Chapter 4 LISTS

All the programs in this file are selected from Ellis Horowitz, Sartaj Sahni, and Susan Anderson-Freed

"Fundamentals of Data Structures in C",

Introduction

Array

successive items locate a fixed distance

disadvantage

- data movements during insertion and deletion
- waste space in storing n ordered lists of varying size

possible solution

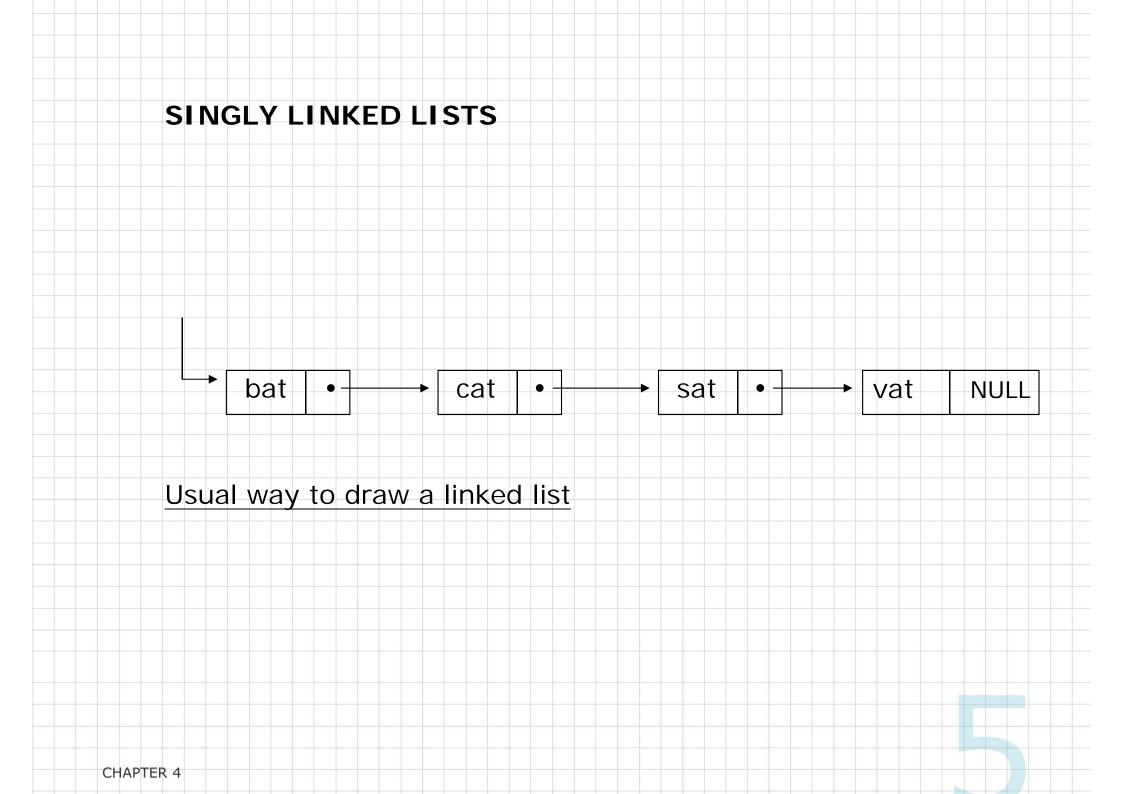
- linked list

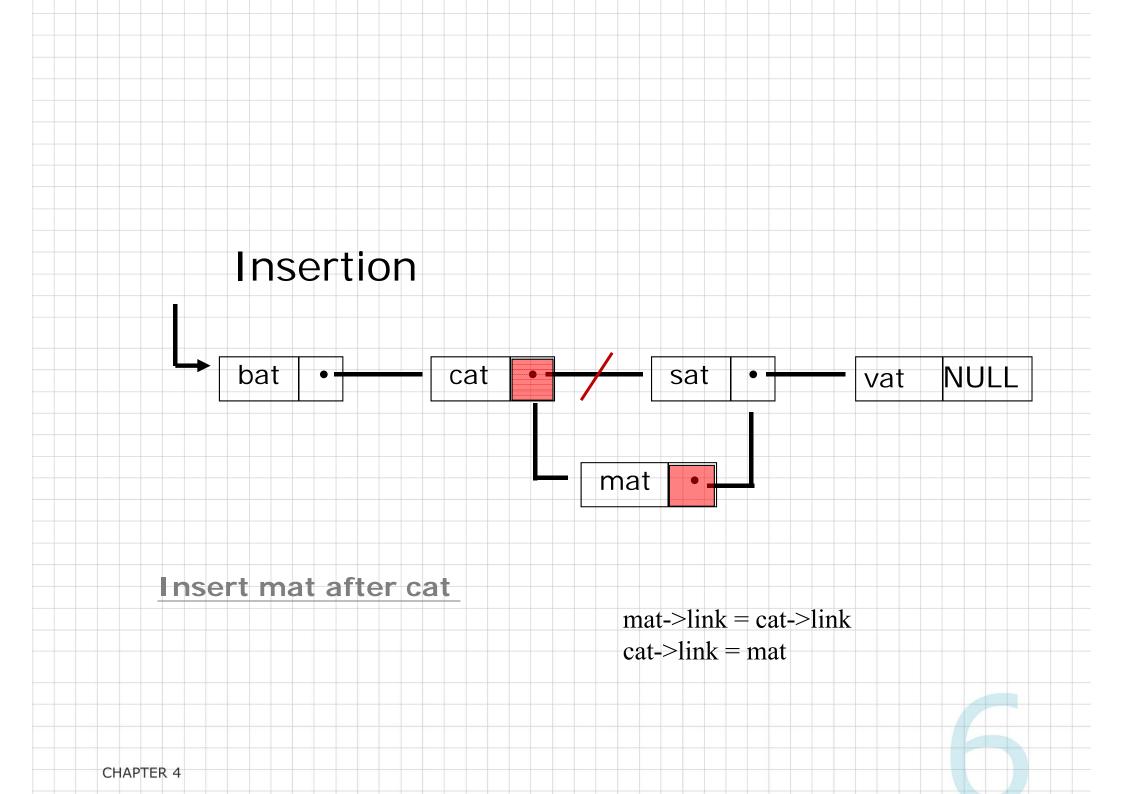
Pointer Can Be Dangerous

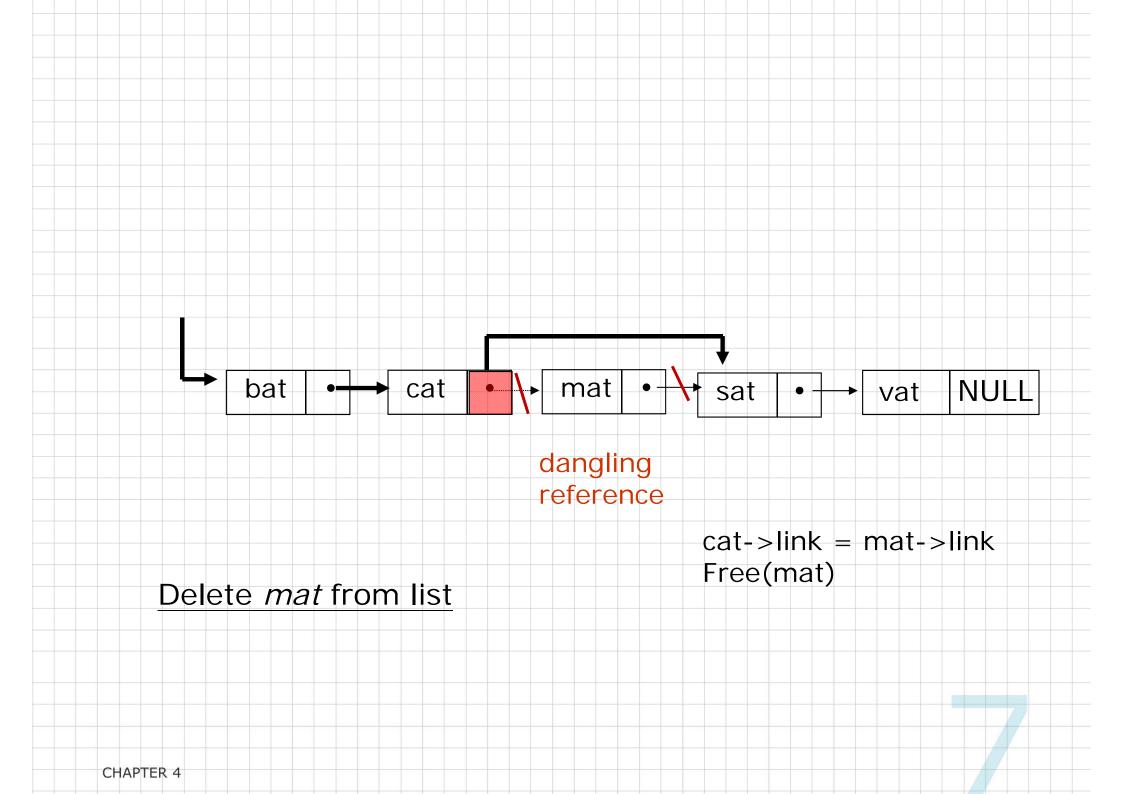
Using Dynamically Allocated Storage

```
int i, *pi;
float f, *pf;
pi = (int *) malloc(sizeof(int));
                                        request memory
pf = (float *) malloc (sizeof(float));
*pi = 1024;
*pf = 3.14;
printf("an integer = %d, a float = %f\n", *pi, *pf);
free(pi);
                        return memory
free(pf);
```

Allocation and deallocation of pointers

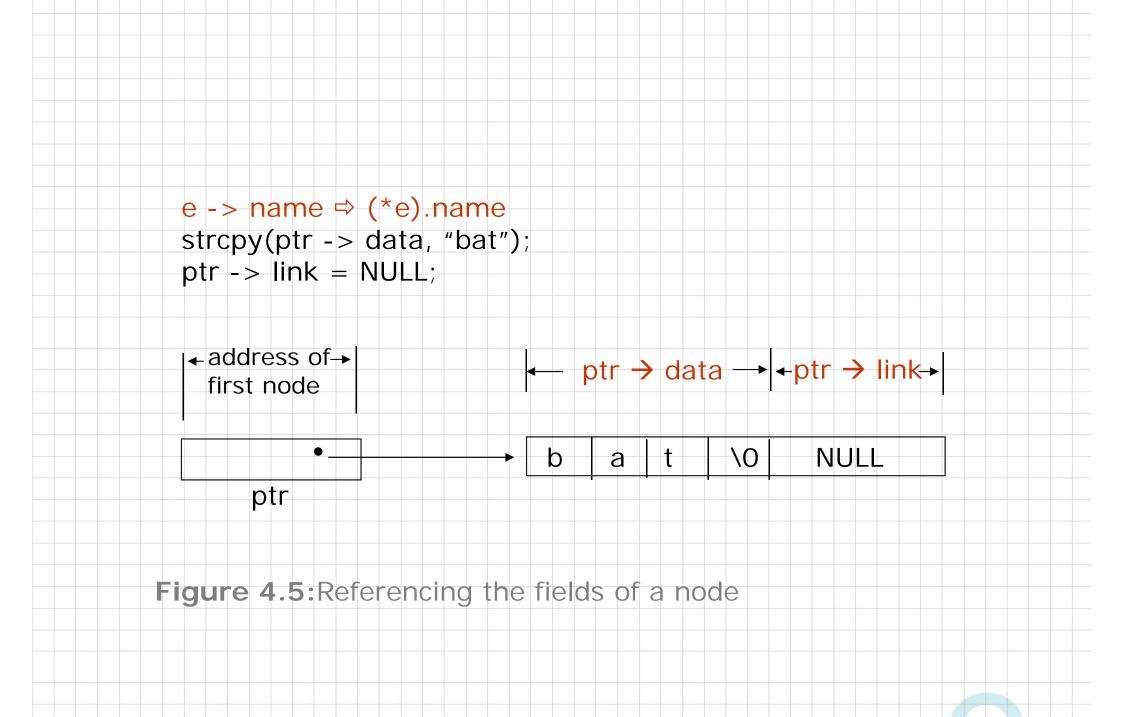






Example 4.1: create a linked list of words

Declaration



Example: create a two-node list

```
list_pointer create2()
 * create a linked list with two nodes */
    list_pointer first, second;
first = (list_pointer) malloc(sizeof(list_node));
second = (list_pointer) malloc(sizeof(list_node));
second -> link = NULL;
    second -> data = 20;
first -> data = 10;
                                            first
    first -> link = second;
    return first;
                                                        10
                                                                                     NULL
```

Program 4.1: Create a tow-node list

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List Insertion:

```
void insert(list_pointer *ptr, list_pointer node)
{
/* insert a new node with data = 50 into the list ptr after
node */
    list_pointer temp;
    temp = (list_pointer) malloc(sizeof(list_node));
    if (IS_FULL(temp)){
        fprintf(stderr, "The memory is full\n");
        exit (1);
    }
}
```

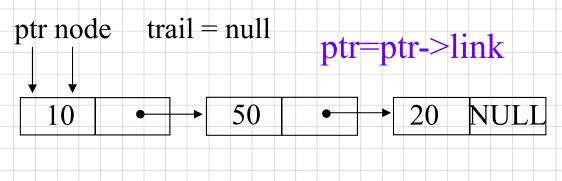
Insert a node after a specific node

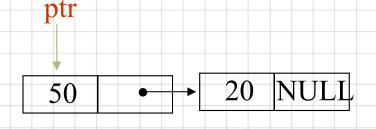
```
temp->data = 50;
  if (*ptr) { //noempty list
      temp->link = node ->link;
node->link = temp;
                                 ptr
 else { //empty list
temp->link = NULL;
                                         10
                                                           20
                                                               NULL
     *ptr =temp;
                                node
                                        temp
```

*Program 4.2: Simple insert into front of list

List Deletion

Delete the first node.

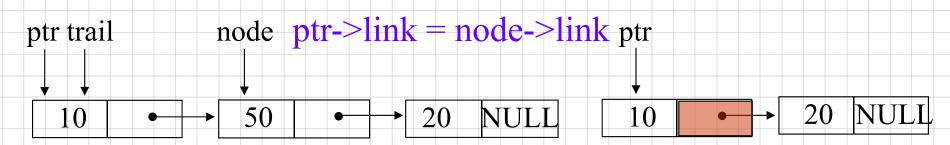




(a) before deletion

(b)after deletion

Delete node other than the first node.



List Deletion

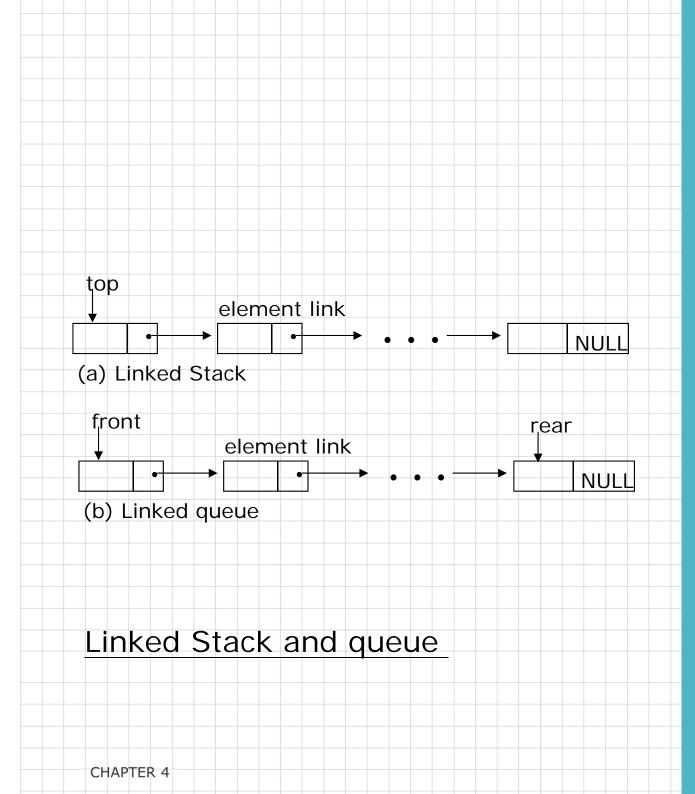
CHAPTER 4

```
void delete(list_pointer *ptr, list_pointer trail,
                                                     list_pointer node)
/* delete node from the list, trail is the preceding node
  ptr is the head of the list */
   if (trail)
                                     trail
                                                  node
     trail->link = node->link;
   else
                                                                    20
                                                                          NULL
                                                     50
                                      10
     *ptr = (*ptr) -> link;
    free(node);
                                                     20
                                                          NULL
                                      10
 ptr, node
               trail = null
                                                  ptr
                                      NULL
                                 20
                                                                  20
                                                                      NULL
  10
                  50
                                                  50
```

Print out a list (traverse a list)

```
void print_list(list_pointer ptr)
{
    printf("The list ocntains: ");
    for (; ptr; ptr = ptr->link)
        printf("%4d", ptr->data);
    printf("\n");
}
```

Program 4.4: Printing a list (p. 155)



DYNAMICALLY LINKED STACKS AND QUEUES

Chapter 4.3

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Represent n stacks

```
#define MAX_STACKS 10 /* maximum number of stacks */
typedef struct {
         int key;
         /* other fields */
        } element;
typedef struct stack *stack_pointer;
typedef struct stack {
        element item;
        stack_pointer link;
stack_pointer top[MAX_STACKS];
```

Represent n queues

Push in the linked stack

```
void add(stack_pointer *top, element item)
 /* add an element to the top of the stack */
  stack_pointer temp =
                 (stack_pointer) malloc (sizeof (stack));
  if (IS_FULL(temp)) {
    fprintf(stderr, "The memory is full\n");
    exit(1);
    temp->item = item;
    temp->link = *top;
    *top= temp;
                           Add to a linked stack
              top
                       element link
                                               NULL
```

CHAPTER 4

pop from the linked stack

```
element delete(stack_pointer *top) {
/* delete an element from the stack */
    stack_pointer temp = *top;
    element item;
    if (IS_EMPTY(temp)) {
        fprintf(stderr, "The stack is empty\n");
    }
}
              exit(1);
   item = temp->item;
*top = temp->link;
free(temp);
return item;
 Delete from a linked stack
                                  ţοp
                                                       element link
```

enqueue in the linked queue

```
void addq(queue_pointer *front, queue_pointer *rear,
element item)
{ /* add an element to the rear of the queue */
  queue_pointer temp =
                 (queue_pointer) malloc(sizeof (queue));
  if (IS_FULL(temp)) {
    fprintf(stderr, " The memory is full\n");
    exit(1);
    temp->item = item;
   temp->link = NULL;
    if (*front)) (*rear) -> link = temp;
    else *front = temp;
    *rear = temp; }
       front
                                       rear
                 element link
                                          NULL
```

CHAPTER 4

dequeue from the linked queue

```
element deleteq(queue_pointer *front) {
/* delete an element from the queue */
  queue_pointer temp = *front;
  element item;
  if (IS_EMPTY(*front)) {
    fprintf(stderr, "The queue is empty\n");
    exit(1);
  item = temp->item;
  *front = temp->link;
   free(temp);
   return item;
    front
             element link
                                       NULL
```

$$A(x) = a_{m-1}x^{e_{m-1}} + a_{m-2}x^{e_{m-2}} + ... + a_0x^{e_0}$$

$$typedef struct poly_node$$

$$*poly_pointer;$$

$$typedef struct poly_node$$

$$int coef;$$

$$int expon;$$

$$poly_pointer link;$$

$$);$$

$$poly_pointer a, b;$$

$$coef expon link$$

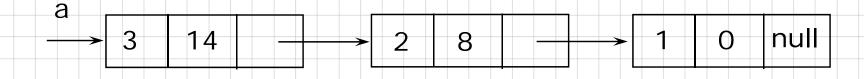
Polynomials

Chapter 4.4

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Examples

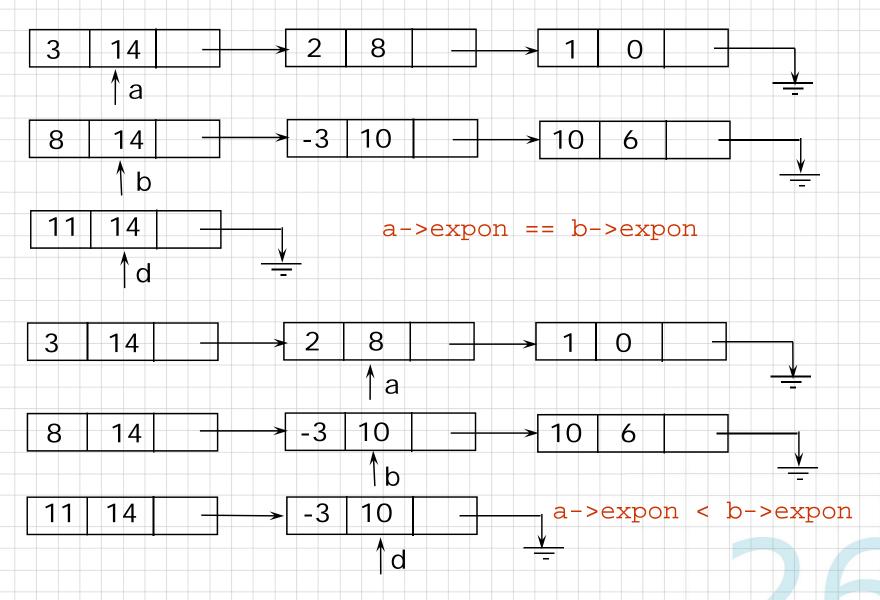
$$a = 3x^{14} + 2x^8 + 1$$



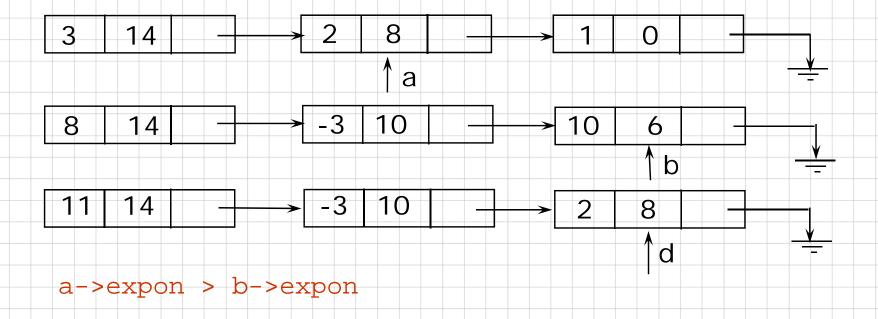
$$b = 8x^{14} - 3x^{10} + 10x^6$$

b 8 14 -3 10 -3 10 6 null

Adding Polynomials



Adding Polynomials (Cont.)



Alogrithm for Adding Polynomials

```
case -1: /* a->expon < b->expon */
            attach(b->coef, b->expon, &rear);
            b= b->link;
            break;
        case 0: /* a->expon == b->expon */
             sum = a - scoef + b - scoef;
             if (sum) attach(sum,a->expon,&rear);
            a = a - \lambda i nk; b = b - \lambda i nk;
            break;
        case 1: /* a->expon > b->expon */
            attach(a->coef, a->expon, &rear);
            a = a - > link;
for (; a; a = a->link)
    attach(a->coef, a->expon, &rear);
for (; b; b=b->link)
    attach(b->coef, b->expon, &rear);
rear->link = NULL;
temp = front; front = front->link; free(temp);
return front;
```

Delete extra initial node.

Attach a Term

```
void attach(float coefficient, int exponent,
             poly_pointer *ptr)
   /* create a new node attaching to the node pointed to
   by ptr. ptr is updated to point to this new node. */
    poly_pointer temp;
    temp = (poly_pointer) malloc(sizeof(poly_node));
    if (IS_FULL(temp)) {
         fprintf(stderr, "The memory is full\n");
         exit(1);
    temp->coef = coefficient;
    temp->expon = exponent;
(*ptr)->link = temp;
*ptr = temp;
                                      14
```

Analysis

- (1) coefficient additions
 - $0 \le additions \le min(m, n)$

where m (n) denotes the number of terms in A (I

(2) exponent comparisons extreme case

```
e_{m-1} > f_{m-1} > e_{m-2} > f_{m-2} > ... > e_0 > f_0
m+n-1 comparisons
```

(3) creation of new nodes extreme case

A Suite for Polynomials e(x) = a(x) * b(x) + d(x)

```
poly_pointer a, b, d, e;

a = read_poly();
b = read_poly();
d = read_poly();
temp = pmult(a, b);
e = padd(temp, d);
print_poly(e);
```

```
read_poly()
print_poly()
padd()
psub()
pmult()
```

temp is used to hold a partial result. By returning the nodes of temp, we may use it to hold other polynomials

Erase Polynomials

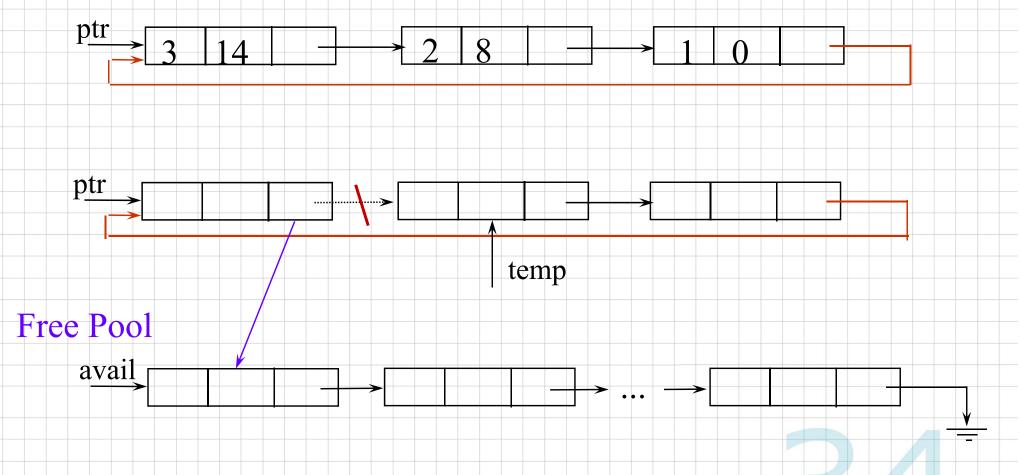
```
void earse(poly_pointer *ptr)

/* erase the polynomial pointed to by ptr */
    poly_pointer temp;
    while (*ptr) {
   temp = *ptr;
   *ptr = (*ptr)->link;
          free(temp);
```

Circularly Linked Lists

circular list vs. chain

CHAPTER 4



Maintain an Available List

```
poly_pointer get_node(void)
 poly_pointer node;
 if (avail) {
    node = avail;
    avail = avail->link:
 else {
    node = (poly_pointer)malloc(sizeof(poly_node));
    if (IS_FULL(node)) {
       printf(stderr, "The memory is full\n");
       exit(1);
 return node;
```

Maintain an Available List (cont.)

Insert ptr to the front of this list

```
void ret_node(poly_pointer ptr)
{
  ptr->link = avail;
  avail = ptr;
}
```

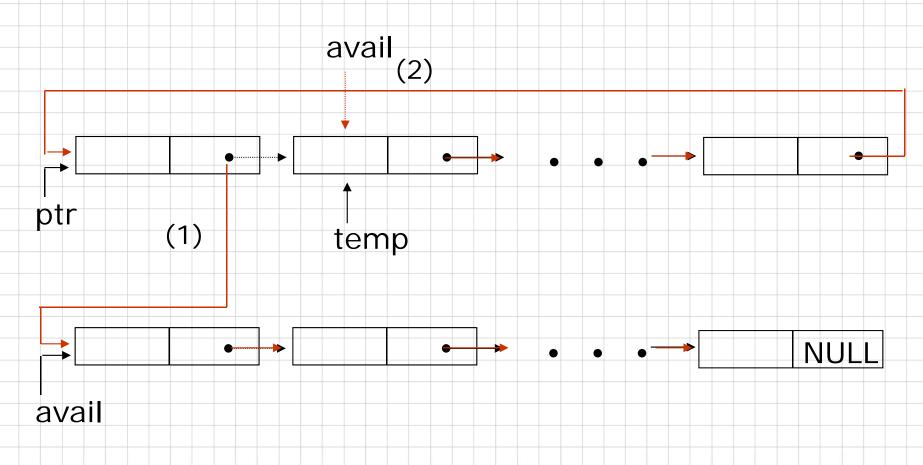
dependent of # of nodes in a list O(n)

Maintain an Available List (cont.)

```
void cerase(poly_pointer *ptr)
{
    poly_pointer temp;
    if (*ptr) {
        temp = (*ptr)->link;
        (*ptr)->link = avail;
        avail = temp;
        *ptr = NULL;
    }
}
```

Independent of # of nodes in a list O(1) constant time

4.4.4 Representing Polynomials As Circularly Linked Lists



Returning a circular list to the avail list

Head Node

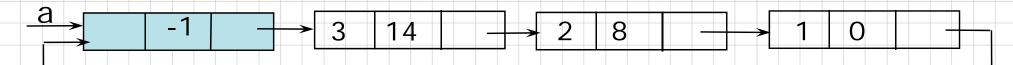
Represent polynomial as circular list.

(1) zero



Zero polynomial

(2) others



$$a = 3x^{14} + 2x^8 + 1$$

Another Padd

```
poly_pointer cpadd(poly_pointer a, poly_pointer b)
  poly_pointer starta, d, lastd;
  int sum, done = FALSE;
  starta = a;
  a = a - > link;
                   Set expon field of head node to -1.
  b = b - > link;
  d = get_node();
  d \rightarrow expon = -1; lastd = d;
  do
    switch (COMPARE(a->expon, b->expon))
      case -1: attach(b->coef, b->expon, &lastd);
                b = b - \sinh i
                break;
```

Another Padd (Continued)

```
case 0: if (starta == a) done = TRUE;
            else {
              sum = a->coef + b->coef;
              if (sum) attach(sum,a->expon,&lastd);
              a = a->link; b = b->link;
            break;
    case 1: attach(a->coef,a->expon,&lastd);
            a = a - > link;
 while (!done);
lastd->link = d; Link last node to first
return d;
```

```
typedef struct list_node
*list_pointer;

typedef struct list_node {
    char data;
    list_pointer link;
};
```

Invert single linked lists

Concatenate two linked lists

Additional List Operations

Chapter 4.5

Pade 171

Invert Single Linked Lists

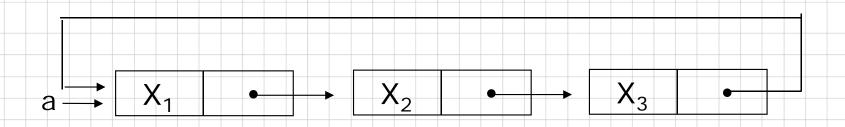
```
list_pointer invert(list_pointer lead)
    list_pointer middle, trail;
    middle = NULL;
    while (lead) {
         trail = middle;
         middle = lead;
         lead = lead->link;
         middle->link = trail;
                            0: null
    return middle;
                            1: lead
                            ≥2: lead-
   Refer to the ppt of Inverting a
       Singly Linked List
```

Concatenate Two Lists

```
list_pointer concatenate(list_pointer
         ptr1, list_pointer ptr2)
  list_pointer temp;
 if (IS_EMPTY(ptr1)) return ptr2;
 else {
   if (!IS_EMPTY(ptr2)) {
    for (temp=ptr1; temp->link; temp=temp->link);
    temp->link = ptr2;
   return ptr1;
                    O(m) where m is # of elements in the first list
CHAPTER 4
```

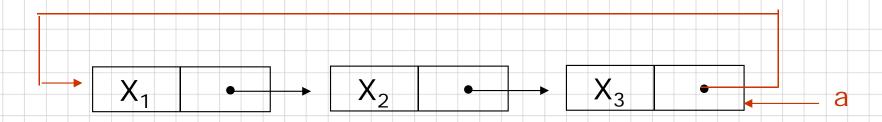
Operations For Circularly Linked List

What happens when we insert a node to the back of a circular linked list?



Problem: move down the whole list.

A possible solution:



Note a pointer points to the last node.

Operations for Circular Linked Lists

```
void insert_front (list_pointer *ptr, list_pointer node)
     if (IS_EMPTY(*ptr)) {
       *ptr= node;
       node->link = node;
     else {
        node->link = (*ptr)->link; (1)
        (*ptr)->link = node; (2)
                                                             ptr
CHAPTER 4
```

Length of a Circular Linked List

```
int length(list_pointer ptr)
   list_pointer temp;
   int count = 0;
   if (ptr) {
      temp = ptr;
      do {
          count++;
          temp = temp->link;
      } while (temp!=ptr);
   return count;
CHAPTER 4
```

Equivalence Relations

A relation over a set, S, is said to be an equivalence relation over S iff it is symmertric, reflexive, and transitive over S.

reflexive, x=x
symmetric, if x=y, then y=x
transitive, if x=y and y=z, then x=z

Examples

```
0 \equiv 4, 3 \equiv 1, 6 \equiv 10, 8 \equiv 9, 7 \equiv 4, 6 \equiv 8, 3 \equiv 5, 2 \equiv 11, 11 \equiv 0
```

three equivalent classes {0,2,4,7,11}; {1,3,5}; {6,8,9,10}

```
Scan and combine operation
```

```
0=4 -> group 1 {0,4}

3=1 => group 2 {3,1}

6=10 => group 3 {6,10}

8=9 => group 4 {8,9}

7=4 => group 1{0,4,7}

6=8 => group 4{6,8,9} + group 3 {6, 10} => group 3 {6, 8, 9, 10}

3=5 => group 2{3,1,5}

2=11 => group 5{2,11}

11=0 => group 5 {2, 11, 0} + group 1{0,4,7} => group 1{0,2,4,7,11}
```

A Rough Algorithm to Find Equivalence Classes

```
void equivalenec()
  initialize;
  while (there are more pairs) {
     read the next pair <i,j>;
     process this pair;
   initialize the output;
   do {
     output a new equivalence class;
   } while (not done);
```

What kinds of data structures are adopted?

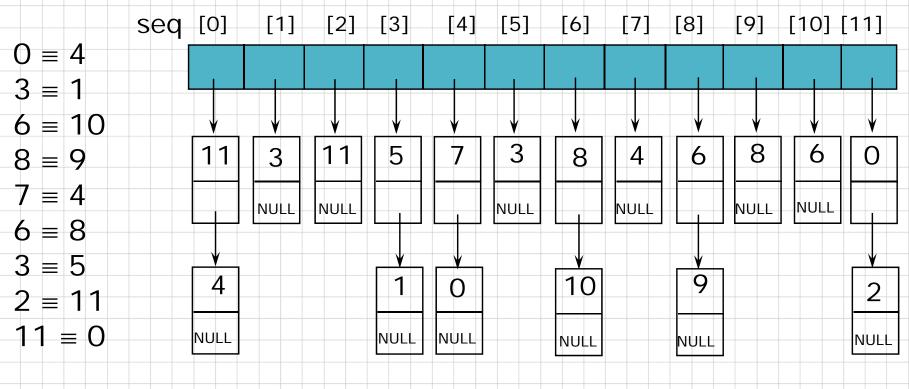
First Refinement

```
void equivalence()
#include <stdio.h>
#include <alloc.h>
#define MAX_SIZE 24
                               initialize seq to NULL and out to TRUE
#define IS_FULL(ptr) (!(ptr))
                               while (there are more pairs) {
#define FALSE 0
                                  read the next pair, <i,j>;
#define TRUE
                                  put j on the seq[i] list;
            direct equivalence
                                  put i on the seq[j] list;
                               for (i=0; i< n; i++)
                                  if (out[i]) {
Compute indirect equivalence
                                     out[i] = FALSE;
using transitivity
                                     output this equivalence class;
```

three equivalent classes

Lists After Pairs are input

$$\{0,2,4,7,11\}; \{1,3,5\}; \{6,8,9,10\}$$



```
typedef struct node *node_pointer;
typedef struct node {
   int data;
   node_pointer link;
};
```

Final Version for Finding Equivalence Classes

```
void main(void)
 short int out[MAX_SIZE];
 node_pointer seg[MAX_SIZE];
 node_pointer x, y, top;
 int i, j, n;
 printf("Enter the size (<= %d) ", MAX_SIZE);
 scanf("%d", &n);
 for (i=0; i< n; i++) {
    out[i] = TRUE; seq[i] = NULL;
 printf("Enter a pair of numbers (-1 -1 to quit): ");
 scanf("%d%d", &i, &j);
```

Phase 1: input the equivalence pairs

```
while (i \ge 0)
    x = (node_pointer) malloc(sizeof(node));
    if (IS_FULL(x))
      fprintf(stderr, "memory is full\n");
        exit(1);
     Insert x to the top of lists seq[i]
    x->data=j; x->link=seq[i]; seq[i]=x;
    if (IS_FULL(x))
      fprintf(stderr, "memory is full\n");
        exit(1);
       Insert x to the top of lists seq[j]
    x->data= i; x->link= seq[j]; seq[j]= x;
    printf("Enter a pair of numbers (-1 -1 to )
          quit): ");
    scanf("%d%d", &i, &j);
```

```
for (i=0; i<n; i++) {
    if (out[i]) {
        printf("\nNew class: %5d", i);
out[i] = FALSE;
        x = seq[i]; top = NULL;
for (;;) {
    while (x) {
                  élse x = x->link;
                                             pop
             if (!top) break;
x = seq[top->data]; top = top->link;
```

Phase 2: output the equivalence classes

4.7 Sparse Matrices

inadequates of sequential schemes

- (1) # of nonzero terms will vary after some matrix computation
- (2) matrix just represents intermediate results

new scheme

Each column (row): a circular linked list with a head node

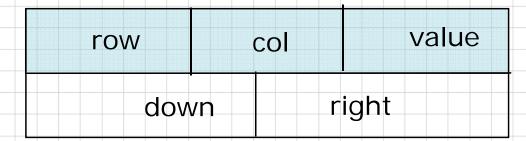
Revisit Sparse Matrices

of head nodes = max{# of rows, # of columns}

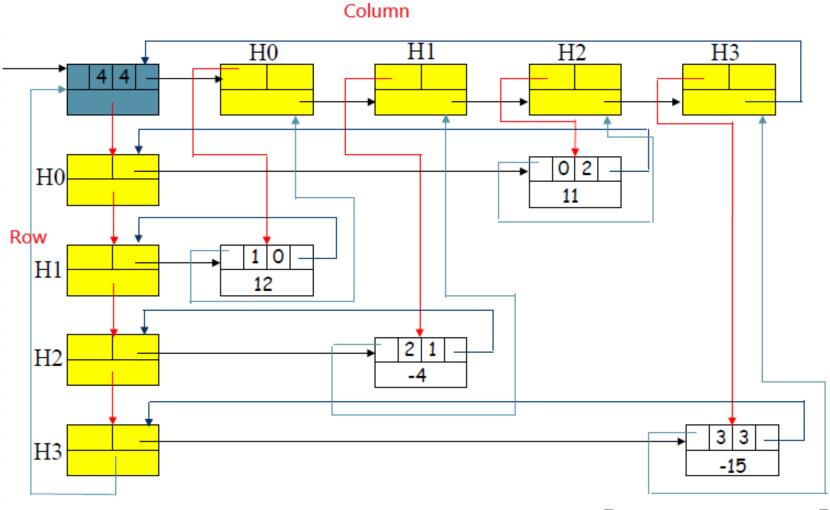
head node



element node



Linked Representation for Matrix



For an n by m sparse matrix with r nonzero terms, the number of nodes needed is max{n, m} + r + 1.

 $\begin{bmatrix} 0 & 0 & 11 & 0 \\ 12 & 0 & 0 & 0 \\ 0 & -4 & 0 & 0 \\ 0 & 0 & 0 & -15 \end{bmatrix}$

```
#define MAX_SIZE 50 /* size of largest matrix */
typedef enum {head, entry} tagfield;
typedef struct matrix_node *matrix_pointer;
typedef struct entry_node {
    int row;
    int col;
                      int value;
typedef struct matrix_node {
    matrix_pointer down;
    matrix_pointer right;
    tagfield tag;
                      union {
                                         matrix_pointer next; entry_node entry;
                       } u;
 matrix_pointer hdnode[MAX_SIZE];
```

	[0]		[2]
[0]	4	4	4
	0	2	11
[2]	1	0	12
[3]	2	1	-4
	3	3	-15

Figure 4.20: Sample input for sparse matrix (p.182)

Read in a Matrix

```
matrix_pointer mread(void)
{
/* read in a matrix and set up its linked
list. An global array hdnode is used */
  int num_rows, num_cols, num_terms;
  int num_heads, i;
  int row, col, value, current_row;
  matrix_pointer temp, last, node;

  printf("Enter the number of rows, columns and number of nonzero terms: ");
```

```
scanf("%d%d%d", &num_rows, &num_cols,
        &num terms);
 num heads =
 (num_cols>num_rows)? num_cols : num_rows;
 /* set up head node for the list of head
     nodes */
 node->u.entry.row = num_rows;
node->u.entry.col = num_cols;
 if (!num_heads) node->right = node;
else { /* initialize the head nodes */
for (i=0; i<num_heads; i++) {</pre>
      term= new_node();
hdnode[i] = temp;
hdnode[i] ->tag = head;
hdnode[i] ->right = temp;
      hdnode[i]->u.next = temp;
```

O(max(n,m))

```
current_row= 0;     last= hdnode[0];
for (i=\overline{0}; i< num\_terms; i++)
  printf("Enter row, column and value:");
  scanf("%d%d%d", &row, &col, &value);
  if (row>current_row) {
    last->right= hdnode[current_row];
    current_row= row; last=hdnode[row];
  temp = new node();
  temp->tag=entry; temp->u.entry.row=row;
  temp->u.entry.col = col;
  temp->u.entry.value = value;
  last->right = temp;/*link to row list */
  last= temp;
  /* link to column list */
  hdnode[col]->u.next->down = temp;
  hdnode[col]=>u.next = temp;
```

利用next field 存放column的last node

```
/*close last row */
  last->right = hdnode[current_row];
/* close all column lists */
  for (i=0; i<num_cols; i++)
    hdnode[i]->u.next->down = hdnode[i];
  /* link all head nodes together */
  for (i=0; i<num_heads-1; i++)
    hdnode[i]->u.next = hdnode[i+1];
  hdnode[num_heads-1]->u.next= node;
  node->right = hdnode[0];
return node;
            O(max{#_rows, #_cols} + #_terms)
```

Write out a Matrix

```
void mwrite(matrix_pointer node)
{ /* print out the matrix in row major form */
  int i;
  matrix_pointer temp, head = node->right;
  printf("\n num_rows = %d, num_cols= %d\n",
         node->u.entry.row,node->u.entry.col);
  printf("The matrix by row, column, and
         value:\n\n'');
  for (i=0; i < node->u.entry.row; i++) { O(\#_rows+\#_terms)
    for (temp=head->right;temp!=head;temp=temp->right)
     printf("%5d%5d%5d\n", temp->u.entry.row,
           temp->u.entry.col, temp->u.entry.value);
    head= head->u.next; /* next row */
```

Erase a Matrix Free the entry and head nodes by row.

```
void merase(matrix_pointer *node)
  int i, num_heads;
  matrix_pointer x, y, head = (*node)->right;
for (i=0; i<(*node)->u.entry.row; i++) {
    y=head->right;
    while (y!=head) {
       x = y; y = y - \dot{y}; free(x);
    \dot{x}= head; head= head->u.next; free(x);
  \dot{y} = \text{head};
  while (y!=*node) {
    x = y; y = y-\lambda i.next; free(x);
  free(*node); *node = NULL;
  O(#_rows+#_cols+#_terms)
```

Alternative: 利用Fig 4.14的技巧, 把一列資料erase (constant time)

Doubly Linked List

Move in forward and backward direction.

Singly linked list (in one direction only)

How to get the preceding node during deletion or insertion?

Using 2 pointers

Node in doubly linked list left link field (llink) data field (item) right link field (rlink)

Doubly Linked Lists

```
typedef struct node *node_pointer;
typedef struct node {
                                          ptr
                                        = ptr->rlink->llink
   node_pointer llink;
                                        = ptr->llink->rlink
   element item;
   node_pointer rlink;
              head node
                                   right
               rlink
     llink
          item
```

An Empty Node

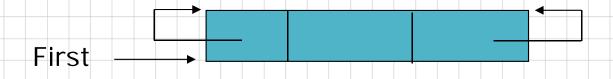


Figure 4.22: Empty doubly linked circular list with head node (p.188)

A Node Insertion

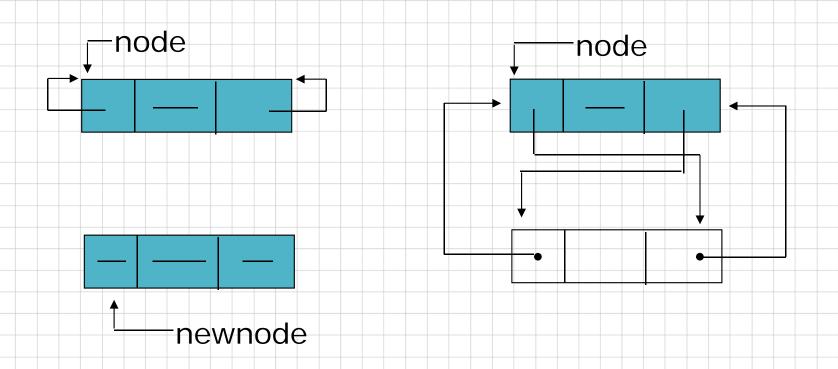


Figure 4.23: Insertion into an empty doubly linked circular list (p.189)

Insert

```
void dinsert(node_pointer node, node_pointer newnode)
      (1) newnode->llink = node;
      (2) newnode->rlink = node->rlink;
(3) node->rlink->llink = newnode;
(4) node->rlink = newnode;
                 head node
                  rlink
      llink
             item
                 (4)
```

Delete

```
void ddelete(node_pointer node, node_pointer deleted)
     if (node==deleted) printf("Deletion of head node
                                    not permitted.\n");
     else
          (1) deleted->llink->rlink= deleted->rlink;
(2) deleted->rlink->llink= deleted->llink;
               free(deleted);
               head node
```