# cgroup v2 block io note

## latency方式io控制器的实测可用条件:

- ext4
  - read ok
  - o write ok
- f2fs
  - 。 read ok (PS: 注意设置正确的blk device, 注意读不同的文件)
  - o write ok
  - 。 PS: rq\_qos似乎会降低磁盘吞吐量?
  - 。 TODO: 需要设计一个测试,记录request层每秒的吞吐量,证实有没有降低?

# 代码调用栈

#### readahead 调用栈 (from page\_fault)

/*******/ //CX
do_page_cache_readahead[157]
//CPU: 0 PID: 128 Comm: mdev Not tainted 4.19.100+ #14
//Hardware name: QEMU Standard PC (i440FX + PIIX, 1996), BIOS 1.10.2-1ubuntu1 04/01/2014
//Call Trace:
dump_stack+0x71/0x97
do_page_cache_readahead+0xbd/0x1c3
? find_get_entry+0x1e/0x176
? pagecache_get_page+0x2d/0x2e8
filemap_fault+0x255/0x661
ext4_filemap_fault+0x31/0x44do_fault+0x34/0xe0handle_mm_fault+0x106b/0x1535
handle_mm_fault+0xe0/0x24e
do_page_fault+0x3eb/0x57d
do_page_fault+0x30/0xe8
? page_fault+0x8/0x30
page_fault+0x1e/0x30

#### readahead 调用栈 (from syscall)

```
/*********////CX____
__do_page_cache_readahead[157] //CPU: 0 PID: 125 Comm: init Not tainted 4.19.100+ #14
```

//Hardware name: QEMU Standard PC (i440FX + PIIX, 1996), BIOS 1.10.2-1ubuntu1 04/01/2014 //Call Trace:

dump\_stack+0x71/0x97 \_\_do\_page\_cache\_readahead+0xbd/0x1c3 ? d\_absolute\_path+0x6b/0x9c ondemand\_readahead+0x1b5/0x2be page\_cache\_sync\_readahead+0xaf/0xd4 generic\_file\_read\_iter+0x305/0xa84 ext4\_file\_read\_iter+0x53/0xdf \_\_vfs\_read+0x13f/0x16f vfs\_read+0x91/0x13d kernel\_read+0x31/0x42 prepare\_binprm+0xfa/0x1cf \_\_do\_execve\_file+0x52a/0x818 do\_execve+0x2d/0x2f \_\_x64\_sys\_execve+0x2b/0x32 do\_syscall\_64+0x5c/0x128 entry\_SYSCALL\_64\_after\_hwframe+0x44/0xa9

# blk\_cgroup\_congest 调用栈

/\*\*\*\*\*\*\*\*\*\*\*/ dump\_stack+0x71/0x97
blk\_cgroup\_congested mem\_cgroup\_throttle\_swaprate+0x1d/0x160
mem\_cgroup\_try\_charge\_delay+0x37/0x43
\_\_handle\_mm\_fault+0xad1/0x1535
handle\_mm\_fault+0xe0/0x24e
\_\_do\_page\_fault+0x3eb/0x57d
do\_page\_fault+0x30/0xe8
? page\_fault+0x8/0x30
page\_fault+0x1e/0x30

## mem\_cgroup\_throttle\_swaprate 调用栈

## iolatency done bio 调用栈

/\*\*\*\*\*\*\*\*/ iolatency\_check\_latencies
blkcg\_iolatency\_done\_bio rq\_qos\_done\_bio bio\_endio req\_bio\_endio blk\_update\_request
scsi\_end\_request scsi\_io\_completion scsi\_finish\_command scsi\_softirq\_done blk\_done\_softirq
\_\_do\_softirq invoke\_softirq irq\_exit exiting\_irq do\_IRQ common\_interrupt

## vfs\_read blkcg\_iolatency\_throttle 调用栈

/\*\*\*\*\*\*\*\*/ \_\_blkcg\_iolatency\_throttle blkcg\_iolatency\_throttle rq\_qos\_throttle blk\_queue\_bio generic\_make\_request submit\_bio ext4\_mpage\_readpages ext4\_readpages read\_pages \_\_do\_page\_cache\_readahead ra\_submit

ondemand\_readahead page\_cache\_async\_readahead generic\_file\_buffered\_read generic\_file\_read\_iter ext4\_file\_read\_iter call\_read\_iter new\_sync\_read \_\_vfs\_read vfs\_read

#### vfs\_read to lkcg\_iolatency\_throttle 调用栈

/\*\*\*\*\*\*\*/ blkcg\_iolatency\_throttle(struct rq\_qos \* rgos, struct bio \* bio, spinlock\_t \* lock) (/home/panard/linux-4.19.100/block/blk-iolatency.c:436) rq\_qos\_throttle(struct request\_queue \* q, struct bio \* bio, spinlock\_t \* lock) (/home/panard/linux-4.19.100/block/blk-rq-qos.c:77) blk queue bio(struct request queue \* q, struct bio \* bio) (/home/panard/linux-4.19.100/block/blk-core.c:2060) generic make request(struct bio \* bio) (/home/panard/linux-4.19.100/block/blk-core.c:2464) submit\_bio(struct bio \* bio) (/home/panard/linux-4.19.100/block/blk-core.c:2573) ext4 mpage readpages(struct address space \* mapping, struct list head \* pages, struct page \* page, unsigned int nr pages, bool is readahead) (/home/panard/linux-4.19.100/fs/ext4/readpage.c:293) ext4\_readpages(struct file \* file, struct address\_space \* mapping, struct list\_head \* pages, unsigned int nr\_pages) (/home/panard/linux-4.19.100/fs/ext4/inode.c:3364) read pages(struct address space \* mapping, struct file \* filp, struct list\_head \* pages, unsigned int nr\_pages, gfp\_t gfp) (/home/panard/linux-4.19.100/mm/readahead.c:123) \_\_do\_page\_cache\_readahead(struct address\_space \* mapping, struct file \* filp, unsigned long offset, unsigned long nr to read, unsigned long lookahead size) (/home/panard/linux-4.19.100/mm/readahead.c:215) ra\_submit() (/home/panard/linux-4.19.100/mm/internal.h:66) do sync mmap readahead() (/home/panard/linux-4.19.100/mm/filemap.c:2467) filemap\_fault(struct vm\_fault \* vmf) (/home/panard/linux-4.19.100/mm/filemap.c:2543) ext4\_filemap\_fault(struct vm\_fault \* vmf) (/home/panard/linux-4.19.100/fs/ext4/inode.c:6353) do fault(struct vm fault \* vmf) (/home/panard/linux-4.19.100/mm/memory.c:3269) do\_read\_fault() (/home/panard/linux-4.19.100/mm/memory.c:3681) do\_fault() (/home/panard/linux-4.19.100/mm/memory.c:3810) handle\_pte\_fault() (/home/panard/linux-4.19.100/mm/memory.c:4041) \_\_handle\_mm\_fault(struct vm\_area\_struct \* vma, unsigned long address, unsigned int flags) (/home/panard/linux-4.19.100/mm/memory.c:4165) handle\_mm\_fault(struct vm\_area\_struct \* vma, unsigned long address, unsigned int flags) (/home/panard/linux-4.19.100/mm/memory.c:4202) \_\_do\_page\_fault(struct pt\_regs \* regs, unsigned long error\_code, unsigned long address) (/home/panard/linux-4.19.100/arch/x86/mm/fault.c:1390) do\_page\_fault(struct pt\_regs \* regs, unsigned long error\_code) (/home/panard/linux-4.19.100/arch/x86/mm/fault.c:1465) page fault() (/home/panard/linux-4.19.100/arch/x86/entry/entry\_64.S:1204) [Unknown/Just-In-Time compiled code] (Unknown Source:0) irq\_stack\_union (Unknown Source:0) [Unknown/Just-In-Time compiled code] (Unknown Source:0)

# f2fs: (cgroup io部分与f2fs的实现有关联)

#### vfs\_read to blkcg\_iolatency\_throttle 调用栈

```
/*************************
blkcg_iolatency_throttle(struct rq_qos * rqos, struct bio * bio, spinlock_t * lock) (/
rq_qos_throttle(struct request_queue * q, struct bio * bio, spinlock_t * lock) (/home/
blk_queue_bio(struct request_queue * q, struct bio * bio) (/home/panard/linux-4.19.100
generic_make_request(struct bio * bio) (/home/panard/linux-4.19.100/block/blk-core.c:2
submit_bio(struct bio * bio) (/home/panard/linux-4.19.100/block/blk-core.c:2573)
 _submit_bio() (/home/panard/linux-4.19.100/fs/f2fs/data.c:307)
f2fs_mpage_readpages(struct address_space * mapping, struct list_head * pages, struct
f2fs_read_data_pages(struct file * file, struct address_space * mapping, struct list_h
read_pages(struct address_space * mapping, struct file * filp, struct list_head * page
 _do_page_cache_readahead(struct address_space * mapping, struct file * filp, <mark>unsigned</mark>
ra_submit() (/home/panard/linux-4.19.100/mm/internal.h:66)
do_sync_mmap_readahead() (/home/panard/linux-4.19.100/mm/filemap.c:2467)
filemap_fault(struct vm_fault * vmf) (/home/panard/linux-4.19.100/mm/filemap.c:2543)
f2fs_filemap_fault(struct vm_fault * vmf) (/home/panard/linux-4.19.100/fs/f2fs/file.c:
__do_fault(struct vm_fault * vmf) (/home/panard/linux-4.19.100/mm/memory.c:3269)
do_read_fault() (/home/panard/linux-4.19.100/mm/memory.c:3681)
do_fault() (/home/panard/linux-4.19.100/mm/memory.c:3810)
handle_pte_fault() (/home/panard/linux-4.19.100/mm/memory.c:4041)
__handle_mm_fault(struct vm_area_struct * vma, unsigned long address, unsigned int fla
handle_mm_fault(struct vm_area_struct * vma, unsigned long address, unsigned int flags
 _do_page_fault(struct pt_regs * regs, unsigned long error_code, unsigned long address
do_page_fault(struct pt_regs * regs, unsigned long error_code) (/home/panard/linux-4.1
page_fault() (/home/panard/linux-4.19.100/arch/x86/entry/entry_64.S:1204)
[Unknown/Just-In-Time compiled code] (Unknown Source:0)
irq_stack_union (Unknown Source:0)
[Unknown/Just-In-Time compiled code] (Unknown Source:0)
```

# f2fs read流程 到 generic\_make\_request 调用栈

```
generic_make_request(struct bio * bio) (/home/panard/linux-4.19.100/block/blk-core.c:2
submit_bio(struct bio * bio) (/home/panard/linux-4.19.100/block/blk-core.c:2573)
  submit_bio() (/home/panard/linux-4.19.100/fs/f2fs/data.c:307)
f2fs_mpage_readpages(struct address_space * mapping, struct list_head * pages, struct
f2fs_read_data_pages(struct file * file, struct address_space * mapping, struct list_h
read_pages(struct address_space * mapping, struct file * filp, struct list_head * page
 _do_page_cache_readahead(struct address_space * mapping, struct file * filp, unsign<mark>ed</mark>
ra_submit() (/home/panard/linux-4.19.100/mm/internal.h:66)
ondemand_readahead(struct address_space * mapping, struct file_ra_state * ra, struct f
page_cache_async_readahead(struct address_space * mapping, struct file_ra_state * ra,
generic_file_buffered_read(ssize_t written) (/home/panard/linux-4.19.100/mm/filemap.c:
generic_file_read_iter(struct kiocb * iocb, struct iov_iter * iter) (/home/panard/linu
call_read_iter() (/home/panard/linux-4.19.100/include/linux/fs.h:1814)
new_sync_read() (/home/panard/linux-4.19.100/fs/read_write.c:406)
 _vfs_read(struct file * file, <mark>char *</mark> buf, <mark>size_t</mark> count, loff_t * pos) (/home/panard/l
vfs_read(struct file * file, char * buf, size_t count, loff_t * pos) (/home/panard/lin
ksys_read(unsigned int fd, char * buf, size_t count) (/home/panard/linux-4.19.100/fs/r
__do_sys_read() (/home/panard/linux-4.19.100/fs/read_write.c:589)
__se_sys_read() (/home/panard/linux-4.19.100/fs/read_write.c:587)
```

\_\_x64\_sys\_read(const struct pt\_regs \* regs) (/home/panard/linux-4.19.100/fs/read\_write do\_syscall\_64(unsigned long nr, struct pt\_regs \* regs) (/home/panard/linux-4.19.100/ar entry\_SYSCALL\_64() (/home/panard/linux-4.19.100/arch/x86/entry/entry\_64.S:238) [Unknown/Just-In-Time compiled code] (Unknown Source:0)

## write

#### writeback 与readahead区别

Write函数通过调用系统调用接口,将数据从应用层copy到内核层,所以write会触发内核态/用户态切换。当数据到达page cache后,内核并不会立即把数据往下传递。而是返回用户空间。数据什么时候写入硬盘,有内核IO调度决定,所以write是一个异步调用。这一点和read不同,read调用是先检查page cache里面是否有数据,如果有,就取出来返回用户,如果没有,就同步传递下去并等待有数据,再返回用户,所以read是一个同步过程。当然你也可以把write的异步过程改成同步过程,就是在open文件的时候带上O\_SYNC标记。

#### O\_DIRECT & RAW flag (权重方式的io控制器只能作用于direct io)

O\_DIRECT 和 RAW设备最根本的区别是O\_DIRECT是基于文件系统的,也就是在应用层来看,其操作对象是文件句柄,内核和文件层来看,其操作是基于inode和数据块,这些概念都是和ext2/3的文件系统相关,写到磁盘上最终是ext3文件。

而RAW设备写是没有文件系统概念,操作的是扇区号,操作对象是扇区,写出来的东西不一定是ext3文件(如果按照ext3规则写就是ext3文件)。

一般基于O\_DIRECT来设计优化自己的文件模块,是不满系统的cache和调度策略,自己在应用层实现这些,来制定自己特有的业务特色文件读写。但是写出来的东西是ext3文件,该磁盘卸下来,mount到其他任何linux系统上,都可以查看。

而基于RAW设备的设计系统,一般是不满现有ext3的诸多缺陷,设计自己的文件系统。自己设计文件布局和索引方式。举个极端例子:把整个磁盘做一个文件来写,不要索引。这样没有inode限制,没有文件大小限制,磁盘有多大,文件就能多大。这样的磁盘卸下来,mount到其他linux系统上,是无法识别其数据的。

两者都要通过驱动层读写;在系统引导启动,还处于实模式的时候,可以通过bios接口读写raw设备。

#### 写操作层次关系

#### 具体文件系统层

当具体文件系统层(像ext2/3/4等,我称之为具体文件系统)接到写IO请求时,会判断该IO是否具有DIRECT1

static ssize\_t ext4\_file\_write(struct kiocb \*iocb,const struct iovec \*iov, unsignedlong nr\_segs, loff\_t pos) { structinode \*inode = file\_inode(iocb->ki\_filp); ssize\_tret; ... if(unlikely(iocb->ki\_filp->f\_flags & O\_DIRECT)) ret =ext4\_file\_dio\_write(iocb, iov, nr\_segs, pos); else ret =generic\_file\_aio\_write(iocb, iov, nr\_segs, pos);

returnret; }

#### 直接IO:

控制流,若进入直接IO,则调用具体文件系统层的直接IO处理函数,然后调用通过submit\_bio函数将IO提交至数据流,数据依旧存放在用户态缓存中,并不需要将数据复制到pagecache中,减少了数据复制次数。

#### 缓存IO:

控制流,若进入BufferIO,则调用具体文件系统的write\_begin进入的准备:比如空间分配,缓存映射(涉及是数据流,数据从用户态复制到内核态page cache中。

#### Page Cache层

Page Cache是文件数据在内存中的副本,因此Page Cache管理与内存管理系统和文件系统都相关:一方面Pa

#### 通用块层

通用块层:由于绝大多数情况的IO操作是跟块设备打交道,所以Linux在此提供了一个类似vfs层的块设备操作是 无论是DirectIO还是BufferIO,最后都会通过submit\_bio()将IO请求提交到通用块层,通过generic\_mal submit\_bio函数通过generic\_make\_request转发bio,generic\_make\_request是一个循环,其通过每个 Generic\_make\_request的执行上下文可能有两种,一种是用户上下文,另一种为pdflush所在的内核线程上 在通用块层,提供了一个通用的请求队列压栈方法:blk\_queue\_bio。在初始化一个有queue块设备驱动的时间

#### IO调度层

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IO调度层:因为绝大多数的块设备都是类似磁盘这样的设备,所以有必要根据这类设备的特点以及应用的不同特到目前为止,文件系统(pdflush或者address\_space\_operations)发下来的bio已经merge到request如果为sync bio,那么直接调用\_\_\_generic\_unplug\_device,否则需要在unplug\_timer的软中断上下文F

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queue是块设备的驱动程序提供的一个请求队列。make\_request\_fn函数将bio放入请求队列中进行调度处理。 举例中的sda是一个scsi设备,在scsi middle level将scsi\_request\_fn函数注册到了queue队列的req

#### 设备驱动层

设备驱动程序要做的事情就是从request\_queue里面取出请求,然后操作硬件设备,逐个去执行这些请求。除了

#### 以sisc设备为例:

接下来的过程实际上和具体的scsi总线操作相关了。在scsi\_request\_fn函数中会扫描request队列,通过elv\_next\_request函数从队列中获取一个request。在elv\_next\_request函数中通过scsi总线层注册的q->prep\_rq\_fn(scsi层注册为scsi\_prep\_fn)函数将具体的request转换成scsi驱动所能认识的scsi command。获取一个request之后,scsi\_request\_fn函数直接调用scsi\_dispatch\_cmd函数将scsi command发送给一个具体的scsi host。到这一步,有一个问题:scsi command具体转发给那个scsi host呢?秘密就在于q->queuedata中,在为sda设备分配queue队列时,已经指定了sda块设备与底层的scsi设备(scsidevice)之间的关系,他们的关系是通过request queue维护的。

在scsi\_dispatch\_cmd函数中,通过scsi host的接口方法queuecommand将scsi command发送给scsi l 在SCSi中断下半部中,调用scsi command结束的回调函数,这个函数往往为scsi\_done,在scsi\_done函数。经设备驱动层,将数据复制到disk cache中。

#### f2fs文件系统write操作到generic\_make\_request的调用栈

generic\_make\_request(struct bio \* bio) (/home/panard/linux-4.19.100/block/blk-core.c:2 submit\_bio(struct bio \* bio) (/home/panard/linux-4.19.100/block/blk-core.c:2573) \_\_submit\_bio() (/home/panard/linux-4.19.100/fs/f2fs/data.c:307) \_\_submit\_merged\_bio(struct f2fs\_bio\_info \* io) (/home/panard/linux-4.19.100/fs/f2fs/da f2fs\_submit\_page\_write(struct f2fs\_io\_info \* fio) (/home/panard/linux-4.19.100/fs/f2fs do\_write\_page(struct f2fs\_summary \* sum, struct f2fs\_io\_info \* fio) (/home/panard/linu f2fs\_outplace\_write\_data(struct dnode\_of\_data \* dn, struct f2fs\_io\_info \* fio) (/home/ f2fs\_do\_write\_data\_page(struct f2fs\_io\_info \* fio) (/home/panard/linux-4.19.100/fs/f2f \_\_write\_data\_page(struct page \* page, bool \* submitted, struct writeback\_control \* wbc f2fs\_write\_cache\_pages(struct address\_space \* mapping, struct writeback\_control \* wbc, \_f2fs\_write\_data\_pages() (/home/panard/linux-4.19.100/fs/f2fs/data.c:2217) f2fs\_write\_data\_pages(struct address\_space \* mapping, struct writeback\_control \* wbc) do\_writepages(struct address\_space \* mapping, struct writeback\_control \* wbc) (/home/p \_writeback\_single\_inode(struct inode \* inode, struct writeback\_control \* wbc) (/home/ writeback\_sb\_inodes(struct super\_block \* sb, struct bdi\_writeback \* wb, struct wb\_writ \_\_writeback\_inodes\_wb(struct bdi\_writeback \* wb, struct wb\_writeback\_work \* work) (/ho wb\_writeback(struct bdi\_writeback \* wb, struct wb\_writeback\_work \* work) (/home/panard

```
wb_check_start_all() (/home/panard/linux-4.19.100/fs/fs-writeback.c:1936)
wb_do_writeback() (/home/panard/linux-4.19.100/fs/fs-writeback.c:1962)
wb_workfn(struct work_struct * work) (/home/panard/linux-4.19.100/fs/fs-writeback.c:19
process_one_work(struct worker * worker, struct work_struct * work) (/home/panard/linu
worker_thread(void * __worker) (/home/panard/linux-4.19.100/kernel/workqueue.c:2296)
kthread(void * _create) (/home/panard/linux-4.19.100/kernel/kthread.c:246)
ret_from_fork() (/home/panard/linux-4.19.100/arch/x86/entry/entry_64.S:415)
[Unknown/Just-In-Time compiled code] (Unknown Source:0)
```

# f2fs write cgroup weight/f2fs文件系统写操作权重方式io控制器调用栈

```
\verb|cfq_group_served| (struct cfq_data * cfqd, struct cfq_group * cfqg, struct cfq_queue * cfq_group_served| (struct cfq_data * cfqd, struct cfq_group_served) | (struct cfq_g
__cfq_slice_expired(struct cfq_data * cfqd, struct cfq_queue * cfqq, bool timed_out) (
cfq_slice_expired() (/home/panard/linux-4.19.100/block/cfq-iosched.c:2756)
cfq_select_queue() (/home/panard/linux-4.19.100/block/cfq-iosched.c:3388)
cfq_dispatch_requests(struct request_queue * q, int force) (/home/panard/linux-4.19.10
elv_next_request() (/home/panard/linux-4.19.100/block/blk-core.c:2857)
blk_peek_request(struct request_queue * q) (/home/panard/linux-4.19.100/block/blk-core
scsi_request_fn(struct request_queue * q) (/home/panard/linux-4.19.100/drivers/scsi/sc
  _blk_run_queue_uncond() (/home/panard/linux-4.19.100/block/blk-core.c:471)
__blk_run_queue(struct request_queue * q) (/home/panard/linux-4.19.100/block/blk-core.
queue_unplugged(struct request_queue * q, unsigned int depth, bool from_schedule) (/ho
blk_flush_plug_list(struct blk_plug * plug, bool from_schedule) (/home/panard/linux-4.
blk_queue_bio(struct request_queue * q, struct bio * bio) (/home/panard/linux-4.19.100
generic_make_request(struct bio * bio) (/home/panard/linux-4.19.100/block/blk-core.c:2
submit_bio(struct bio * bio) (/home/panard/linux-4.19.100/block/blk-core.c:2573)
__submit_bio() (/home/panard/linux-4.19.100/fs/f2fs/data.c:307)
__submit_merged_bio(struct f2fs_bio_info * io) (/home/panard/linux-4.19.100/fs/f2fs/da
f2fs_submit_page_write(struct f2fs_io_info * fio) (/home/panard/linux-4.19.100/fs/f2fs
do_write_page(struct f2fs_summary * sum, struct f2fs_io_info * fio) (/home/panard/linu
f2fs_outplace_write_data(struct dnode_of_data * dn, struct f2fs_io_info * fio) (/home/
f2fs_do_write_data_page(struct f2fs_io_info * fio) (/home/panard/linux-4.19.100/fs/f2f
__write_data_page(struct page * page, bool * submitted, struct writeback_control * wbc
f2fs_write_cache_pages(struct address_space * mapping, struct writeback_control * wbc,
__f2fs_write_data_pages() (/home/panard/linux-4.19.100/fs/f2fs/data.c:2217)
f2fs_write_data_pages(struct address_space * mapping, struct writeback_control * wbc)
do_writepages(struct address_space * mapping, struct writeback_control * wbc) (/home/p
  _writeback_single_inode(struct inode * inode, struct writeback_control * wbc) (/home/
writeback_sb_inodes(struct super_block * sb, struct bdi_writeback * wb, struct wb_writ
__writeback_inodes_wb(struct bdi_writeback * wb, struct wb_writeback_work * work) (/ho
wb_writeback(struct bdi_writeback * wb, struct wb_writeback_work * work) (/home/panard
wb_check_background_flush() (/home/panard/linux-4.19.100/fs/fs-writeback.c:1880)
wb_do_writeback() (/home/panard/linux-4.19.100/fs/fs-writeback.c:1968)
wb_workfn(struct work_struct * work) (/home/panard/linux-4.19.100/fs/fs-writeback.c:19
process_one_work(struct worker * worker, struct work_struct * work) (/home/panard/linu
worker_thread(void * __worker) (/home/panard/linux-4.19.100/kernel/workqueue.c:2296)
kthread(void * _create) (/home/panard/linux-4.19.100/kernel/kthread.c:246)
```

```
ret_from_fork() (/home/panard/linux-4.19.100/arch/x86/entry/entry_64.S:415)
[Unknown/Just-In-Time compiled code] (Unknown Source:0)
```

```
balance_dirty_pages_ratelimited
        f2fs_update_time
        f2fs_put_page
        set_page_dirty
    f2fs_write_end
    flush_dcache_page
    iov_iter_copy_from_user_atomic
        f2fs_wait_on_page_writeback
            __up_read
        __do_map_lock
        if f2fs_has_inline_data return 0;
        prepare_write_begin //we already allocated all the blocks, so we don't need to
            find_get_entry
        pagecache_get_page(struct address_space * mapping, unsigned long offset, int f
        f2fs_pagecache_get_page() (/home/panard/linux-4.19.100/fs/f2fs/f2fs.h:2044)
    f2fs_write_begin(struct file * file, struct address_space * mapping, loff_t pos, u
generic_perform_write(struct file * file, struct iov_iter * i, loff_t pos) (/home/pana
 _generic_file_write_iter(struct kiocb * iocb, struct iov_iter * from) (/home/panard/l
f2fs_file_write_iter(struct kiocb * iocb, struct iov_iter * from) (/home/panard/linux-
call_write_iter() (/home/panard/linux-4.19.100/include/linux/fs.h:1820)
new_sync_write() (/home/panard/linux-4.19.100/fs/read_write.c:474)
__vfs_write(struct file * file, const char * p, size_t count, loff_t * pos) (/home/pan
vfs_write(struct file * file, const char * buf, size_t count, loff_t * pos) (/home/pan
ksys_write(unsigned int fd, const char * buf, size_t count) (/home/panard/linux-4.19.1
__do_sys_write() (/home/panard/linux-4.19.100/fs/read_write.c:611)
__se_sys_write() (/home/panard/linux-4.19.100/fs/read_write.c:608)
__x64_sys_write(const struct pt_regs * regs) (/home/panard/linux-4.19.100/fs/read_writ
do_syscall_64(unsigned long nr, struct pt_regs * regs) (/home/panard/linux-4.19.100/ar
entry_SYSCALL_64() (/home/panard/linux-4.19.100/arch/x86/entry/entry_64.S:238)
```

#### 伪代码,结构体&函数栈

```
struct address_space {
                                       /* owner: inode, block_device */若为空, swap are
   struct inode
                           *host;
                                       /* cached pages */
   struct radix_tree_root i_pages;
   struct rb_root_cached
                           i_mmap;
                                       /* tree of private and shared mappings */
   /* Protected by the
                           i_pages lock */
                                       /* number of total pages */
   unsigned long
                           nrpages;
   pgoff t
                           writeback_index;/* writeback starts here */
                                                   /* methods */
   const struct address_space_operations *a_ops;
   /*指向操作函数表 (struct address_space_operations) ,每个后备存储都要实现这个函数表,比如exi
   unsigned long
                                       /* error bits */
                           flags;
                                          /* for use by the address_space */
   spinlock_t
                           private_lock;
                                       /* implicit gfp mask for allocations */
   gfp_t
                           gfp_mask;
   struct list_head
                           private_list;
                                          /* for use by the address_space */
```

```
*private_data; /* ditto */
   void
    errseq_t
                          wb_err;
};
const struct address_space_operations f2fs_dblock_aops = {
              = f2fs_read_data_page,
    /* readpage()首先会调用find_get_page(mapping, index)在page cache中寻找请求的数据,mappi
    .writepage = f2fs_write_data_page,
    /* 对于文件映射 (host指向一个inode对象) , page每次修改后都会调用SetPageDirty (page) 将page标
    .write_begin
                = f2fs_write_begin,
    .write_end = f2fs_write_end,
    .set_page_dirty = f2fs_set_data_page_dirty,
};
f2fs_read_inline_data : ... ? f2fs_mpage_readpages
f2fs_write_data_page
    --> __write_data_page
        --> f2fs_has_inline_data
           no inline data:
           --> f2fs_do_write_data_page ???
               --> encrypt_one_page
               --> set_page_writeback
                   -->test_set_page_writeback (page-writeback.c)
                       --> lock_page_memcg(page)
                       --> TestSetPageWriteback
               --> f2fs_inplace_write_data
                   --> f2fs_submit_page_bio
                   --> update_device_state
                   --> f2fs_update_iostat
               --> inode_dec_dirty_pages
               --> f2fs_submit_merged_write_cond
               --> clear_inode_flag
               --> f2fs_remove_dirty_inode
           has inline data:
           --> f2fs_write_inline_data
               --> set_new_dnode
               --> f2fs_wait_on_page_writeback
               --> set_page_dirty
/*当一个block被读入内存或者等待写入块设备时,保存在buffer中,一个buffer对应一个block*/
struct buffer_head {
    unsigned long b_state;
                             /* buffer state bitmap (see above) 如是否dirty,是否正在
    struct buffer_head *b_this_page;/* circular list of page's buffers */
                             /* the page this bh is mapped to 指向buffer所在的page (判
    struct page *b_page;
                         /* start block number */
    sector_t b_blocknr;
    size_t b_size;
                         /* size of mapping 表示block的大小。该block起始于b_data,终止于
                          /* pointer to data within the page 直接指向block所在位置 (在b_
    char *b_data;
    struct block_device *b_bdev; /* 指向buffer对应的block所在的块设备对象 */
    void *b_private;
                         /* reserved for b_end_io */
    struct list_head b_assoc_buffers; /* associated with another mapping */
    struct address_space *b_assoc_map; /* mapping this buffer is associated with */
```

```
atomic_t b_count; /* users using this buffer_head buffer的使用计数,调用get_bh(
};
/* 表示一个正在进行的块 I/0操作 */
struct bio {
    sector_t bi_sector; /* associated sector on disk */
    struct bio *bi_next; /* list of requests */
    struct block_device *bi_bdev; /* associated block device */
    unsigned long bi_flags; /* status and command flags */
    unsigned long bi_rw; /* read or write? */
    unsigned short bi_vcnt; /* number of bio_vecs off */ /* 是bi_io_vec数组的长度 */
    unsigned short bi_idx; /* current index in bi_io_vec */ /* 当前正在进行I/0操作bio_vec
    unsigned short bi_phys_segments; /* number of segments */
    unsigned int bi_size; /* I/O count */
    unsigned int bi_seg_front_size; /* size of first segment */
    unsigned int bi_seg_back_size; /* size of last segment */
    unsigned int bi_max_vecs; /* maximum bio_vecs possible */
    unsigned int bi_comp_cpu; /* completion CPU */
    atomic_t bi_cnt; /* usage counter */
    struct bio_vec *bi_io_vec; /* bio_vec list */ /* bi_io_vec是一个指针,指向一个bio_vec
    /* bi_io_vec数组使得bio结构体可以支持在一次I/O操作中,使用多个在内存中不连续的segment,这个又[
    bio_end_io_t *bi_end_io; /* I/O completion method */
    void *bi_private; /* owner-private method */
    bio_destructor_t *bi_destructor; /* destructor method */
    struct bio_vec bi_inline_vecs[0]; /* inline bio vectors */
};
struct bio_vec {
    struct page *bv_page; //向segment所在的物理page
    unsigned int bv_len; //segment的大小 (字节)
    unsigned int bv_offset;//segment起始点在page中的偏移量。
};
```

# latency方式io控制器调用栈

#### author comments

```
Block rq-qos base io controller

This works similar to wbt with a few exceptions
```

- It's bio based, so the latency covers the whole block layer in addition to the actua -它是基于bio的,所以除了实际的io,latency覆盖了整个block层。
- We will throttle all IO that comes in here if we need to.
  -如果需要的话,我们会限制所有进入这里的IO。
- We use the mean latency over the 100ms window. This is because writes can be partic -我们使用100毫秒窗口的平均延迟。这是因为写入速度特别快,这可能会让我们错误地感觉到其他工作负载对受
- By default there's no throttling, we set the queue\_depth to INT\_MAX so that we can h -默认情况下没有限制,我们将队列深度设置为INT\_MAX,这样我们就可以拥有尽可能多的未完成的bio。只有在

The hierarchy works like the cpu controller does, we track the latency at every config 层次结构的工作方式与cpu控制器类似,我们跟踪每个配置节点的延迟,每个配置节点都有自己独立的队列深度。这意

Consider the following

"a" and "b" have no target, but their combined io under "fast" cannot exceed an averag "a"和"b"没有目标,但它们在"fast"下的组合io不能超过平均5毫秒的延迟,如果超过,我们将限制"slow"组。对

In this example "fast", "slow", and "normal" will be the only groups actually accounti 在本例中, "fast"、"slow"和"normal"将是唯一实际计算io延迟的组。每次提交时,我们都必须从继承人那里走到

There are 2 ways we throttle IO.

- 1) Queue depth throttling. As we throttle down we will adjust the maximum number of I
- 1) 队列深度限制。当我们减速时,我们将调整允许在飞行中使用的IO的最大数量。从(u64)-1到1。如果组只为自
- 2) Induced delay throttling. This is for the case that a group is generating IO that
- 2) 诱导延迟节流。这是针对一个组正在生成IO的情况,该IO必须由根cg发出以避免优先级反转。所以想想REQ\_ME<sup>-</sup>

```
total_time += min_lat_nsec - actual_io_completion
and then at throttle time will do
throttle_time = min(total_time, NSEC_PER_SEC)
```

This induced delay will throttle back the activity that is generating the root cg issu 这种诱导的延迟将抑制生成根cg发出的io的活动,不管是一些元数据密集型操作还是组使用了太多内存,以至于将我们

#### latency ops:

https://md2pdf.netlify.app/

```
static struct rq_qos_ops blkcg_iolatency_ops = {
    .throttle = blkcg_iolatency_throttle,
    .done_bio = blkcg_iolatency_done_bio,
    .exit = blkcg_iolatency_exit,
};
```

#### blk\_iolatency\_init (PS: add loop device)

```
timer_setup (blkiolatency_timer_fn)
blkcg_activate_policy (blkcg_policy_iolatency)
rq_qos_add (blkcg_iolatency_ops)
blk_iolatency_init(struct request_queue * q) (/root/linux-4.19.100/block/blk-iolatency
blkcg_init_queue(struct request_queue * q) (/root/linux-4.19.100/block/blk-cgroup.c:12
```

12/30

```
blk_alloc_queue_node(gfp_t gfp_mask, int node_id, spinlock_t * lock) (/root/linux-4.19
blk_mq_init_queue(struct blk_mq_tag_set * set) (/root/linux-4.19.100/block/blk-mq.c:24
loop_add(struct loop_device ** l, int i) (/root/linux-4.19.100/drivers/block/loop.c:19
loop_init() (/root/linux-4.19.100/drivers/block/loop.c:2236)
do_one_initcall(initcall_t fn) (/root/linux-4.19.100/init/main.c:883)
do_initcall_level() (/root/linux-4.19.100/init/main.c:951)
do_initcalls() (/root/linux-4.19.100/init/main.c:959)
do_basic_setup() (/root/linux-4.19.100/init/main.c:977)
kernel_init_freeable() (/root/linux-4.19.100/init/main.c:1144)
kernel_init(void * unused) (/root/linux-4.19.100/init/main.c:1061)
ret_from_fork() (/root/linux-4.19.100/arch/x86/entry/entry_64.S:415)
[Unknown/Just-In-Time compiled code] (Unknown Source:0)
```

#### blk\_iolatency\_init (PS: scsi\_add\_device)

```
timer_setup (blkiolatency_timer_fn)
    blkcg_activate_policy (blkcg_policy_iolatency)
    rq_qos_add (blkcg_iolatency_ops)
blk_iolatency_init(struct request_queue * q) (/root/linux-4.19.100/block/blk-iolatency
blkcg_init_queue(struct request_queue * q) (/root/linux-4.19.100/block/blk-cgroup.c:12
blk_alloc_queue_node(gfp_t gfp_mask, int node_id, spinlock_t * lock) (/root/linux-4.19
scsi_old_alloc_queue(struct scsi_device * sdev) (/root/linux-4.19.100/drivers/scsi/scs
scsi_alloc_sdev(struct scsi_target * starget, u64 lun, void * hostdata) (/root/linux-4
scsi_probe_and_add_lun(struct scsi_target * starget, u64 lun, blist_flags_t * bflagsp,
__scsi_add_device(struct Scsi_Host * shost, uint channel, uint id, u64 lun, void * hos
ata_scsi_scan_host(struct ata_port * ap, int sync) (/root/linux-4.19.100/drivers/ata/l
async_port_probe(void * data, async_cookie_t cookie) (/root/linux-4.19.100/drivers/ata
async_run_entry_fn(struct work_struct * work) (/root/linux-4.19.100/kernel/async.c:127
process_one_work(struct worker * worker, struct work_struct * work) (/root/linux-4.19.
worker_thread(void * __worker) (/root/linux-4.19.100/kernel/workqueue.c:2296)
kthread(void * _create) (/root/linux-4.19.100/kernel/kthread.c:246)
ret_from_fork() (/root/linux-4.19.100/arch/x86/entry/entry_64.S:415)
[Unknown/Just-In-Time compiled code] (Unknown Source:0)
```

#### iolatency\_set\_min\_lat\_nsec 设置latency target

```
iolat->min_lat_nsec = val; //target*1000 (ns)
    iolat->cur_win_nsec = max_t(u64, val << 4, BLKIOLATENCY_MIN_WIN_SIZE) /* 100ms <--
    iolat->cur_win_nsec
iolatency_set_min_lat_nsec(struct blkcg_gq * blkg, u64 val) (/root/linux-4.19.100/bloc
iolatency_set_limit(struct kernfs_open_file * of, char * buf, size_t nbytes, loff_t of
cgroup_file_write(struct kernfs_open_file * of, char * buf, size_t nbytes, loff_t off)
kernfs_fop_write(struct file * file, const char * user_buf, size_t count, loff_t * ppo
    _vfs_write(struct file * file, const char * p, size_t count, loff_t * pos) (/root/lin
vfs_write(struct file * file, const char * buf, size_t count, loff_t * pos) (/root/lin
ksys_write(unsigned int fd, const char * buf, size_t count) (/root/linux-4.19.100/fs/r
    _do_sys_write() (/root/linux-4.19.100/fs/read_write.c:611)
    _se_sys_write() (/root/linux-4.19.100/fs/read_write.c:608)
    _x64_sys_write(const struct pt_regs * regs) (/root/linux-4.19.100/fs/read_write.c:608)
```

do\_syscall\_64(unsigned long nr, struct pt\_regs \* regs) (/root/linux-4.19.100/arch/x86/ entry\_SYSCALL\_64() (/root/linux-4.19.100/arch/x86/entry/entry\_64.S:238)

#### rq\_qos\_throttle 调用栈

```
__blkcg_iolatency_throttle
    bio_issue_init
blkcg_iolatency_throttle
rq_qos_throttle(struct request_queue * q, struct bio * bio, spinlock_t * lock) (/root/
blk_queue_bio(struct request_queue * q, struct bio * bio) (/root/linux-4.19.100/block/
generic_make_request(struct bio * bio) (/root/linux-4.19.100/block/blk-core.c:2471)
submit_bio(struct bio * bio) (/root/linux-4.19.100/block/blk-core.c:2580)
 _submit_bio() (/root/linux-4.19.100/fs/f2fs/data.c:307)
f2fs_mpage_readpages(struct address_space * mapping, struct list_head * pages, struct
f2fs_read_data_pages(struct file * file, struct address_space * mapping, struct list_h
read_pages(struct address_space * mapping, struct file * filp, struct list_head * page
__do_page_cache_readahead(struct address_space * mapping, struct file * filp, unsigned
ra_submit() (/root/linux-4.19.100/mm/internal.h:66)
do_sync_mmap_readahead() (/root/linux-4.19.100/mm/filemap.c:2467)
filemap_fault(struct vm_fault * vmf) (/root/linux-4.19.100/mm/filemap.c:2543)
f2fs_filemap_fault(struct vm_fault * vmf) (/root/linux-4.19.100/fs/f2fs/file.c:42)
__do_fault(struct vm_fault * vmf) (/root/linux-4.19.100/mm/memory.c:3269)
do_read_fault() (/root/linux-4.19.100/mm/memory.c:3681)
do_fault() (/root/linux-4.19.100/mm/memory.c:3810)
handle_pte_fault() (/root/linux-4.19.100/mm/memory.c:4041)
__handle_mm_fault(struct vm_area_struct * vma, unsigned long address, unsigned int fla
handle_mm_fault(struct vm_area_struct * vma, unsigned long address, unsigned int flags
__do_page_fault(struct pt_regs * regs, unsigned long error_code, unsigned long address
do_page_fault(struct pt_regs * regs, unsigned long error_code) (/root/linux-4.19.100/a
page_fault() (/root/linux-4.19.100/arch/x86/entry/entry_64.S:1204)
[Unknown/Just-In-Time compiled code] (Unknown Source:0)
irq_stack_union (Unknown Source:0)
[Unknown/Just-In-Time compiled code] (Unknown Source:0)
```

#### blkcg\_iolatency\_done\_bio 调用栈

```
iolatency_record_time
blkcg_iolatency_done_bio(struct rq_qos * rqos, struct bio * bio) (/root/linux-4.19.100
rq_qos_done_bio(struct request_queue * q, struct bio * bio) (/root/linux-4.19.100/bloc
bio_endio(struct bio * bio) (/root/linux-4.19.100/block/bio.c:1755)
req_bio_endio() (/root/linux-4.19.100/block/blk-core.c:285)
blk_update_request(struct request * req, blk_status_t error, unsigned int nr_bytes) (/
scsi_end_request(struct request * req, blk_status_t error, unsigned int bytes, unsigne
scsi_io_completion(struct scsi_cmnd * cmd, unsigned int good_bytes) (/root/linux-4.19.
scsi_finish_command(struct scsi_cmnd * cmd) (/root/linux-4.19.100/drivers/scsi/scsi.c:
scsi_softirq_done(struct request * rq) (/root/linux-4.19.100/drivers/scsi/scsi_lib.c:1
blk_done_softirq(struct softirq_action * h) (/root/linux-4.19.100/block/blk-softirq.c:
__do_softirq() (/root/linux-4.19.100/kernel/softirq.c:292)
```

```
invoke_softirq() (/root/linux-4.19.100/kernel/softirq.c:372)
irq_exit() (/root/linux-4.19.100/kernel/softirq.c:412)
exiting_irq() (/root/linux-4.19.100/arch/x86/include/asm/apic.h:536)
do_IRQ(struct pt_regs * regs) (/root/linux-4.19.100/arch/x86/kernel/irq.c:258)
common_interrupt() (/root/linux-4.19.100/arch/x86/entry/entry_64.S:670)
init_thread_union (Unknown Source:0)
[Unknown/Just-In-Time compiled code] (Unknown Source:0)
```

#### blk\_clear\_congested 调用栈

```
blk_clear_congested(struct request_list * rl, int sync) (/root/linux-4.19.100/block/bl
__freed_request(struct request_list * rl, int sync) (/root/linux-4.19.100/block/blk-co
freed\_request(struct\ request\_list\ ^*\ rl,\ bool\ sync,\ req\_flags\_t\ rq\_flags)\ (/root/linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux-linux
 __blk_put_request(struct request_queue * q, struct request * req) (/root/linux-4.19.10
blk_finish_request(struct request * req, blk_status_t error) (/root/linux-4.19.100/blo
scsi_end_request(struct request * req, blk_status_t error, unsigned int bytes, unsigne
scsi_io_completion(struct scsi_cmnd * cmd, unsigned int good_bytes) (/root/linux-4.19.
scsi_finish_command(struct scsi_cmnd * cmd) (/root/linux-4.19.100/drivers/scsi/scsi.c:
scsi_softirq_done(struct request * rq) (/root/linux-4.19.100/drivers/scsi/scsi_lib.c:1
blk_done_softirq(struct softirq_action * h) (/root/linux-4.19.100/block/blk-softirq.c:
   _do_softirq() (/root/linux-4.19.100/kernel/softirq.c:292)
invoke_softirq() (/root/linux-4.19.100/kernel/softirq.c:372)
irg_exit() (/root/linux-4.19.100/kernel/softirg.c:412)
exiting_irq() (/root/linux-4.19.100/arch/x86/include/asm/apic.h:536)
do_IRQ(struct pt_regs * regs) (/root/linux-4.19.100/arch/x86/kernel/irq.c:258)
common_interrupt() (/root/linux-4.19.100/arch/x86/entry/entry_64.S:670)
[Unknown/Just-In-Time compiled code] (Unknown Source:0)
```

#### \_\_mark\_inode\_dirty 调用栈

```
_mark_inode_dirty(struct inode * inode, int flags) (/root/linux-4.19.100/fs/fs-writeb
mark_inode_dirty_sync() (/root/linux-4.19.100/include/linux/fs.h:2087)
dquot_free_space() (/root/linux-4.19.100/include/linux/quotaops.h:383)
dquot_free_block() (/root/linux-4.19.100/include/linux/quotaops.h:393)
f2fs_i_blocks_write() (/root/linux-4.19.100/fs/f2fs/f2fs.h:2390)
dec_valid_block_count() (/root/linux-4.19.100/fs/f2fs/f2fs.h:1773)
f2fs_truncate_data_blocks_range(struct dnode_of_data * dn, int count) (/root/linux-4.1
f2fs_truncate_data_blocks(struct dnode_of_data * dn) (/root/linux-4.19.100/fs/f2fs/fil
truncate_dnode(struct dnode_of_data * dn) (/root/linux-4.19.100/fs/f2fs/node.c:880)
truncate_nodes(struct dnode_of_data * dn, unsigned int nofs, int ofs, int depth) (/roo
f2fs_truncate_inode_blocks(struct inode * inode, unsigned long from) (/root/linux-4.19
f2fs_truncate_blocks(struct inode * inode, u64 from, bool lock) (/root/linux-4.19.100/
f2fs_truncate(struct inode * inode) (/root/linux-4.19.100/fs/f2fs/file.c:685)
f2fs_setattr(struct dentry * dentry, struct iattr * attr) (/root/linux-4.19.100/fs/f2f
notify_change(struct dentry * dentry, struct iattr * attr, struct inode ** delegated_i
do_truncate(struct dentry * dentry, loff_t length, unsigned int time_attrs, struct fil
handle_truncate() (/root/linux-4.19.100/fs/namei.c:3009)
do_last() (/root/linux-4.19.100/fs/namei.c:3427)
```

```
path_openat(struct nameidata * nd, const struct open_flags * op, unsigned int flags) (
do_filp_open(int dfd, struct filename * pathname, const struct open_flags * op) (/root
do_sys_open(int dfd, const char * filename, int flags, umode_t mode) (/root/linux-4.19
__do_sys_openat(int flags) (/root/linux-4.19.100/fs/open.c:1115)
__se_sys_openat() (/root/linux-4.19.100/fs/open.c:1109)
__x64_sys_openat(const struct pt_regs * regs) (/root/linux-4.19.100/fs/open.c:1109)
do_syscall_64(unsigned long nr, struct pt_regs * regs) (/root/linux-4.19.100/arch/x86/entry_SYSCALL_64() (/root/linux-4.19.100/arch/x86/entry/entry_64.S:238)
[Unknown/Just-In-Time compiled code] (Unknown Source:0)
```

# blkcg\_iolatency\_throttle 调用栈

```
\verb|blkcg_iolatency_throttle| (struct rq_qos * rqos, struct bio * bio, spinlock_t * lock) (/ lock_t * 
rq_qos_throttle(struct request_queue * q, struct bio * bio, spinlock_t * lock) (/root/
blk_queue_bio(struct request_queue * q, struct bio * bio) (/root/linux-4.19.100/block/
generic_make_request(struct bio * bio) (/root/linux-4.19.100/block/blk-core.c:2471)
submit_bio(struct bio * bio) (/root/linux-4.19.100/block/blk-core.c:2580)
__submit_bio() (/root/linux-4.19.100/fs/f2fs/data.c:307)
__submit_merged_bio(struct f2fs_bio_info * io) (/root/linux-4.19.100/fs/f2fs/data.c:32
f2fs_submit_page_write(struct f2fs_io_info * fio) (/root/linux-4.19.100/fs/f2fs/data.c
do_write_page(struct f2fs_summary * sum, struct f2fs_io_info * fio) (/root/linux-4.19.
f2fs_outplace_write_data(struct dnode_of_data * dn, struct f2fs_io_info * fio) (/root/
f2fs_do_write_data_page(struct f2fs_io_info * fio) (/root/linux-4.19.100/fs/f2fs/data.
   _write_data_page(struct page * page, <mark>bool</mark> * submitted, struct writeback_control * wbc
f2fs_write_cache_pages(struct address_space * mapping, struct writeback_control * wbc,
__f2fs_write_data_pages() (/root/linux-4.19.100/fs/f2fs/data.c:2217)
f2fs_write_data_pages(struct address_space * mapping, struct writeback_control * wbc)
do_writepages(struct address_space * mapping, struct writeback_control * wbc) (/root/l
  _writeback_single_inode(struct inode * inode, struct writeback_control * wbc) (/root/
writeback_sb_inodes(struct super_block * sb, struct bdi_writeback * wb, struct wb_writ
  _writeback_inodes_wb(struct bdi_writeback * wb, struct wb_writeback_work * work) (/ro
wb_writeback(struct bdi_writeback * wb, struct wb_writeback_work * work) (/root/linux-
wb_check_start_all() (/root/linux-4.19.100/fs/fs-writeback.c:1954)
wb_do_writeback() (/root/linux-4.19.100/fs/fs-writeback.c:1981)
wb_workfn(struct work_struct * work) (/root/linux-4.19.100/fs/fs-writeback.c:2015)
process_one_work(struct worker * worker, struct work_struct * work) (/root/linux-4.19.
worker_thread(void * __worker) (/root/linux-4.19.100/kernel/workqueue.c:2296)
kthread(void * _create) (/root/linux-4.19.100/kernel/kthread.c:246)
ret_from_fork() (/root/linux-4.19.100/arch/x86/entry/entry_64.S:415)
[Unknown/Just-In-Time compiled code] (Unknown Source:0)
```

#### blkcg\_maybe\_throttle\_blkg 调用栈

```
exit_to_usermode_loop() {
    mem_cgroup_handle_over_high();
    blkcg_maybe_throttle_current() {
        /* blkcg_maybe_throttle_current[1750] --- */
        kthread_blkcg();
```

```
blkg_lookup_slowpath();
ktime_get();
/* blkcg_maybe_throttle_blkg[1689] --- */
blkcg_scale_delay();
blk_put_queue();
}

// TODO: 测试是否会调用到blkcg_maybe_throttle_blkg, 已经添加printk, 检查是否有delay_nsec输出
```

#### qos 调用栈 (多核cpu虚拟机trace文件截取)

PS: 多核cpu与单核cpu在执行同样的cgroup io动作时,流程似有差异,低优先级研究

```
__brelse();
 _getblk_gfp() {
 __find_get_block() {
    pagecache_get_page() {
      find_get_entry();
    }
  }
 _cond_resched() {
    rcu_all_qs();
 __find_get_block() {
   pagecache_get_page() {
      find_get_entry();
    }
  }
  pagecache_get_page() {
    find_get_entry();
    __page_cache_alloc() {
      alloc_pages_current() {
        get_task_policy();
        policy_nodemask();
        policy_node();
        __alloc_pages_nodemask() {
          _cond_resched() {
            rcu_all_qs();
          }
          get_page_from_freelist() {
            __inc_numa_state();
             __inc_numa_state();
          }
        }
      }
    }
    add_to_page_cache_lru() {
      __add_to_page_cache_locked() {
        PageHuge();
        mem_cgroup_try_charge() {
          get_mem_cgroup_from_mm() {
            mem_cgroup_from_task();
```

```
}
        try_charge() {
          do_memsw_account();
          page_counter_try_charge() {
            propagate_protected_usage();
            propagate_protected_usage();
          }
          refill_stock() {
            drain_stock() {
              page_counter_uncharge() {
                page_counter_cancel() {
                  propagate_protected_usage();
                }
                page_counter_cancel() {
                  propagate_protected_usage();
                }
              do_memsw_account();
          }
        }
      }
      _raw_spin_lock_irq();
      page_cache_tree_insert() {
        shmem_mapping();
        workingset_update_node();
      }
      __inc_node_page_state() {
        __inc_node_state();
      mem_cgroup_commit_charge() {
        mem_cgroup_charge_statistics();
        memcg_check_events();
        do_memsw_account();
      }
   lru_cache_add() {
      __lru_cache_add();
    }
  }
}
alloc_page_buffers() {
  get_mem_cgroup_from_page();
  alloc_buffer_head() {
    kmem_cache_alloc() {
      _cond_resched() {
        rcu_all_qs();
      }
      should_failslab();
      memcg_kmem_get_cache();
      memcg_kmem_put_cache();
    recalc_bh_state();
  }
  set_bh_page();
```

```
7/17/2020
                                          md2pdf - Markdown to PDF
    }
    _raw_spin_lock();
    init_page_buffers() {
      I_BDEV();
    }
    unlock_page();
    __find_get_block() {
      pagecache_get_page() {
        find_get_entry();
        mark_page_accessed() {
          workingset_activation();
        }
      }
      _raw_spin_lock();
      __brelse();
    }
  ll_rw_block() { // low level access to block devices
    submit_bh() {
      submit_bh_wbc(struct buffer_head *bh, struct writeback_control *wbc) {
        bio_alloc_bioset() {
          mempool_alloc() {
            _cond_resched() {
              rcu_all_qs();
            }
            mempool_alloc_slab() {
              kmem_cache_alloc() {
                should_failslab();
                memcg_kmem_put_cache();
              }
            }
          bio_init();
        bio_add_page() {
          __bio_try_merge_page();
          __bio_add_page();
        }
        guard_bio_eod() { // 这允许我们在设备的奇数个最后扇区上执行IO偶数,即使块大小是物理扇区大小{
          __disk_get_part(bio->bi_disk, bio->bi_partno);
        submit_bio(struct bio *bio) { // fs/buffer.c:3094 //submit a bio to the block de
          generic_make_request() { //block/block-core.c:2415
            blk_queue_enter(); // try to increase q->q_usage_counter
            generic_make_request_checks() { //check device exist, sector num, read only.
              _cond_resched() {
                rcu_all_qs();
              }
              should_fail_bio();
              __disk_get_part();
              blk_partition_remap(); // CX____ add
              /* block_bio_remap: 8,0 RA 58724688 + 8 <- (8,2) 58720592 */
              create_io_context(GFP_ATOMIC, q->node); // CX____ add
              blkcg_bio_issue_check () {
                  blkcg = bio_blkcg(bio); {
```

```
kthread_blkcg();
      }
      bio_associate_blkcg();
      blkg_lookup_slowpath();
      blk_throtl_bio() {
          /* blk_throtl_bio[2148] --- */
          throtl_update_latency_buckets(td);
          blk_throtl_assoc_bio(tg, bio);
          blk_throtl_update_idletime(tg);
          throtl_downgrade_check(tg);
          throtl_upgrade_check(tg);
          if (throtl_can_upgrade(td, tg)) {
              throtl_upgrade_state(td);
              goto again;
          throtl_charge_bio(tg, bio) {
              bio_set_flag(bio, BIO_THROTTLED);
          }
          throtl_trim_slice(tg, rw);
          throtl_add_bio_tg(bio, qn, tg);
          if (tg->flags & THROTL_TG_WAS_EMPTY) {
              tg_update_disptime(tg);
              throtl_schedule_next_dispatch(tg->service_queue.parent_sq, tru
          }
      }
      blkg_rwstat_add();
  }
  trace_block_bio_queue(); /* block_bio_queue: 8,0 RA 58724688 + 8 [bash] */
}
/*int blk_init_allocated_queue(struct request_queue *q)
      {
          blk_queue_make_request(q, blk_queue_bio);
*/
q->make_request_fn /* aka */ blk_queue_bio() {
  blk_queue_bounce();
  blk_queue_split();
  bio_integrity_prep();
  blk_attempt_plug_merge();
  _raw_spin_lock_irq();
  elv_merge() {
    elv_rqhash_find();
  }
  rq_qos_throttle() {
    blkcg_iolatency_throttle() {
      /* blkcg_iolatency_throttle[439] --- */
      bio_associate_blkcg();
      blkg_lookup_slowpath();
      ktime_get();
      /* CX____ bio_issue_init[139]: issue->value=18014516950275448 */
      bio_associate_blkg();
      /* blkcg_iolatency_throttle[458] --- bio->bi_issue=137408 */
      // TODO: here change the scale, max_depth
      check_scale_change();
```

```
//__blkcg_iolatency_throttle();
/* __blkcg_iolatency_throttle[207] --- use_delay ^^^ use=0 */
// TODO: This set's the notify_resume for the task to check and see if
blkcg_schedule_throttle() {
 blk_get_queue();
 /* blkcg_schedule_throttle[1807] --- current->throttle_queue.id=8 */
 blk_put_queue(); //减小引用?
 current->throttle_queue = q;
   if (use_memdelay)
       current->use_memdelay = use_memdelay;
   set_notify_resume(current); // 向指定的进程设置了一个TIF_NOTIFY_RESUME
}
// TODO: 重要,增加了inflight
iolatency_may_queue() {
   rq_wait_inc_below() {
       /* rq_wait_inc_below[30] --- set rq_weit->inflight to 1 */
       /* atomic_inc_below[13] --- */
   }
}
/* 将iolat->rg_wait->wait添加到等待队列 */
prepare_to_wait_exclusive(TASK_UNINTERRUPTIBLE) {
 _raw_spin_lock_irqsave();
 set_current_state(WQ_FLAG_EXCLUSIVE);
 _raw_spin_unlock_irqrestore();
iolatency_may_queue () {
   rq_wait_inc_below() {
       /* rq_wait_inc_below[30] --- set rq_weit->inflight to 1 */
       /* atomic_inc_below[13] --- */
   }
}
/* __blkcq_iolatency_throttle[237] --- io_schedule() */
// TODO: 打印io_schedule()耗时,看当前进程被wait的时间
io_schedule() {
   current->in_iowait = 1;
   /* 当task发生iowait的时候,内核对他们的处理方法是将task切换出去,让可运行的t
   /* 内核在调用io_schedule,io_schedule_timeout前都会设置task运行状态TASK
   /* http://m.elecfans.com/article/611049.html */
   io_schedule_prepare();
     blk_schedule_flush_plug(current);
 schedule() {
   rcu_note_context_switch() {
      rcu_sched_qs();
   }
   _raw_spin_lock();
   update_rq_clock();
   deactivate_task() {
     dequeue_task_fair() {
       dequeue_entity() {
         update_curr() {
           update_min_vruntime();
           cpuacct_charge();
           __cgroup_account_cputime() {
             cgroup_rstat_updated();
           }
```

```
}
      __update_load_avg_se();
      __update_load_avg_cfs_rq();
      clear_buddies();
      account_entity_dequeue();
      update_cfs_group();
      update_min_vruntime();
    }
    dequeue_entity() {
      update_curr() {
        update_min_vruntime();
      }
      __update_load_avg_se();
      __update_load_avg_cfs_rq();
      clear_buddies();
      account_entity_dequeue();
      update_cfs_group() {
        reweight_entity();
      }
      update_min_vruntime();
    hrtick_update();
  }
 _delayacct_blkio_start() {
  ktime_get();
}
pick_next_task_fair() {
  check_cfs_rq_runtime();
  pick_next_entity() {
    __pick_first_entity();
    clear_buddies();
  }
  pick_next_entity() {
    __pick_first_entity();
    clear_buddies();
  }
  put_prev_entity() {
    check_cfs_rq_runtime();
    check_spread();
  }
  set_next_entity() {
    __update_load_avg_se() {
      decay_load();
      decay_load();
      decay_load();
      __accumulate_pelt_segments() {
        decay_load();
        decay_load();
      }
    }
     _update_load_avg_cfs_rq() {
      decay_load();
      decay_load();
      decay_load();
```

```
__accumulate_pelt_segments() {
                          decay_load();
                          decay_load();
                        }
                      }
                    }
                    put_prev_entity() {
                      check_cfs_rq_runtime();
                      check_spread();
                    }
                    set_next_entity() {
                      __update_load_avg_se() {
                        decay_load();
                        decay_load();
                        decay_load();
                        __accumulate_pelt_segments() {
                          decay_load();
                          decay_load();
                        }
                       _update_load_avg_cfs_rq();
                    }
                  }
                  switch_mm_irqs_off() {
                    load_new_mm_cr3();
   sleep-2186
               => systemd-368
-----
0) systemd-387 => sleep-2186
                  finish_task_switch();
                } /* schedule */
              } /* io_schedule */
              _raw_spin_lock_irq();
              prepare_to_wait_exclusive() {
                _raw_spin_lock_irqsave();
                _raw_spin_unlock_irqrestore();
              }
              // iolatency_may_queue
              rq_wait_inc_below() {
                /* rq_wait_inc_below[30] --- set rq_weit->inflight to 1 */
                /* atomic_inc_below[13] --- */
              }
              //break;
              finish_wait() {
                _raw_spin_lock_irqsave();
                _raw_spin_unlock_irqrestore();
            } /* blkcg_iolatency_throttle */
          } /* rq_qos_throttle */
          get_request() {
            blkg_lookup_slowpath();
```

```
elv_may_queue();
  mempool_alloc() {
    _cond_resched() {
      rcu_all_qs();
    }
    alloc_request_size() {
      __kmalloc_node() {
        kmalloc_slab();
        should_failslab();
        memcg_kmem_put_cache();
      }
      scsi_old_init_rq() {
        kmem_cache_alloc_node() {
          should_failslab();
          memcg_kmem_put_cache();
        }
      }
    }
  blk_rq_init() {
    ktime_get();
  }
  elv_set_request();
  /* block_getrq: 8,0 RA 58724688 + 8 [bash] */
}
rq_qos_track();
blk_init_request_from_bio() {
  blk_rq_bio_prep();
}
_raw_spin_lock_irq();
add_acct_request() {
    blk_account_io_start() {
    disk_map_sector_rcu();
    part_round_stats();
    part_inc_in_flight();
    }
     _elv_add_request() {
    /* block_rq_insert: 8,0 RA 4096 () 58724688 + 8 [bash] */
    elv_rqhash_add();
    noop_add_request();
    }
}
__blk_run_queue() {
  scsi_request_fn() {
    blk_peek_request() {
      noop_dispatch() {
        elv_dispatch_sort() {
          elv_rqhash_del();
        }
      /* block_rq_issue: 8,0 RA 4096 () 58724688 + 8 [bash] */
      scsi_prep_fn() {
        scsi_prep_state_check();
        get_device();
        scsi_init_command() {
```

```
scsi_initialize_rq() {
        scsi_req_init();
      }
      init_timer_key();
      scsi_add_cmd_to_list();
    }
    scsi_setup_cmnd() {
      sd_init_command() {
        scsi_init_io() {
          scsi_init_sgtable() {
            mempool_alloc() {
              mempool_alloc_slab() {
                kmem_cache_alloc() {
                  should_failslab();
                  memcg_kmem_put_cache();
                }
              }
            }
            blk_rq_map_sg();
          }
        }
      }
    }
  }
blk_queue_start_tag() {
  blk_start_request() {
    ktime_get();
    rq_qos_issue();
    blk_add_timer() {
      round_jiffies_up() {
        round_jiffies_common();
      }
      blk_rq_timeout() {
        round_jiffies_up() {
          round_jiffies_common();
        }
      }
  }
}
scsi_init_cmd_errh();
scsi_dispatch_cmd() {
  scsi_log_send();
  ata_scsi_queuecmd() {
    _raw_spin_lock_irqsave();
    ata_scsi_find_dev() {
      __ata_scsi_find_dev() {
        ata_find_dev();
      }
    }
    ata_scsi_translate() {
      ata_qc_new_init();
      ata_sg_init();
      ata_scsi_rw_xlat() {
```

```
ata_check_nblocks();
                         ata_build_rw_tf();
                      }
                      ahci_pmp_qc_defer() {
                         ata_std_qc_defer();
                      }
                      ata_qc_issue() {
                         dma_direct_map_sg() {
                           check_addr();
                         }
                         ahci_qc_prep() {
                           ata_tf_to_fis();
                           ahci_fill_cmd_slot();
                         ahci_qc_issue();
                      }
                     _raw_spin_unlock_irqrestore();
                  }
                }
                _raw_spin_lock_irq();
                blk_peek_request() {
                  noop_dispatch();
                }
              }
            }
          } /* blk_queue_bio */
          blk_queue_exit();
        } /* generic_make_request */
      } /* submit_bio */
    } /* submit_bh_wbc */
  } /* submit_bh */
} /* ll_rw_block */
```

#### structs 相关结构体

```
struct blkcg {
    struct cgroup_subsys_state css;
    spinlock_t
                        lock;
    struct radix_tree_root
                                 blkg_tree;
    struct blkcg_gq __rcu
                                 *blkg_hint;
    struct hlist_head
                             blkg_list;
    struct blkcg_policy_data
                                 *cpd[BLKCG_MAX_POLS];
    struct list_head
                             all_blkcgs_node;
#ifdef CONFIG_CGROUP_WRITEBACK
    struct list head
                             cgwb_list;
    refcount t
                        cgwb_refcnt;
#endif
};
```

```
struct blkcg_gq {
    /* Pointer to the associated request_queue */
    struct request_queue *q;
    struct list_head
                         q_node;
    struct hlist_node
                          blkcg_node;
    struct blkcg
                           *blkcg;
    * Each blkg gets congested separately and the congestion state is
    * propagated to the matching bdi_writeback_congested.
    struct bdi_writeback_congested *wb_congested;
    /* all non-root blkcg_gq's are guaranteed to have access to parent */
    struct blkcg_gg
                           *parent;
    /* request allocation list for this blkcg-q pair */
    struct request_list
                        rl:
    /* reference count */
    atomic_t
                       refcnt;
    /* is this blkg online? protected by both blkcg and q locks */
    bool
                       online;
    struct blkg_rwstat
                         stat_bytes;
    struct blkg_rwstat stat_ios;
    struct blkg_policy_data
                              *pd[BLKCG_MAX_POLS];
    struct rcu head
                           rcu_head;
    atomic_t
                       use_delay;
    atomic64_t
                      delay_nsec;
                       delay_start;
    atomic64_t
   u64
                  last_delay;
   int
                   last_use;
};
struct blk_iolatency {
    struct rq_qos rqos;
    struct timer_list timer;
    atomic_t enabled;
};
struct bio {
   struct bio
                   *bi_next; /* request queue link */
    struct gendisk
                       *bi_disk;
                                  /* bottom bits req flags,
    unsigned int
                       bi_opf;
                        * top bits REQ_OP. Use
                        * accessors.
                        */
    unsigned short
                       bi_flags; /* status, etc and bvec pool number */
    unsigned short
                       bi_ioprio;
```

```
unsigned short
                        bi_write_hint;
    blk_status_t
                        bi_status;
                bi_partno;
    /* Number of segments in this BIO after
    * physical address coalescing is performed.
    */
    unsigned int
                        bi_phys_segments;
     * To keep track of the max segment size, we account for the
    * sizes of the first and last mergeable segments in this bio.
     * /
    unsigned int
                        bi_seg_front_size;
    unsigned int
                        bi_seg_back_size;
    struct bvec_iter bi_iter;
    atomic t
                   __bi_remaining;
    bio_end_io_t
                        *bi_end_io;
                    *bi_private;
#ifdef CONFIG_BLK_CGROUP
    /*
     * Optional ioc and css associated with this bio. Put on bio
     * release. Read comment on top of bio_associate_current().
    */
    struct io_context *bi_ioc;
    struct cgroup_subsys_state *bi_css;
                      *bi_blkg;
    struct blkcg_gq
    struct bio_issue bi_issue;
#endif
};
struct iolatency_grp {
    struct blkg_policy_data pd;
    struct blk_rq_stat __percpu *stats;
    struct blk_iolatency *blkiolat;
    struct rq_depth rq_depth;
    struct rq_wait rq_wait;
    atomic64_t window_start;
    atomic_t scale_cookie;
    u64 min_lat_nsec;
    u64 cur_win_nsec;
    /* total running average of our io latency. */
    u64 lat_avg;
    /* Our current number of IO's for the last summation. */
    u64 nr_samples;
    struct child_latency_info child_lat;
};
```

#### actions

#### rq to blkg path

blk\_qc\_t generic\_make\_request(struct bio \*bio) struct request\_queue \*q = bio->bi\_disk->queue; blk\_queue\_bio(struct request\_queue \*q, struct bio \*bio) rq\_qos\_throttle(struct request\_queue \*q, struct bio \*bio, spinlock\_t \*lock) blkcg\_iolatency\_throttle(struct rq\_qos \*rqos, struct bio \*bio, spinlock\_t \*lock)

## latency qos与wbt的差异

latency qos与wbt实现很相似,但属于两种不同的东西, 启用latency时wbt并未启用,后续列出详细差异

#### **TODO**

- [x] 尝试是否能找到适合google pixel 2xl的4.15以上的内核,若能找到,编译一个可开启cgroup v2 的版本并测试 . //不可行,仅找到4.4内核
- []继续研究组迁移时的逻辑,若与推测相符,提出修改思路 . //重点放在快速改变队列深度,以及队列深度对读写速度的影响
- []根据对iolatency的已有研究,尝试移植到4.9版本内核上去 // 优先级低, 目前看来可行性不高
- [x] 补充generic\_make\_request与设备驱动之间的联系 //已经理清,需要总结
- [x] 什么情况下走了wbt, 什么情况下走了delay & io\_schedule // wbt 与iolatency为不同的 rq\_qos, 各自通过rq\_qos\_add来生效, wbt->BLK\_WBT; 二者实现原理相似,互不依赖.
- [x] 研究一下latency\_unknown由来 , 优先级低 //属于wbt, skip
- [] 研究队列深度对组写入速度的影响 //在emmc, ssd, hdd和不周文件系统均有差异,需要分开研究
- [x] 研究一种在一些位置打印出组信息,队列信息,队列深度的方法 //队列深度未完成,rq和rqos 队列深度
- []研究组切换时,能否实时调整组深度,和实时的组深度变化
- []研究generic\_perform\_request与pagecache的路径关系
- [x] 研究wb\_timer\_fn与时间窗口是否对应 //wbt, skip
- [x] emmc流程需要抓log整理,找出一些方法指针的对应 //暂时放弃,seattletmo为ufs, 与x86一样走scsi. 需注意可能存在mq
- [x] 找出所有的request\_queue //确定一个磁盘设备一个rq
- [x] request\_queue 与cgroup的对应关系,是否一对一 //no, gendisk<->request\_queue
- [x] 在某一时刻打出rq里有多少的request
- [x] 在throttle方法中打出当前的rg所属组 //不可行, rg为已经没有组归属, 仅有一个root\_blkg
- [x] 参考cfq在权重改变时的重排方法 //latency 并不需要,并未在电梯层使用树

- [] 在一个组上, depth, inflight代表什么意义
- [x] 跟踪一个bio,弄明白bio的完成机制,是在放入rq时完结,还是被下发到驱动才完结 //done mostly
- [] single core不走io\_scheduler()路径, why?与smp区别在哪里? //低优先级研究
- [] 走到io\_scheduler的具体条件是什么? // on going
- [] nr\_samples, samples\_thresh, deepth, scale\_cookie
- [] plug unplug
- [] depth & use\_delay //全部为随机4k时,应用depth, 没有限流效果 //use\_delay在哪里执行的 delay? //if (iolat->rq\_depth.max\_depth == 1 && direction < 0), use\_delay //blkcg\_use\_delay --> atomic\_inc(&blkg->blkcg->css.cgroup->congestion\_count); //depth的调整速度是关键,目前看来 最需要修改的是这一块。
- [] block io 在memory部分的逻辑,cgroup v2相比v1最重要的改变之一就是关联了cache