Tel. 2402 7227

www.giovannini.edu.uy

Eduardo Acevedo 1462

Curso 2019 Lógica

Práctico 4: deducción natural en prop

1. Realice los agregados que estime pertinentes para que los siguientes árboles sean elementos de DER. Asimismo, provea juicios que correspondan con dichas pruebas i.e. dé $\Gamma \subseteq PROP$ y $\varphi \in PROP$ tales que cada razonamiento (árbol) pruebe que $\Gamma \vdash \varphi$.

(a) (b)
$$\frac{\frac{[g]}{b \to g} \quad [a \leftrightarrow b \to g] \quad [\neg a \land b]}{\frac{a}{-\frac{1}{2}g}} \qquad \frac{\frac{[c]}{d \to c}}{\frac{\frac{1}{2}c}{-\frac{1}{2}c}} \qquad \frac{\frac{1}{2}c}{\frac{1}{2}c}$$

$$\frac{(a) \quad (b) \quad (b) \quad (c) \quad (c)$$

(c)
$$\frac{[a \wedge \neg b]}{\frac{a}{a}} \quad [a \rightarrow b] \quad [a \wedge \neg b] \\ \frac{\underline{b} \qquad \qquad \neg b}{\underline{\neg b}}$$

$$\underline{[a \wedge \neg b]} \qquad \frac{\underline{\bot}}{\neg (a \rightarrow b)}$$

$$\underline{a \wedge \neg (a \rightarrow b)} \qquad \neg (a \wedge \neg (a \rightarrow b))$$

$$\underline{-(a \wedge \neg b)} \qquad \qquad (d)$$

$$\frac{[f] \quad [f \to p]}{\frac{p}{p \lor s}} \\
\frac{\frac{p}{p \lor s}}{\neg f \lor p \lor s} \quad [\neg (\neg f \lor p \lor s)] \\
\frac{\frac{\bot}{\neg (f \to p)}}{\neg (f \to p)} \quad \neg (f \to p) \to s \\
\frac{\frac{s}{p \lor s}}{\neg f \lor p \lor s} \quad [\neg (\neg f \lor p \lor s)] \\
\frac{\frac{\bot}{\neg f}}{\neg f \lor p \lor s} \quad [\neg (\neg f \lor p \lor s)] \\
\frac{\bot}{\neg f \lor p \lor s} \\

(p) \quad \frac{[p] \quad p \leftrightarrow \neg p}{\neg p} \quad [\neg p] \quad [\neg p]$$

(e)
$$\frac{[p] \quad p \leftrightarrow \neg p}{\neg p} \qquad \frac{[\neg p] \quad p \leftrightarrow \neg p}{p} \qquad \frac{\bot}{p} \qquad p \leftrightarrow \neg p}{\frac{\bot}{p} \qquad p \leftrightarrow \neg p}$$

Tel. 2402 7227

www.giovannini.edu.uy

Eduardo Acevedo 1462

Lógica

Curso 2019

2. Construya derivaciones que justifiquen los siguientes juicios.

(a)
$$\vdash \varphi \rightarrow \varphi$$

(b)
$$\vdash \bot \rightarrow \varphi$$

(c)
$$\vdash \neg(\varphi \land \neg\varphi)$$

(d) **[PP2007]**
$$\varphi \lor \varphi \vdash \varphi$$

(e)
$$\alpha \vee \beta \vdash \beta \vee \alpha$$

(f)
$$\neg \varphi \vdash \varphi \rightarrow (\bot \lor \neg \bot)^1$$

(g) **[ED2005]**
$$\vdash \varphi \lor \neg \varphi$$
 (*)

(h)
$$\neg(\varphi \rightarrow \psi) \rightarrow \sigma \vdash \neg \varphi \lor (\psi \lor \sigma)$$
 [**EF2014**]

$$(i) ~~ \clubsuit \vdash \neg\bot \leftrightarrow \neg\neg\neg\bot$$

3. Demuestre que:

(a) Si
$$\vdash \varphi$$
 entonces $\vdash \psi \lor \varphi$.

(c) Si
$$\vdash \varphi$$
 y $\vdash \psi$ entonces $\vdash \varphi \land (\psi \lor \sigma)$

(b) Si
$$\vdash \varphi$$
 entonces $\vdash \psi \rightarrow \varphi$.

(d)
$$\clubsuit \vdash \varphi$$
 si y solamente si $\vdash \varphi \leftrightarrow \neg \bot$

Ejercicios integradores.

Construya derivaciones que justifiquen los siguientes juicios.

1. **[PP1998]**
$$(\varphi \rightarrow \psi) \vdash ((\psi \rightarrow \sigma) \land \neg \sigma) \rightarrow \neg \varphi)$$

2. **[PP1999]**
$$\varphi \wedge \sigma \rightarrow \psi$$
, $\varphi \rightarrow \sigma \vdash \neg \psi \rightarrow \neg \varphi$

3. **[PP2000]**
$$\varphi \to (\sigma \lor \psi) \vdash \neg \sigma \to (\varphi \to \psi)$$

4. **[PP2001]**
$$(\neg \beta \rightarrow \neg \alpha) \land \alpha \vdash \beta$$

5. **[PP2002]**
$$\neg \psi \rightarrow \neg \varphi \vdash \varphi \rightarrow ((\psi \rightarrow \neg \varphi) \rightarrow \sigma)$$

6. **[PP2003]**
$$\neg(\alpha \land \neg(\alpha \rightarrow \beta)) \vdash \neg(\alpha \land \neg\beta)$$

7. **[PP2004]**
$$\vdash \neg(\alpha \lor \beta) \leftrightarrow (\neg \alpha \land \neg \beta)$$

[Ley de DeMorgan]

8. **[PP2005]**
$$(\alpha \land \neg \beta) \lor (\neg \alpha \land \beta) \vdash \neg(\alpha \leftrightarrow \beta)$$

9. **[PP2006]**
$$\neg \beta \rightarrow \gamma$$
, $\alpha \lor \beta \vdash g$

(*)

10. **[PP2007]**
$$\vdash (\neg \sigma \rightarrow (\neg \varphi \land \neg \psi)) \rightarrow ((\varphi \lor \psi) \rightarrow \sigma)$$

11. **[PP2008]**
$$\vdash (a \rightarrow b) \land (c \rightarrow b) \rightarrow (\neg \neg a \lor c) \rightarrow b$$

12. **[PP2009]**
$$\vdash \neg \alpha \land (\beta \rightarrow \gamma) \rightarrow ((\alpha \lor \beta) \rightarrow \gamma)$$

13. **[PP2010]**
$$\neg ((\psi \rightarrow \varphi) \lor \alpha) \vdash \neg \varphi \lor \neg \beta$$

¹Sugerencia: notar que $\vdash \neg \bot$.

Tel. 2402 7227

www.giovannini.edu.uy

Eduardo Acevedo 1462

14. **[PP2010]**
$$((\alpha \lor \beta) \to (\alpha \land \beta)) \leftrightarrow (\alpha \leftrightarrow \beta)$$

15. **[PP2011]**
$$\vdash ((\alpha \lor \beta) \to (\beta \land \neg \alpha)) \to \neg \alpha$$

16. **[PP2011]**
$$\neg (\varphi \land \psi) \leftrightarrow (\neg \varphi \lor \neg \neg \psi)$$

17. **[PP2012]**
$$(q \rightarrow \neg p) \rightarrow p \vdash \neg (p \rightarrow \neg q) \lor p$$

18. **[PP2012]**
$$\vdash ((p \rightarrow q) \rightarrow p) \rightarrow \neg(\neg p \land \neg(p \land q))$$

19. **[PP2013]**
$$\varphi \rightarrow \psi \vdash \neg \neg \psi \lor \neg \varphi$$

20. **[PP2013]**
$$\neg \varphi$$
, $\neg \alpha \lor \neg \beta \vdash \sigma \Rightarrow \neg \varphi \vdash \neg (\alpha \land \beta) \to \sigma$

21. **[PP2013]**
$$(p_0 \to p_1) \to \neg p_1 \vdash \neg(\neg(\neg p_1)) \lor \neg(p_0 \to p_1)$$

22. **[PP2014]**
$$(\alpha \rightarrow \beta) \rightarrow \gamma \vdash \gamma \lor \neg \beta$$

23. **[PP2014]**
$$\vdash (\alpha \leftrightarrow (\beta \rightarrow \gamma)) \rightarrow (\neg \alpha \land \beta \rightarrow \neg \gamma)$$

24. **[PP2015]**
$$\vdash ((\alpha \rightarrow \gamma) \land (\beta \rightarrow \gamma)) \leftrightarrow (\alpha \lor \beta \rightarrow \gamma)$$

25. **[PP2015]**
$$\vdash (\alpha \rightarrow \beta \rightarrow \gamma) \rightarrow \neg(\neg \gamma \lor \neg \beta) \rightarrow \gamma$$

26. **[PP2016]**
$$\neg (\beta \lor \alpha) \vdash \alpha \land \beta$$

27. **[PP2016]**
$$(\varphi \lor \psi) \leftrightarrow (\varphi \lor (\neg \varphi \land \psi))$$

28. **[PP2017]**
$$\gamma \to \alpha$$
, $\neg \to (\neg \beta \lor \alpha) \vdash \neg \alpha \to (\neg \beta \land \neg \gamma)$

29. **[PP2017]**
$$p_2 \to \neg p_1$$
, $p_3 \to p_1$, $\neg p_1 \to (\neg p_2 \lor p_1) \vdash \neg p_1 \to (\neg p_2 \land \neg p_3)$

30. **[PP2017]**
$$p_3 \to p_1$$
, $\neg p_1 \to (\neg p_2 \lor p_1) \vdash (p_2 \lor p_3) \to p_1$ [Riesgo de salud]

31. **[PP2018]**
$$p \leftrightarrow \neg p \vdash \bot$$

32. **[PP2018]**
$$p \lor q$$
, $p \to r$, $q \to s \vdash t \to \neg(\neg r \land \neg s)$

33. **[ED2008]**
$$\neg p \land q \rightarrow r$$
, $r \lor s \rightarrow \neg p$, $s \vdash \neg q \lor r$

34. **[ED2010]**
$$\alpha \vee \beta \leftrightarrow \beta \vdash \alpha \leftrightarrow \alpha \wedge \beta$$

35. **[ED2011]**
$$\neg p \land \neg q \rightarrow r \vdash (p \lor q) \lor r$$

36. **[EF2012]**
$$(\neg \alpha \lor \beta) \lor \gamma \vdash \neg \gamma \land \alpha \rightarrow \neg \neg \beta$$

Tel. 2402 7227

www.giovannini.edu.uy

Eduardo Acevedo 1462

Lógica Curso 2019

Soluciones propuestas al Práctico 4

A continuación se presenta un esbozo de las soluciones de los ejercicios, es decir, las mismas no están completas: falta agregar el inciso indicador de la cancelación de las hipótesis destacadas entre corchetes. Se alienta al lector a realizar tales agregados a efectos de profundizar la comprensión de dichas pruebas.

Asimismo, a efectos de no sobrecargar la notación se ha optado por emplear letras del alfabeto latino en lugar de aquellas del alfabeto griego.

- 1. Vistos en clase.
- 2. (i)

$$\frac{[\bot] \quad [\neg\bot]}{\overset{\bot}{\bot} RAA} E \neg \qquad \frac{\overset{[\bot]}{\bot} RAA}{\overset{\bot}{\neg\bot} I \neg} \qquad [\neg(\neg\bot)]}{\overset{\bot}{\neg(\neg(\neg\bot))} I \neg} E \neg \\
\frac{\overset{\bot}{\bot} RAA}{\xrightarrow{\neg\bot} I \neg} \qquad \xrightarrow{\neg\bot} I \rightarrow$$

Otra forma más sencilla:

$$\frac{\left[\bot\right]}{\neg\bot} \, I \neg \qquad \frac{\left[\neg\bot\right] \quad \left[\neg(\neg\bot)\right]}{\neg(\neg(\neg\bot))} \, I \neg}_{\neg\bot \leftrightarrow \neg(\neg(\neg\bot))} \, I \leftrightarrow$$

3. Vistos en clase.

Ejercicios integradores.

1.

$$\frac{[a] \quad a \to b}{b} E \to \frac{[(b \to c) \land \neg c]}{b \to c} E \land_{1} \quad \underbrace{[(b \to c) \land \neg c]}_{\neg c} E \land_{2}$$

$$\frac{\frac{\bot}{\neg a} I \neg}{(b \to c) \land \neg c \to \neg a} I \to$$

Tel. 2402 7227

www.giovannini.edu.uy

Eduardo Acevedo 1462

Lógica Curso 2019

3.

$$\frac{[a] \quad a \to b \lor c}{b \lor c} E \to \frac{[b] \quad [\neg b]}{\frac{1}{c} E \bot} E \neg \frac{[c] \quad [\neg b]}{\frac{c \land \neg b}{c} E \land_{1}} I \land \\
\frac{\frac{c}{a \to c} I \to}{\neg b \to a \to c} I \to$$

4.

$$\frac{(\neg b \to \neg a) \land a}{\underbrace{a}} E \land_2 \underbrace{ \begin{bmatrix} \neg b \end{bmatrix}} \underbrace{ \frac{(\neg b \to \neg a) \land a}{\neg b \to \neg a}}_{\underline{-a}} E \xrightarrow{} \underbrace{E} \land_1$$

$$\frac{\bot}{b} RAA$$

5.

$$\underbrace{ \begin{array}{c|c} [a] & [a \to \neg b] \\ \hline -b & E \neg \\ \hline \\ \underline{\frac{\bot}{\neg a} I \neg} & \neg a \to \neg b \\ \hline \hline & \underline{-b} E \neg \\ \hline \\ \underline{\frac{\bot}{c} E \bot} \\ (a \to \neg b) \to c I \to \\ \hline \\ b \to (a \to \neg b) \to c I \to \\ \hline \end{array} } E \to$$

6.

$$\frac{[a \wedge \neg b]}{\frac{a}{a}} E \wedge_1 \quad [a \to b]}{\frac{b}{b}} E \to \frac{[a \wedge \neg b]}{\neg b} E \wedge_2$$

$$\frac{[a \wedge \neg b]}{\frac{a}{a}} E \wedge_1 \qquad \frac{\bot}{\neg (a \to b)} I \neg$$

$$\frac{a \wedge \neg (a \to b)}{a \wedge \neg (a \to b)} I \neg$$

$$\frac{\bot}{\neg (a \wedge \neg b)} I \neg$$

$$\underbrace{ \begin{bmatrix} \neg a \end{bmatrix} \quad [\neg a \rightarrow \neg b \wedge \neg c]}_{ \begin{array}{c} \underline{\neg b \wedge \neg c} \\ \overline{\neg b \wedge \neg c} \\ E \wedge 1 \\ \end{array}}_{ \begin{array}{c} \underline{b \vee c} \\ \end{array}} E \rightarrow \underbrace{ \begin{array}{c} [\neg a] \quad [\neg a \rightarrow \neg b \wedge \neg c] \\ \underline{\neg b \wedge \neg c} \\ E \wedge 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}} E \rightarrow \underbrace{ \begin{array}{c} \underline{\neg b \wedge \neg c} \\ \underline{\neg b \wedge \neg c} \\ E \wedge 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ E \neg 2 \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg c} \\ \end{array}}_{ \begin{array}{c} \underline{b \wedge \neg$$

Tel. 2402 7227

www.giovannini.edu.uy Eduardo Acevedo 1462

Curso 2019 Lógica

$\frac{1}{a} E_{\wedge 2}$ $\frac{[-a \wedge -b]}{-b} E_{\wedge 2} \frac{[a]}{a \vee b} I^{\vee_1} \frac{[-a \wedge b]}{[-a \vee b]} E_{\wedge} \frac{[b]}{a \vee b} I^{\vee_2} \frac{[-a \vee b]}{[-a \vee b]} E_{\wedge}$ $\frac{1}{a \vee b} I^{\vee_2} \frac{[-a \vee b]}{[-a \vee b]} E_{\wedge}$ $\frac{1}{a \vee b} I^{\vee_2} \frac{[-a \vee b]}{[-a \vee b]} I^{\vee_2}$ $\frac{1}{a \vee b} I^{\vee_2}$		$\frac{ g }{[-(g \vee \neg g)]} E_{\neg}$ $\frac{\perp}{a} RAA$ $\frac{\perp}{g \vee \neg g} RAA$ $\frac{\perp}{a \vee g} RAA$ $\frac{ g }{g \vee \neg g} E_{\wedge} I_{\wedge}$ $\frac{ g }{a \vee g} E_{\wedge}$ $\frac{ g }{a \vee g} E_{\wedge}$ $\frac{ g }{a \vee g} E_{\wedge}$ $\frac{ g }{g} E_{\wedge}$ $\frac{ g }{g} E_{\wedge}$ $\frac{ g }{g} RAA$ $\frac{ g }{g} RAA$
$ \frac{[a]}{[a \lor b]} \frac{[\neg a \land \neg b]}{b} E \land_1 [b] \frac{[\neg a \land \neg b]}{\neg b} E \land_2 $ $ \frac{\bot}{b} E \bot \qquad \frac{\bot}{b} E \bot \qquad \frac{\bot}{b} E \bot \qquad \frac{\bot}{b} E \bot \qquad \frac{\bot}{a \lor b} I \rightharpoonup $	$\frac{[a \land \neg b]}{a} E^{\land 1} [a \leftrightarrow b]$ $(a \land \neg b) \lor \neg a \land b$	$\frac{[g]}{g \vee \neg g} I^{\vee_1} [\neg (g \vee \neg g)]} E_{\neg}$ $\frac{\bot}{g \vee \neg g} I^{\vee_2}$ $[\neg b] \neg b \rightarrow g E \rightarrow$ $\frac{\bot}{g} RAA$ $\frac{\bot}{a \vee b} RAA$

Tel. 2402 7227

www.giovannini.edu.uy

Eduardo Acevedo 1462

Tel. 2402 7227

www.giovannini.edu.uy

Eduardo Acevedo 1462

Lógica Curso 2019

12.

$$\underbrace{ \begin{bmatrix} a \end{bmatrix} \quad \frac{ \begin{bmatrix} \neg a \wedge (b \to g) \end{bmatrix}}{\neg a} E \land_1 }_{ \begin{matrix} \underline{b} \\ \end{matrix} E \bot} \quad \underbrace{ \begin{bmatrix} b \end{bmatrix}}_{ \begin{matrix} E \lor \end{matrix} } E \lor \quad \underbrace{ \begin{bmatrix} \neg a \wedge (b \to g) \end{bmatrix}}_{ \begin{matrix} b \to g \\ \end{matrix} E \land_2 }$$

$$\underbrace{ \begin{matrix} \underline{b} \\ \end{matrix} \quad \underbrace{ \begin{matrix} g \\ a \lor b \to g \end{matrix}}_{ \begin{matrix} E \lor \end{matrix} } I \to }_{ \begin{matrix} \underline{a} \land (b \to g) \end{matrix} \to g } E \land_2$$

13.

$$\frac{\frac{[c]}{d \to c} I \to}{\frac{(d \to c) \vee a}{I \vee_1} \qquad \neg ((d \to c) \vee a)} E \neg \frac{\frac{\bot}{\neg c} I \neg}{\neg c \vee \neg b} I \vee_1$$

15.

$$\frac{\frac{[a]}{a \vee b} I \vee_{1} \quad [a \vee b \to b \wedge \neg a]}{\frac{b \wedge \neg a}{\neg a} E \wedge_{2}} E \to \frac{\frac{\bot}{\neg a} I \neg}{(a \vee b \to b \wedge \neg a) \to \neg a} I \to$$

17.

$$\frac{\frac{[\neg p]}{q \to \neg p} I \to (q \to \neg p) \to p}{\frac{p}{\frac{\bot}{p} RAA} E \to [\neg p]} E \to \frac{\frac{\bot}{p} RAA}{\neg (p \to \neg q) \lor p} I \lor_{2}$$

$$\frac{[p] \quad \frac{[\neg p \land \neg (p \land q)]}{\neg p} E \land 1}{\frac{\frac{1}{q} E \bot}{p \rightarrow q} I \rightarrow} E \land 1$$

$$\frac{p}{\frac{[(p \rightarrow q) \rightarrow p]}{p} E \rightarrow} \frac{[\neg p \land \neg (p \land q)]}{\neg p} E \land 1$$

$$\frac{p}{\frac{1}{q} (\neg p \land \neg (p \land q))} I \neg$$

$$\frac{((p \rightarrow q) \rightarrow p) \rightarrow \neg (\neg p \land \neg (p \land q))}{((p \rightarrow q) \rightarrow p) \rightarrow \neg (\neg p \land \neg (p \land q))} I \rightarrow$$

www.giovannini.edu.uy

Eduardo Acevedo 1462



Tel. 2402 7227

www.giovannini.edu.uy

Eduardo Acevedo 1462

$$\frac{[f] \ [\neg p]}{f \land \neg p} \ I \land \frac{[f] \ [\neg p]}{\neg f} \ I \land \frac{[f] \ [\neg p]}{\neg f} \ I \land \frac{[f] \ [\neg p]}{\neg f} \ E \land \frac{[f] \ [\neg p]}{\neg f \lor \neg (\neg p)} \ I \land \frac{[f] \ [\neg p]}{\neg f \lor \neg (\neg p)} \ I \land \frac{[f] \ [\neg p] \ [f] \ [\neg f] \ [f] \ [f$$

Tel. 2402 7227

www.giovannini.edu.uy

Eduardo Acevedo 1462

Lógica

Curso 2019

19.

$$\frac{[f] \quad f \to p}{\frac{p}{p}} E \to [\neg p] E \neg \frac{\frac{\bot}{\neg f} I \neg}{\frac{\bot}{\neg (\neg p) \lor \neg f} I \lor_2} [\neg (\neg (\neg p) \lor \neg f)]} E \neg \frac{\frac{\bot}{\neg (\neg p)} I \neg}{\frac{\bot}{\neg (\neg p) \lor \neg f} I \lor_1} [\neg (\neg (\neg p) \lor \neg f)]} E \neg \frac{\bot}{\neg (\neg p) \lor \neg f} RAA$$
(tautología dem. en [PP2014])

20.

22.

$$\frac{[\neg b]}{g \vee \neg b} I \vee_2 \qquad [\neg (g \vee \neg b)] E \neg$$

$$\frac{\frac{1}{b} RAA}{a \to b} I \to \qquad (a \to b) \to g E \to$$

$$\frac{g}{g \vee \neg b} I \vee_1 \qquad [\neg (g \vee \neg b)] E \neg$$

$$\frac{\frac{1}{b} RAA}{a \to b} I \to \qquad [\neg (g \vee \neg b)] E \neg$$

$$\frac{1}{b} RAA \qquad [\neg (g \vee \neg b)] E \neg$$

$$\frac{1}{b} RAA \qquad [\neg (g \vee \neg b)] E \neg$$

$$\frac{1}{b} RAA \qquad [\neg (g \vee \neg b)] E \neg$$

$$\frac{1}{b} RAA \qquad [\neg (g \vee \neg b)] E \neg$$

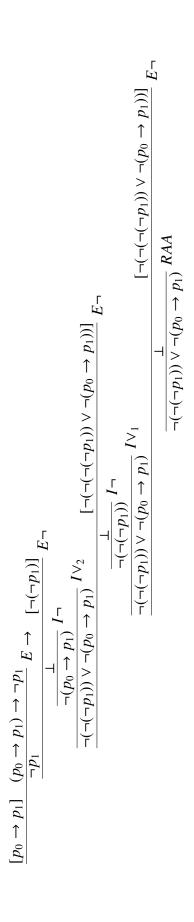
$$\frac{\frac{[g]}{b \to g} I \to \frac{[a \leftrightarrow b \to g]}{a} E \leftrightarrow_1 \frac{[\neg a \land b]}{\neg a} E \land_1}{\frac{\frac{1}{\neg g} I \neg}{\neg a \land b \to \neg g} I \to} E \land_1$$

Tel. 2402 7227

www.giovannini.edu.uy

Eduardo Acevedo 1462

Curso 2019 Lógica



www.giovannini.edu.uy

Eduardo Acevedo 1462

$$\frac{[a]}{a \lor b} I^{\vee_1} \xrightarrow{[a \lor b \to g]} E \to \frac{[b]}{a \lor b} I^{\vee_2} \xrightarrow{[a \lor b \to g]} E \to \frac{[a \lor b \to g]}{[a \lor b]} E \to \frac{[a \lor b]}{[a \lor b]} E \to \frac{[a \lor g) \land (b \to g)]}{[a \lor b]} E \to \frac{[a \lor g) \land (b \to g)]}{[a \lor b]} E \to \frac{[a \lor b]}{[a \lor$$

Tel. 2402 7227

www.giovannini.edu.uy

Eduardo Acevedo 1462

Lógica Curso 2019

25.

$$\frac{\frac{[\neg g]}{\neg g \vee \neg b} I \vee_1 \qquad [\neg (\neg g \vee \neg b)]}{\frac{\frac{1}{g} RAA}{\neg (\neg g \vee \neg b) \rightarrow g} I \rightarrow} E \neg$$

$$\frac{(a \rightarrow b \rightarrow g) \rightarrow \neg (\neg g \vee \neg b) \rightarrow g}{(a \rightarrow b \rightarrow g) \rightarrow \neg (\neg g \vee \neg b) \rightarrow g} I \rightarrow$$

26.

28.

$$\frac{[\neg a] \quad \neg a \rightarrow \neg b \vee a}{\neg b \vee a} E \rightarrow [\neg b] \quad \frac{[a] \quad [\neg a]}{\neg b} E \rightarrow [g] \quad g \rightarrow a \\
E \rightarrow [\neg a] E \rightarrow [\neg a]$$

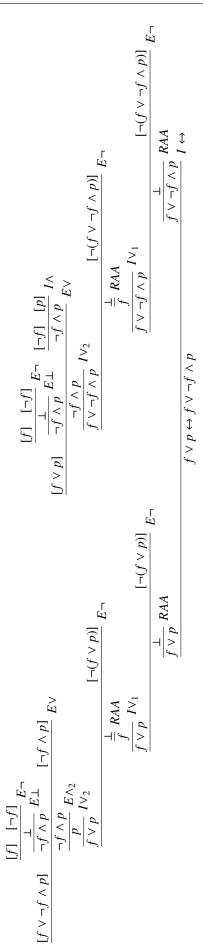
$$\frac{[p] \quad p \to r}{\frac{r}{E}} E \to \frac{[\neg r \land \neg s]}{\neg r} E \land 1$$

$$\frac{[p] \qquad \frac{\bot}{\neg p} I \neg}{\frac{\bot}{q} E \bot} \qquad [q] E \lor \qquad q \to s E \to \frac{[\neg r \land \neg s]}{\neg s} E \land 2$$

$$\frac{s}{\frac{\bot}{\neg (\neg r \land \neg s)} I \neg} I \rightarrow$$

www.giovannini.edu.uy

Eduardo Acevedo 1462

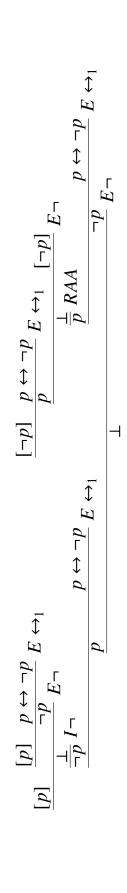


Tel. 2402 7227

www.giovannini.edu.uy

Eduardo Acevedo 1462

$\frac{[p_3] p_3 \to p_1}{p_1} E \to \frac{[\neg p_1]}{\neg p_3} E \to \frac{\bot}{I_{\bullet}}$	
$ \frac{[p_2] p_2 \to \neg p_1}{\neg p_1} E \to \neg p_1 \to \neg p_2 \lor p_1} E \to \frac{\neg p_2 \lor p_1}{\neg p_2 \lor p_1} E \to \frac{\bot}{\neg p_2} E \to \frac{\bot}{\neg p_2} I \neg p_2 $	$\frac{\neg p_2 \land \neg p_3}{\neg p_1 \rightarrow \neg p_2 \land \neg p_3} I \rightarrow$
[<i>p</i> ₂]	



Tel. 2402 7227

www.giovannini.edu.uy

Eduardo Acevedo 1462

Lógica Curso 2019

33.

$$\frac{\frac{s}{r \vee s} I \vee_{2} \quad r \vee s \rightarrow \neg p}{\frac{\neg p}{p} E \rightarrow \quad [q]} I \wedge \quad \neg p \wedge q \rightarrow r} E \rightarrow \frac{\frac{\neg p \wedge q}{\neg q \vee r} I \vee_{2}}{\frac{\neg q}{\neg q} I \neg} E \rightarrow \frac{\frac{\bot}{\neg q} I \neg}{\neg q \vee r} I \vee_{1}} E \rightarrow \frac{\frac{\bot}{\neg q} I \neg}{\frac{\neg q}{\neg q} \vee r} E \rightarrow \frac{\bot}{\neg q} I \neg}{\frac{\bot}{\neg q} \vee r} E \rightarrow \frac{[\neg (\neg q \vee r)]}{\neg q} E \rightarrow \frac{[\neg ($$

34.

$$\underbrace{ \begin{bmatrix} a \\ \hline a \lor b \end{bmatrix}}_{ \underline{a} \lor b} I \lor_{1} \quad \underline{a \lor b \leftrightarrow b}}_{ \underline{a} \lor b} E \leftrightarrow_{1}$$

$$\underbrace{ \begin{bmatrix} a \land b \\ \hline a \\ \hline a \leftrightarrow a \land b \end{bmatrix}}_{ \underline{a} \lor b} I \land$$

Última actualización: 5 de abril de 2019.

www.giovannini.edu.uy

Eduardo Acevedo 1462

$$\frac{[p]}{\frac{p \vee q}{p}} I^{\vee_1} \underbrace{[\neg (p \vee q)]}_{-p \wedge \neg q} E_{\neg} \underbrace{\frac{[q]}{\frac{p}{p} \vee q}}_{-p \wedge \neg q} I^{\vee_2} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge \neg q} E_{\neg} \underbrace{\frac{\bot}{\frac{p \vee q}{p \vee q}} I^{\vee_1}}_{-p \wedge \neg q} \underbrace{\frac{\bot}{\frac{p \vee q}{p \vee q}} R^{AA}}_{-p \wedge \neg q} \underbrace{\frac{\bot}{\frac{p \vee q}{p \vee q} N^{\vee_1}}}_{-p \wedge \neg q} \underbrace{[\neg (p \vee q) \vee r)]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]]}_{-p \wedge q \wedge r} E_{\neg} E_{\neg} \underbrace{[\neg (p \vee q) \vee r]}_{-p \wedge r} E_{\neg} E_{\neg}$$