

BLayout Documentation

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API

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Types

```
typedef size_t blsize;

struct blayout {
    blsize nmemb;
    blsize size;
    blsize align;
};
```

- `blsize` is the API's size type. It's `size_t` by default. You may change this type by modifying BLayout's header. A `signed` type is also valid. You'd have to change `BL_SIZEMAX` accordingly (see [below](#)).
- `blayout` describes a single memory allocation request for an object:
 - `nmemb` is the number of elements (like `calloc()`'s first argument),
 - `size` is the size (in bytes) of each element/type (like `calloc()`'s second argument),
 - `align` is the alignment¹ of the object's type

Constants

```
#define BL_SIZEMAX    SIZE_MAX
#define BL_ALIGNMENT  alignof(max_align_t)
```

- `BL_SIZEMAX` is the equivalent to `size_t`'s `SIZE_MAX` and is equal to that by default. You may override this, but the header assumes that it's **greater** than 0.
- `BL_ALIGNMENT` is never used internally. It's equal to the maximum alignment among C's scalar types. It's provided as a convenience when calling `blcalc()` (see [below](#)). This, too, can be overridden.

*Note: To override these, either modify BLayout's header or `#define` them **before** including `blayout.h`.*

Macros

```
#define BL_API      static
#define BL_ASSERT  assert
#define BL_INLINE  inline
#define BL_DEBUG    0
```

- `BL_API` is currently only used as a visual aid, do **not** try to change it.

¹ **alignment** is always assumed to be valid: (1) it denotes *byte* boundaries and (2) is a power of 2.

- BLayout can use assertions through the `BL_ASSERT` macro to enforce API contracts and prevent footguns. You can override this macro if you use a custom `assert()` function. See `BL_DEBUG` below if you want to disable assertions.
- Every function is `inline` (C99 [semantics](#)) through the `BL_INLINE` macro. This is so that you can workaround C's deficiencies, if you so wish.
- `BL_DEBUG` takes three possible values:
 - 0, where BLayout will use **no** assertions (see above) and, in addition, will take advantage of compiler-specific optimization hints (e.g. `attribute(nonnull(...))`). This is the default.
 - 1, where BLayout will use **some** assertions **and** optimization hints.
 - 2, where BLayout will use **all** assertions **and no** optimization hints.

Functions

```
BL_API blsize blcalc(blsize align,
                    ptrdiff_t offs,
                    blsize n,
                    const struct blayout *lays,
                    blsize prev_size);

BL_API void *blnext(void *ptr, blsize curr_size, blsize next_align);
BL_API void *blprev(void *ptr, blsize prev_size, blsize prev_align);

BL_API blsize blsizeof(const struct blayout *l);
```

- `blcalc()` returns the minimum size needed to contiguously allocate multiple objects. The function assumes that all arguments are valid and within bounds. If wrap-around is detected when computing the size, 0 is returned instead, indicating error.
 - `align` is the default alignment² your allocator supports. In case you already have an allocated block, pass the block's alignment. `BL_ALIGNMENT` should work with `malloc()` and with any memory block allocated by it.
 - `offs` is used in case you already have a block and want to allocate starting from an offset into that block. Pass 0 otherwise.
 - `n` is the number of layouts. Should be **greater** than 0.
 - `lays` is an array of length `n` containing layouts,
 - `prev_size` is used to chain multiple `blcalc()` calls. When first invoking, 0 must be passed, otherwise the result of the previous `blcalc()` call must be passed, assuming the call succeeded and a **non-0** value was returned. `align` and `offs` must not change across any chained calls.
- `blnext()` allocates the next object in a **left-to-right** manner, where:
 - `ptr` is a pointer to the current allocated object. When first invoking, pass a pointer to your block (or `block + offs` if `blcalc()` was passed a non-zero offset). It's assumed to **not** be `NULL` and thus the function doesn't check for this.
 - `curr_size` is the size of the current object (see `blsizeof()`) and is assumed to be valid. When first invoking, pass 0.
 - `next_align` is the alignment of the next object's type and is assumed to be valid. When first invoking, pass the alignment of the **first** object's type.
- `blprev()` is like `blnext()`, but allocates and returns the previous object, in a **right-to-left** manner. That means you should allocate in **reverse** order, starting with **last** object.
 - `ptr` is a pointer to the current allocated object. When first invoking, pass a pointer to the **end** of your block. It's assumed to **not** be `NULL` and thus the function doesn't check for this.
 - `prev_size` is the size of the previous object (see `blsizeof()`) and is assumed to be valid. When first invoking, pass the size of the **last** object.
 - `prev_align` is the alignment of the previous object's type and is assumed to be valid. When first invoking pass the alignment of the **last** object's type.
- `blsizeof()` returns the total size (in bytes) of an object described by its layout. Effectively, it multiplies `blayout.nmemb` with `blayout.size`. It's provided as a convenience.
 - `l` is the pointer to the aforementioned layout.
 1. *Caveat: Padding due to alignment is **not** taken into account.*
 2. *Caveat: Potential integer overflow is **not** checked. The layout is assumed to be correct. `blcalc()` already checks for this.*

Usage

1. `blcalc()` with `blnext()`
2. `blprev()`
3. `blnext()` VS. `blprev()`

²[alignment](#) is always assumed to be valid: (1) it denotes byte boundaries and (2) is a power of 2.

blcalc() with blnext()

```
/*
 * Request:
 *   I. One integer, naturally aligned.
 *   II. Two floats, naturally aligned.
 *
 * The order of the `lays` array is important! `blcalc()` takes the order into
 * account when computing its result. And it does that for the simple reason
 * that the whole point of this header is to allow the programmer to lay out
 * their objects in memory exactly how they want. Hence, we preserve the order,
 * because it might be important, we wouldn't know. If the order is _not_
 * important to _you_, this detail doesn't impair you.
 *
 * This order also defines how `blnext()` (and `blprev()`; see below) should be
 * called. The functions don't check for this, the burden, unfortunately, falls
 * onto the programmer.
 */
const struct blayout lays[] = {{1, sizeof(int), alignof(int) },
                               {2, sizeof(float), alignof(float)}};

size_t size = blcalc(BL_ALIGNMENT, /* Going to use the default alignment. */
                    0, /* Allocating from the `0`th position. */
                    2, /* The number of layouts. */
                    lays, /* The array of layouts describing our objects. */
                    0); /* We aren't chaining `blcalc()` calls. */

if (size == 0) {
    fprintf(stderr, "blcalc() error\n");
    return 1;
}

void *block = malloc(size);
if (block == NULL) {
    fprintf(stderr, "malloc() error\n");
    return 1;
}

/* The first object. */
int *i = blnext(block, /* First allocation: pass the block. */
                0, /* First allocation: pass `0`. */
                lays[0].align); /* Pass the alignment of the first object's type (`int`). */
assert(i != NULL); /* Guaranteed to _not_ be `NULL`. */

/* Continue with the next object... */
float *f = blnext(i, /* Pass the current allocated object. */
                 blsizeof(&lays[0]), /* Pass the size of the current allocated object. */
                 lays[1].align); /* Pass the alignment of the next object's type (`float`). */
assert(f != NULL); /* Likewise. */

/* Use `i` and `f` normally. */
*i = 42;
f[0] = 2.71;
f[1] = 3.14;
printf("*i=%d f[0]=%f f[1]=%f\n", *i, f[0], f[1]);

/* Cleanup. */
free(block); /* `i` and `f` are guaranteed to be cleaned up with a _single_ `free()`. */
return 0;

/*
 * Alternatively, if the alignment of the _first_ object's type is
 * _less-or-equal_ to your block's alignment, then the pointer to the first
 * allocated object is equivalent to a pointer to the block.
 *
 * In our case:
 *
 * 1. The first object's type has alignment alignof(int).
 * 2. BL_ALIGNMENT is the default alignment of malloc() (which we also
 *    used when calling blcalc()).
 * 3. Thus, the block returned to us by malloc() (block) has that
 *    alignment.
 * 4. malloc() must, as mandated by the C standard, be able to return a
 *    suitably aligned pointer for _every_ naturally aligned type. That
 *    includes our case: int.
 * 5. Since block has alignment BL_ALIGNMENT and that alignment _must_ be
 *    suitable for int, we can conclude that BL_ALIGNMENT >= alignof(int)
 *    or, equivalently, alignof(int) <= BL_ALIGNMENT.
 */
```

```

/*
 * Note that any type can be over-aligned. This is fine, because a greater
 * alignment guarantees natural alignment. Or, in the words of the standard[1]:
 *
 * In general, the concept "correctly aligned" is transitive: if a pointer
 * to type A is correctly aligned for a pointer to type B, which in turn is
 * correctly aligned for a pointer to type C, then a pointer to type A is
 * correctly aligned for a pointer to type C.
 *
 * [1]: <https://port70.net/~nsz/c/c11/n1570.html#note68>
 */
static_assert(alignof(int) <= BL_ALIGNMENT, "incorrect alignment"); /* Just to be sure. */
free(i); /* Same as `free(block);` _in our case_. See the above comments why and when this holds
        true. */
return 0;

/*
 * The above observation also implies that we could skip the first call to
 * `blnext()` and allocate the first object (`i`) like so:
 */
int *i = block; /* Okay, as long as `alignof(int) <= BL_ALIGNMENT`. */

/* With the rest of the code being identical... */

```

blprev()

Usage of `blprev()` is similar to the usage of `blnext()` with these notable differences:

1. You must allocate in **reverse** order.
2. In order to cleanup safely, you **must** retain a pointer to your block.

Let's see with a similar example:

```

/*
 * ... Same layouts and boilerplate as the `blnext()` example...
 *
 * Assume `block` of size `size` has already been allocated as above.
 */
size_t size;
void *block;

/*
 * Allocate in **reverse** order, starting with the **last** object. `blprev()`
 * works in a _right-to-left_ manner; we have to pass the _end_ of our block
 * in the first allocation.
 */
void *end = (char *)block + size;
float *f = blprev(end, /* First allocation: pass the end of the block. */
                  blsizeof(&lays[1]), /* First allocation: pass the size of the last object. */
                  lays[1].align); /* First allocation: pass the alignment of the last object's
                                type. */
assert(f != NULL); /* Always true, same as before. */

/* Continue with the previous object... */
int *i = blprev(f, /* Pass the current allocated object. */
               blsizeof(&lays[0]), /* Pass the size of the previous object. */
               lays[0].align); /* Pass the alignment of the previous object's type. */
assert(i != NULL); /* Likewise. */

*i = 1337;
f[0] = 1.41;
f[1] = 1.61;
printf("i=%d f[0]=%f f[1]=%f\n", *i, f[0], f[1]);

/*
 * Cleanup!
 *
 * NOTE: You can _not_ pass `i` or `f` in the case of `blprev()`! You must pass
 * the pointer returned to you by `malloc()`!
 */
free(block);
return 0;

//free(i); /* XXX: Don't do this _ever_! */
//free(f); /* XXX: Or this either! */

```

blnext () vs. blprev ()

When should you use the one or the other? Given the limitations of `blprev()`, shouldn't you always use `blnext()`? It depends. Firstly, you indeed **can only use one**³ of the two to allocate from a single memory region. Given that, which to pick?

Let's assume your layouts array is `{{1, 1, 1}, {1, 2, 2}}`. Meaning:

1. One (1) type of size 1 (bytes) at an 1-byte alignment boundary, and
2. One (1) type of size 2 (bytes) at an 2-byte alignment boundary.

Assuming your block is 16-byte aligned, `blcalc()` (correctly) returns 4 for the above array.

If your block sits at address 16 then, using `blnext()`, the bytes would be laid out like this:

```
Address | 16 17 18 19
Bytes   | X0   Y0 Y1
```

But, if you were to use `blprev()`, allocating from the end of the block, they'd be laid out like this:

```
Address | 16 17 18 19
Bytes   |     X0 Y0 Y1
```

Where x_i is the i th byte of object x , respectively y .

Thus, the two methods give different results. This example also makes clear why you need to retain a pointer to your block in the case of `blprev()`: `X0` doesn't sit at the (original) address 16!

Always use `blnext()`, unless you desire the special properties of `blprev()` and the limitations don't affect you. Here's two scenarios where that could be true:

- You always carry around a pointer to your block, so using `blprev()` has no extra burden. *Note that `blprev()` is 2-3 machine instructions shorter than `blnext()`, under `x86_64-sysv` and also depending on compiler and optimization options (Related).*
- You depend on the layout `blprev()` gives. This happens when giving a “header” to a “payload”, just like `malloc()` does.

```
struct header {
    int id;
};

struct payload {
    size_t size;
    unsigned char data[]; /* Flexible array member:
    <https://gustedt.wordpress.com/2011/03/14/flexible-array-member/> */
};

struct payload *new_payload(my_ctx *ctx, size_t payload_size)
{
    const struct layout lays[] = {
        {1, sizeof(struct header), alignof(struct header)},
        {1, offsetof(struct payload, data) + payload_size, alignof(struct payload)}
    };
    size_t size = blcalc(/* ... */, 0, 2, lays, 0);
    if (size == 0)
        abort();

    void *block_end = /* ... */;
    struct payload *p = blprev(block_end, blsizeof(&lays[1]), lays[1].align);
    p->size = payload_size;
    memset(p->data, 0, payload_size);

    struct header *h = blprev(p, blsizeof(&lays[0]), lays[0].align);
    h->id = /* ... */;

    return p;
}

void recycle_payload(my_ctx *ctx, struct payload *p)
{
    struct header *h = blprev(p, sizeof(*h), alignof(struct header)); /* Okay. */
    int id = h->id;
    /* Do something with `id`... */
}
```

³Which also means you can't use one function to retrieve objects in reverse order, if you allocated using the other.

The above example **cannot** work with `blnext()`, it **only works with `blprev()`**. Even if you only used `blnext()` to allocate the objects, once you got to retrieve the header in:

```
struct header *h = blprev(p, sizeof(*h), alignof(struct header)); /* I'm sure this is
    perfectly fine, what could possibly go wr */
int id = h->id; /* Kaboom! */
```

your computer would explode. You can't do this! *Even if it works*, you can't depend on this if you allocate using `blnext()`. Use `blprev()` instead. Remember that the two functions *lay out objects differently*.

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