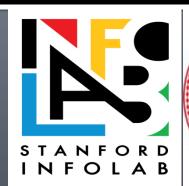
# PageRank

Random Surfers on the Web Transition Matrix of the Web Dead Ends and Spider Traps Topic-Specific PageRank Hubs and Authorities

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#### Intuition — (1)

- Web pages are important if people visit them a lot.
- But we can't watch everybody using the Web.
- A good surrogate for visiting pages is to assume people follow links randomly.
- Leads to random surfer/walker model:
  - Start at a random page and follow random out-links repeatedly, from whatever page you are at.
- PageRank of a page = limiting probability of being at that page.

#### Intuition — (2)

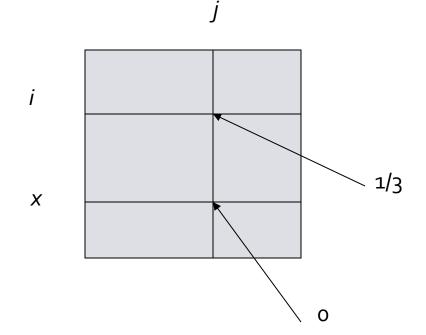
- Solve the recursive equations: "importance of a page = its share of the importance of each of its predecessor pages."
  - Equivalent to the random-surfer definition of PageRank.

#### Transition Matrix of the Web

- Number the pages 1, 2,... .
  - Page i corresponds to row and column i.
- M [i, j] = 1/n if page j links to n pages, including page i; 0 if j does not link to i.
  - M [i, j] is the probability a surfer will next be at page i if it is now at page j.
  - Or it is the share of j's importance that i receives.
  - Note: no matter how many links to page i you find on page j, give i only 1/n<sup>th</sup> of j's importance.

#### **Example: Transition Matrix**

Suppose page *j* links to 3 pages, including *i* but not *x*.



Called a (*column*) *stochastic matrix* = "all columns sum to 1."

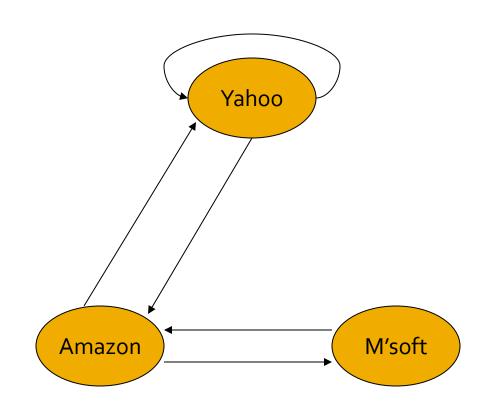
#### Random Walks on the Web

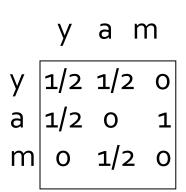
- Suppose v is a vector whose i th component is the probability that a random surfer is at page i at a certain time.
- If a surfer chooses a successor page from page i at random, the probability distribution for surfers is then given by the vector Mv.

#### Random Walks - (2)

- Starting from any vector u, the limit M (M (...M (M u ) ...)) is the long-term distribution of the surfers.
- The math: limiting distribution = principal eigenvector of M = PageRank.
  - Note: If v is the limit of MM...Mu, then v satisfies the equation v = Mv, so v is an eigenvector of M with eigenvalue 1.

## Example: The Web in 1839





# Solving The Equations

- Because there are no constant terms, the equations v = Mv do not have a unique solution.
  - Example: doubling each component of solution v yields another solution.
- In Web-sized examples, we cannot solve by Gaussian elimination anyway; we need to use relaxation (= iterative solution).

## Simulating a Random Walk

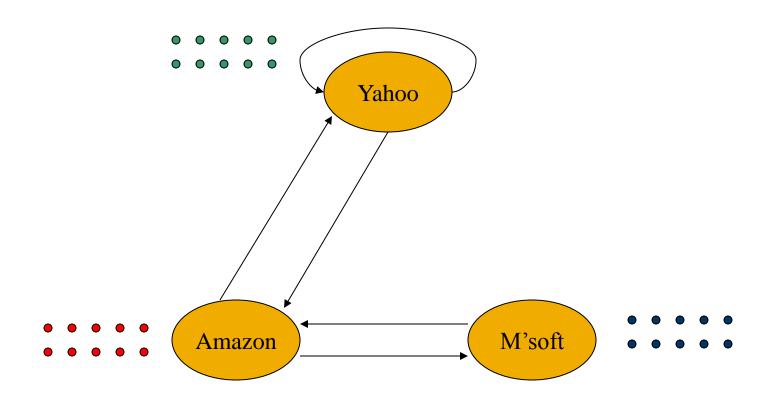
- Start with the vector u = [1, 1,..., 1] representing the idea that each Web page is given one unit of importance.
  - Note: it is more common to start with each vector element = 1/N, where N is the number of Web pages and to keep the sum of the elements at 1.
  - Question for thought: Why such small values?
- Repeatedly apply the matrix M to u, allowing the importance to flow like a random walk.
- About 50 iterations is sufficient to estimate the limiting solution.

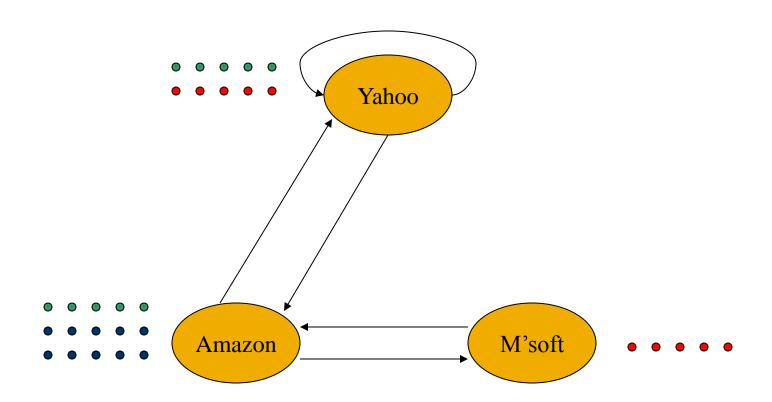
## Example: Iterating Equations

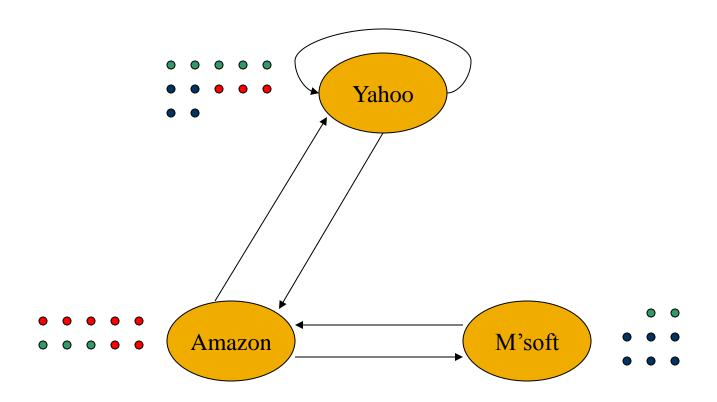
• Equations  $\mathbf{v} = M\mathbf{v}$ :

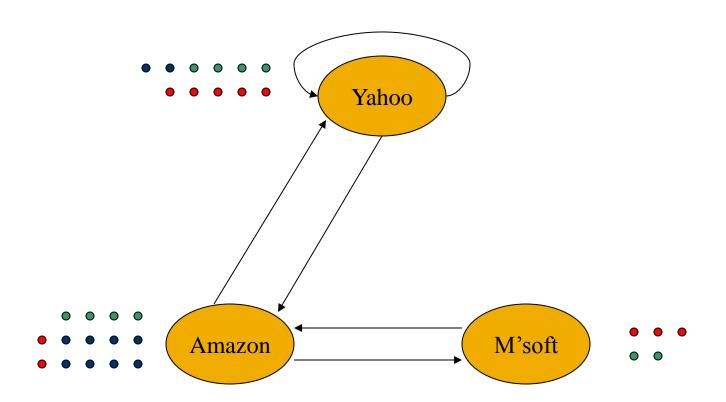
$$y = y/2 + a/2$$
  
 $a = y/2 + m$   
 $m = a/2$ 

Note: "=" is really "assignment," but in parallel

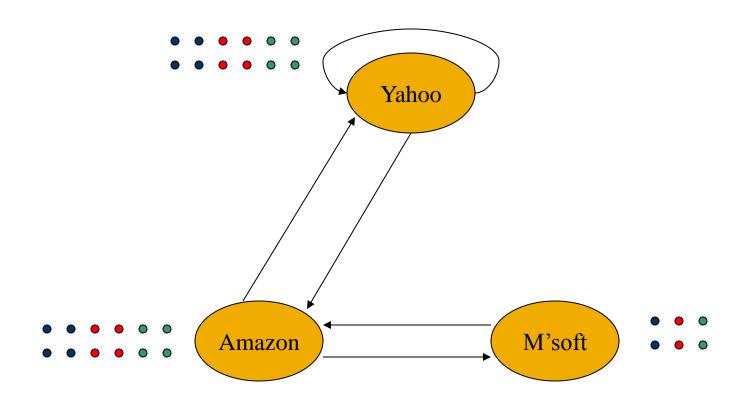








# In the Limit ...

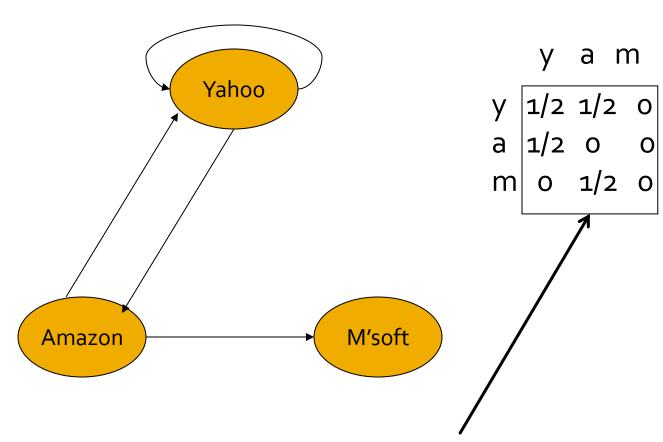


# The Web Is More Complex Than That

Dead Ends
Spider Traps
Taxation Policies

#### Real-World Problems

- Some pages are dead ends (have no links out).
  - Such a page causes importance to leak out, or surfers to disappear.
- Other groups of pages are spider traps (all outlinks are within the group).
  - Eventually spider traps absorb all importance; all surfers get stuck in the trap.



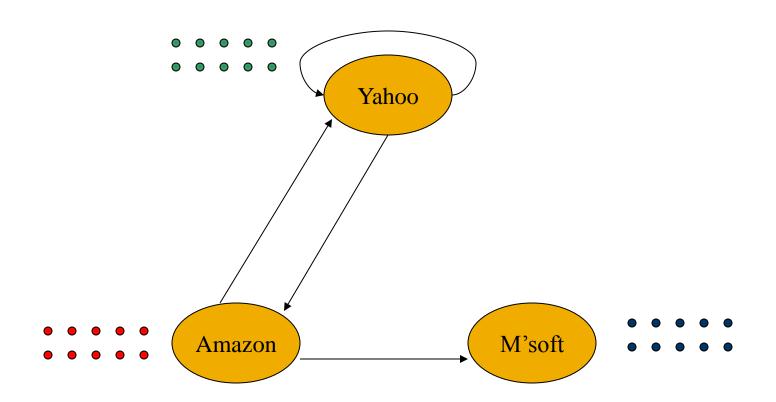
A (*column*) *substochastic matrix* = "all columns sum to at most 1."

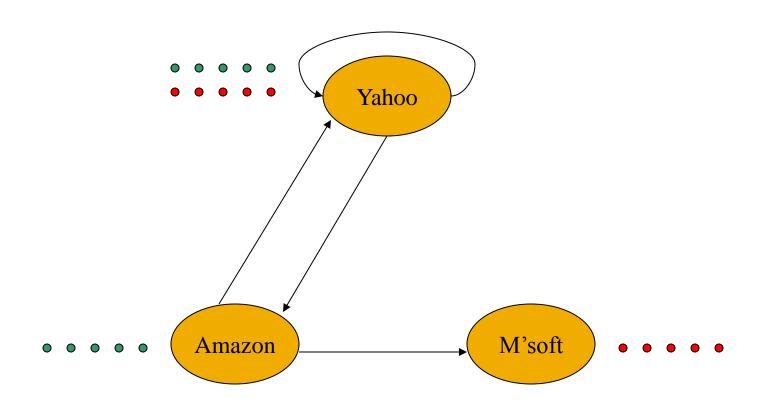
#### **Example: Effect of Dead Ends**

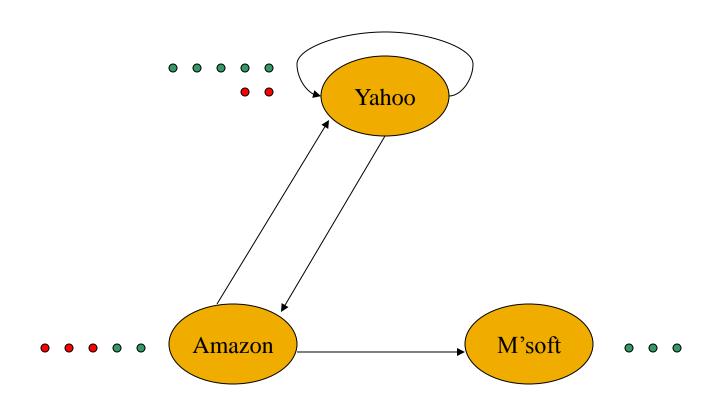
• Equations  $\mathbf{v} = M\mathbf{v}$ :

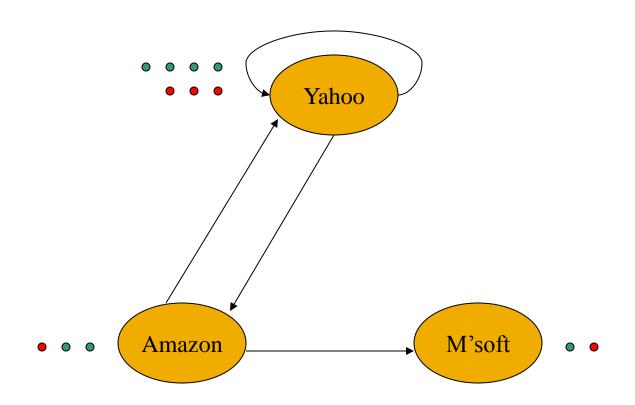
$$y = y/2 + a/2$$
  
 $a = y/2$   
 $m = a/2$ 

$$y$$
 1 1 3/4 5/8 0
a = 1 1/2 1/2 3/8 ... 0
m 1 1/2 1/4 1/4 0

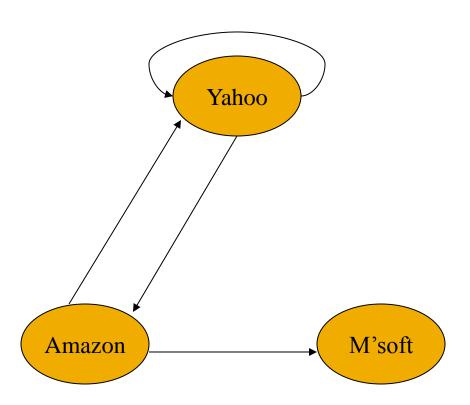




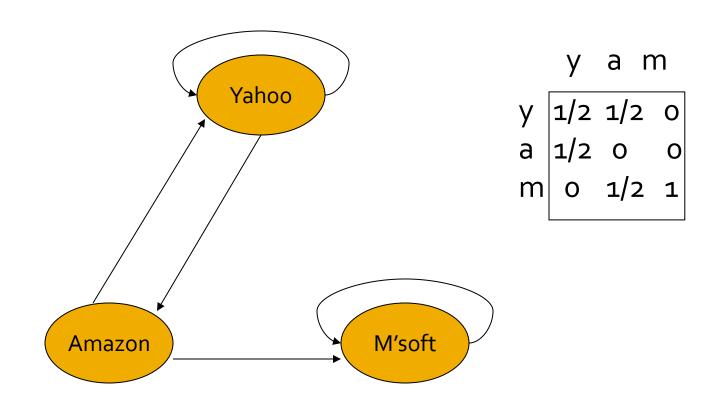




# In the Limit ...



## M'soft Becomes Spider Trap



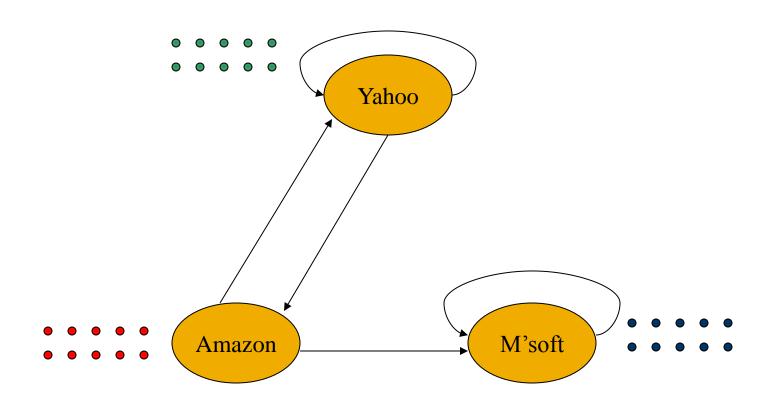
## Example: Effect of Spider Trap

• Equations  $\mathbf{v} = M\mathbf{v}$ :

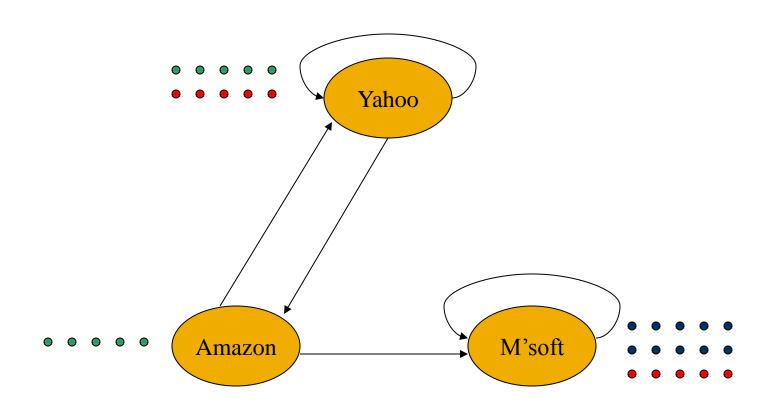
$$y = y/2 + a/2$$
  
 $a = y/2$   
 $m = a/2 + m$ 

$$y$$
 1 1 3/4 5/8 0
 $a = 1$  1/2 1/2 3/8 ... 0
 $m$  1 3/2 7/4 2

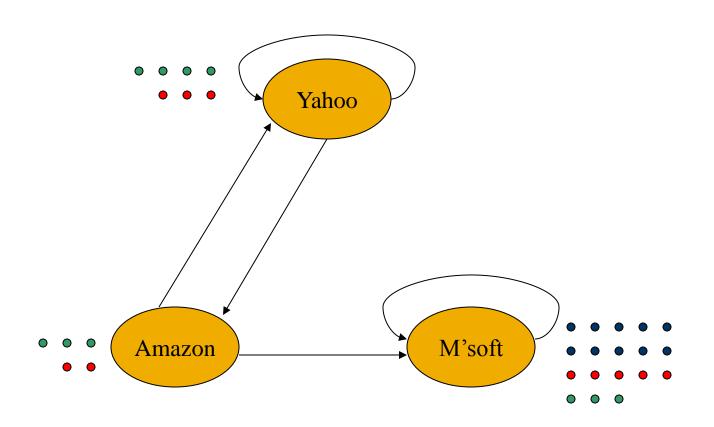
# Microsoft Becomes a Spider Trap



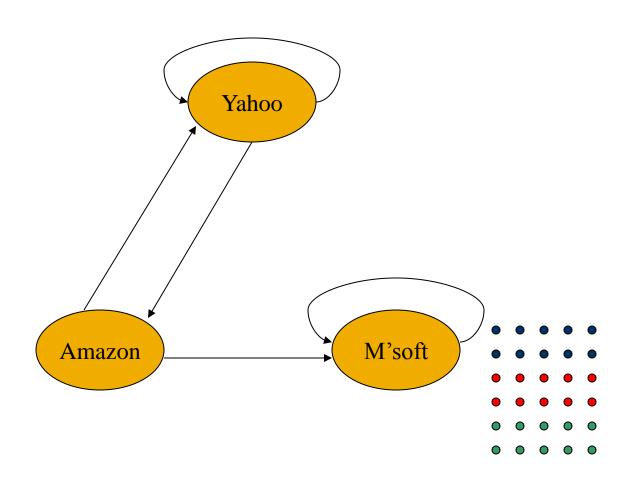
## Microsoft Becomes a Spider Trap



## Microsoft Becomes a Spider Trap



# In the Limit ...



## PageRank Solution to Traps, Etc.

- "Tax" each page a fixed percentage at each iteration.
- Add a fixed constant > 0 to all pages.
  - Optional but useful: add exactly enough to balance the loss (tax + PageRank of dead ends).
- Models a random walk with a fixed probability of leaving the system, and a fixed number of new surfers injected, at random pages, into the system at each step.

#### Example: Microsoft is a Spider Trap; 20% Tax

• Equations v = 0.8(Mv) + 0.2:

$$y = 0.8(y/2 + a/2) + 0.2$$

$$a = 0.8(y/2) + 0.2$$

$$m = 0.8(a/2 + m) + 0.2$$

Note: amount injected is chosen to balance the tax. If we started with 1/3 for each rather than 1, the 0.2 would be replaced by 0.0667.

# Topic-Specific PageRank

Focusing on Specific Pages
Teleport Sets
Interpretation as a Random Walk

## Topic-Specific Page Rank

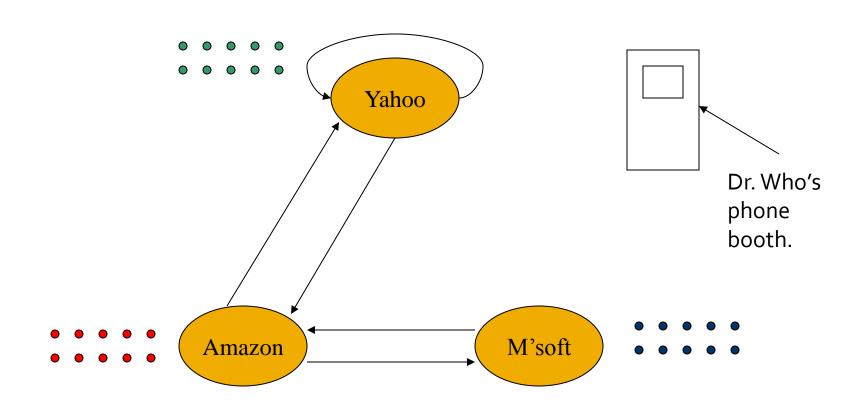
- Goal: Evaluate Web pages not just by popularity, but also by relevance to a particular topic, e.g. "sports" or "history."
- Allows search queries to be answered based on interests of the user.
- Example: Search query [jaguar] wants different pages depending on whether you are interested in automobiles, nature, or sports.
  - Might discover interests by recent browsing history, bookmarks, e.g.

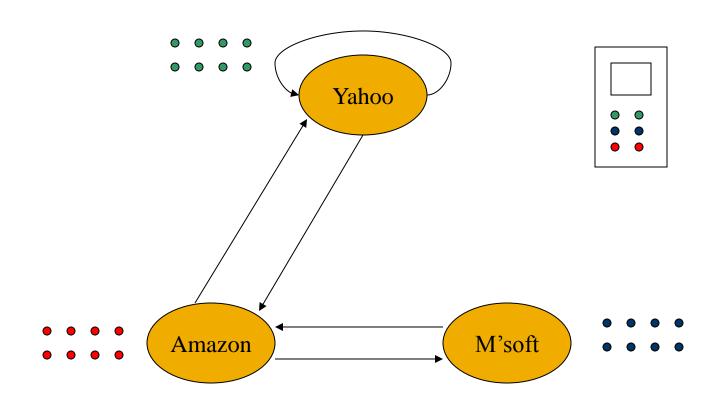
#### **Teleport Sets**

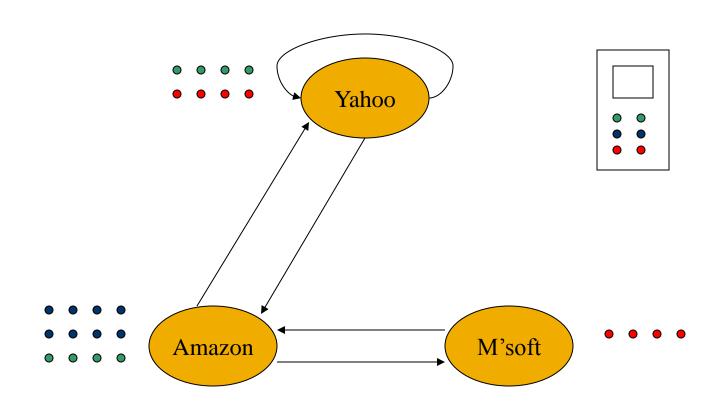
- Assume each surfer has a small probability of "teleporting" at any tick.
- Teleport can go to:
  - 1. Any page with equal probability.
    - As in the "taxation" scheme.
  - 2. A set of "relevant" pages (teleport set).
    - For topic-specific PageRank.
  - Note: can also inject surfers to compensate for surfers lost at dead ends.
    - Or imagine a surfer always teleports from a dead end.

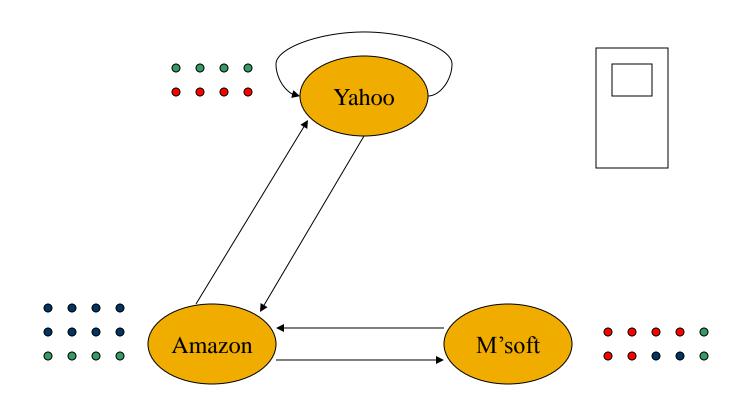
## Example: Topic = Software

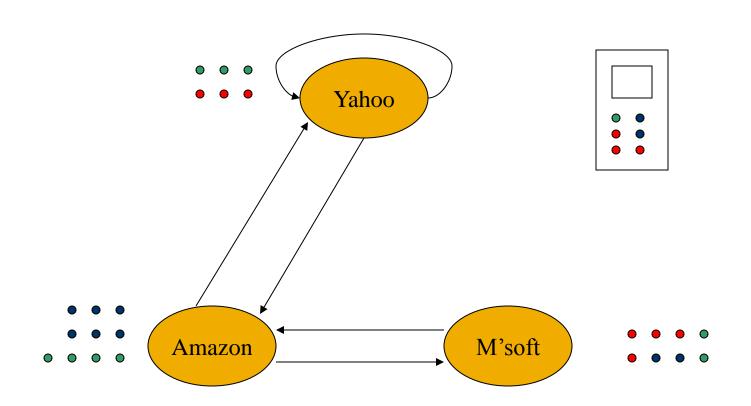
- Only Microsoft is in the teleport set.
- Assume 20% "tax."
  - I.e., probability of a teleport is 20%.

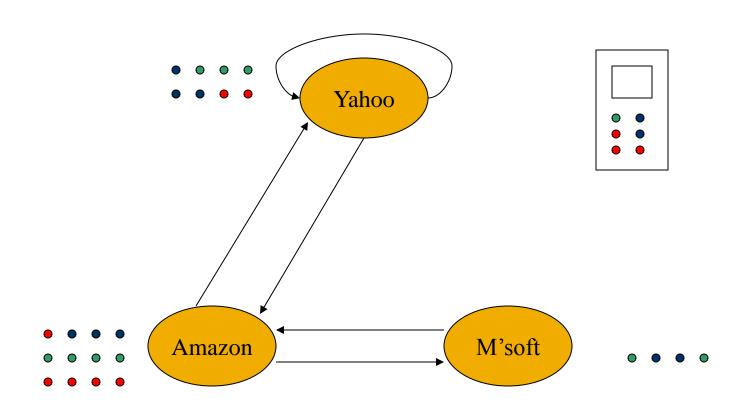


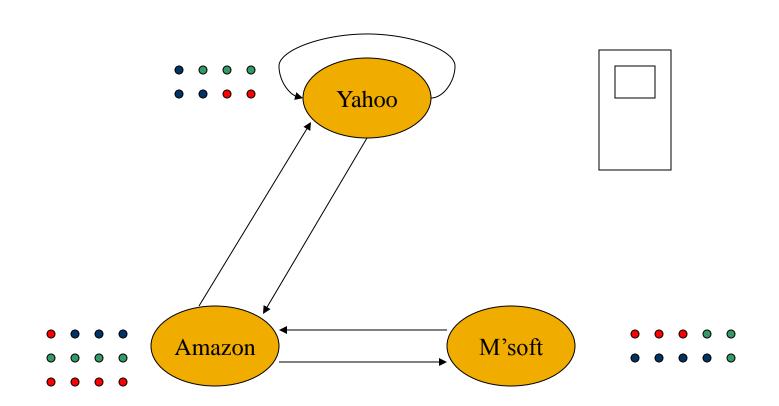












## Picking the Teleport Set

- One option is to choose the pages belonging to the topic in Open Directory.
- 2. Another option is to "learn," from a training set (which could be Open Directory), the typical words in pages belonging to the topic; use pages heavy in those words as the teleport set.

### **Application: Link Spam**

- Spam farmers create networks of millions of pages designed to focus PageRank on a few undeserving pages.
  - We'll discuss this "technology" next time.
- To minimize their influence, use a teleport set consisting of trusted pages only.
  - Example: home pages of universities.

## HITS

Hubs
Authorities
Solving the Implied Recursion

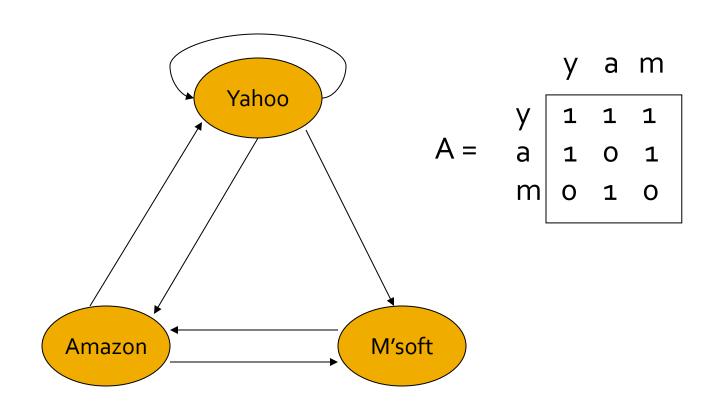
#### Hubs and Authorities ("HITS")

- Mutually recursive definition:
  - A hub links to many authorities;
  - An authority is linked to by many hubs.
- Authorities turn out to be places where information can be found.
  - Example: course home pages.
- Hubs tell where the authorities are.
  - Example: departmental course-listing page.

#### Transition Matrix A

- HITS uses a matrix A[i, j] = 1 if page i links to page j, 0 if not.
- $A^T$ , the transpose of A, is similar to the PageRank matrix M, but  $A^T$  has 1's where M has fractions.
- Also, HITS uses column vectors h and a representing the degrees to which each page is a hub or authority, respectively.
- Computation of h and a is similar to the iterative way we compute PageRank.

### **Example: H&A Transition Matrix**



#### Using Matrix A for HITS

- Powers of A and  $A^T$  have elements whose values grow exponentially, so we need scale factors  $\lambda$  and  $\mu$ .
- Let h and a be column vectors measuring the "hubbiness" and authority of each page.
- Equations:  $\mathbf{h} = \lambda A \mathbf{a}$ ;  $\mathbf{a} = \mu A^T \mathbf{h}$ .
  - Hubbiness = scaled sum of authorities of successor pages (out-links).
  - Authority = scaled sum of hubbiness of predecessor pages (in-links).

#### Consequences of Basic Equations

- From  $\mathbf{h} = \lambda A \mathbf{a}$ ;  $\mathbf{a} = \mu A^T \mathbf{h}$  we can derive:
  - $\mathbf{h} = \lambda \mu A A^T \mathbf{h}$
  - $a = \lambda \mu A^T A a$
- We could compute h and a by iteration, assuming initially each page has one unit of hubbiness and one unit of authority.
- Technically, these equations let you solve for λμ as well as h and a.
- In practice, you don't fix  $\lambda\mu$ , but rather scale the result at each iteration.
  - Example: scale to keep largest value at 1.

#### Scale Doesn't Matter

- Remember: it is only the direction of the vectors, or the relative hubbiness and authority of Web pages that matters.
- As for PageRank, the only reason to worry about scale is so you don't get overflows or underflows in the values as you iterate.

## Example: Iterating H&A

$$\mathbf{a} = \lambda \mu A^T A \mathbf{a}$$
;  $\mathbf{h} = \lambda \mu A A^T \mathbf{h}$ 

$$A = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}$$

$$A^{T} = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 0 & 1 \\ 1 & 1 & 0 \end{bmatrix}$$

$$AA^{T} = \begin{bmatrix} 3 & 2 & 1 \\ 2 & 2 & 0 \\ 1 & 0 & 1 \end{bmatrix}$$

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$$A^{T}A = \begin{bmatrix} 1 & 1 & 0 \\ 1 & 2 & 1 \\ 2 & 1$$

#### A Better Way to Solve HITS

- Start with h = [1,1,...,1]; multiply by A<sup>T</sup> to get first a; scale so largest component = 1; then multiply by A to get next h, and repeat until approximate convergence.
- You may be tempted to compute AA<sup>T</sup> and A<sup>T</sup>A
  first, then iterate multiplication by these
  matrices, as for PageRank.
- Question for thought: Why is the separate calculations of h and a actually less efficient than the method suggested on this slide?