## Computer Science 112 Data Structures

Lecture 11:

**Trees and Binary Search Trees** 

#### Midterm Exam

- This Sunday, March 1
- 3 4:20 pm
- Know which recitation you are registered for
- Rooms:

Section 20 (Tues, 3:35 pm): Tillet 257

Section 21 (Tues, 5:15 pm): Engineering B120

Section 22 (Tues, 6:65 pm): Engineering B120

#### Midterm Exam

- Closed book
- Closed notes
- No electronics
  - EG: no phone, no calculator, no smart watch
- Do bring photo ID with legible picture

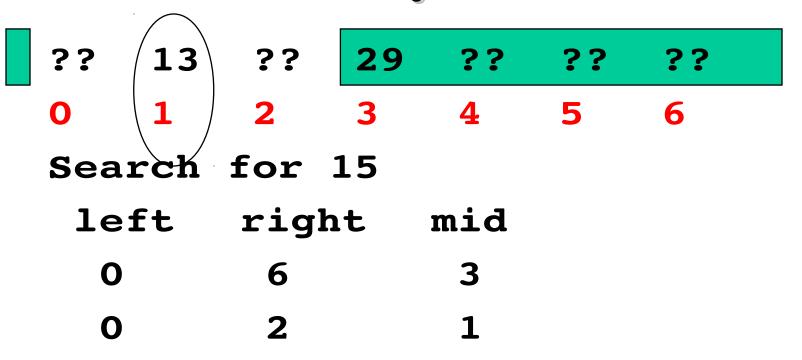
#### Midterm Exam

- Topics: Everything in lecture and assigned reading, through Binary Search
  - Will <u>not</u> cover Binary Search Trees

```
?? ?? ?? ?? ?? ??
0 1 2 3 4 5 6
Search for 15
left right mid
0 6
```

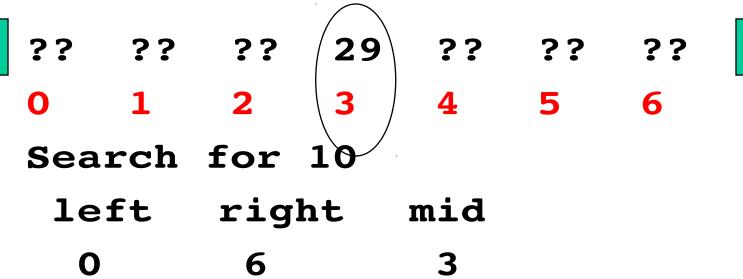
```
?? ?? ?? 29 ?? ?? ??
0 1 2 3 4 5 6
Search for 15
left right mid
0 6 3
```

```
?? ?? ?? 29 ?? ?? ??
0 1 2 3 4 5 6
Search for 15
left right mid
0 6 3
0 2
```

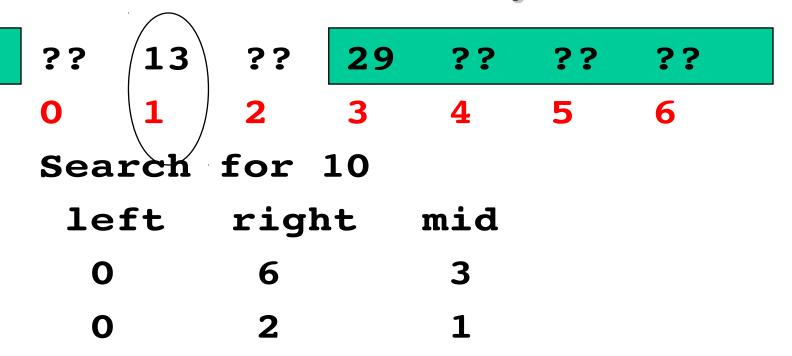


```
?? 13
                       ??
Search for 15
left right mid
        6
```

```
?? ?? ?? ?? ?? ??
0 1 2 3 4 5 6
Search for 10
left right mid
0 6
```



```
?? ?? ?? 29 ?? ?? ??
0 1 2 3 4 5 6
Search for 10
left right mid
0 6 3
0 2
```



??	13	??	29	??	??	??
0	1	2	3	4	5	6
Sea	rch	for	10			
le	left		right			
0	0		6			
0	0		2			
0	0					

12	13	??	29	??	??	??
<b>O</b>	1	2 for	3	4	5	6
Sea	rch	for	10			
left		right		mid		
0		6		3		
		O		3		
0		2		1		

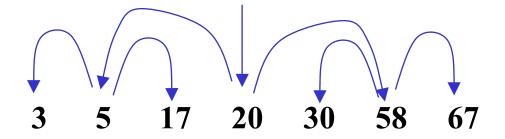
12	13	??	29	??	??	??	
0	1	2	3	4	5	6	
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0		0		0			
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			See	Binary	Searc	ch.java	l
CS112: Slide	es for Prof. Stei	inberg's lecture		Recurs e-trees.odp	siveB	inSearc 17	h.j

## **How Many Comparisons?**

- How many elements of the array do we have to compare with the target?
  - For each element we look at, size of remaining region is cut in half
  - When remaining region is 1 element, we are done
  - O(log(n))

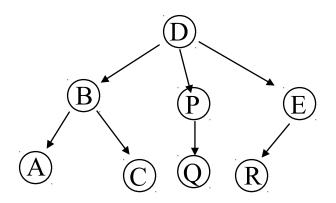
## **Binary Search Trees**

- Why can't we do binary search on a linked list?
  - can't jump to middle
- Lets keep a pointer to the middle!
- Then what?



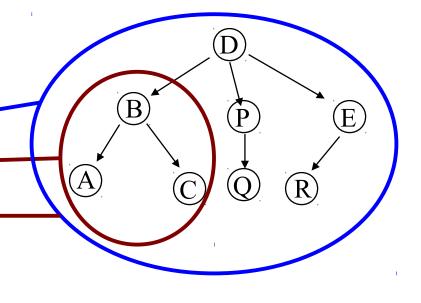
#### **New: Tree Terminology**

- Nodes (vertices) and arcs (edges)
- Relationships:
  - Parent and Child
    - E is a child of D
    - D is the parent of E

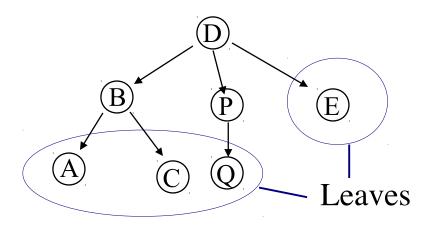


- Nodes and arcs (edges)
- Relationships:
  - Parent and Child
    - E is a child of D
    - D is the parent of E
  - Root and Subtree
    - D is the root of this tree

    - The root of the subtreeis B

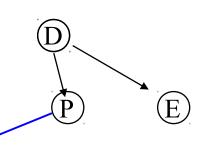


- Root has no parents
- All nodes except the root have a single parent
- Leaf node has no children
- There is exactly one path from root to any node



Height of tree
Depth of a node
depth = 1
depth = 2
depth = 3

- What is the height of this tree?
  - What is the depth of this node? –



- What is the height of this tree?
  - What is the depth of this node?



• What is the height of the empty tree?

## Binary tree

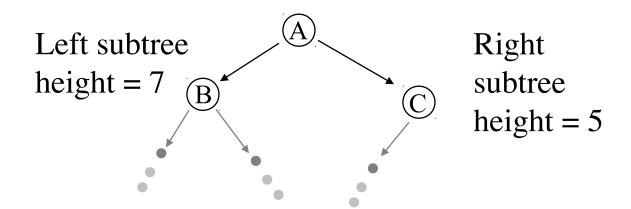
- each node has at most 2 subtrees
  - left and right subtreeroot of left (right) subtree is calledthe left (right) child

#### **Recursive Data Structures**

- Recursive definition of a binary tree
  - empty (i.e. null)
  - not empty
    - data at the root
    - a left subtree, which is a binary tree
    - a right subtree, which is a binary tree

## height

What is the height of the whole tree?



# Recursive functions height

#### **Recursive functions**

• Common form of function on a tree is recursive

f(tree):	
if (tre	ee = = null) return
else	return (data, f(tree.lst), f(tree.rst))
	is a value and a function

# Recursive functions height

```
height(tree):

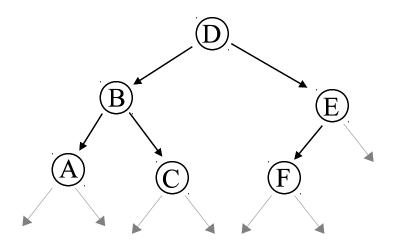
if (tree = = null) return (-1)

else return (1 + max) ( height (tree.lst),

height (tree.rst))
```

## Recursive functions nodeCount

# Recursive functions nodeCount



#### Recursive functions Sum

## Recursive functions has0

```
has0(tree):

if (tree = = null) return

else return (tree.data,

has0( tree.lst ),

has0( tree.rst ))
```

# Recursive functions has0

```
has0(tree):

if (tree = = null) return false

else return or (tree.data = = 0,

has0(tree.lst),

has0(tree.rst))
```

## Binary Search Tree

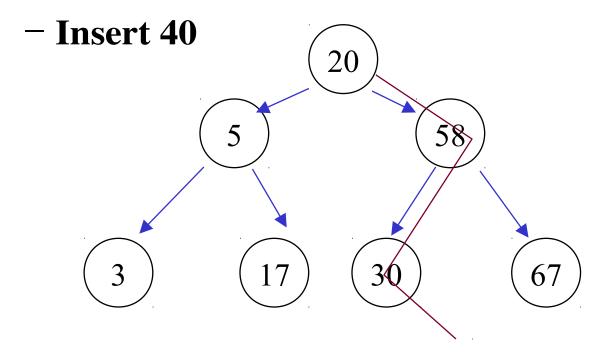
- data at a node is > any data in left subtree
- data at a node is < any data in right subtree</li>
- Therefore, to print a BST in data order:
  - Print left subtree in data order
  - Print data
  - Print right subtree in data order

#### Search

- Searching a BST is easy
  - if node = null, search fails
  - if node.data equals target, found
  - if target < node.data, search on left subtree
  - else search on right subtree

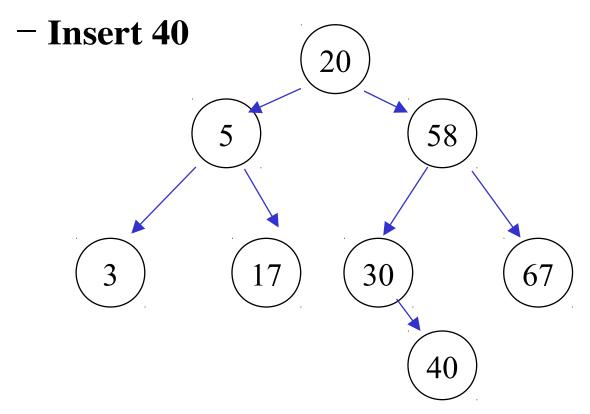
#### **Insert**

Search, fail, insert where failed



#### **Insert**

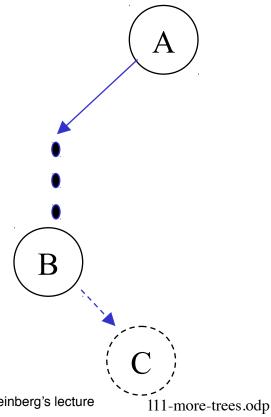
Search, fail, insert where failed



#### **Delete**

- Three cases
  - node to delete has no children => delete it
  - node to delete has 1 child => replace node with child
  - node to delete has 2 children

 Observation: for node with left child, inorder predecessor has no right child



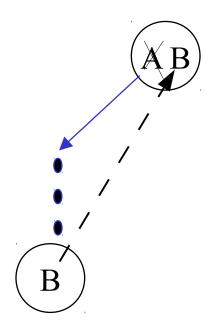
If C exists, C > B and A > C

So B cannot be inorder predecessor of A

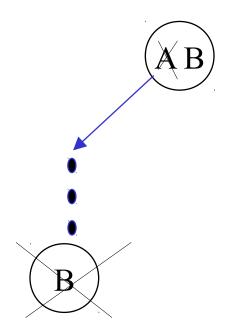
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- Replace data at node with data of inorder predecessor
- Delete inorder predecessor (which must have either 0 or 1 child)

Replace data at node with data of inorder predecessor



• Delete inorder predecessor (which must have either 0 or 1 child)



See BSTNode.java, BST.java, and BSTApp.java

## Repeated Keys

• What do you do if you can have two nodes with the same data?

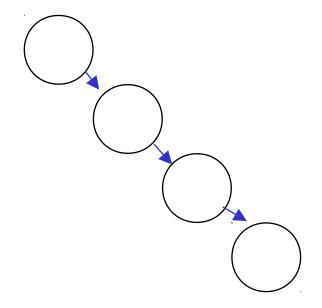
## **Cost of using BST**

Search, insert delete: O(depth)

- What is depth of tree?
  - with n nodes, best depth is log n
  - but worst depth is n

## **Binary Search Trees**

• Problem: insertion & deletion can give tree of any shape - even



#### **Binary Search Trees**

- Problem: insertion & deletion can give tree of any shape
- Solution: AVL trees