

# **Computer Science 112**

## **Data Structures**

### **Lecture 22:**

### **Graphs:**

### **Breadth First Search**

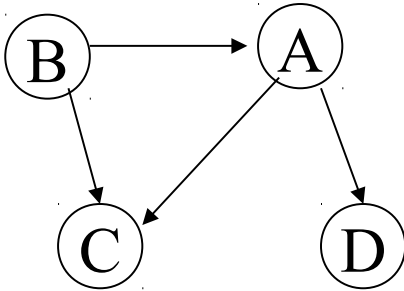
### **Shortest Path**

# Review: Graph Traversals

## Depth First Traversal

- **for each vertex  $v$  in graph:**  
    **call  $\text{dfsG}(v)$**
- **$\text{dfsG}(v)$ :**  
    **if ( $\text{marked}(v)$ ) return;**  
    **visit  $v$ ;**  
    **mark  $v$ ;**  
    **for each  $vn$  in  $\text{neighbors}(v)$**   
         **$\text{dfsG}(vn)$**

# Graph Traversals



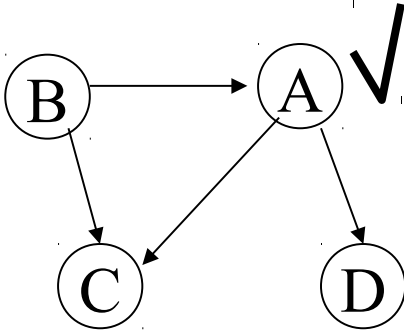
**Driver**

$\mathbf{v} = \langle \mathbf{A} \rangle$

**dfsG**

$\mathbf{v} = \langle \mathbf{A} \rangle$

# Graph Traversals



**Driver**

$\mathbf{v} = \langle \mathbf{A} \rangle$

**dfsG**

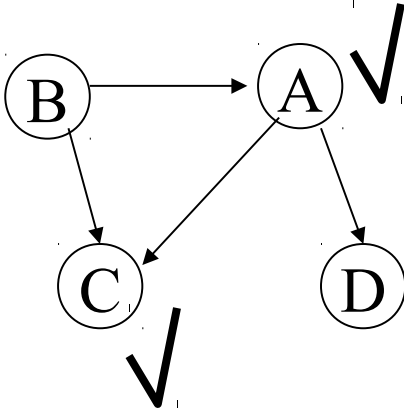
$\mathbf{v} = \langle \mathbf{A} \rangle$

$\mathbf{vn} = \langle \mathbf{C} \rangle$

**dfsG**

$\mathbf{v} = \langle \mathbf{C} \rangle$

# Graph Traversals



**Driver**

$v = \langle A \rangle$

**dfsG**

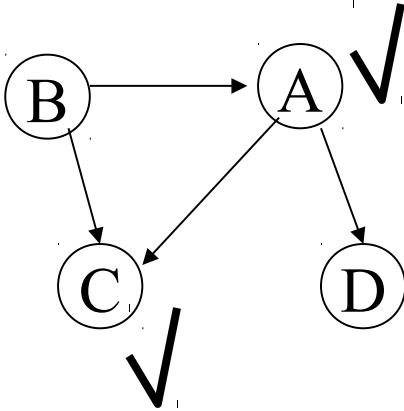
$v = \langle A \rangle$

$vn = \langle C \rangle$

**dfsG**

$v = \langle C \rangle$

# Graph Traversals



**Driver**

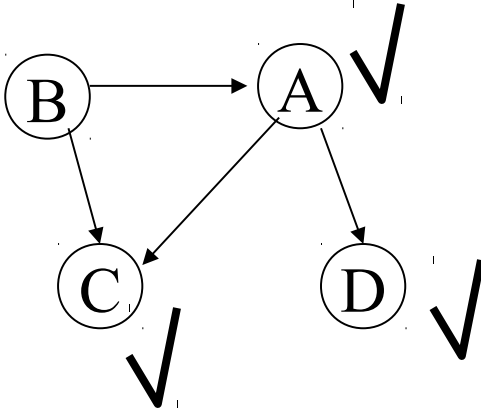
$\mathbf{v} = \langle \mathbf{A} \rangle$

**dfsG**

$\mathbf{v} = \langle \mathbf{A} \rangle$

$\mathbf{vn} = \langle \mathbf{D} \rangle$

# Graph Traversals



**Driver**

**$v = \langle A \rangle$**

**dfsG**

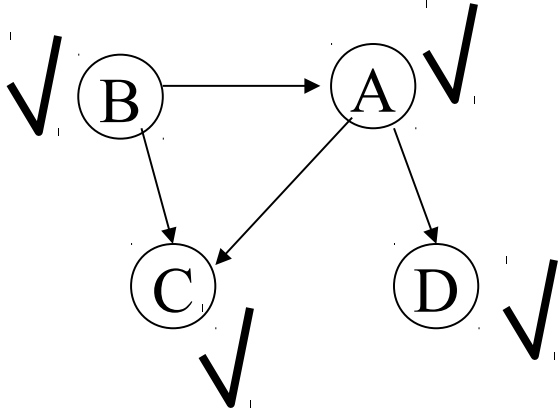
**$v = \langle A \rangle$**

**$vn = \langle D \rangle$**

**dfsG**

**$v = \langle D \rangle$**

# Graph Traversals



**Driver**

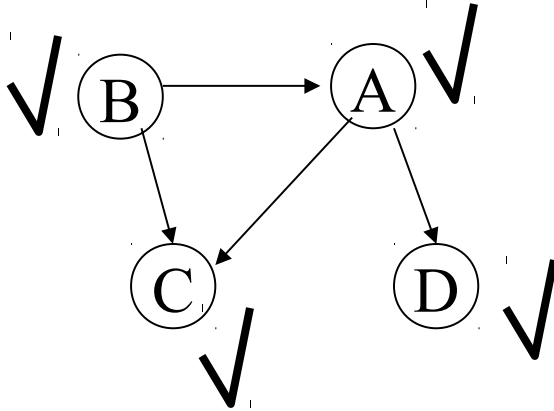
$\mathbf{v} = \langle \mathbf{B} \rangle$

**dfsG**

$\mathbf{v} = \langle \mathbf{B} \rangle$



# Graph Traversals



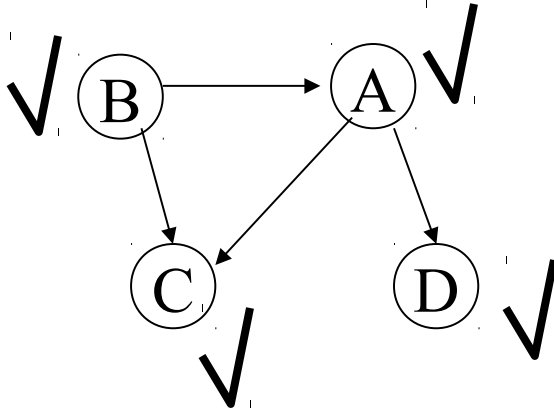
**Driver**

$v = \langle C \rangle$

**dfsG**

$v = \langle C \rangle$

# Graph Traversals



**Driver**

**$v = \langle D \rangle$**

**dfsG**

**$v = \langle D \rangle$**

# Graph Traversals

- **Time:**
  - Visit each vertex
  - inspect each edge **$O(n + e)$   $n$  vertices,  $e$  edges**

# Uses of DFS Traversal

- **Connected Components**
  - See GraphCC.java
- **Topsort**
  - See GraphTS.java

# Review: Topological Sort

- **Acyclic Digraph  $\Leftrightarrow$  partial order**
- **Topsort: find total order consistent with partial order**

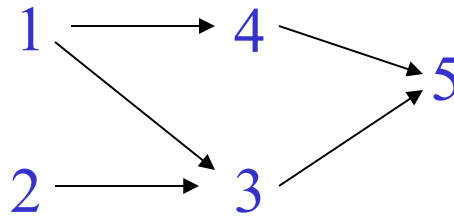
1    **a=1;**

2    **b=2;**

3    **c=a\*b;**

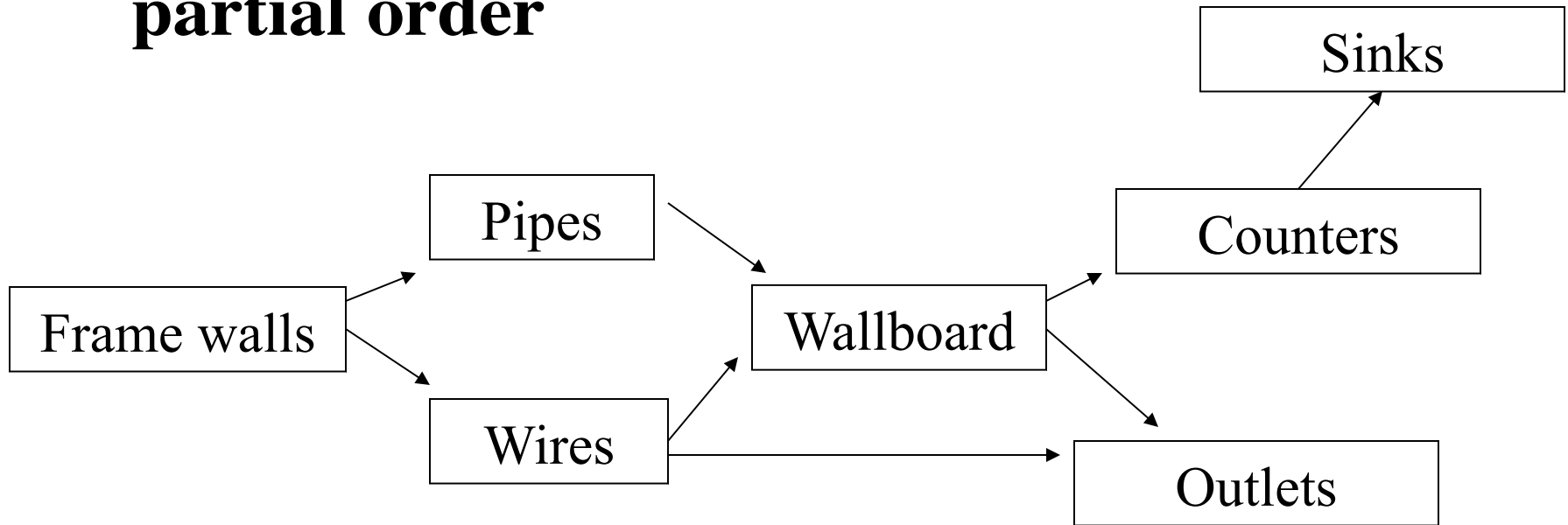
4    **d=a+4;**

5    **c=c+d**



# Topological Sort

- **Acyclic Digraph  $\Leftrightarrow$  partial order**
- **Topsort: find total order consistent with partial order**



# Topsort Algorithms

- **Most work by assigning numbers to vertices**
  - **topsorted order = numerical order**
- **Depth first**

# DFS Topsort Algorithm

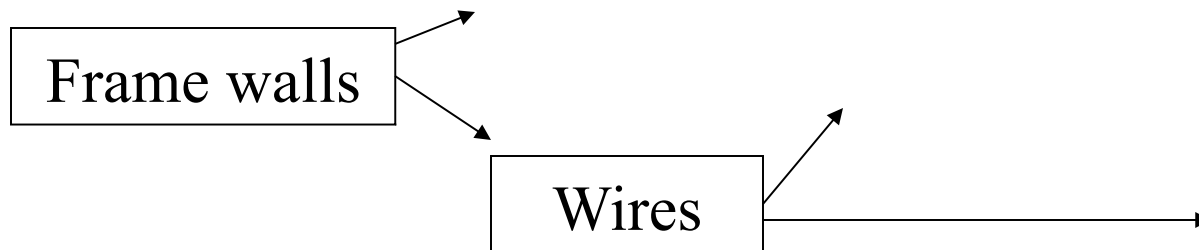
- **Algorithm:**
  - Do DFS
  - Number vertices as you leave them
- **Problem:** leave vertex *after* leave reachable vertices, but needs number *smaller* than reachable vertices
  - **Solution:** number from largest to smallest numbers
- **See GraphTS.java**



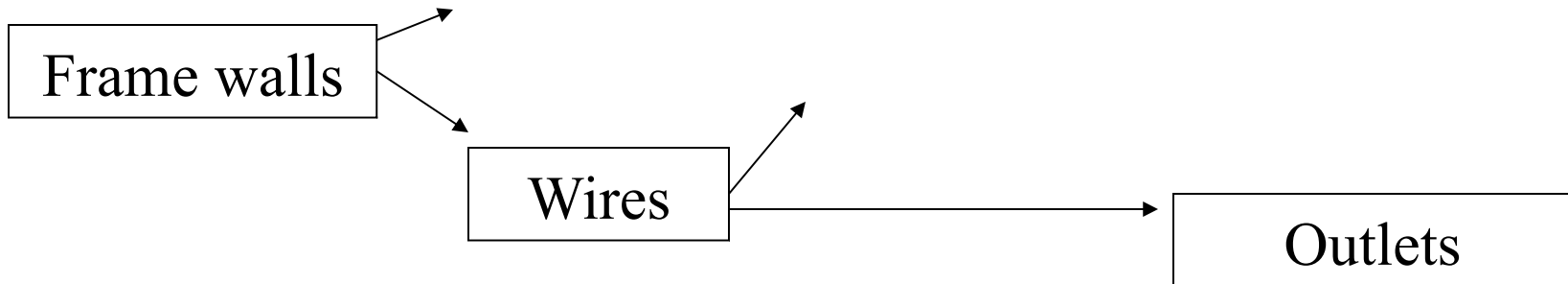
# New: Topsort Example



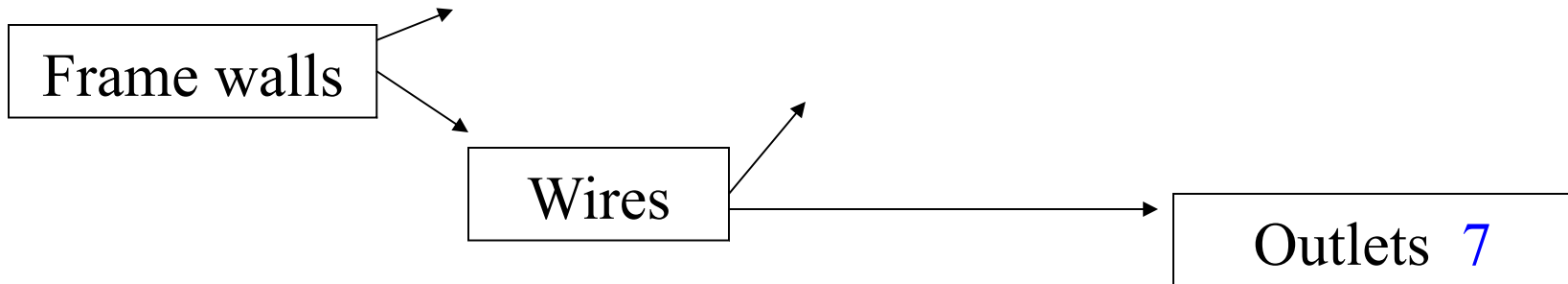
# New: Topsort Example



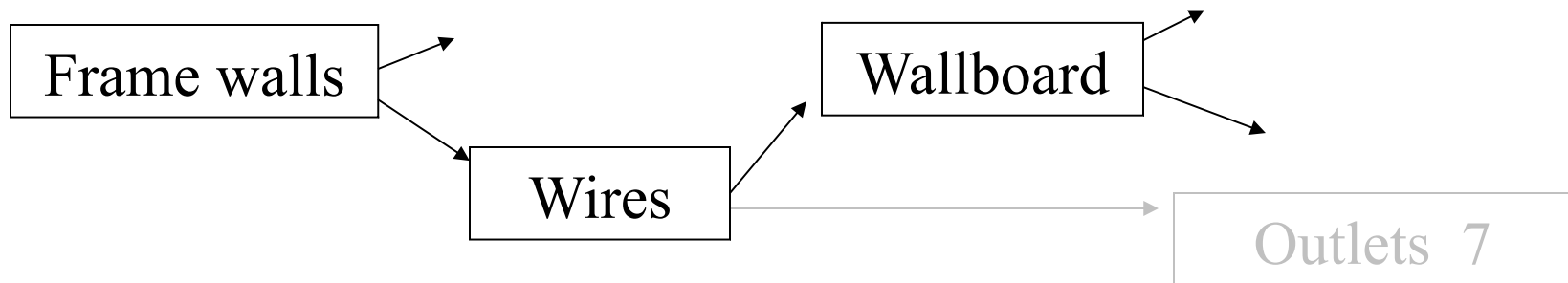
# Topsort Example



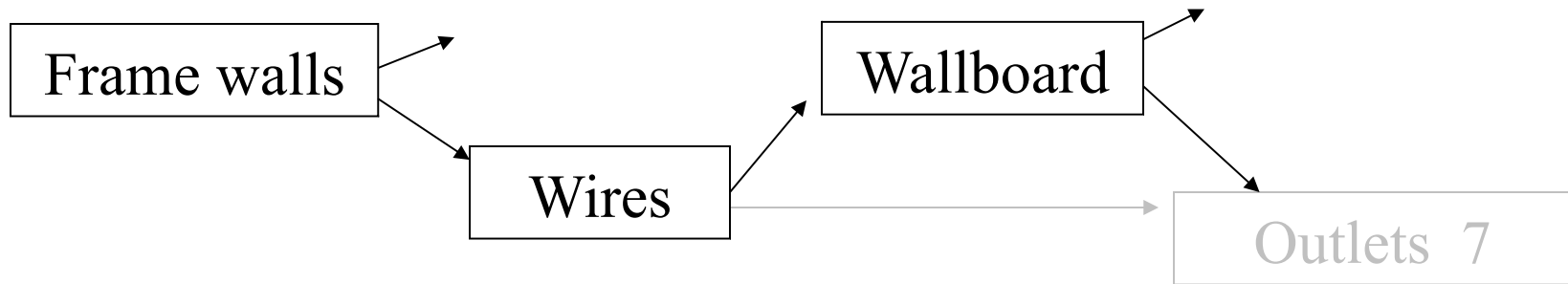
# Topsort Example



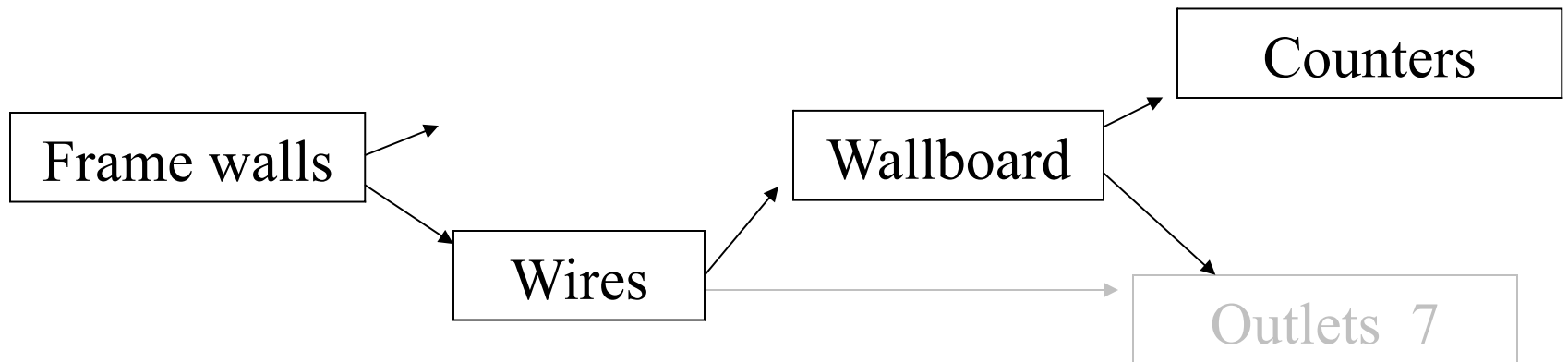
# Topsort Example



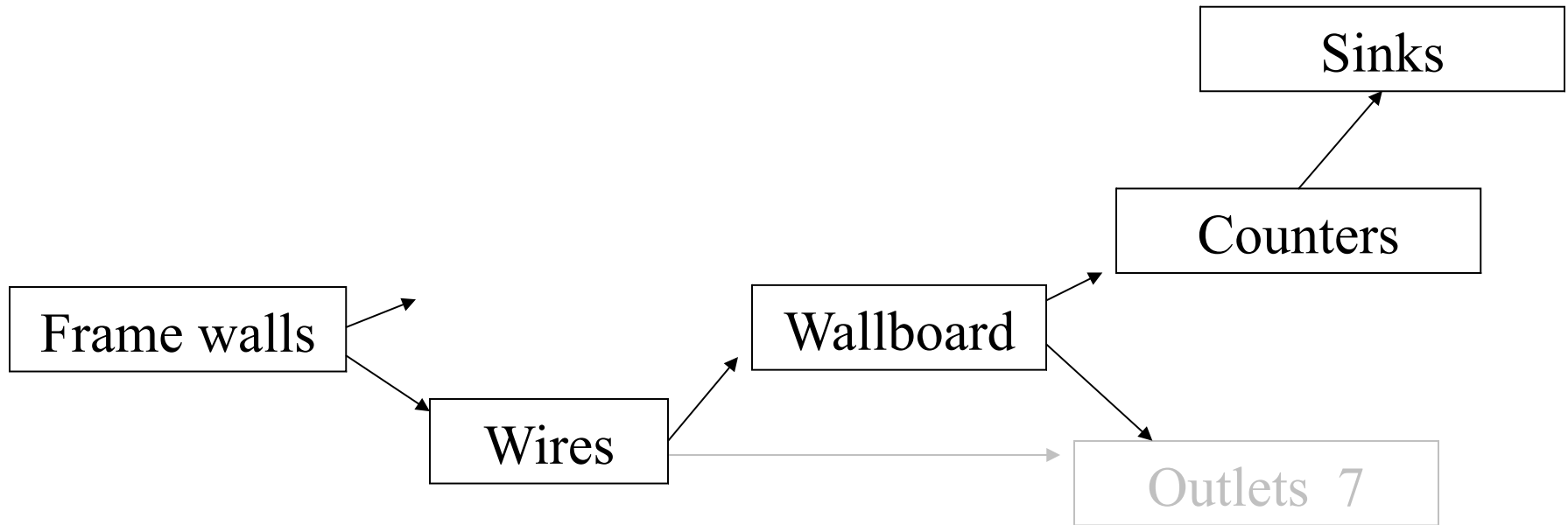
# Topsort Example



# Topsort Example

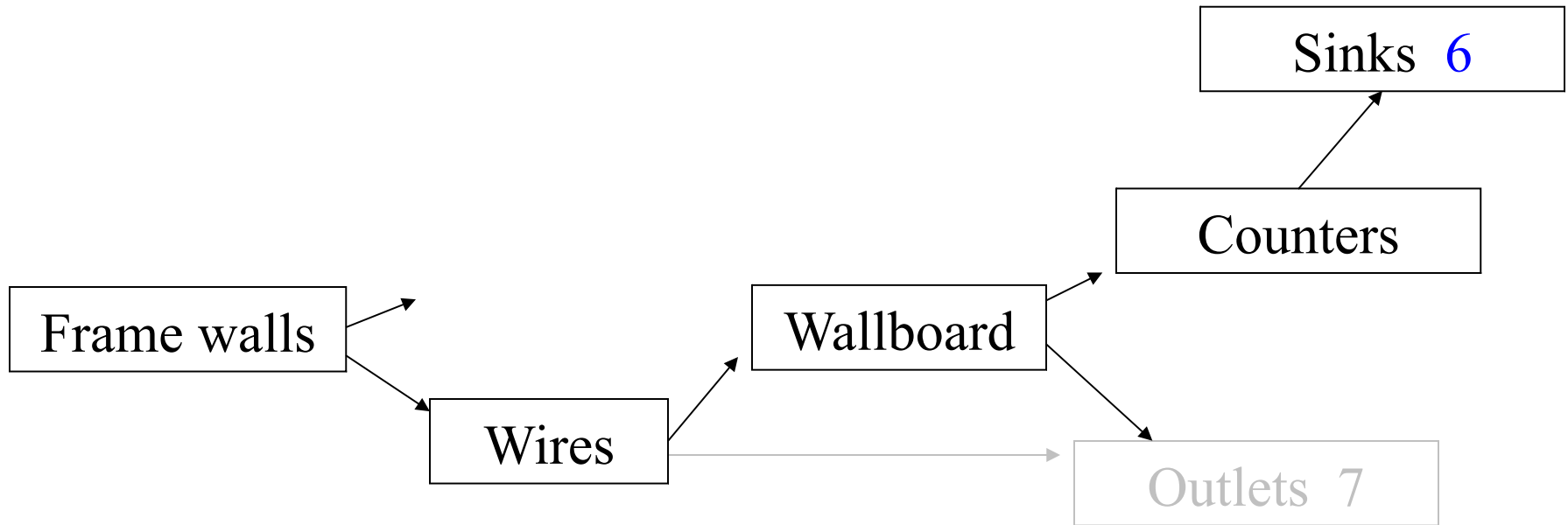


# Topsort Example

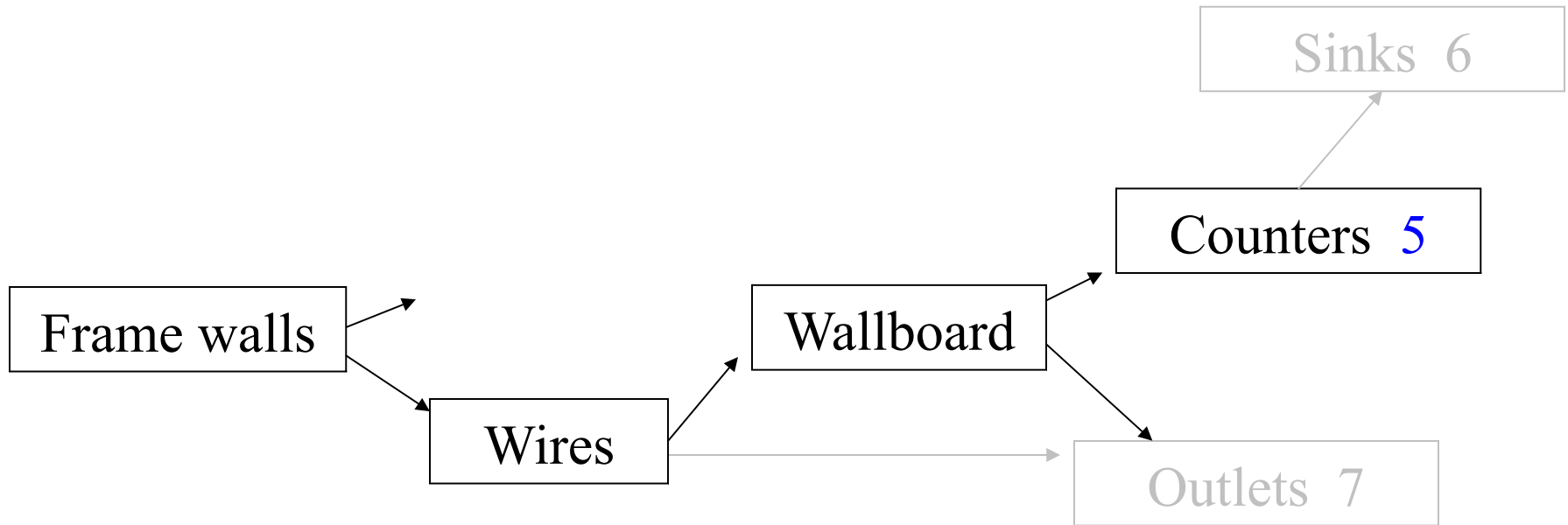




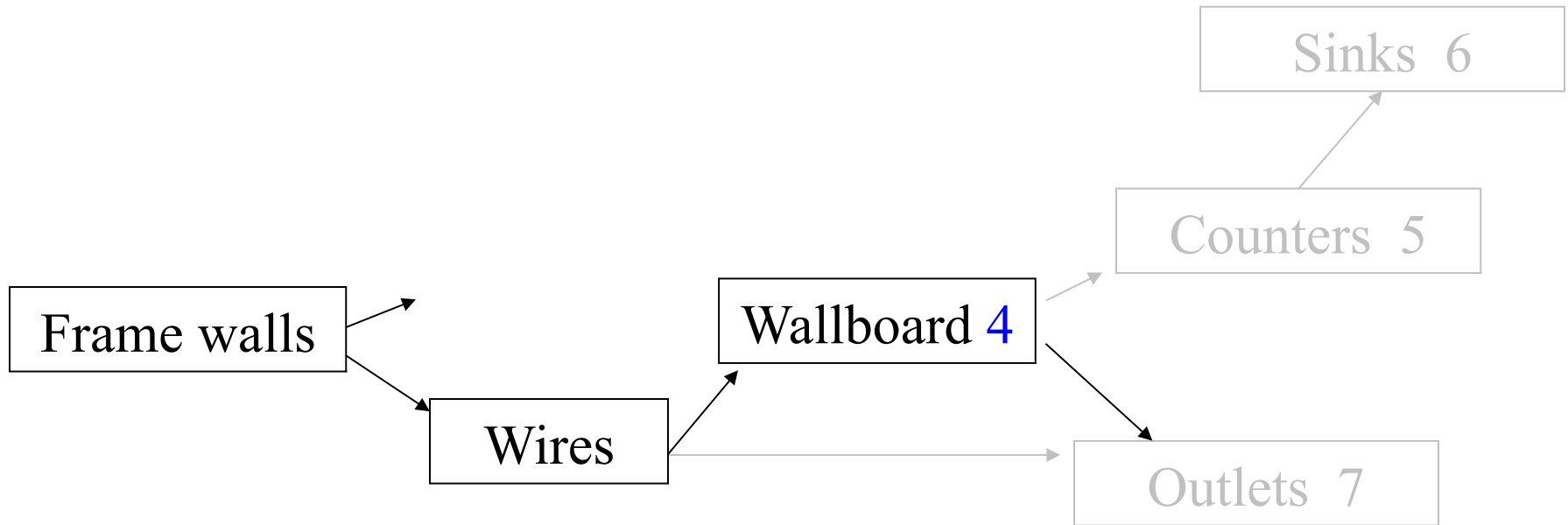
# Topsort Example



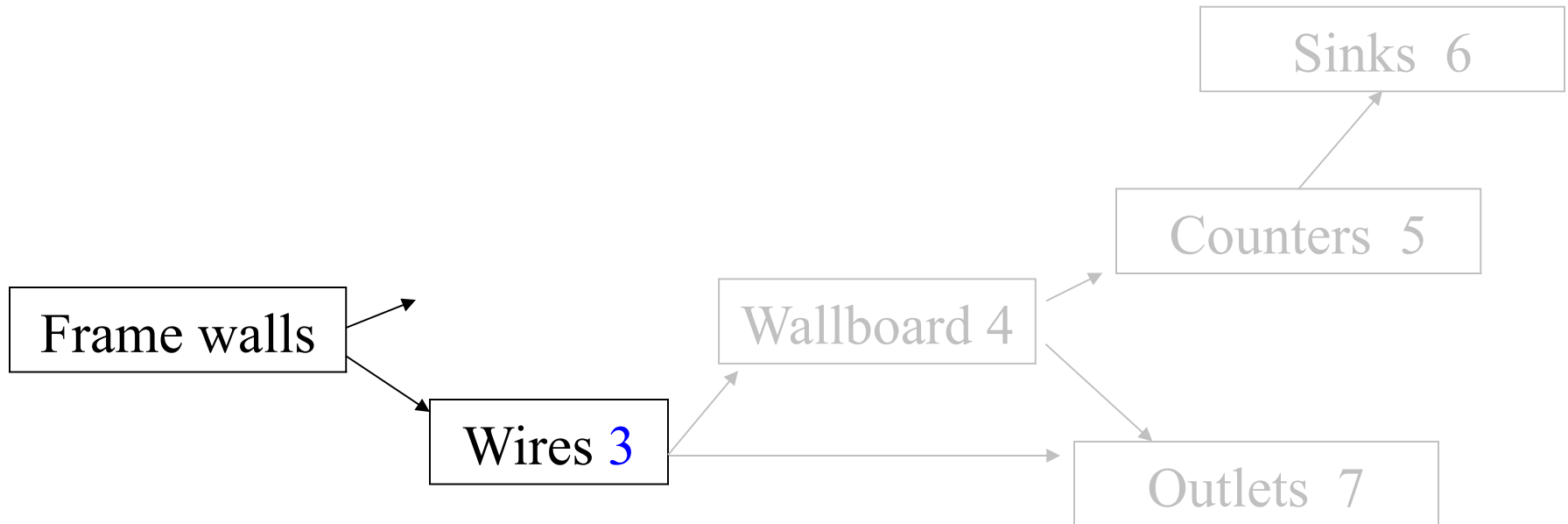
# Topsort Example



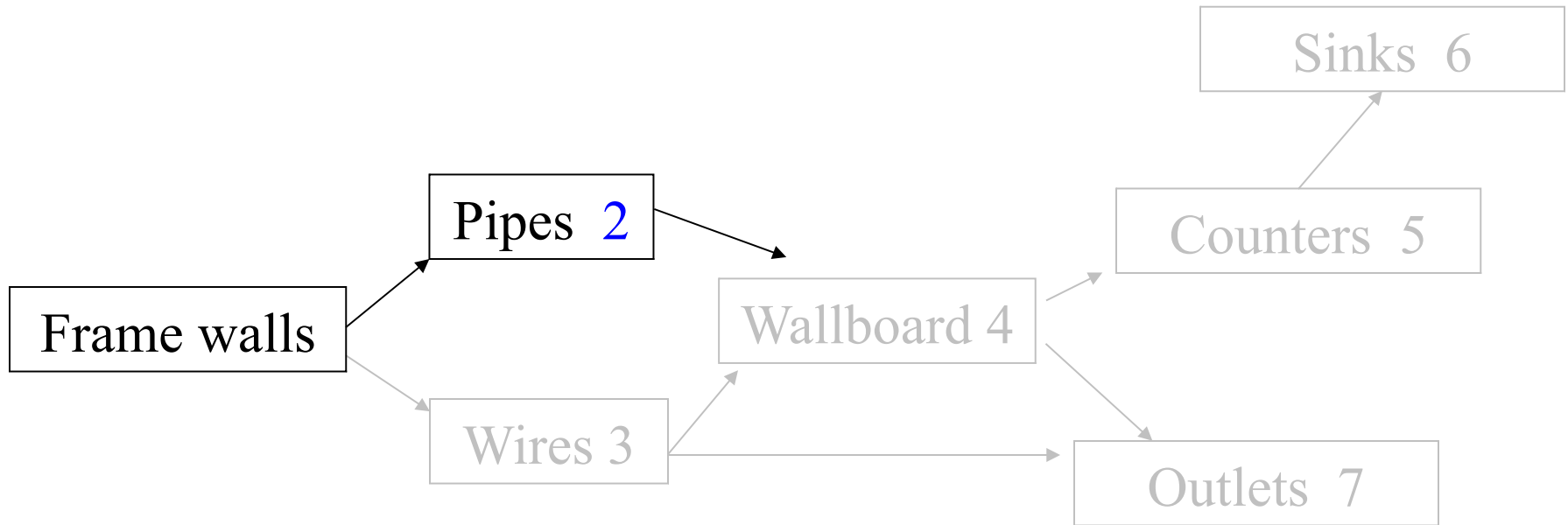
# Topsort Example



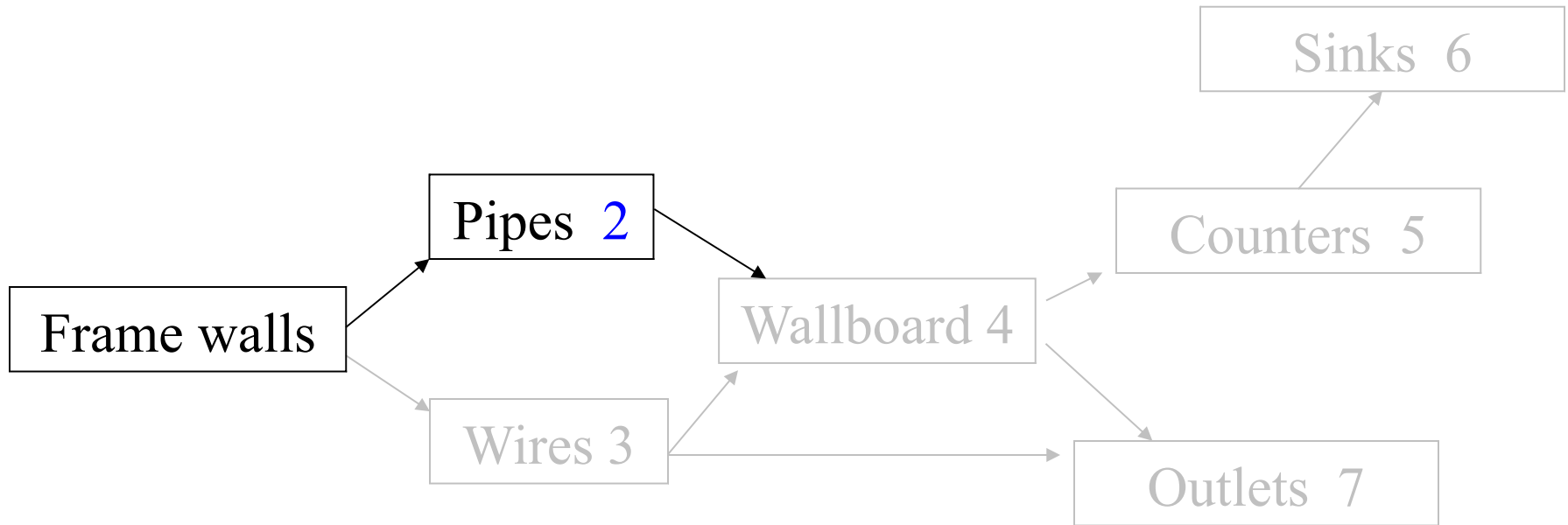
# Topsort Example



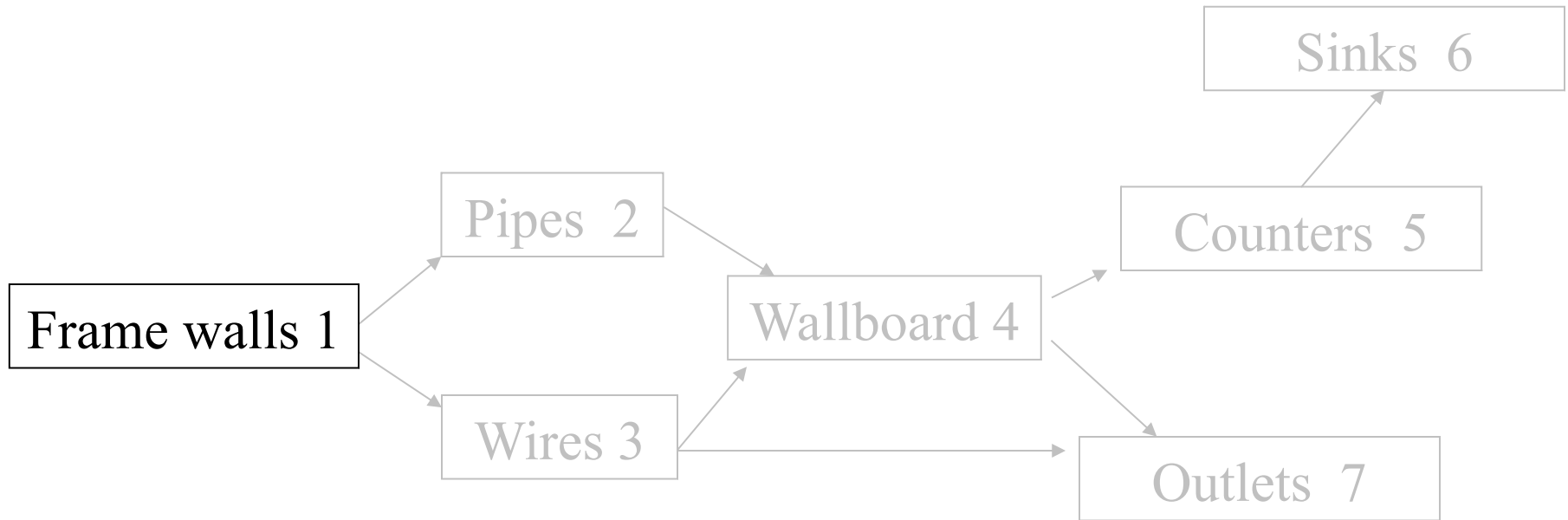
# Topsort Example



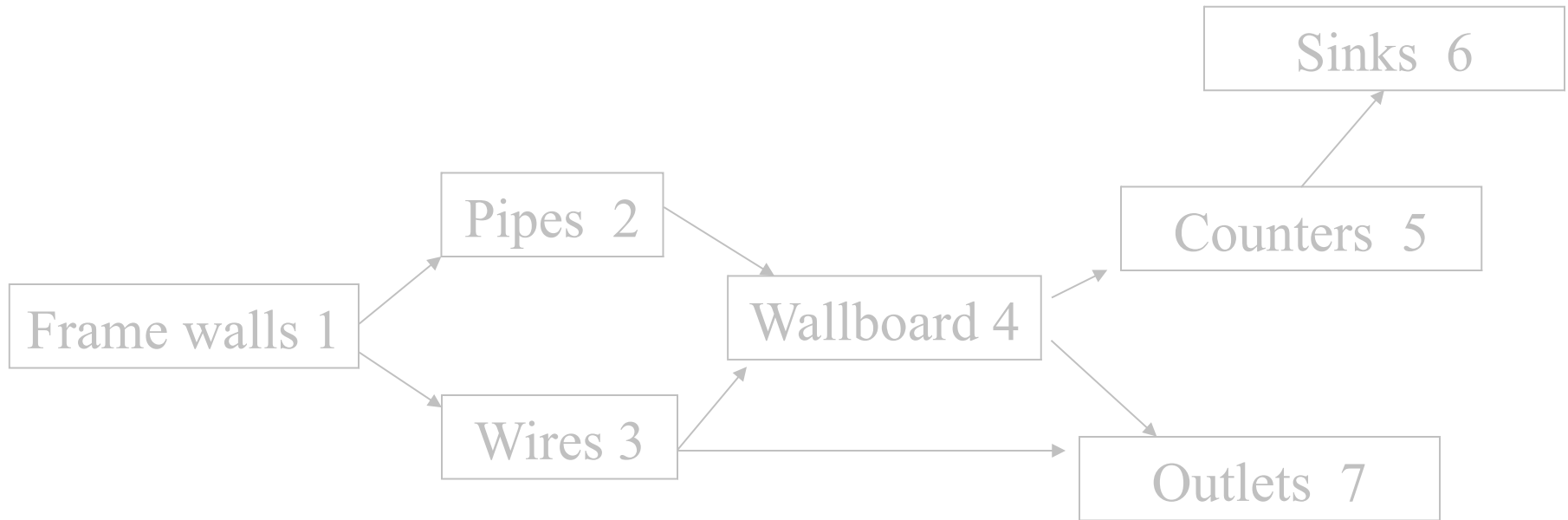
# Topsort Example



# Topsort Example

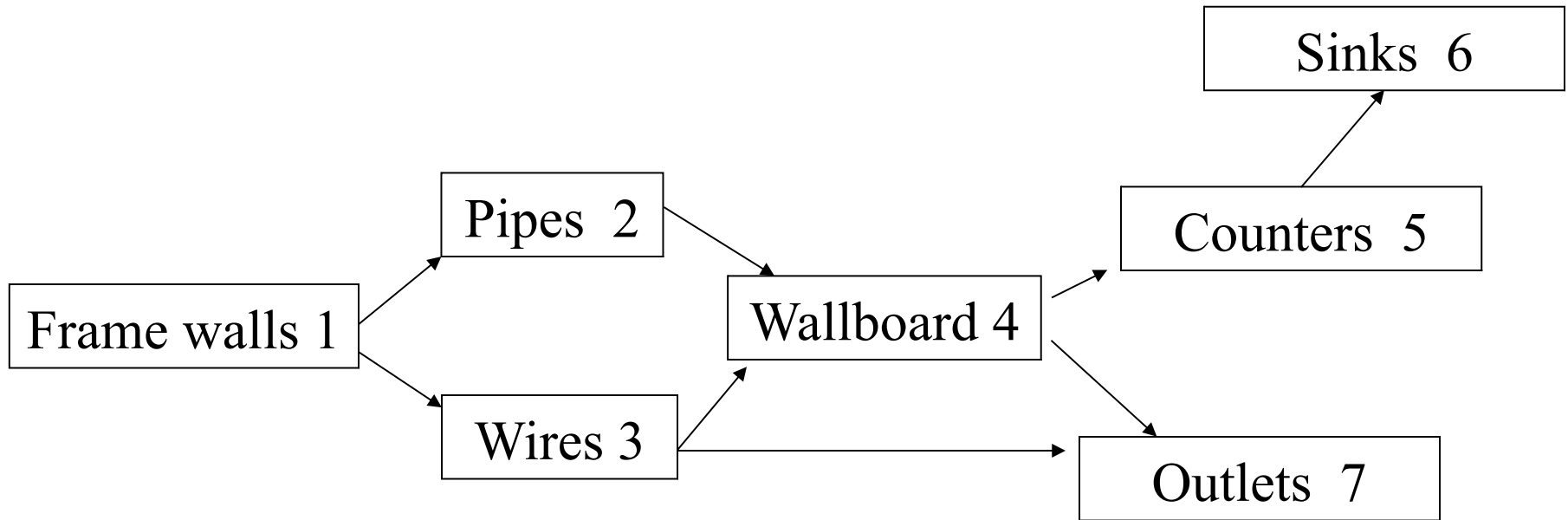


# Topsort Example





# Topsort Example

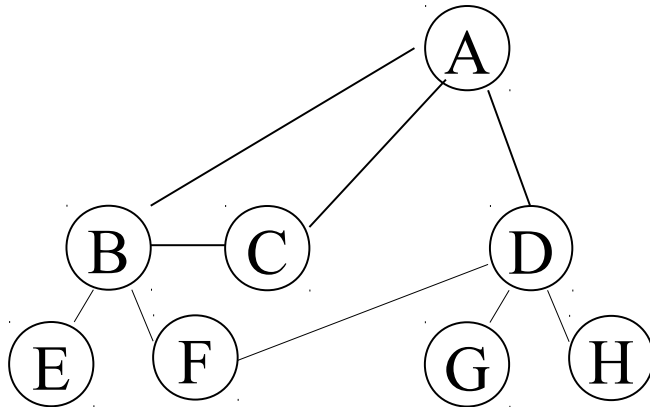


# Topsort Cost

- **DFS + numbering**  
 **$= O(n+e) + O(n) = O(n+e)$**

# New: Breadth First Search

- Like breadth first search on tree



- **Breadth First:** A B C D E F G H
- **Depth First:** A B E F D G H C

# Breadth First Algorithm

- **dfsG(v):**  
visit and mark v  
enqueue v  
while not queue.empty( )  
    dequeue into w  
    for each neighbor n of w:  
        if n not visited:  
            visit and mark n  
            enqueue n

# Breadth First Cost

- **like depth-first:**
  - **visit every vertex**
  - **cross every edge**

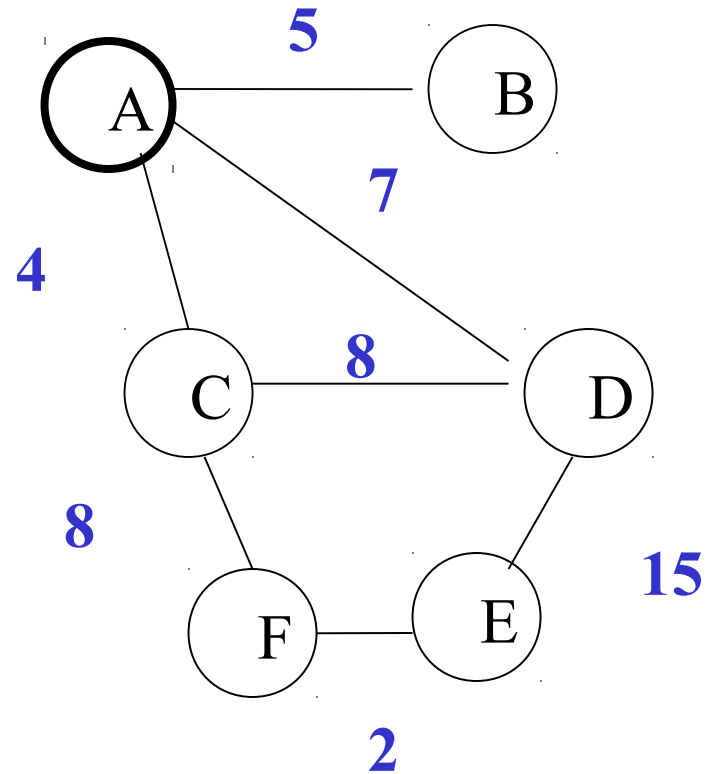
**$O(n+e)$**

# Shortest Path

- **weighted digraph**
  - **weights are all  $> 0$**
- **“length” of a path = sum of weights of arcs on path**
- **given start vertex, end vertex, find shortest path from start to end**

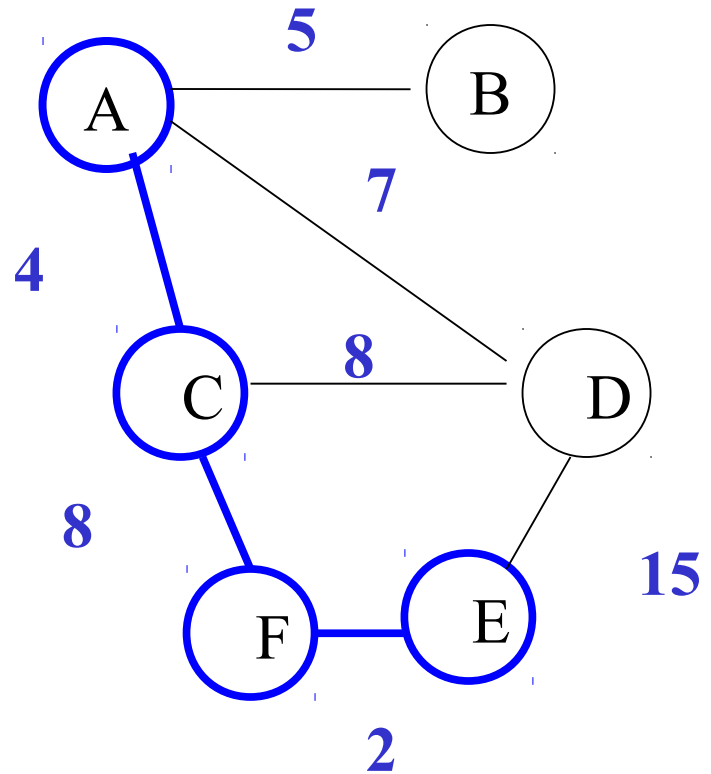
# Shortest Paths

- **What is the shortest path**
  - **from A to E?**
  - **from A to F?**



# Shortest Paths

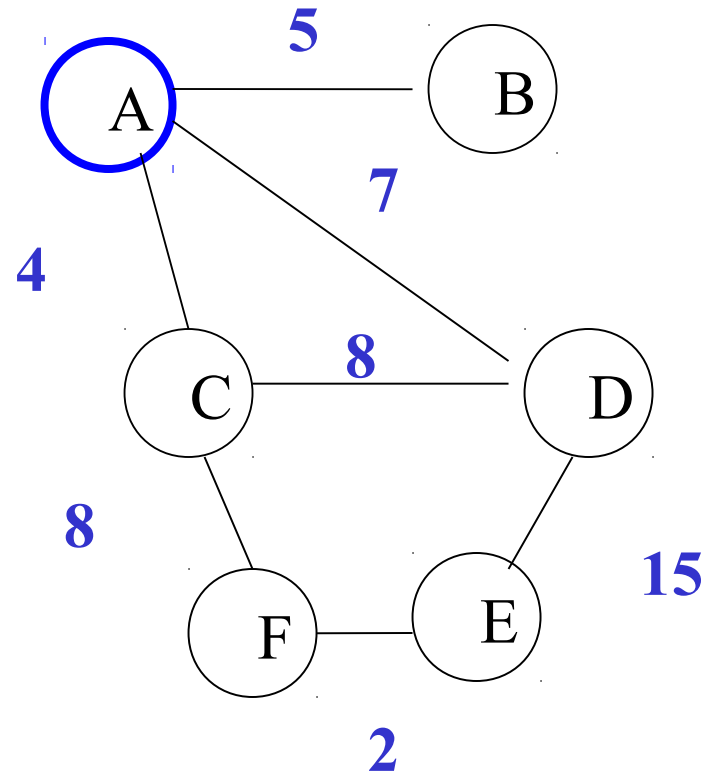
- If a shortest path from A to E runs through F, the part from A to F is a shortest path from A to F





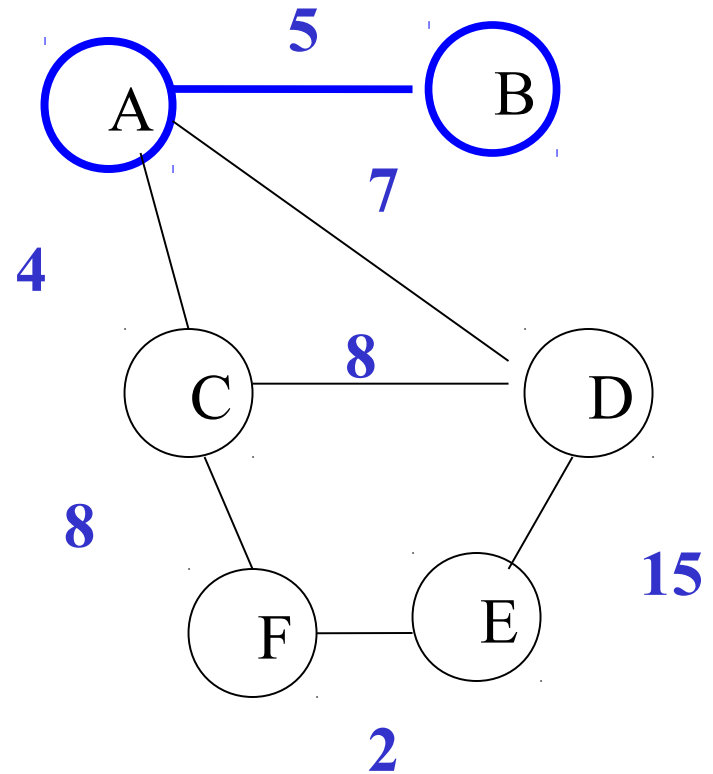
# Dijkstra's Algorithm

- Consider the shortest paths from A to each other vertex.



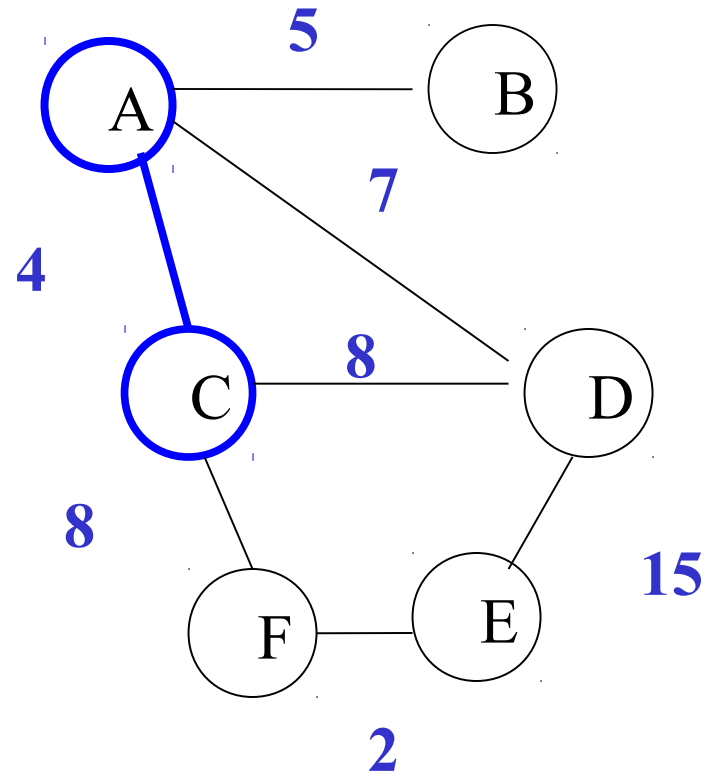
# Dijkstra's Algorithm

- Consider the shortest paths from A to each other vertex.



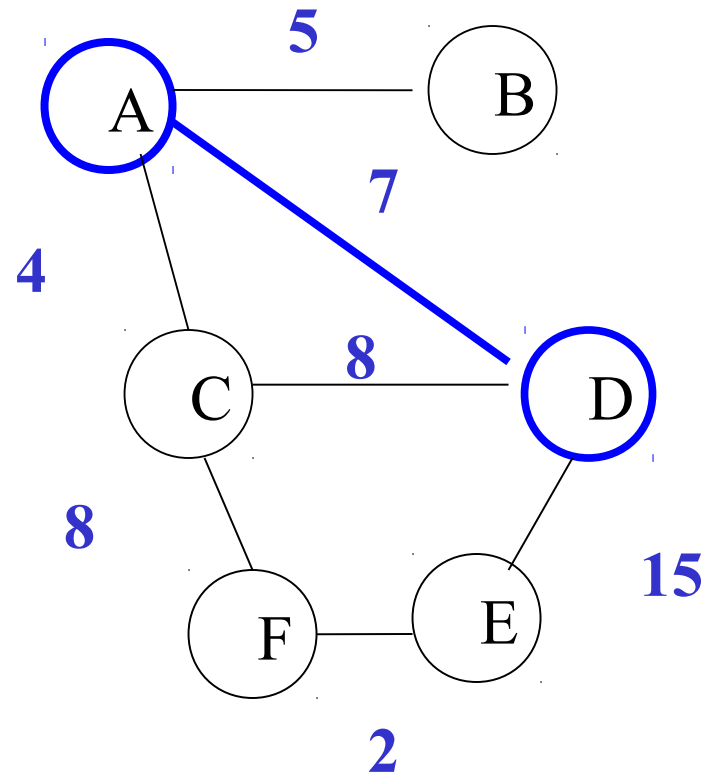
# Dijkstra's Algorithm

- Consider the shortest paths from A to each other vertex.



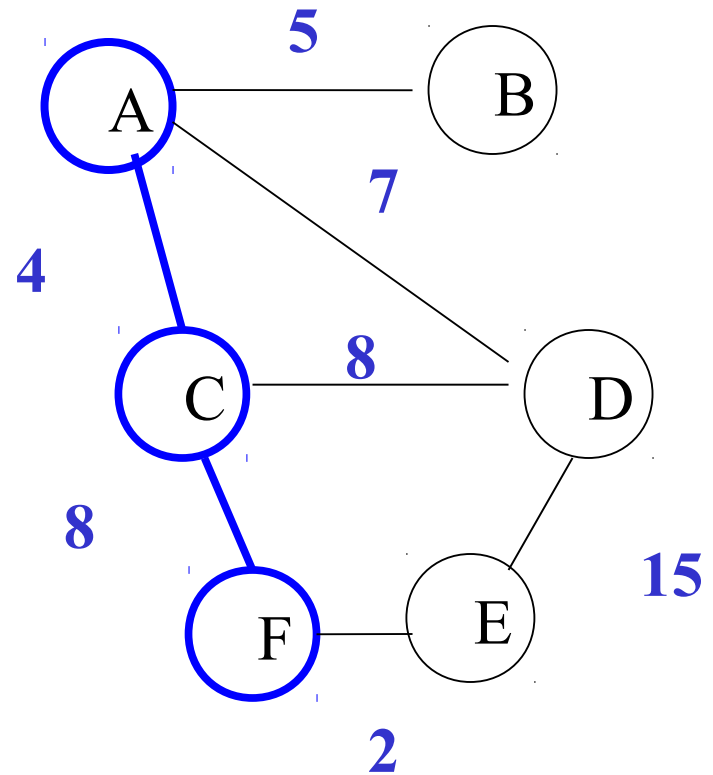
# Dijkstra's Algorithm

- Consider the shortest paths from A to each other vertex.



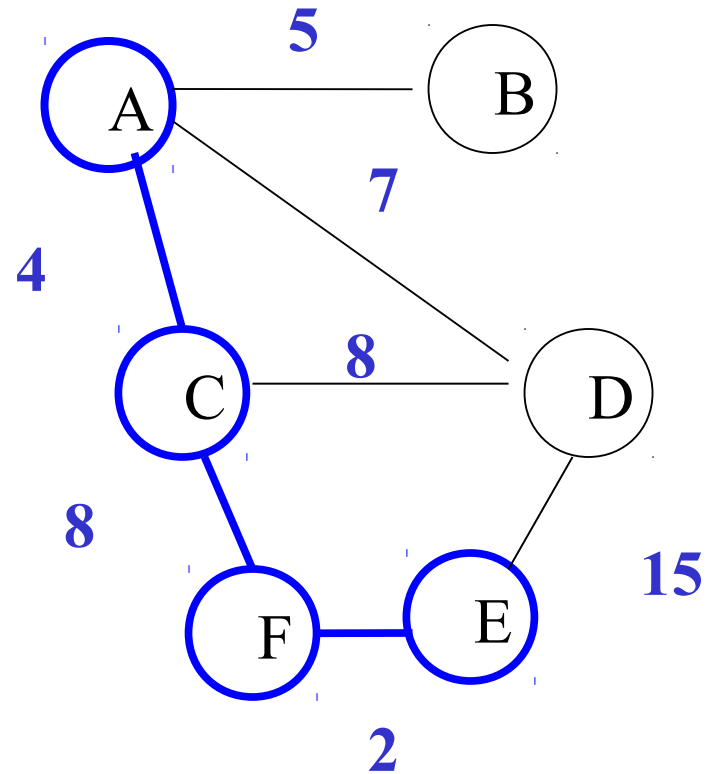
# Dijkstra's Algorithm

- Consider the shortest paths from A to each other vertex.



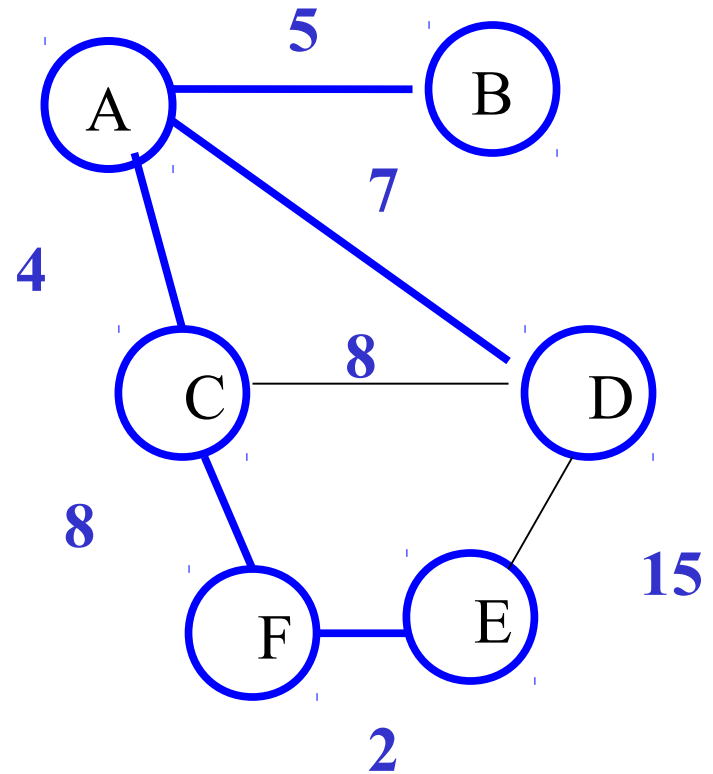
# Dijkstra's Algorithm

- Consider the shortest paths from A to each other vertex.



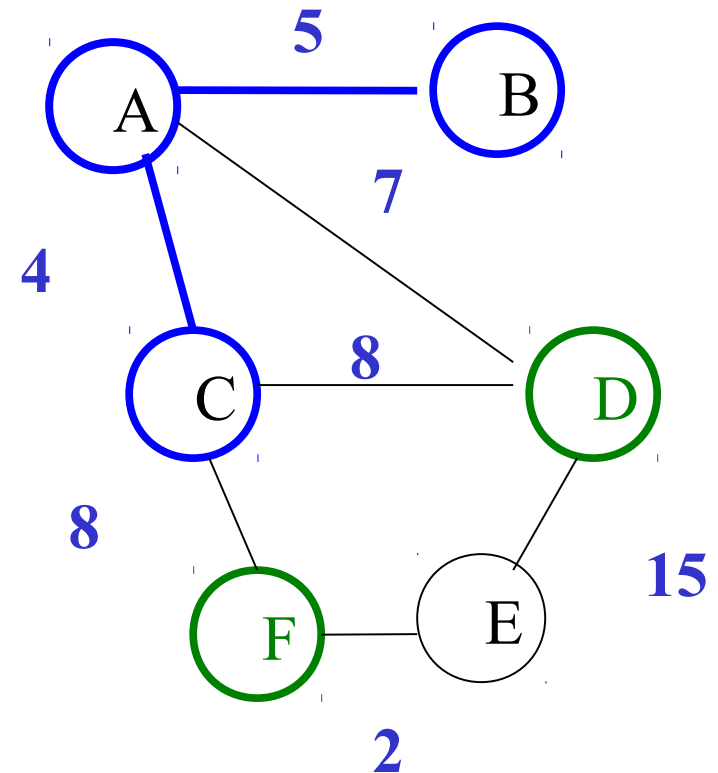
# Dijkstra's Algorithm

- These can be put together to form a tree
- *A Shortest Path Tree*



# Dijkstra's Algorithm

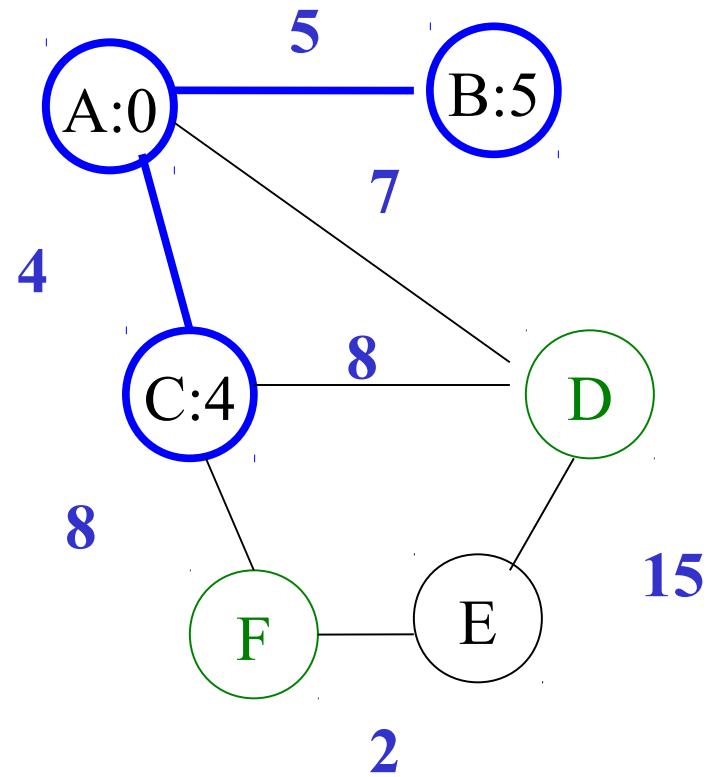
- **Grow a tree of shortest paths from start**
  - grow it one vertex at a time, closest to farthest
- **Fringe:** nodes that are not in the tree yet but have a neighbor in the tree





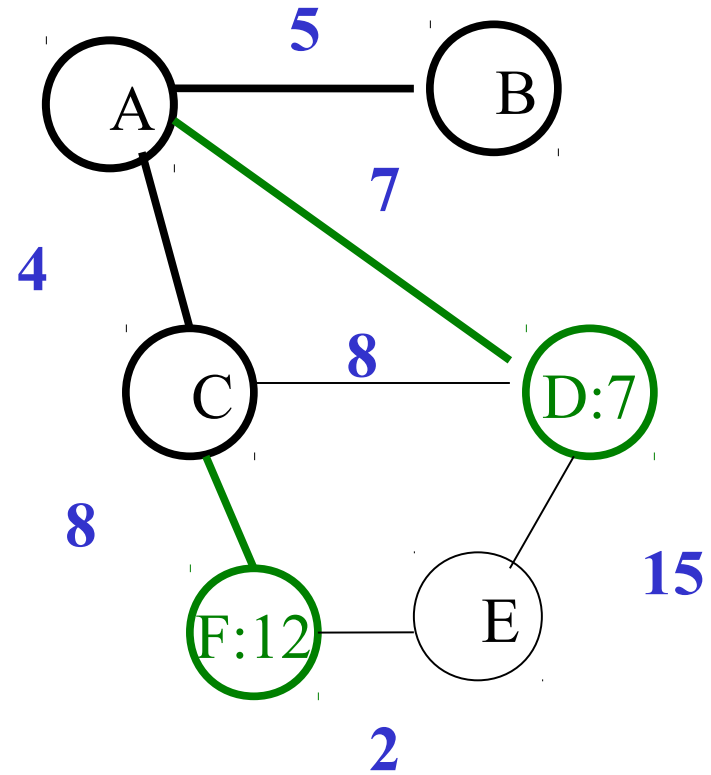
# Dijkstra's Algorithm

- Vertices in the tree have
  - a link: first step on the shortest path back to start
  - a distance: the length of that whole path



# Dijkstra's Algorithm

- Vertices in the fringe have
  - a **link**: an arc to the tree if  $> 1$  of these, use the arc that gives the shortest path back to start
  - a distance: the length of the path using link



## **Algorithm:**

**Put start vertex in the tree**

**While there are any vertices in fringe**

**Let  $v$  be vertex in fringe with  
smallest distance-from-start.**

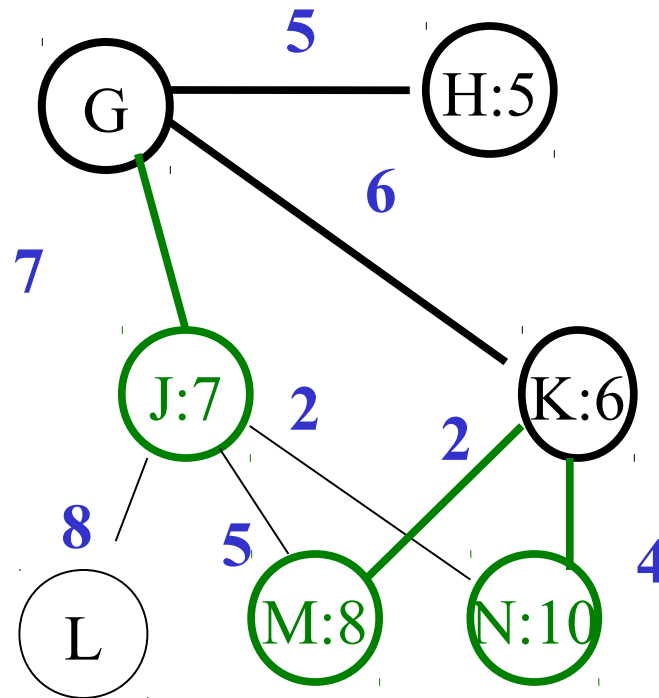
**Put  $v$  in the tree.**

**Update fringe**

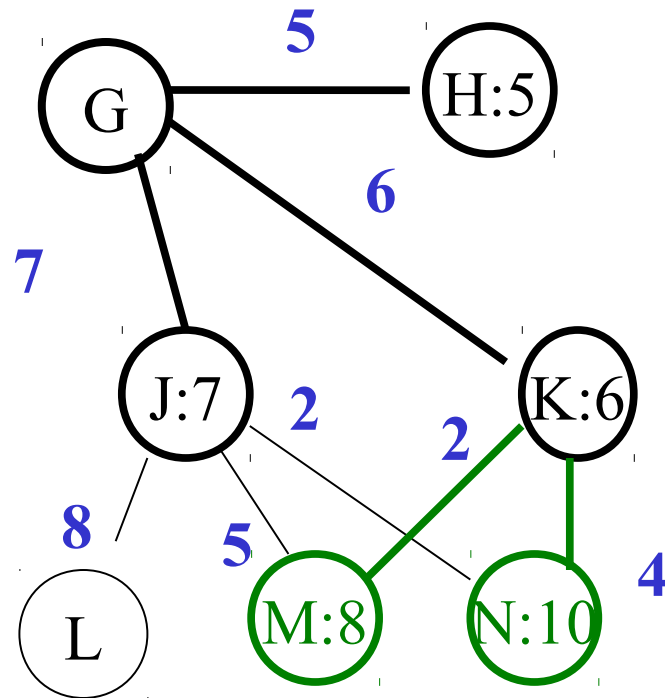
# Update fringe

- **Neighbors of  $v$  that are not in tree or fringe get added to fringe**
- **Neighbors of  $v$  that are in the fringe get checked: would changing link to be  $v$  result in a smaller distance? If so, change link and distance**

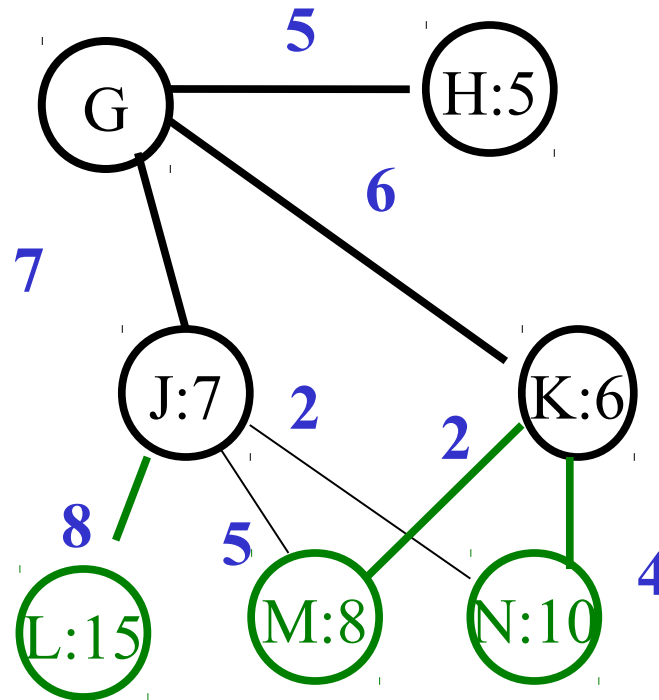
# **J $\rightarrow$ tree, Update fringe**



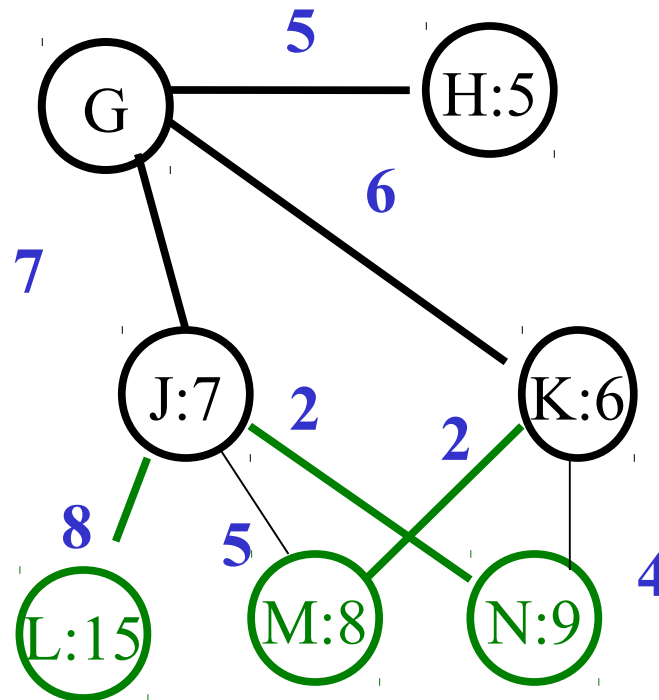
# **J → tree**



# Update fringe: neighbors $\rightarrow$ fringe



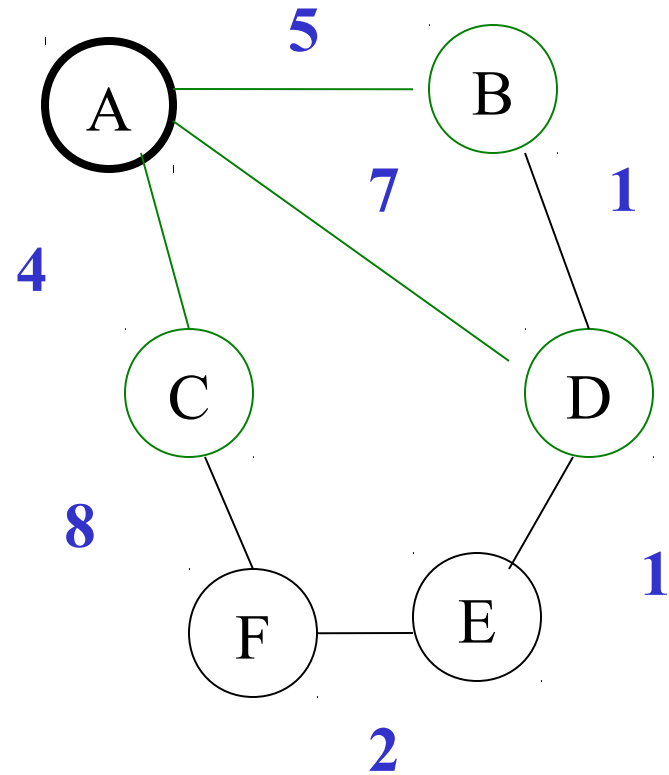
# Update fringe: Check neighbors' links





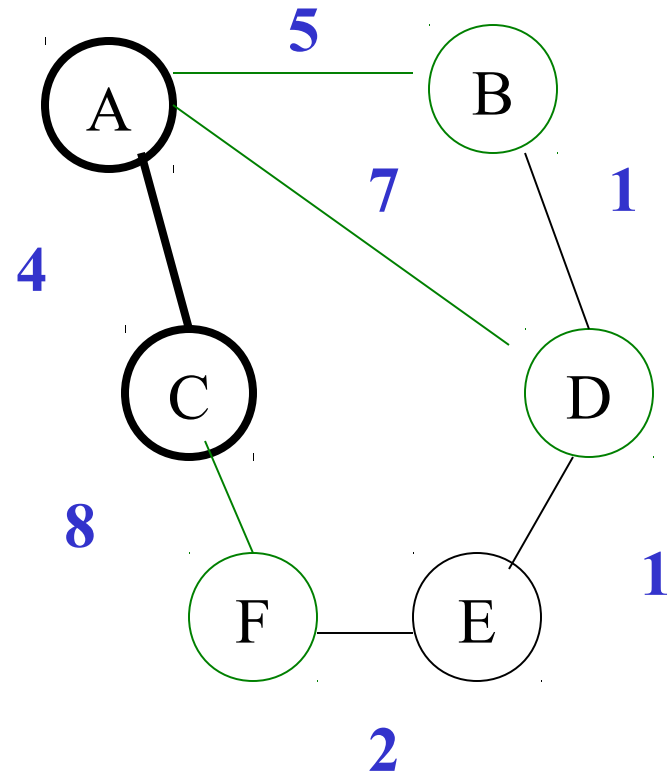
# Example

Node	Status	Link	Distance
A	Tree	--	0
B	Fringe	A	5
C	Fringe	A	4
D	Fringe	A	7
E			
F			



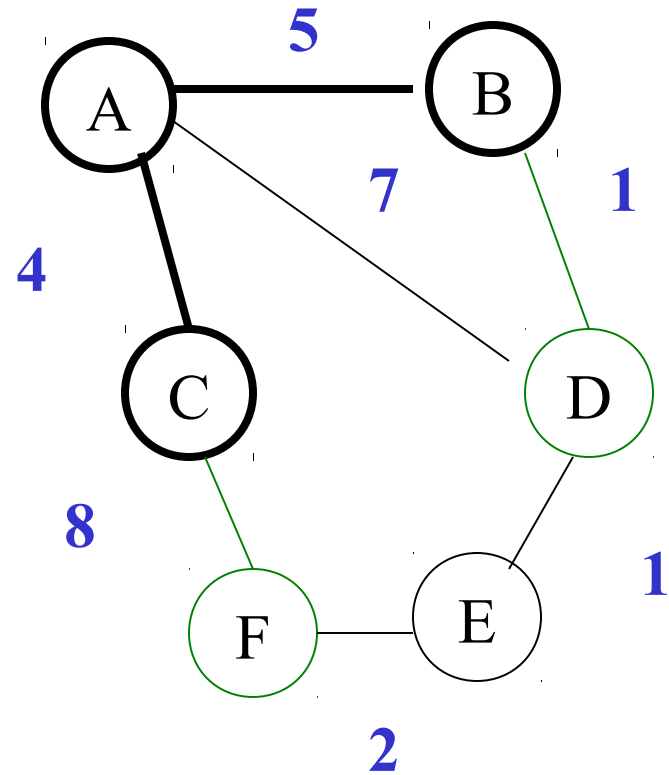
# Example

Node	Status	Link	Distance
A	Tree	--	0
B	Fringe	A	5
C	Tree	A	4
D	Fringe	A	7
E			
F	Fringe	C	12



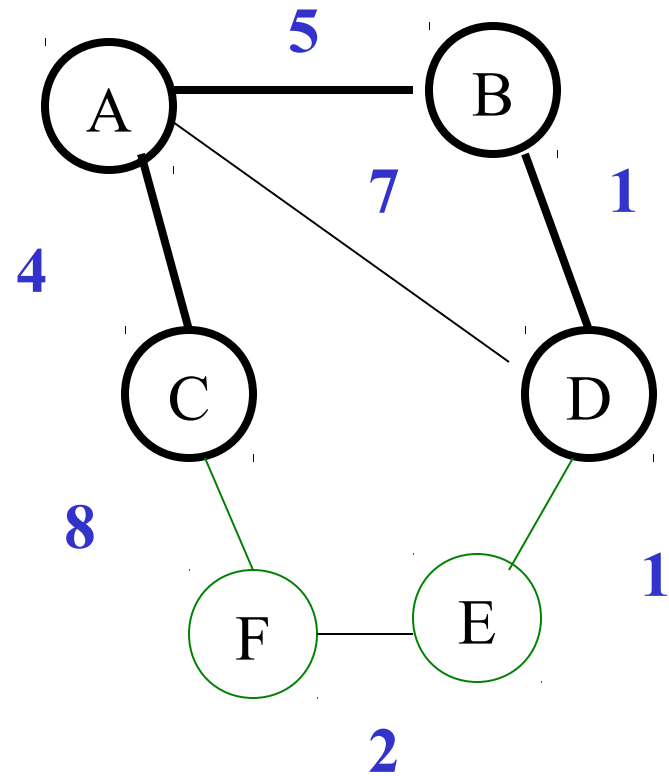
# Example

Node	Status	Link	Distance
A	Tree	--	0
B	Tree	A	5
C	Tree	A	4
D	Fringe	B	6
E			
F	Fringe	C	12



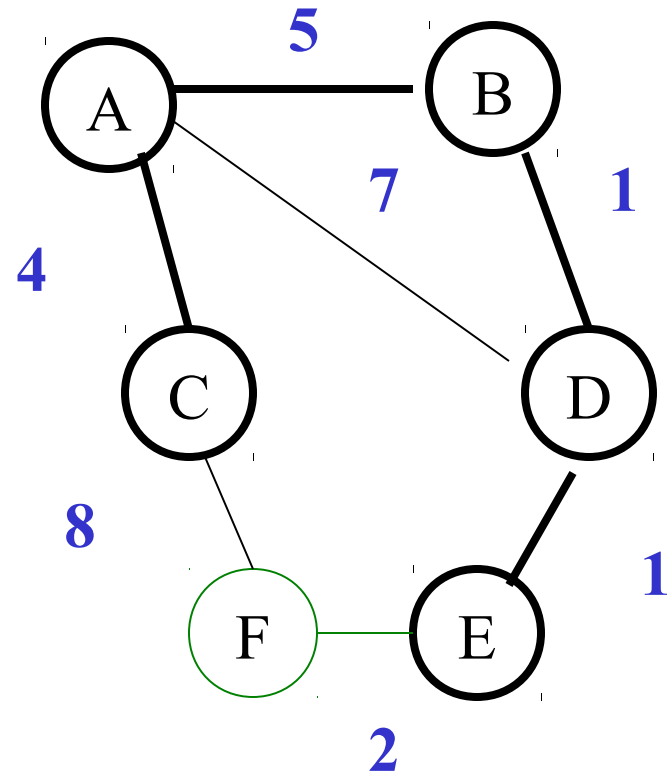
# Example

Node	Status	Link	Distance
A	Tree	--	0
B	Tree	A	5
C	Tree	A	4
D	Tree	B	6
E	Fringe	D	7
F	Fringe	C	12



# Example

Node	Status	Link	Distance
A	Tree	--	0
B	Tree	A	5
C	Tree	A	4
D	Tree	B	6
E	Tree	D	7
F	Fringe	E	9



# Example

Node	Status	Link	Distance
A	Tree	--	0
B	Tree	A	5
C	Tree	A	4
D	Tree	B	6
E	Tree	D	7
F	Tree	E	9

