



SCHOOL OF COMPUTER SCIENCE AND ENGINEERING

Fall Semester 2023-2024

CONTINUOUS ASSESSMENT TEST – II

(Common Question Paper & Key)

Programme Name & Branch : B.Tech – All Specialization

Course Name & code : BCSE302L – Database Systems

SLOT : A1+TA1 & A2+TA2

Exam Duration : 90 Min.

Maximum Marks: 50

ANSWER ALL THE QUESTIONS(5X10=50 Marks)

Q.No	Question	Module	Max Mark	CO	BL
1	<p>a. Consider the following relation and functional dependencies. $R(A,B,C,D,E,F)$, $F = \{ A \rightarrow BC, AC \rightarrow DEF, F \rightarrow AB \}$</p> <p>i. List all of the candidate keys (1 mark)</p> <p>Answer:</p> <p>$A^+ = A$ $= ABC \quad A \rightarrow BC$ $= ABCDEF \quad AC \rightarrow DEF$</p> <p>$F^+ = F$ $= FAB \quad F \rightarrow AB$ $= FABC \quad A \rightarrow BC$ $= FABCDE \quad AC \rightarrow DEF$</p> <p>Candidate Keys are A and F.</p> <p>ii. Give the minimal cover for F. (5 mark)</p> <p>Answer:</p> <p>$A \rightarrow BC$:</p> <p>Case 1: B can be extraneous?</p>	3	6+4	CO2	BL3

$F' = \{A \rightarrow C, AC \rightarrow DEF, F \rightarrow AB\}$

A^+ on $F' = ACDEFB$ contains attribute B, So B is
extraneous.

$F = \{A \rightarrow C, AC \rightarrow DEF, F \rightarrow AB\}$

$AC \rightarrow DEF$:

Case 1.1: A is extraneous?

$C^+ = C$ does not contains attribute DEF, So A is not

Extraneous

Case 1.2: C is extraneous?

$A^+ = ADEFB$ does contains attribute DEF, So C is

Extraneous

$F = \{A \rightarrow C, A \rightarrow DEF, F \rightarrow AB\}$

$A \rightarrow DEF$:

Case 1.3: D is extraneous?

$F' = \{A \rightarrow C, A \rightarrow EF, F \rightarrow AB\}$

A^+ on $F' = CEFAB$ not contain D, Hence D is not extra.

Case 1.4: E is extraneous?

$F' = \{A \rightarrow C, A \rightarrow DF, F \rightarrow AB\}$

A^+ on $F' = CDFAB$ not contain E, E is not extra.

Case 1.5: F is extraneous?

$F' = \{A \rightarrow C, A \rightarrow DE, F \rightarrow AB\}$

A^+ on $F' = CDEAB$ does not contain F, F is not extra.

$F \rightarrow AB$:

Case 1.6: A is extraneous?

$F' = \{A \rightarrow C, A \rightarrow DEF, F \rightarrow B\}$

F^+ on $F' = FB$ does not contain A, A is not extra.

Case 1.7: B is extraneous?

$F' = \{A \rightarrow C, A \rightarrow DEF, F \rightarrow A\}$

F^+ on $F' = FACDEF$ not contain B, B is not extra.

$F_c = \{ A \rightarrow C, A \rightarrow DEF, F \rightarrow AB \} \rightarrow \{ A \rightarrow CDEF, F \rightarrow AB \}$

All left hand side are unique and no extraneous attribute in FD.

$A \rightarrow BC$

Case 2: C can be extraneous

$F' = \{ A \rightarrow B, AC \rightarrow DEF, F \rightarrow AB \}$

A^+ on $F' = AB$ not contain attribute C, C is not extra.

So Final minimal cover $F_c = \{ A \rightarrow CDEF, F \rightarrow AB \}$

b. Given a relational schema $R(A, B, C, D, E)$ and a set of functional dependencies P and Q such that:

$P = \{ A \rightarrow B, AB \rightarrow C, D \rightarrow AC, D \rightarrow E \}$, $Q = \{ A \rightarrow BC, D \rightarrow AE \}$

Check whether P and Q are equivalent? (4 mark)

Answer:

To check $P \subseteq Q$

Let's find closure of the left side of each FD of P using FD Q.

1. $A^+ = ABC$ (using $A \rightarrow BC$)
2. $AB^+ = ABC$ (using $A \rightarrow BC$)
3. $D^+ = DAEBC$ (using $D \rightarrow AE$ and $A \rightarrow BC$)
4. $D^+ = DAEBC$ (using $D \rightarrow AE$ and $A \rightarrow BC$)

Now compare closure of each A, AB, D and D calculated using FD Q with the right-hand side of FD P. Closure of each A, AB, D and D has all the attributes which are on the right-hand side of each FD of P. Hence, we can say **P is a subset of Q**

To check $Q \subseteq P$

Using definition of equivalence of FD set, let us determine the right-hand side of the FD set of Q using FD set P.

Given $P = \{ A \rightarrow B, AB \rightarrow C, D \rightarrow AC, D \rightarrow E \}$ and $Q = \{ A \rightarrow BC, D \rightarrow AE \}$

Let us find closure of the left side of each FD of Q using FD P.

1. $A^+ = ABC$ (using $A \rightarrow B$ and $AB \rightarrow C$)
2. $D^+ = DACEB$ (using $D \rightarrow AC, D \rightarrow E$, and $A \rightarrow B$)

	Now compare closure of each A, D calculated using FD P with the right-hand side of FD Q. Closure of each A and D has all the attributes which are on the right-hand side of each FD of Q. Q is a subset of P Hence P is equivalent to Q																															
2	<p>Consider a Student table which has information about the subjects and the activity the students have enrolled in. Subject and Activity are the two attributes that are independent of each other.</p> <p>a. Check whether the given table is in 4NF? If not, partition the table and convert to 4NF. (5 mark)</p> <table><tr><th><u>Student</u></th><th><u>Subject</u></th><th><u>Activity</u></th></tr><tr><td>Jones</td><td>Data Structures</td><td>P01</td></tr><tr><td>Bella</td><td>Algorithms</td><td>P03</td></tr><tr><td>Bella</td><td>Computation</td><td>P06</td></tr><tr><td>Bella</td><td>Algorithms</td><td>P07</td></tr><tr><td>Bella</td><td>Computation</td><td>P08</td></tr><tr><td>Margot</td><td>Database</td><td>P09</td></tr><tr><td>Margot</td><td>Compiler</td><td>P09</td></tr><tr><td>Princy</td><td>Network</td><td>P06</td></tr></table> <p>Answer:</p> <p>The table consists of multivalued dependency since all the column of the table has duplicate values.</p> <p>And the Subject and Activity are independent.</p> <p>Hence</p> <p>Student \twoheadrightarrow Subject and Student \twoheadrightarrow Activity exists in the table.</p> <p>All the attribute of the table act as key.</p> <p>So the Given table is not in 4NF.</p>	<u>Student</u>	<u>Subject</u>	<u>Activity</u>	Jones	Data Structures	P01	Bella	Algorithms	P03	Bella	Computation	P06	Bella	Algorithms	P07	Bella	Computation	P08	Margot	Database	P09	Margot	Compiler	P09	Princy	Network	P06	3	5+5	C02	BL6
<u>Student</u>	<u>Subject</u>	<u>Activity</u>																														
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We decompose the table into R1, R2, and R3 such as

Table R1:

<u>Student</u>	<u>Subject</u>
Jones	Data Structures
Bella	Algorithms
Bella	Computation
Margot	Database
Margot	Compiler
Pricy	Network

Key: {Student, Subject}

Table R2:

<u>Student</u>	<u>Activity</u>
Jones	P01
Bella	P03
Bella	P06
Bella	P07
Bella	P08
Margot	P09
Pricy	P06

Key: {Student, Activity}

Hence the table are in 4NF.

- b. If the partition tables are in 4NF, check whether it is in 5NF? (5 mark)

Answer:

If above partition is considered then it is not in 5NF.

We decompose the table into R1, R2, and R3 such as

Table R1:

<u>Student</u>	<u>Subject</u>
Jones	Data Structures
Bella	Algorithms
Bella	Computation
Margot	Database
Margot	Compiler
Pricy	Network

Key: {Student, Subject}

Table R2:

<u>Student</u>	<u>Activity</u>
Jones	P01
Bella	P03
Bella	P06
Bella	P07
Bella	P08
Margot	P09
Pricy	P06

Key: {Student, Activity}

Table R3:

Key: {Subject, Activity}

<u>Subject</u>	<u>Activity</u>
Data Structures	P01
Algorithms	P03

Computation	P06
Algorithms	P07
Computation	P08
Database	P09
Compiler	P09
Network	P06

Hence the table are in 4NF.

To check R1, R2 and R3 are in 5NF should satisfy

$$R1 \bowtie R3 \bowtie R2 \equiv R$$

We perform $R1 \bowtie_{\text{Subject}} R3$

<u>Student</u>	<u>Subject</u>	<u>Activity</u>
Jones	Data Structures	P01
Bella	Algorithms	P03
Bella	Algorithms	P07
Bella	Computation	P06
Bella	Computation	P08
Margot	Database	P09
Margot	Compiler	P09
Pricy	Network	P06

We perform $(R1 \bowtie_{\text{Subject}} R3) \bowtie_{\text{activity}} R2$

<u>Student</u>	<u>Subject</u>	<u>Activity</u>
Jones	Data Structures	P01
Bella	Algorithms	P03
Bella	Algorithms	P07
Bella	Computation	P06
Bella	Computation	P08

Margot	Database	P09
Margot	Compiler	P09
Princy	Network	P06

$R1 \bowtie R3 \bowtie R2 \equiv R$

Hence R1, R2 and R3 are in 5NF.

3 Consider the below relational database, where the primary keys are underlined.

employee (emp-name, street, city),
works (emp-name, company-name, salary),
company (company-name, city),
manages (emp-name, manager-name)

Give an expression in the relational algebra to express each of the following queries: (5X2=10 marks)

a. Find the employee's name whose salary is greater than 5000.

Answer:

$$\Pi_{emp-name} (\sigma_{salary > 5000} (works))$$

b. Find the names, street address, and cities of residence of all employees who work for Honda Company and earn more than 10,000.

Answer:

$$\Pi_{e.emp-name, e.street, e.city} (\sigma_{company-name="honda" \wedge salary > 10000 \wedge e.emp-name=emp-name} (\rho_e(employee) \times works))$$

c. Find the names of all employees in this database who live in the same city as the company for which they work.

Answer:

$$\Pi_{emp-name} ($$
 4 | 10 | C03 | BL6 |

	<p>$\sigma_{emp-name=w.emp-name \wedge city=c.city \wedge w.company-name=c.company-name}$</p> <p>$(emp \bowtie \rho_w(works) \bowtie \rho_c(company))$</p> <p>d. Find the names of all employees in this database who do not work for Honda Company.</p> <p>Answer:</p> <p>$\Pi_{emp-name}(works) - \Pi_{emp-name}(\sigma_{company="honda"}(works))$</p> <p>e. Find the total salary, minimum salary and maximum salary for each company.</p> <p>Answer:</p> <p>$company - name \mathcal{G}_{sum(salary),min(salary),max(salary)}(works)$</p>				
4	<p>Consider an ordered file with 50,000 records stored on a disk with block size $B = 512$ bytes. File records are of fixed size and are un-spanned, with record length $R = 50$ bytes.</p> <p>a. How many blocks are needed for the file? (1)</p> <p>Answer:</p> <p>Blocking Factor</p> <p>$bfr = \lfloor (B/R) \rfloor = \lfloor (512/50) \rfloor = 10 \text{ record/block}$</p> <p>Block needed for file</p> <p>$b = \lceil (r/bfr) \rceil = \lceil (50000/10) \rceil = 5000 \text{ blocks}$</p> <p>b. How many blocks access is needed if a binary search is performed? (1)</p> <p>Answer:</p> <p>Binary search takes $= \lceil \log_2 b \rceil = \lceil \log_2 5000 \rceil = 13 \text{ block access.}$</p> <p>c. Consider the ordering key field of the file is 8 bytes long, a block pointer, $P = 5$ bytes long, and a primary index is constructed for the file.</p> <p>i. What is the size of the index file? (1)</p> <p>Answer:</p> <p>Size of index $R_i = V + P = 8 + 5 = 13 \text{ bytes.}$</p>	4	10	C03	BL3

ii. What is the blocking factor for the primary index file?
(1)

Answer:

Blocking factor

$$bfr_i = \lfloor (B/R_i) \rfloor = \lfloor (512/13) \rfloor = 39 \text{ record/block}$$

iii. How many numbers of index entries will be there for the primary index file?
(1)

Answer:

No.of index entries = No.of blocks of file b = 5000 entries

iv. How many numbers of index block are needed? And how many block access will be performed to search for a record using the index? (2)

Answer:

$$\text{No. of index block } b_i = \lceil (r_i/bfr_i) \rceil = \lceil (5000/39) \rceil = 129 \text{ blocks}$$

For Linear search = 129 block access in worst case.

$$\text{For Binary search} = \lceil \log_2 b \rceil = \lceil \log_2 129 \rceil = 7+1=8 \text{ block access}$$

d. Consider non-ordering key field of the file is 10 bytes long and the block pointer, P = 5 bytes long, and a secondary index is constructed on the non-ordering key field of the file.

i. What is the size of secondary index?
(1)

Answer:

$$\text{Size of Secondary index } R_i = V+P = 10+5 = 15 \text{ bytes.}$$

ii. What is blocking factor for secondary index?
(1)

Answer:

Blocking factor

$$bfr_i = \lfloor (B/R_i) \rfloor = \lfloor (512/15) \rfloor = 34 \text{ record/block}$$

	<p>iii. How many numbers of index blocks are needed? and how many block access will be performed to search for a record using the secondary index? (1)</p> <p>Answer:</p> <p>No.of index entries r_i = No.of records in file = 50000 entries</p> <p>No. of index block</p> <p>$b_i = \lceil (r_i / bfr_i) \rceil = \lceil (50000 / 34) \rceil = 1471$ blocks</p> <p>For Linear search = 1471 block access in worst case.</p> <p>For Binary search = $\lceil \log_2 b \rceil = \lceil \log_2 1471 \rceil = 11 + 1 = 12$ block access</p>				
5	<p>Find which of these terms refers to Atomicity, Consistency, Isolation, Durability; and explain with example.</p> <ol style="list-style-type: none"> The changes of a successful transaction occurs even if the system fail occurs. The transaction takes place at once or doesn't happen at all. Multiple transactions occur independently without interference. The database must be consistent before and after the transaction. <p>Answer:</p> <ol style="list-style-type: none"> Durability Atomicity Isolation Consistency 	5	10	C04	BL2