

# Network-aware scheduler (NetSched) for containerized applications in an multilevel edge/fog/cloud distributed environment\*

Pablo Jaramillo, Agustín C. Caminero, Rocío Muñoz

National Distance University of Spain (UNED)

Madrid, Spain

pjaramill7@alumno.uned.es, accaminero@scc.uned.es, rmunoz@dia.uned.es

**Abstract**—The use of containers to deploy applications has suffered a huge growth in the last years. Containers allow the easy deployment of applications in such a way they can be isolated from the underlying hardware without the resource consumption and the performance degradation of traditional virtualization technologies (e.g. virtual machines). With this regard, containers can be used in a variety of scenarios, some of them harness distributed multilevel infrastructures (edge, fog and cloud). In these scenarios, the network plays an essential role in the proper execution of the applications, since it is the communication medium for all the agents in such distributed scenario and its poor performance may lead to poor quality of service of the entire system. This paper presents a new technique to perform scheduling of containerized applications on distributed multilevel infrastructures (edge, fog, and cloud). This technique considers the status of the interconnection network as a top-class resource to be considered when performing the scheduling tasks. A performance evaluation using a cluster of microcomputers Raspberry Pi has been conducted as a proof of concept of or work.

**Index Terms**—scheduling algorithm, network, virtualization, container, distributed system, edge, fog, cloud

## I. INTRODUCTION AND MOTIVATION

During the last years, distributed heterogeneous computing resources have arisen. *Cloud computing* can be considered as the application of business models existing in traditional utilities (such as water or power supply) to computing equipment, where users consume virtualized storage or compute resources offered by a provider based on a pay-as-you-go model. The resources providers use are large datacenters with plenty of computing, memory and storage capacity. Along with this, devices at the *edge* of the network can be part of Internet of Things (IoT) deployments. These can be sensors and devices that are part of Smart City infrastructure, lifestyle gadgets like wearables, and smart appliances or smart phones. These devices have been originally designed to sense and generate data streams, and they have some spare capacity (computing,

This work has been funded by UNED through the Independent Thinking-Jóvenes Investigadores 2019 project *efficient scheduling of containers* (FILE, 2019-PUNED-0007), by Ministry of Economy and Competitiveness (Ministerio de Economía y Competitividad, MINECO) through the Red Temática Española de Analítica de Aprendizaje, Spanish Network of Learning Analytics (SNOLA, RED2018-102725-T), and by the Region of Madrid through the project *e-Madrid-CM* (P2018/TCS-4307).

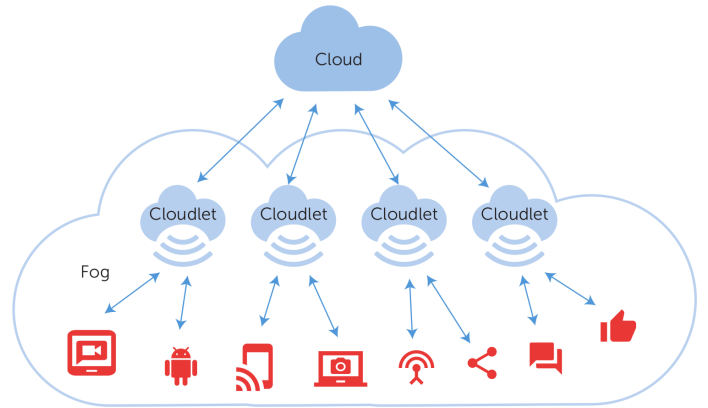


Fig. 1. Multilevel edge/fog/cloud scenario [5].

memory, and storage) which could be harnessed to execute applications [1].

The last term to be coined is *fog computing*, which refers to computing resources that are between the edge and cloud in the network hierarchy, with compute, storage and memory capacities that fall between these layers as well [2]. These edge and fog resources provide the opportunity for low latency processing of the generated data, closer to its source, and on the wide-area network [3]. As a result, there is critical need to understand how this diverse ecosystem of edge, fog and cloud resources can be effectively used by large-scale distributed applications [4]. This multilevel scenario can be seen in Figure 1 [5], where the term *cloudlets* refers to the access points devices extended with computing and storage services to create the fog system.

In such multilevel environment, the network plays an essential role since it is the communication medium between all the actors in the system and its poor performance affects the whole system. Thus, the network should be taken into account when making all kinds of decisions in order to provide Quality of Service (QoS), one of the main ones is scheduling. The term *scheduling* refers to the process of allocating computing resources to an application and mapping constituent components of that application onto those resources, in order to meet certain QoS and resource conservation goals [6].

Virtualization technologies allow the package and execution of applications in isolated environments, independently from the underlying hardware resources. The most traditional virtualization technology is *hardware virtualization*, which creates *virtual machines (VM)*, which are abstractions of hardware on top of which an operating system runs. They require extensive hardware resources to run since VMs run full operating systems. A more recent way of virtualization is *operating system-level virtualization*, which creates *containers* [7]. They use the host operating system and thus they do not require one for each container (as for the VMs), so the hardware requirements to run containers are lower than for VMs. Thanks to this, containers are suitable to run on low cost computers such as Raspberry Pi.

In order to deploy containers, this can be done in a cluster of computers. In this case, there is a need to manage such cluster, and this is done by the *container orchestrator*, one of the best known being Kubernetes [8]. Among the tasks the orchestrator has to perform, one of the most important is the scheduling of the containers into the nodes of the cluster.

Considering all of this, this work proposes an scheduling algorithm for multilevel (edge/fog/cloud) distributed environments in which the container orchestrator uses the status of the network, along with more information such as the memory or CPU status of the nodes, in order to choose the node of the cluster that will run a container – this is the main contribution of this work. Also, this paper presents proof-of-concept implementation of the algorithm on Kubernetes, and a performance evaluation of the algorithm, which has been conducted by deploying Kubernetes on a cluster of micro-computers Raspberry Pi [9]. Raspberry Pi has been chosen because of its low cost and its ability to run docker containers. Kubernetes has been chosen because it has been used for compute containers on Raspberry Pi-class edge devices [10], and used to deploy software on the fly on fog resources [11].

This paper is organized as follows. Section II presents the related work of this paper. Section III presents the scheduling algorithm which is the main contribution of this work. Section IV details the performance evaluation conducted to illustrate the usefulness of our work. Finally, Section V concludes the paper and suggests guidelines for future work.

## II. RELATED WORK

Studies have been conducted that address the effects of the deployment configuration in the performance and robustness of the application. In [12] a performance model of Kubernetes application is obtained with the application of benchmarks. It defends that for applications with more than eight containers, the bandwidth and final throughput is higher if every container is deployed in a separate pod – a *pod* is the basic deployment unit in the system. This suggest that deploying more pods with fewer containers in each pod, is better in terms of performance and bandwidth availability. In [13], from the same authors, the effects of the type of deployment explained above are studied and detailed depending on the nature of the application and its services. The paper makes a distinction between high CPU and

I/O demanding services and high bandwidth demand. At the same time, it makes a distinction between service container, that are executing full time, and job containers, which are created and destroyed with each execution. For CPU and I / O intensive applications it is argued that if there is more containers than nodes in the cluster, the best option is to group all the containers, however, it is important to divide the pods among all possible nodes, so it is recommended to divide the number of containers by the number of nodes available in the cluster and deploy a pod on each node. If there is a smaller number of containers than nodes, a container in each pod and a pod in each node. For the second type of application, intensive use of network resources, it is recommended the same policy applied previously in the case of type containers service, while for job-type containers it is always recommended to deploy a single container per pod. This is due to the fact that the network overload in a pod increases significantly when several containers are executed, since they all share the same IP address, unique in each pod.

These studies are very useful when making decisions about the deployment of an application on infrastructures with container orchestration. Despite all this, more precise adjustments may be required in the process of assignment of a node to each pod. If the application is based on a large number of job type containers that are created and destroyed proportionally with the load of work, it may be the case that the default Kubernetes planner is not enough. This scheduler takes into account resources such as CPU and memory, but it does not take into account the network congestion at the time of assignment or the overhead of the interfaces. This issue has been addressed in [14], where an scheduler that makes its decisions based, only, on the network bandwidth of the cluster resources. This work illustrates its usefulness with an evaluation using Kubernetes on a virtual environment. As opposed to it, our work presents a network-aware scheduler which makes its decisions by combining the main resources, compute, memory and network. Furthermore, this scheduler is evaluated by means of a real cluster made of Raspberry Pi.

## III. CONTAINER SCHEDULING ALGORITHM DEVELOPMENT

The Kubernetes scheduling algorithm, from now on, scheduler, has the only mission of scheduling the just created pods to a cluster node. It is a process executing with the rest of components in the master node. When a new pod is created, an event is created and the Kubernetes scheduler is triggered to start the assignation of the pod. It has three main execution phases: filtering, node ranking and node selection. During the filtering phase, the scheduler discards the nodes that are not suitable to run the pod due to restrictions and taints, that are applicable to each pod and node. During the second phase, the scheduler populates a list containing all the suitable nodes for the pod execution, applying a collection of priority functions to every node, trying to divide the pods among all the available nodes and giving more importance to the ones with less workload. During the last execution phase, the node

with higher priority is selected, and in case of existing more than one, it a random one is choosen. The scheduler is written in Golang, just like the rest of the Kubernetes project. The scheduler performance is a key part of the system, and in some cases it become the bottleneck of the cluster. For this work, the developed algorithm will only implement the last two phases of the original scheduling algorithm, avoiding the filtering phase as it is out of scope in the work.

The scheduler that has been developed makes use of six metrics, cpu usage, memory occupation, number of sent and received network packets, maximun available bandwidth and disk input and ouput. For the node ranking the scheduling algorithm will compare each node metrics, giving, for every metric, a pondered score to the node with best metrics. Each metric score has a different value, being CPU, memory and bandwidth the ones with highest weight. The priority computation function implements this logic, and is executed for every new pod in the cluster. The pseudocode for this logic can be seen in the Algorithm 1. The proposed algorithm has been called *Network-aware scheduler* (NetSched).

---

**Algorithm 1** Network-aware Scheduler (NetSched)

---

```

1: function CALCULATEPRIORITIES
2:   for int  $i = 0$ ;  $i < \text{numNodes}$ ;  $i++$  do
3:     if  $\text{nodeMetrics}[i].\text{cpu} < \text{bestCpu}$  then
4:        $\text{bestCpu} \leftarrow \text{nodeMetrics}[i].\text{cpu}$ 
5:        $\text{bestCpuN} \leftarrow \text{nodeMetrics}[i].\text{name}$ 
6:     if  $\text{nodeMetrics}[i].\text{mem} < \text{bestMem}$  then
7:        $\text{bestMem} \leftarrow \text{nodeMetrics}[i].\text{mem}$ 
8:        $\text{bestMemN} \leftarrow \text{nodeMetrics}[i].\text{name}$ 
9:     if  $\text{nodeMetrics}[i].\text{recPck} < \text{bestRecPck}$  then
10:       $\text{bestRecPck} \leftarrow \text{nodeMetrics}[i].\text{recPck}$ 
11:       $\text{bestRecPckN} \leftarrow \text{nodeMetrics}[i].\text{name}$ 
12:     if  $\text{nodeMetrics}[i].\text{sentPck} < \text{bestSentPck}$  then
13:       $\text{bestSentPck} \leftarrow \text{nodeMetrics}[i].\text{sentPck}$ 
14:       $\text{bestSentPckN} \leftarrow \text{nodeMetrics}[i].\text{name}$ 
15:     if  $\text{nodeMetrics}[i].\text{bwidth} > \text{bestBwidth}$  then
16:       $\text{bestBwidth} \leftarrow \text{nodeMetrics}[i].\text{bwidth}$ 
17:       $\text{bestBwidthN} \leftarrow \text{nodeMetrics}[i].\text{name}$ 
18:     if  $\text{nodeMetrics}[i].\text{diskIO} < \text{bestDiskIO}$  then
19:       $\text{bestDiskIO} \leftarrow \text{nodeMetrics}[i].\text{diskIO}$ 
20:       $\text{bestDiskION} \leftarrow \text{nodeMetrics}[i].\text{name}$ 
21:
22:    $\text{nodePriorities}[\text{bestCpuN}] += 3$ 
23:    $\text{nodePriorities}[\text{bestMemN}] += 2$ 
24:    $\text{nodePriorities}[\text{bestRecPckN}] += 1$ 
25:    $\text{nodePriorities}[\text{bestSentPckN}] += 1$ 
26:    $\text{nodePriorities}[\text{bestBwidthN}] += 2$ 
27:    $\text{nodePriorities}[\text{bestDiskION}] += 1$ 
28:
29:   return  $\text{nodePriorities}$ 

```

---

The algorithm makes use of the mentioned metrics, being nodeMetrics an array with the metrics for each node, cpu, mem, sentPck, recPck, bwidth and diskIO the metrics

taken into account for each node in the priority computation function. For every node, each metric is compared with the same metric for the rest of the nodes, and the best metric is choosen (in the case of cpu, memory, packets sent and received and disk usage metrics, less is better and for bandwidth metric more is better). The node with the best metric for each category is stored in bestCpuN , bestMemN, bestRecPckN, bestSentPckN, bestBwidthN and bestDiskION, so, at the end of the computation, scores are assigned to each of the nodes, returning an array with the nodes ordered by priority of scheduling.

The scheduler code is compiled and packaged in a container and deployed in a pod, and has to be allowed with a role based access account to schedule and allocate new pods. The docker image is available in DockerHub [15].

#### IV. PERFORMANCE EVALUATION

In order to illustrate the NetSched algorithm, a performance evaluation has been conducted in which it is compared with the original scheduling technique of Kubernetes. This is detailed in this section, where the experiments setup is explained in the first place, followed by an analysis of the results obtained.

##### A. Work Environment Preparation

For the testing environment, four Raspberry Pi 3 Model B and one AMD64 Intel Core i7 architecture computer have been used. They have been connected using a five ports Ethernet switch. Figure 2 shows the hardware and software configuration of the cluster. A Kubernetes cluster has been deployed over this hardware infrestructure configuring the AMD computer as the master node and the Raspberries as worker nodes, and tainting the master node to avoid the execution of pods on it. Prometheus [16] has also been installed in the cluster, so we are able to monitor the cluster resources usage, CPU and memory occupation, pods distribution and network metrics.

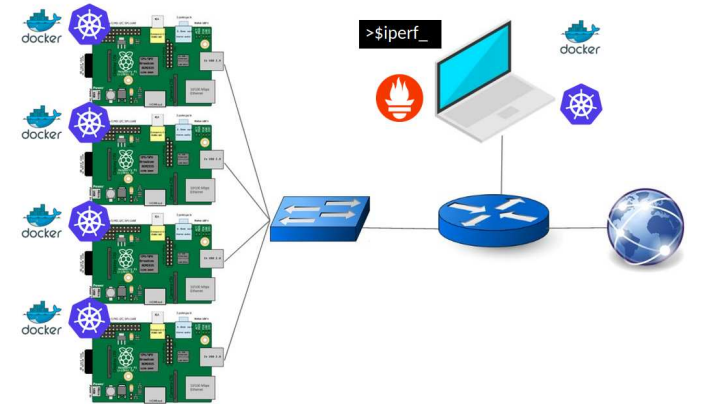


Fig. 2. Cluster hardware and software distribution

iPerf tool [17] which allows network monitoring between two or more nodes has also been installed and configured in the cluster, to collect information and metrics about the available

bandwidth between each node, packet loss and delay. With this two monitoring tools we are able to know the state of the cluster and feed the container scheduling algorithm which will select the best node for a certain pod taking all this data in account.

### B. Experiments and results

For the performance evaluation, the official benchmarking tools of the Kubernetes project have been used. This tool kit include different test suites in order to evaluate the performance and reliability of the cluster. For the performance of the Raspberry cluster and the scheduling algorithm the suite clusterloader2 [18] has been used. This suite creates a variable number of deployments, composed of a single pod, which contain a Docker container with a Kubernetes custom image inside.

Each container starts running and executes the command pause indefinitely. This suite is especially suitable for the performance evaluation of the project since it applies a considerable workload on the cluster and the developed scheduler will be in charge of assigning the pods to the most suitable nodes based on the parameters exported by Prometheus and iperf deployment.

The clusterloader2 suite, after creating all the deployments specified in its configuration, waits for all the containers to be in the Running state and then begins to eliminate the deployments. Throughout this process, this work presents results on the API responsiveness, scheduler latencies, and network congestion. Thanks to this data, the behavior of the cluster with the developed scheduler can be analyzed. Three different tests have been performed with each of the scheduling algorithms, varying in each one the total number of pods deployed in the cluster. For the first test, ninety pods have been configured to be planned by the algorithm, for the second one, one hundred ten and for the third one hundred and thirty.

1) *API responsiveness*: The API responsivity measures the latencies between the execution of an operation on the Kubernetes API server and the resolution of the response. It is an indicator to take into account during the performance evaluation since the API server works together between the master node and the slave nodes to evaluate and resolve the requests, so the response times offer a good approximation to know the performance of the cluster during the execution of the suite.

Figure 3 show the average response times in milliseconds for each of the operations available on etcd during the executions with both algorithms and 90 pods. The results with bigger deployments are similar so they are not shown on this short paper. They are divided between the operations of creation, deletion, obtaining, listing, listing with counting and updating. As can be seen, the most demanding operations for the cluster are those of creation and get, due to the high number of pods housed in the cluster at the same time. It can also be seen that these are quite high response times, this is due to the nature of the cluster, being composed of low-cost computers such as Raspberry Pi, its resources are limited. They do not

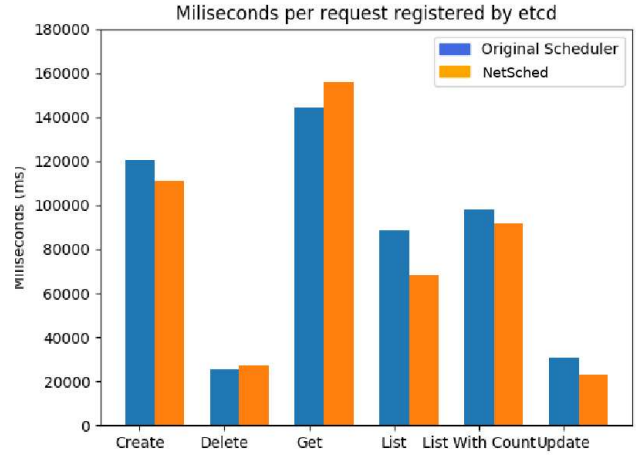


Fig. 3. API responsiveness, 90 pods

have high speed storage and the CPU has limitations when executing all the requests that the suite requires. This data allows to know the limitations of the cluster when analyzing possible applications of this in other areas. If the response times of both algorithms are compared, the similarity between the data obtained can be seen. This indicates that the developed scheduler has come into operation and has made an assignment of the nodes in a manner similar to that made by the original scheduler, without overloading any node, so the data collected is very similar to each other.

2) *Scheduler latencies*: The executed suite also offers the latency data between the state transitions of each pod, that is, the time that passes from being in one state to the next. For this evaluation, the latency between the Pending and Scheduled states has been taken into account.

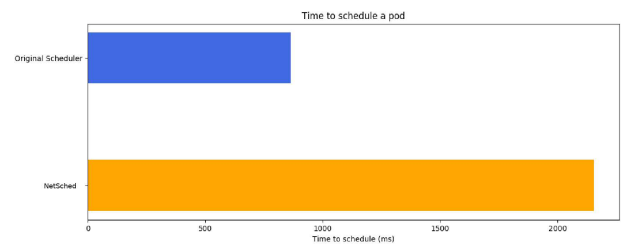


Fig. 4. Schedule time per pod

Figure 4 shows that the original Kubernetes scheduling algorithm is faster than NetSched in pod assignment than the developed algorithm. This measurement is the same in all three executions. This has been foreseen from the beginning and it is because NetSched is not optimized for use in production. It makes several requests to Prometheus node exporters and has to load and parse the JSON file with the information about the bandwidth at each execution, that is to say every time a pod is assigned. The original scheduler obtains the CPU and

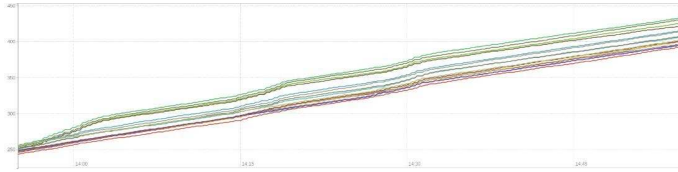


Fig. 5. CPU active time, Original Scheduler

memory metrics directly from the kubelet component, which is deployed on each node. This supposes a greater latency for the developed scheduler and although it has been developed in the Golang language, which offers greater performance than Python, it is an important point of improvement.

3) *Network Congestion Evaluation*: To test the efficiency of NetSched in the detection and selection of nodes in the event of a network congestion, a high demand for network resources has been simulated in three of the worker nodes, limiting their available bandwidth to a minimum by means of use of the netem tool [19]. Once configured, the clusterloader2 suite has been re-executed with both NetSched and the original algorithm, obtaining the following results. Figures 5 and 6 show the CPU usage time of each of the cores of the worker nodes.

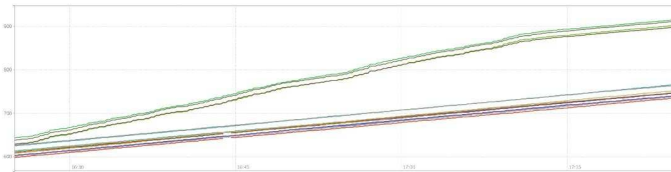


Fig. 6. CPU active time, NetSched

In Figure 5, the execution with the original Kubernetes algorithm shows a similar use of the CPU of each node, indicating that the pods have been assigned in the usual way, without taking into account that there is a network congestion. In Figure 6 the CPU usage timelines that have increased belong to the unmodified worker node. Figure 7 shows the assignation made to each worker node by both algorithms. It can be seen that NetSched, thanks to the constant monitoring of the network implemented with netperf, has been able to detect the low bandwidth available by the three modified worker nodes and has assigned the majority of the workload to the remaining worker node. In this way it is proved that NetSched can adapt to unfavorable situations of the network and act accordingly.

4) *Sending data through an unfavorable network*: The last test carried out to evaluate the favorable use cases for NetSched has consisted of sending data through a network in an unfavorable situation, similar to the previous case, but in this case the clusterloader2 suite has not been executed, but a benchmark developed specifically for this work. Three of the worker nodes have been modified by limiting their bandwidth to a minimum, leaving only one worker node in normal











	Scheduler	
	Original	NetSched
 	19	90
 	25	0
 	22	0
 	24	0
 	0	0

Fig. 7. Pods distribution by nodes in each execution

state. The benchmark developed has consisted of configuring a deployment of a variable number of containers with the image of iperf and programming each of them to perform a network test consisting of sending data. When creating the cluster deployment, the pods are assigned to the node that the running planning algorithm orders. Two tests have been carried out, the first with five containers and the second with ten containers, each sending 100 megabytes. Both tests have been performed on both algorithms. The results are presented below.











	Scheduler			
	Original 500MB	NetSched 500MB	Original 1GB	NetSched 1GB
 	1	5	1	10
 	2	0	4	0
 	1	0	2	0
 	1	0	3	0
 	0	0	0	0

Fig. 8. Pods distribution and features by nodes in each execution

Figure 8 shows the distribution of the pods made by each scheduling algorithm in both executions. As can be seen, when an unfavorable situation in the network is detected in any node, the pods are not assigned it, preventing low performance and network problems in the application.

Figures 9 and 10 show the total execution time of the tests performed, which is equal to the allocation time of the scheduler added to the total delivery time of all scheduled pods, ending when the last container ends Sending data The total test time with the work algorithm, including the planning time, is affected by the allocation latency of the developed scheduler.



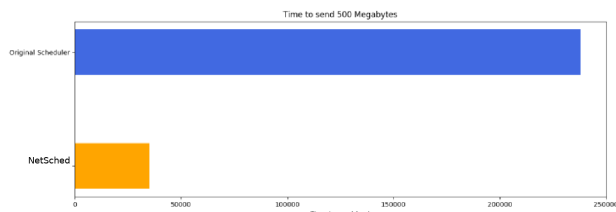


Fig. 9. Test execution duration. 500MB

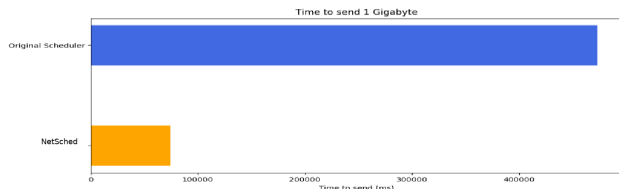


Fig. 10. Test execution duration. 1GB

Despite this latency, NetSched has a clear advantage over the original scheduler, which is not able to detect unfavorable network situations. In the case of such a scenario, the original scheduler assigns the pods without taking into account the state of the network, which has a significant influence on the total sending time.

## V. CONCLUSIONS AND FUTURE WORK

Distributed systems spanning infrastructures at different levels (edge/fog/cloud) have gained popularity in the recent years, and an appropriate scheduling technique is needed in order to provide QoS to the running applications. This paper proposes a network-aware scheduling algorithm that includes the network as a first class resource (along with CPU or memory) in order to make scheduling decisions. This algorithm, called NetSched, has been tested in several experiments, conducted after deploying a real scenario based on Raspberry Pi and Kubernetes. These experiments show that NetSched outperforms the original scheduling algorithm of Kubernetes when the network is overloaded. Thanks to NetSched, Kubernetes is able to detect such congestion and act accordingly, assigning pods to nodes with less network congestion. When there is no network congestion, the original scheduler outperforms NetSched since it has been tuned to be used in production, for instance, it avoids communication with other entities (e.g. iperf or Prometheus) to retrieve information to make the scheduling decisions. As for future work,

## REFERENCES

- [1] R. Ghosh, S. P. R. Komma, and Y. Simmhan, "Adaptive energy-aware scheduling of dynamic event analytics across edge and cloud resources," in *18th IEEE/ACM Intl. Symposium on Cluster, Cloud and Grid Computing (CCGRID)*, Washington DC, USA, 2018.
- [2] Y. Simmhan, "Big data and fog computing," in *Encyclopedia of Big Data Technologies*, S. Sakr and A. Y. Zomaya, Eds. Springer, 2019.

- [3] I. Stojmenovic, "Fog computing: A cloud to the ground support for smart things and machine-to-machine networks," *2014 Australasian Telecommunication Networks and Applications Conference (ATNAC)*, pp. 117–122, 2014.
- [4] P. Varshney and Y. Simmhan, "Characterizing application scheduling on edge, fog, and cloud computing resources," *Software: Practice and Experience (in press)*, vol. n/a, no. n/a.
- [5] L. F. Bittencourt, J. D. Montes, R. Buyya, O. F. Rana, and M. Parashar, "Mobility-aware application scheduling in fog computing," *IEEE Cloud Computing*, vol. 4, no. 2, pp. 26–35, 2017.
- [6] M. L. Pinedo, *Scheduling: Theory, Algorithms, and Systems*. Springer, 2008.
- [7] R. Buyya, C. Vecchiola, and S. T. Selvi, *Mastering Cloud Computing: Foundations and Applications Programming*, 1st ed. San Francisco, CA, USA: Morgan Kaufmann Publishers Inc., 2013.
- [8] D. Bernstein, "Containers and cloud: From lxc to docker to kubernetes," *IEEE Cloud Computing*, vol. 1, no. 3, pp. 81–84, 2014.
- [9] V. Vujovic and M. Maksimovic, "Raspberry pi as a sensor web node for home automation," *Computers & Electrical Engineering*, vol. 44, pp. 153–171, 2015.
- [10] P. Tsai, H. Hong, A. Cheng, and C. Hsu, "Distributed analytics in fog computing platforms using tensorflow and kubernetes," in *Proc. of the 19th Asia-Pacific Network Operations and Management Symposium (APNOMS)*, Seoul, Korea, 2017.
- [11] H. Hong, P. Tsai, and C. Hsu, "Dynamic module deployment in a fog computing platform," in *Proc. of 18th Asia-Pacific Network Operations and Management Symposium (APNOMS)*, Kanazawa, Japan, 2016.
- [12] V. Medel, J. BaÑsares, U. Arronategui, and O. Rana, "Modelling performance and resource management in kubernetes," *IEEE Computer/ACM International Conference on Utility and Cloud Computing*, p. 257–262, 2016.
- [13] —, "Characterising resource management performance in kubernetes," *IEEE Computer/ACM International Conference on Utility and Cloud Computing*, pp. 20–23, 2018.
- [14] J. Santos, T. Wauters, B. Volckaert, and F. De Turck, "Towards network-aware resource provisioning in kubernetes for fog computing applications," in *2019 IEEE Conference on Network Softwarization (NetSoft)*, Paris, France, 2019.
- [15] "Network-aware scheduler docker image available at docker hub," Web page at [https://hub.docker.com/r/pablojara/network\\_aware\\_scheduler](https://hub.docker.com/r/pablojara/network_aware_scheduler), Date of last access: January 9, 2020.
- [16] "Prometheus: Monitoring system and time series database," Web page at <https://prometheus.io/docs/introduction/overview/>, Date of last access: January 9, 2020.
- [17] "iperf - the tcp, udp and sctp network bandwidth tool," Web page at <https://prometheus.io/docs/introduction/overview/>, Date of last access: January 9, 2020.
- [18] "Kubernetes: Cluster loader performance test," Web page at <https://github.com/kubernetes/perf-tests/tree/master/clusterloader2/>, Date of last access: January 9, 2020.
- [19] "netem," Web page at <https://wiki.linuxfoundation.org/networking/netem>, Date of last access: January 9, 2020.