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## Práctico 4: deducción natural en prop

1. Realice los agregados que estime pertinentes para que los siguientes árboles sean elementos de DER. Asimismo, provea juicios que correspondan con dichas pruebas i.e. dé  $\Gamma \subseteq PROP$  y  $\varphi \in PROP$  tales que cada razonamiento (árbol) pruebe que  $\Gamma \vdash \varphi$ .

(a) 
$$\frac{[g]}{b \to g} \underbrace{[a \leftrightarrow b \to g]} \underbrace{[\neg a \land b]}_{\neg a} \underbrace{\frac{[c]}{d \to c}}_{\underline{d \to c}}$$

$$\underbrace{\frac{\frac{\bot}{\neg g}}{\neg a \land b \to \neg g}}_{\overline{(a \leftrightarrow b \to g)} \to \neg a \land b \to \neg g} \underbrace{\frac{\bot}{\neg c}}_{\overline{\neg c} \lor \neg b}$$

(c) 
$$\frac{[a \wedge \neg b]}{a} \quad [a \rightarrow b] \quad [a \wedge \neg b] \\ \underline{b} \quad \neg b \\ \underline{a} \quad \boxed{\frac{b}{\neg b}} \\ \underline{a} \quad \boxed{\frac{\bot}{\neg (a \rightarrow b)}} \\ \underline{a} \wedge \neg (a \rightarrow b) \quad \neg (a \wedge \neg (a \rightarrow b)) \\ \underline{a} \wedge \neg (a \wedge \neg b) \quad \boxed{\frac{\bot}{\neg (a \wedge \neg b)}}$$
(d) 
$$\underline{-\frac{\bot}{\neg (a \wedge \neg b)}}$$

$$\frac{[f] \quad [f \to p]}{\frac{p}{p \lor s}}$$

$$\frac{\frac{1}{\neg (f \to p)} \quad \neg (f \to p) \to s}{\frac{\frac{s}{p \lor s}}{\neg (f \to p)} \quad \neg (f \to p) \to s}$$

$$\frac{\frac{s}{p \lor s}}{\neg f \lor p \lor s} \quad [\neg (\neg f \lor p \lor s)]$$

$$\frac{\frac{\bot}{\neg f}}{\neg f \lor p \lor s} \quad [\neg (\neg f \lor p \lor s)]$$

$$\frac{\bot}{\neg f \lor p \lor s}$$

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2. Construya derivaciones que justifiquen los siguientes juicios.

(a) 
$$\varphi \vdash \varphi$$

(b) 
$$\vdash \varphi \rightarrow \varphi$$

(c) 
$$\vdash \bot \rightarrow \varphi$$

(d) 
$$\vdash \neg(\varphi \land \neg\varphi)$$

(e) **[PP2007]** 
$$\varphi \lor \varphi \vdash \varphi$$

(f) 
$$\alpha \vee \beta \vdash \beta \vee \alpha$$

(g) 
$$\neg \varphi \vdash \varphi \rightarrow (\bot \lor \neg \bot)^1$$

(h) **[ED2005]** 
$$\vdash \varphi \lor \neg \varphi$$
 (\*)

(i) 
$$\neg(\varphi \rightarrow \psi) \rightarrow \sigma \vdash \neg \varphi \lor (\psi \lor \sigma)$$
 [EF2014]

$$(j) \clubsuit \vdash \neg \bot \leftrightarrow \neg \neg \neg \bot$$

3. Demuestre que:

(a) Si 
$$\vdash \varphi$$
 entonces  $\vdash \psi \lor \varphi$ .

(c) Si 
$$\vdash \varphi$$
 y  $\vdash \psi$  entonces  $\vdash \varphi \land (\psi \lor \sigma)$ 

(b) Si 
$$\vdash \varphi$$
 entonces  $\vdash \psi \rightarrow \varphi$ .

(d) 
$$\clubsuit \vdash \varphi$$
 si y solamente si  $\vdash \varphi \leftrightarrow \neg \bot$ 

Ejercicios integradores.

Construya derivaciones que justifiquen los siguientes juicios.

1. **[PP1998]** 
$$(\varphi \rightarrow \psi) \vdash ((\psi \rightarrow \sigma) \land \neg \sigma) \rightarrow \neg \varphi)$$

2. **[PP1999]** 
$$\varphi \wedge \sigma \rightarrow \psi$$
,  $\varphi \rightarrow \sigma \vdash \neg \psi \rightarrow \neg \varphi$ 

3. **[PP2000]** 
$$\varphi \to (\sigma \lor \psi) \vdash \neg \sigma \to (\varphi \to \psi)$$

4. **[PP2001]** 
$$(\neg \beta \rightarrow \neg \alpha) \land \alpha \vdash \beta$$

5. **[PP2002]** 
$$\neg \psi \rightarrow \neg \varphi \vdash \varphi \rightarrow ((\psi \rightarrow \neg \varphi) \rightarrow \sigma)$$

6. **[PP2003]** 
$$\neg(\alpha \land \neg(\alpha \rightarrow \beta)) \vdash \neg(\alpha \land \neg\beta)$$

7. **[PP2004]** 
$$\vdash \neg(\alpha \lor \beta) \leftrightarrow (\neg \alpha \land \neg \beta)$$

[Ley de DeMorgan]

8. **[PP2005]** 
$$(\alpha \land \neg \beta) \lor (\neg \alpha \land \beta) \vdash \neg(\alpha \leftrightarrow \beta)$$

9. **[PP2006]** 
$$\neg \beta \rightarrow \gamma$$
,  $\alpha \lor \beta \vdash g$  (\*)

10. **[PP2007]** 
$$\vdash (\neg \sigma \rightarrow (\neg \varphi \land \neg \psi)) \rightarrow ((\varphi \lor \psi) \rightarrow \sigma)$$

11. **[PP2008]** 
$$\vdash (a \rightarrow b) \land (c \rightarrow b) \rightarrow (\neg \neg a \lor c) \rightarrow b$$

12. **[PP2009]** 
$$\vdash \neg \alpha \land (\beta \rightarrow \gamma) \rightarrow ((\alpha \lor \beta) \rightarrow \gamma)$$

 $<sup>^{1}</sup>$ Sugerencia: notar que  $\vdash \neg \bot$ .

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13. **[PP2010]** 
$$\neg ((\psi \rightarrow \varphi) \lor \alpha) \vdash \neg \varphi \lor \neg \beta$$

14. **[PP2010]** 
$$((\alpha \lor \beta) \to (\alpha \land \beta)) \leftrightarrow (\alpha \leftrightarrow \beta)$$

15. **[PP2011]** 
$$\vdash ((\alpha \lor \beta) \to (\beta \land \neg \alpha)) \to \neg \alpha$$

16. **[PP2011]** 
$$\neg(\varphi \land \psi) \leftrightarrow (\neg \varphi \lor \neg \neg \psi)$$

17. **[PP2012]** 
$$(q \rightarrow \neg p) \rightarrow p \vdash \neg (p \rightarrow \neg q) \lor p$$

18. **[PP2012]** 
$$\vdash ((p \rightarrow q) \rightarrow p) \rightarrow \neg(\neg p \land \neg(p \land q))$$

19. **[PP2013**] 
$$\varphi \rightarrow \psi \vdash \neg \neg \psi \lor \neg \varphi$$

20. **[PP2013]** 
$$\neg \varphi$$
,  $\neg \alpha \lor \neg \beta \vdash \sigma \Rightarrow \neg \varphi \vdash \neg (\alpha \land \beta) \to \sigma$ 

21. **[PP2013]** 
$$(p_0 \to p_1) \to \neg p_1 \vdash \neg (\neg (\neg p_1)) \lor \neg (p_0 \to p_1)$$

22. **[PP2014]** 
$$(\alpha \rightarrow \beta) \rightarrow \gamma \vdash \gamma \lor \neg \beta$$

23. **[PP2014]** 
$$\vdash (\alpha \leftrightarrow (\beta \rightarrow \gamma)) \rightarrow (\neg \alpha \land \beta \rightarrow \neg \gamma)$$

24. **[PP2015]** 
$$\vdash ((\alpha \rightarrow \gamma) \land (\beta \rightarrow \gamma)) \leftrightarrow (\alpha \lor \beta \rightarrow \gamma)$$

25. **[PP2015]** 
$$\vdash (\alpha \rightarrow \beta \rightarrow \gamma) \rightarrow \neg(\neg \gamma \lor \neg \beta) \rightarrow \gamma$$

26. **[PP2016]** 
$$\neg (\beta \lor \alpha) \vdash \alpha \land \beta$$

27. **[PP2016]** 
$$(\varphi \lor \psi) \leftrightarrow (\varphi \lor (\neg \varphi \land \psi))$$

28. **[PP2017]** 
$$\gamma \to \alpha$$
,  $\neg \to (\neg \beta \lor \alpha) \vdash \neg \alpha \to (\neg \beta \land \neg \gamma)$ 

29. **[PP2017]** 
$$p_2 \to \neg p_1, \quad p_3 \to p_1, \quad \neg p_1 \to (\neg p_2 \lor p_1) \vdash \neg p_1 \to (\neg p_2 \land \neg p_3)$$

30. **[PP2017**] 
$$p_3 \to p_1$$
,  $\neg p_1 \to (\neg p_2 \lor p_1) \vdash (p_2 \lor p_3) \to p_1$  [Riesgo de salud]

31. **[PP2018]** 
$$p \leftrightarrow \neg p \vdash \bot$$

32. **[PP2018]** 
$$p \lor q$$
,  $p \to r$ ,  $q \to s \vdash t \to \neg(\neg r \land \neg s)$ 

33. **[ED2008]** 
$$\neg p \land q \rightarrow r$$
,  $r \lor s \rightarrow \neg p$ ,  $s \vdash \neg q \lor r$ 

34. **[ED2010]** 
$$\alpha \vee \beta \leftrightarrow \beta \vdash \alpha \leftrightarrow \alpha \wedge \beta$$

35. **[ED2011]** 
$$\neg p \land \neg q \rightarrow r \vdash (p \lor q) \lor r$$

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36. **[EF2012]** 
$$(\neg \alpha \lor \beta) \lor \gamma \vdash \neg \gamma \land \alpha \rightarrow \neg \neg \beta$$

37. 
$$a \lor b$$
,  $c \lor \neg b \vdash a \lor c$ 

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# Soluciones propuestas al Práctico 4

A continuación se presenta un esbozo de las soluciones de los ejercicios, es decir, las mismas no están completas: falta agregar el inciso indicador de la cancelación de las hipótesis destacadas entre corchetes. Se alienta al lector a realizar tales agregados a efectos de profundizar la comprensión de dichas pruebas.

Asimismo, a efectos de no sobrecargar la notación se ha optado por emplear letras del alfabeto latino en lugar de aquellas del alfabeto griego.

- 1. Vistos en clase.
- 2. (i)

$$\begin{array}{c|c} & \underbrace{ \begin{bmatrix} \bot \end{bmatrix} & \begin{bmatrix} \neg \bot \end{bmatrix}}_{E \neg} E \neg \\ & \frac{\bot}{L} RAA \\ & \frac{\bot}{\neg \bot} I \neg & \begin{bmatrix} \neg (\neg \bot) \end{bmatrix}}_{\neg (\neg (\neg \bot))} E \neg \\ & \frac{\bot}{\neg (\neg (\neg \bot))} I \rightarrow \\ & \frac{\bot}{\neg \bot} I \neg & \frac{\bot}{\neg (\neg (\neg \bot))} I \rightarrow \\ \end{array}$$

Otra forma más sencilla:

$$\frac{[\bot]}{\neg\bot} I \neg \qquad \frac{[\neg\bot] \quad [\neg(\neg\bot)]}{\neg(\neg(\neg\bot))} I \neg}_{\neg\bot \leftrightarrow \neg(\neg(\neg\bot))} I \leftrightarrow$$

3. Vistos en clase.

#### Ejercicios integradores.

1.

$$\frac{[a] \quad a \to b}{b} E \to \frac{[(b \to c) \land \neg c]}{b \to c} E \land 1 \qquad \underbrace{[(b \to c) \land \neg c]}_{\neg c} E \land 2$$

$$\frac{\frac{\bot}{\neg a} I \neg}{(b \to c) \land \neg c \to \neg a} I \to$$

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3.

$$\frac{[a] \quad a \to b \lor c}{b \lor c} E \to \frac{[b] \quad [\neg b]}{\frac{1}{c} E \bot} E \neg \frac{[c] \quad [\neg b]}{\frac{c \land \neg b}{c} E \land_{1}} I \land \\
\frac{\frac{c}{a \to c} I \to}{\neg b \to a \to c} I \to$$

4.

$$\frac{(\neg b \to \neg a) \land a}{\underbrace{a}} E \land_2 \underbrace{ \begin{bmatrix} \neg b \end{bmatrix}} \underbrace{ \begin{array}{c} (\neg b \to \neg a) \land a \\ \neg b \to \neg a \end{array}} E \land_1 \\ \underline{\frac{\bot}{b}} RAA \\ \end{array}$$

5.

$$\underbrace{ \begin{array}{c} [b] \quad \overline{a} \quad [a \rightarrow \neg b] \\ \overline{-b} \quad E \neg \\ \hline \\ \underline{\frac{\bot}{\neg a} I \neg} \quad \overline{a} \rightarrow \neg b \\ \hline \\ \underline{\frac{\bot}{\neg a} I \neg} \quad \overline{b} \quad E \neg \\ \hline \\ \underline{\frac{\bot}{c} E \bot} \\ \underline{(a \rightarrow \neg b) \rightarrow c} I \rightarrow \\ \hline \\ b \rightarrow (a \rightarrow \neg b) \rightarrow c I \rightarrow \\ \end{array} } E \rightarrow$$

6.

$$\frac{[a \wedge \neg b]}{\frac{a}{a}} E \wedge_1 \quad [a \to b] E \to \frac{[a \wedge \neg b]}{\neg b} E \wedge_2$$

$$\frac{[a \wedge \neg b]}{\frac{a}{a}} E \wedge_1 \qquad \frac{\bot}{\neg (a \to b)} I \neg$$

$$\frac{a \wedge \neg (a \to b)}{\neg (a \wedge \neg b)} I \neg$$

$$\frac{-(a \wedge \neg (a \to b))}{\neg (a \wedge \neg (a \to b))} E \neg$$

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 $\frac{[b]}{a \vee b} I^{\bigvee_2} \quad [\neg(a \vee b)]$  $\frac{[a]}{a \vee b} \ I^{\bigvee_1} \quad [\neg(a \vee b)]$  $\frac{\perp}{\neg a} I_{\neg}$  $\neg (a \lor b) \leftrightarrow \neg a \land \neg b$  $[\neg a \land \neg b] \to E_{\wedge 2}$  $[\neg \underline{a} \land \neg b] E_{\wedge_2}$  $\Gamma$  $\neg (a \lor b)$  $p_{\parallel}$ [q] $[\neg a \land \neg b] \to E_{\land 1}$ a $E_{\perp}$ q[a]

 $\overline{a}$   $E^{\perp}$  $E \leftrightarrow_1$  $E \lor$  $[\neg a \land b] E \land_2 [a \leftrightarrow b]$  $[a \land \neg b] _{E \land_2}$  $\bar{q}_{\Gamma}$  $\neg(a \leftrightarrow b)$  $E \leftrightarrow_1$  $\underline{[a \land \neg b]} \ E \land_1 \quad [a \leftrightarrow b]$  $(a \land \neg b) \lor \neg a \land b$ 

 $a \lor b \to g \to E \to$ RAA**⊣**|∞  $\frac{g \wedge \neg g}{\hat{E} \wedge \hat{e}}$  $VI = \overline{[8]}$  $\infty$  $[\neg(\underline{g} \vee \neg g)]$  $\frac{1\delta J}{g\sqrt{\neg g}} I^{V_1} \left[\neg (g\vee \neg g)\right] E_{\neg}$ E $\overline{8} \leftarrow q$  $[q \vdash]$ 

 $\infty$ 

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$$\begin{array}{c|c} \overline{[\neg a]} & \overline{[\neg (\neg a)]} & E \neg & \overline{[(a \to b) \land (c \to b)]} & E \land 1 \\ \hline \frac{\bot}{a} & RAA & a \to b & E \to \\ \hline [\neg (\neg a) \lor c] & b & c \to b \\ \hline \hline \neg (\neg a) \lor c \to b & I \to \\ \hline \hline (a \to b) \land (c \to b) \to \neg (\neg a) \lor c \to b & I \to \\ \hline \hline (a \to b) \land (c \to b) \to \neg (\neg a) \lor c \to b & I \to \\ \hline \end{array} \right.$$

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12.

$$\underbrace{ \begin{bmatrix} a \end{bmatrix} \quad \frac{ \begin{bmatrix} \neg a \wedge (b \to g) \end{bmatrix}}{\neg a} E \wedge_1 }_{ \begin{matrix} \frac{1}{b} E \bot \end{matrix}} E \neg_ \\ \underbrace{ \begin{matrix} b \end{matrix}}_{ \begin{matrix} b \end{matrix}} E \bot \qquad \underbrace{ \begin{matrix} [b] \\ b \end{matrix}}_{ \begin{matrix} b \end{matrix}} E \lor \qquad \underbrace{ \begin{matrix} [\neg a \wedge (b \to g)] \\ b \to g \end{matrix}}_{ \begin{matrix} b \end{matrix}} E \wedge_2 \\ \underbrace{ \begin{matrix} \frac{g}{a \vee b \to g} I \to \end{matrix}}_{ \begin{matrix} \neg a \wedge (b \to g) \end{bmatrix}}_{ \begin{matrix} \neg a \wedge (b \to g) \end{bmatrix}}_{ \begin{matrix} a \vee b \to g \end{matrix}} E \wedge_2$$

13.

$$\frac{\frac{[c]}{d \to c} I \to}{\frac{(d \to c) \vee a}{I \vee_1} \qquad \neg ((d \to c) \vee a)} E \neg \frac{\frac{\bot}{\neg c} I \neg}{\neg c \vee \neg b} I \vee_1$$

15.

$$\frac{\frac{[a]}{a \vee b} I \vee_{1} \quad [a \vee b \to b \wedge \neg a]}{\frac{b \wedge \neg a}{\neg a} E \wedge_{2}} E \to \frac{\frac{\bot}{\neg a} I \neg}{(a \vee b \to b \wedge \neg a) \to \neg a} I \to$$

17.

$$\frac{\frac{[\neg p]}{q \to \neg p} I \to (q \to \neg p) \to p}{\frac{p}{\frac{\bot}{p} RAA} E \to [\neg p]} E \to \frac{\frac{\bot}{p} RAA}{\neg (p \to \neg q) \lor p} I \lor_{2}$$

$$\frac{[p] \quad \frac{[\neg p \land \neg (p \land q)]}{\neg p} E \neg}{\frac{\frac{1}{q} E \bot}{p \rightarrow q} I \rightarrow} E \land_{1}$$

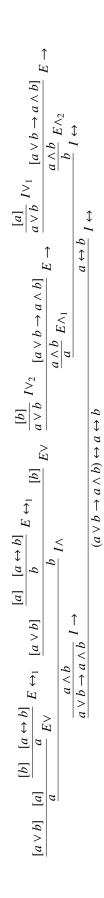
$$\frac{p}{\frac{p}{\neg (\neg p \land \neg (p \land q))} E \rightarrow} \frac{[(p \rightarrow q) \rightarrow p]}{p} E \rightarrow \frac{[\neg p \land \neg (p \land q)]}{\neg p} E \land_{1}$$

$$\frac{p}{\frac{1}{\neg (\neg p \land \neg (p \land q))} I \neg} E \land_{1}$$

$$\frac{(p \rightarrow q) \rightarrow p) \rightarrow \neg (\neg p \land \neg (p \land q))}{((p \rightarrow q) \rightarrow p) \rightarrow \neg (\neg p \land \neg (p \land q))} I \rightarrow$$

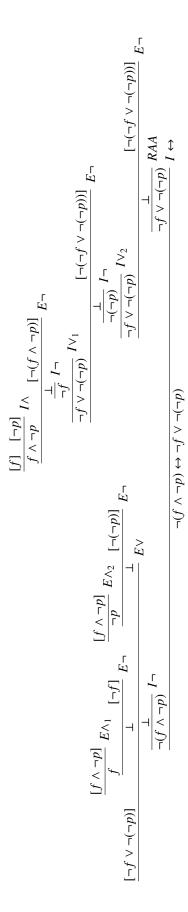
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19.

$$\frac{[f] \quad f \to p}{\frac{p}{p}} E \to [\neg p] E \neg$$

$$\frac{\frac{\bot}{\neg f} I \neg}{\frac{\neg (\neg p) \lor \neg f}{\neg (\neg p) \lor \neg f}} I \lor_2 \qquad [\neg (\neg (\neg p) \lor \neg f)]} E \neg$$

$$\frac{\frac{\bot}{\neg (\neg p)} I \neg}{\frac{\neg (\neg p) \lor \neg f}{\neg (\neg p) \lor \neg f}} I \lor_1 \qquad [\neg (\neg (\neg p) \lor \neg f)]} E \neg$$

$$\frac{\bot}{\neg (\neg p) \lor \neg f} RAA$$
(tautología dem. en [PP2014])

20.

22.

$$\frac{\left[\neg b\right]}{g \vee \neg b} I \vee_{2} \qquad \left[\neg (g \vee \neg b)\right] E \neg$$

$$\frac{\frac{1}{b} RAA}{a \to b} I \to \qquad (a \to b) \to g E \to$$

$$\frac{g}{g \vee \neg b} I \vee_{1} \qquad \left[\neg (g \vee \neg b)\right] E \neg$$

$$\frac{1}{g \vee \neg b} RAA$$

$$\frac{\left[g\right]}{b \to g} I \to \qquad \left[a \leftrightarrow b \to g\right] E \leftrightarrow_{1} \qquad \left[\neg a \wedge b\right] \neg a = E \wedge_{1}$$

$$\frac{\frac{[g]}{b \to g} I \to \frac{[a \leftrightarrow b \to g]}{a} E \leftrightarrow_1 \frac{[\neg a \land b]}{\neg a} E \land_1}{\frac{\frac{1}{\neg g} I \neg}{\neg a \land b \to \neg g} I \to} E \land_1$$

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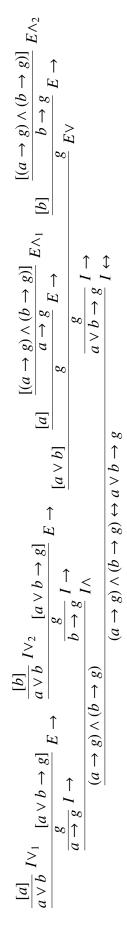
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 $[\neg(\neg(\neg(\neg(p_1)) \lor \neg(p_0 \to p_1))]$  $[\neg(\neg(\neg(\neg p_1)) \lor \neg(p_0 \to p_1))]$  $E \rightarrow [\neg(\neg p_1)]$  $\neg(\neg(\neg p_1)) \lor \neg(p_0 \to p_1)$  $\vec{a}(p_0 \to p_1)$  $[p_0 \to p_1] \quad (p_0 \to p_1) \to \neg p_1$ 

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25.

$$\frac{\frac{[\neg g]}{\neg g \vee \neg b} I \vee_1 \qquad [\neg (\neg g \vee \neg b)]}{\frac{\frac{1}{g} RAA}{\neg (\neg g \vee \neg b) \to g} I \to} E \neg$$

$$\frac{(a \to b \to g) \to \neg (\neg g \vee \neg b) \to g}{(a \to b \to g) \to \neg (\neg g \vee \neg b) \to g} I \to$$

26.

28.

$$\frac{[\neg a] \quad \neg a \rightarrow \neg b \vee a}{\neg b \vee a} E \rightarrow [\neg b] \quad \frac{[a] \quad [\neg a]}{\neg b} E \rightarrow [g] \quad g \rightarrow a \\
E \rightarrow [\neg a] E \rightarrow [\neg a]$$

$$\frac{[p] \quad p \to r}{r} E \to \frac{[\neg r \land \neg s]}{\neg r} E \land 1$$

$$\frac{[p] \qquad \frac{\bot}{\neg p} I \neg}{\frac{\bot}{q} E \bot} \qquad \qquad [q] E \lor \qquad q \to s E \to \frac{[\neg r \land \neg s]}{\neg s} E \land 2$$

$$\frac{g}{\frac{\bot}{\neg (\neg r \land \neg s)} I \neg} \qquad \qquad E \land 2$$

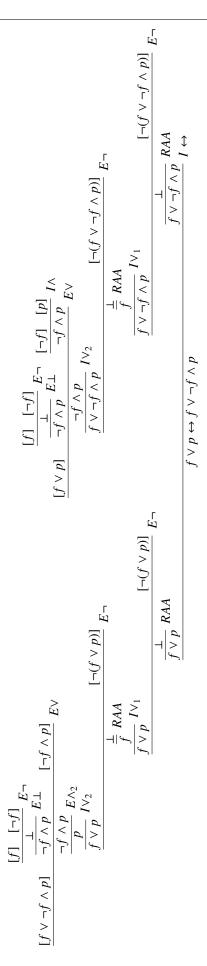
$$\frac{\bot}{\tau \to \neg (\neg r \land \neg s)} I \to$$

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$$\frac{[p_2]}{\neg p_1} \xrightarrow{P_2} \xrightarrow{P_1} E \to \frac{[p_1]}{\neg p_2} \lor p_1 = \frac{[p_1]}{\neg p_2} E \to \frac{[p_1]}{\neg p_2} E \to \frac{[p_1]}{\neg p_2} E \to \frac{[p_1]}{\neg p_2} E \to \frac{[p_2]}{\neg p_2} E \to \frac{[p_3]}{\neg p_2} E \to \frac{[p_3]}{\neg p_2} E \to \frac{[p_3]}{\neg p_3} E \to \frac{[p_$$

29.

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33.

$$\frac{\frac{s}{r \vee s} I \vee_{2} \qquad r \vee s \rightarrow \neg p}{\frac{\neg p}{p} E \rightarrow \qquad [q]} I \wedge \qquad \neg p \wedge q \rightarrow r} E \rightarrow \frac{\frac{r}{\neg q \vee r} I \vee_{2}}{\frac{\frac{\bot}{\neg q} I}{\neg q \vee r} I \vee_{1}} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I \vee_{1}}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee r} RAA} E \rightarrow \frac{\frac{\bot}{\neg q \vee r} I}{\frac{\bot}{\neg q \vee$$

34.

$$\underbrace{\frac{[a]}{a \lor b} I \lor_1}_{\substack{a \lor b \Leftrightarrow b \\ a \Leftrightarrow a \land b}} E \leftrightarrow_1 \underbrace{\frac{[a]}{a \lor b} I \lor_1}_{\substack{a \lor b \Leftrightarrow b \\ a \Leftrightarrow a \land b}} E \leftrightarrow_1$$

37.

$$\frac{a \lor b \quad [a]}{\frac{1}{a}} \frac{E \bot}{E \lor} E \neg$$

$$\frac{a \lor b \quad [a]}{\frac{a}{a}} \frac{E \bot}{E \lor}$$

$$\frac{a \lor c}{\frac{a \lor c}{a}} I \lor_{1} \qquad [\neg (a \lor c)] E \neg$$

$$\frac{b}{b} RAA \qquad [\neg b] E \neg$$

$$\frac{c \lor \neg b \quad [c]}{\frac{c}{a} \lor c} I \lor_{2} \qquad [\neg (a \lor c)] E \neg$$

$$\frac{c}{a \lor c} I \lor_{2} \qquad [\neg (a \lor c)] E \neg$$

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