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CS 314 – Principles of Programming Languages

Spring, 2016

Midterm Exam 1 - Answers

- Please write your name and net ID readably above.
- **Do not open this exam** until everyone has an exam and the instructor tells you to begin.
- There are 4 pages in this exam, including this one. Make sure you have them all.
- This exam is closed book closed notes closed electronics.
- You must put your cellphone, tablet, iPod, or other electronic devices in a backpack, etc, and leave it out of reach. The only exception is that you can use a watch that only has time-related functions (e.g. not a calculator watch, not a "smart watch").
- Write clearly if we can't read or can't find your answer your, your answer is wrong.
- Make clear what is your answer versus intermediate work.

I	/20
II	/20
III	/20
IV	/20
V	/20
VI	/20
Total	/120

I. Suppose we define an "id-line" as a name, followed by '%', followed by a domain, where a "name" is a string of one or more letters and digits and a "domain" is a sequence of one or more names separated by dots.

For instance, the following strings <u>are</u> id-lines:

```
19 chaplin 20 % films.funny.haha
adog % pets
```

but these are <u>not</u> (b) indicates a space character):

```
cat% (no domain)
cat% meow. (no name after the dot)
dog% petbit (space is not a letter or a digit)
c&d% fight (& is not a letter or a digit)
```

Finish the following grammar so that it defines the language id-lines. Be sure to <u>underline</u> terminal symbols. You may use symbols { } and [] from EBNF if you want to. You may add extra lines at the bottom if you want to.

```
ID-LINE \rightarrow NAME \frac{\%}{2} DOMAIN

NAME \rightarrow ALPHANUM | NAME ALPHANUM

DOMAIN \rightarrow NAME | NAME \underline{.} DOMAIN

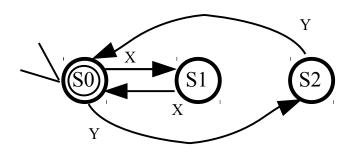
ALPHANUM \rightarrow \underline{a} | \underline{b} | \dots | \underline{z} | \underline{0} | \underline{1} | \underline{2} | \underline{3} | \underline{4} | \underline{5} | \underline{6} | \underline{7} | \underline{8} | \underline{9} | \underline{0}
```

II. Define a "simple name" as a sequence of one or more 'a's and/or 'b's and define a "simple domain" as a sequence of one or more simple names separated by '.'s. Write a Regular Expression that defines the language of simple domains. Note that these are not id-lines and so will not have any '%' characters.

```
(a \mid b)+((a \mid b)+)* or many equivalent REs
```

III.

Consider the FA defined by the diagram below:



State S0 is the start state and also the only accepting state.

- A. Is this FSA deterministic or non-deterministic? Why? deterministic: no epsilon moves and for every state/input char pair there is at most one next state
- B. For each if the following strings, circle Yes if it is accepted by the FSA above and No if it is not:
 - i. XXYY
- No

ii. YYXX

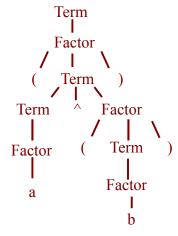
- No
- iii. X Y Y X Yes
- No
- iv. Y X X Y Yes
- No

- v. XXXX
-))
- III. Consider the following grammar, G1:

(Terminals are <u>underlined</u>. Term is the start symbol.)

Term => Term
$$\underline{v}$$
 Term | Term $\hat{}$ Factor | Factor Factor => $\underline{a} \mid \underline{b} \mid \underline{c} \mid (\text{Term})$

A. Draw a parse tree to show that (a^(b)) is in the language of G1



B. Show a string that is in the language of G1 and whose parse is ambiguous. (You do **not** have to show any parse trees, just the string,)

a v b v c (and many others)

For the Scheme questions below, all repetition must be done by recursion (including the recursion implicit in functions like map and assoc). You **may** write and use additional functions if you wish. You may **not** use do or any function whose name ends in '!', e.g. you may not use set!. You **may** use any other function in R5RS Scheme including the following:

Expression	Value	Expression	Value
(map sqrt '(9 1 4))	(3 1 2)	(reverse '(a (b c) d))	(d (b c) a)
(member 'a '(b c a d a))	(a d a)	(list (+ 2 3) '(a))	(5 (a))
(member 'x '(b c a d))	#f	(cons (+ 2 3) '(a))	(5 a)
(assoc 'x '((a b) (x y)(q r)))	(x y)	(append '(a b) '(c d))	(a b c d)

(null? x) is true if x is the empty list (), (eq? x y) is like Java's x = y,

```
(-x y) is x - y, similarly for + and *
```

Many equivalent answers for the following questions

(cons x (n-repeats x (- n 1))))

IV. Define the function n-repeats. (n-repeats x n) returns a list of length n, all of whose elements are x. E.g., (n-repeats '(a b) 3) should return ((a b)(a b)(a b)). (define (n-repeats x n) (if (= n 0) '())

Define the function fact-tr. (fact-tr n) is n factorial. Fact-tr and any helper functions **must be tail-recursive** if they are recursive at all.

```
(define (fact-tr n)

(fact-tr-h n 1))

(define (fact-tr-h n so-far)

(if (= n 0)

so-far

(fact-tr-h (- n 1)(* so-far n))))
```

Define (map-pairs fn lst), where fn is a 2-argument function, e.g. +, and lst is a list of 2-element sub-lists, e.g., ((1 2) (3 4)(5 6)). map-pair calls the function on each sub-list and returns a list of the results. E.g., (map-pairs + '((1 2)(3 4)(5 6))) returns the list (3 7 11). Note that there are correct answers that call map and also those that do not. Either way is ok. NOT calling map is probably easier.

```
(define (map-pairs fn lst)

(if (null? lst) '( ) ....

(cons (fn (caar lst) (cadar lst)) ← OR → (cons (apply fn (car lst))

(map-pairs1 fn (cdr lst)))) ....

OR (define (map-pairs2 fn lst)

(map fn

(map car lst)

(map cadr lst)))
```