

Homotopy levels

We want to say things like "U is not a set".

- A set is something whose $=$ -types don't have structure.

Def. A type T 's h-level is 0 if

$$\text{hlevel } 0 T := \sum_{t:T} \prod_{s:T} s = t.$$

A type T 's h-level is n if

$$\text{hlevel } n T := \prod_{s,t:T} \text{hlevel } n s = t.$$

We've defined a function $\text{hlevel} : \mathbb{N} \rightarrow \text{Type} \rightarrow \text{Type}$.

h-level 0. AKA contractible, is Contr

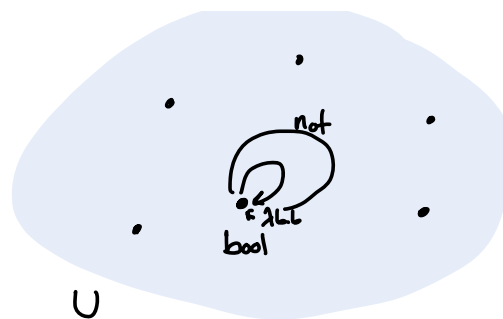
Ex. $\mathbb{1}$ is contractible.

Most boring.

h-level 1. AKA propositions, is Prop

Fact. Equivalent to $\prod_{x,y:P} x = y$.

Ex. $\emptyset, \mathbb{1}$ are propositions



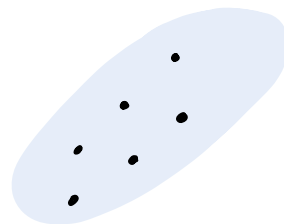
Ex. In fact, any contractible type is a proposition.

Ex. If a proposition is inhabited, it is contractible.

→ So roughly, a proposition is $=$ to \top or \perp .

So these behave like logical propositions where Σ behaves like \exists , etc.

h-level 2. AKA sets, is Set



Fact. bool , \mathbb{N} are propositions

h-level 3. AKA groupoids



Fact. Type has h-level at least 3.

Ex. If a type T has h-level n , then it has h-level $n+1$.

Equivalence

Sometimes we want types to be propositions (no structure). Sometimes we're interested in structure.

Given $f: A \rightarrow B$, want a proposition $\text{isEquiv}(f)$.

The type $\sum_{g: B \rightarrow A} fg = 1 \times gf = 1$ is not a proposition.

→ Could ask for adjoint equivalence, or equivalently:

Def. A function $f: A \rightarrow B$ is an equivalence if:
 $\text{isEquiv}(f) := \prod_{b:B} \text{isContr} \left(\sum_{a:A} f a = b \right).$

$\nwarrow \quad \nearrow$
 $\downarrow = \text{fiber}$

Write

$$A \simeq B := \sum_{f:A \rightarrow B} \text{isEquiv}(f).$$

Fact. For every type A , $A \simeq A$, so we can define
 $\text{id to equiv} : A = B \rightarrow A \simeq B.$

Def. The univalence axiom asserts
 $\text{ua} : \text{isEquiv}(\text{id to equiv}).$

Univalence for logics and sets

Def. $\text{Prop} := \sum_{P:\text{Type}} \text{isProp}(P)$ (Since isProp is a proposition
 $\text{isProp} \leq \text{Type}$)

Fact. The univalence axiom implies
 $(P =_{\text{Prop}} Q) \simeq (P \leftrightarrow Q).$

Def. $\text{Set} := \sum_{S:\text{Type}} \text{isSet}(S)$

Fact. The univalence axiom implies

$$(P =_{{\text{Set}}} Q) \simeq (P \cong Q).$$

$$\begin{aligned} \text{Def. } \text{Grp} := & \sum_{G: \text{Set}} \sum_{e: G} \sum_{\substack{m: G \rightarrow G \\ -G}} \sum_{\substack{i: G \\ -G}} \prod_{x: G} (m(e, x) = x) \times (m(x, e) = x) \\ & \times \prod_{x, y, z: G} ((xy)z = x(yz)) \\ & \times \prod_{x: G} (m(ix, x) = \text{id } x \\ & \quad (m(x, ix) = e)). \end{aligned}$$

Fact. The univalence axiom implies

$$(G =_{{\text{Grp}}} H) \simeq (G \cong H)$$