

Análisis de la herramienta TLA+ Proof System

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EFD-sentences

Let \mathcal{L} be an algebraic first-order language.

An *Equational Function Definition Sentence* (*EFD for short*) is a sentence of the form

$$\varphi = \forall \vec{x} \exists! \vec{z} \bigwedge_{i=1}^k p_i(\vec{x}, \vec{z}) = q_i(\vec{x}, \vec{z}).$$

where p_i and q_i are \mathcal{L} -terms.

Algebraic Functions

Suppose φ is an EFD, and $\mathbf{A} \models \varphi$.

The *function defined by φ* is the map

$$[\varphi]^{\mathbf{A}} : \mathbf{A}^n \rightarrow \mathbf{A}^m$$

given by

$[\varphi]^{\mathbf{A}}(\vec{a}) = \text{unique } \vec{b} \text{ such that}$

$$\bigwedge_{i=1}^k p_i(\vec{a}, \vec{b}) = q_i(\vec{a}, \vec{b})$$

We denote by $[\varphi]_i^{\mathbf{A}}$ the coordinates of this map, ie, $[\varphi]_i^{\mathbf{A}} = \pi_i \circ [\varphi]^{\mathbf{A}}$

We use $E(\varphi)$ (resp. $U(\varphi)$) to denote the *existence sentence of φ* (resp. *unicity sentence of φ*):

$$\begin{aligned} E(\varphi) &= \forall \vec{x} \exists \vec{z} \bigwedge_{i=1}^k p_i(\vec{x}, \vec{z}) = q_i(\vec{x}, \vec{z}) \\ U(\varphi) &= \forall \vec{x} \forall \vec{z} \forall \vec{w} \bigwedge_{i=1}^k p_i(\vec{x}, \vec{z}) = q_i(\vec{x}, \vec{z}) \\ &\quad \wedge (\bigwedge_{i=1}^k p_i(\vec{x}, \vec{w}) = q_i(\vec{x}, \vec{w})) \rightarrow (\vec{z} = \vec{w}) \end{aligned}$$

Definition

A sentence of the form $[\varphi]_i^{\mathbf{A}}$ is called an *algebraic function over \mathbf{A}*

An *algebraically expandable (AE) class* is a class \mathcal{C} of algebras definable by EFDs,

i.e.: $\mathcal{C} = \text{Mod}(\Sigma)$, for some set Σ of EFDs

Definition

We say that **B** is a *daughter* of **A**
(or **A** is a *mother* of **B**) if

$$\mathbf{B} \in H(\mathbf{A}) \cap IS(\mathbf{A})$$

Remark

If

- ▶ φ is an EFD,
- ▶ $\mathbf{A} \models \varphi$ and
- ▶ **B** is a daughter of **A**

Then

$$\mathbf{B} \models \varphi$$

Global Products

Definition

A subdirect product

$$\mathbf{A} \subseteq_{SD} \prod_{i \in I} \{\mathbf{A}_i : i \in I\}$$

is *global* if there is a topology τ over I such that

- ▶ $\tau(E(x, y) : x, y \in \mathbf{A}) \subseteq \tau$
- ▶ If $x \in \prod_{i \in I} \{\mathbf{A}_i : i \in I\}$ and $E(x, y) \in \tau$ for all $y \in \mathbf{A}$, then $x \in \mathbf{A}$

Lemma

If φ is an EFD, $\{A_i : i \in I\} \models \varphi$ and \mathbf{A} is a global product of the algebras $\{\mathbf{A}_i : i \in I\}$ then $\mathbf{A} \models \varphi$

The case of \mathcal{D}_{01ab}

Lemma

Every lattice in \mathcal{D}_{01ab} can be represented as a global subdirect product with factors in the subclass \mathcal{G} of all two- and three-element members

Starting with Maximal Subclasses

$$Mod(\varphi) = \mathcal{G} - \{\mathbf{A}\}$$

$$Mod(\psi) = \mathcal{G} - \{\mathbf{B}\}$$

$$Mod(\varphi \wedge \psi) = \mathcal{G} - \{\mathbf{A}, \mathbf{B}\}$$

$\mathcal{G} - \{\mathbf{3}_{MM}\}$ is not AE

Proposition

If \mathcal{C} is an AE subclass of \mathcal{D}_{01ab} containing $\mathcal{G}' = \{\mathbf{2}_{00}, \mathbf{2}_{11}, \mathbf{3}_{0M}, \mathbf{3}_{M0}, \mathbf{3}_{M1}, \mathbf{3}_{1M}\}$ and any member \mathbf{L} of \mathcal{G} with $\{a^{\mathbf{L}}, b^{\mathbf{L}}\} = \{0, 1\}$, then $\mathbf{3}_{MM} \in \mathcal{C}$

If there existed an EFD φ such that $\mathcal{G}' \models \varphi$
and $\mathbf{3}_{MM} \not\models \varphi \dots$

1) We can assume

$$\varphi = \exists! z \bigwedge_{i=1}^k p_i(z) = q_i(z).$$

- Use the *translation property* that holds

in $(\mathcal{D}_{01ab})_{SI}$:

$$\mathbf{2}_{ab} \models \forall x y (x = y) \vee (z = w) \Leftrightarrow$$

$$z|_{x \vee y}^{x \vee y} = w|_{x \wedge y}^{x \vee y}$$

$$\text{where } x|_u^v = (x \vee u) \wedge v$$

2) Programming techniques...

... to see that φ should be

$$\exists! z (a \wedge b \wedge z = 0) \wedge (a \vee b \vee z = 1) \wedge \varepsilon(z)$$

for some set of equations $\varepsilon(z)$

3) Finally, more programming...

... to conclude that there is no possible choice for $\varepsilon(z)$.

But

$$\exists! z (a \wedge b \wedge z = 0) \wedge (a \vee b \vee z = 1)$$

holds for no \mathbf{L} in \mathcal{G} with $\{a^{\mathbf{L}}, b^{\mathbf{L}}\} = \{0, 1\}$

Maximal Classes Axiomatizations

φ	$Mod_{\mathcal{G}}(\varphi)$
$\exists!z (a \vee b) \wedge z = 0$	$\mathcal{G} - \{\mathbf{3}_{00}, \mathbf{2}_{00}\}$
$\exists!z (a \wedge b) \vee z = 1$	$\mathcal{G} - \{\mathbf{3}_{11}, \mathbf{2}_{11}\}$
$\forall x \exists!z (x _{a \vee b}^1 \wedge z = a \vee b) \wedge (x _{a \vee b}^1 \vee z = 1)$	$\mathcal{G} - \{\mathbf{3}_{00}\}$
$\forall x \exists!z (x _a^{a \vee b} \wedge z = a) \wedge (x _a^{a \vee b} \vee z = a \vee b)$	$\mathcal{G} - \{\mathbf{3}_{01}\}$
$\forall x \exists!z (x _0^{a \wedge b} \wedge z = 0) \wedge (x _0^{a \wedge b} \vee z = a \wedge b)$	$\mathcal{G} - \{\mathbf{3}_{11}\}$
$\forall x \exists!z (x _b^{a \vee b} \wedge z = b) \wedge (x _b^{a \vee b} \vee z = a \vee b)$	$\mathcal{G} - \{\mathbf{3}_{10}\}$
$\exists!z (a _b^1 \wedge z = b) \wedge (a _b^1 \vee z = 1)$	$\mathcal{G} - \{\mathbf{3}_{M0}\}$
$\exists!z (a _0^b \wedge z = 0) \wedge (a _0^b \vee z = b)$	$\mathcal{G} - \{\mathbf{3}_{M1}\}$
$\exists!z (b _a^1 \wedge z = a) \wedge (b _a^1 \vee z = 1)$	$\mathcal{G} - \{\mathbf{3}_{0M}\}$
$\exists!z (b _0^a \wedge z = 0) \wedge (b _0^a \vee z = a)$	$\mathcal{G} - \{\mathbf{3}_{1M}\}$

Proposition

Suppose \mathcal{C} is a subclass of \mathcal{G} closed under the relationship of daughtership that contains

$$\{2_{01}, 2_{10}, 3_{MM}\}$$

Then there is an EFD φ such that

$$\mathcal{C} = \text{Mod}(\varphi) \cap \mathcal{G}$$

References



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