1. Estructura óptima del problema

$$e\left[i,j\right] = \left\{ \begin{array}{ll} q_{i-1} & ; & j=i-1 \\ & \displaystyle \min_{i \leq r \leq j} \left\{ e\left[i,r-1\right] + e\left[r+1,j\right] + w\left(i,j\right) \right\} & ; & i \leq j \end{array} \right.$$

$$w\left(i,j\right) = \sum_{l=i}^{j} p_{l} + \sum_{l=i-1}^{j} q_{l}$$

2. Algoritmo recurrente

```
1: procedure Build_OptBinTree(P,Q)
2: return Build_OptBinTree_Rec(P,Q,1,|P|)
3: end procedure
```

```
1: procedure BUILD OPTBINTREE Rec(P, Q, i, j)
        if j = i - 1 then
 2:
            return Q[i-1]
 3:
        else
            w \leftarrow \text{Dummy Weight}(P, Q, i, j)
 5:
 6:
            for r \leftarrow i to j do
 7:
                v_l \leftarrow \text{Build OptBinTree Rec}(P, Q, i, r - 1)
 8:
                v_r \leftarrow \text{Build\_OptBinTree\_Rec}\left(P, Q, r + 1, j\right)
 9:
10:
                v \leftarrow v_l + v_r + w
                if v < e then
11:
                    e \leftarrow v
12:
                end if
13:
14:
            end for
            \mathbf{return}\ e
15:
16:
        end if
17: end procedure
```

```
1: \operatorname{procedure} \operatorname{Dummy}_{-}\operatorname{Weight}(P,Q,i,j)
2: w \leftarrow Q[i-1]
3: \operatorname{for} l \leftarrow i \operatorname{to} j \operatorname{do}
4: w \leftarrow w + P[l] + Q[l]
5: \operatorname{end} \operatorname{for}
6: \operatorname{return} w
7: \operatorname{end} \operatorname{procedure}
```

3. Algoritmo recurrente «memoizado»

```
1: procedure BUILD OPTBINTREE(P,Q)
          let W be a matrix [1..|P|] \times [1..|P|]
          let M be a matrix [0..|P|] \times [0..|P|]
 3:
 4:
          W \leftarrow 0
          M \leftarrow 0
 5:
          for i \leftarrow 1 to |P| do
 6:
               W\left[i,i\right] \leftarrow Q\left[i-1\right] + P\left[i\right] + Q\left[i\right]
 7:
               M[i,i] \leftarrow Q[i-1]
 8:
               for j \leftarrow i + 1 to |P| do
 9:
                    W\left[i,j\right] \leftarrow W\left[i,j-1\right] + P\left[j\right] + Q\left[j\right]
10:
                    M\left[i,j\right] \leftarrow \infty
11:
               end for
12:
          end for
13:
          \mathbf{return}\ \mathtt{Build\_OptBinTree\_Rec}\left(P,Q,1,\left|P\right|,M,W\right)
14:
15: end procedure
```

```
1: procedure BUILD OPTBINTREE Rec(P, Q, i, j, M, W)
        if M[i,j] = \infty then
2:
           if j = i - 1 then
3:
               M[i,j] \leftarrow Q[i-1]
4:
            else
               for r \leftarrow i to j do
                   v_l \leftarrow \text{Build OptBinTree Rec}(P, Q, i, r - 1, M, W)
                   v_r \leftarrow \text{Build OptBinTree Rec}(P, Q, r+1, j, M, W)
8:
                   v \leftarrow v_l + v_r + W[i, j]
9:
                   if v < M[i, j] then
10:
                       M\left[i,j\right] \leftarrow v
11:
                   end if
12:
               end for
13:
           end if
14:
        end if
15:
        return M[i,j]
16:
17: end procedure
```

La tabla de memoización tiene la siguiente forma:

M[i,j]	0	1	2	3	4	5
0	0	0	0	0	0	0
1	Q[0]	M[1,1]	M[1,2]	M[1,3]	M[1,4]	M[1,5]
2	0	Q[1]	M[2,2]	M[2,3]	M[2,4]	M[2,5]
3	0	0	Q[2]	M[3,3]	M[3,4]	M[3,5]
4	0	0	0	Q[3]	M[4,4]	M[4,5]
5	0	0	0	0	Q[4]	M[5,5]

Y w puede ser pre-calculado:

$W\left[i,j ight]$	1	2	3	4	5
1	Q[0] + P[1] + Q[1]	W[1,1] + P[2] + Q[2]	W[1,2] + P[3] + Q[3]	W[1,3] + P[4] + Q[4]	W[1,4] + P[5] + Q[5]
2	0	$Q\left[1\right] + P\left[2\right] + Q\left[2\right]$	W[2,2] + P[3] + Q[3]	W[2,3] + P[4] + Q[4]	W[2,4] + P[5] + Q[5]
3	0	0	Q[2] + P[3] + Q[3]	W[3,3] + P[4] + Q[4]	W[3,4] + P[5] + Q[5]
4	0	0	0	Q[3] + P[4] + Q[4]	W[4,4] + P[5] + Q[5]
5	0	0	0	0	$Q\left[4\right] + P\left[5\right] + Q\left[5\right]$

4. Algoritmo «bottom-up» con la solución

```
1: procedure BUILD OPTBINTREE(P,Q)
         let W be a matrix [1..|P|] \times [1..|P|]
 2:
         let M be a matrix [0..|P|] \times [0..|P|]
 3:
         let R be a matrix [1..|P|] \times [1..|P|]
 4:
         W \leftarrow 0
         M \leftarrow 0
 6:
         R \leftarrow 0
 7:
 8:
         for i \leftarrow 1 to |P| do
             W\left[i,i\right] \leftarrow Q\left[i-1\right] + P\left[i\right] + Q\left[i\right]
 9:
             M[i,i] \leftarrow Q[i-1]
10:
             for j \leftarrow i + 1 to |P| do
11:
                 W[i, j] \leftarrow W[i, j-1] + P[j] + Q[j]
12:
                  M[i,j] \leftarrow \infty
13:
             end for
14:
         end for
15:
         for l \leftarrow 1 to |P| do
16:
             for i \leftarrow 1 to |P| - l + 1 do
17:
                  j \leftarrow i + l - 1
18:
19:
                  for r \leftarrow i to j do
                      v \leftarrow M[i, r - 1] + M[r + 1, j] + W[i, j]
20:
                      if v < M[i,j] then
21:
                           M[i,j] \leftarrow v
22:
                           R[i,j] \leftarrow r
23:
                      end if
24:
                  end for
25:
             end for
26:
         end for
27:
         return R
29: end procedure
```