



# UNIT 1

INTRODUCTION TO THE ALGORITHMIC

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## EXERCISE 3

Analyze the efficiency of the following code:

```

fun Calculo(x,y,z: entero) dev valor:entero
var i,j,t: entero
valor ← 0
Desde i ← x hasta y Hacer valor ← valor + i fdesde
si (valor ÷ (x+y)) <= 1 entonces Devolver z
si no
    t ← x + ((y-x) ÷ 2)  { ÷ es la división entera }
    Desde i ← x hasta y Hacer
        Desde j ← (3*x) hasta (3*y) Hacer
            valor ← valor + Mínimo(i,j)
        fdesde
    fdesde
valor ← valor + 4*Calculo(t,y,valor)
Devolver valor
fsi
ffun
    
```

$$n = y - x$$

We're using  $y - x$  as the value of  $n$  because the amount of loop iterations is  $y - x$ , except the inner one, which is  $3n$ .

$$T(n) = 1 + \sum_{i=1}^n \sum_{i=1}^{3n} 1 + T\left(\frac{n}{2}\right) + 1$$

$$y - t = y - x - \left(\frac{y-x}{2}\right) = \frac{-y+x+2y-2x}{2} = \frac{y-x}{2} = \frac{n}{2}$$

After seeing that  $y - t$  is equal to  $\frac{n}{2}$ , the recursive call can be replaced by  $T(\frac{n}{2})$ .

$$n = 2^k \quad T(n) = x^k$$

$$T(n) = 3n^2 + n + 4 + T\left(\frac{n}{2}\right)$$

$$x^k = 3 * 2^{2k} + 2^k + 4 + x^{k-1}$$

$$x^k - x^{k-1} = 3 * 2^{2k} + 2^k + 4$$

Homogeneous equation:

$$x^{k-1}(x - 1) = 0 \quad x^H = A * 1^k$$

$$\text{roots: } x = 1$$

Particular equation:

$$x^{k-1}(x - 1) = 3 * 2^{2k} + 2^k + 4 \quad x^P = B * 4^k + C * 2^k + D * k * 1^k$$

$$\text{roots: } x = 4$$

$$x = 2$$

$$x = 1$$

Finally:

$$x = A + 4^k B + 2^k C + kD \quad n = 2^k \rightarrow k = \log_2(n)$$

$$T(n) = x = A + Bn^2 + Cn + D \log_2(n)$$

Complexity of **O(n<sup>2</sup>)**

## EXERCISE 5

Program a function to determine if a number received as parameter is prime. Analyze the efficiency and complexity.

```
1  def isPrime(num):
2      b = True
3      for i in range(2, num):
4          if num % i == 0:
5              b = False
6      return b
7
8  print(isPrime(12))
```

$$T(n) = 1 + \sum_{i=1}^n 1 + \max\{1,0\} = 1 + \sum_{i=1}^n 1 + 1 = 1 + 2n$$

Complexity of **O(n)**

## EXERCISE 8

Program a recursive procedure to obtain the inverse number of a given one.

Example:  $627 \rightarrow 726$

Analyze the efficiency and complexity.

```
1  def inverseNumber(n):
2      numberInv = ''
3      if (int(n/10) == 0):
4          numero = str(n%10);
5          numberInv += numero
6      else:
7          numero = str(n%10);
8          numberInv += numero + inverseNumber(int(n/10))
9      return numberInv
10
11 print(inverseNumber(678))
```

After seeing that  $n/10$  is equal to  $n-1$ , the recursive call can be replaced by  $T(n-1)$ .

$$T(n) = 1 + 1 + \max\{1, 1 + T(n-1)\} = 2 + 1 + T(n-1)$$

$$T(n) = 3 + T(n-1)$$

$$T(n) = x^n$$

$$x^n = 3 + x^{n-1}$$

$$x^n - x^{n-1} = 3$$

$$x^{n-1}(x-1) = 3$$

Homogeneous equation:

$$x^{k-1}(x-1) = 0 \quad x^H = A * 1^k$$

roots:  $x = 1$

Particular equation:

$$x^{k-1}(x-1) = 3 \quad x^p = B * n * 1^k$$

roots:  $x = 1$

We multiply by  $n$  because the value in both roots is 1

Finally:

$$T(n) = x = A + Bn$$

Complexity of **O(n)**

## EXERCISE 6: Extra Exercise

Program a function to determine if a number received as parameter is perfect. Analyze the efficiency and complexity.

```
1  def isPerfect(num):
2      n = 0
3      for i in range(1, num):
4          if num % i == 0:
5              n = n + i
6      return n == num
7
8  print(isPerfect(6))
```

$$T(n) = 1 + \sum_{i=1}^n 1 + \max\{1, 0\} = 1 + \sum_{i=1}^n 1 + 1 = 1 + 2n$$

Complexity of **O(n)**

## EXERCISE 7: Extra Exercise

Write a program which ask a positive number to the user (N) and obtain how many prime numbers there are between 1 and that number N, and how many perfects between 1 and N. Analyze the efficiency and complexity.

```

1  def isPrime(num):
2      b = True
3      for i in range(2, num):
4          if num % i == 0:
5              b = False
6      return b
7
8  def isPerfect(num):
9      n = 0
10     for i in range(1, num):
11         if num % i == 0:
12             n = n + i
13     return n == num
14
15
16  n = int(input("Write down a number: "))
17  nPrimes = 0
18  nPerfects = 0
19  for i in range (1, n):
20      if (isPrime(i) == True):
21          nPrimes = nPrimes + 1
22      if (isPerfect(i) == True):
23          nPerfects = nPerfects + 1
24  print("Between 1 and", n, ":\nNumber of Perfects: ", nPerfects, "\nNumber of Primes: ", nPrimes)
25

```

$$T(n) = 1 + 1 + 1 + \underbrace{\sum_{i=1}^n (1 + (\sum_{i=1}^n 1 + 1))}_{\text{isPrime(i)}} + \underbrace{1 + 1 + (\sum_{i=1}^n 1 + 1) + 1}_{\text{isPerfect(i)}} + 1$$

$$T(n) = 4 + \sum_{i=1}^n (4 + 2n + 2n)$$

$$T(n) = 4 + \sum_{i=1}^n (4 + 4n)$$

$$T(n) = 4n + 4n^2 + 4$$

Complexity of **O(n<sup>2</sup>)**