



#### **Outline**

#### Homework to do:

Hardware

#### **ISA**

New instructions(CSR instructions, mret, WFI)

#### **IPs**

- DRAM Wrapper
- ROM Wrapper (storing booting code)
- > DMA
- Watch Dog timer (clocked by 20 MHz)
- Software & Bootloader
  - Sorting, Gray Scale
  - Timer interrupt
  - DMA interrupt
  - Booting(from DRAM to IM & DM)
- EDA tool verification
  - Spyglass CDC check
- Submission rule





## Problem 1 - Hardware

Specification





#### **New instructions - Overview**

To support trap handler, there are 8 new instructions in HW3:

- CSR Instructions (6) see unprivileged ISA
  - Register operand
    - CSRRW (Atomic Read/Write CSR)
    - CSRRS (Atomic Read and Set Bits in CSR)
    - CSRRC (Atomic Read and Clear Bits in CSR)
  - Immediate operand
    - CSRRWI (Atomic Read/Write CSR)
    - CSRRSI (Atomic Read and Set Bits in CSR)
    - CSRRCI (Atomic Read and Clear Bits in CSR)
- Trap-Return Instructions (1) see privileged ISA
  - Machine Mode return from trap
    - **MRET**
- Interrupt-Management Instructions (1) see privileged ISA
  - Wait for Interrupt







# ¥B €

## **New instructions – Corresponding registers**

- □ CSRs to be implemented see privileged ISA
- ☐ There are 3 privilege levels defined by RISCV
  - Machine level highest priority, which is to be implemented in HW3
  - Supervisor level OS usually operates in this level
  - User level user applications usually run on this level
- □ 6 machine mode CSRs to be implemented:

Address	Privilege	Name	Description
0x300	М	mstatus	Machine status register
0x304	М	mie	Machine interrupt-enable register
0x305	М	mtvec	Machine Trap-Vector Base-Address register
0x341	М	mepc	Machine exception program counter
0x344	М	mip	Machine interrupt pending register

☐ See CSR appendix for details of m-mode CSRs



## **New instructions - Description**

☐ Description of CSR instructions in unprivileged ISA

31 20	19 15	14 12	11 7	6 0		
imm[11:0]	rs1	funct3	rd	opcode	Mnemonic	Description
csr	rs1	001	rd	1110011	CSRRW	rd = csr, if #rd != 0 csr = rs1
csr	rs1	010	rd	1110011	CSRRS	rd = csr, if #rd != 0 csr = csr   rs1, if that csr bit is writable and #rs1 != 0
csr	rs1	011	rd	1110011	CSRRC	rd = csr, if #rd != 0 csr = csr & (~rs1), if that csr bit is writable and #rs1 != 0
csr	uimm[4:0]	101	rd	1110011	CSRRWI	rd = csr, if #rd != 0 csr = uimm(zero-extend)
csr	uimm[4:0]	110	rd	1110011	CSRRSI	rd = csr, if #rd != 0 csr = csr   uimm(zero-extend), if that csr bit is writable and uimm != 0
csr	uimm[4:0]	111	rd	1110011	CSRRCI	rd = csr, if #rd != 0 csr = csr & (~uimm(zero- extend)), if that csr bit is writable and uimm != 0







## **New instructions - Description**

Description of instructions in privileged ISA

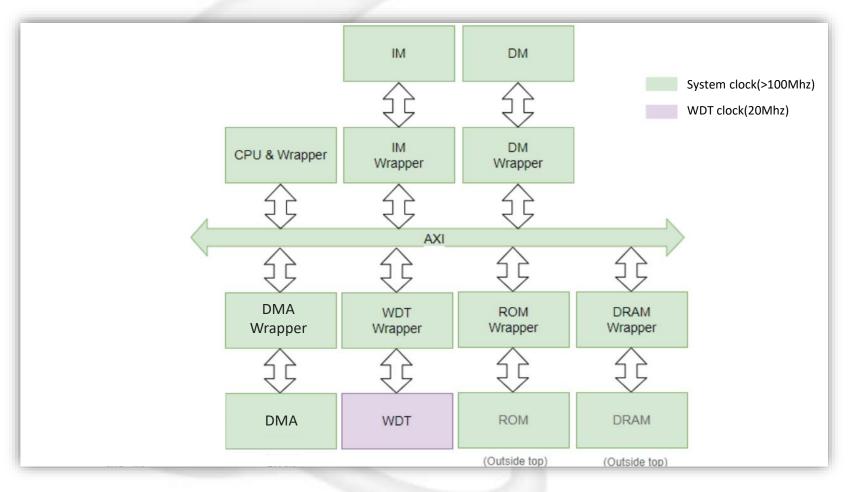
31	20	19 15	14 12	11 7	6 0		
imm	11:0]	rs1	funct3	rd	opcode	Mnemonic	Description
0011000	00010	00000	000	00000	1110011	MRET	Return from traps in Machine Mode

31	20	19	15	14	12	11	7	6 0		
imm[1	1:0]	rs1		fun	ct3	1	rd	opcode	Mnemonic	Description
0001000	00101	0000	00	00	00	00	0000	1110011	WFI	Wait for interrupt



#### **New IPs**

System Architecture in HW3







### **New IPs - Slave Configuration**

#### **Table: Master configuration**

NAME	Number
CPU_M0	Master 0
CPU_M1	Master 1
DMA	Master 2

#### **Table: Slave configuration**

NAME	Number	Start address	End address	Property
ROM	Slave 0	0x0000_0000	0x0000_1FFF	memory
IM	Slave 1	0x0001_0000	0x0001_FFFF	memory
DM	Slave 2	0x0002_0000	0x0002_FFFF	memory
DMA	Slave 3	0x1002_0000	0x1002_0400	I/O
WDT	Slave 4	0x1001_0000	0x1001_03FF	I/O
DRAM	Slave 5	0x2000_0000	0x201F_FFFF	memory





#### **New IPs - DRAM**

- Memory slower then SRAM, storing programs to load in HW3
- Excluded outside of top module, need to write FSM to control

	System signals					
	CK	input	1	System clock		
	RST	input	1	System reset (active high)		
			Memor	y ports		
	CSn	input	1	DRAM Chip Select		
	CSII	input	1	(active low)		
	WEn	input	4	DRAM Write Enable		
	VV EII	input	4	(active low)		
	RASn	input	1	DRAM Row Access Strobe		
	KASII	input		(active low)		
DRAM	CASn	input	1	DRAM Column Access Strobe		
	CASII			(active low)		
	A	input	11	DRAM Address input		
	D	input	32	DRAM data input		
	Q	output	32	DRAM data output		
	VALID	output	1	DRAM data output valid		
	Memory space					
	Memory_byte0	reg	8	Size: [0:2097151]		
	Memory_byte1	reg	8	Size: [0:2097151]		
	Memory_byte2	reg	8	Size: [0:2097151]		
	Memory_byte3	reg	8	Size: [0:2097151]		

See DRAM appendix for more details





#### New IPs - ROM

**ROM** 

- Read only memory, excluded outside of top module
- Used to store booting program

000		System s	signals				
CK	input	1	System clock				
	Memory ports						
DO	output	32	ROM data output				
OE	input	1	Output enable (active high)				
CS	input	1	Chip select (active high)				
A	input	12	ROM address input				
		Memory	Space				
Memory_byte0	reg	8	Size: [0:4095]				
Memory_byte1	reg	8	Size: [0:4095]				
Memory_byte2	reg	8	Size: [0:4095]				
Memory_byte3	reg	8	Size: [0:4095]				
,							





## **New IPs – DMA (1/2)**

DMA is a technology that allows hardware devices to read from and write to system memory directly, bypassing the CPU.

Module	Specification						
	Name	Signal	Bits	Function Explanation			
	clk	Input	1	System clock			
	rst	Input	1	System reset(active high)			
DNAA	DMAEN	Input	1	Enable the DMA			
DMA	DMASRC	Input	32	Source address of DMA			
	DMADST	Input	32	Destination address of DMA			
W	DMALEN	Input	32	Total length of the data			
	DMA_interrupt	Output	1	DMA interrupt			





## New IPs -DMA(2/2)

- During booting, the CPU first writes the source address, destination address, and data length into the corresponding registers of the DMA.
- After completing the register setup, the CPU will enable DMAEN, enter WFI, and the DMA will begin its data transfer operation.
- After the DMA completes the data transfer, it will stop functioning and send a DMA interrupt to the CPU.
- After completing the ISR, the booting program needs to set DMAEN to 0.

Address	Mapping
0X10020100	DMAEN
0X10020200	DMASRC
0X10020300	DMADST
0X10020400	DMALEN





## New IPs - Watch Dog Timer(1/2)

- □ Watch Dog timer is set by CPU and is used to count for a given time.
- Once time up, WDT will send timeout signal to interrupt CPU.
- CPU control signals come from "clk" domain, and counting register is in "clk2" domain.

Module	Specifications					
	Name	Signal	Bits	Function explanation		
	clk	input	1	System clock		
	rst	input	1	System reset (active high)		
	clk2	input	1	WDT clock		
WDT	rst2	input	1	WDT reset (active high)		
	WDEN	input	1	Enable the watchdog timer		
	WDLIVE	input	1	Restart the watchdog timer		
	WTOCNT	input	32	Watchdog timeout count		
	WTO	output	1	watchdog timeout		







## New IPs - Watch Dog Timer(2/2)

- CPU can write WTOCNT to set timer value.
- ☐ When WDT is enabled(WDEN = 1), the "clk2" domain counter starts counting.
- □ When WDT get restart signal (WDLIVE = 1), it will reset counter to 0.
- ☐ When counter value exceeds WTOCNT, WDT will assert timeout interrupt(WTO = 1).
- Write non-zero value in 0x1001\_0100 or 0x1001\_0200 to enable WDEN or WDLIVE

Address	Mapping
0x1001_0100	WDEN
0x1001_0200	WDLIVE
0x1001_0300	WTOCNT





## Problem 2 – Software & Bootloader

Specification





## **Verification(1/7)**

#### Software list:

- prog0
  - → 測試45個instruction (助教提供)
- prog1
  - Sort algorithm of half-word
- prog2
  - Gray scale
- Prog3
  - → Timer interrupt (助教提供)
- □ Prog4
  - → Timer interrupt with clock uncertainty (助教提供)

#### Firmware:

- Bootloader
  - → Load software from DRAM to IM/DM





## **Verification (2/7)**

#### **Bootloader introduction:**

- The first program that CPU will execute after reset state
- Used to load main programs or OS to faster memories.

#### Booting in HW3

- The booting program is stored in ROM
- Booting program will move instr & data from DRAM to IM and DM

```
extern unsigned int dram i start;
extern unsigned int dram i end;
extern unsigned int imem start;
extern unsigned int
                      sdata start;
extern unsigned int
                      sdata end;
extern unsigned int
                      sdata paddr start;
extern unsigned int
                      data start;
extern unsigned int
                      data end;
extern unsigned int
                      data paddr start;
```

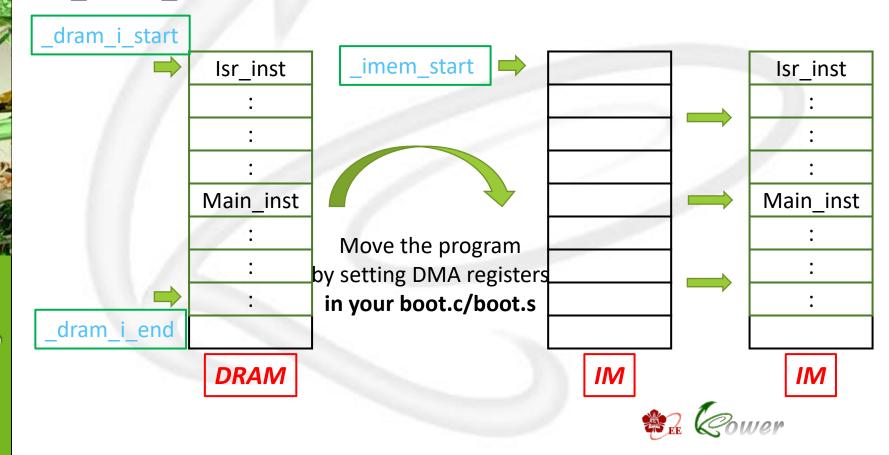




## Verification (3/7)

Booting(1/3) - Move instructions from DRAM to IM:

- \_dram\_i\_start = instruction start address in DRAM.
- \_dram\_i\_end = instruction end address in DRAM.
- \_imem\_start = instruction start address in IM

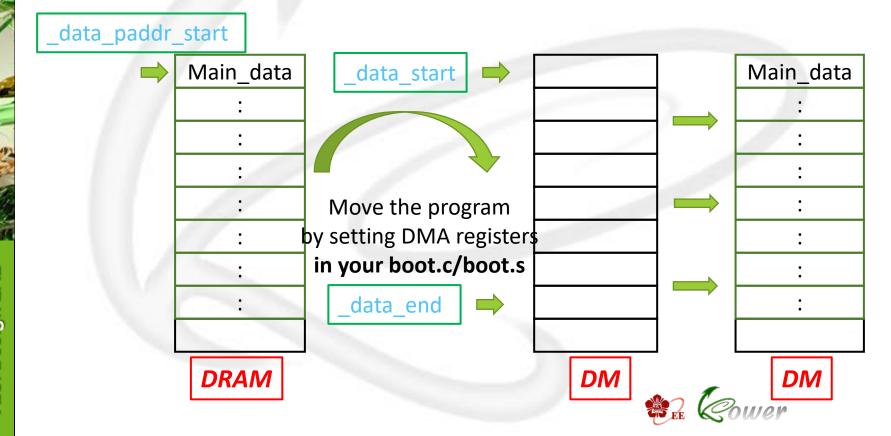




## **Verification (4/7)**

Booting(2/3) - Move data from DRAM to DM:

- \_data\_start = main data start address in DM.
- \_data\_end = main data end address in DM.
- \_data\_paddr\_start = main data start address in DRAM

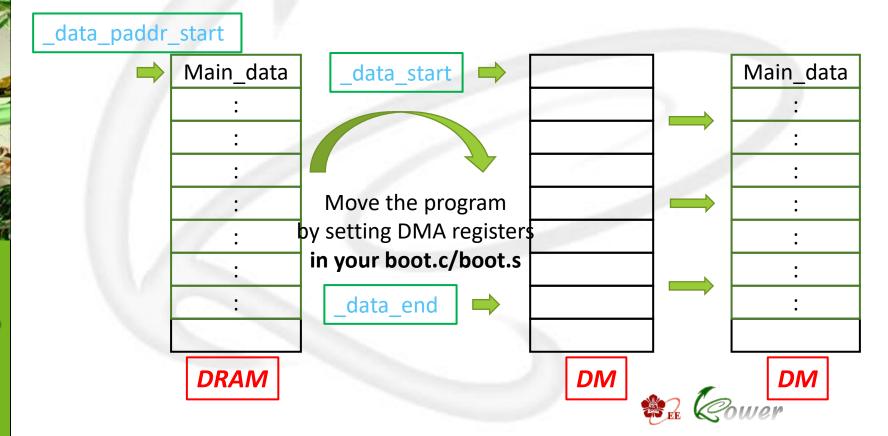




## **Verification (5/7)**

Booting(3/3) - Move data from DRAM to DM:

- \_sdata\_start = main data start address in DM.
- \_sdata\_end = main data end address in DM.
- \_sdata\_paddr\_start = main data start address in DRAM





## **Verification(6/7)**

Grayscale in prog2



```
Advanced VLSI System Design(Graduate Level) Fall 2024
...

12

1f

25

12

1y

y = 0.11*b + 0.59*g + 0.3*r

= 0.11*25 + 0.59*g + 0.3*r

= 0.11*b + 0.59*g + 0.3*r

= 0.11*b + 0.59*g + 0.3*r

y

y

y = 0.11*b + 0.59*g + 0.3*r

= 0.11*25 + 0.59*1f + 0.3* 12
```

M Size of BMP file (byte) The number of bits per pixel

Header

(54 bytes)

BMP檔案解析參考資料:

Header

(54 bytes)

## **Verification(7/7)**

- Prog3 & prog4 will test timer interrupt
- Prog4 will test CDC functions
  - Prog3: "clk" domain is not delayed
  - Prog4: "clk" domain is delayed
- WDT interrupt:
  - CPU will start executing program and enter dead loop
  - When WDT assert timer interrupt, CPU will be reset
  - After CPU is reset, it will periodically reset timer



Timer interrupt will reset CPU, make CPU reboot







## Problem 3 – Spyglass CDC check





## Spyglass CDC check

- Use Spyglass to do CDC check on the SoC
- Steps:
  - 1. Modify Spyglass.sgdc
    - Set initial value of CDC flip-flops, fifo memory...etc
  - Open GUI
    - Type "make spyglass" in HW3
  - 3. Modify SoC according to the advice given by Spyglass
    - The detail description of warnings and errors can be found in Spyglass CDC Rules reference guide.

```
Message

☐ Message Tree ( Total: 91, Displayed: 88, Waived: 3 )
☐ Info (79)
☐ Warning (9)
☐ ##1 Ac_cdc01a (3): Checks data loss for multi-flop or sync cell or qualifier synchronized clock domain crossings
☐ ##1 Ac_conv01 (1): Checks sequential convergence of same-domain signals synchronized in the same destination domain
☐ ##1 M Ac_datahold01a (5): Checks the functional synchronization of synchronized data crossings
```

See Spyglass CDC appendix for more details





## **Report Requirements**

- Proper explanation of your design is required for full credits.
- Block diagrams shall be drawn to depict your designs.
- Show your screenshots of the waveforms and the simulation results on the terminal for the different test cases in your report and illustrate the correctness of your results
- Explain your codes of prog1, prog2.
- Explain your codes of boot.c.
- Report your Superlint coverage





## **Report Requirements**

- Report and show screenshots of your prog0 to prog4 simulation time after synthesis and total cell area of your design.
- Explain how the interrupt mechanism work
- Show your screenshots of the Spyglass CDC reports and explain why your CDC circuit can work correctly





## **Submission rule**





## Report

- □ 請使用附在檔案內的Submission Cover
- □ 請勿將code貼在.docx內 (program的程式可截圖 說明)
  - → 請將.sv包在壓縮檔內,不可截圖於.docx中
- 需要Summary及Lessons learned(Summary table請放在第二頁,清楚列出有完成以及沒完成的部分)
- □ 若兩人為一組,須寫出貢獻度(貢獻度請放第二 頁)
  - Ex: A(N26071234) 55%, B(N26075678) 45%
  - → Total 100%
  - → 自己一組則不用寫





## **Specification**

□ Module name須符合下表要求

Category	Name		
	File	Module	Instance
RTL	top.sv	top	TOP
Gate-Level	top_syn.v	top	TOP
RTL	SRAM_wrapper.sv	SRAM_wrapper	IM1
RTL	SRAM_wrapper.sv	SRAM_wrapper	DM1
RTL	TS1N16ADFPCLLLVTA512X45M4SWSH OD.sv	SRAM	i_SRAM
Behavior	ROM.v	ROM	i_ROM
Behavior	DRAM.v	DRAM	i_DRAM

□ 需按照要求命名,以免testbench抓不到正確的名稱



## 繳交檔案 (1/2)

- □ 依照檔案結構壓縮成 ".tar" 格式
  - → 在Homework主資料夾(N260XXXXX)使用make tar產生的tar檔即可符合要求
- □ 檔案結構請依照作業說明
- □ 請勿附上檔案結構內未要求繳交的檔案
  - → 在Homework主資料夾(N260XXXXX)使用make clean即可刪除不必要的檔案
- □ 請務必確認繳交檔案可以在SoC實驗室的工作站下compile,且功能正常
- □ 無法compile將直接以0分計算
- □ 請勿使用generator產生code再修改
- □ 禁止抄襲



## 繳交檔案 (2/2)

- □ 一組只需一個人上傳作業到Moodle
  - → 兩人以上都上傳會斟酌扣分
- □ 壓縮檔、主資料夾名稱、Report名稱都要為上傳者的學號, StudentID檔案則個別寫上組員的學號。
  - → Ex: A(N26101234)負責上傳,組員為B(N26105678)
  - → N26101234.tar (壓縮檔) N26101234 (主資料夾) N26101234.docx (Report, Cover寫上兩者的學號)





## 繳交期限

- □ 2024/12/4 (三) 14:00前上傳
  - → 不接受遲交,請務必注意時間
  - → Moodle只會留存你最後一次上傳的檔案,檔名只要是「N26XXXXXXX.tar」即可,不需要加上版本號





## Thanks for your participation and attendance!!

