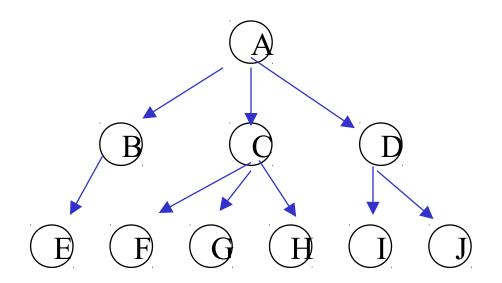
Computer Science 112 Data Structures

Lecture 14:

Trees

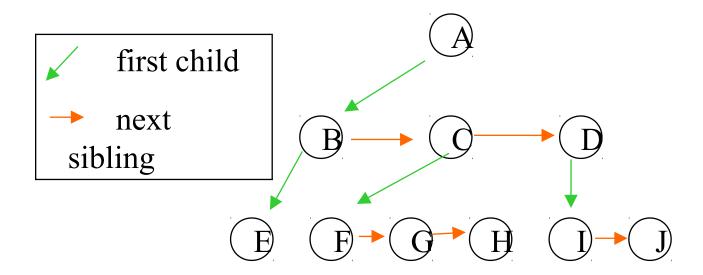
General Trees

- Each node has an arbitrary number of children
- Problem: representation of a node



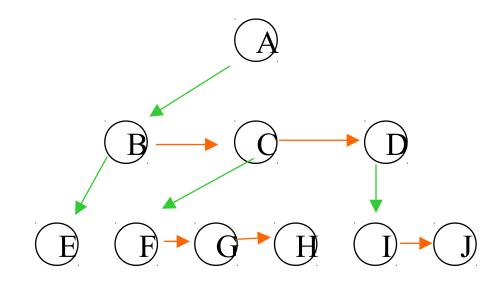
General Trees

- Each node has an arbitrary number of children
- Problem: representation
- Solution: linked list of children



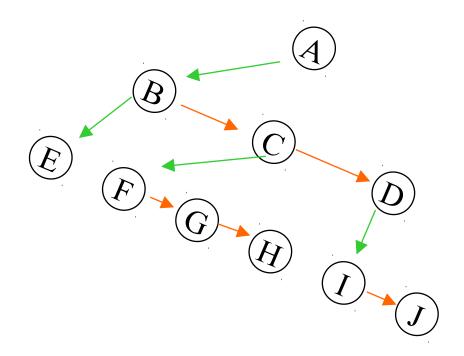
General Tree as Binary

- First child <=> Left child
- Next sib <=> Right child



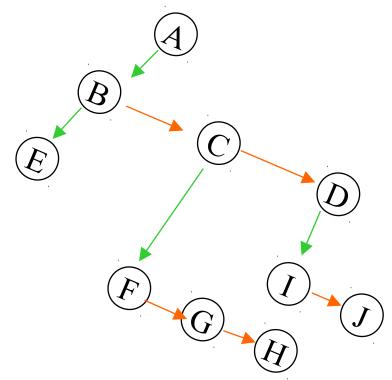
General Tree as Binary

- First child <=> Left child
- Next sib <=> Right child



General Tree as Binary

- First child <=> Left child
- Next sib <=> Right child

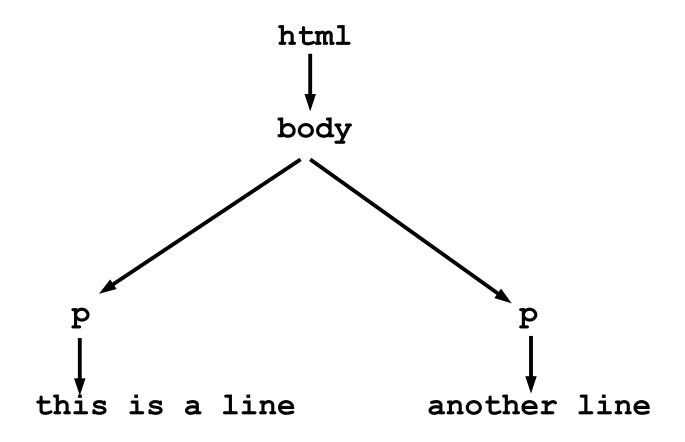


6

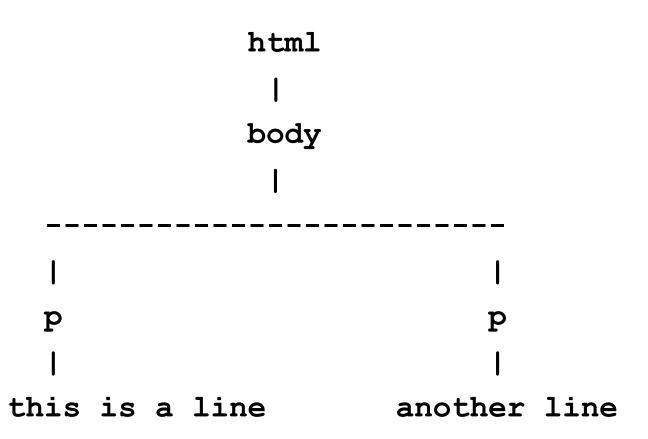
DOM tree in assignment

```
<html>
<body>
>
this is a line
>
another line
</body>
</html>
```

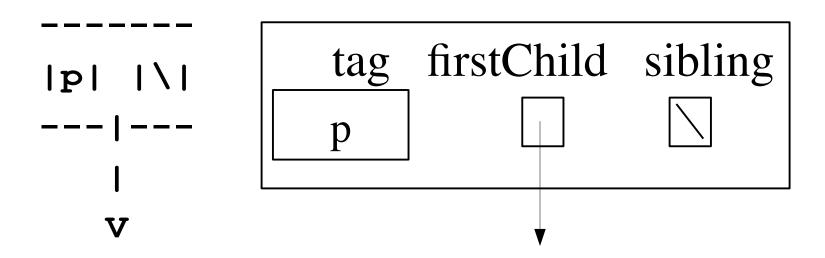
The tree as we think about it



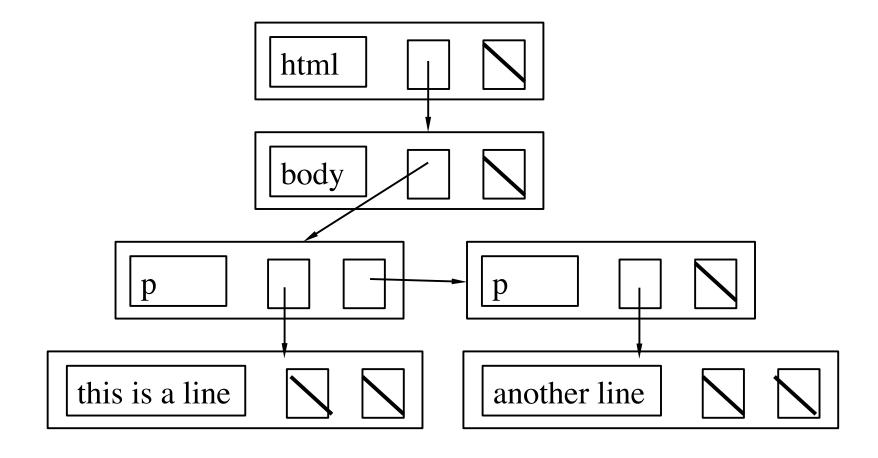
The tree as we drew it



TagNodes

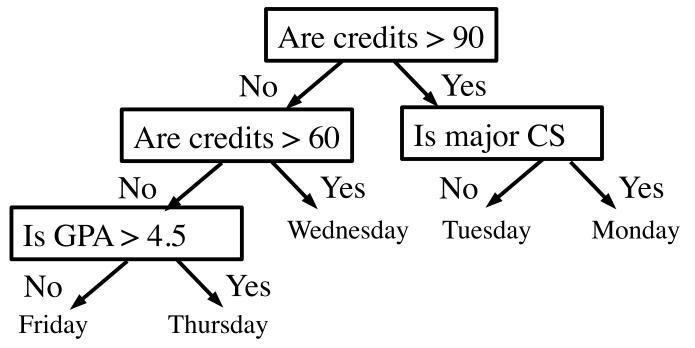


Tree as TagNodes



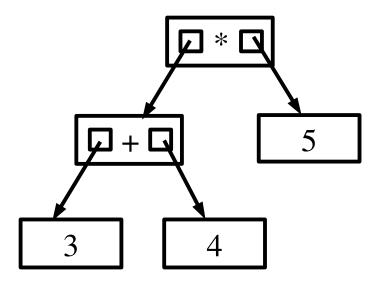
Some uses of Binary Trees

• Decision trees: e.g., when can you register



Some uses of Binary Trees

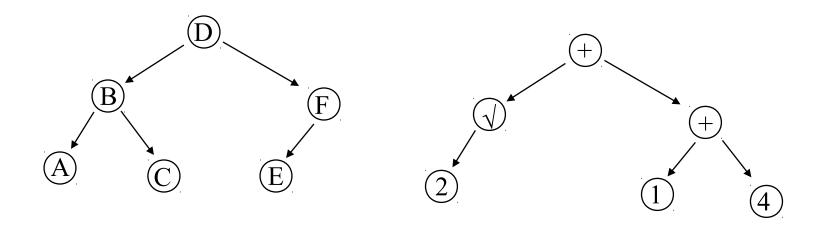
• Expression Trees: (3+4)*5



Tree Traversals

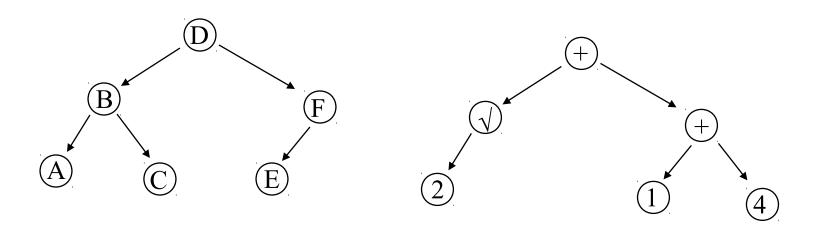
- · "Visit" each node in the tree
 - "visit" = do something, e.g. print
- In some systematic order
 - eg "inorder":
 visit all nodes in left subtree
 then visit node
 then visit all nodes in right subtree

Traversals



InOrder ____

Traversals



InOrder

A B C D E F

 $2\sqrt{14}$

See TreeNode.java

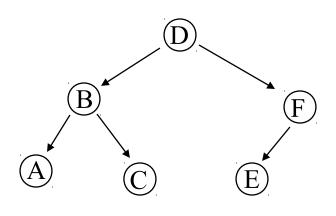
Recursive Traversals

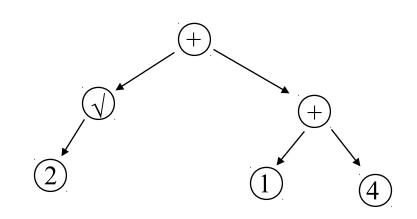
```
inOrderPrint(tree):
    if (tree == null)
        return
    inOrderPrint(tree.lst)
    print(tree.node)
    inOrderPrint(tree.rst)
```

```
preOrderPrint(tree):
    if (tree == null)
        return
    print(tree.node)
    preOrderPrint(tree.left)
    preOrderPrint(tree.right)
```

```
postOrderPrint(tree):
    if (tree == null)
        return
    postOrderPrint(tree.lst)
    postOrderPrint(tree.rst)
    print(tree.node)
```

Traversals





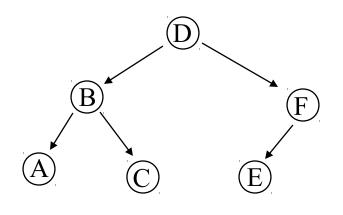
InOrder ABCDEF

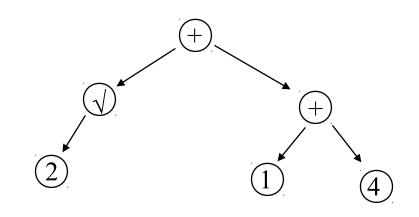
 $2 \sqrt{+1} + 4$

PreOrder

PostOrder

Traversals





InOrder

ABCDEF

 $2 \sqrt{+1} + 4$

PreOrder D B A C F E

 $+ \sqrt{2} + 1 4$

PostOrder

ACBEFD

 $2 \sqrt{14} + +$

From here to end of lecture is optional

Will not be on any exam

Non-recursive Traversals: Not the books approach

- Problem: If a node has more than one child
 - can't work on all children, grandchildren, ...
 at once
 - have to store children that have been found but not processed
- Solution: store in a stack or queue

Stack-based Traversal

```
s.push(root);
while(! s.isEmpty()){
 next = s.pop();
 if (next != null){
   print next.data;
   s.push next.rightSubTree;
   s.push next.leftSubTree;
}}
```

Queue-based Traversal

```
q.enqueue(root);
while(! q.isEmpty()){
 next = q.dequeue();
 if (next != null){
   print next.data;
   q.enqueue next.leftSubTree;
   q.enqueue next.rightSubTree;
```

Breadth vs Depth first

- Stack: depth first
 - do all children before anything else
- Queue: breadth first
 - do all at same level before anything else

Size of Stack / Queue

- Stack: path from root to leaf: O(height)
- Queue: entire level: O(2height)
 - That's a lot!
 - Solution: Iterative Deepening

Iterative Deepening

```
print all nodes at depth d:
  idfs(tree, d)
      if d == 0
         print tree.data
      else idfs(tree.lst, d-1)
           idfs(tree.rst, d-1)
• Try all depths
 for(j=0; j<maxDepth; j++)</pre>
      idfs(tree, j)
```

Iterative Deepening

- How much extra work?
 - How many leaves in complete binary tree of depth d? 2^d
 - How many non-leaves: 2^d-1
- Time overhead: roughly a factor of 2

Stack based traversal: The book's approach

- What to store?
- Mimic stack of invocation records from recursive version
 - TreeNode node;
 - int milestone;

See TreeNode.java