

Divide and Conquer: Problem 1

We have a vector $V [1..N]$ formed by integers, so that all of them are different, and that they are ordered in an increasing way. It is said that a vector of these characteristics is coincident if it has at least one position such that it is equal to the value that contains the vector in that position. For example, in the vector

1	2	3	4	5	6	7	8
-14	-6	3	6	16	18	27	43

it can be seen that $V [3] = 3$; therefore, this vector is coincident.

Design a Divide and Conquer algorithm that determines in an order of efficiency not greater than $O(\log n)$ if a vector is coincident, receiving as data the vector and its size.

Divide and Conquer: Problem 1

```

const MaxSize= ...
type vector = array[1.. MaxSize] of integer
fun IsCoincident(I/O xi, xf: integer; I/O v: vector) ret e: boolean
var m: integer
    e ← false
    if ( xi == xf ) then                                //Basic case
        if ( v(xi) == xi) then e ← true elif
    else
        m = ( xi + xf ) % 2
        if ( v(m) == m) then e ← true
        else
            if ( v(m) < m) then                          //Upper half
                xi = m
            else                                          //Lower half
                xf = m
            elif
                e = IsCoincident (xi, xf, v)
        elif
    return e
ffun

```

(1)
(1)
~~(1+1)~~
(1)
(1+1)
(1)
(1)
//Lower half
~~(1)~~
 $T\left(\frac{n}{2}\right)$
(1)

$$T(n) = 7 + T\left(\frac{n}{2}\right)$$

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fun <i>IsCoincident</i> (I/O xi, xf: integer ; I/O v: vector) ret e: boolean		
var m: integer		
e ← false		(1)
if (xi == xf) then //Basic case		(1)
if (v(xi) == xi) then e ← true elif		(1+1)
else		
m = (xi + xf) % 2		(1)
if (v(m) == m) then e ← true		(1+1)
else		
if (v(m) < m) then //Upper half		(1)
xi = m+1		(1)
else //Lower half		
xf = m-1		(1)
elif		
e = <i>IsCoincident</i> (xi, xf, v)		$T\left(\frac{n}{2}\right)$
elif		
return e		(1)
ffun		

$$T(n) = 7 + T\left(\frac{n}{2}\right)$$

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- ▶ Step 2: Polynomial equation

$$n = 2^k \Leftrightarrow k = \log_2 n \rightarrow x^k - x^{k-1} = 7$$

- ▶ Step 3: Solution of the homogeneous part

$$x^{k-1} \cdot (x - 1) = 0 \rightarrow \text{Raíz: } (x - 1)^1$$

- ▶ Step 4: Solution of the non homogeneous part

$$7 = 7k^0 \cdot 1^k \rightarrow \text{Raíz: } (x - 1)^{0+1}$$

- ▶ Step 5: Combination of roots

$$(x - 1)^2$$

$$T(k) = c_1 + c_2 \cdot k \rightarrow T(n) = c_1 + c_2 \cdot \log n$$

$$T(n) \in \mathcal{O}(\log n)$$

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$$T(n) = 7 + T\left(\frac{n}{2}\right)$$

- ▶ Is a typical Divide and Conquer expression with $p=0$, $b=2$ and $k=1$. Then:

$$k = b^p$$

- ▶ Therefore:

$$T(n) \in \mathcal{O}(n^p \log n) \in \mathcal{O}(\log n)$$