

Project Proposal

Proof-of-Replication Implementation



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1. Background

1.1. Choice of Topic

My area of interest in blockchain consensus is Proof-of-Storage, specifically Proof-of-Replication (PoRep) and Proof-of-Spacetime (PoSt), which Filecoin has been researching and developing. In this class project, I plan to implement the PoRep consensus based on Filecoin's research. If time permits, I will extend the PoRep to PoSt.

1.2. Theory

Proof-of-Replication: a consensus gives proof to the user (verifier V) that some data D has been replicated and stored by the server (prover P) in physical storage.

Proof-of-Spacetime: a consensus ensures the storage provider is storing data at a specific time of challenge.

| Filecoin PoRep protocol | Filecoin PoSt protocol |
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| <p>Setup</p> <ul style="list-style-type: none"> INPUTS: <ul style="list-style-type: none"> prover key pair (pk_p, sk_p) prover SEAL key pk_{SEAL} data \mathcal{D} OUTPUTS: replica \mathcal{R}, Merkle root rt of \mathcal{R}, proof π_{SEAL} <ol style="list-style-type: none"> 1) Compute $h_{\mathcal{D}} := CRH(\mathcal{D})$ 2) Compute $\mathcal{R} := Seal^r(\mathcal{D}, sk_p)$ 3) Compute $rt := MerkleCRH(\mathcal{R})$ 4) Set $\vec{x} := (pk_p, h_{\mathcal{D}}, rt)$ 5) Set $\vec{w} := (sk_p, \mathcal{D})$ 6) Compute $\pi_{SEAL} := SCIP.Prove(pk_{SEAL}, \vec{x}, \vec{w})$ 7) Output $\mathcal{R}, rt, \pi_{SEAL}$ <p>Prove</p> <ul style="list-style-type: none"> INPUTS: <ul style="list-style-type: none"> prover <i>Proof-of-Storage</i> key pk_{POS} replica \mathcal{R} random challenge c OUTPUTS: a proof π_{POS} <ol style="list-style-type: none"> 1) Compute $rt := MerkleCRH(\mathcal{R})$ 2) Compute $path :=$ Merkle path from rt to leaf \mathcal{R}_c 3) Set $\vec{x} := (rt, c)$ 4) Set $\vec{w} := (path, \mathcal{R}_c)$ 5) Compute $\pi_{POS} := SCIP.Prove(pk_{POS}, \vec{x}, \vec{w})$ 6) Output π_{POS} | <p>Setup</p> <ul style="list-style-type: none"> INPUTS: <ul style="list-style-type: none"> prover key pair (pk_p, sk_p) prover POST key pair pk_{POST} some data \mathcal{D} OUTPUTS: replica \mathcal{R}, Merkle root rt of \mathcal{R}, proof π_{SEAL} <ol style="list-style-type: none"> 1) Compute $\mathcal{R}, rt, \pi_{SEAL} := PoRep.Setup(pk_p, sk_p, pk_{SEAL}, \mathcal{D})$ 2) Output $\mathcal{R}, rt, \pi_{SEAL}$ <p>Prove</p> <ul style="list-style-type: none"> INPUTS: <ul style="list-style-type: none"> prover PoSt key pk_{POST} replica \mathcal{R} random challenge c time parameter t OUTPUTS: a proof π_{POST} <ol style="list-style-type: none"> 1) Set $\pi_{POST} := \perp$ 2) Compute $rt := MerkleCRH(\mathcal{R})$ 3) For $i = 0 \dots t$: <ol style="list-style-type: none"> a) Set $c' := CRH(\pi_{POST} c i)$ b) Compute $\pi_{POS} := PoRep.Prove(pk_{POS}, \mathcal{R}, c')$ c) Set $\vec{x} := (rt, c, i)$ d) Set $\vec{w} := (\pi_{POS}, \pi_{POST})$ e) Compute $\pi_{POST} := SCIP.Prove(pk_{POST}, \vec{x}, \vec{w})$ 4) Output π_{POST} |

Verify

- INPUTS:

- prover public key, pk_p
- verifier SEAL and POS keys vk_{SEAL}, vk_{POS}
- hash of data \mathcal{D} , $h_{\mathcal{D}}$
- Merkle root of replica \mathcal{R} , rt
- random challenge, c
- tuple of proofs, (π_{SEAL}, π_{POS})

- OUTPUTS: bit b , equals 1 if proofs are valid

- 1) Set $\vec{x}_1 := (pk_p, h_{\mathcal{D}}, rt)$
- 2) Compute $b_1 := \text{SCIP.Verify}(vk_{SEAL}, \vec{x}_1, \pi_{SEAL})$
- 3) Set $\vec{x}_2 := (rt, c)$
- 4) Compute $b_2 := \text{SCIP.Verify}(vk_{POS}, \vec{x}_2, \pi_{POS})$
- 5) Output $b_1 \wedge b_2$

Verify

- INPUTS:

- prover *public key* pk_p
- verifier SEAL and POST keys vk_{SEAL}, vk_{POST}
- hash of some data $h_{\mathcal{D}}$
- Merkle root of some replica rt
- random challenge c
- time parameter t
- tuple of proofs (π_{SEAL}, π_{POST})

- OUTPUTS: bit b , equals 1 if proofs are valid

- 1) Set $\vec{x}_1 := (pk_p, h_{\mathcal{D}}, rt)$
- 2) Compute $b_1 := \text{SCIP.Verify}(vk_{SEAL}, \vec{x}_1, \pi_{SEAL})$
- 3) Set $\vec{x}_2 := (rt, c, t)$
- 4) Compute $b_2 := \text{SCIP.Verify}(vk_{POST}, \vec{x}_2, \pi_{POST})$
- 5) Output $b_1 \wedge b_2$

Figure 1: PoRep & PoSt protocols (Source [1])

1.3. Filecoin Protocol Overview

Filecoin protocol consists of Storage Market and Retrieval Market networks, with Client, Storage Miner, and Retrieval Miner.

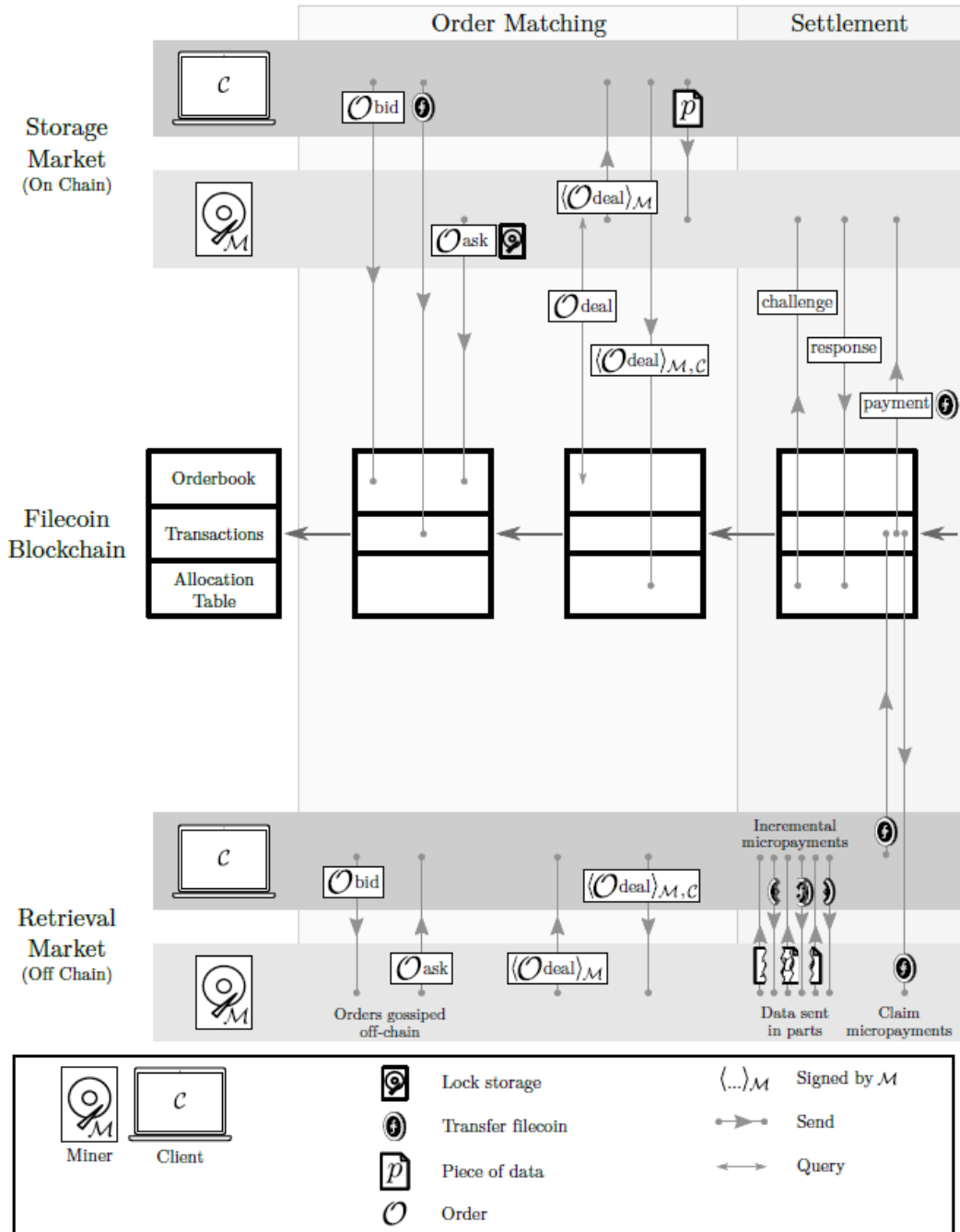


Figure 2: Filecoin Protocol (Source [1])

2. Implementation

In this class project, I plan to implement the simplified Storage Market, where the Client pays the Storage Miner to store their data and the Storage Miner can prove their action with Proof-of-Replication. During the development, some specific parts be simplified for easier implementation.

Put Protocol

Market

AddOrders

- INPUTS: list of orders $\mathcal{O}^1..\mathcal{O}^n$
 - OUTPUTS: bit b , equals 1 if successful
- 1) Set $\mathbf{tx}_{\text{order}} := (\mathcal{O}^1, \dots, \mathcal{O}^n)$
 - 2) Submit $\mathbf{tx}_{\text{order}}$ to \mathcal{L}
 - 3) Wait for $\mathbf{tx}_{\text{order}}$ to be included in \mathcal{L}
 - 4) Output 1 on success, 0 otherwise

MatchOrders

- INPUTS:
 - the current Storage Market OrderBook
 - query order to match \mathcal{O}^q
 - OUTPUTS: matching orders $\mathcal{O}^1..\mathcal{O}^n$
- 1) Match each \mathcal{O}^i in OrderBook such that:
 - a) If \mathcal{O}^q is an *ask* order:
 - i) Check if \mathcal{O}^i is *bid* order
 - ii) Check $\mathcal{O}^i.\text{price} \geq \mathcal{O}^q.\text{price}$
 - iii) Check $\mathcal{O}^i.\text{size} \leq \mathcal{O}^q.\text{space}$
 - b) If \mathcal{O}^q is a *bid* order:
 - i) Check if \mathcal{O}^i is *ask* order
 - ii) Check $\mathcal{O}^i.\text{price} \leq \mathcal{O}^q.\text{price}$
 - iii) Check $\mathcal{O}^i.\text{space} \geq \mathcal{O}^q.\text{size}$
 - 2) Output matched orders $\mathcal{O}^1..\mathcal{O}^n$

Exchange

SendPiece

- INPUTS:
 - an *ask* order \mathcal{O}_{ask}
 - a *bid* order \mathcal{O}_{bid}
 - a piece p
 - OUTPUTS: a *deal* order $\mathcal{O}_{\text{deal}}$ signed by \mathcal{M}_i
- 1) Get identity of \mathcal{M}_i from \mathcal{O}_{ask} signature
 - 2) Send $(\mathcal{O}_{\text{ask}}, \mathcal{O}_{\text{bid}}, p)$ to \mathcal{M}_i
 - 3) Receive $\mathcal{O}_{\text{deal}}$ signed by \mathcal{M}_i
 - 4) Check if $\mathcal{O}_{\text{deal}}$ is valid according to Definition 5.2
 - 5) Output $\mathcal{O}_{\text{deal}}$

ReceivePiece

- INPUTS:
 - signing key for \mathcal{M}_j .
 - current orderbook OrderBook
 - *ask* order \mathcal{O}_{ask}
 - *bid* order \mathcal{O}_{bid}
 - piece p
 - OUTPUTS: *deal* order $\mathcal{O}_{\text{deal}}$ signed by \mathcal{C}_i and \mathcal{M}_j
- 1) Check if \mathcal{O}_{bid} is valid:
 - a) Check if \mathcal{O}_{bid} is in OrderBook
 - b) Check if \mathcal{O}_{bid} is not referenced by other active $\mathcal{O}_{\text{deal}}$
 - c) Check if $\mathcal{O}_{\text{bid}}.\text{size}$ is equal to $|p|$
 - d) Check if \mathcal{O} is signed by \mathcal{M}_i
 - 2) Store p locally
 - 3) Set $\mathcal{O}_{\text{deal}} := (\mathcal{O}_{\text{ask}}, \mathcal{O}_{\text{bid}}, \mathcal{H}(p))_{\mathcal{M}_i}$
 - 4) Get identity of \mathcal{C}_j from \mathcal{O}_{bid}
 - 5) Send $\mathcal{O}_{\text{deal}}$ to \mathcal{C}_j
 - 6) Output $\mathcal{O}_{\text{deal}}$

Figure 3: Storage Market Protocol (Source [1])

Manage Protocol

Network

AssignOrders

- INPUTS:
 - deal orders $\mathcal{O}_{\text{deal}}^1 \dots \mathcal{O}_{\text{deal}}^n$
 - allocation table allocTable
 - OUTPUTS: updated allocation table $\text{allocTable}'$
- 1) Copy allocTable in $\text{allocTable}'$
 - 2) For each order $\mathcal{O}_{\text{deal}}^i$:
 - a) Check if $\mathcal{O}_{\text{deal}}^i$ is valid according to Definition 5.2
 - b) Get \mathcal{M}_j from $\mathcal{O}_{\text{deal}}^i$ signature
 - c) Add details from $\mathcal{O}_{\text{deal}}^i$ to $\text{allocTable}'$
 - 3) Output $\text{allocTable}'$

RepairOrders

- INPUTS:
 - current time t
 - current ledger \mathcal{L}
 - table of storage allocations allocTable
 - OUTPUTS: orders to repair $\mathcal{O}_{\text{deal}}^1 \dots \mathcal{O}_{\text{deal}}^n$, updated allocation table allocTable
- 1) For each allocEntry in allocTable :
 - a) If $t < \text{allocEntry.last} + \Delta_{\text{proof}}$: skip
 - b) Update $\text{allocEntry.last} = t$
 - c) Check if π is in $\mathcal{L}_{t-\Delta_{\text{proof}}:t}$ and $\text{PoSt.Verify}(\pi)$
 - d) On success: update $\text{allocEntry.missing} = 0$
 - e) On failure:
 - i) update $\text{allocEntry.missing}++$
 - ii) penalize collateral from \mathcal{M}_i 's pledge
 - f) If $\text{allocEntry.missing} > \Delta_{\text{fault}}$ then set all the orders from the current sector as failed orders
 - 2) Output failed orders $\mathcal{O}_{\text{deal}}^1 \dots \mathcal{O}_{\text{deal}}^n$ and $\text{allocTable}'$.

Miner

PledgeSector

- INPUTS:
 - current allocation table allocTable
 - pledge request pledge
 - OUTPUTS: $\text{allocTable}'$
- 1) Copy allocTable to $\text{allocTable}'$
 - 2) Set $\text{tx}_{\text{pledge}} := (\text{pledge})$
 - 3) Submit $\text{tx}_{\text{pledge}}$ to \mathcal{L}
 - 4) Wait for $\text{tx}_{\text{pledge}}$ to be included in \mathcal{L}
 - 5) Add new sector of size pledge.size in $\text{allocTable}'$
 - 6) Output $\text{allocTable}'$

SealSector

- INPUTS:
 - miner public/private key pair \mathcal{M}
 - sector index j
 - allocation table allocTable
 - OUTPUTS: a proof π_{SEAL} , a root hash rt
- 1) Find all the pieces $p_1 \dots p_n$ in sector S_j in the allocTable
 - 2) Set $\mathcal{D} := p_1 | p_2 | \dots | p_n$
 - 3) Compute $(\mathcal{R}, \text{rt}, \pi_{\text{SEAL}}) := \text{PoSt.Setup}(\mathcal{M}, \text{pk}_{\text{SEAL}}, \mathcal{D})$
 - 4) Output $\pi_{\text{SEAL}}, \text{rt}$

ProveSector

- INPUTS:
 - miner public/private key pair \mathcal{M}
 - sector index j
 - challenge c
 - OUTPUTS: a proof π_{POS}
- 1) Find \mathcal{R} for sector j
 - 2) Compute $\pi_{\text{POST}} := \text{PoSt.Prove}(\text{pk}_{\text{POST}}, \mathcal{R}, c, \Delta_{\text{proof}})$
 - 3) Output π_{POST}

Figure 4: Manage Protocol (Source [1])

References

1. Filecoin: A decentralized storage network. Technical report, Protocol Labs, August 2017.