

Computer Science 112

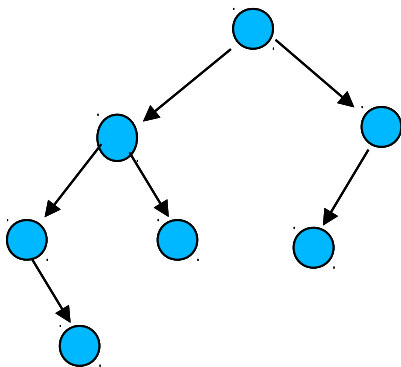
Data Structures

Lecture 13:

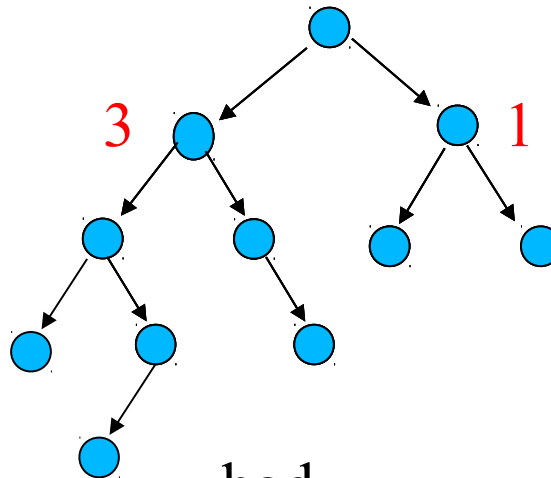
More AVL Trees

Review: AVL Trees

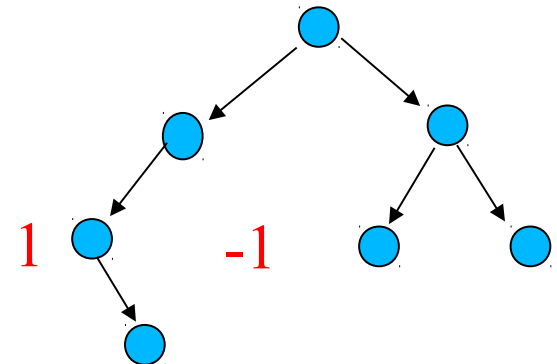
- **Binary Search Tree**
- **Almost balanced**
 - **At every node, subtree heights same ± 1**



good



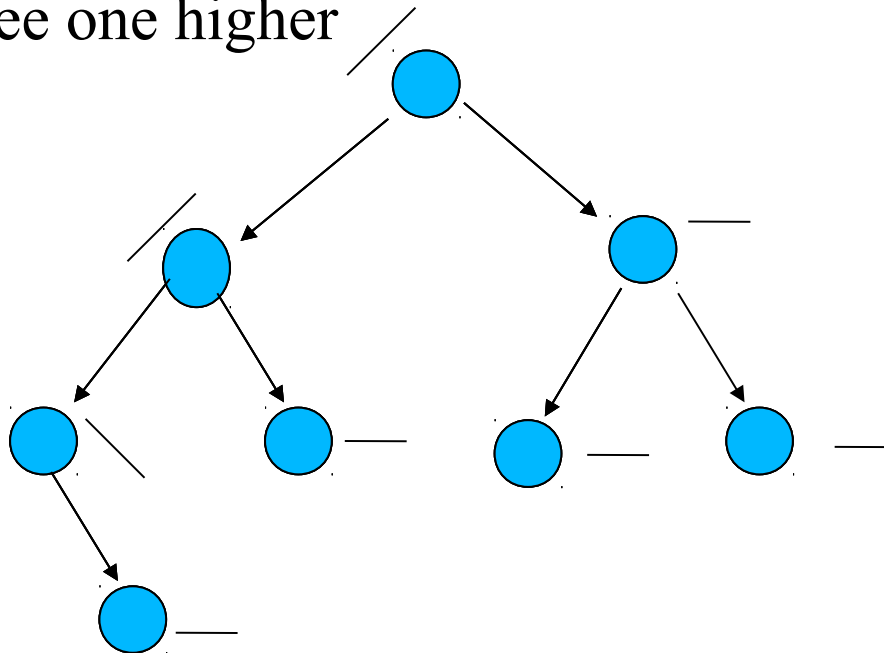
bad



bad
2

Labeling an AVL Tree

- **Label each node as**
 - left & right subtrees equally high
 - \ right subtree one higher
 - / left subtree one higher

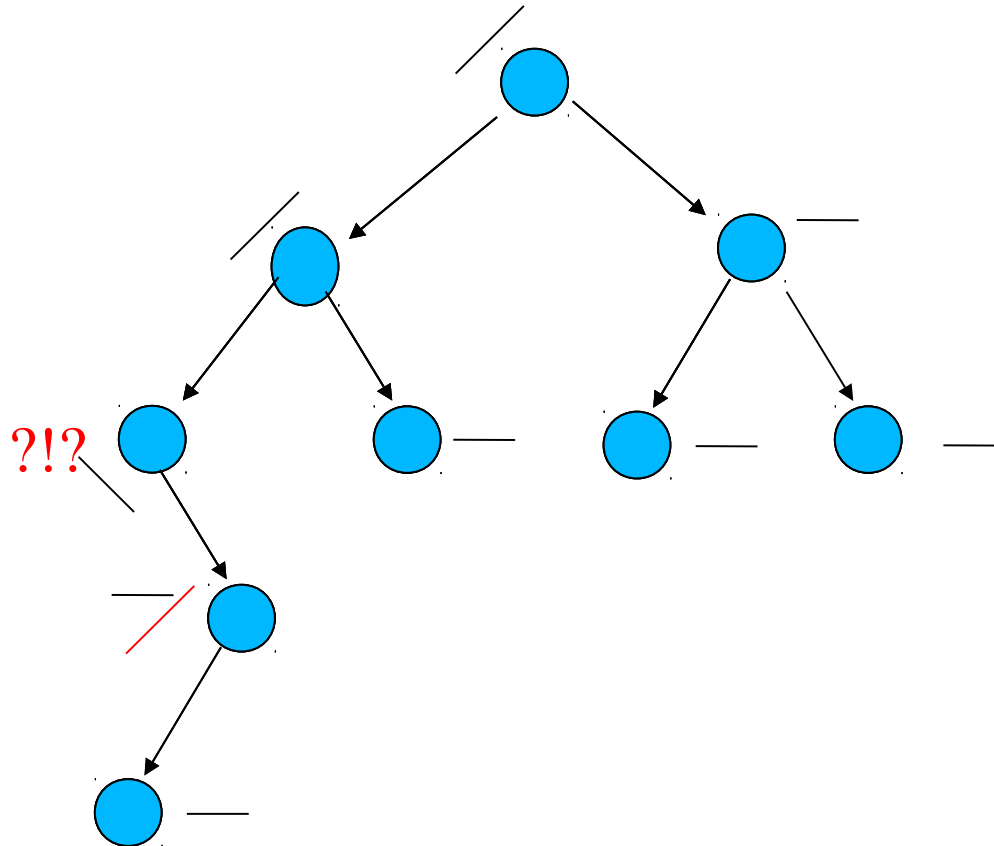


AVL Node

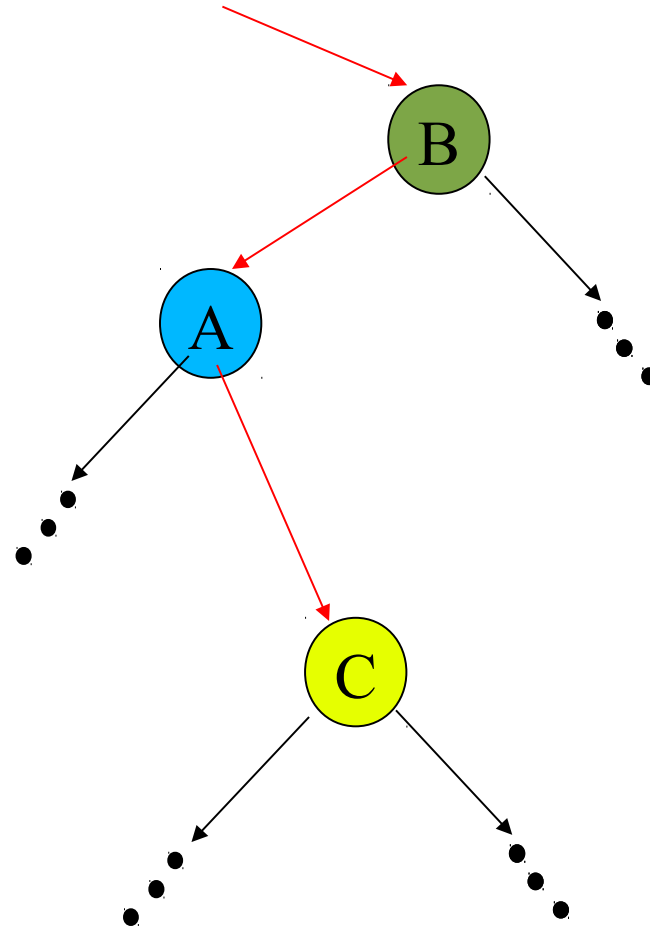
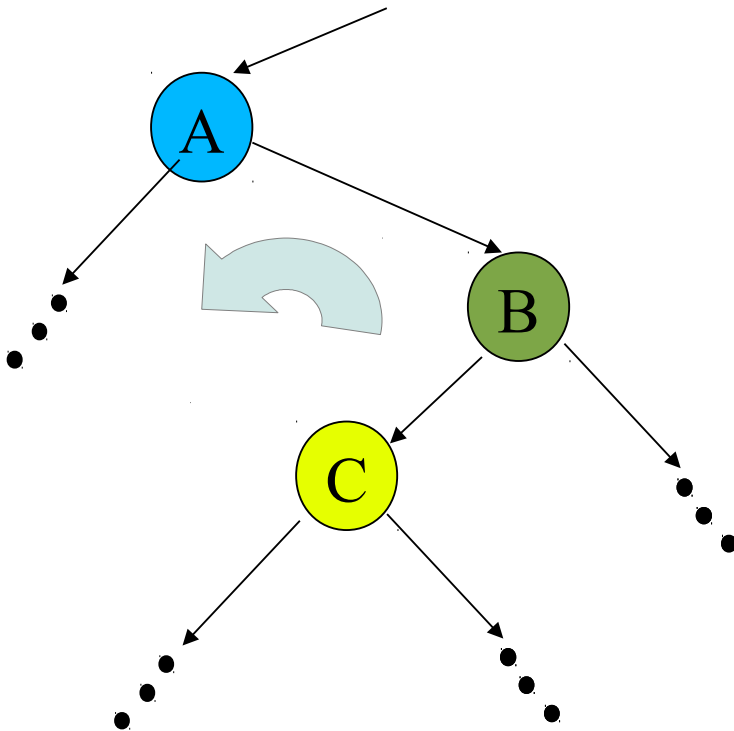
- **As in BSTNode**
 - left subtree, right subtree, data
- **Also:**
 - label
 - parent (may be convenient)
 - height (may be convenient)

Rebalancing

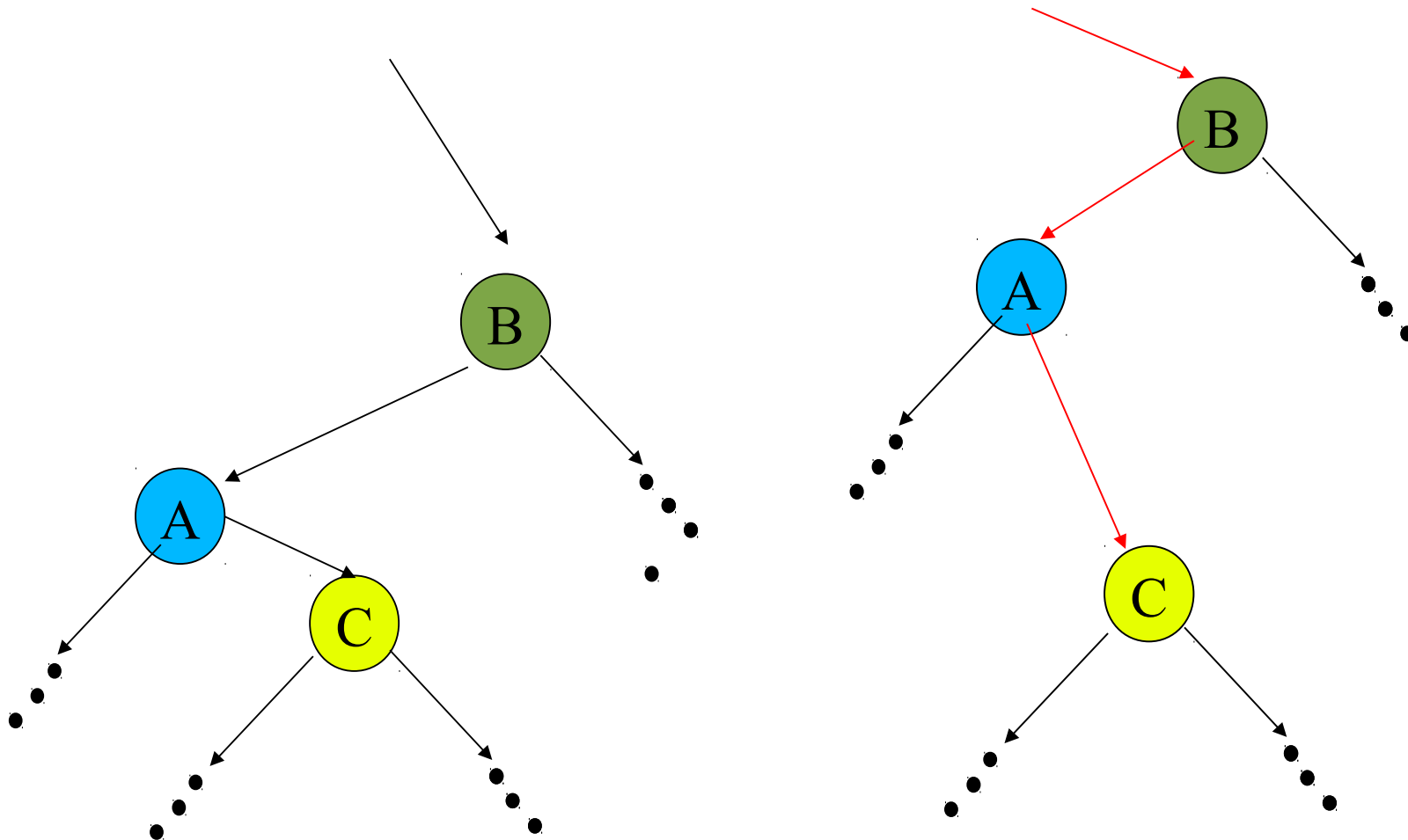
- **Problem: insert/delete -> not balanced**



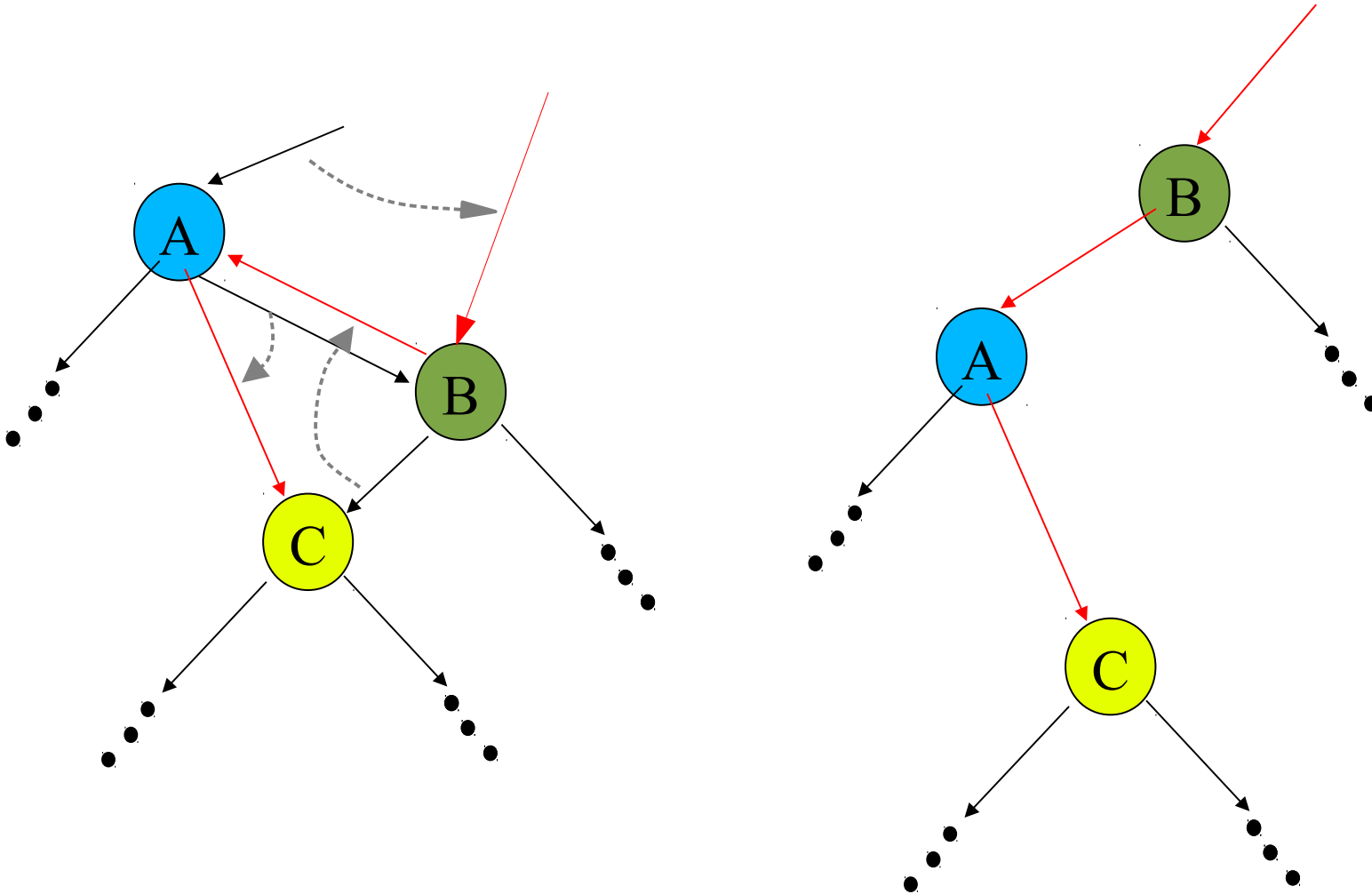
Rotation



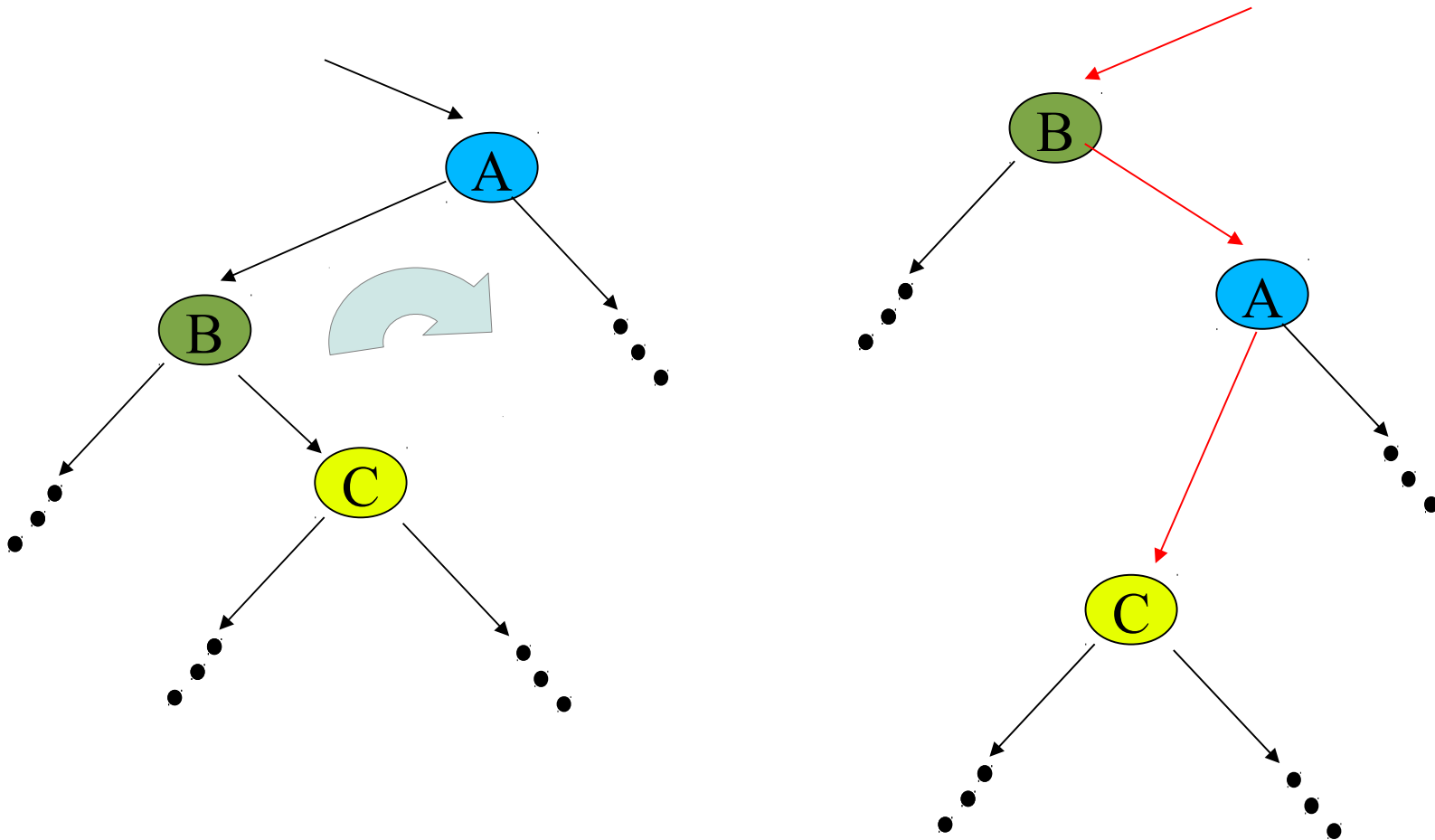
Rotation



Rotation

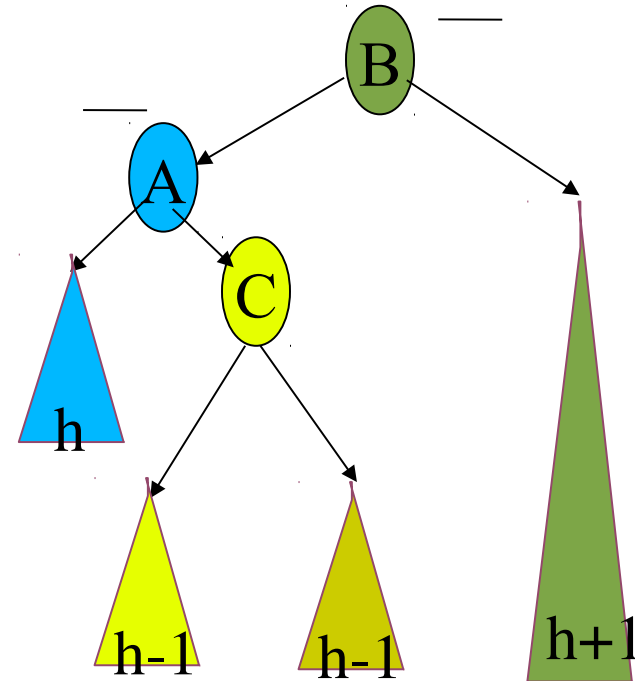
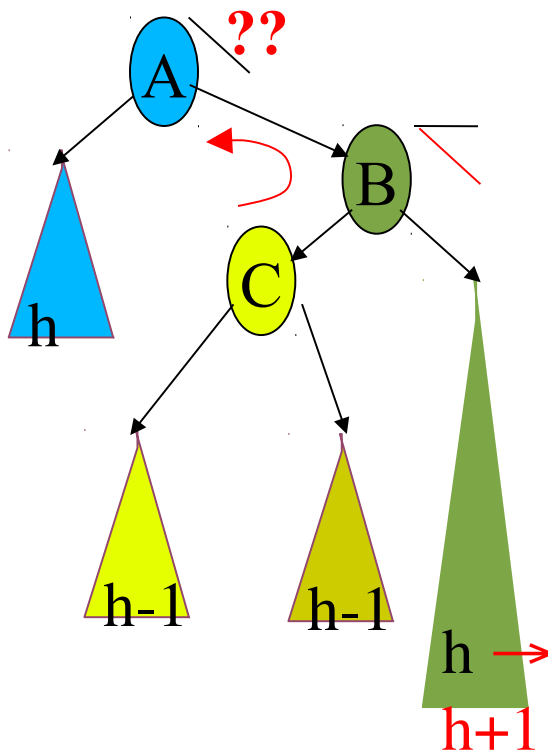


Mirror Image Rotation



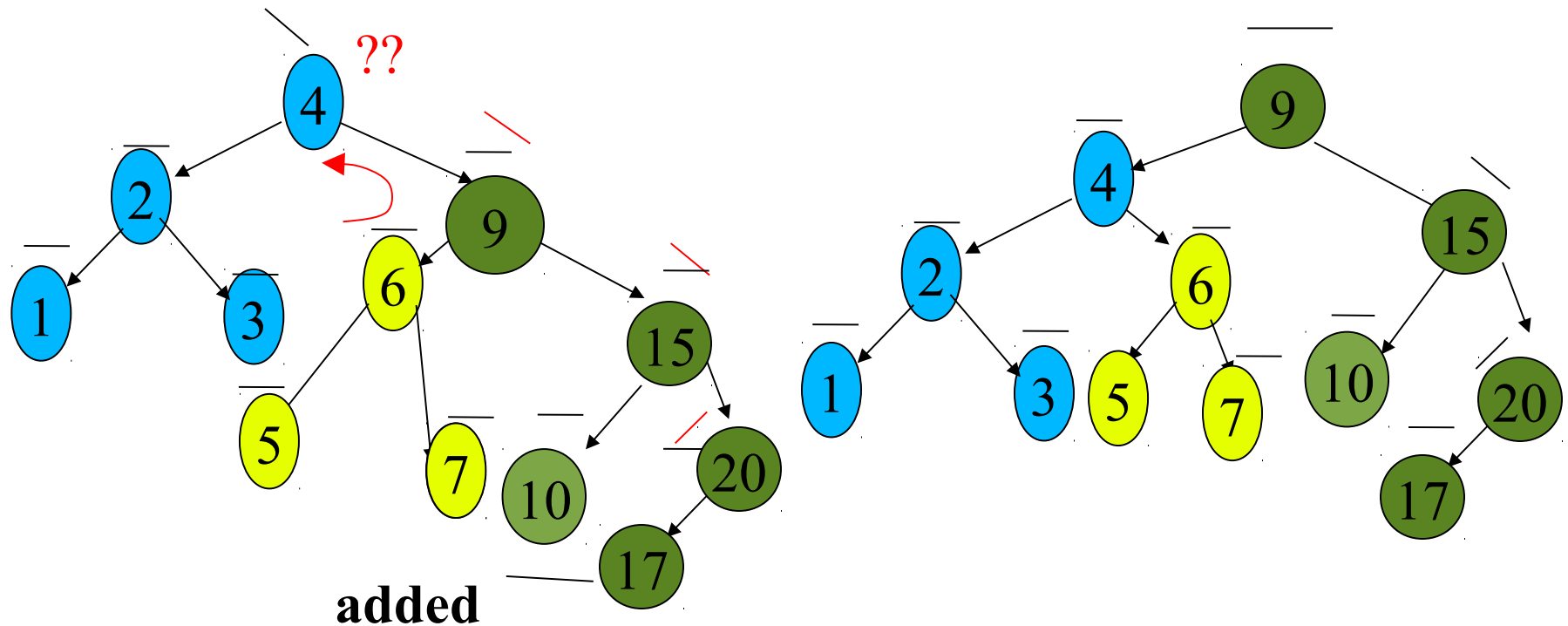
Rebalancing Using Rotation

- **Case 1: Highside child of A has same label as A**



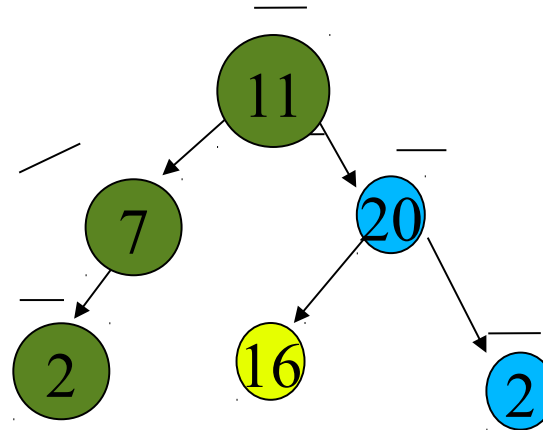
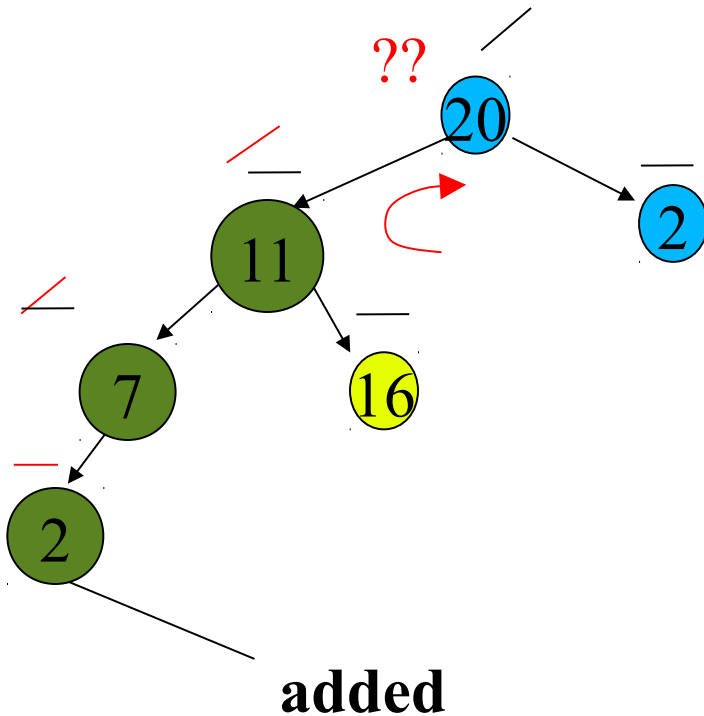
Rebalancing Using Rotation

- Case 1



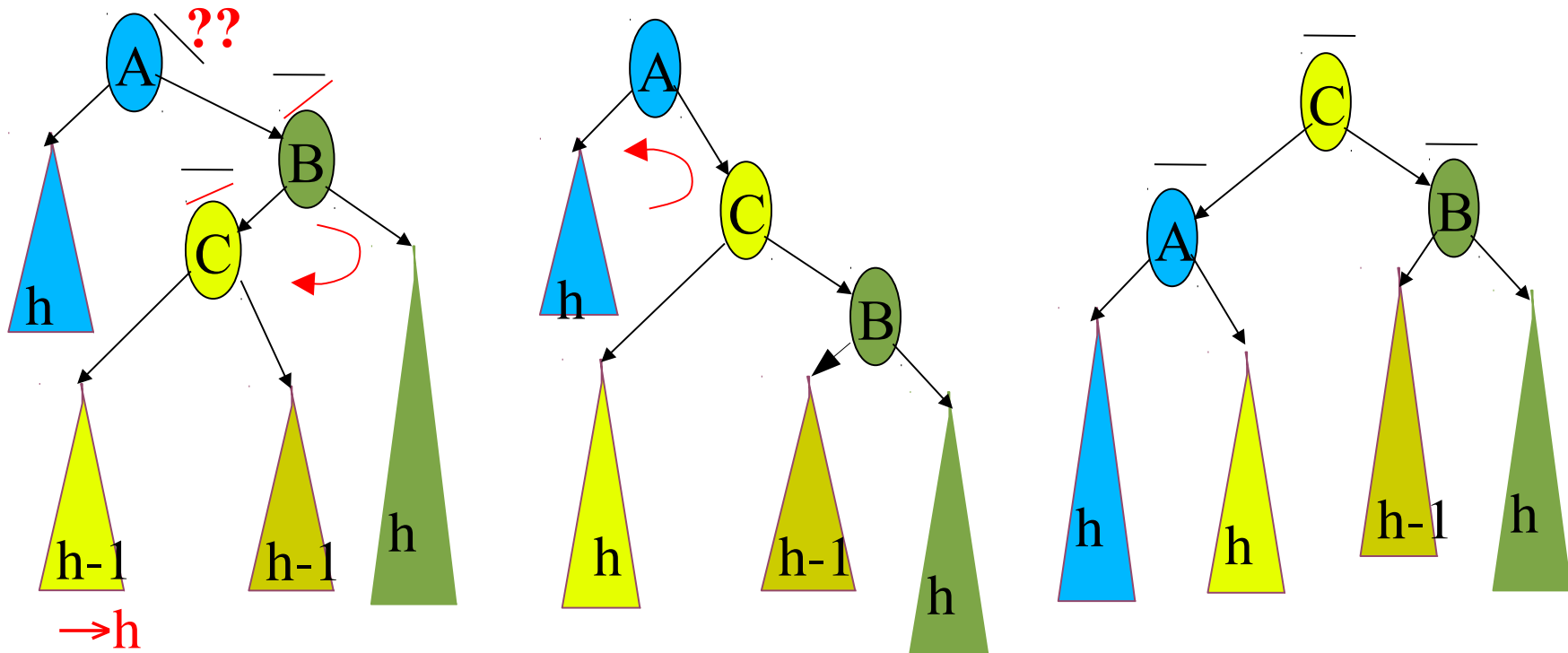
Rebalancing Using Rotation

- Case 1, mirror image



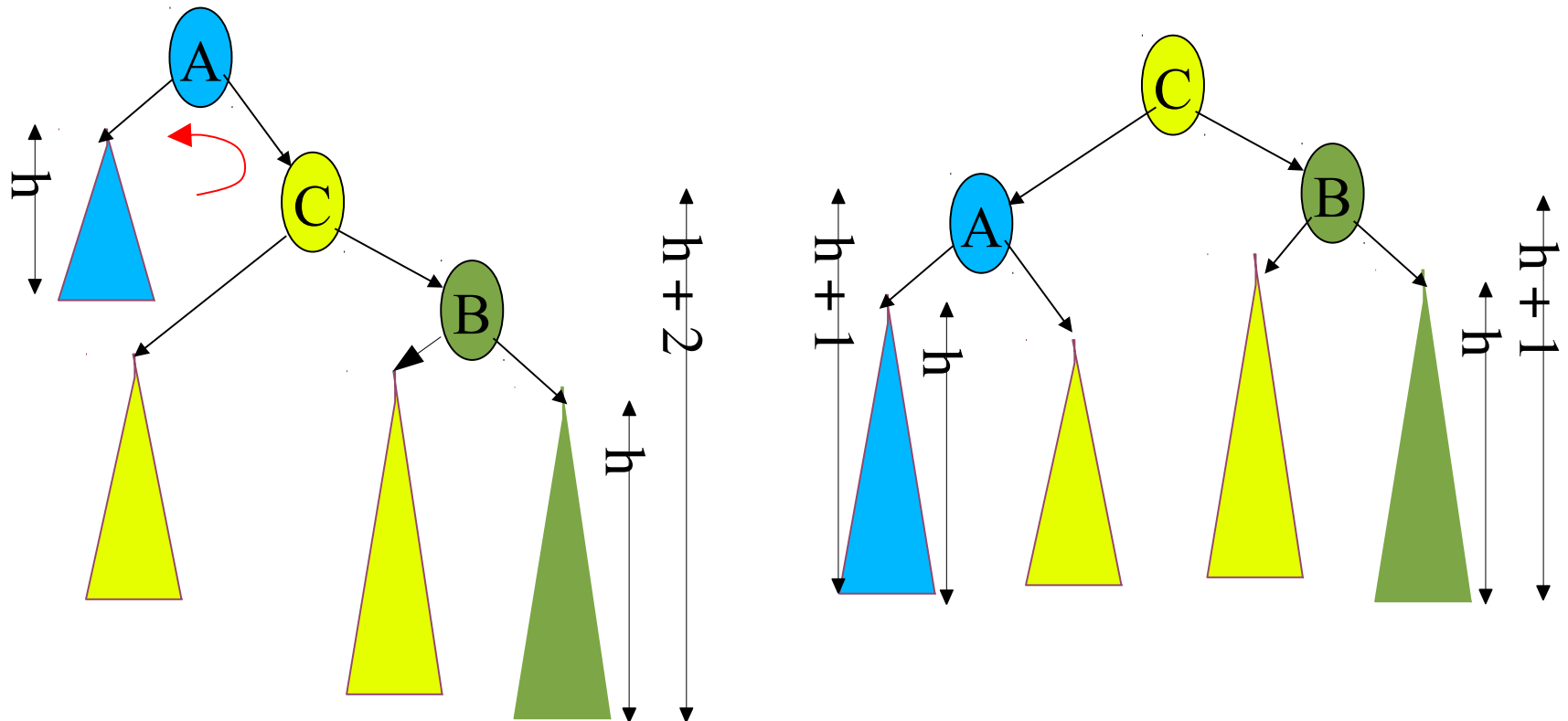
Rebalancing Using Rotation

- **Case 2: Highside child of A has opposite label from A**
 - **Two rotations: BC, then AC**



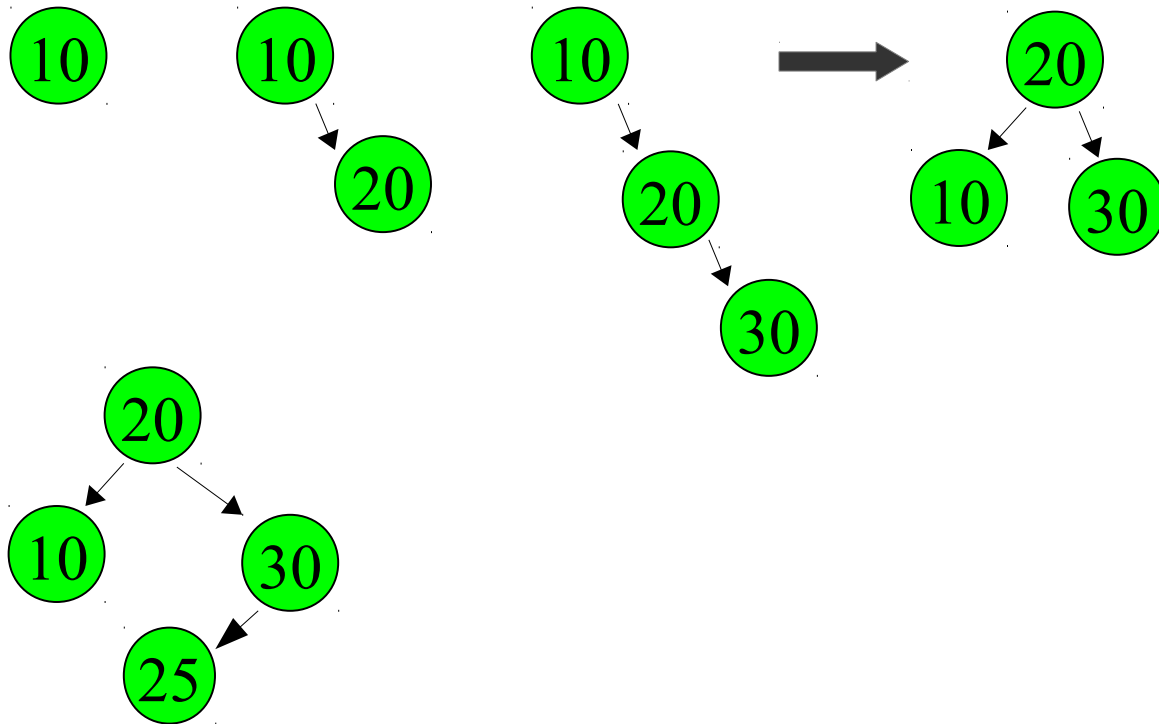
Rebalancing

- **Case 2: Highside child of A has opposite label from A**
 - Two rotations: BC, then AC



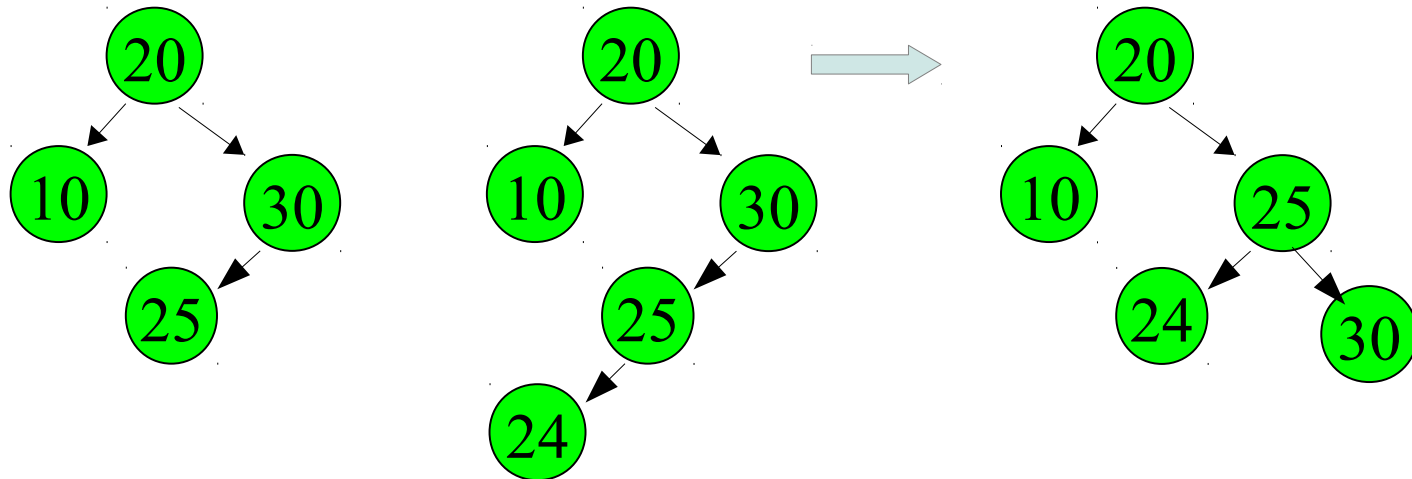
Example

- insert 10, 20, 30, 25, ...



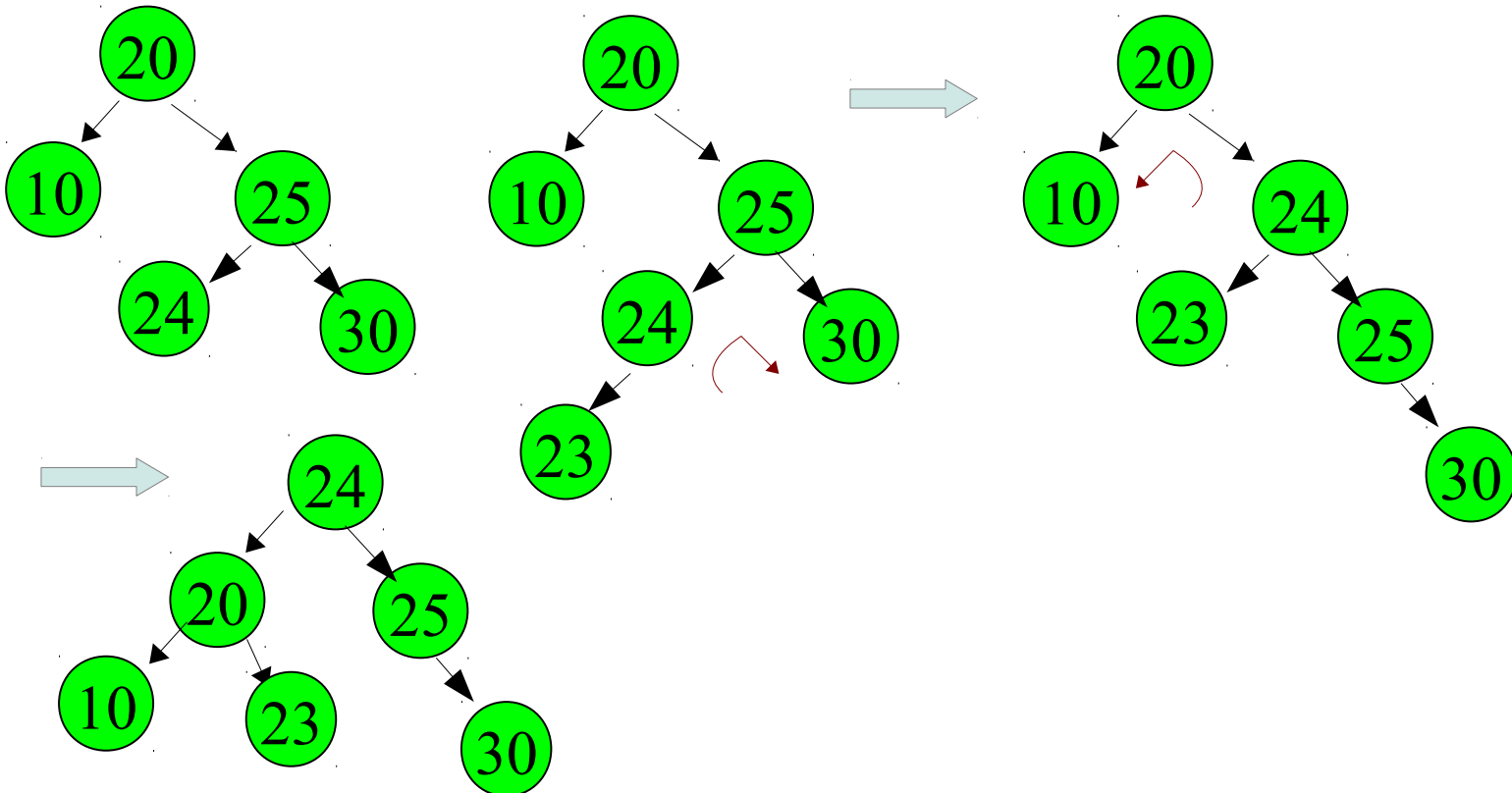
Example

- insert ..., 24, ...



Example

- insert ..., 23, ...



Big-O for insert

- **climb tree adjusting balance factors:**
 worst case: $O(\text{height of tree})$
- **rotate once or twice**
 worst case: $O(1)$
- **worst case total: $O(\text{height})$**
- **worst case height: $O(\log(\text{number of nodes}))$**
- **worst case insert: $O(\log(\text{number of nodes}))$**

Big-O for AVL Tree

- **insert:** as in BST then rebalance
 $O(\log(n))$
- **search:** just like Binary Search Tree
 $O(\log(n))$
- **Delete:** delete as in BST, then rebalance
 $O(\log(n))$