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Computation of Clustered Argumentation Frameworks via Boolean Satisfiability

MASTER'S THESIS

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Abstract

English abstract of your thesis

Kurzfassung

Deutsche Kurzfassung der Abschlussarbeit

Acknowledgements

Thanks to everyone who made this thesis possible

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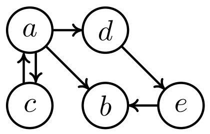
A	small
example	table

Table 2.1: A table caption for the list of tables.

3 Examples

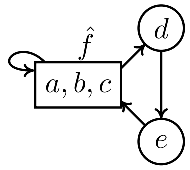
3.1 Basic AF

3.1.1 Concrete AF



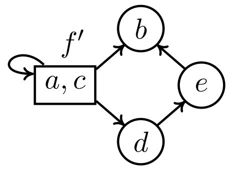
Stable Sets: $\{\}$, $\{a, e\}$, $\{b, c, d\}$

3.1.2 Abstract AF



Stable Sets: $\{\}$, $\{\hat{f}, e\}$, $\{\hat{f}, d\}$
concrete with main abstract $\rightarrow$ `FAITHFUL`

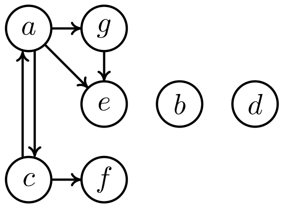
3.1.3 Abstract AF with Concretized Argument b



3.2 Basic Example

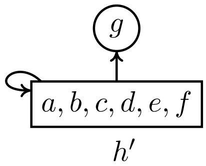
3.2.1 Concrete AF

Let $X = (ARG, ATT)$ be a concrete AF with the following arguments and attacks. Then the stable sets $ST(X)$ would be $\{\}, \{a, b, d, f\}, \{b, c, d, g\}$.



3.2.2 Abstract AF

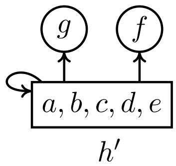
If we now abstract the concrete AF X to X', we obtain the following stable sets $\{\}, \{h'\}, \{h', g\}, \{g\}$. This would lead to a spurious abstraction, due to set $\{g\}$.



Now let X' be the input to our `CONCRETIZER` program and we parse as concretizer list the argument f.

3.2.3 Concretized Abstract AF (f)

We obtain the following AF X'' with the following stable sets $\{\}$, $\{h'\}$, $\{f, g, h'\}$, $\{g, h'\}$, $\{f, h'\}$, $\{f, g\}$. Which would lead to a spurious abstraction, due to the sets $\{h'\}$, $\{f, g, h'\}$, $\{f, g\}$.



3.2.4 Concretizing until Faithfulness

Since we want to obtain a faithful abstraction of the AF X' with the concretized argument f, we create all possible combinations of further concretization. Therefore, we need the spurious sets of X'' i.e. $\{h'\}$, $\{f, g, h'\}$, $\{f, g\}$. Since we are in the stable semantics, the depth of the concretizer search is 2 (i.e. if an argument x is spurious, we investigate all its attackers, and the attacker of the attackers and the same for the defenders (=the arguments which x attacks)).

Pre Filtering

The spurious sets of X'' can also have clusters in the sets. Since we relate to the attackers and defenders of the concrete AF X we can filter them out (because the concrete AF has no clusters). We then obtain the following sets: $\{f, g\}$, $\{f, g\}$, which can be reduced to $\{f, g\}$.

Attacker and Defender Depth 2

We now iterate over the filtered sets and check for each attacker a, the attackers of the attacker a_x. We also check, if a or a_x is in a cluster, because if they are not, we can not concretize them. Furthermore we add all the elements c from the concretizer list (if not already present) and create the following list of sets: $[\{a, c\}$, $\{a, a_0\, c\}$, $\{a, a_1\, c\}$, ..., $\{a, a_n, c\}]$ which in the current example would lead to the following list: $[\{c, f\}, \{a, c, f\}, \{a, f, g\}, \{a, c, f, g\}]$. The exact same is done with the defenders, where the list is $[\{e, f, g\}]$.

Combining Sets

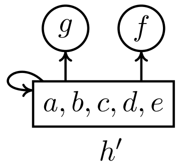
We now create each possible combination out of the 5 lists. This leads to a total of $\sum_{k=1}^{5} \binom{5}{k} = 31$ solutions. Since f.e. the combination $\{c, f\}$ and and $\{a, c, f\}$ are already covered in $\{a, c, f\}$ we remove the duplicates and obtain the following seven sets: $\{c, f\}$, $\{a, f\}$, $\{e, f\}$, $\{a, c, f\}$, $\{c, e, f\}$, $\{a, e, f\}$, $\{a, c, e, f\}$.

3.3 Problem

This approach works well for very small AF. But once we have more spurious sets, the list of the combinations 3.2.4 grows vastly. I had one instance, were 11 spurious sets led to 120 combinations which would then lead to $\sum_{k=1}^{120} \binom{120}{k}$ combinations, which is simply not feasible. Since conflict free sets produce a lot of sets, this case is not abstract and quite common.

3.4 Thoughts

If we consider the "larger" combinations first and once they result into faithfulness, we reduce the search to the selected set and try to concretize further each argument one by one. This would return a faithful solution. But I am not sure if it holds, that if the "larger" concretization AF is spurious, its fragmentation has to be spurious as well. To explain further what I mean: Let's take the previous example, where we tried to concretize the argument f.



Since this was spurious, we created the concretizer list $\{c, f\}$, $\{a, f\}$, $\{e, f\}$, $\{a, c, f\}$, $\{c, e, f\}$, $\{a, e, f\}$, $\{a, c, e, f\}$. Instead of creating the complete concretizer list (which is not feasible for a large amount of solutions as explained before) we produce a single set that contains all the unique singletons of the combinations, so in this example $\{a, c, e, f\}$. This is faithful in this case, so we focus only on this set and try to concretize its combinations. So in this case: $\{a, f\}$, $\{c, f\}$, $\{e, f\}$, $\{a, c, f\}$, $\{a, e, f\}$, $\{c, e, f\}$. If one of these is faithful, we found a better solution than the "larger" one. If all of these combinations are spurious, we just return the "larger" one.

For this approach I would simply compare each spurious solution, create one list of all the unique arguments. Instead of creating the combination list, I extend the spurious solution list with the attackers (and its attackers) and defenders (and its defenders).

Bibliography

[Knu97] Donald Ervin Knuth. *The Art of Computer Programming, Volume I: Fundamental Algorithms, 3rd Edition.* Addison-Wesley, 1997.