Elliptic Curve Digital Signature Algorithm

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Public Key Cryptography

- ► Two keys (called a keypair)
 - ▶ Public: the whole world can know this
 - Private: only the owner knows this
- Four operations
 - ▶ Sign(Message, Private) \rightarrow Signature
 - ▶ Verify(Message, Signature, Public) \rightarrow Boolean
 - lacktriangle Encrypt(Message, Public) ightarrow Cyphertext
 - $\blacktriangleright \ \, {\tt Decrypt(Cyphertext, Private)} \to {\tt Message}$
- ► ECDSA is a widely-used (SSH, SSL/TLS) signature algorithm
 - Supports Sign and Verify

Cryptography and the Swarm

- Original motivation: the universal dataplane
 - Storage component of SwarmOS
 - ► Large network (10¹⁰ nodes)
 - ▶ Many low power sensors (μW)
 - Some high power servers
- Goal: all storage is fault-tolerant
 - Byzantine fault tolerence algorithm
 - Sign every message on the sensors
- Problem: ECDSA is mJ per signature!

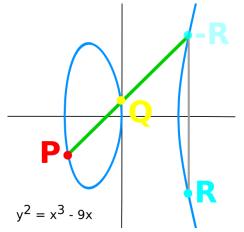
The Discrete Logarithm Problem

- $\triangleright b^k = g$
 - Difficult in one direction (given b and g find k)
 - ► Easy in another direction (given b and k find g)
- At the center of many public-key crypto schemes
- Only asymetric for some groups
 - ▶ $(\mathbb{Z}_p)^{\times}$: Integers modulo some prime, repeating multiplication
 - ightharpoonup GF(p): Elliptic curves modulo some prime, repeating addition
- Integer fields have a sub-linear algroithm

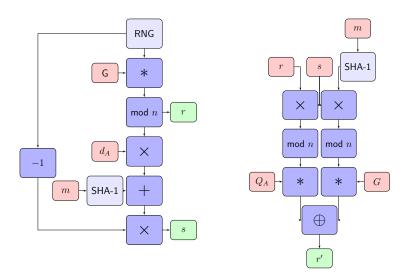
Security Level	DSA	ECDSA	Ratio
80	1024	160	3:1
112	2048	224	6:1
128	3072	256	10:1
192	7680	384	32:1
256	15360	512	64:1

Elliptic Curve Cryptography

- ▶ Elliptic curve: the set of points that satisfy $y^2 = x^3 + ax + b$
 - lacktriangle Need a finite field ightarrow everything's modular
 - lacktriangle Needs to be difficult ightarrow modulo a large number
- Add and multiply defined to satsify a field



ECDSA



Point Operations

- ECDSA is essentially point multiply and some cleanup
 - ▶ Point Multiply: O(million) cycles
 - Cleanup: O(thousand) cycles
- Point multiplication is defined as repeated point doubling
- Point doubling is defined as

$$\left(\frac{(3P_x^2+a)}{(2P_y)}-P_x\right)^2*\left(\frac{(3P_x^2+a)}{(2P_y)}\right)-P_y$$

- ▶ Simplifies to 5 modular multiplies and one modular divide
 - ► Modular operations are modulo a large (O(256-bit)) number

The Design Space

- A whole bunch of math
 - Data types: Points, Modular integers, and integers
 - Arithmatic operations: Add, Subtract, Multiply, Invert
- Everything boils down to integer arithmetic and control logic
 - ▶ What should be in software, what should be in hardware?
- Unfortunately, it's too big!
 - ▶ 2¹² software configurations!

Hardware-Software Cotuning

- OpenSSL is the industry standard, but it's difficult to hack on
- ▶ Wrote our own ECDSA implementation in C++
- Generates software for a family of ECDSA accelerators

```
palmer palmer-caldesktop rocket-ecc $ time make check | tail -n0

real 12m1.390s

user 63m20.796s

sys 7m37.179s

palmer palmer-caldesktop rocket-ecc $ ptest | tail -n4

NRUN 6895

NPASS 6895

NFAIL 0

NEROR 0
```

The (Reduced) Design Space

Point Multiplication

- Control logic
- Point Addition
 - Control logic
 - Modular multiply
 - Control logic
 - Integer shifts, adds
 - Modular inverse
 - Control logic
 - Integer shifts, adds

- Two interesting hardware configurations
 - Point multiply hardware block
 - Modular multiply, inverse hardware blocks

Results

Platform	Power	Speed	Area
	mJ/op	op/sec	mm^2
OpenSSL (45nm)	20	1000	
x86 (45 nm)	4000	5	
Rocket	800	0.05	
Virtex 2 (90nm)	4000	250	
Virtex 6 (45nm)	500	4000	
Mod Mul	200	0.3	0.04
$Mod\ Mul + Div$	2	20	0.10
Point $Add+Dbl$	0.5	100	0.31
Point Mul	0.06	400	0.35

Future Work

- ► Time-space tradeoffs
 - Montgomery multipliers
 - Projective point representation
- Parallel software, multiple signatures in flight
- Hard-coded curve parameters
- ▶ 300MHz clock