

Hashmap

- » Hashmap - Introduction
- «/» Highest Frequency Character
- «/» Get Common Elements - 1
- «/» Get Common Elements - 2
- «/» Longest Consecutive Sequence Of Elements
- «/» Write Hashmap

Heap/Priority Queue

- » Heaps - Introduction And Usage
- «/» K Largest Elements
- » Efficient Heap Constructor
- «/» Write Priority Queue Using Heap
- «/» Sort K-sorted Array
- » Heap - Comparable V/s Comparator
- «/» Merge K Sorted Lists
- «/» Median Priority Queue

#HashMap

Insert → put $\Rightarrow O(1)$

Update

Delete \rightarrow remove $\Rightarrow O(1)$

Read \rightarrow get $\Rightarrow O(1)$

If key exist then
return value
else return null

Display

size $\Rightarrow O(1)$

contains key \rightarrow find $\Rightarrow O(1)$

keyset \rightarrow keys

Key \rightarrow value

e.g. IPL teams \rightarrow trophies count
String \rightarrow int

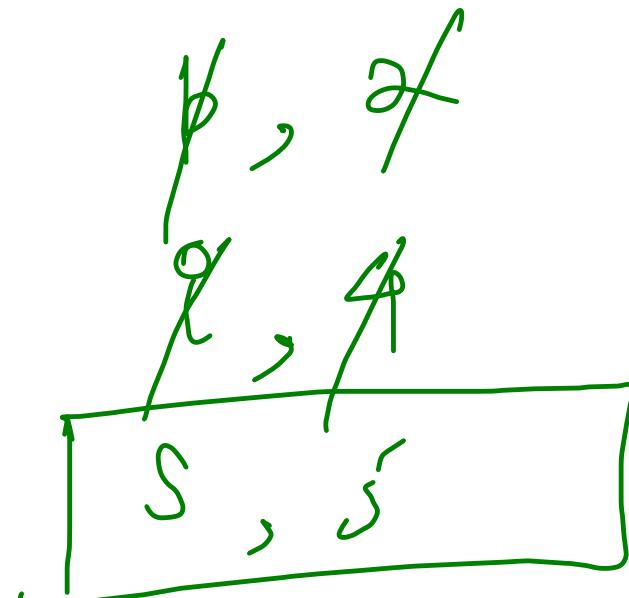
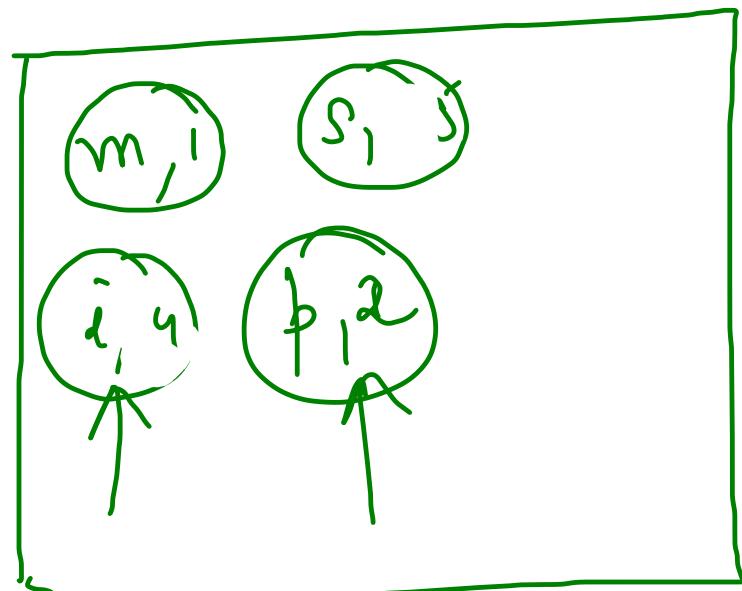
{ CSK \rightarrow 4
MI \rightarrow 5
SRH \rightarrow 2
DC \rightarrow 0

e.g. countries \rightarrow population
String \rightarrow integer

{ India \rightarrow 130
USA \rightarrow 90
China \rightarrow 200

Highest Frequency Character

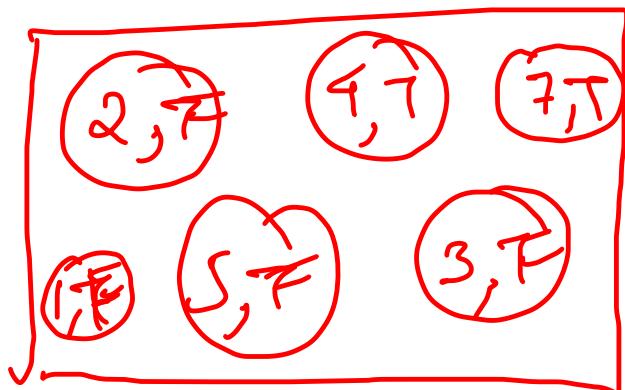
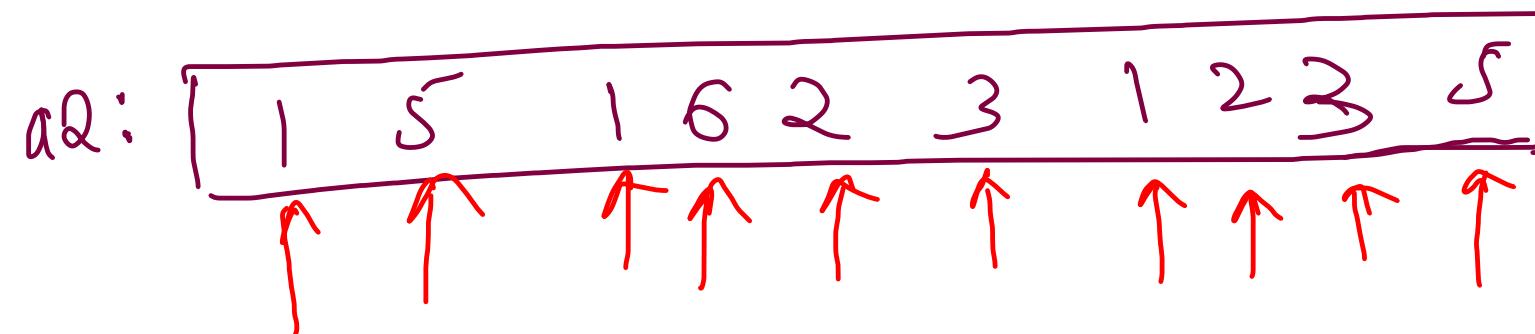
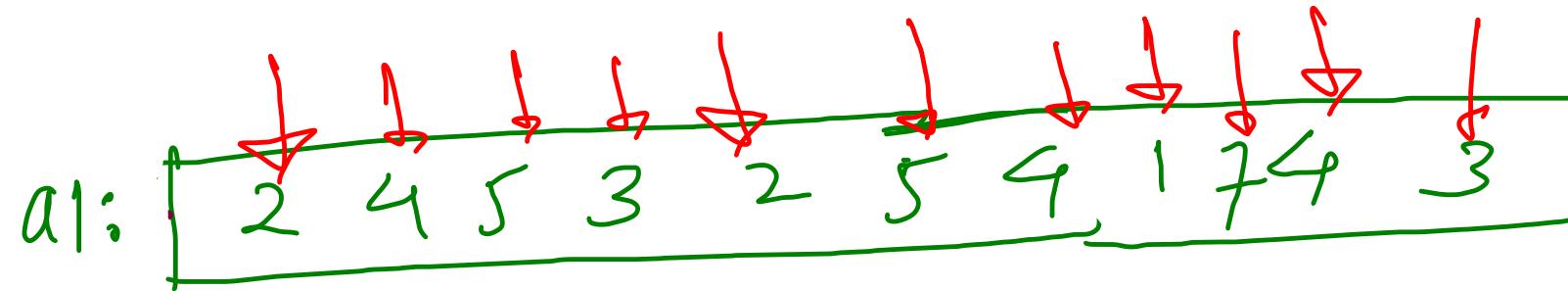
m i s s o s i p b o x f



```
HashMap<Character, Integer> freq = new HashMap<>();  
  
for(int i=0; i<str.length(); i++){  
    char ch = str.charAt(i);  
    if(freq.containsKey(ch)){  
        int oldFreq = freq.get(ch);  
        freq.put(ch, oldFreq + 1);  
    }  
    else {  
        freq.put(ch, 1);  
    }  
  
    char ch = str.charAt(0);  
    int maxFreq = freq.get(ch);  
  
    for(Character key: freq.keySet()){  
        int currFreq = freq.get(key);  
  
        if(currFreq > maxFreq){  
            ch = key;  
            maxFreq = currFreq;  
        }  
    }  
  
    System.out.println(ch);
```

The code is annotated with time complexities:
- The first loop is labeled $O(N)$.
- The nested loop is labeled $O(\text{Keys})$.
- The final \maxFreq assignment is labeled $= O(256)$.

Get Common Elements ↗



$O(n)$

$\langle \text{Integer}, \text{Boolean} \rangle$

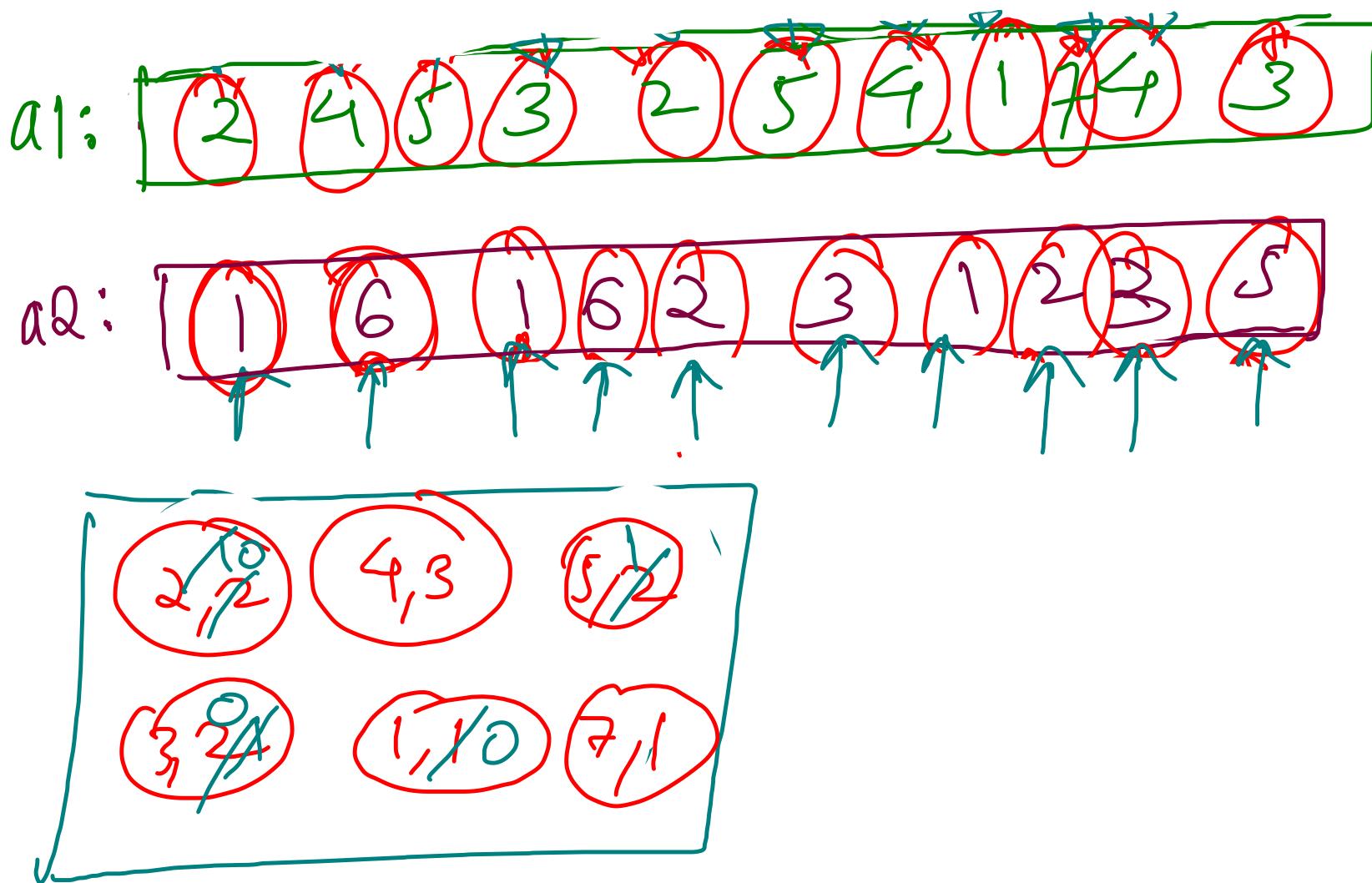
```
HashMap<Integer, Boolean> hm = new HashMap<>();
for(int i=0; i<n1; i++)
    hm.put(arr1[i], true);

for(int i=0; i<n2; i++){
    if(hm.containsKey(arr2[i]) && hm.get(arr2[i])){
        System.out.println(arr2[i]);
        hm.put(arr2[i], false);
    }
}
```

$\mathcal{O}(n_1 + n_2)$

1, 5, 2, 3

Get Common Elements - 2

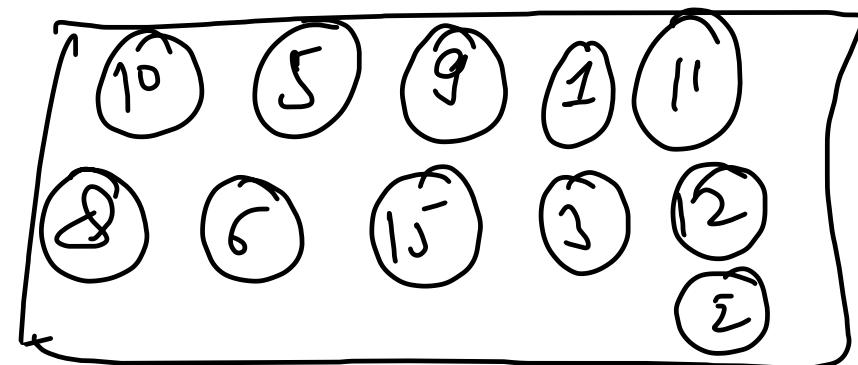
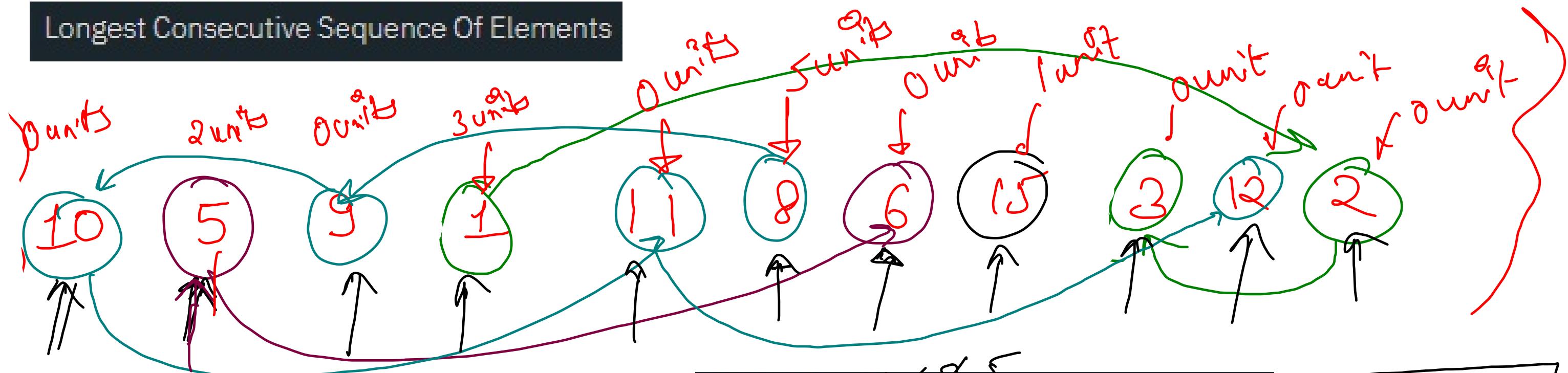


```
HashMap<Integer, Integer> hm = new HashMap<>();
for(int i=0; i<n1; i++){
    if(hm.containsKey(arr1[i])){
        hm.put(arr1[i], hm.get(arr1[i]) + 1);
    } else {
        hm.put(arr1[i], 1);
    }
}

for(int i=0; i<n2; i++){
    if(hm.containsKey(arr2[i]) && hm.get(arr2[i]) > 0){
        System.out.println(arr2[i]);
        hm.put(arr2[i], hm.get(arr2[i]) - 1);
    }
}
```

1, 2, 3, 2, 3, 5

Longest Consecutive Sequence Of Elements



```

int maxChain = 0;
int startingPt = 0;

for(Integer key: hm.keySet()){
    if(hm.containsKey(key - 1) == false){
        // chain starting pt
        int length = 1;
        while(hm.containsKey(key + length) == true){
            length++;
        }
        if(length > maxChain){
            maxChain = length;
            startingPt = key;
        }
    }
}

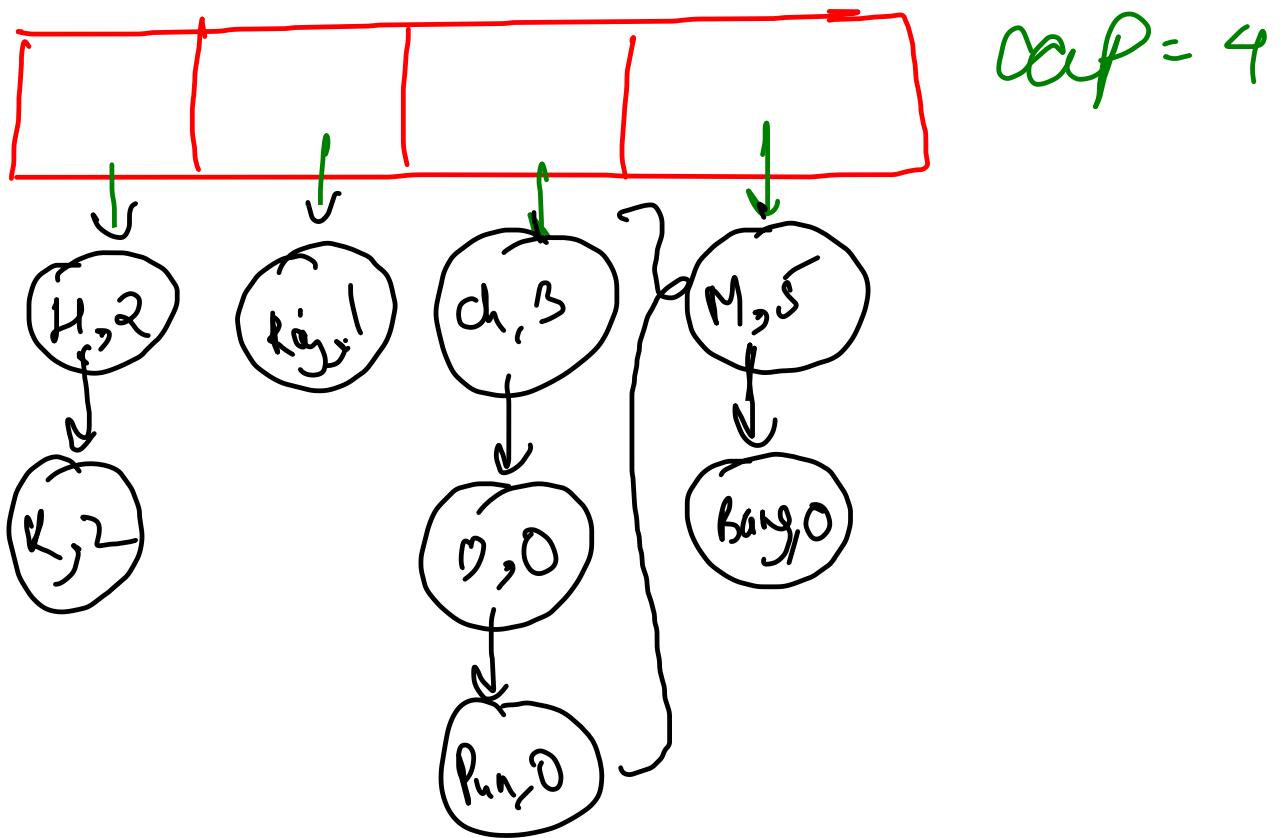
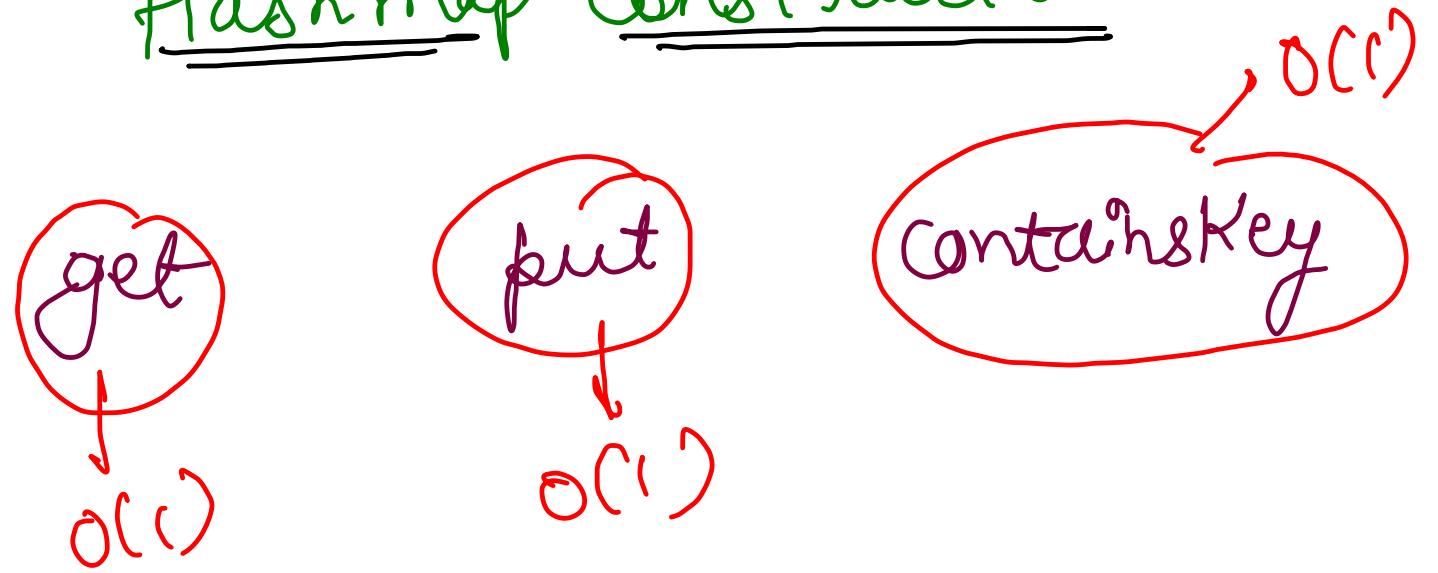
```

$2 \rightarrow 3$
 $8 \rightarrow 9 \rightarrow 10 \rightarrow 11 \rightarrow 12$

$5 \rightarrow 6$
 15

$O(N)$ Time

HashMap Construction



keySet

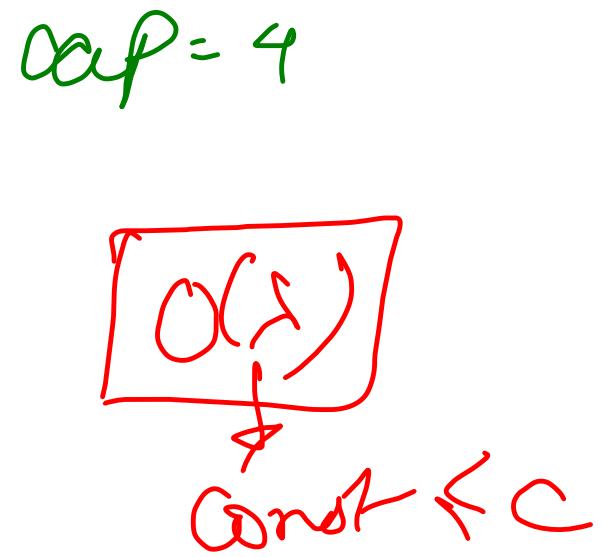
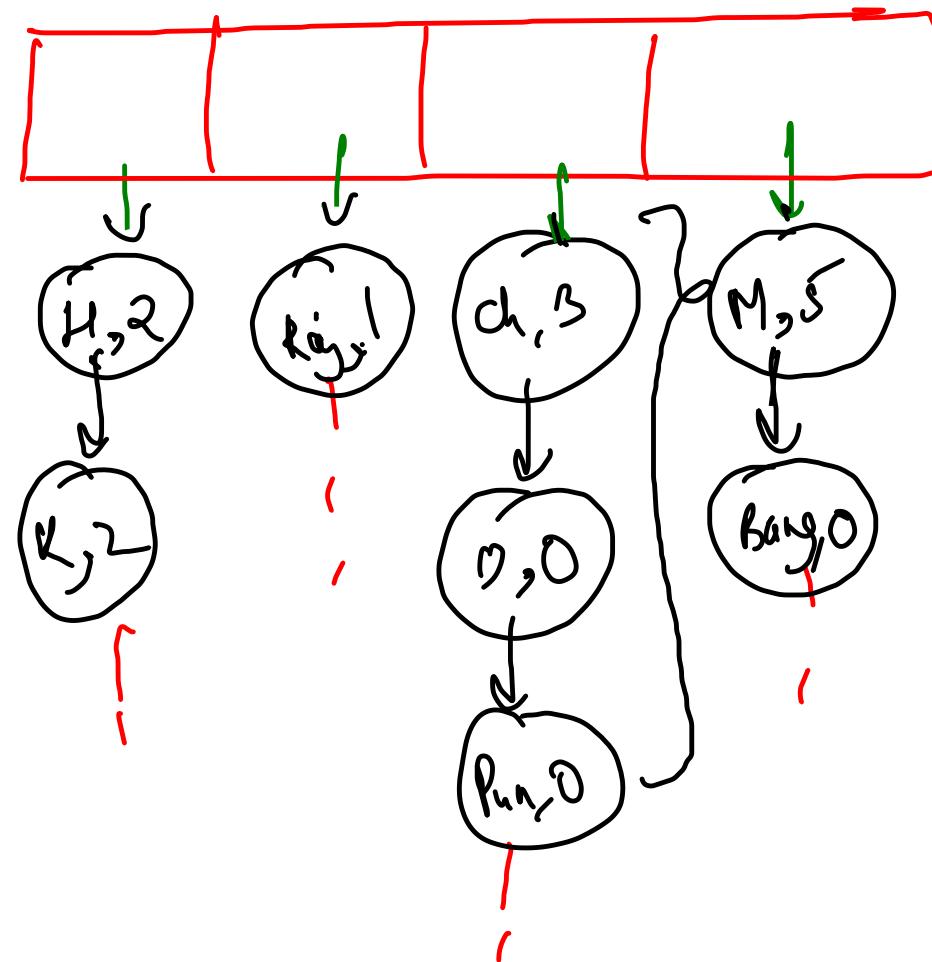
size

display

hashfunction

→ Two keys can have same hashfn

→ A given key will never have two hashfn values.



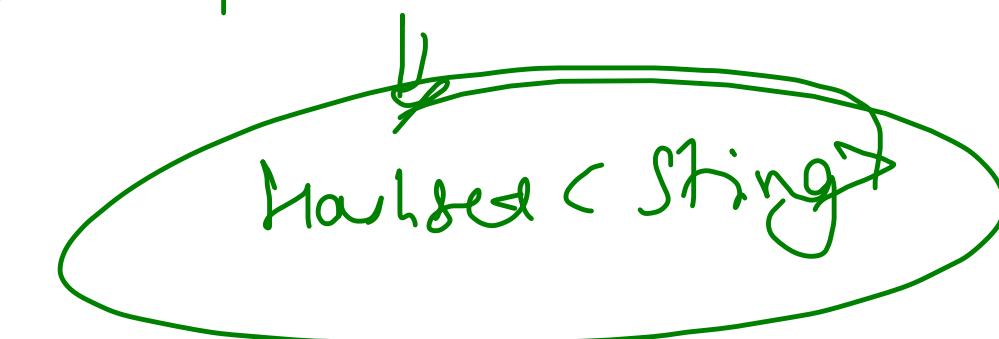
loading factor $\frac{1}{4}$

= $\frac{\text{no of nodes}}{\text{in one bucket}}$

= $\frac{\text{total nodes}}{\text{no of buckets}}$

$$= \frac{8}{4} = 2$$

HashMap < String, Boolean >



```

private LinkedList<HMNode>[] buckets;
private int noOfNodes;
private int noOfBuckets;

public HashMap(){
    noOfBuckets = 4;
    noOfNodes = 0;
    buckets = new LinkedList[noOfBuckets];

    for(int i=0; i<noOfBuckets; i++){
        buckets[i] = new LinkedList<>();
    }
}

```

```

public void put(K key, V value) throws Exception {
    // O(1)
    int bucketId = getBucketId(key);
    HMNode data = getData(bucketId, key);

    if(data == null){
        // Insertion
        HMNode node = new HMNode(key, value);
        buckets[bucketId].addLast(node);
        noOfNodes++;
    } else {
        // Updation
        data.value = value;
    }
}

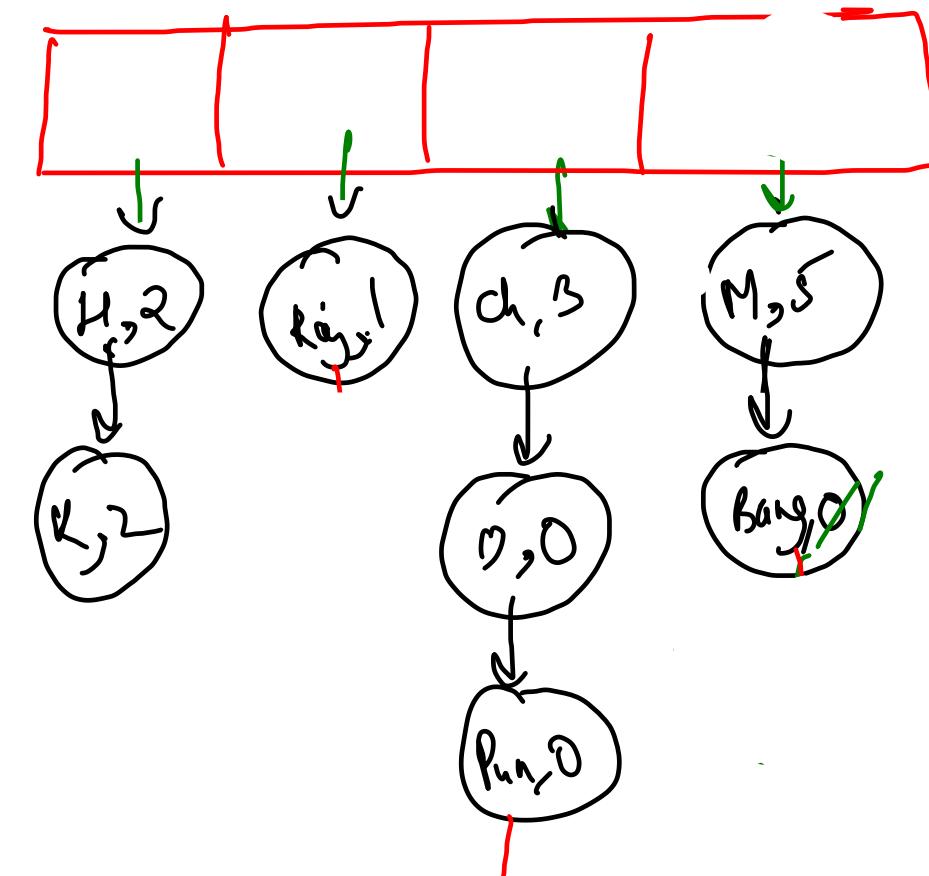
```

```

public int getBucketId(K key) throws Exception{
    // O(1)
    int hashCode = key.hashCode();
    int bucketId = (Math.abs(hashCode)) % noOfBuckets;
    return bucketId;
}

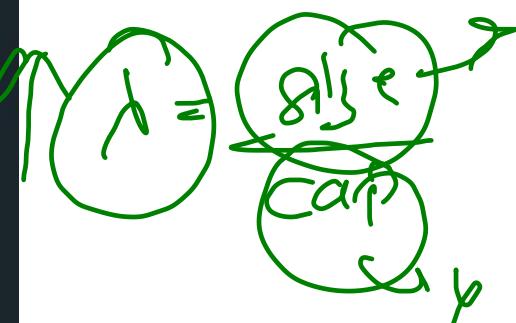
private HMNode getData(int bucketId, K key) throws Exception{
    for(HMNode node: buckets[bucketId]){
        if(node.key.equals(key) == true){
            return node; // data already exist
        }
    }
    return null; // data not exist
}

```



$\text{Cap} = 9$

$\text{Size} = 18$

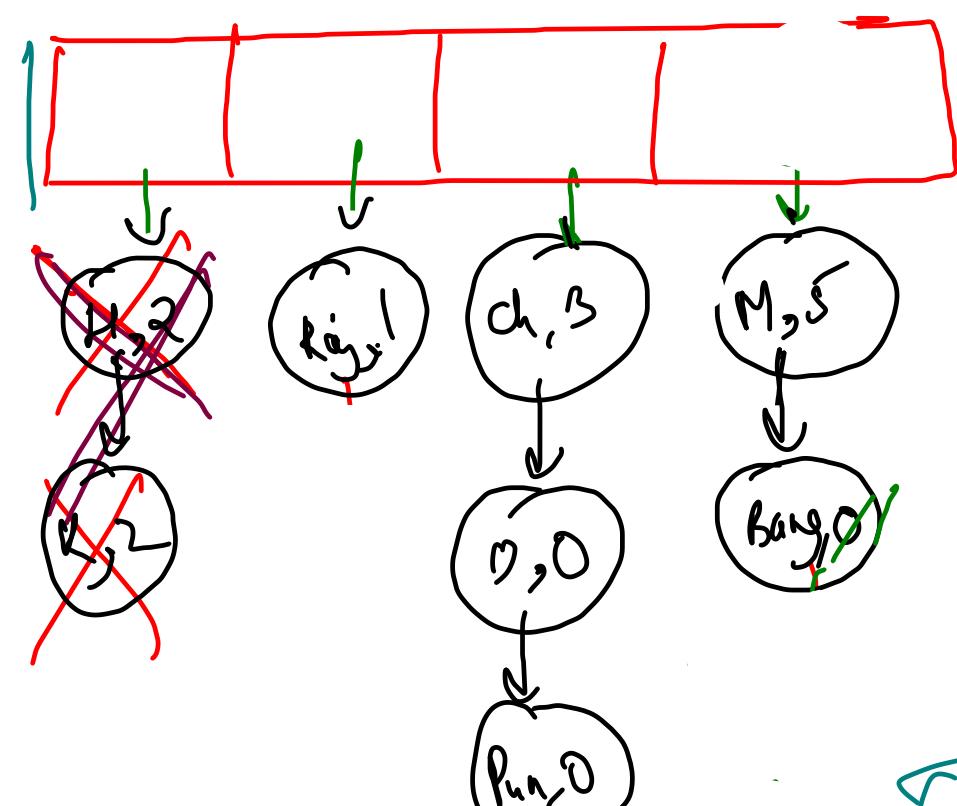


$O(2)$

$O(C)$

Ahm

$1/4 = O(1,2,3)$

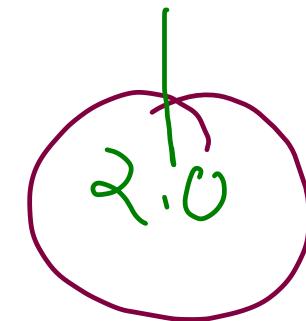


size = 8
cap = 9

$O(N)$

Rehashing

$$\lambda = \frac{8}{9} \approx 0.88$$

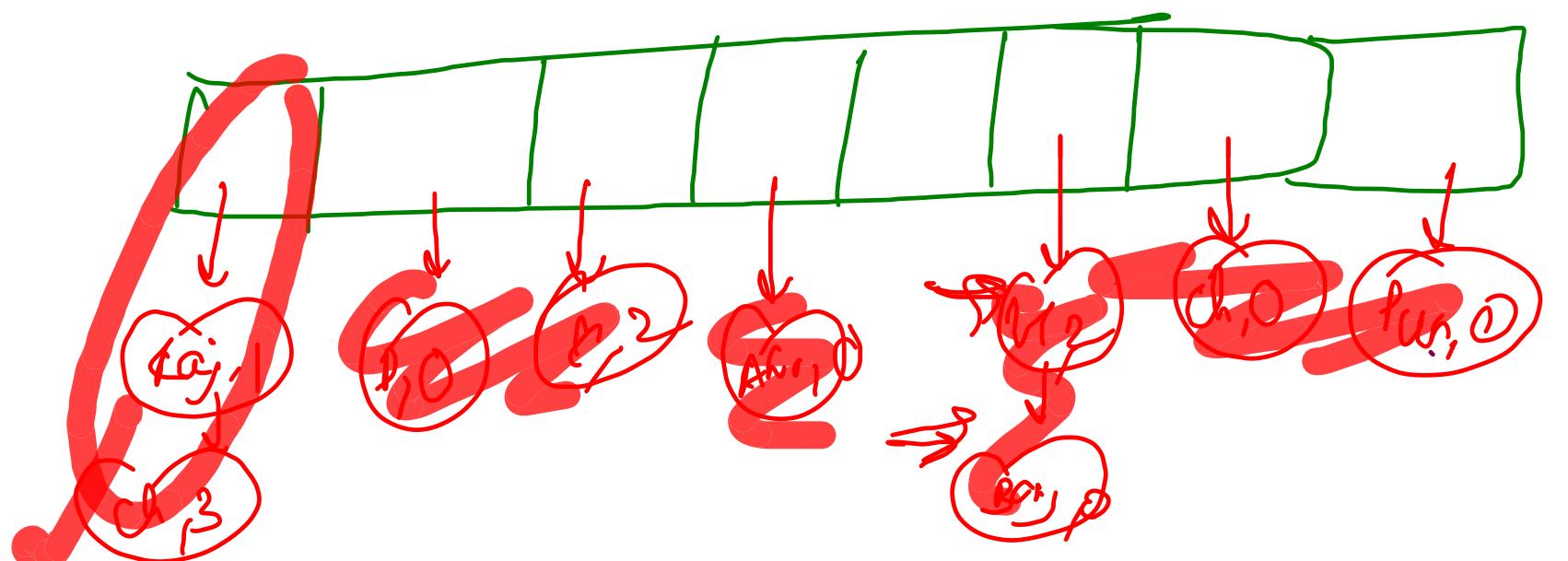


~~$O(c)$~~

hashcode = Hyp
 $(h1) \mod 4$



$$\lambda = \frac{9}{8} = 1.125$$



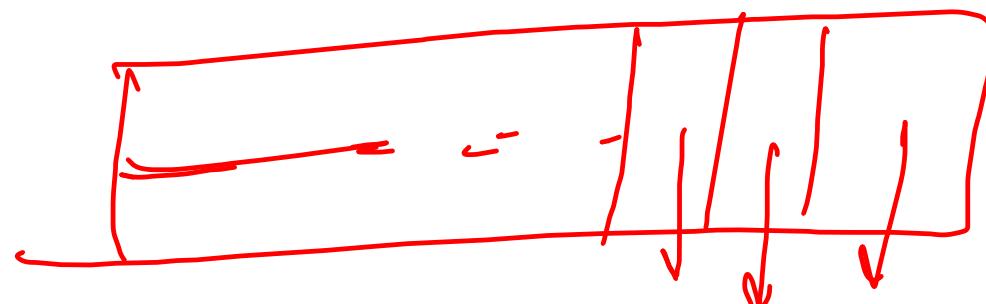
Time Comp

- Normal
- Rehashing

→ worst case

loading factors

Collisions $O(w)$



9^{..};
931.
92/1

C - 999

~~Collisions~~

10/11

93

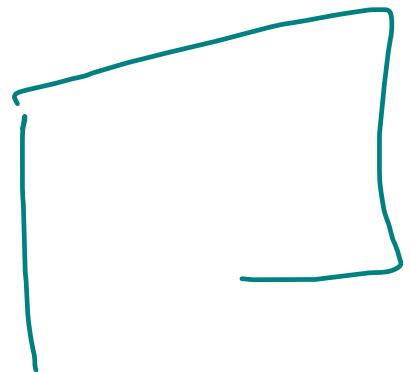
9117(889)

1 2 3 5 6 0 7

A simple red line drawing of a face. The face has a large, irregular oval shape. It features two large, circular eyes with thick outlines and simple curved lines for pupils. A small, open mouth is positioned at the bottom center. The drawing is done with a single continuous red line.

A simple line drawing of a person's head and shoulders, facing right. The head is large and rounded, with a small tuft of hair at the top. The eyes are represented by two curved lines, and the mouth is a simple U-shape. The body is a simple line extending from the right side of the head.

← OCD



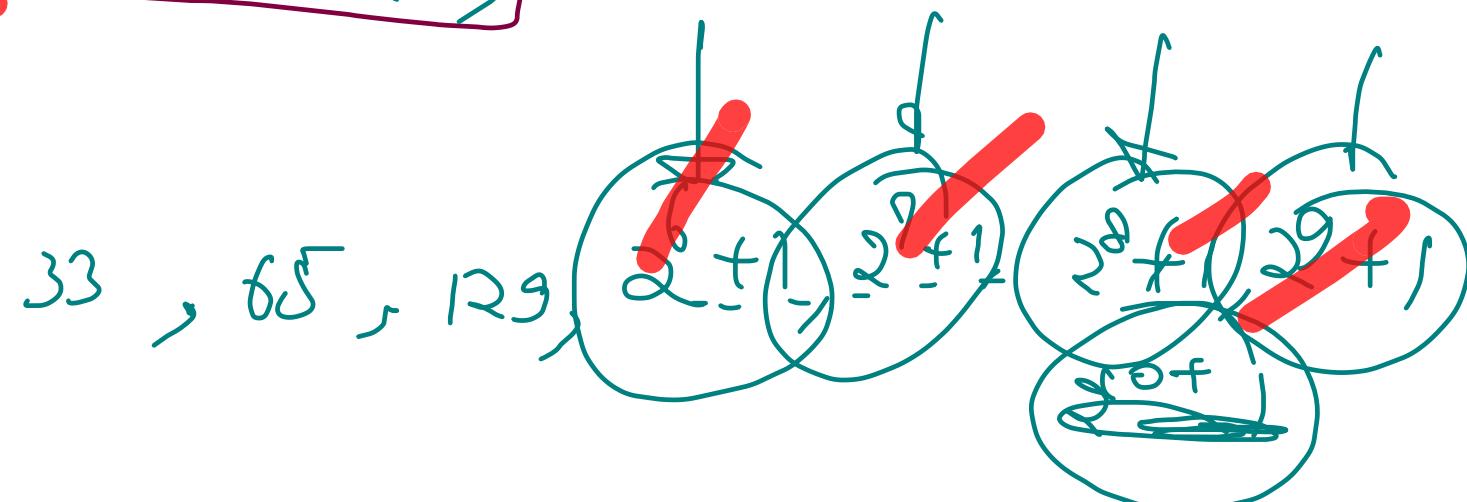
~~1st $\rightarrow O(1)$~~
~~2nd $\rightarrow O(n)O(1)$~~
~~3rd $\rightarrow O(1)$~~
~~4th $\rightarrow O(1)$~~
~~5th $\rightarrow O(8)$~~
~~6th $\rightarrow O(1)$~~
~~7th $\rightarrow O(1)$~~

~~8 $\rightarrow O(1)$~~
~~9 $\rightarrow O(16)$~~
~~10 $\rightarrow O(1)$~~
~~11 $\rightarrow O(32)$~~

Amortized

~~1 + 1 + 1 + 1 + 8 + 1 + ...~~
~~+ 16 + ...~~
~~+ 32 + ...~~

$$2^{10} + 1 = O(n)$$



✓ String \Rightarrow

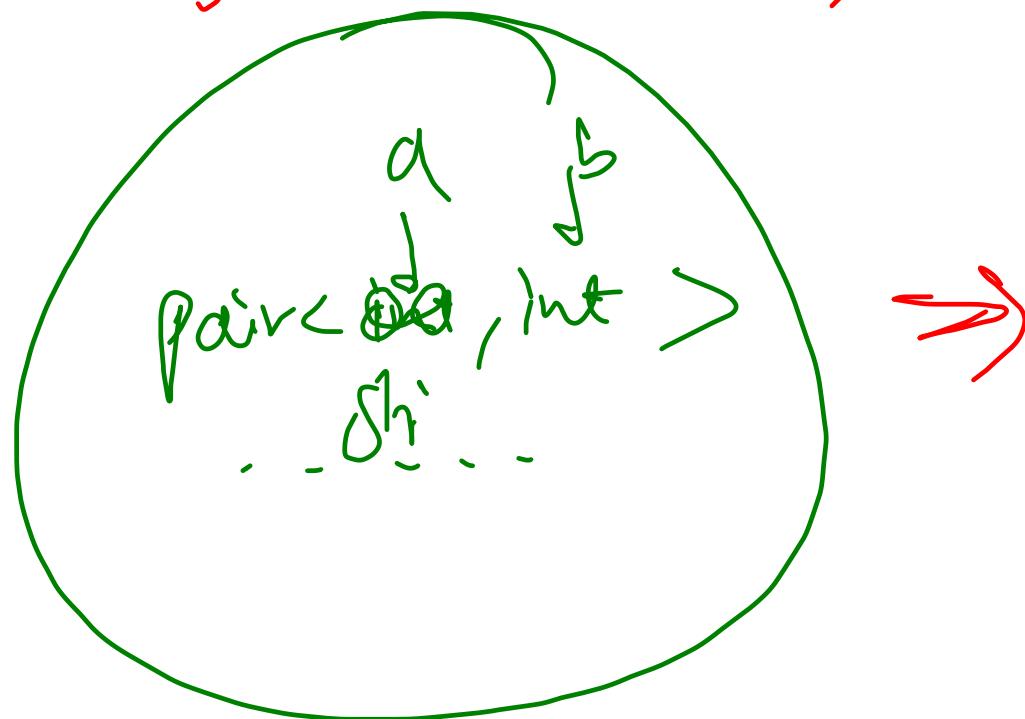
bepcoder
0 1 2 3 4 5 6 7

$O \times P$
 $+ 1 \times e$

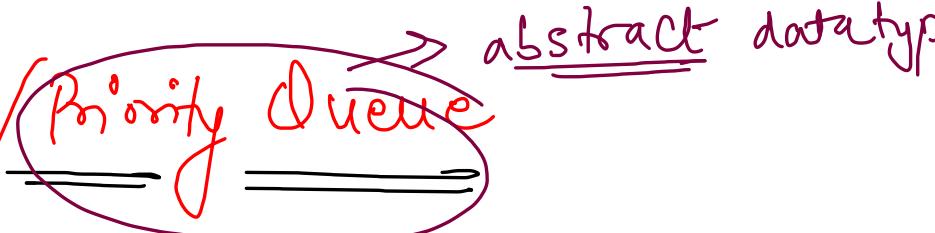
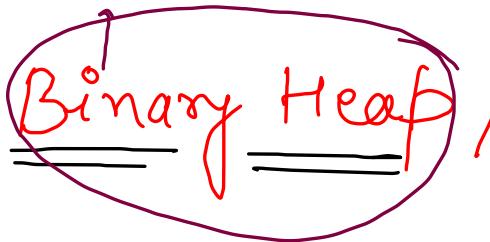
$+ 2 \times P + 3 \times C$
 $+ \dots) / f$

✓ int $\Rightarrow R_1, R_2$

$\rightarrow (\text{key } \times P_1) \times P_2$



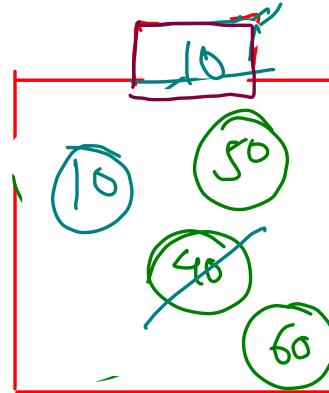
Data Structure



Priority Queue

- add → $O(\log n)$
- removal → $O(\log n)$
- peek → $O(1)$

f+ Min Heap



- pq.remove()
- pq.remove()
- pq.remove()
- pq.remove()

$$n! = n^n$$

Student
→ roll no
→ marks
→ height

- size { $O(1)$ }
- display { $O(n)$ }

{30, 50, 40, 20, 60}

(10)

Inserting n elements

$$\begin{aligned} & \log 1 + \log 2 + \log 3 \\ & + \dots \log n \\ & - \log (1 \times 2 \times 3 \times \dots \times n) \\ & = \log(n!) \\ & = \log(n^n) = n \log n \end{aligned}$$

for inserting
n elements
one by one

Deleting n elements

$$\begin{aligned} & \log n + \log(n-1) + \log(n-2) \\ & + \dots \log 1 \quad \text{for removing } n \text{ elements} \\ & = n \log n \end{aligned}$$

Binary Heap :- Array

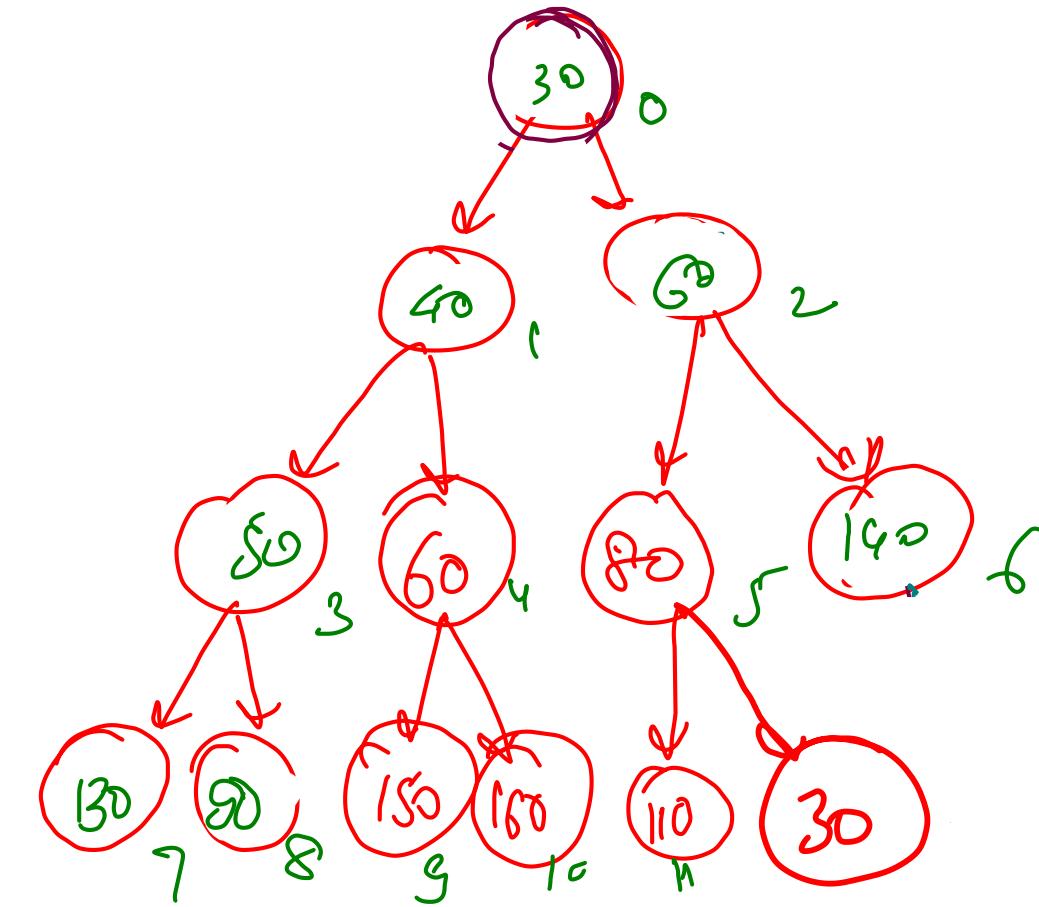
(\Rightarrow complete Binary Tree) \Rightarrow

\checkmark {Heap Order Property}

→ left child :- $2^i + 1$

→ right child :- $2^i + 2$

→ parent :- $(i-1)/2$



$$N \log(N) = \underline{N} \underline{\log N}$$

heapSort

5

10

20

30

40

60

Quick Sort :- ~~Stack~~

Merge Sort :- extra space

Heap Sort :- O(1) extra space

```

public void upheapify(int idx){
    int parIdx = (idx - 1) / 2;

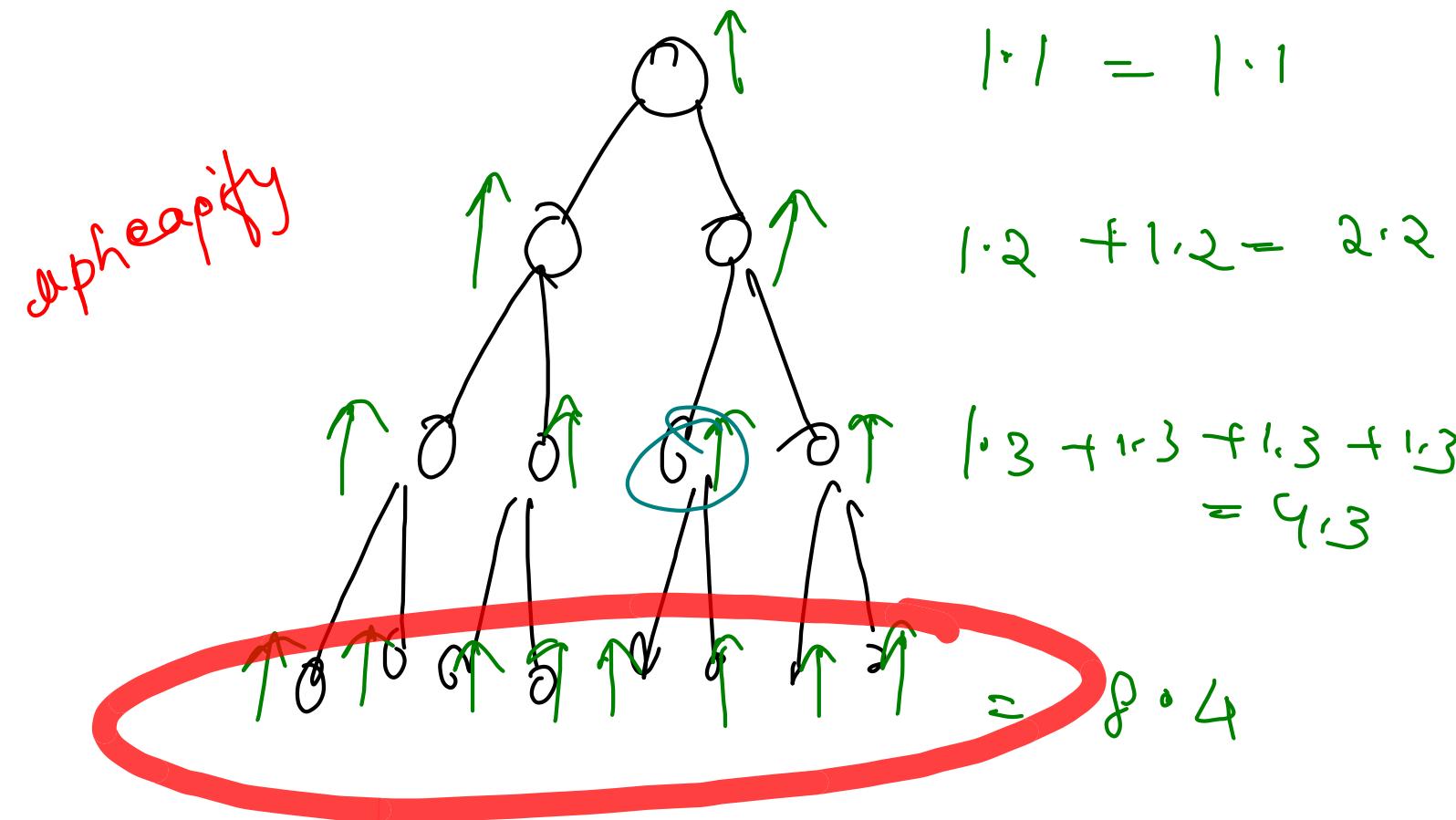
    if(isSmaller(idx, parIdx)){
        swap(idx, parIdx);
        upheapify(parIdx);
    }
}

```

```

public void add(int val) {
    // O(log n)
    data.add(val);
    upheapify(size() - 1);
}

```



insertion of
N nodes

$$\left\{ 2^0 \cdot 1 + 2^1 \cdot 2 + 2^2 \cdot 3 + 2^3 \cdot 4 + 2^4 \cdot 5 + 2^5 \cdot 6 \dots \right\} = N \log N$$

3yoda
nodes,
3yada
room
↓
 $O(n \log n)$
3yada
T.C.

```

public void downheapify(int idx){
    int min = idx;
    int left = 2 * idx + 1;
    int right = 2 * idx + 2;

    if(left < size() && isSmaller(left, min)){
        min = left;
    }

    if(right < size() && isSmaller(right, min)){
        min = right;
    }

    if(min != idx){
        swap(idx, min);
        downheapify(min);
    }
}

```

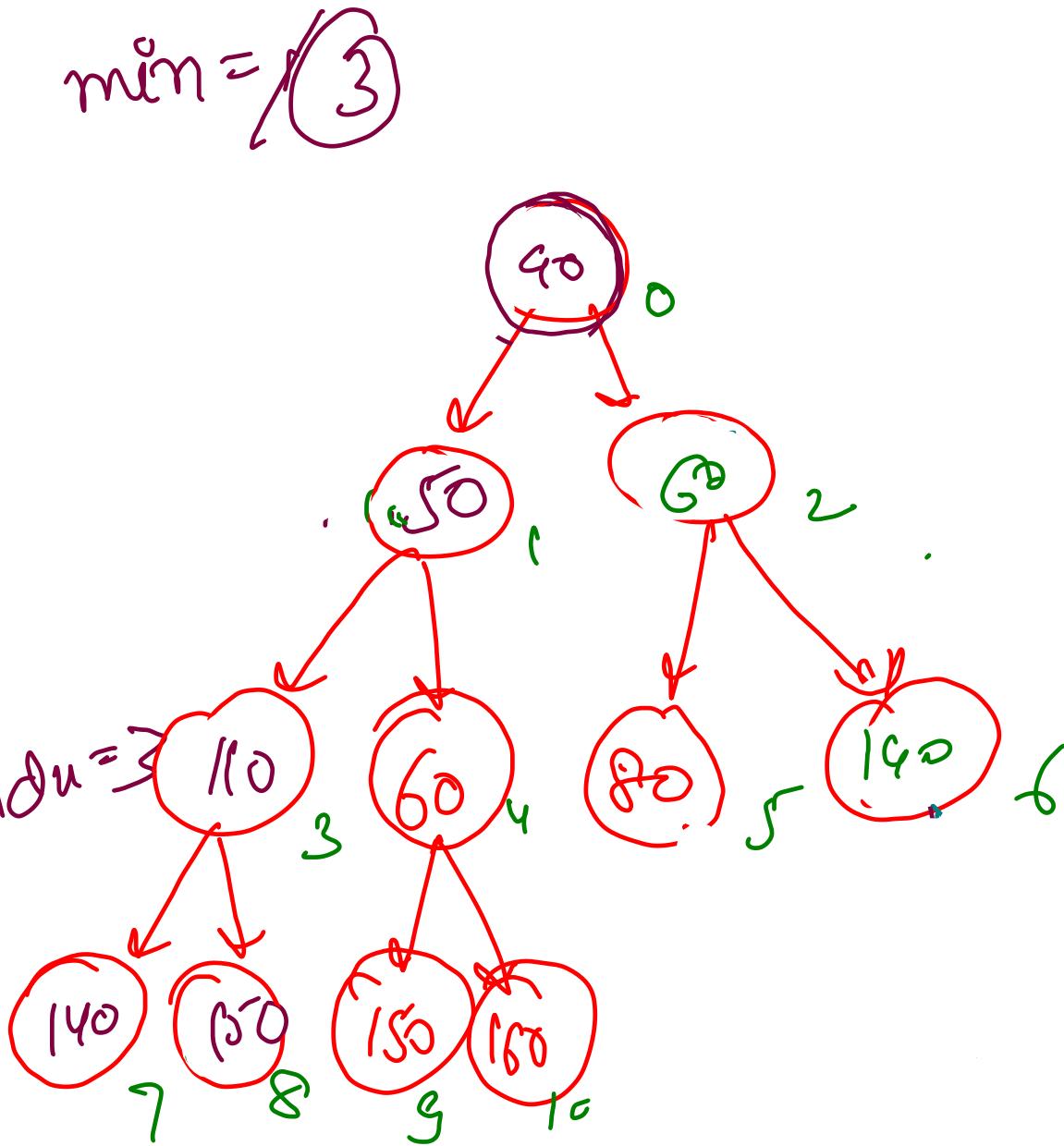
```

public int remove() {
    // O(log n)
    if(size() == 0){
        System.out.println("Underflow");
        return -1;
    }

    int val = peek();
    swap(0, size() - 1);
    data.remove(size() - 1);

    downheapify(0);
    return val;
}

```



opti nodes, zyada kaam

0. 2⁰ + 4

+ 2¹ + 3

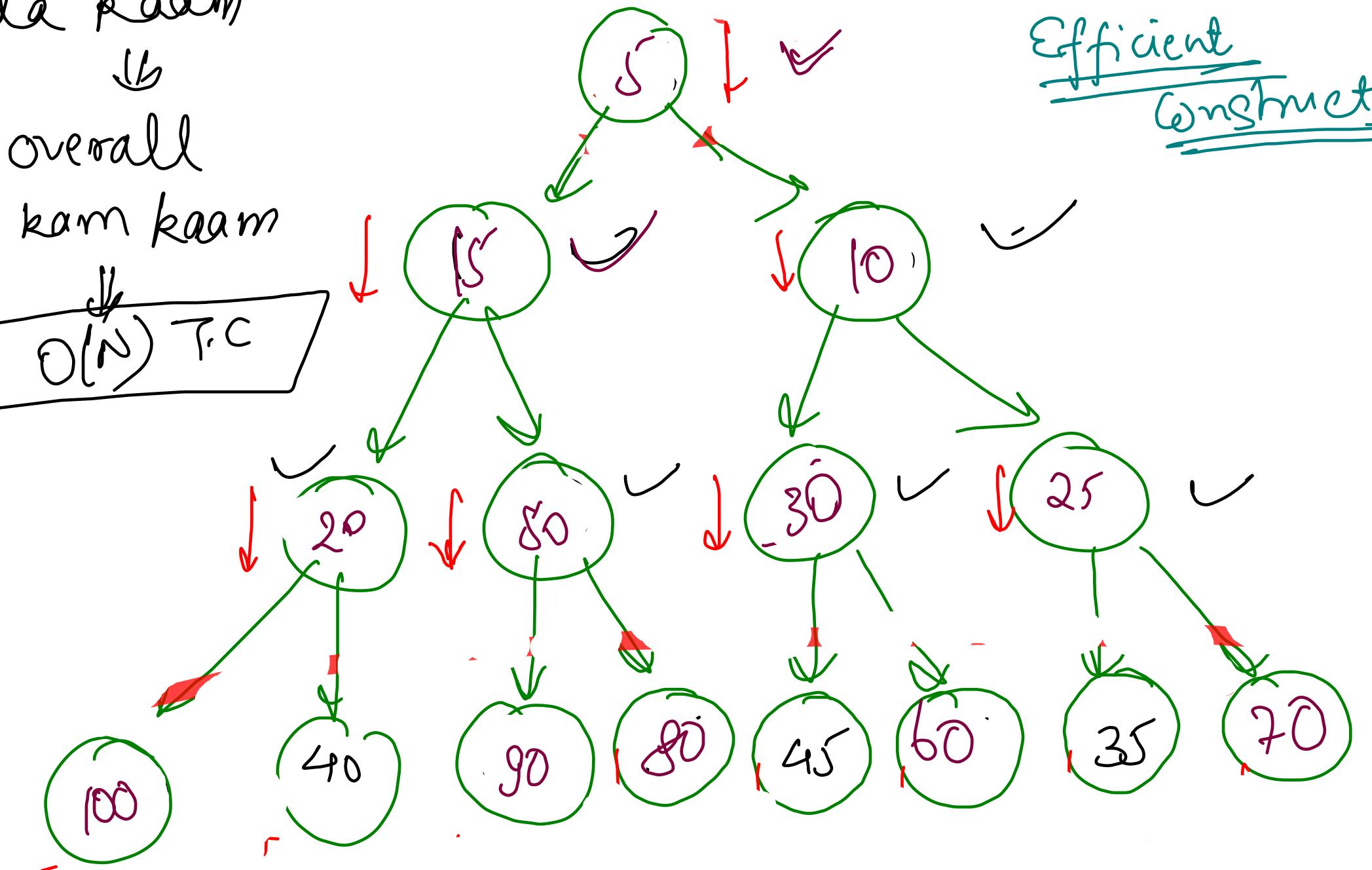
+ 2² + 2

+ 2³ + 1

overall

kam kaam

$O(N)$ T.C



Efficient Constructor

AGP

$$T(n) = 2^0 \cdot h + 2^1 \cdot (h-1) + 2^2 \cdot (h-2) + 2^3 \cdot (h-3) + \dots + 2^{h-1} \cdot 1$$

$$+ 2^k T(n) = 2^0 \cdot h + 2^1 \cdot (h-1) + 2^2 \cdot (h-2) + \dots + 2^{h-1} \cdot 1$$

$$T(n) = -2^0 \cdot h + \{ 2^1 \cdot 1 + 2^2 \cdot 1 + 2^3 \cdot 1 + \dots + 2^{h-1} \cdot 1 \} + 2^h \cdot 1$$

$$T(n) = (2^h - 1) - 2^0 \cdot h$$

$$= 2^{\log_2 n} - 1 - \log_2 n$$

$$= n - \log n - 1$$

$$\boxed{h = \log_2 n}$$

$$\boxed{O(n)}$$

```
// O(n) for inserting n elements -> per element O(1)
public PriorityQueue(int[] arr){
    data = new ArrayList<>();
    for(int val: arr){
        data.add(val);
        size++;
    }
    for(int i=(size() - 1)/2; i>=0; i--){
        downheapify(i);
    }
}
```

Interview Prep

10 questions \Rightarrow 3-4 Ques

\Rightarrow 3-4 Ques Algo's

\Rightarrow 1-2 complete

~~360~~
~~660~~

3 monthly

5 classes * 4 weekly

= 20 classes

$20 \times 3 = 60$ classes

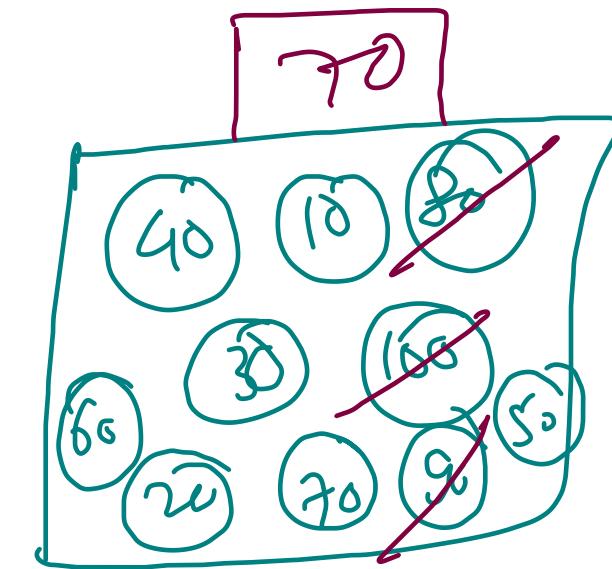
MTW

60 & 6 = 360

$\{ 40, 10, 80, 30, 100, 60, 20, 70, 90, 50 \}$ $k=3$

Approach 1 \rightarrow Max Heap

$\{ 100, 30, 80 \}$

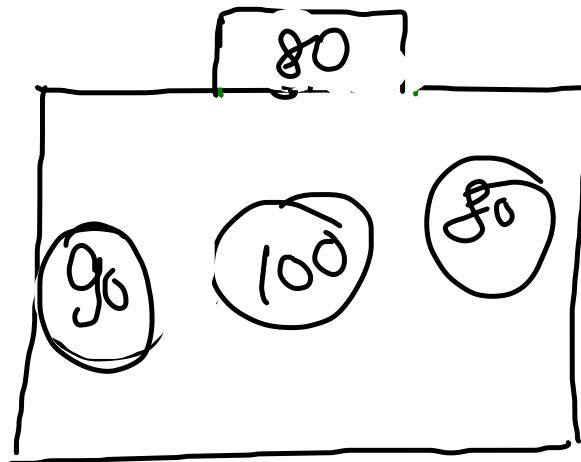


$\mathcal{O}(N + K \log N)$

```
public static void approach1(ArrayList<Integer> arr, int k){  
    PriorityQueue<Integer> pq = new PriorityQueue<>(Collections.reverseOrder());  
  
    // O(n)  
    pq.addAll(arr);  
  
    // O(k log N)  
    ArrayList<Integer> res = new ArrayList<>();  
    while(k-- > 0)  
        res.add(pq.remove());  
  
    for(int i=res.size() - 1; i>=0; i--){  
        System.out.println(res.get(i));  
    }  
}
```

$\{ 40, 10, 80, 30, 100, 60, 20, 70, 90, 50 \}$ $\underline{k=3}$

f min heap



$$O(k \log k + (n-k) \log k) = O(n \log k)$$

```
// O(n log k)
public static void approach2(ArrayList<Integer> arr, int k){
    // Min heap
    PriorityQueue<Integer> pq = new PriorityQueue<>();

    // O(k logk)
    for(int i=0; i<k; i++){
        pq.add(arr.get(i));
    }

    // O((n-k) logk)
    for(int i=k; i<arr.size(); i++){

        if(arr.get(i) > pq.peek()){
            pq.remove();
            pq.add(arr.get(i));
        }
    }

    while(pq.size() > 0){
        System.out.println(pq.remove());
    }
}
```