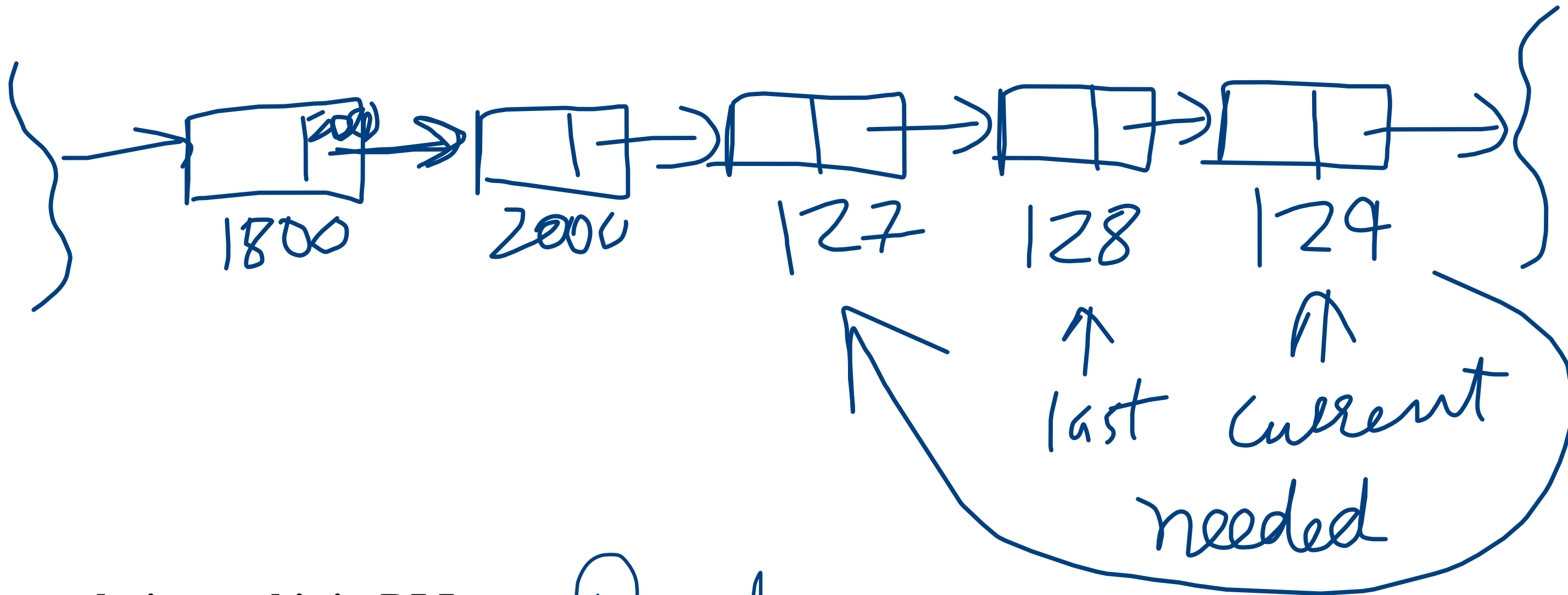
```

void print_list() {
    if (root == null) //no root
        System.out.println("List is
empty");
    else {
        Node t = root; //1
        do{
            System.out.print("| " +
t.data + " |->");
            t = t.next;
        }while (t != root);
    }
}

```

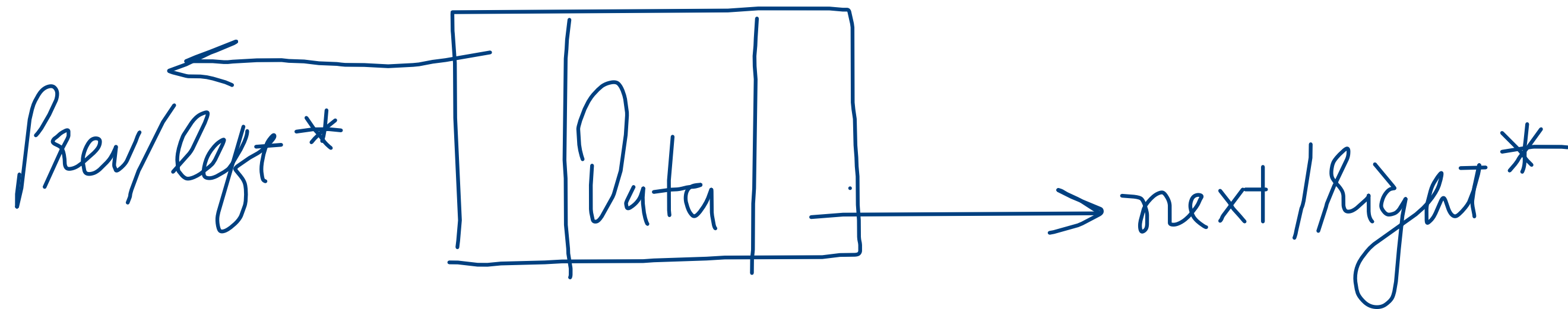
start from root print
till root not reached

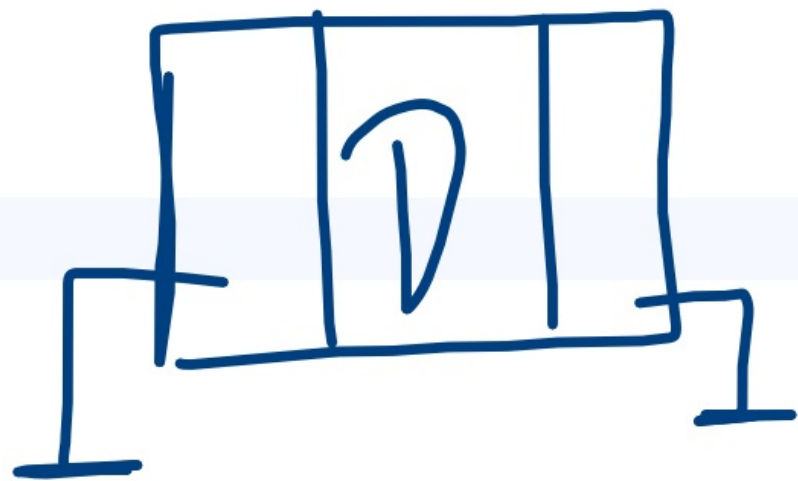


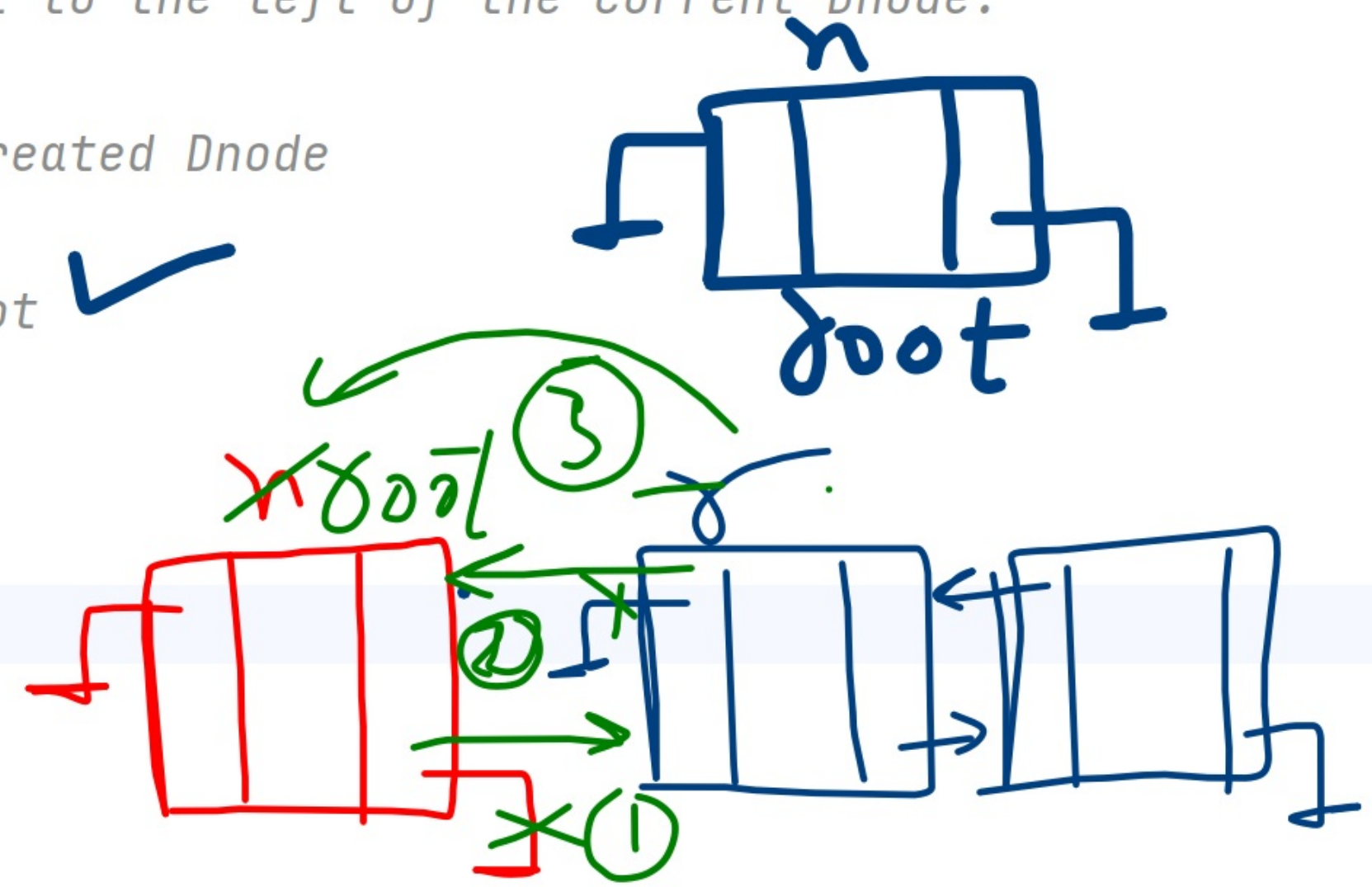


solution to this is **DLL**

Node



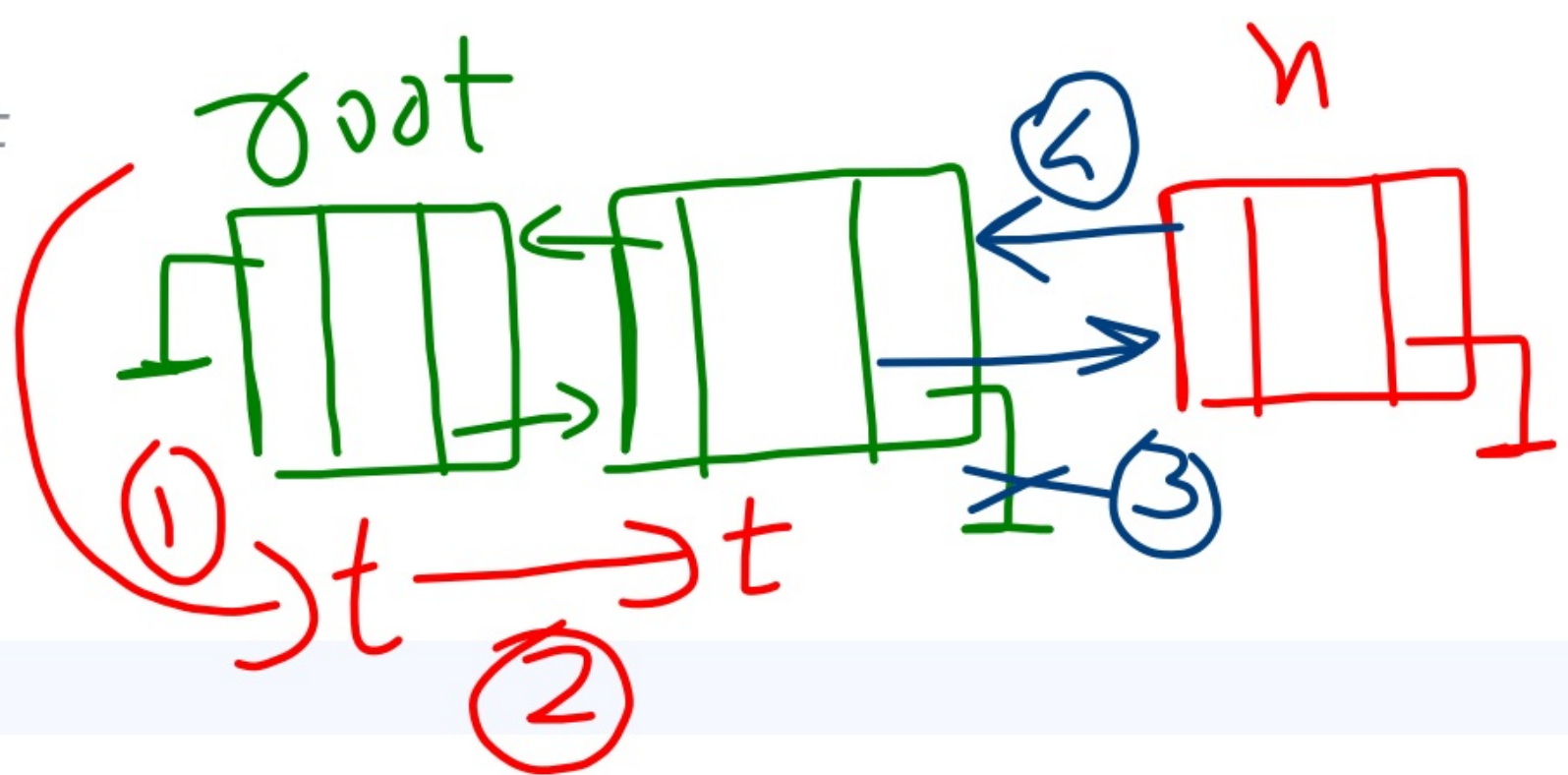




```

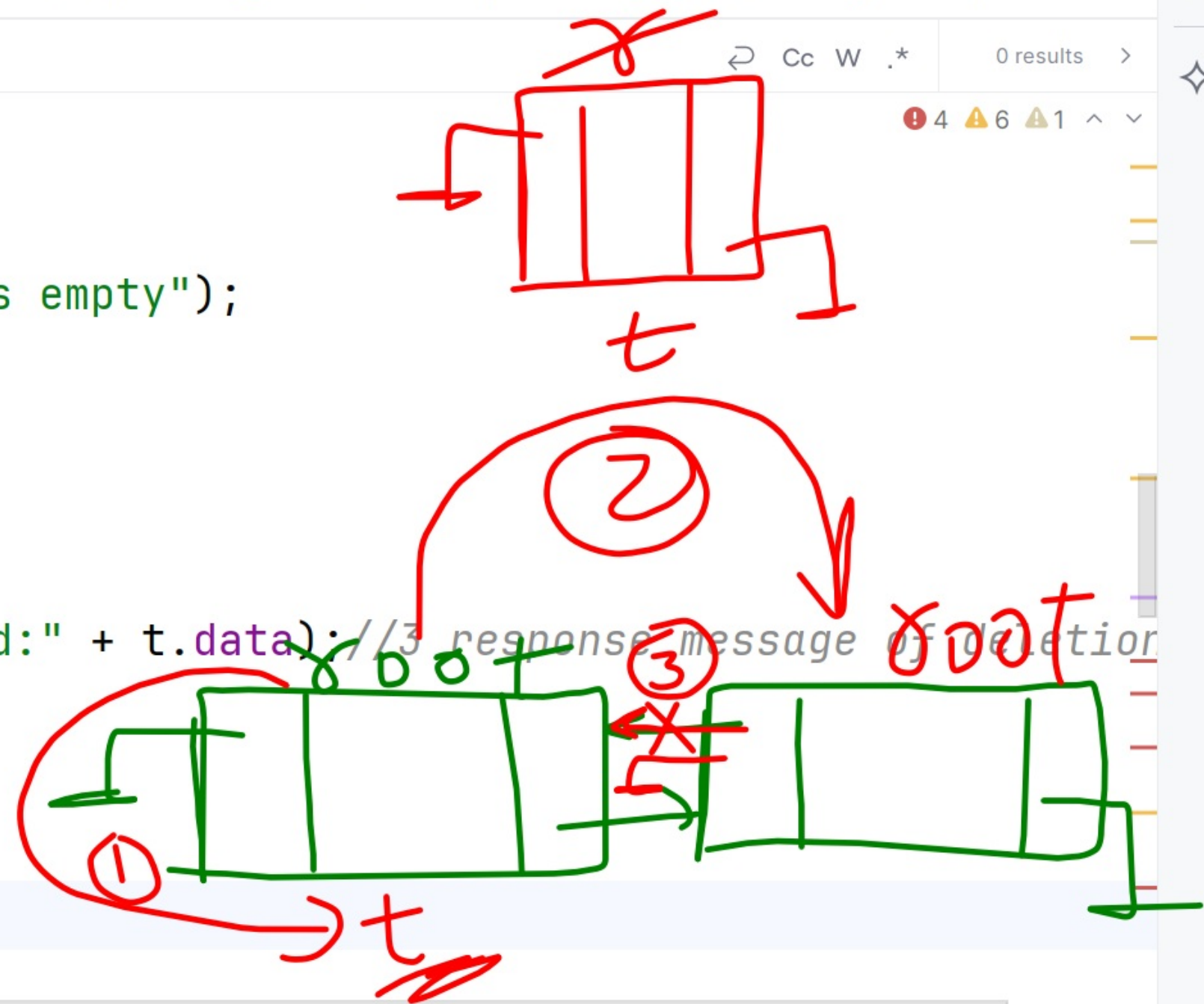
19 void insert_right(int data) {
20     Dnode n = new Dnode(data); // created Dnode
21     if (root == null) // no root
22         root = n; // 1st becomes root
23     else {
24         Dnode t = root; // 1
25         while (t.right != null) // 2
26             t = t.right;
27         t.right = n; // 3 connected
28         n.left = t; // 4
29     }
30 }
31
32 void delete_left() {

```

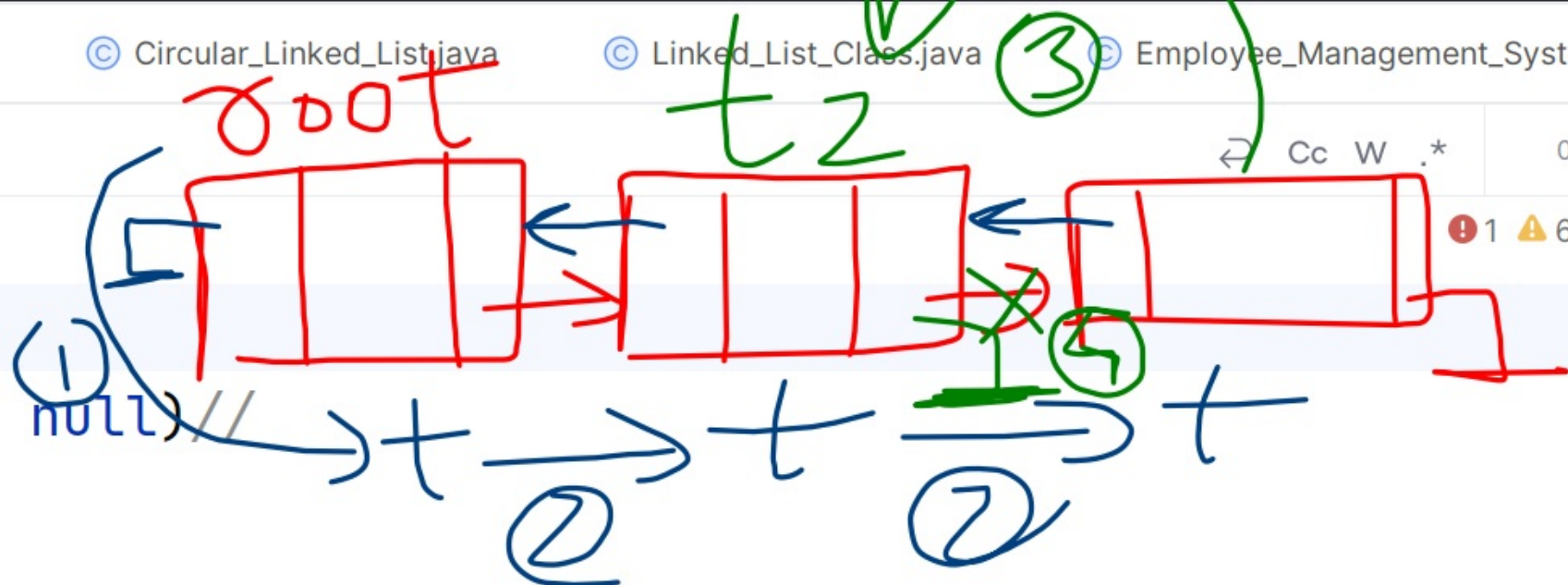



```
void delete_left() {  
    if (root == null) //no root  
        System.out.println("List is empty");  
    else {  
        Dnode t = root; //1  
        root = root.right; //2  
        root.left = null; //3  
        System.out.println("Deleted:" + t.data); //3 response message of deletion  
    }  
}
```

```
void delete_right() {  
    if (root == null) //no root
```




```
System.out.println("Deleted:" + t.data); //3 response message of deletion
```



FileEditViewNavigateCodeRefactorBuildRunToolsVCSWindowHelp

Current File

Doubly_Linked_List.javaNode.javaDnode.javaCircular_Linked_List.javaLinked_List_Class.javaEmployee_Management_System

void print_list_rev() {
if (root == null) // no root
System.out.println("List is empty");
else {
Dnode t = root; // 1
while (t.right != null)
t = t.right;
while (t != null)
{
System.out.print("<-| " + t.data + "|>");
t = t.left;
}
}

①

10 ↔ 20 ↔ 30

②

t → t → t

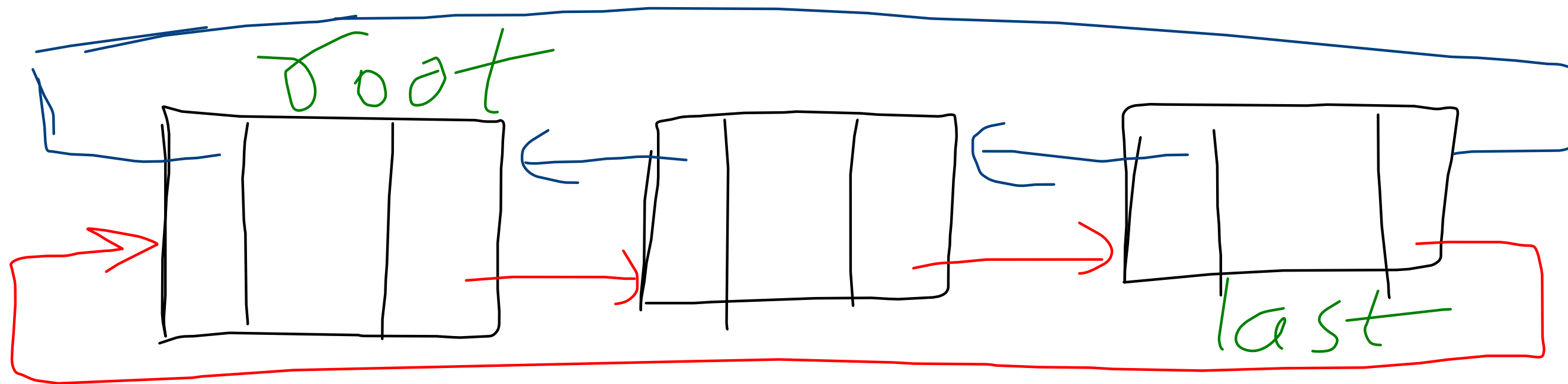
③

t ← t ← t

// stop

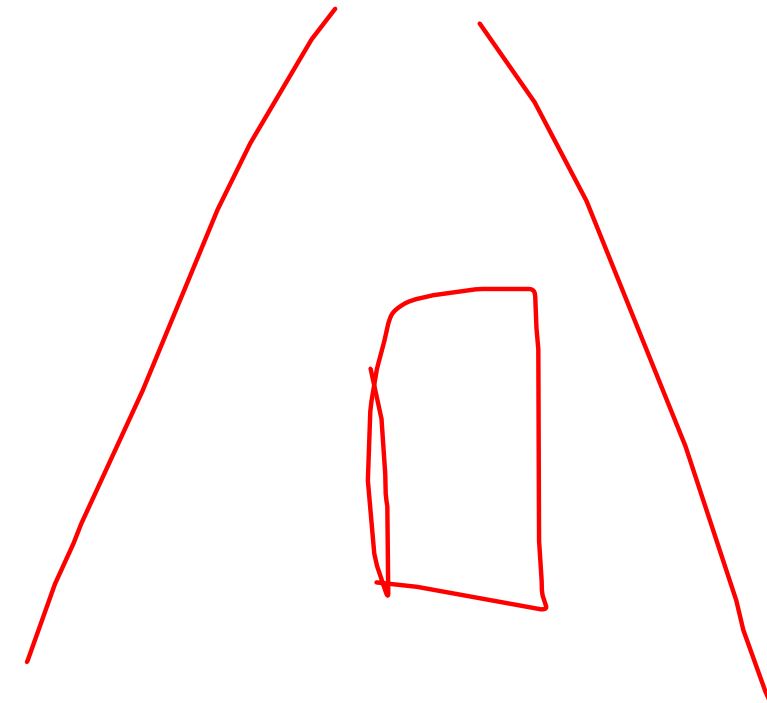
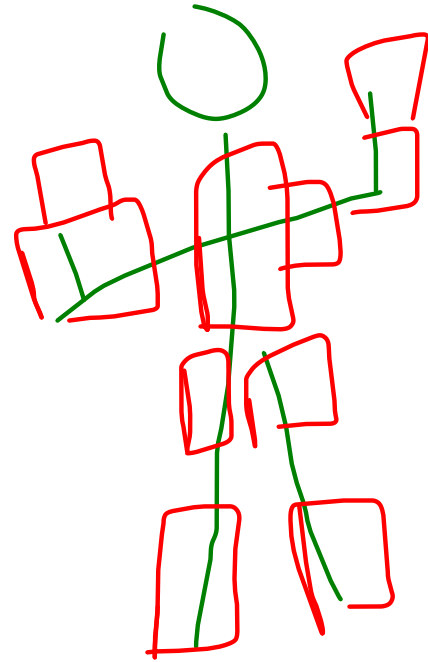
30 -> 20 -> 10

Data_Structures_VITA > src > Linked_List_Examples > Doubly_Linked_List > print_list_rev87:14CRLFUTF-84 spaces10:0914-11-2025



DCLL

Application



inorder:LPR

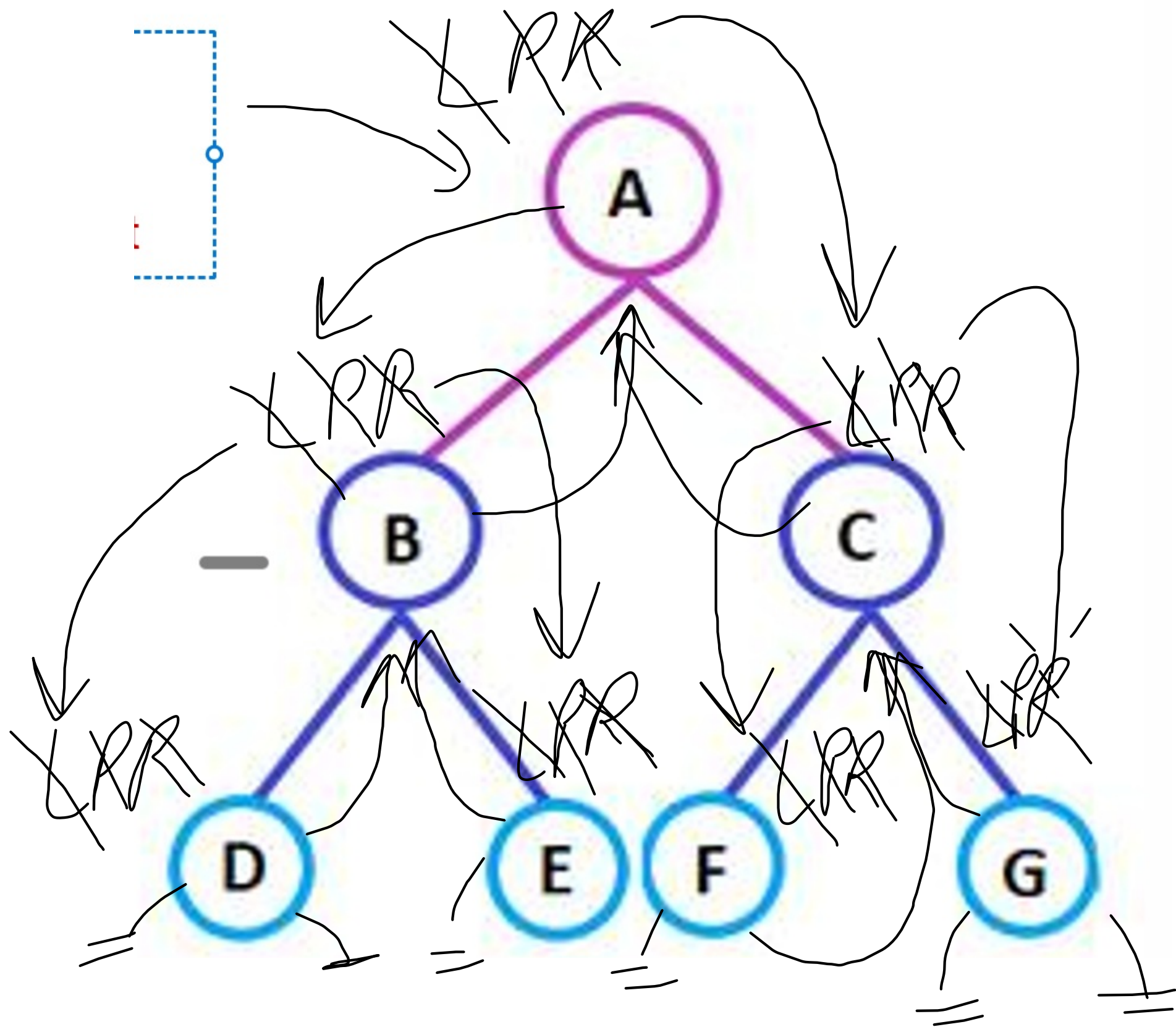
when u cut L:go to left

when u cut R:go to right

when u cut P:print

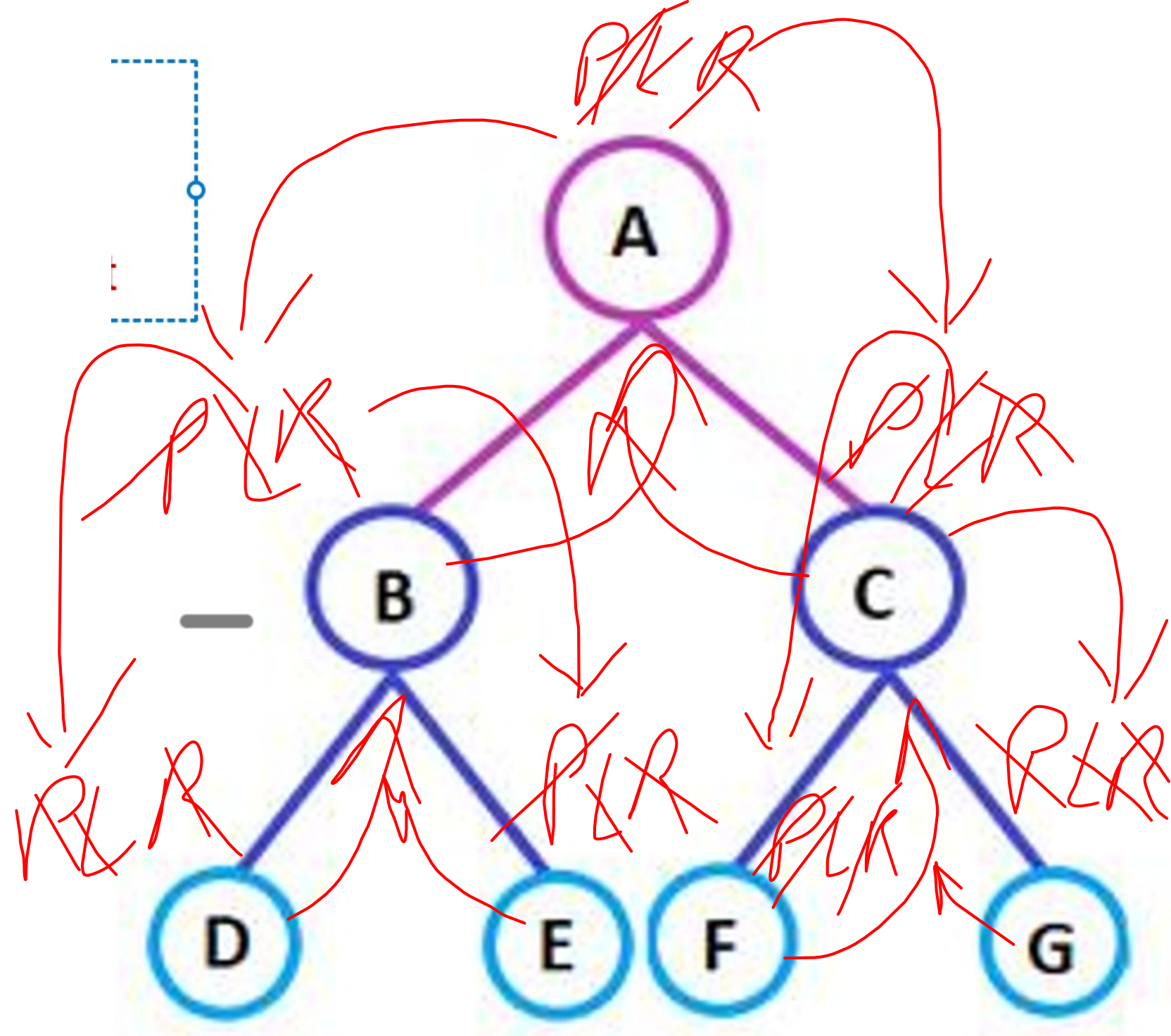
when all done go to parent

D, B, E, A, F, C, G,



when all done go to parent

~~AB~~ D E C F G



Post-order:LRP

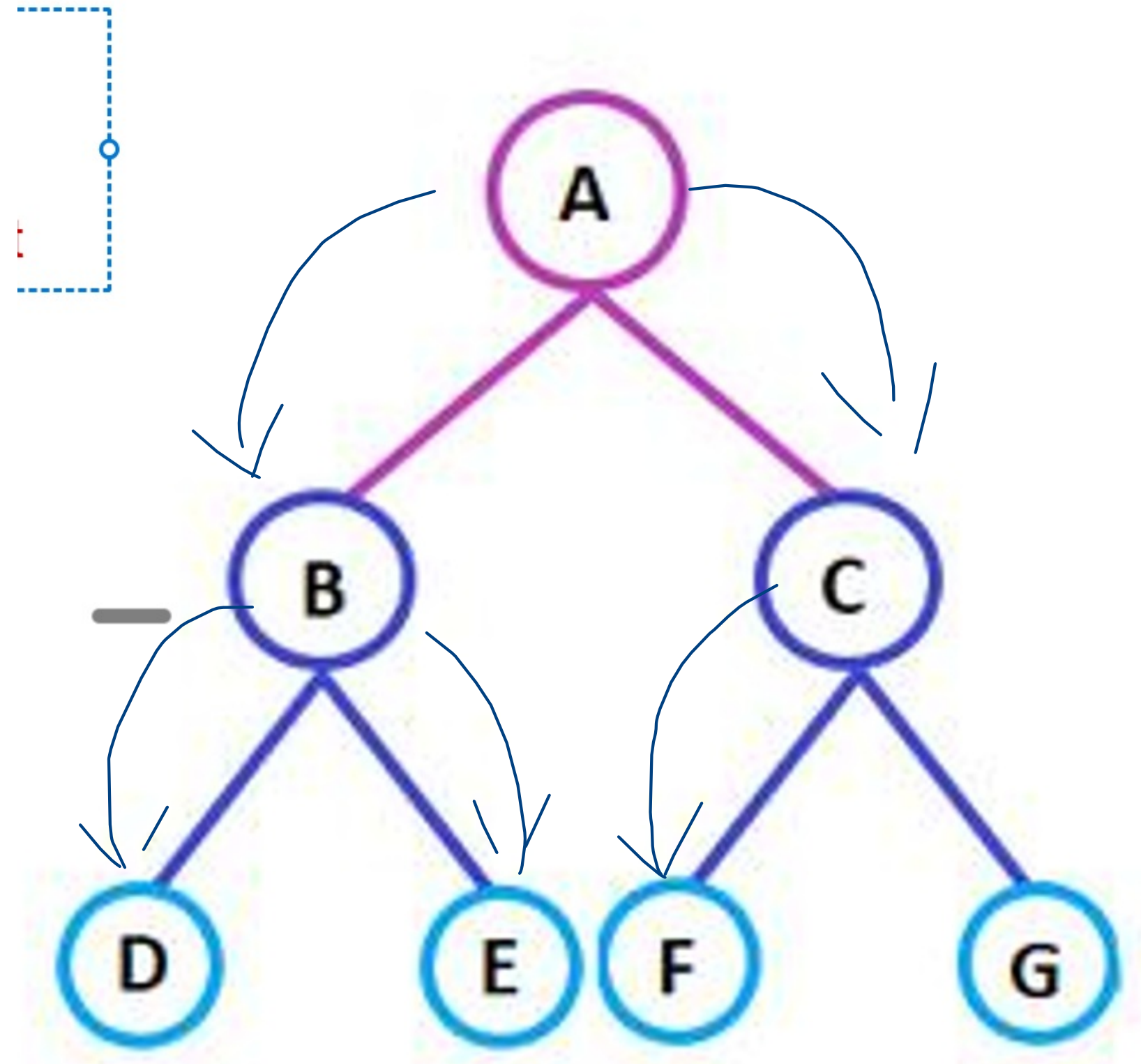
when u cut L:go to left

when u cut R:go to right

when u cut P:print

when all done go to parent

DEBFGCA



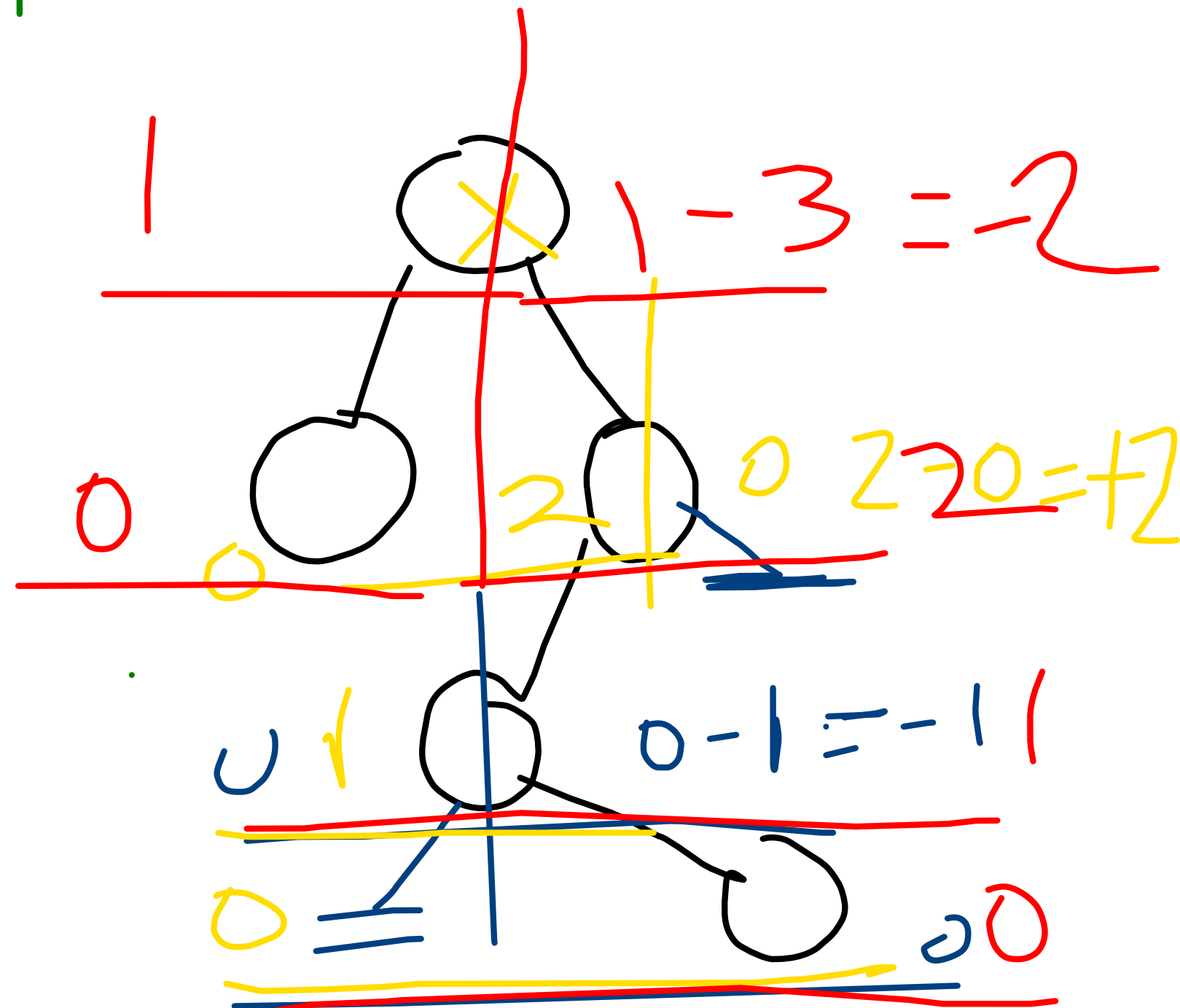
$$1-2=-1$$

$$3-1=2$$

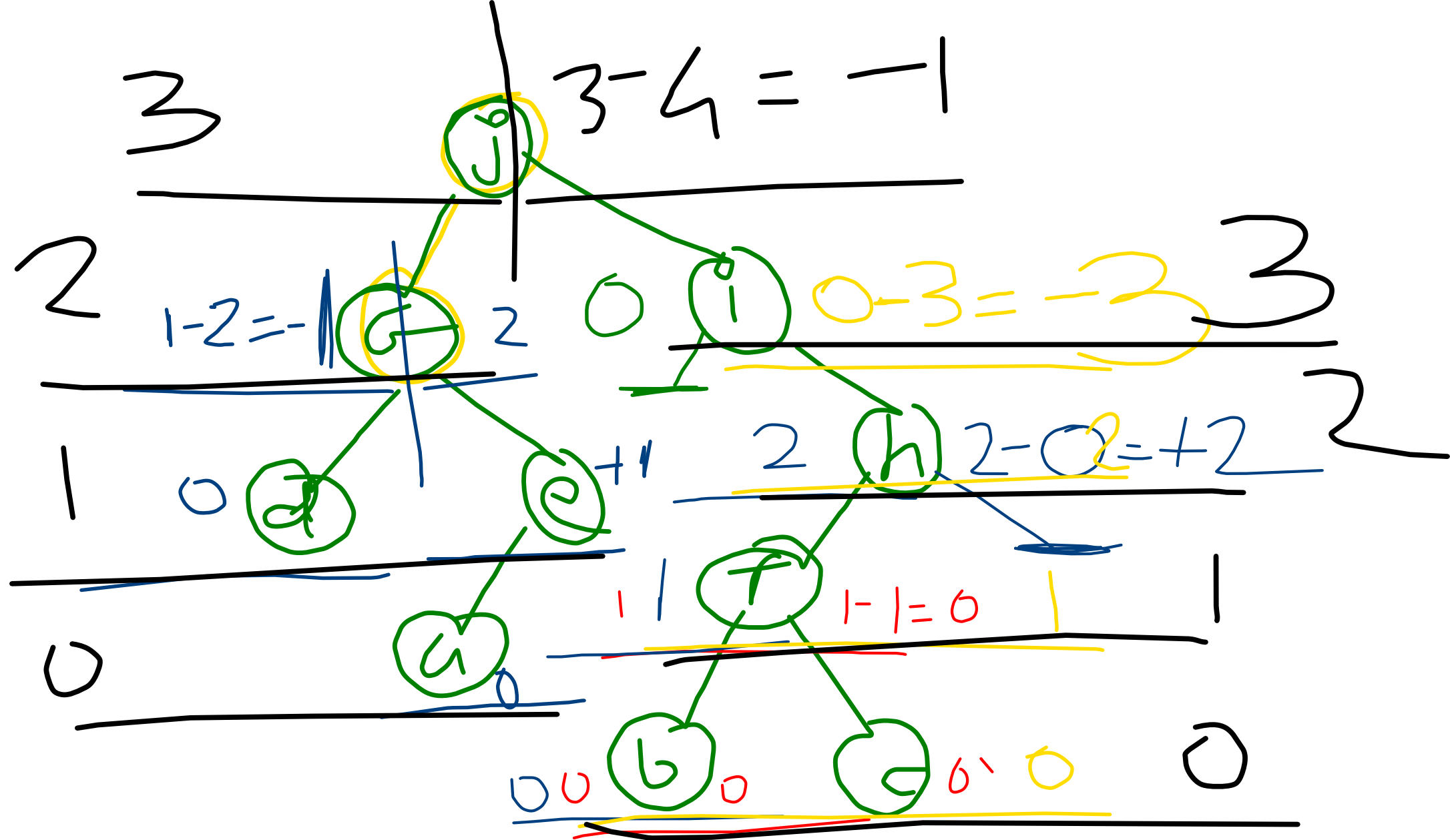


$$0-1-0=+1$$

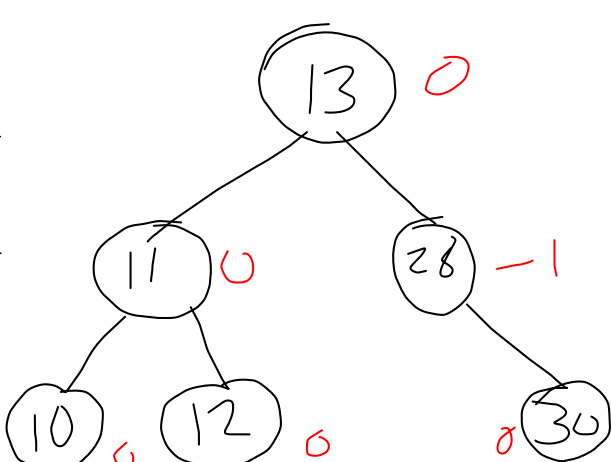
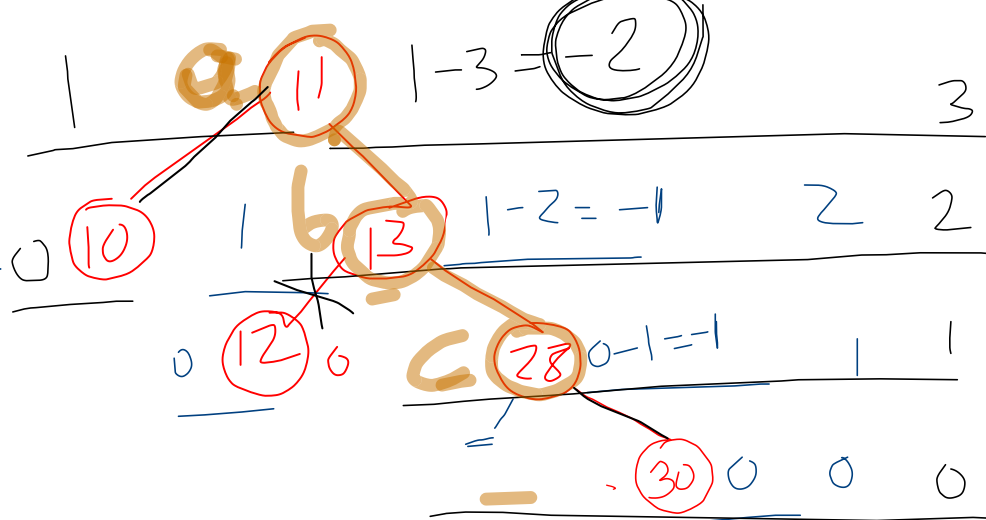
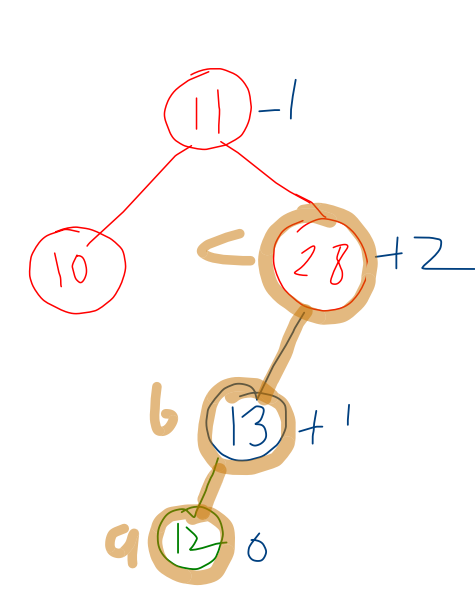
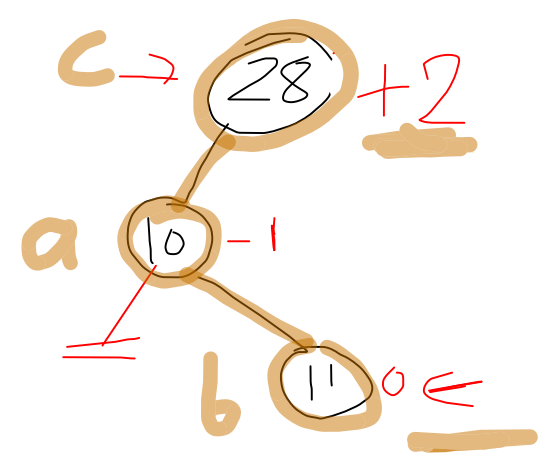
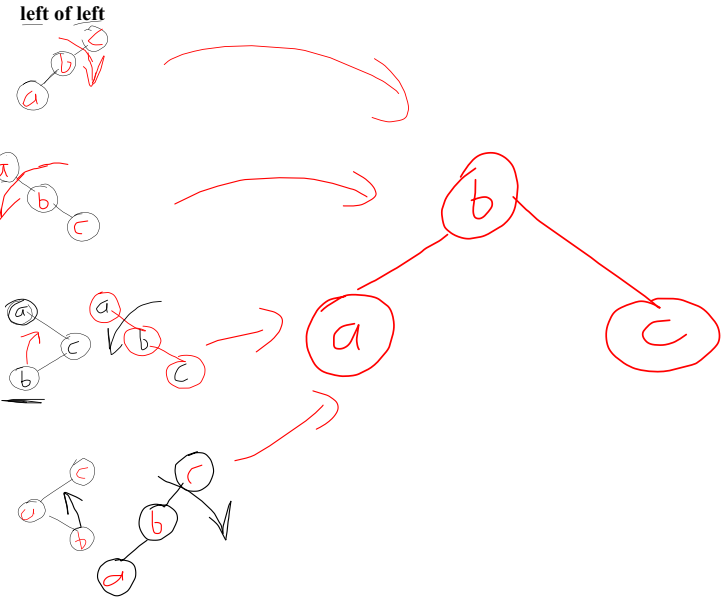
BF = HOLS - HORS



$$0-1=-1$$

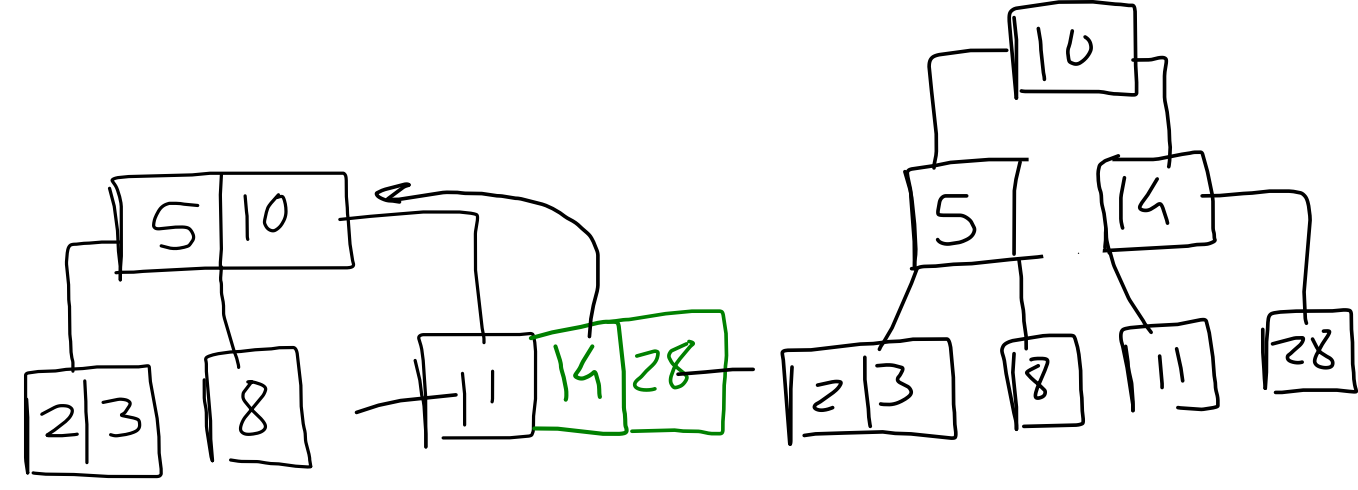


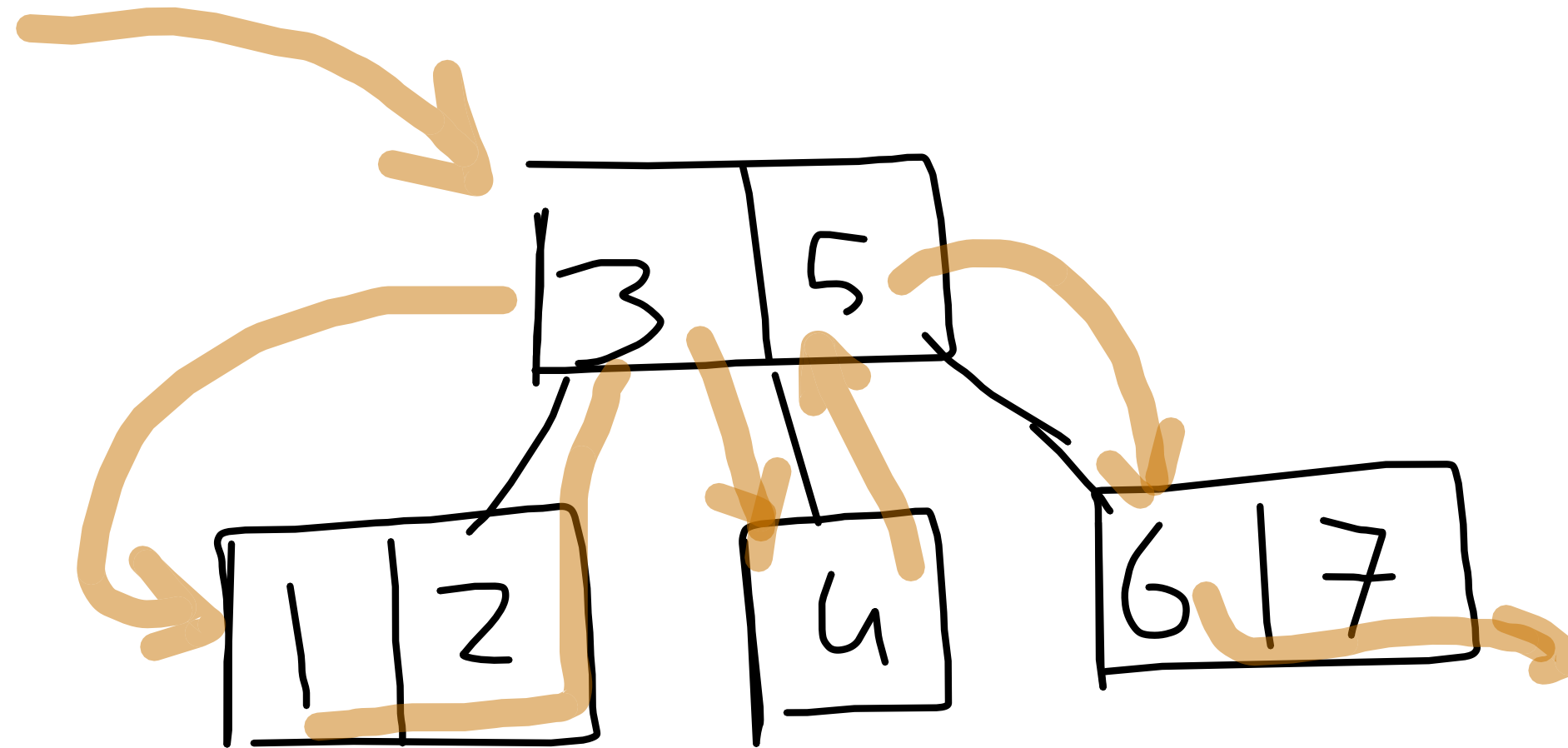
1. In AVL, insertion is done one node at a time in a BST fashion.
2. After every insertion, we calculate balance factor from bottom to top.
3. The first point where the balance factor is out of the ± 1 range (stop), see case and follow either of four types of rotation.
(look where insertion is done and where unbalanced is)



M=3 (min $3/2=1$, max $3-1=2$, split=3)

5, 10, 2, 3, 8, 11, 14, 28

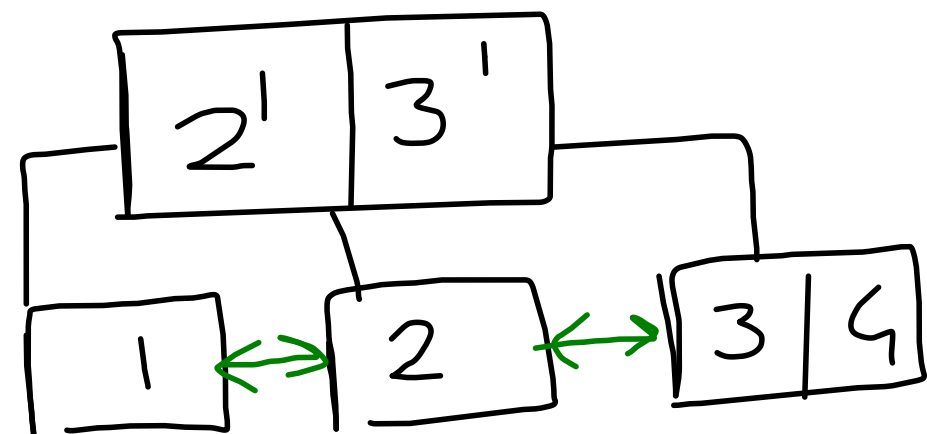
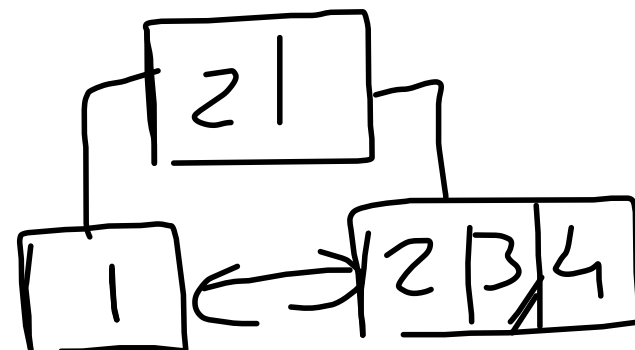
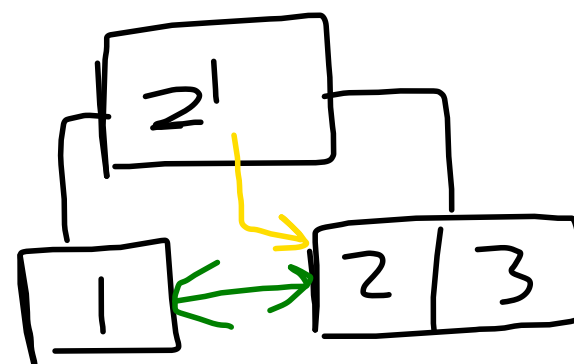
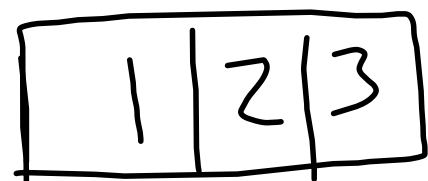




B-tree is excellent for storage but behaves in a very slow manner for sequential access of data because data is scattered among N nodes.

Sol: B+

$m=3$



1347%3000
1347