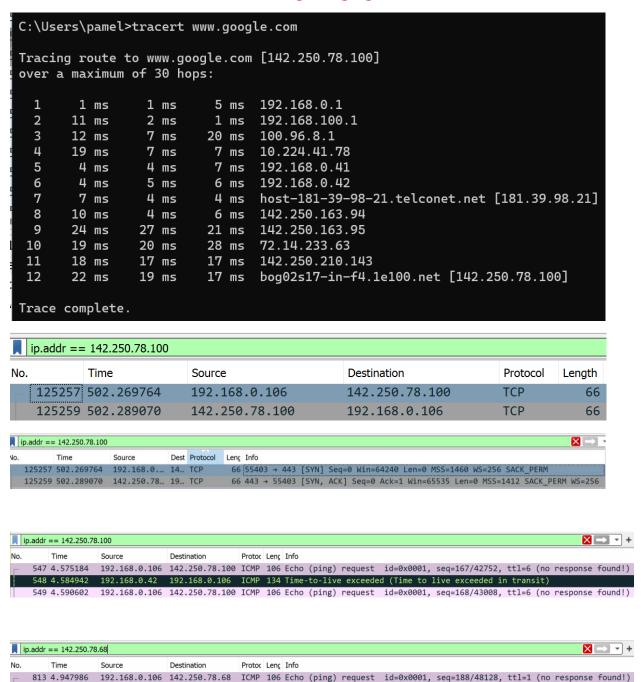
#### Pamela Mena - 00211564

### Universidad San Francisco de Quito

#### Tarea 03

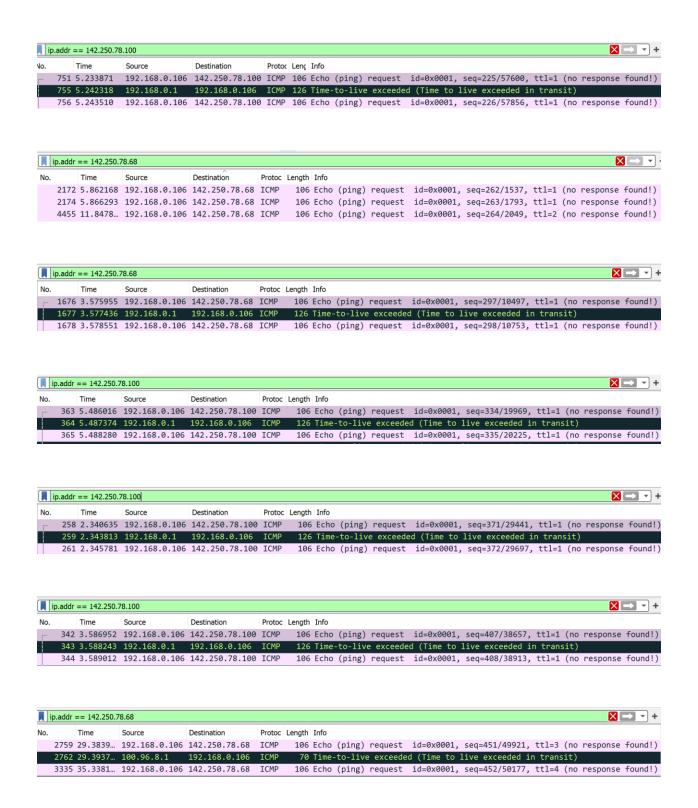
- 1. Read the following Wireshark tutorial and use it to capture traffic from the following scenarios. Use screenshots to show your results.
  - a. Run 10 traceroute commands against google.com.

192.168.0.106

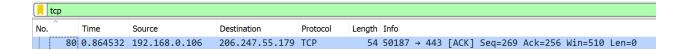


ICMP 126 Time-to-live exceeded (Time to live exceeded in transit)

815 4.950529 192.168.0.106 142.250.78.68 ICMP 106 Echo (ping) request id=0x0001, seq=189/48384, ttl=1 (no response found!)



b. Watch a video from youtube.com. Capture the TCP handshake, and the congestion window.



Transmission Control Protocol, Src Port: 50187, Dst Port: 443, Seq: 269, Ack: 256, Len: 0

Source Port: 50187
Destination Port: 443
[Stream index: 0]

[Conversation completeness: Incomplete (12)]

[TCP Segment Len: 0]

Sequence Number: 269 (relative sequence number)

Sequence Number (raw): 1701516990

[Next Sequence Number: 269 (relative sequence number)]

Acknowledgment Number: 256 (relative ack number)

Acknowledgment number (raw): 540568757 0101 .... = Header Length: 20 bytes (5)

> Flags: 0x010 (ACK)

Window: 510

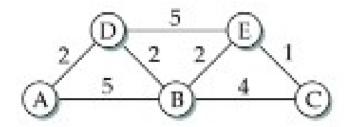
[Calculated window size: 510]

[Window size scaling factor: -1 (unknown)]

Checksum: 0x173c [unverified] [Checksum Status: Unverified]

Urgent Pointer: 0
> [Timestamps]

# 2. Use Dijkstra's to get the routing tables for nodes A, B and E.

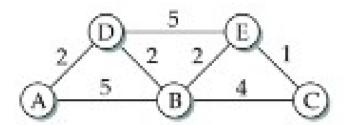


## Routing para nodo A:

Nodo	Costo	Salto
A	0	A
В	5	B - A
С	9	B - A
D	2	D - A
Е	7	D - A

## Routing para nodo B:

Nodo	Costo	Salto
A	5	A - B
В	0	В
С	4	C - B
D	2	D-B
Е	2	E - B



## Routing para nodo E:

Nodo	Costo	Salto
A	7	D-E
В	2	B-E
С	1	C - E
D	5	D-E
Е	0	Е

3. Suppose a host wants to establish the reliability of a link by sending packets and measuring the percentage that are received; routers, for example, do this. Explain the difficulty of doing this over a TCP connection.

Cuando los paquetes no llegan a su destino a través de una conexión TCP, generalmente se debe a congestión en lugar de problemas de confiabilidad del enlace. Identificar la causa de un paquete perdido puede ser difícil. La asignación insuficiente de recursos puede resultar en dos problemas principales: posible injusticia y un estado de bloqueo conocido como "deadlock". Primero, exploremos el "deadlock". Si todos los buffers de un router están llenos de paquetes, el router no podrá aceptar nuevos marcos. Aún más problemático, podría ignorar marcos con ACKs que podrían liberar algunos de esos buffers. Por ejemplo, si dos routers vecinos, A y B, se envían paquetes entre sí y ambos están esperando que el otro reciba un paquete, ninguno puede avanzar. Esta situación se conoce como un "deadlock".

4. Consider a simple congestion control algorithm that uses linear increase and multiplicative decrease (no slow start). Assume the congestion window size is in units of packets rather than bytes, and it is one packet initially.

- a. Give a detailed sketch of this algorithm.
- b. Assume the delay is latency only, and that when a group of packets is sent, only a single ACK is returned.
- c. Plot the congestion window as a function of RTT for the situation in which the following packets are lost: 9, 25, 30, 38 and 50. For simplicity, assume a perfect timeout mechanism that detects a lost packet exactly 1 RTT after it is transmitted.

ACK: cwnd = cwnd +1 /cwnd

Timeat: cwnd = min (1, cwnd/2)

Packet 9 bot and cwnd - 2:

Packet 25 lost and cwnd - 3:

Packet 30 lost and cwnd - 2:

Packet 38 lost and awnd - 2:

Packet 50 lost and cwnd - 2:

