

# Graph algorithms

- **Graph Traversal (Graph Searching)**

- Breadth-first search
- Depth-first search

- **Shortest-Path Algorithm**

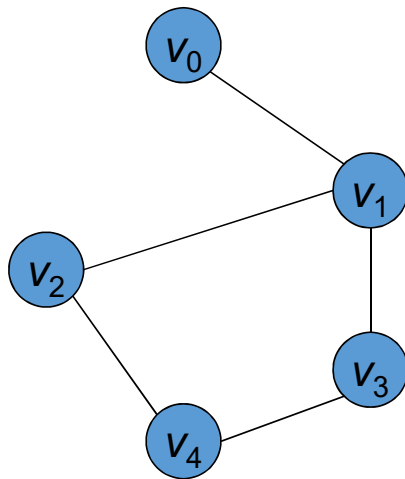
- Dijkstra's algorithm

- **Minimum Spanning Tree**

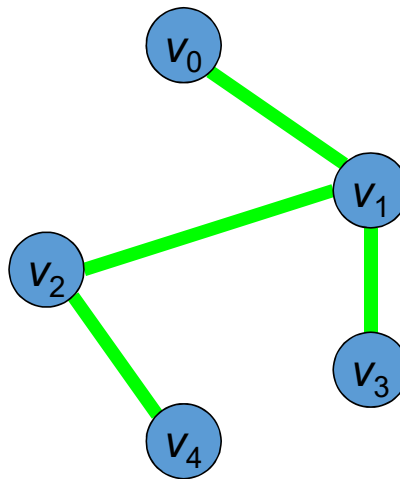
- Prim's Algorithm
- Kruskal's Algorithm

# Spanning Tree

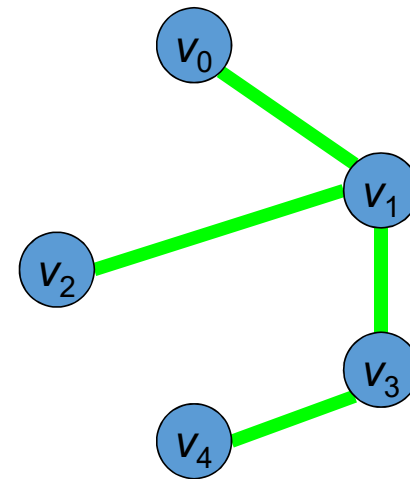
- A **spanning tree** (ST) of an *undirected* graph is a *tree* which contains *all vertices* and *some edges* of the graph.



Graph  $G$



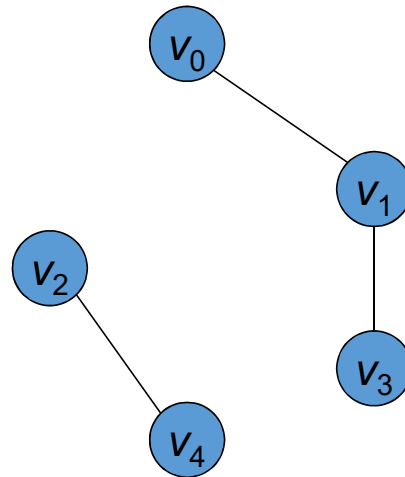
A spanning tree of  $G$



*Another* spanning tree of  $G$

# Spanning Tree and Connectivity

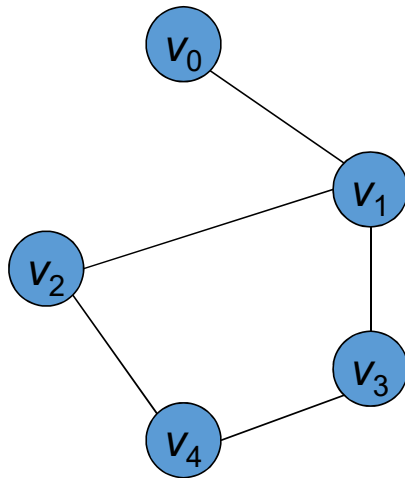
- An *undirected* graph has a spanning tree if and only if the graph is *connected*.



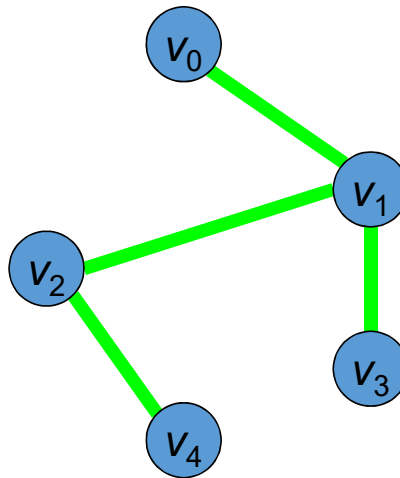
A graph that has *no* spanning trees

# Spanning Tree

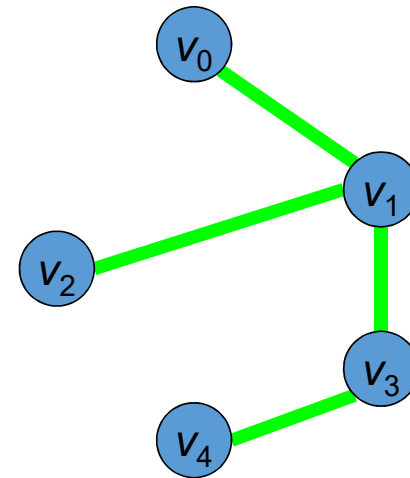
- For a graph with  $n$  vertices, its spanning tree always has exactly  $n - 1$  edges.



Graph  $G$



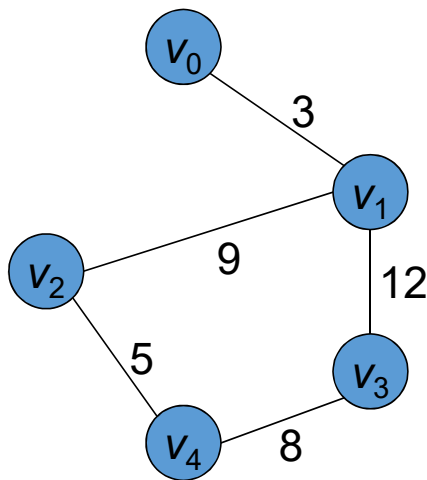
A spanning tree of  $G$



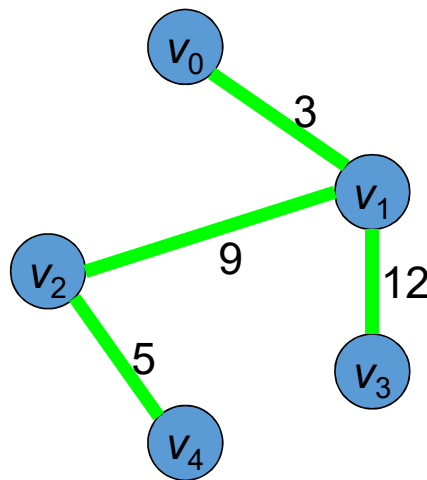
*Another* spanning tree of  $G$

# Minimum Spanning Tree

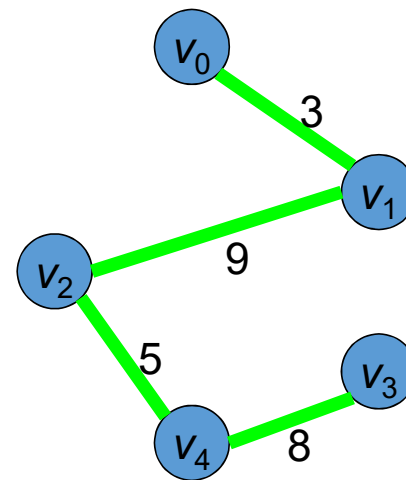
- A **minimum spanning tree** (MST) of an **undirected weighted** graph is a spanning tree whose sum of all weights is minimum.



Graph  $G$



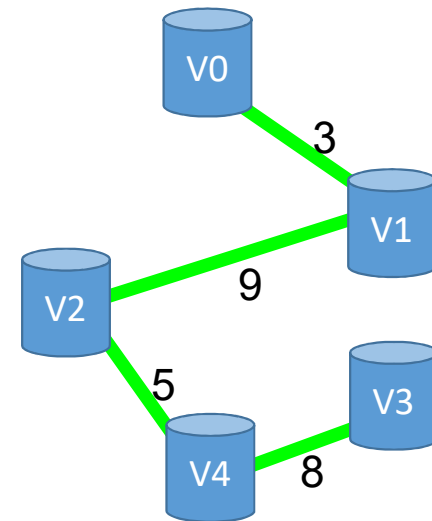
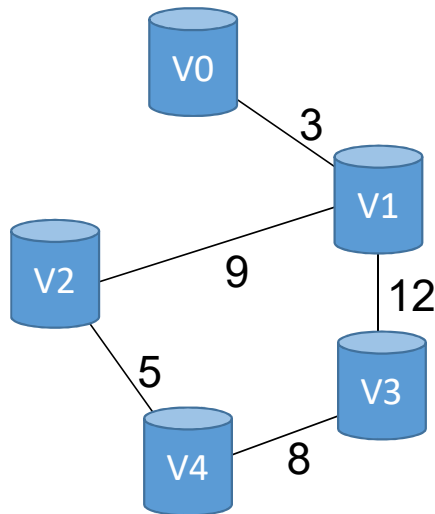
A spanning tree of  $G$   
(total weight =  $3+9+12+5 = 29$ )



A MST of  $G$   
(total weight =  $3+9+5+8 = 25$ )

# Application : Minimum Spanning Tree

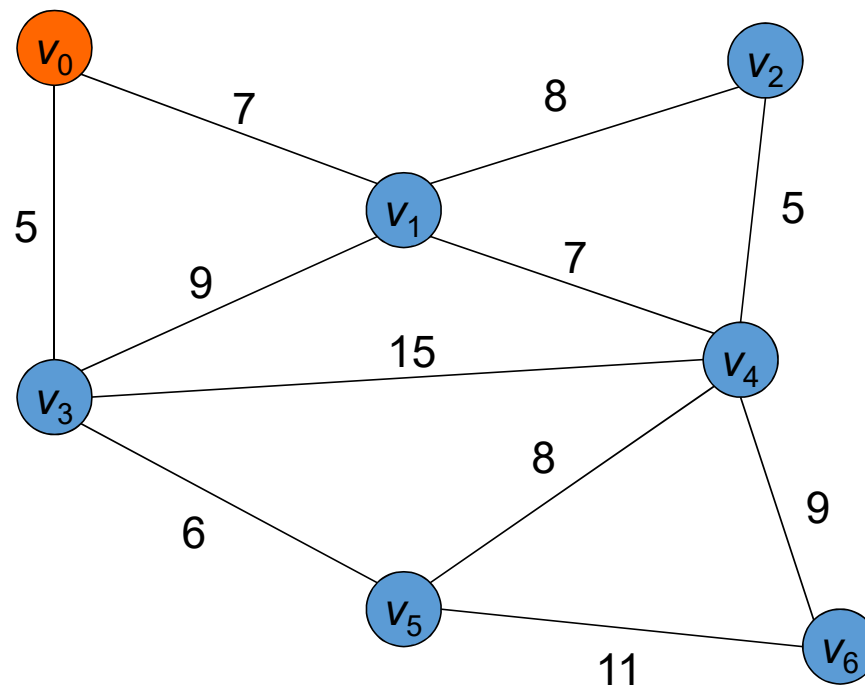
- Find the cheapest route to connect all computers/devices.



# MST Algorithms

- Two common algorithms for finding MSTs.
  - *Prim's algorithm*
    - Build tree to span all vertices
  - *Kruskal's algorithm*
    - From “forest” to tree

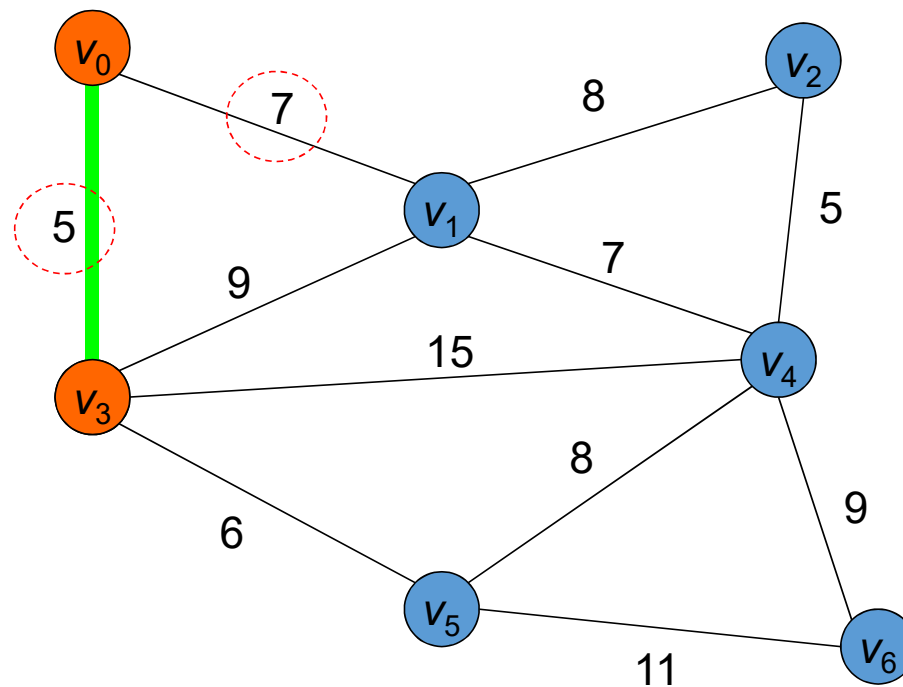
# Prim's Algorithm



Start from any one vertex, say,  $v_0$ .  
(Consider  $v_0$  as a tree with one node.)

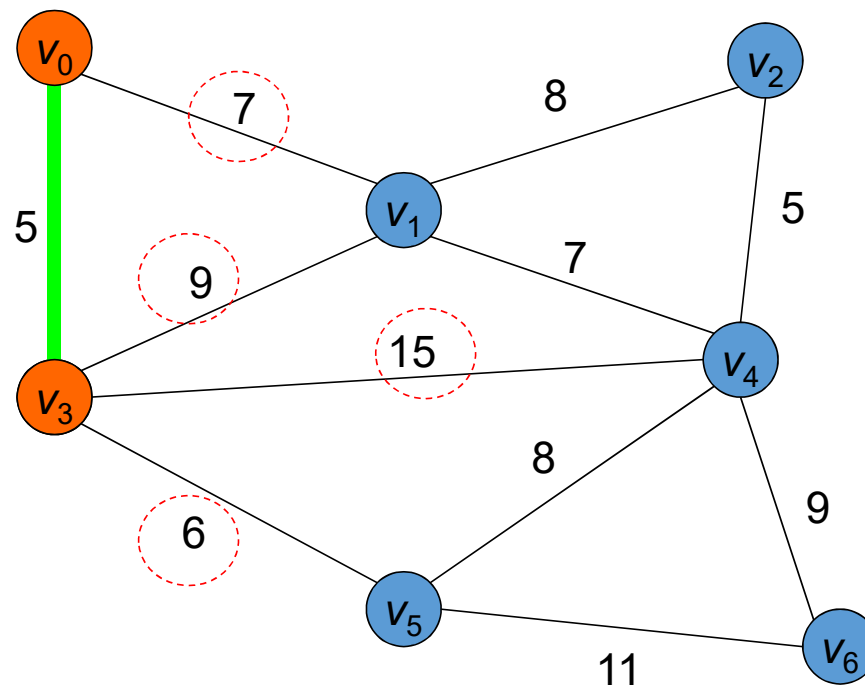


# Prim's Algorithm



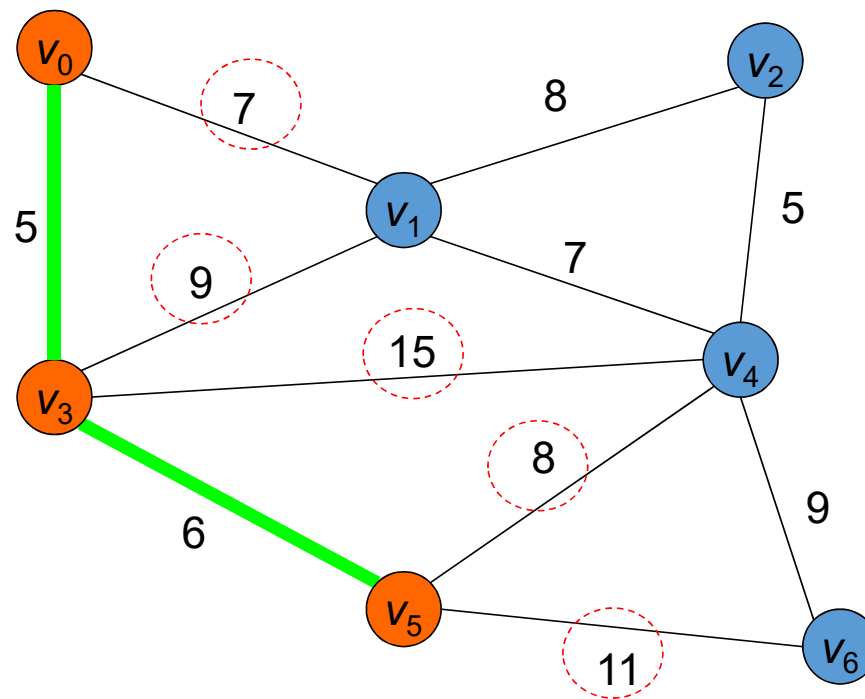
Pick the edge around the current tree whose weight is the smallest and link to an unvisited node.

# Prim's Algorithm



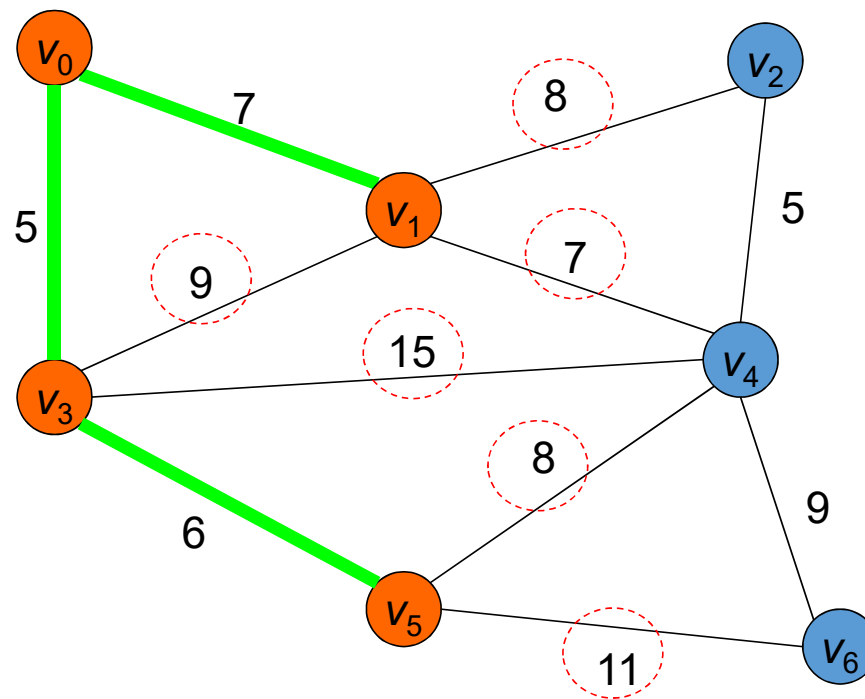
Pick the edge around the current tree whose weight is the smallest and link to an unvisited node.

# Prim's Algorithm



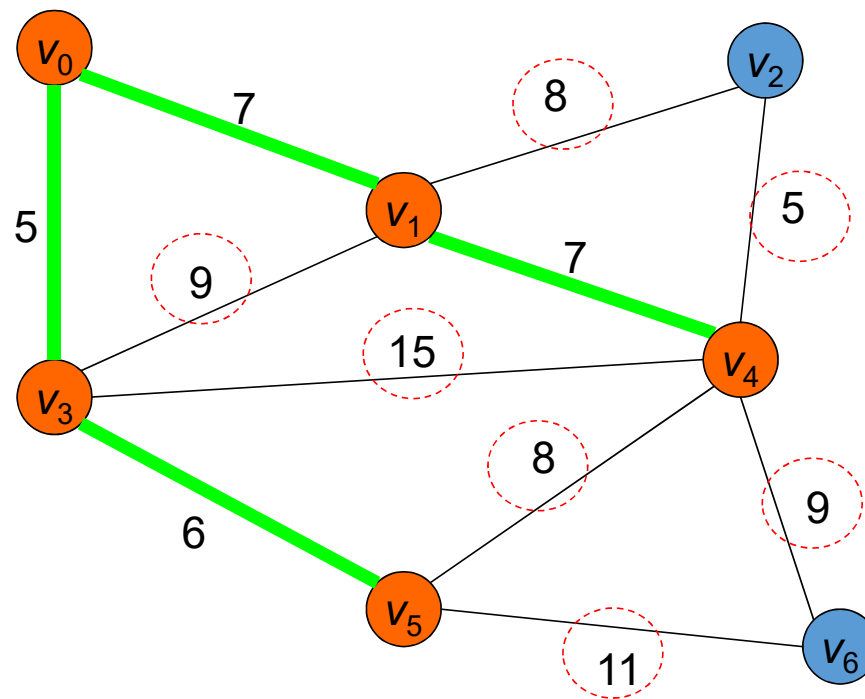
Pick the edge around the current tree whose weight is the smallest and link to an unvisited node.

# Prim's Algorithm



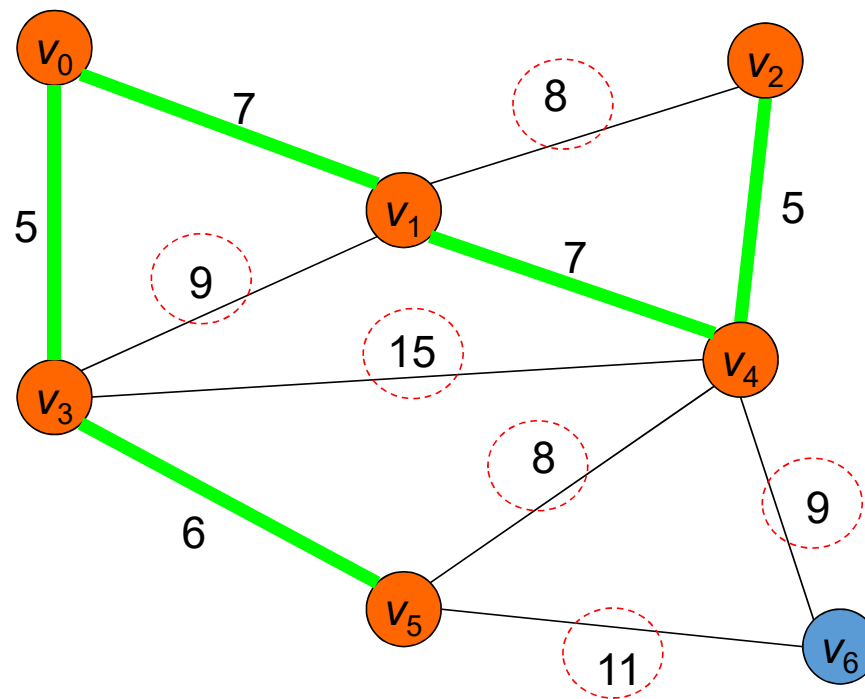
Pick the edge around the current tree whose weight is the smallest and link to an unvisited node.

# Prim's Algorithm



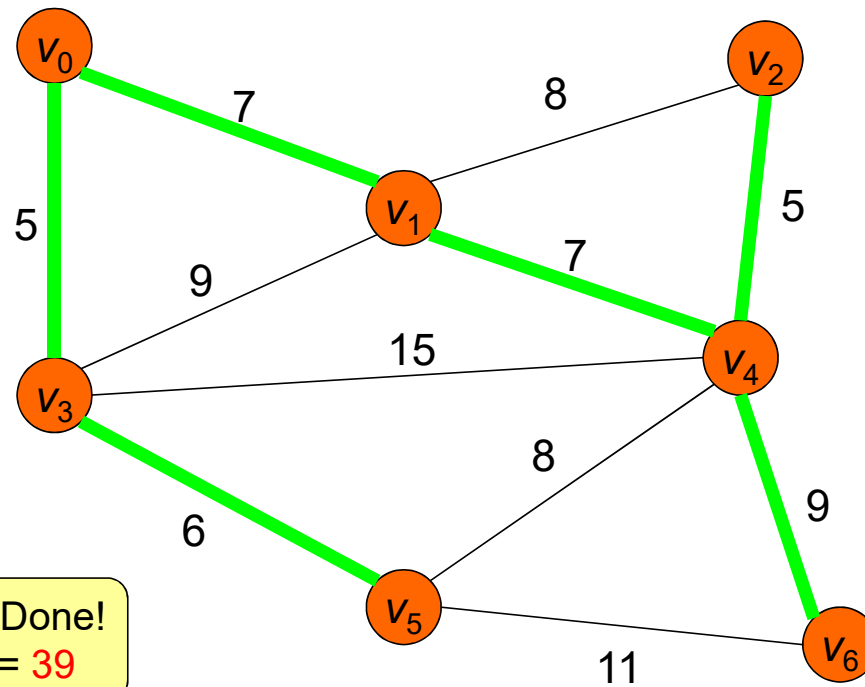
Pick the edge around the current tree whose weight is the smallest and link to an unvisited node.

# Prim's Algorithm



Pick the edge around the current tree whose weight is the smallest and link to an unvisited node.

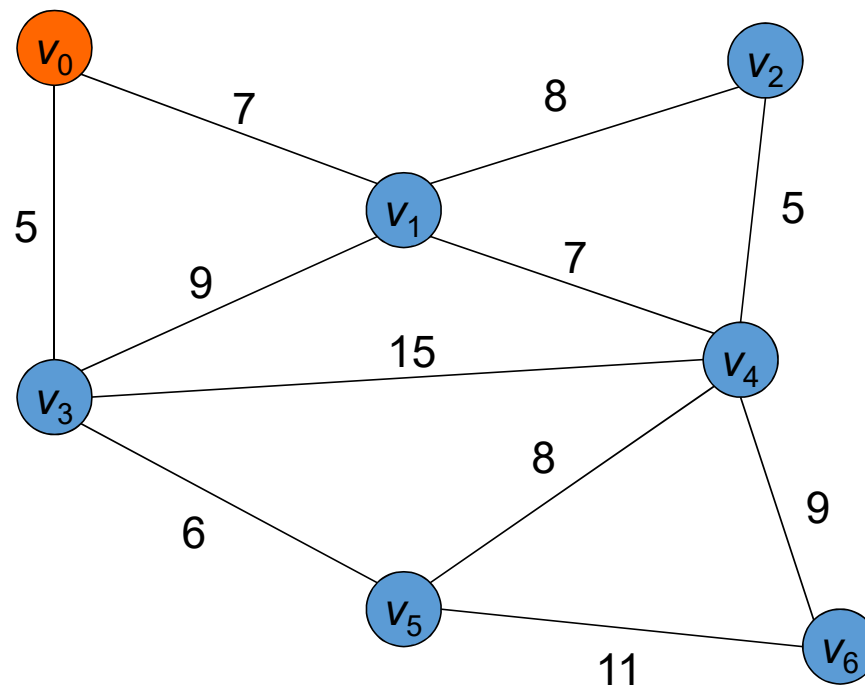
# Prim's Algorithm



Spans all vertices. Done!  
Total tree weight = 39

Pick the edge around the current tree whose weight is the smallest and link to an unvisited node.

# Prim's Algorithm

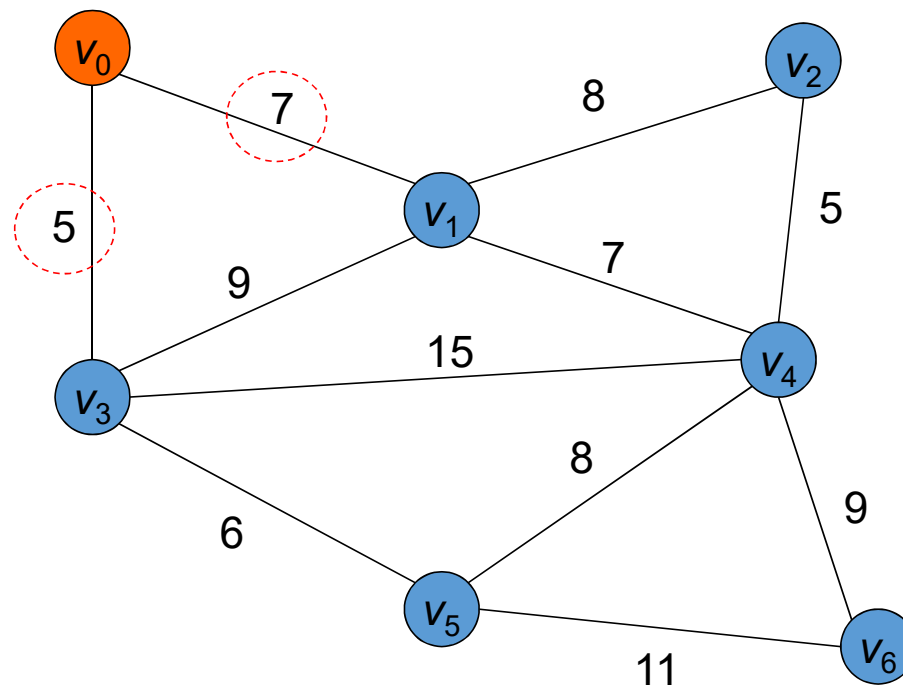


| vertex | distance | previous |
|--------|----------|----------|
| $v_0$  | 0        | 0        |
| $v_1$  | $\infty$ | 0        |
| $v_2$  | $\infty$ | 0        |
| $v_3$  | $\infty$ | 0        |
| $v_4$  | $\infty$ | 0        |
| $v_5$  | $\infty$ | 0        |
| $v_6$  | $\infty$ | 0        |

Start from any one vertex, say,  $v_0$ .  
(Consider  $v_0$  as a tree with one node.)



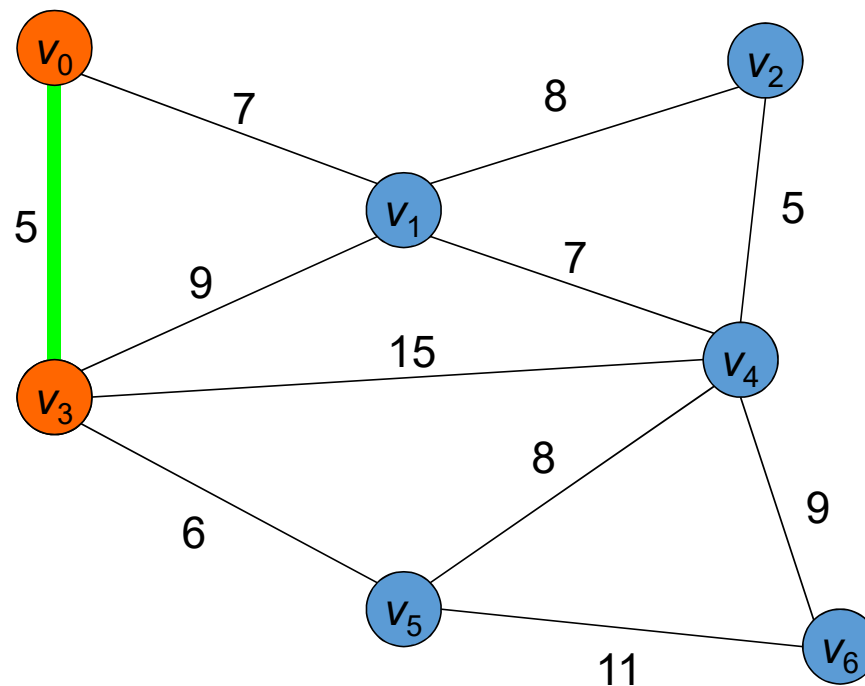
# Prim's Algorithm



| vertex | distance | previous |
|--------|----------|----------|
| $v_0$  | 0        | 0        |
| $v_1$  | 7        | $v_0$    |
| $v_2$  | $\infty$ | 0        |
| $v_3$  | 5        | $v_0$    |
| $v_4$  | $\infty$ | 0        |
| $v_5$  | $\infty$ | 0        |
| $v_6$  | $\infty$ | 0        |

Update the distance to all nodes adjacent to  $v_0$

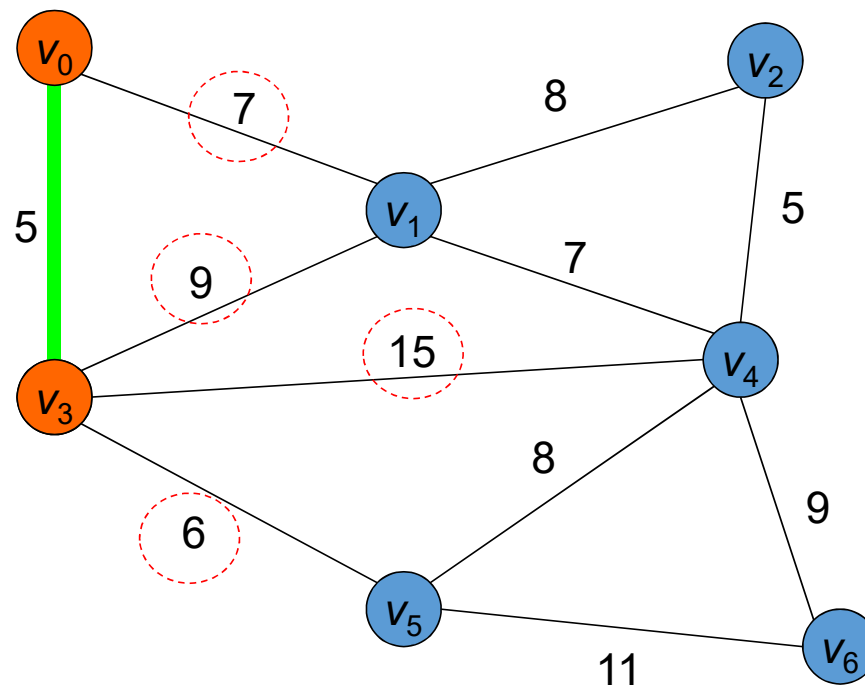
# Prim's Algorithm



| vertex | distance | previous |
|--------|----------|----------|
| $v_0$  | 0        | 0        |
| $v_1$  | 7        | $v_0$    |
| $v_2$  | $\infty$ | 0        |
| $v_3$  | 5        | $v_0$    |
| $v_4$  | $\infty$ | 0        |
| $v_5$  | $\infty$ | 0        |
| $v_6$  | $\infty$ | 0        |

Pick the edge around the current tree whose weight is the smallest and link to an unvisited node.

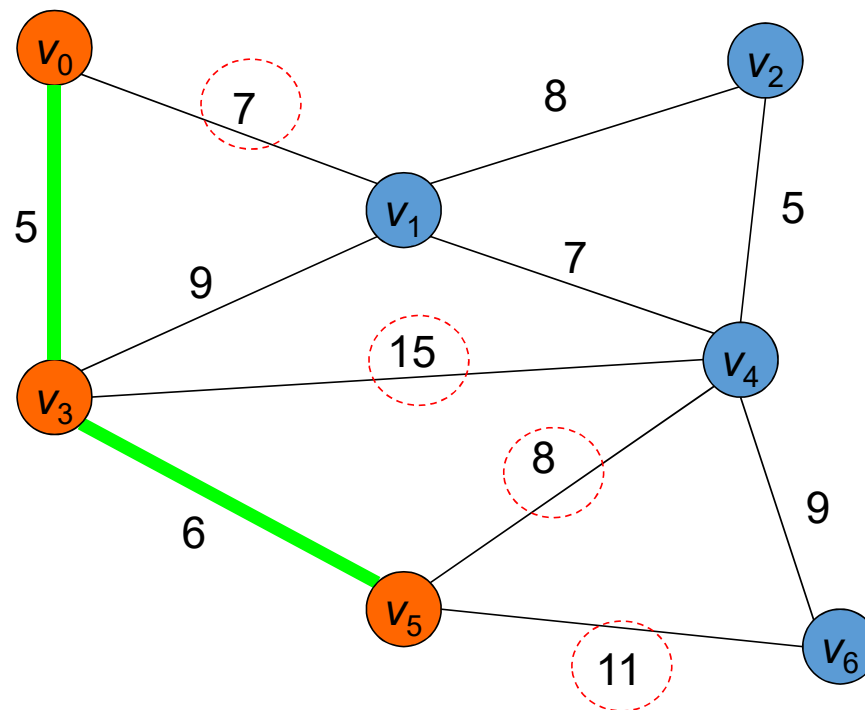
# Prim's Algorithm



| vertex | distance | previous |
|--------|----------|----------|
| $v_0$  | 0        | 0        |
| $v_1$  | Min(7,9) | $v_0$    |
| $v_2$  | $\infty$ | 0        |
| $v_3$  | 5        | $v_0$    |
| $v_4$  | 15       | $v_3$    |
| $v_5$  | 6        | $v_3$    |
| $v_6$  | $\infty$ | 0        |

Update the distance to all nodes adjacent to  $v_0$

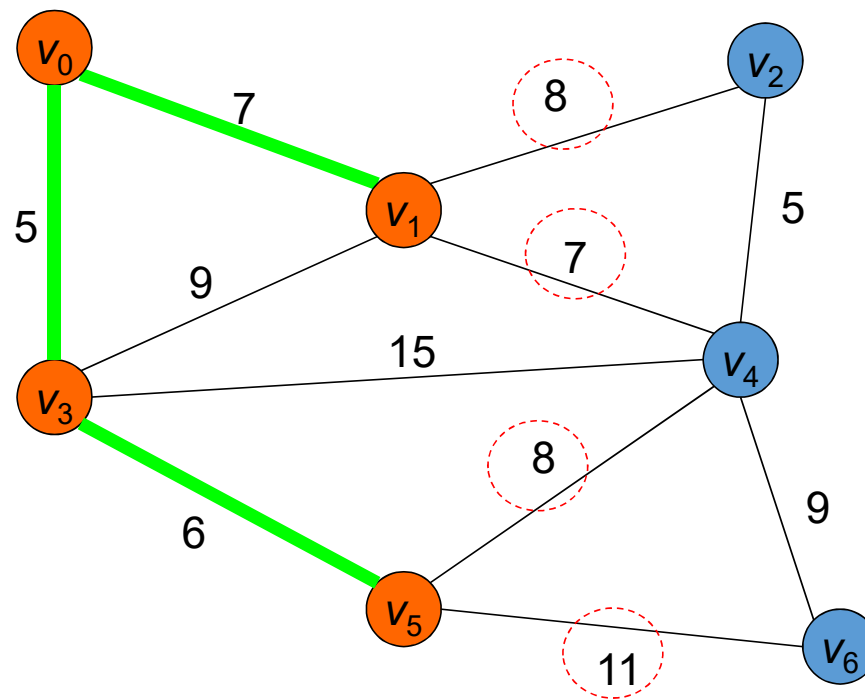
# Prim's Algorithm



| vertex | distance  | previous |
|--------|-----------|----------|
| $v_0$  | 0         | 0        |
| $v_1$  | 7         | $v_0$    |
| $v_2$  | $\infty$  | 0        |
| $v_3$  | 5         | $v_0$    |
| $v_4$  | Min(15,8) | $v_5$    |
| $v_5$  | 6         | $v_3$    |
| $v_6$  | 11        | $v_5$    |

Pick the edge around the current tree whose weight is the smallest and link to an unvisited node.

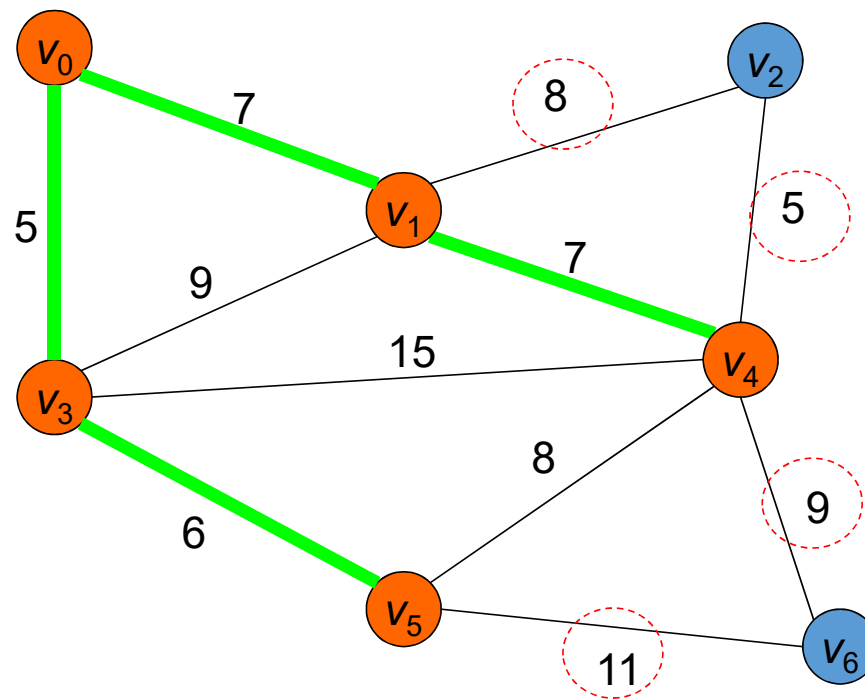
# Prim's Algorithm



| vertex | distance | previous |
|--------|----------|----------|
| $v_0$  | 0        | 0        |
| $v_1$  | 7        | $v_0$    |
| $v_2$  | 8        | $v_1$    |
| $v_3$  | 5        | $v_0$    |
| $v_4$  | Min(8,7) | $v_1$    |
| $v_5$  | 6        | $v_3$    |
| $v_6$  | 11       | $v_5$    |

Pick the edge around the current tree whose weight is the smallest and link to an unvisited node.

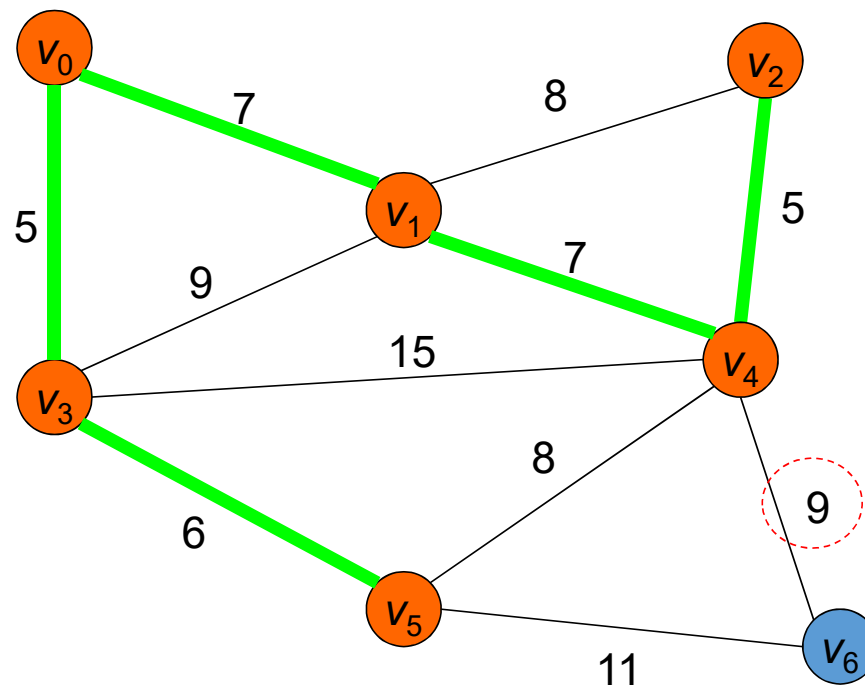
# Prim's Algorithm



| vertex | distance  | previous |
|--------|-----------|----------|
| $v_0$  | 0         | 0        |
| $v_1$  | 7         | $v_0$    |
| $v_2$  | Min(8,5)  | $v_4$    |
| $v_3$  | 5         | $v_0$    |
| $v_4$  | 7         | $v_1$    |
| $v_5$  | 6         | $v_3$    |
| $v_6$  | Min(11,9) | $v_4$    |

Pick the edge around the current tree whose weight is the smallest and link to an unvisited node.

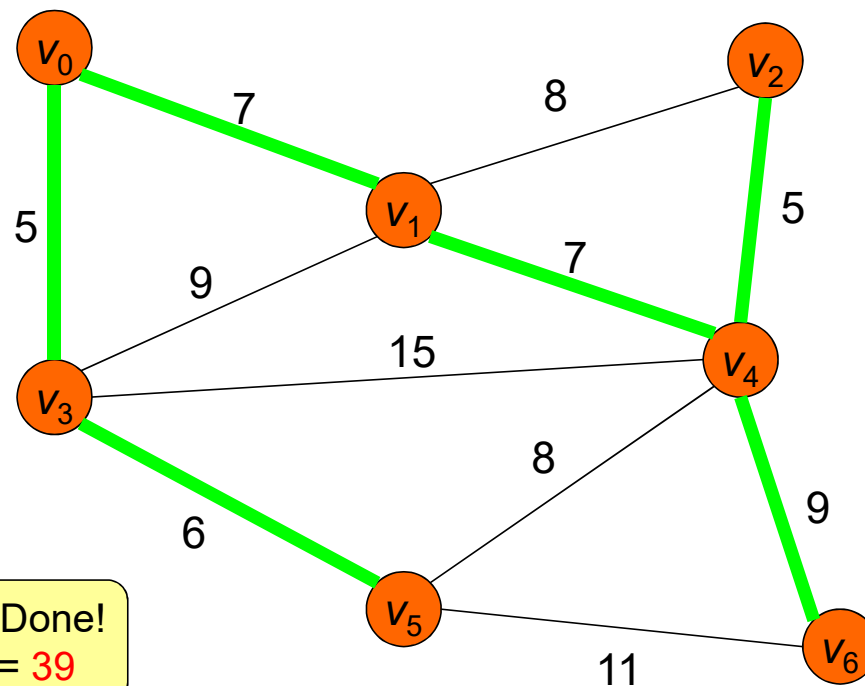
# Prim's Algorithm



| vertex | distance | previous |
|--------|----------|----------|
| $v_0$  | 0        | 0        |
| $v_1$  | 7        | $v_0$    |
| $v_2$  | 5        | $v_4$    |
| $v_3$  | 5        | $v_0$    |
| $v_4$  | 7        | $v_1$    |
| $v_5$  | 6        | $v_3$    |
| $v_6$  | 9        | $v_4$    |

Pick the edge around the current tree whose weight is the smallest and link to an unvisited node.

# Prim's Algorithm



| vertex | distance | previous |
|--------|----------|----------|
| $v_0$  | 0        | 0        |
| $v_1$  | 7        | $v_0$    |
| $v_2$  | 5        | $v_4$    |
| $v_3$  | 5        | $v_0$    |
| $v_4$  | 7        | $v_1$    |
| $v_5$  | 6        | $v_3$    |
| $v_6$  | 9        | $v_4$    |

Spans all vertices. Done!  
Total tree weight = 39

Pick the edge around the current tree whose weight is the smallest and link to an unvisited node.



# Prim's Algorithm

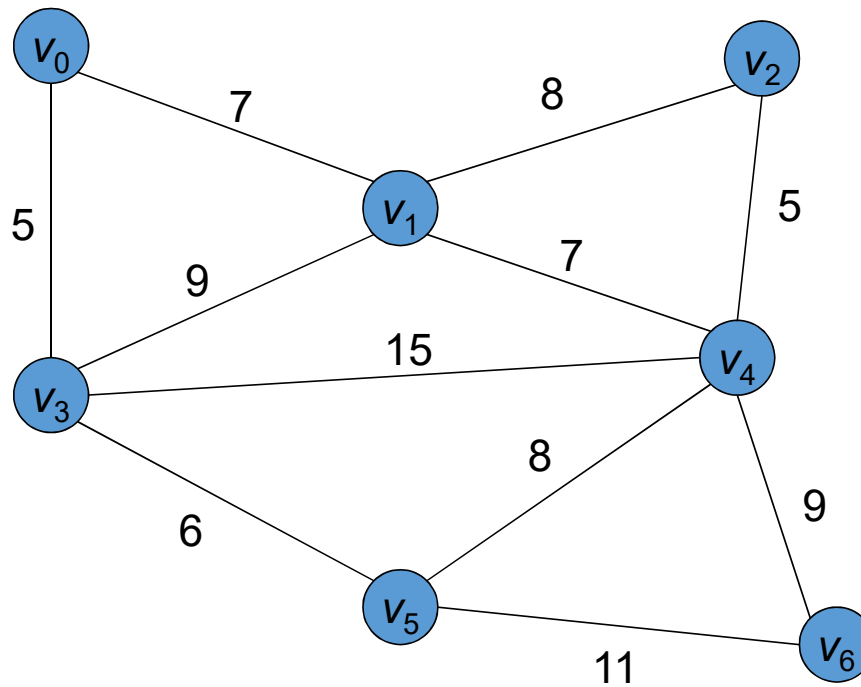
- Initialize the current tree  $T$  to have any one vertex.
- While  $T$  has fewer than  $n - 1$  edges
  - Pick the edge around  $T$  whose weight is the smallest and link to an unvisited node.
  - (When no such edge is found, graph is disconnected and no MST exists.)

# Kruskal's Algorithm

Increasingly  
sorted by weights

| Edge         | Weight |
|--------------|--------|
| $(v_0, v_3)$ | 5      |
| $(v_2, v_4)$ | 5      |
| $(v_3, v_5)$ | 6      |
| $(v_0, v_1)$ | 7      |
| $(v_1, v_4)$ | 7      |
| $(v_1, v_2)$ | 8      |
| $(v_4, v_5)$ | 8      |
| $(v_1, v_3)$ | 9      |
| $(v_4, v_6)$ | 9      |
| $(v_5, v_6)$ | 11     |
| $(v_3, v_4)$ | 15     |

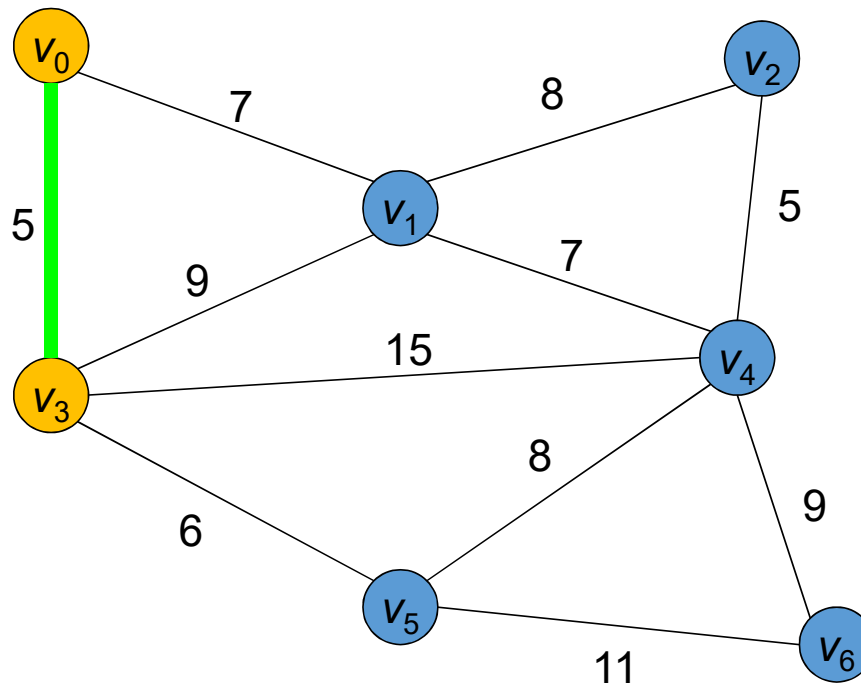
Pick edge one by one



# Kruskal's Algorithm

| Edge         | Weight |
|--------------|--------|
| $(v_0, v_3)$ | 5      |
| $(v_2, v_4)$ | 5      |
| $(v_3, v_5)$ | 6      |
| $(v_0, v_1)$ | 7      |
| $(v_1, v_4)$ | 7      |
| $(v_1, v_2)$ | 8      |
| $(v_4, v_5)$ | 8      |
| $(v_1, v_3)$ | 9      |
| $(v_4, v_6)$ | 9      |
| $(v_5, v_6)$ | 11     |
| $(v_3, v_4)$ | 15     |

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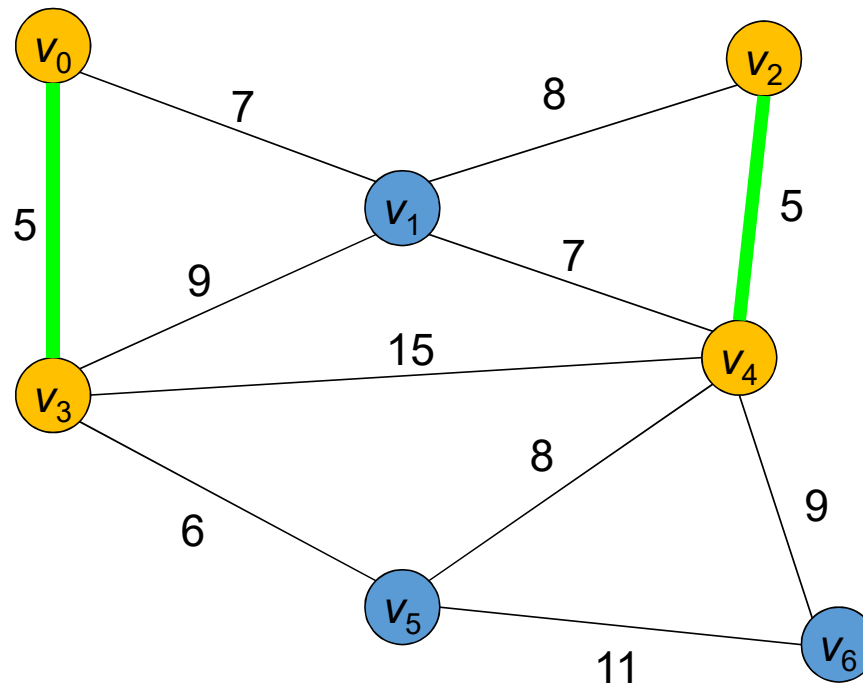


# Kruskal's Algorithm

| Edge         | Weight |
|--------------|--------|
| $(v_0, v_3)$ | 5      |
| $(v_2, v_4)$ | 5      |
| $(v_3, v_5)$ | 6      |
| $(v_0, v_1)$ | 7      |
| $(v_1, v_4)$ | 7      |
| $(v_1, v_2)$ | 8      |
| $(v_4, v_5)$ | 8      |
| $(v_1, v_3)$ | 9      |
| $(v_4, v_6)$ | 9      |
| $(v_5, v_6)$ | 11     |
| $(v_3, v_4)$ | 15     |

✓

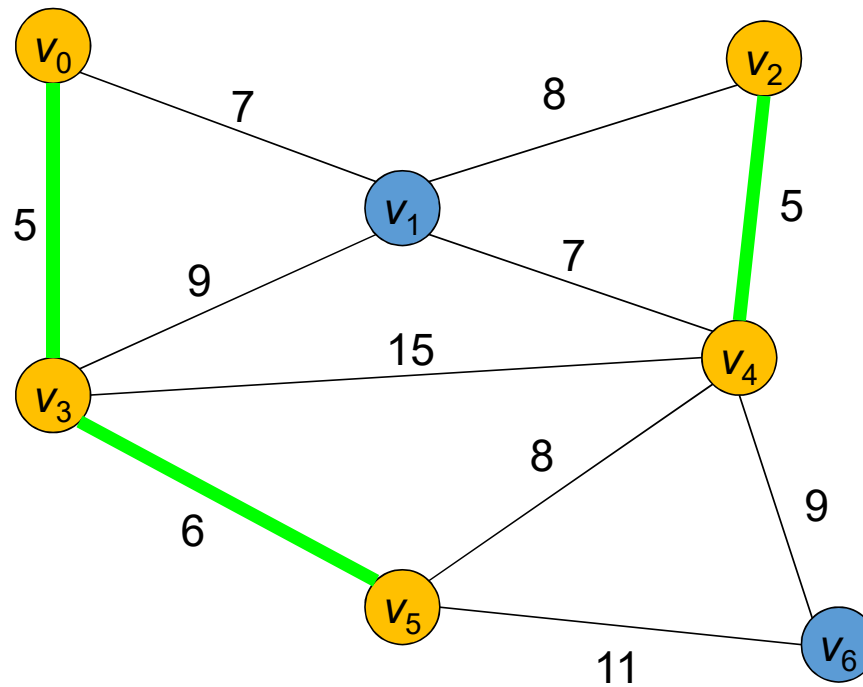
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# Kruskal's Algorithm

| Edge         | Weight |
|--------------|--------|
| $(v_0, v_3)$ | 5      |
| $(v_2, v_4)$ | 5      |
| $(v_3, v_5)$ | 6      |
| $(v_0, v_1)$ | 7      |
| $(v_1, v_4)$ | 7      |
| $(v_1, v_2)$ | 8      |
| $(v_4, v_5)$ | 8      |
| $(v_1, v_3)$ | 9      |
| $(v_4, v_6)$ | 9      |
| $(v_5, v_6)$ | 11     |
| $(v_3, v_4)$ | 15     |

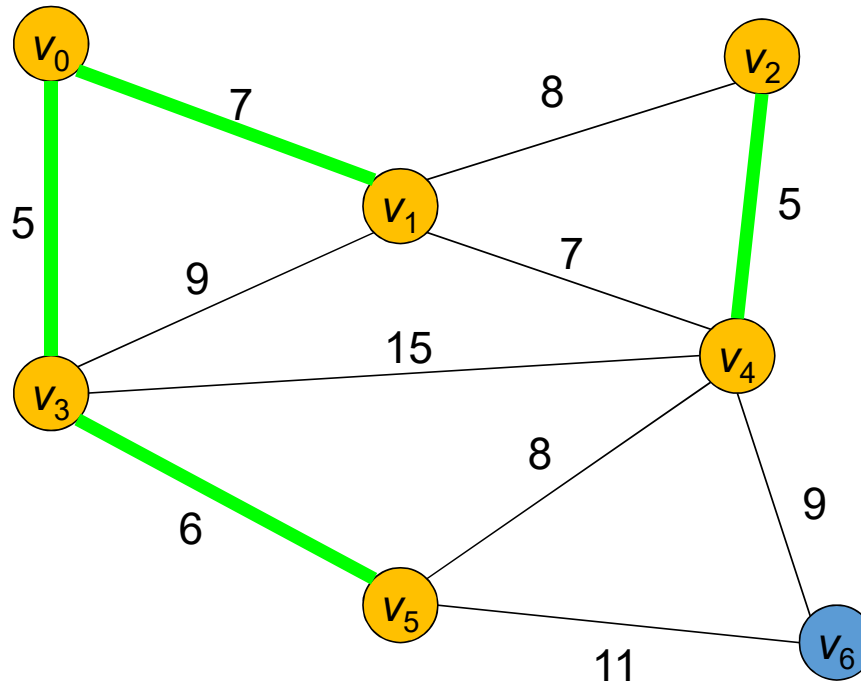
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# Kruskal's Algorithm

| Edge         | Weight |
|--------------|--------|
| $(v_0, v_3)$ | 5      |
| $(v_2, v_4)$ | 5      |
| $(v_3, v_5)$ | 6      |
| $(v_0, v_1)$ | 7      |
| $(v_1, v_4)$ | 7      |
| $(v_1, v_2)$ | 8      |
| $(v_4, v_5)$ | 8      |
| $(v_1, v_3)$ | 9      |
| $(v_4, v_6)$ | 9      |
| $(v_5, v_6)$ | 11     |
| $(v_3, v_4)$ | 15     |

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# Kruskal's Algorithm

| Edge         | Weight |
|--------------|--------|
| $(v_0, v_3)$ | 5      |
| $(v_2, v_4)$ | 5      |
| $(v_3, v_5)$ | 6      |
| $(v_0, v_1)$ | 7      |
| $(v_1, v_4)$ | 7      |
| $(v_1, v_2)$ | 8      |
| $(v_4, v_5)$ | 8      |
| $(v_1, v_3)$ | 9      |
| $(v_4, v_6)$ | 9      |
| $(v_5, v_6)$ | 11     |
| $(v_3, v_4)$ | 15     |

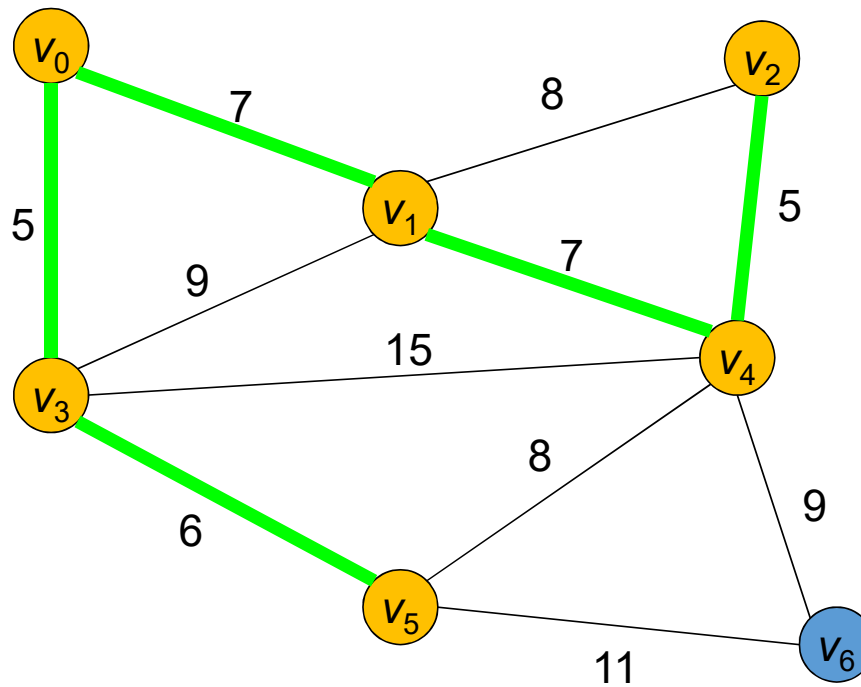
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# Kruskal's Algorithm

| Edge         | Weight |
|--------------|--------|
| $(v_0, v_3)$ | 5      |
| $(v_2, v_4)$ | 5      |
| $(v_3, v_5)$ | 6      |
| $(v_0, v_1)$ | 7      |
| $(v_1, v_4)$ | 7      |
| $(v_1, v_2)$ | 8      |
| $(v_4, v_5)$ | 8      |
| $(v_1, v_3)$ | 9      |
| $(v_4, v_6)$ | 9      |
| $(v_5, v_6)$ | 11     |
| $(v_3, v_4)$ | 15     |

✓

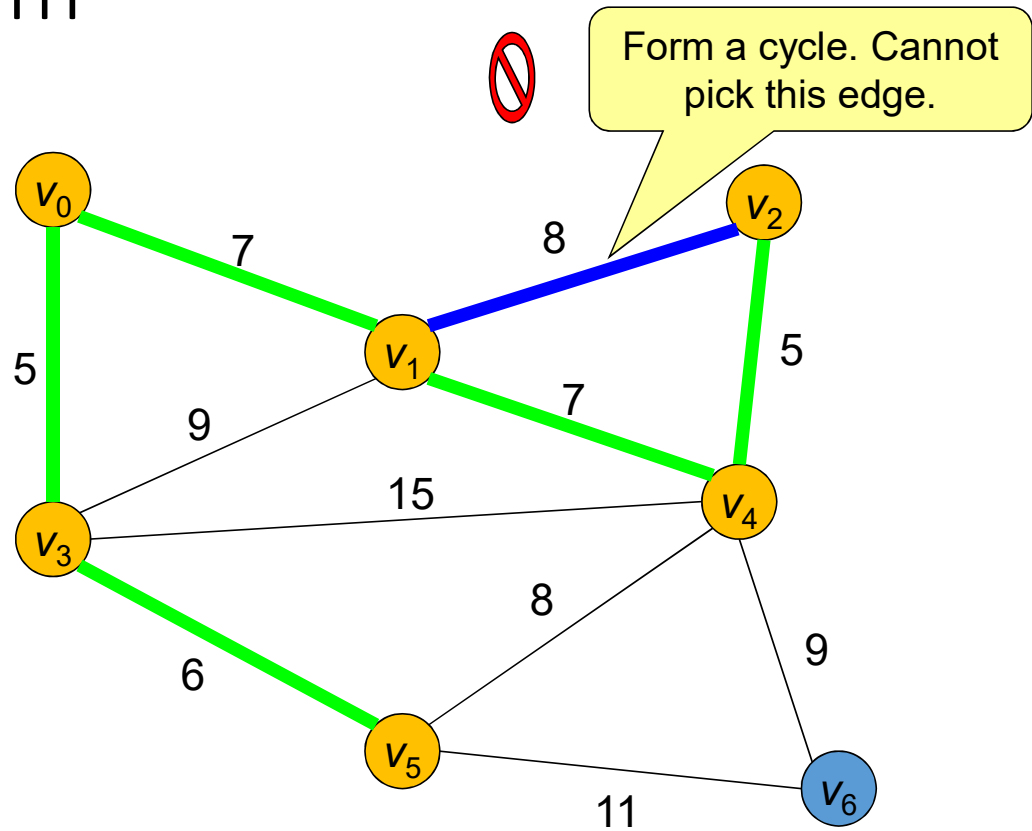
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✗





# Kruskal's Algorithm

| Edge         | Weight |
|--------------|--------|
| $(v_0, v_3)$ | 5      |
| $(v_2, v_4)$ | 5      |
| $(v_3, v_5)$ | 6      |
| $(v_0, v_1)$ | 7      |
| $(v_1, v_4)$ | 7      |
| $(v_1, v_2)$ | 8      |
| $(v_4, v_5)$ | 8      |
| $(v_1, v_3)$ | 9      |
| $(v_4, v_6)$ | 9      |
| $(v_5, v_6)$ | 11     |
| $(v_3, v_4)$ | 15     |

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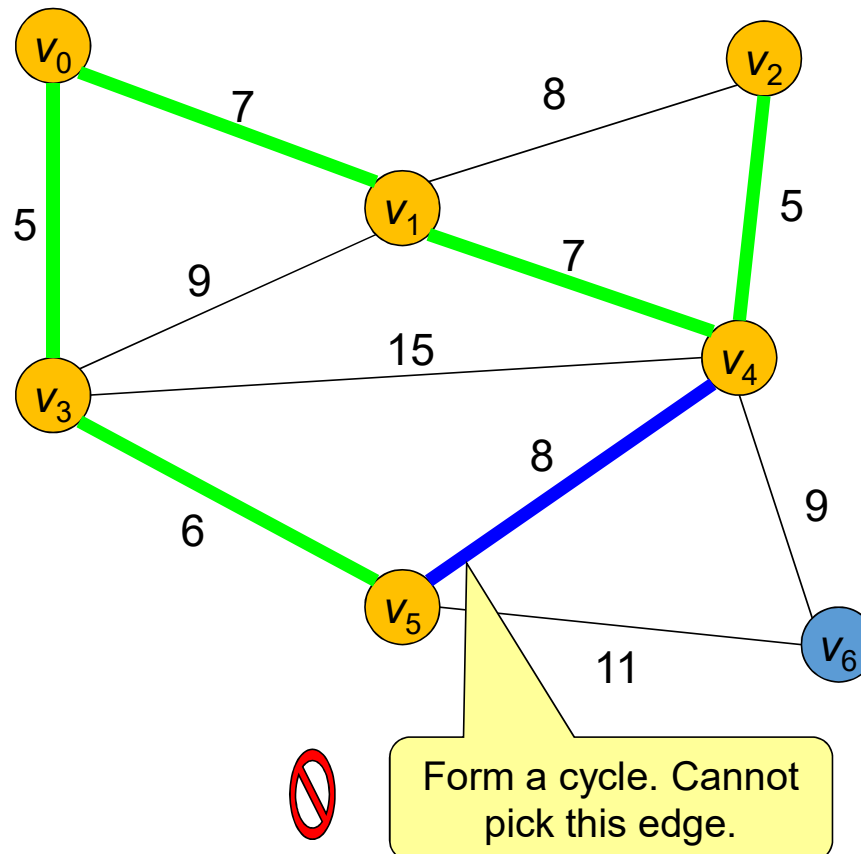
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✗



# Kruskal's Algorithm

| Edge         | Weight |
|--------------|--------|
| $(v_0, v_3)$ | 5      |
| $(v_2, v_4)$ | 5      |
| $(v_3, v_5)$ | 6      |
| $(v_0, v_1)$ | 7      |
| $(v_1, v_4)$ | 7      |
| $(v_1, v_2)$ | 8      |
| $(v_4, v_5)$ | 8      |
| $(v_1, v_3)$ | 9      |
| $(v_4, v_6)$ | 9      |
| $(v_5, v_6)$ | 11     |
| $(v_3, v_4)$ | 15     |

✓

✓

✓

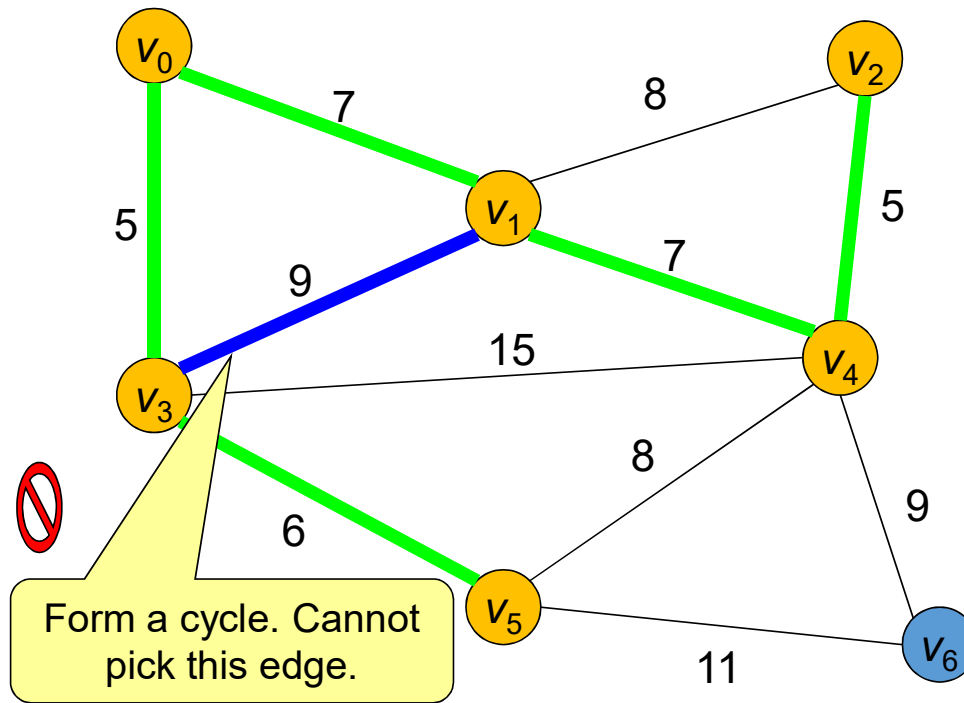
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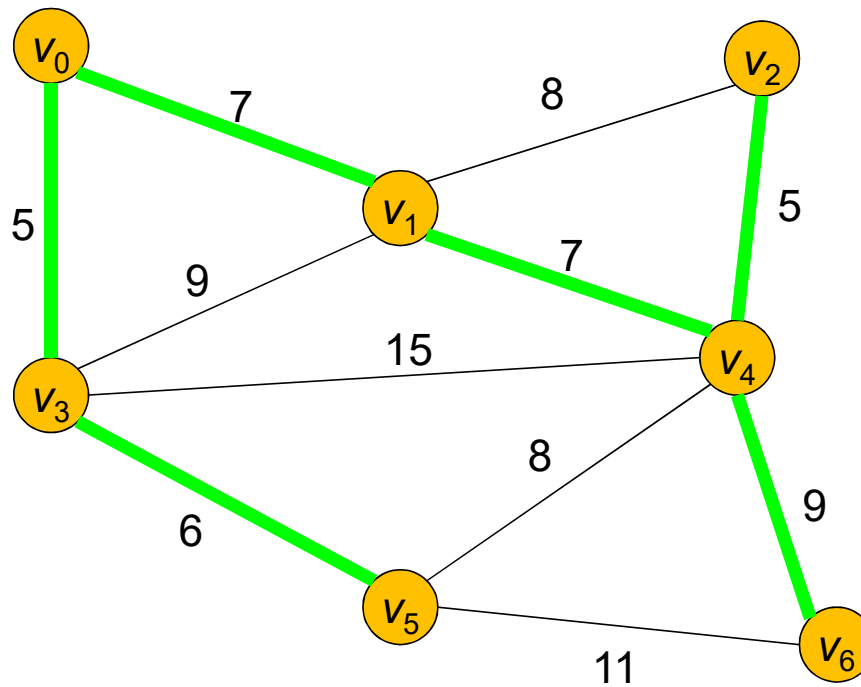
✗

✗



# Kruskal's Algorithm

| Edge         | Weight |   |
|--------------|--------|---|
| $(v_0, v_3)$ | 5      | ✓ |
| $(v_2, v_4)$ | 5      | ✓ |
| $(v_3, v_5)$ | 6      | ✓ |
| $(v_0, v_1)$ | 7      | ✓ |
| $(v_1, v_4)$ | 7      | ✓ |
| $(v_1, v_2)$ | 8      | ✗ |
| $(v_4, v_5)$ | 8      | ✗ |
| $(v_1, v_3)$ | 9      | ✗ |
| $(v_4, v_6)$ | 9      | ✓ |
| $(v_5, v_6)$ | 11     |   |
| $(v_3, v_4)$ | 15     |   |



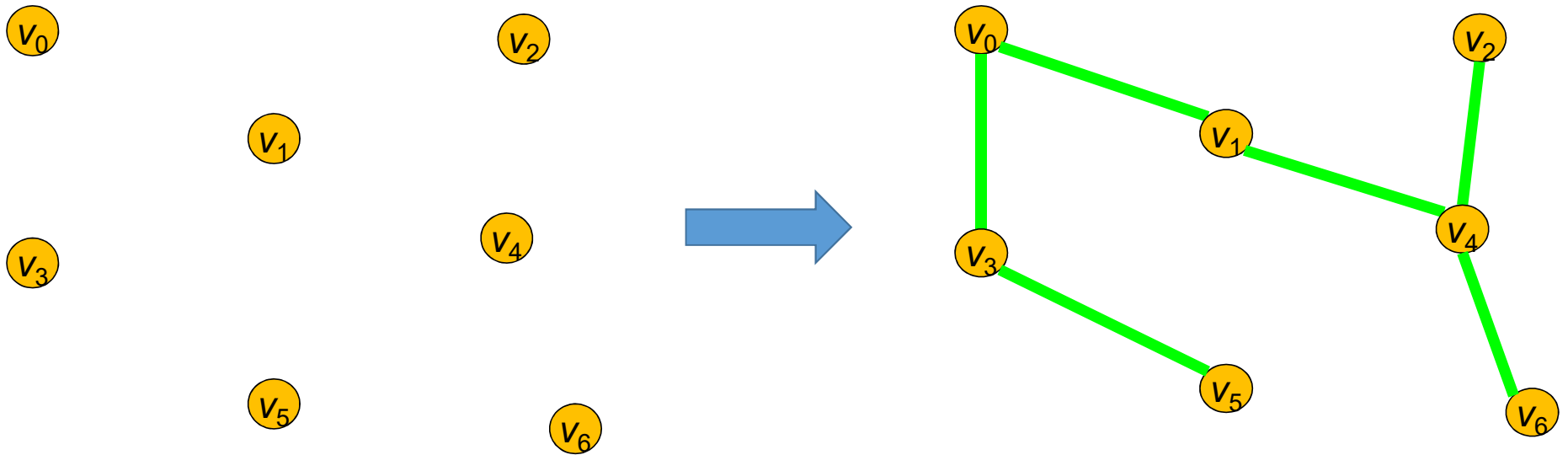
Spans all vertices. Done!  
Total tree weight = 39

# Kruskal's Algorithm

- Increasingly sort the edges by weights.
- For each edge  $e$  in sorted order
  - If  $e$  does not form a cycle with the already picked edges, then
    - Pick  $e$ . (Done when  $n - 1$  edges are picked.)
  - Else
    - Discard  $e$ .
- If fewer than  $n - 1$  edges are picked, then
  - Graph is disconnected. No MST exists.

# Kruskal's Algorithm

- From “forest” to tree



# Kruskal's vs Prim's Algorithms

- Order of edges picked to solution

| <u>Kruskal</u> | <u>Prim</u>  |
|----------------|--------------|
| $(v_0, v_3)$   | $(v_0, v_3)$ |
| $(v_2, v_4)$   | $(v_3, v_5)$ |
| $(v_3, v_5)$   | $(v_0, v_1)$ |
| $(v_0, v_1)$   | $(v_1, v_4)$ |
| $(v_1, v_4)$   | $(v_2, v_4)$ |
| $(v_4, v_6)$   | $(v_4, v_6)$ |

- Kruskal's algorithm repeatedly merges multiple trees (forest) until only one tree is left.
- Prim's algorithm repeatedly grows one tree until all vertices are spanned.