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# MIPS ASSEMBLY PROGRAMMING LANGUAGE PART V

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# SYSCALL. EXITING PROGRAMS

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- The 'syscall' instruction suspends the execution of your program and transfers control to the operating system.
- The operating system then looks at the contents of register \$v0 to determine what it is that your program is asking it to do.
- Here's the proper way of telling MARS to stop executing your program:

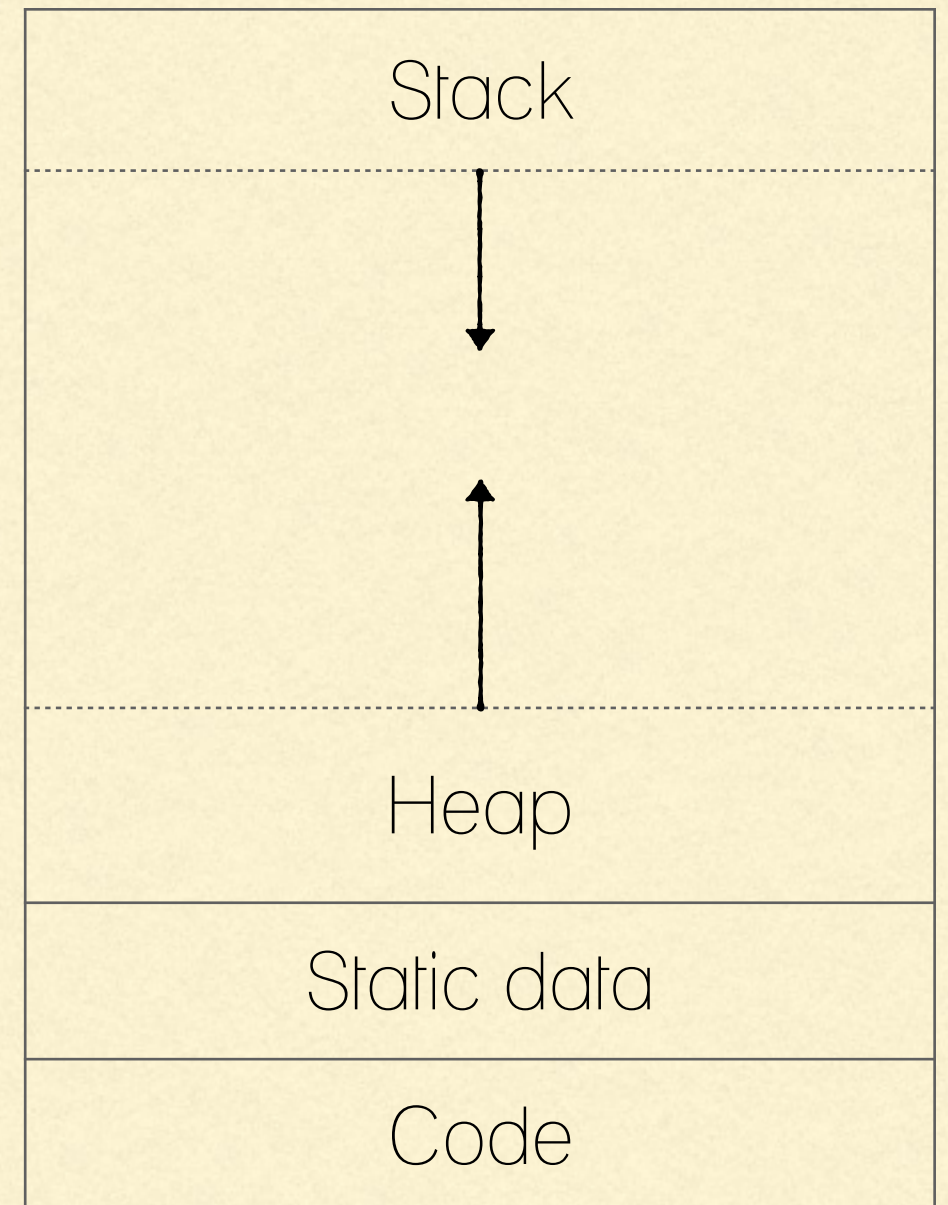
```
addi $v0, $zero, 10
```

```
syscall
```



# C/MIPS MEMORY MANAGEMENT

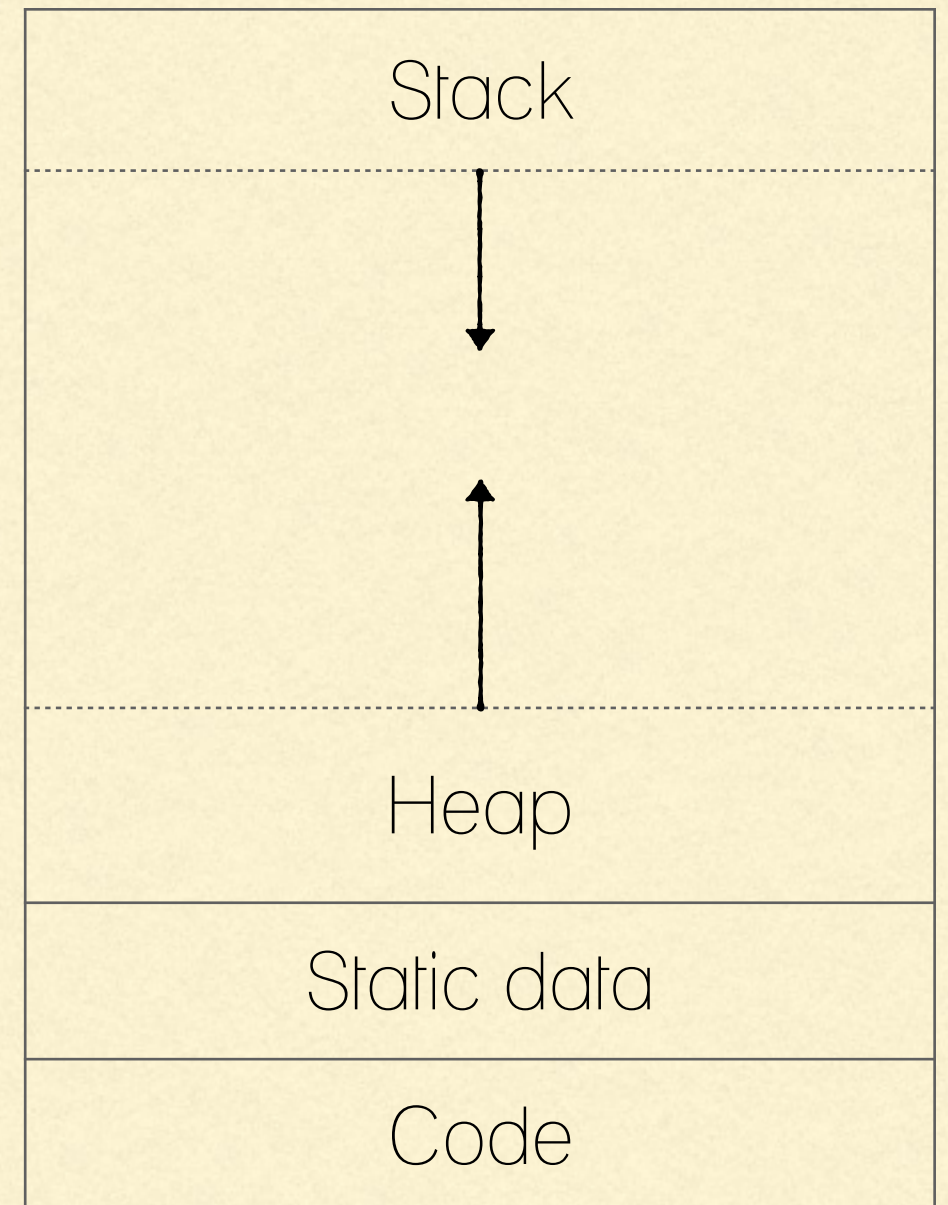
- Programs's address space contains 4 regions;
  - Stack; local variables inside functions, grows downward
  - Heap; space requested for dynamic data
  - Static data; variables declared outside functions, does not grow or shrink
  - Code; loaded when program starts — does not change





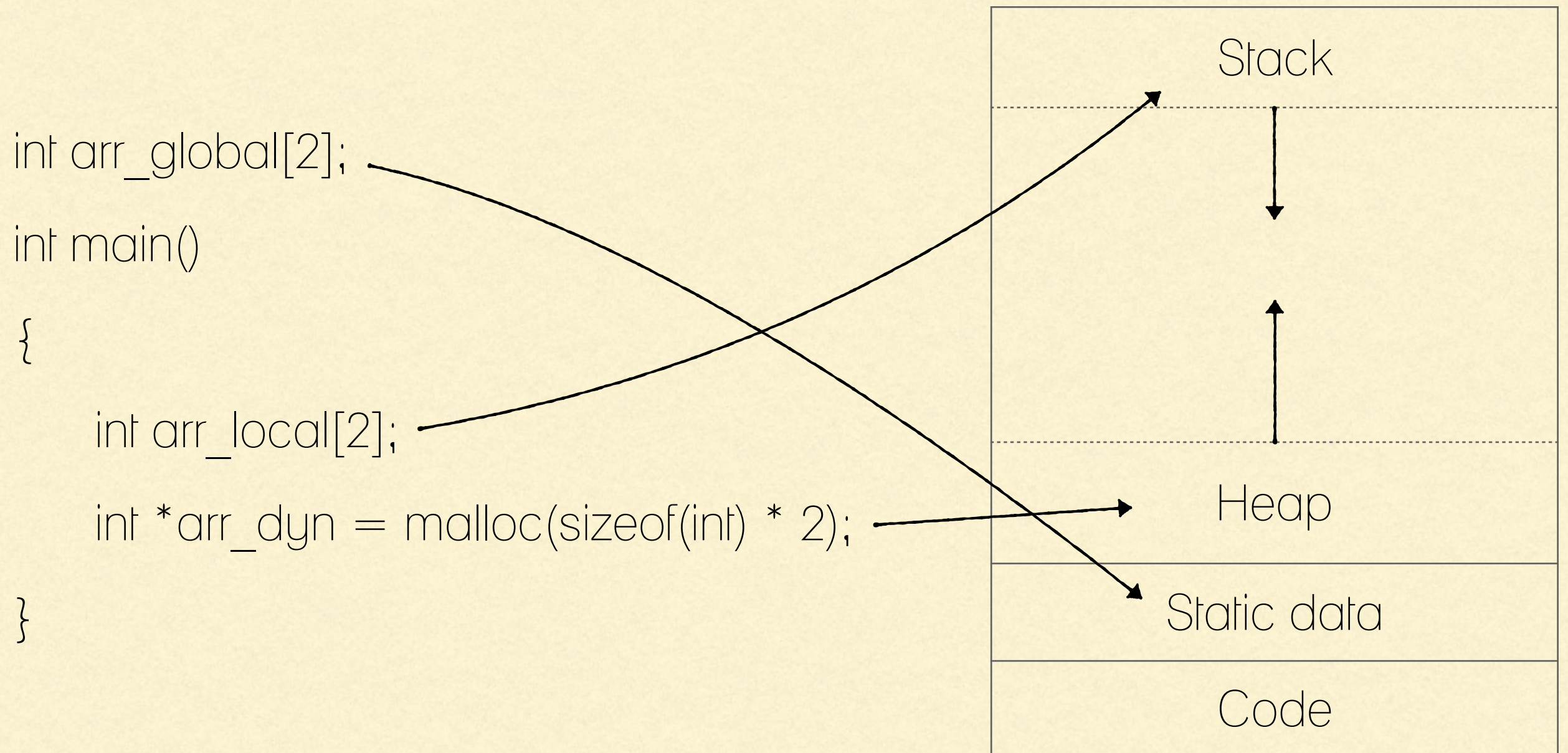
# C/MIPS MEMORY MANAGEMENT

```
int arr_global[2];  
  
int main()  
{  
    int arr_local[2];  
    int *arr_dyn = malloc(sizeof(int) * 2);  
}
```





# C/MIPS MEMORY MANAGEMENT





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# MIPS ASSEMBLER DIRECTIVES

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- Directives are commands that are part of the assembler syntax but are not related to the instruction set.
  - All assembler directives begin with a period (.)
- `.data` and `.text` are both directives to the assembler.
  - `.data` tells the assembler that the upcoming section is considered data (static data section, or global data)
  - `.text` tells the assembler that the upcoming section is considered assembly language instructions



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# DATA DECLARATION

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- A single declaration consists of:
  1. a label (identifier), followed by a colon. Each label corresponds to a unique address in memory, which the assembler determines.
  2. a storage type. MIPS has a weak sense of type, but it's the closest term to describe what's going on
  3. data values. Data values must be of the correct type (see next slide; 'var1' is numeric variable, 'array1' is an array of characters)



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# DATA DECLARATION; EXAMPLES

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- `var1: .word 3`
  - Create a single integer variable and assign it the value 3
- `array1: .byte 'a', 'b'`
  - Create a 2-element character array
- `array2: .space 40`
  - Allocate 40 consecutive bytes, with store uninitialized (could be used as a 40-element character array, or as a 10-element integer array)



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# ALLOCATING SPACE IN STATIC DATA

---

```
int arr[] = [2, 5, 10];
```

```
void main()
```

```
{
```

What's here?

```
}
```

```
.data
```

```
arr: .word 2, 5, 10
```

```
.text
```

```
addi $s0, $zero, 4
```

```
addi $s1, $zero, 10
```

```
sw $s0, arr($zero)
```

```
addi $t0, $zero, 4
```

```
sw $s1, arr($t0)
```

```
# Static data section
```

```
# Space for 3 words
```

```
# Instruction section
```



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# ALLOCATING SPACE IN STATIC DATA

---

```
int arr[] = [2, 5, 10];
```

```
void main()
```

```
{
```

```
    arr[0] = 4;
```

```
    arr[1] = 10;
```

```
}
```

```
.data
```

```
    arr: .word 2, 5, 10
```

```
.text
```

```
    addi $s0, $zero, 4
```

```
    addi $s1, $zero, 10
```

```
    sw $s0, arr($zero)
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```
    addi $t0, $zero, 4
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    sw $s1, arr($t0)
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# Static data section
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# Instruction section
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# CALLING A FUNCTION

---

1. Put parameters in a place where function can access them
2. Transfer control to function
3. Acquire (local) storage resources needed for function
4. Perform desired task of the function
5. Put result value in a place where calling program can access it and restore any registers you used
6. Return control to point of origin, since a function can be called from several points in a program



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# MIPS FUNCTION CALL CONVENTIONS

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1. Registers faster than memory, so use them
2. \$a0 to \$a3; four argument registers to pass parameters
3. \$v0–\$v1; two value registers to return values
4. \$ra; one return address register to return to the point of origin



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# INSTRUCTION SUPPORT FOR FUNCTIONS

---

In main(), we call `sum(a, b); /* a in $s0, b in $s1 */`

```
int sum(int x, int y)
{
    return x + y;
}
```

---

Address (shown in decimal)

0996  
1000  
1004  
1008  
1012  
1016

...  
2000  
2004  
2008

In MIPS, all instructions are 4 bytes, and stored in memory just like data. So here we show the addresses of where the programs are stored.



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0996      main:

1000      ...

1004      ...

1008      ...

1012      ...

1016      ...

...

2000      sum:

2004      ...

2008      ...



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0996      main:

1000      ...

1004      ...

1008      ...

1012      ...

1016      ...

...

2000      sum:

2004      ...

2008      ...



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Address (shown in decimal)

0996	main:	
1000	add \$a0, \$s0, \$zero	# \$a0 = \$s0 (a)
1004	add \$a1, \$s1, \$zero	# \$a1 = \$s1 (b)
1008	...	
1012	...	
1016	...	
...		
2000	sum:	
2004	...	
2008	...	



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2000	sum:	
2004	...	
2008	...	



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...		
2000	sum:	
2004	add \$v0, \$a0, \$a1	
2008	...	



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...		
2000	sum:	
2004	add \$v0, \$a0, \$a1	
2008	...	



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```
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---

Address (shown in decimal)

0996	main:	
1000	add \$a0, \$s0, \$zero	# \$a0 = \$s0 (a)
1004	add \$a1, \$s1, \$zero	# \$a1 = \$s1 (b)
1008	addi \$ra, \$zero, 1016	
1012	j sum	# jump to sum
1016		
...		
2000	sum:	
2004	add \$v0, \$a0, \$a1	
2008	jr \$ra	



# INSTRUCTION SUPPORT FOR FUNCTIONS

In main(), we call sum(a, b); /\* a in \$s0, b in \$s1 \*/

```
int sum(int x, int y)
{
    return x + y;
}
```

Question;  
Why use jr here? Why not use j (num)?

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1008	addi \$ra, \$zero, 1016	
1012	j sum	# jump to sum
1016		
...		
2000	sum:	
2004	add \$v0, \$a0, \$a1	
2008	jr \$ra	



---

# INSTRUCTION SUPPORT FOR FUNCTIONS

---

- Single instruction to jump and save return address:
  - jump and link (jal)
  - Before:
    - 1008 addi \$ra, \$zero, 1016    # \$ra = 1016
    - 1012 j sum    # goto sum
    - 1016 ...
  - Now:
    - 1008 jal sum    # \$ra = 1012, goto sum



# INSTRUCTION SUPPORT FOR FUNCTIONS

- Single instruction to jump and save return address:

- jump and link (jal)

- Before:

- 1008 addi \$ra, \$zero, 1016    # \$ra = 1016

- 1012 j sum    # goto sum

- 1016 ...

Why 1012 not 1016?

- Now:

- 1008 jal sum    # \$ra = 1012, goto sum