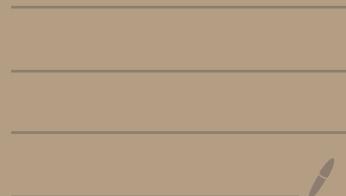


# Category Theory:

## a crash course



§1 Categories, functors, & Naturality

"To understand a structure, it is necessary to understand the morphisms that preserve it"

	Objects	morphisms	gouven
Sets	sets	functions	
Grp	groups	group homomorphisms	
Vect	vector spaces	linear maps	
Top	topological spaces	continuous fns	
	posets	order-preserving fns	
	.	.	
	:	:	
	?	?	
Cat	categories	functors	
	:	:	
	:	:	
	?	?	

Def: A category is (loosely) a collection of objects & a collection of morphisms between objects

(i) a collection of objects  $\text{Ob}(\mathcal{C})$

(ii) for any  $X \in \mathcal{C}$ , a set

$\text{Hom}_{\mathcal{C}}(X, Y) :=$  set of morphisms from  $X$  to  $Y$

(sometimes called "arrows")

Compatibility

$$X \xrightarrow{f} Y$$

(iii) for any  $X, Y, Z \in \text{Ob}(\mathcal{C})$ ,  $\text{Hom}_{\mathcal{C}}(X, Y) \times \text{Hom}_{\mathcal{C}}(Y, Z) \rightarrow \text{Hom}_{\mathcal{C}}(X, Z)$

$$X \xrightarrow{f} Y \xrightarrow{g} Z$$

$g \circ f = \text{composition}$

$$\text{Hom}_{\mathcal{C}}(X, Y) \times \text{Hom}_{\mathcal{C}}(Y, Z) \rightarrow \text{Hom}_{\mathcal{C}}(X, Z)$$

$$(X \xrightarrow{f} Y, Y \xrightarrow{g} Z) \mapsto X \xrightarrow{g \circ f} Z$$

space  
of matrices

$$\text{Hom}_{\mathcal{C}}(\mathbb{R}^n, \mathbb{R}^m) = M_{m \times n}(\mathbb{R})$$

## Compatibility

(iv) function composition is associative

$$h \circ (g \circ f) = (\text{h o g}) \circ f$$

$$X \xrightarrow{f} Y, Y \xrightarrow{g} Z, Z \xrightarrow{h} W$$

(v) identity morphisms  $\text{id}_X$

for any  $X \in \text{Ob}(C)$ ,  $\exists \text{id}_X \in \text{Hom}_C(X, X)$

$$\text{for my } e.g. e \cdot g = g \cdot e = g \quad g \in G$$

for any  $f: X \rightarrow Y$ ,  
 $g: Z \rightarrow X$ ,

Equivalent

$$\left\{ \begin{array}{l} f \circ \text{id}_X = f \\ g = \text{id}_Z \circ g \end{array} \right.$$

when are two objects "the same"?

Def: two objects  $X$  &  $Y$  of  $C$

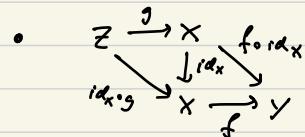
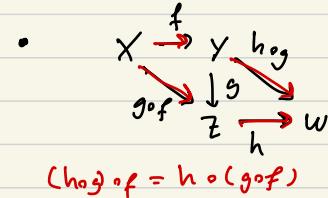
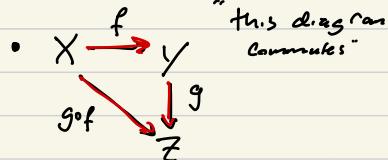
are isomorphic if  $\exists$  a pair of  
 morphisms

$$\begin{matrix} f \\ X \xleftrightarrow{\hspace{1cm}} Y \\ g \end{matrix}$$

such that  $\left\{ \begin{array}{l} g \circ f = \text{id}_X \\ f \circ g = \text{id}_Y \end{array} \right.$

$$\begin{matrix} f \\ X \xleftrightarrow{\hspace{1cm}} Y \\ g \end{matrix} \quad g = f^{-1}, \text{ etc.}$$

in terms of "commutative diagrams"



Examples: Set, Grp, Ab, Top, Vect<sub>K</sub>, Vect<sup>fd</sup><sub>K</sub>, Mor, Met, ...  
Banach Analytic Manifolds

Ban Ana Man

Ex: Let  $\mathcal{C}$  be a category. its opposite category,  $\mathcal{C}^{\text{op}}$ , is given by:

$$(i) \quad \text{Ob}(\mathcal{C}^{\text{op}}) = \text{Ob}(\mathcal{C})$$

$$(ii) \quad \text{Hom}_{\mathcal{C}^{\text{op}}}(X, Y) = \text{Hom}_{\mathcal{C}}(Y, X)$$

$$\begin{array}{ccc} & f & \\ X & \xrightarrow{\hspace{2cm}} & Y \\ & X \leftarrow Y & \text{in } \mathcal{C}^{\text{op}} \end{array}$$

Let  $R$  be a ring w/ unity,

$R\text{-mod}$  = category of left  $R$ -modules

the category of right  $R$ -modules

is  $(R\text{-mod})^{\text{op}}$

$\left. \begin{array}{l} \text{in } \mathcal{D}_X\text{-modules} \\ \text{sheaf of linear diff operators} \\ \text{w/ holomorphic coefficients} \\ \text{on a complex manifold} \end{array} \right\}$

Example

$$\Rightarrow \mathcal{C} = \text{Set}, \emptyset = \text{emptyset} \in \text{Ob}(\text{Set})$$

for any set  $X$ ,  $\emptyset \subseteq X$ .

Categorically: for any  $X$ ,  $\exists! \emptyset \hookrightarrow X$  { if  $Y$  satisfies this property for any  $X$ , then  $Y \cong \emptyset$  isomorphic }

$$\text{Hom}_{\text{Set}}(\emptyset, X) = \{\emptyset \rightarrow X\}$$

Exercise 1: Show that  $\emptyset$  is the only set with this property (up to isomorphism)

$X$  is an initial object in  $\mathcal{C}$

if<sup>2</sup>

$X$  is a terminal object in  $\mathcal{C}$

$$A = \emptyset$$

$$\begin{aligned} e_1, e_2 &\in G \\ e_1 \cdot e_2 &= e_2 \end{aligned}$$

suppose a set  
 $A$  has the same property as  $\emptyset$   
 $\Rightarrow$  for every  $X \in \text{Ob}(\text{Set})$ ,  $\exists A \rightarrow X$

$$A = B \text{ iff } A \subseteq B \text{ & } B \subseteq A$$

• by assumption,  
 $\emptyset \subseteq A \quad \emptyset \hookrightarrow A$   
but also, know  $\exists A \xrightarrow{\sim} \emptyset$

$$\begin{aligned} \text{Hom}(A, \emptyset) &\neq \{\text{id}_A: A \rightarrow \emptyset\} \\ (A \neq \emptyset) \end{aligned}$$

Example: "the" singleton set  $\{a\} \in \text{Ob}(\text{set})$ ,  
a set containing only one element.  
are these different?

$$A = \{a\} \xrightarrow{\sim} \{\text{chair}\} = B$$

uniquely  
isomorphic!

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ a & \mapsto \text{chair} & \text{chair} \mapsto a \end{array} \quad \begin{array}{ccc} B & \xrightarrow{g} & A \end{array}$$

the singleton set is the unique (up to isomorphism)

set  $\{*\}$  satisfying:

$$\text{for any set } X, \exists! X \xrightarrow{\alpha} \{*\}$$

$$\text{Horn}_{\text{Set}}(X, \{*\}) = \left\{ \begin{array}{l} X \xrightarrow{\alpha} \{*\} \\ X \xrightarrow{\beta} \bullet \end{array} \right\}$$

terminal  
object

Exercise 2: formulate this property in terms an arbitrary category  $\mathcal{C}$ , & prove it uniquely defines this object.

$Z$  is terminal in  $\mathcal{C}$   
if for every  $Y \in \text{Ob}(\mathcal{C})$ ,  
 $\exists! Y \rightarrow Z$

Exercise 3: for  $\text{Vect}_k$ , is there an initial object or a terminal object?

$\Rightarrow$  yes

$\emptyset \in \text{Vect}_k \dots \text{and?}$

any v-space  $V$  contains a zero element  $0 \in V$

$\emptyset \rightarrow V$  is always linear

on the other hand,  $\exists! V \rightarrow \emptyset$

$$\hat{v} \mapsto 0$$

$\emptyset$  is initial & terminal!  
↑  
"zero object"

in  $\text{Grp}$ ,  $\{\epsilon\}$  is a zero object!

partially ordered set

Example

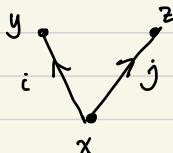
let  $(X, \leq)$  be a poset.

Consider as a category vvv:

$$\text{Ob}(X, \leq) = X$$

for any two  $x, y \in X$ ,

$$\text{Hom}_{(X, \leq)}(x, y) = \begin{cases} \{x \rightarrow y\}, & \text{if } x \leq y \\ \emptyset, & \text{else} \end{cases}$$



$$\text{Hom}(X, y) = \{i\}$$

$$\text{Hom}(y, z) = \emptyset$$

$$\text{Hom}(x, z) = \{j\}$$

$$R = (\mathbb{Z}, +, \cdot) \text{ is}$$

a Ring

$$(\mathbb{Q}, +, \cdot) \text{ is}$$

a field

$\text{Vect}_{\mathbb{Q}}$  is "nice"

$R\text{-mod}$  is "nice"

Let  $X$  be any set.

$P(X) = 2^X =$  set of subsets  
of  $X$ ,  
partially ordered by inclusion.  
↓  
Complete Boolean  
algebra

Let  $X$  be a topological space,

$\text{Op}(X) =$  set of open subsets of  $X$ ,  
partially ordered by inclusion.

comes up all the time in algebraic

geometry, for Topos theory,

Grothendieck topology . . .

New!

Exercise:

Show, for any set  $X$ ,  
that there is a bijection

$$P(X) \cong \text{Hom}_{\text{Set}}(X, \{0, 1\})$$

(power set)  
of  $X$

$$\begin{array}{ccc} X & \xrightarrow{\quad \sim \quad} & X \rightarrow \{0, 1\} \\ \downarrow \text{id}_X & & x \mapsto \begin{cases} 1, & \text{if } x \in A \\ 0, & \text{if } x \notin A \end{cases} \\ A & & \text{characteristic function of } A \end{array}$$

In Set,  $\{0, 1\}$  is called

a subobject classifier

Ex:

Monoids.

like groups, but w/o inverses  
 $(\mathbb{Z}, +)$  v.  $(\mathbb{N}, +)$   
group monoid

monoid multiplication  
by  $\otimes$

$$M$$

$$\begin{array}{ccc} M & \xrightarrow{x} & M \\ & \otimes & \downarrow y \\ & y \otimes x = x \otimes y & \end{array}$$

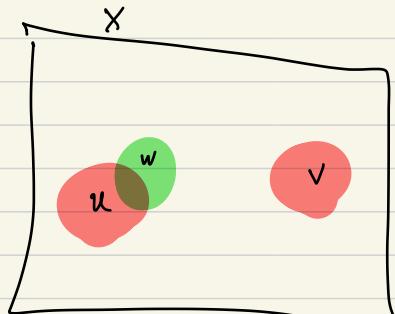
$$\begin{array}{c} M \times M \rightarrow M \\ \{(x, y) \mapsto x \otimes y\} \end{array}$$

turn into a category:  $(M)$

only one object,  $\text{Ob}(M) = \{*\}$

$$\text{Hom}_M(M, M) = M$$

$$x \in M \Leftrightarrow x : (M) \rightarrow (M)$$



U  $\hookrightarrow X$

V  $\hookrightarrow$

Exercise 4: prove the statement

"a group is a category with one element, where every morphism is an isomorphism"

what is a "structure-preserving" morphism between categories?

$$\mathcal{C} \xrightarrow{F} \mathcal{D}$$

F should define  
a function

?

"Functor"

(i)  $F: \text{Ob}(\mathcal{C}) \rightarrow \text{Ob}(\mathcal{D})$

$$X \mapsto F(X)$$

(ii)  $(X \xrightarrow{f} Y) \mapsto (F(X) \xrightarrow{F(f)} F(Y))$

$\text{in } \mathcal{C} \qquad \qquad \text{in } \mathcal{D}$

$$\text{Hom}_{\mathcal{C}}(X, Y) \xrightarrow[F]{f_X} \text{Hom}_{\mathcal{D}}(F(X), F(Y))$$

respect composition

(iii)  $F(gof) = F(g) \circ F(f)$

$F(gof) = F(f) \circ F(g)$  covariant ⚡ contravariant

(iv)  $F(id_X) = id_{F(X)}$

"Every sufficiently good analogy

yearns to become a functor"

-Baez

Category theory := mathematics  
of analogies

→ Grothendieck topology  
on a category

$$\mathbb{N} = (\mathbb{N}, +)$$

$\mathbb{Z}_{\geq 0} \quad \cup \quad 0$

natural numbers under addition

$$m: \mathbb{N} \rightarrow \mathbb{N} \in \text{Hom}_{\mathbb{N}}(\mathbb{N}, \mathbb{N})$$

$$n \mapsto n+m$$

$$k, m \in \mathbb{N}$$

$$k \cdot m: \mathbb{N} \xrightarrow{\quad} \mathbb{N} \xrightarrow{k} \mathbb{N}$$

$$n \mapsto (n+m) \mapsto (n+m)+k$$

$$2: \mathbb{N} \rightarrow \mathbb{N}$$

$$n \mapsto n+2$$

$$0, 1 \notin \text{Im}(2: \mathbb{N} \rightarrow \mathbb{N})$$


---

$$\mathbb{Z} \xrightarrow{+2} \mathbb{Z}$$

$$n \mapsto n+2$$

Surjective  
& injective.

$$g \in G$$

$$g: G \rightarrow G$$

$$h \mapsto h \cdot g$$

$$g^{-1}: G \rightarrow G$$

$$h \mapsto h \cdot g^{-1}$$

$g \circ g^{-1}$  are inverses

Ex: for any category  $\mathcal{C}$ , there's an identity functor:  $\text{id}_{\mathcal{C}} : \mathcal{C} \rightarrow \mathcal{C}$

$$x \mapsto x$$

$$f \mapsto f$$

"forgetful functor"

Ex:  $U: \text{Grp} \rightarrow \text{Set}$

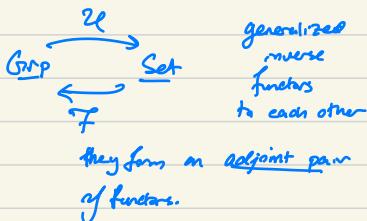
$$(G, \cdot) \mapsto G$$

forget the binary operation!

$F: \text{Set} \rightarrow \text{Grp}$

$$S \mapsto \text{Free}(S) = \left\{ \begin{array}{l} \text{free group on the} \\ \text{set } S \end{array} \right\}$$

look this up,  
if hard, talk about it later!



Ex:  $\mathcal{C} = \text{Vect}_k$  dual vector space functor.

$$D: \underline{\text{Vect}}_k^{op} \rightarrow \underline{\text{Vect}}_k$$

$$V \mapsto V^* = \text{Hom}_k(V, k)$$

$\text{Hom}_k(V, k)$

$$\text{Hom}_k(W, k) \rightarrow \text{Hom}_k(V, k)$$

$$(w \mapsto k) \mapsto (v \mapsto \delta)$$

$$(\delta \circ f) = f^*(\delta)$$

Exercise: describe  $D^2: \underline{\text{Vect}}_k \rightarrow \underline{\text{Vect}}_k$   
functor  
as "double dual"

$$D^2(V) = D(D(V)) = \text{Hom}_k(\text{Hom}_k(V, k), k)$$

& its relationship w/ the identity functor

$$\text{Id}_{\underline{\text{Vect}}_k}: \underline{\text{Vect}}_k \rightarrow \underline{\text{Vect}}_k$$

first problem session  
next Sunday

for any  $V \in \text{Vect}_k$ ,  
define a function

$$V \mapsto V^{**} = D^2(V)$$

$$\text{Id}_{\underline{\text{Vect}}_k} \rightarrow D^2$$

" $\delta$  is a function between functors"

$\delta$  is a natural transformation

restricted to subcategory

FinVect<sub>k</sub>,  $\delta$  is an isomorphism.

but not in general!  
why?

$$\begin{aligned} V^* &= \text{Hom}_k(V, k) \\ \xi: V &\rightarrow k \\ V \xrightarrow{\delta_V} V^{**} &= \text{Hom}_k(V^*, k) = \text{Hom}_k(\text{Hom}_k(V, k), k) \\ \vec{v} \mapsto \left\{ \begin{array}{l} \xi: V \rightarrow k \\ \vec{v}^* \end{array} \right\} &\mapsto \left\{ \begin{array}{l} \xi(\vec{v}) \in k \\ \vec{v}^* \end{array} \right\} \end{aligned}$$

$\delta_V(\vec{v})$  = "Evaluate the functional"  
at  $\vec{v}$

Lecture 2: Natural transformations  
&  
The Yoneda Lemma

given two categories  $\mathcal{C}$  &  $\mathcal{D}$ ,  
a functor  $F: \mathcal{C} \rightarrow \mathcal{D}$

→ a structure-preserving morphism  
between categories

defines  $F: \text{Ob}(\mathcal{C}) \rightarrow \text{Ob}(\mathcal{D})$   
 $X \mapsto F(X)$

$$\text{Hom}_{\mathcal{C}}(X, Y) \mapsto \text{Hom}_{\mathcal{D}}(F(X), F(Y))$$

$$\left( \begin{matrix} f \\ X \rightarrow Y \end{matrix} \right) \mapsto \left( \begin{matrix} F(f) \\ F(X) \rightarrow F(Y) \end{matrix} \right)$$

→ a functor  $F: \mathcal{C} \rightarrow \mathcal{D}$   
is full if, for any  $X, Y \in \text{Ob}(\mathcal{C})$ ,  
the function  
 $\text{Hom}_{\mathcal{C}}(X, Y) \rightarrow \text{Hom}_{\mathcal{D}}(F(X), F(Y))$   
is surjective.

$F: \mathcal{C} \rightarrow \mathcal{D}$  is faithful if

$$\text{Hom}_{\mathcal{C}}(X, Y) \rightarrow \text{Hom}_{\mathcal{D}}(F(X), F(Y))$$

is injective.

→  $F$  is fully faithful if  
" "  
is a bijection.

$$\text{Hom}_{\mathcal{C}}(X, Y) \cong \text{Hom}_{\mathcal{D}}(F(X), F(Y))$$

Ex: (full)  
subcategories:

$$\begin{array}{ccc} \underline{\text{FinSet}} & \xhookrightarrow{l} & \underline{\text{Sets}} \\ x & \mapsto & x \\ x \xrightarrow{f} y & \mapsto & x \xrightarrow{f} y \end{array}$$

↑ inclusion functor

$$\text{Hom}_{\underline{\text{FinSets}}}(X, Y) \cong \text{Hom}_{\underline{\text{Set}}}(X, Y)$$

$\underline{\text{FinVect}_k}$  = finite dimensional vector spaces over a field  $k$ .

$\underline{\text{Vect}_k}$

$$\text{Hom}_{\underline{\text{FinVect}_k}}(V, W) \cong \text{Hom}_{\underline{\text{Vect}_k}}(V, W).$$

### Natural Transformations

given  $F, G: \mathcal{C} \rightarrow \mathcal{D}$ , what is a

morphism  $\eta: F \rightarrow G$  ?

associate, for  $X \in \text{Ob}(\mathcal{C})$ , a morphism in  $\mathcal{D}$ :

$$\eta_X: F(X) \rightarrow G(X)$$

The component of  $\eta$  over  $X$

for any morphism  $X \xrightarrow{f} Y$ , a

commutative diagram:

$$G(f) \circ \eta_X = \eta_Y \circ F(f)$$

$$\begin{array}{ccc} & \eta_X & \\ F(X) & \xrightarrow{\quad} & G(X) \\ F(f) \downarrow & & \downarrow G(f) \\ F(Y) & \xrightarrow{\quad} & G(Y) \\ & \eta_Y & \end{array}$$

this diagram commutes

identity functor

$$\text{id}_{\underline{\text{Vect}}_k} = \mathbb{1}_k : \underline{\text{Vect}}_k \rightarrow \underline{\text{Vect}}_k$$



$$D = \underline{\text{Vect}}_k^{\text{op}} \rightarrow \underline{\text{Vect}}_k$$

$$V \mapsto \text{Hom}_{\underline{\text{Vect}}_k}(V, k) = D(V)$$

↙ a natural transformation  $\mathbb{1}_k \rightarrow D$

but this does work for  $D^2$ , double dual.

$$\exists \quad \mathbb{1}_k \rightarrow D^2$$

$$(\text{id}_{\underline{\text{Vect}}_k} \xrightarrow{\sim} D^2_{\underline{\text{Vect}}_k}) \text{ isomorphism}$$

→ a finite-dim. vector space is  
naturally isomorphic to its double-dual.

⇒ Yoneda Lemma:

for any  $X \in \text{Ob}(\mathcal{C})$ ,  $\text{Hom}_{\mathcal{C}}(X, -) \in \text{Ob}(\underline{\text{Set}})$

$$h_X := \text{Hom}_{\mathcal{C}}(X, -) : \mathcal{C} \rightarrow \underline{\text{Set}}$$
$$y \mapsto \text{Hom}_{\mathcal{C}}(X, y)$$

$$h_X(g) = \text{Hom}_{\mathcal{C}}(X, y) \rightarrow \text{Hom}_{\mathcal{C}}(X, z)$$

$$(x \xrightarrow{g} y) \mapsto (x \xrightarrow{g \circ f} z)$$

$$\begin{matrix} x & \xrightarrow{g \circ f} & z \\ y & \xrightarrow{g} & z \end{matrix}$$

$$h_X(g) = g_* \text{ "pushforward along } g$$

$$h^x := \text{Hom}_{\mathcal{C}}(-, x) : \mathcal{C}^{\text{op}} \rightarrow \text{Set}$$

$$y \mapsto \text{Hom}_{\mathcal{C}}(y, x) = h^x(y)$$

$$\begin{array}{ccc} y & \xrightarrow{g} & z \\ f \downarrow & \nearrow g^{-1} & \\ x & & \end{array}$$

$$h^x(z) \xrightarrow{g^*} h^x(y) = h^x(y \xrightarrow{g} z)$$

$$\text{Hom}_{\mathcal{C}}(y, x) \leftarrow \text{Hom}_{\mathcal{C}}(z, x)$$

$$\begin{array}{ccc} y & \xrightarrow{f} & x \\ g \downarrow & \nearrow f^{-1} & \\ z & & \end{array}$$

$$z \xrightarrow{f^{-1}} x$$

$$h^x(g)(\ell) := f \circ g$$

$\Rightarrow$  what does a natural transformation

$$\eta : h^x \rightarrow G$$

look like, for some  $G : \mathcal{C}^{\text{op}} \rightarrow \underline{\text{Set}}$ ?

This is context of Yoneda lemma.

$\Rightarrow$  denote  $[\mathcal{C}^{\text{op}}, \underline{\text{Set}}]$  the category  
of contravariant functors from  $\mathcal{C}$  to  $\underline{\text{Set}}$ .

Yoneda

$$\text{Hom}_{[\mathcal{C}^{\text{op}}, \underline{\text{Set}}]}(h^x, G) \cong G(x)$$

$$(y \in \text{Hom}_{\mathcal{C}}(-, x) \rightarrow G) \mapsto \exists \in G(x)$$

In particular, if  $G = h^y$ , there is a bijection.

$$\text{Hom}_{[\mathcal{C}^{\text{op}}, \underline{\text{Set}}]}(h^x, h^y) \cong h^y(x) \\ = \text{Hom}_{\mathcal{C}}(X, Y)$$

left as an exercise ...

Stop now and prove it yourself!

Want a way to associate

$$(\gamma: h^x \rightarrow G) \mapsto ? \in G(x)$$

$$\begin{array}{c} \exists \\ id_x \in \text{Hom}_{\mathcal{C}}(X, X) = h^x(X) \xrightarrow{\gamma_X} (G(X)) \\ \text{char} \\ id_X \leftarrow \gamma_{X(id_X)} = ?: x_i \end{array}$$

Exercise 1: Show that  $\exists := \gamma_{X(id_X)} \in G(X)$

Uniquely determines the component

$$\gamma_y: h^x(y) \rightarrow G(y)$$

for any  $y \in \text{ob}(\mathcal{C})$ .

$$h^x \xrightarrow{\gamma} G \mapsto \gamma_{X(id_X)} \in G(X)$$

Exercise 2: Show this is a bijection!

$$\Rightarrow \text{Hom}_{[\mathcal{C}^{\text{op}}, \underline{\text{Set}}]}(h^x, h^y) \cong \text{Hom}_{\mathcal{C}}(x, y)$$

$$h^{\wedge}: \mathcal{C} \rightarrow [\mathcal{C}^{\text{op}}, \underline{\text{Set}}]$$

$$X \mapsto h^X := \text{Hom}_{\mathcal{C}}(-, X)$$

by Yoneda,  $h^{\wedge}$  is fully faithful (embedding)

$h^{\wedge}$  embeds  $\mathcal{C}$  as a full subcategory  
of  $[\mathcal{C}^{\text{op}}, \underline{\text{Set}}]$ .

next class, we'll see  $[\mathcal{C}^{\text{op}}, \underline{\text{Set}}]$  is much  
nicer than  $\mathcal{C}$

analogy w/ analogies,

{ Smooth Fns }  $\rightarrow$  { Distributions }

$$f \mapsto \{ g \mapsto \int_X f g \}$$

analogous  
to  
 $\mathcal{C} \hookrightarrow \{ \text{Set}^{\text{op}}, \text{Set} \}$

$$X \mapsto \{ Y \mapsto \text{Hom}(Y, X) \}$$

Yoneda gives a full subcategory of  
functions  $\mathcal{C}^{\text{op}} \rightarrow \text{Set}$  isomorphic to  $\mathcal{C}^{\text{op}}$

for some  $X$ :

$\rightarrow$  if  $G \cong h^X$  for some  $X$ ,  
call  $G$  a representable functor.

$\rightarrow$  in problem session 1;

for any set  $X$ ,

$$\mathcal{P}(X) = \text{power set of } X$$

$$= \{ \text{Set} \hookrightarrow \text{Subsets of } X \}$$

in Set,  $\mathcal{P}(X)$  is a  
subobject classifier

by previous,

$$\mathcal{P}(X) \cong \text{Hom}_{\text{Set}}(X, \mathcal{P}, \text{id})$$

$$A \mapsto \{ X_A(x) = \begin{cases} 1, & \text{if } x \in A \\ 0, & \text{if } x \notin A. \end{cases} \}$$

$\rightarrow$  For any category  $\mathcal{C}$ ,

&  $X \in \text{Ob}(\mathcal{C})$

$$\text{Sub}_{\mathcal{C}}(X) = \{ \text{Subobjects of } \mathcal{C} \hookrightarrow X \}$$

$$\text{Sub}_{\mathcal{C}} : \mathcal{C}^{\text{op}} \rightarrow \text{Set}$$

$$\mathcal{P} : \underline{\text{Set}}^{\text{op}} \rightarrow \underline{\text{Set}}$$

$$x \xrightarrow{\exists} \text{determines } P(y) \xrightarrow{\exists} P(X)$$

$P_{12}^{10,11}$  is a representable functor!

whether or not  $\text{Sub}_{\mathcal{C}}$  is representable

is equivalent to  $\mathcal{C}$

having a subobject classifier  $\mathcal{P} \in \text{Ob}(\mathcal{C})$

important in categorical  
logic, topos theory

Exercise 3: define & prove a covariant

version of Yoneda lemma

$$h_X : \text{Hom}_{\mathcal{C}, \text{Set}}(X, -) \rightarrow F \text{ in } [\mathcal{C}, \text{Set}]$$

covariant functors

$$\text{Hom}_{[\mathcal{C}, \text{Set}]}(h_X, F) \cong F(X)$$

& analogue of yoneda embedding:

$$h^{\vee} : \mathcal{C}^{\text{op}} \rightarrow [\mathcal{C}, \text{Set}]$$

### 3. Examples

The central point about examples of representable functors is:

*Representable functors are ubiquitous.*

To a fair extent, category theory is all about representable functors and the other universal constructions: Kan extensions, adjoint functors, limits, which are all special cases of representable functors – and representable functors are special cases of these.

Listing examples of representable functors in category theory is much like listing examples of integrals in analysis: one can and does fill books with these. (In fact, that analogy has more to it than meets the casual eye: see coend for more).

Keeping that in mind, we do list some special cases and special classes of examples that are useful to know. But any list is necessarily wildly incomplete.

Let  $\mathcal{C}$  be a Category

$X \in \text{ob}(\mathcal{C})$

$$\left\{ \begin{array}{l} h^X := \text{Hom}_{\mathcal{C}}(-, X) : \mathcal{C}^{\text{op}} \rightarrow \underline{\text{Set}} \\ h^X(Y) := \text{Hom}_{\mathcal{C}}(Y, X) \\ h_X = \text{Hom}_{\mathcal{C}}(X, -) : \mathcal{C} \rightarrow \underline{\text{Set}} \end{array} \right.$$

$$h^X(Y \xrightarrow{f} Z) = h^X(Z) \xrightarrow{f^*} h^X(Y)$$

$(Z \xrightarrow{g} X) \mapsto (Y \xrightarrow{g \circ f} X)$

$f$  pullback along  $f$

$y \xrightarrow{f} z$   
 $\downarrow g$   
 $x$

$$h_X(Y \xrightarrow{f} Z) = h_X(Y) \xrightarrow{f_*} h_X(Z)$$

$(X \xrightarrow{g} Y) \mapsto (X \xrightarrow{g \circ f} Z)$

$x \downarrow g$   
 $y \xrightarrow{f} z$   
 $f$  pushforward along  $f$

$$D: \underline{\text{Vect}}_k^{\text{op}} \rightarrow \underline{\text{Vect}}_k$$

(Abelian Category)

$$V \mapsto \text{Hom}_k(V, k) = h^k(V)$$

$$G: \mathcal{C}^{\text{op}} \rightarrow \underline{\text{Set}}$$

Yoneda

$$\text{Hom}_{[\mathcal{C}^{\text{op}}, \underline{\text{Set}}]}(h^x, G_x) \cong G(x)$$

morphisms

between  
functors

$G$  is representable

if  $G \cong h^x$

- $\mathcal{C}$  is small
- if  $\text{Ob}(\mathcal{C})$  is a set
- $\mathcal{C}$  is locally small
- if  $\text{Hom}_{\mathcal{C}}(X, Y)$  is a set

$$G: \mathcal{C}^{\text{op}} \rightarrow \underline{\text{Set}}$$

sheaves  
of sets, etc?  
enriched categories?

$$\text{Hom}_{[\mathcal{C}^{\text{op}}, \underline{\text{Set}}]}(h^x, h^y) \xrightarrow{\cong} \text{Hom}_{\mathcal{C}}(X, Y)$$

(Yoneda)

$$h^y(X) \cong \text{Hom}_{\mathcal{C}}(X, Y)$$

def

$$h^*: \mathcal{C} \rightarrow [\mathcal{C}^{\text{op}}, \underline{\text{Set}}]$$

$X \mapsto h^X$

$$\text{Hom}_{[\mathcal{C}, \underline{\text{Set}}]}(h_X, h_Y) \cong h_Y(X) = \text{Hom}_{\mathcal{C}}(Y, X)$$

$$h^*: \mathcal{C}^{\text{op}} \rightarrow [\mathcal{C}, \underline{\text{Set}}]$$

$$X \mapsto h_X$$

## Lecture 3: Universal Properties

"A universal property is a property that characterizes (up to isomorphism) the uniquely result of some construction"

Initial objects      } Uniquely characterized  
terminal objects    } by some property.

- in Set,  $\emptyset$  (the empty set)  
is "the" initial object.

$\emptyset \subseteq X$  for any set  $X$ .  
 $\emptyset \rightarrow X \Leftrightarrow$  "is a subset of every set"

- the singleton set  $\{\#\}$  is terminal in cat. of sets.

{all} vs. {choose}

For any set  $X$ ,  $\exists! X \rightarrow \{\#\}$   
 $x \mapsto *$

"admits a map from my set"  
 $\hookrightarrow X \rightarrow \{\#\}$

### Examples:

- free groups / vector spaces
- quotient groups / quotient vector spaces / quotient topological spaces
- (co)Kernels of group hom's / (co)Kernels of linear maps between  $v$ -spaces
- direct products, direct sums ...
- completion of a metric space (e.g.  $\mathbb{R}$  from  $\mathbb{Q}$ ,  $\mathbb{Z}_p$  ( $p$ -adics) from  $\mathbb{Q}$ )
- tensor products
- fiber products (pullbacks)

; Ubiquitous in math!

Exercise 1: find more

Example: the Kernel of a group homomorphism

$$G \xrightarrow{\varphi} H \quad \text{a grp hom.}$$

$\left[ \begin{array}{l} \varphi \text{ is a function between} \\ \text{the underlying sets of } G \text{ & } H \\ \text{s.t. } \forall g_1, g_2 \in G, \\ \varphi(g_1 * g_2) = \underset{H}{\varphi(g_1)} * \varphi(g_2) \end{array} \right]$

$$V \xrightarrow{\psi} W \quad k\text{-linear map of} \\ v \text{ spaces}/k$$

The Kernel of  $\varphi$ ,  $\text{Ker } \varphi$ , is  
the (normal) subgroup of  $G$  defined by  
usual defn.

$$\text{Ker } \varphi := \{g \in G \mid \varphi(g) = 0_H\}.$$

$$\mathbb{Z} \xrightarrow{\varphi} \mathbb{Z}/n\mathbb{Z}$$

$$m \mapsto m + n\mathbb{Z}$$

$$\text{Ker } \varphi = n\mathbb{Z} := \{nm \in \mathbb{Z} \mid m \in \mathbb{Z}\}.$$

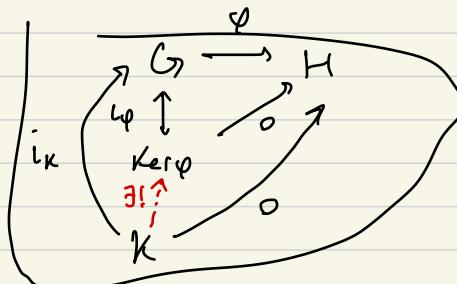
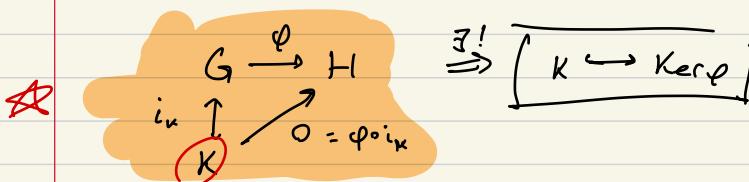
→ "Ker  $\varphi$  is the largest <sup>(normal)</sup> subgroup of  $G$  that is annihilated by  $\varphi$ "

commutative diagram

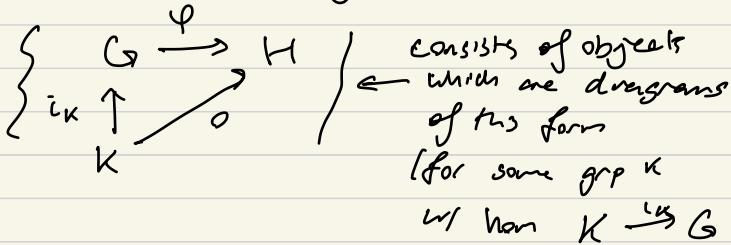
$$\begin{array}{ccc} G & \xrightarrow{\varphi} & H \\ i_\varphi \uparrow & & \searrow \\ \text{Ker } \varphi & & 0 = \varphi \circ i_\varphi \end{array}$$

$$\iota_K : K \hookrightarrow G$$

means: if  $K \trianglelefteq G$  is any other normal subgroup such that  $\phi_{|K} = 0$  (i.e.,  $\phi \circ \iota_K = 0$ ), then  $K \subseteq \text{Ker } \phi$



Define a category of diagrams



A morphism of such diagrams  
is given by

$$\left( \begin{array}{ccc} G & \xrightarrow{\phi} & H \\ i_1 \uparrow & \nearrow 0 & \\ K_1 & & \end{array} \right) \xrightarrow{\psi} \left( \begin{array}{ccc} G & \xrightarrow{\phi} & H \\ i_2 \uparrow & \nearrow 0 & \\ K_2 & & \end{array} \right)$$

Exercise 2: Check this is actually a category! Dip

Given by a group hom.

$$\psi: K_1 \rightarrow K_2$$

such that the following diagram commutes

$$\begin{array}{ccc} G & \xrightarrow{\phi} & H \\ i_2 \uparrow & \nearrow o & \downarrow \\ K_2 & & \\ \downarrow \psi & & \\ K_1 & & \end{array}$$

punchline:  $\begin{array}{ccc} G & \xrightarrow{\phi} & H \\ \psi \uparrow & \nearrow o & \\ \text{Ker } \phi & & \end{array}$  is the terminal object in the category of diagrams  $\Rightarrow$

Exercise 3: (part I) for  $K \xrightarrow{i_K} G$ , didn't require that  $K$  is a subgroup/normal subgroup of  $G$ .

Q: why does this still work?

Hint: think about Image ( $\phi$ ).

$$\begin{array}{ccc} G & \xrightarrow{\phi} & H \\ i_K \uparrow & \nearrow o & \\ K & & \end{array} \quad \text{vs} \quad \begin{array}{ccc} G & \xrightarrow{\phi} & H \\ \uparrow & \nearrow o & \\ \text{Im}(\phi) & & \end{array} ?$$

Exercise 3. (part II)

What's (if it exists) the property defining  $\text{Im}(\phi)$ ?

Similarly, for cokernels:

$$G \xrightarrow{\varphi} H$$

$$\text{coker } (\varphi) = \text{coker}(\varphi) := H / \text{Im } (\varphi)$$

coker  $\varphi$  is the smallest quotient group of  $H$  s.t.

next lecture,  
we will axiomatize  
this idea

$$\begin{array}{ccc} G & \xrightarrow{\varphi} & H \\ \downarrow & \lrcorner & \downarrow j_{\varphi} \\ \text{coker } \varphi & = & H / \text{Im } (\varphi) \end{array} \quad \leftarrow \begin{array}{l} \text{this is the initial object} \\ \text{in } \mathcal{D}^{\varphi}. \end{array}$$

cat. of diagrams:  $\mathcal{D}^{\varphi}$ .

$$\left\{ \begin{array}{ccc} G & \xrightarrow{\varphi} & H \\ \downarrow & \lrcorner & \downarrow j_L \\ \text{---} & & \end{array} \right\}$$

Example: from topology

form category  $\text{Top}$  of topological spaces & continuous func between them.

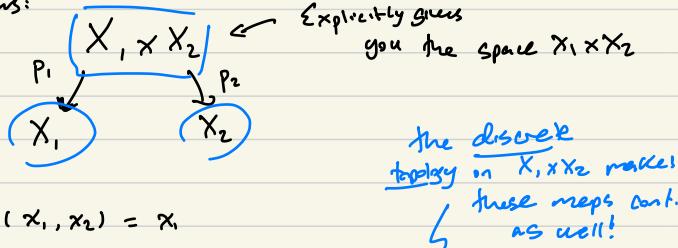
Let  $X_1$  &  $X_2$  be top. sp. Can form the product space  $\underline{X_1 \times X_2}$

as a set,

this is cartesian product.

with topology given

as follows:



$$p_1(x_1, x_2) = x_1$$

$$p_2(x_1, x_2) = x_2$$

the topology on  $\underline{X_1 \times X_2}$  is the coarsest topology for which  $p_1$  &  $p_2$  are continuous.

the discrete topology in  $X_1 \times X_2$  makes these maps cont.  
as well!

If  $U_i \subset X_i$  is open, then  $p_i^{-1}(U_i)$  is open in  $X_1 \times X_2$

universal properties identify objects that

are the "best"/"most optimal" at that property

→ e.g. biggest object such that, smallest / finest / coarsest, etc..

define category of diagrams of top. spaces

$$\begin{array}{ccc} & Y & \text{top. space} \\ f_1 \swarrow & \downarrow & \searrow f_2 \\ X_1 & X_2 & \text{Cont.} \end{array}$$

$$\begin{array}{ccc} & Y & \xrightarrow{\psi} Z \\ f_1 \curvearrowleft & \downarrow \psi & \downarrow \\ & Z & \\ g_1 \swarrow & \downarrow g_2 & \searrow f_2 \\ X_1 & X_2 & \text{Cont.} \\ \text{s.t. the diagram commutes.} \end{array}$$

Then,

$$\begin{array}{ccc} & X_1 \times X_2 & \\ p_1 \swarrow & \downarrow p_2 & \\ X_1 & & X_2 \end{array}$$

is the terminal object in this category.

The whole diagram matters,  
not just  $X_1 \times X_2$ .

$$\begin{array}{ccc} & Y & \\ f_1 \swarrow & \downarrow f & \searrow f_2 \\ X_1 & X_1 \times X_2 & X_2 \\ \downarrow p_1 & \downarrow p_2 & \downarrow \\ X_1 & & X_2 \end{array}$$

$p_i \circ f = f_i$

Exercise 4: If I reverse some arrows  
what am I defining?

$$\begin{array}{ccc} & X_1 & X_2 \\ f_1 \swarrow & \downarrow g_1 & \searrow g_2 \\ & Z & \\ \downarrow \psi & & \downarrow \\ & Y & \end{array}$$

Exercise 5:

Quotient vector spaces

Quotient groups

Quotient topological spaces

Seem similar...

Identify the universal property,

find the cat of diagrams,

& compute initial object / etc.