

Tutorial 05: Sorting

CSCI2520 - DATA STRUCTURES AND APPLICATIONS

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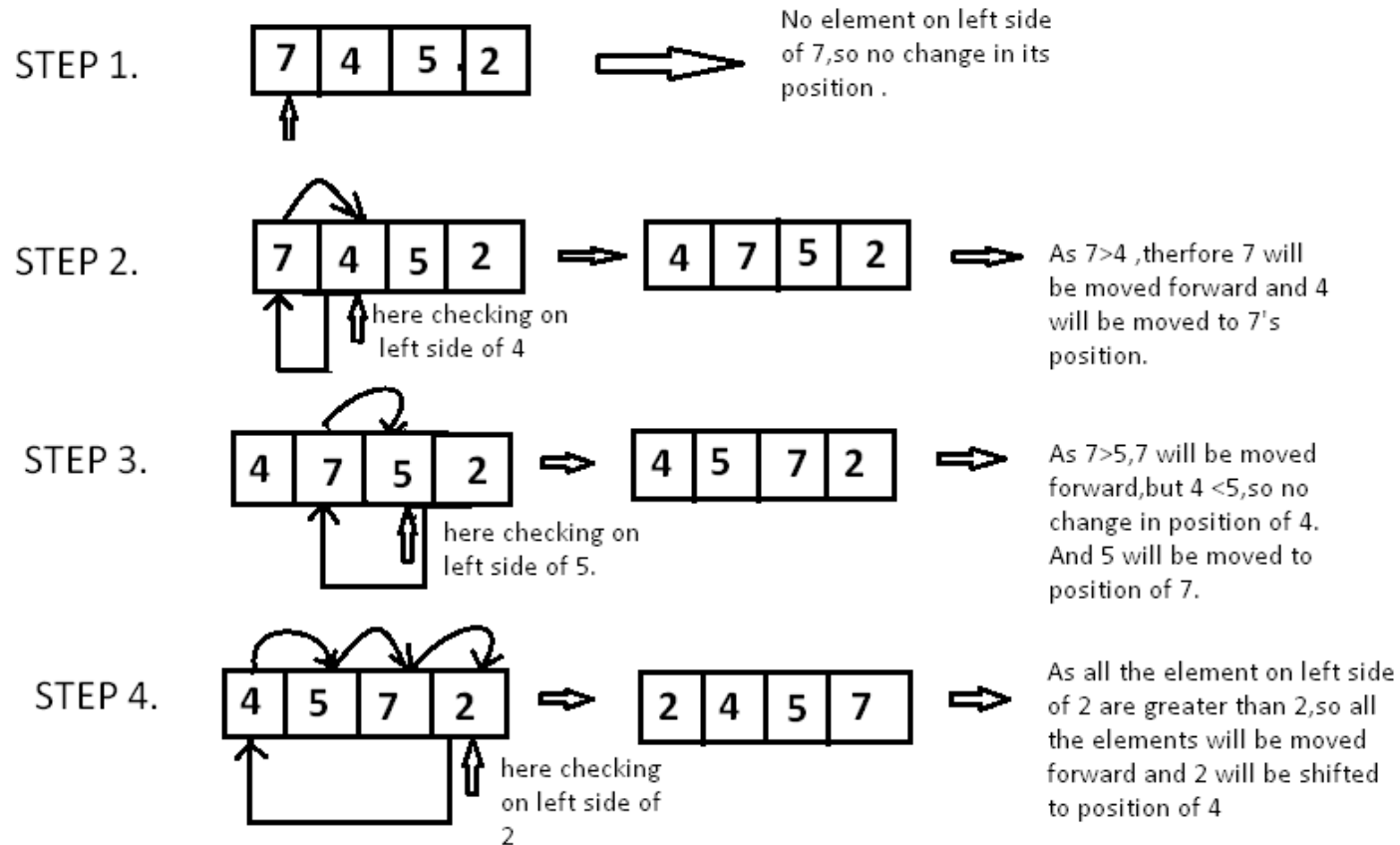
Outlines

1. Several Sorting Algorithms

- Insertion Sort
- Selection Sort
- Merge Sort
- Quicksort
- Counting Sort
- Radix Sort

2. Exercises

Insertion Sort $O(N^2)$



Selection Sort $O(N^2)$

The algorithm goes through each array **position** and **selects** a suitable value for that position.

Initial:	56	25	37	58	95	19	73	30
Round 0:	19	25	37	58	95	56	73	30
Round 1:	19	25	37	58	95	56	73	30
Round 2:	19	25	30	58	95	56	73	37
Round 3:	19	25	30	37	95	56	73	58
Round 4:	19	25	30	37	56	95	73	58
Round 5:	19	25	30	37	56	58	73	95
Round 6:	19	25	30	37	56	58	73	95

selects a value for:

`array[0]` 56 : swap with 19

`array[1]` 25 : no swap

`array[2]` 37 : swap with 30

`array[3]` 58 : swap with 37

`array[4]` 95 : swap with 56

`array[5]` 95 : swap with 58

`array[6]` 73 : no swap

Goes through each array *position* and *selects* a suitable value for that position.

Merge Sort $O(N \log N)$

Divide-and-Conquer

- Divide the array into two (or more) subarrays
- Sort each subarray (Conquer)
- Merge them into one (in a smart way!)

Pseudo Code:

MERGE_SORT(A, p, r)

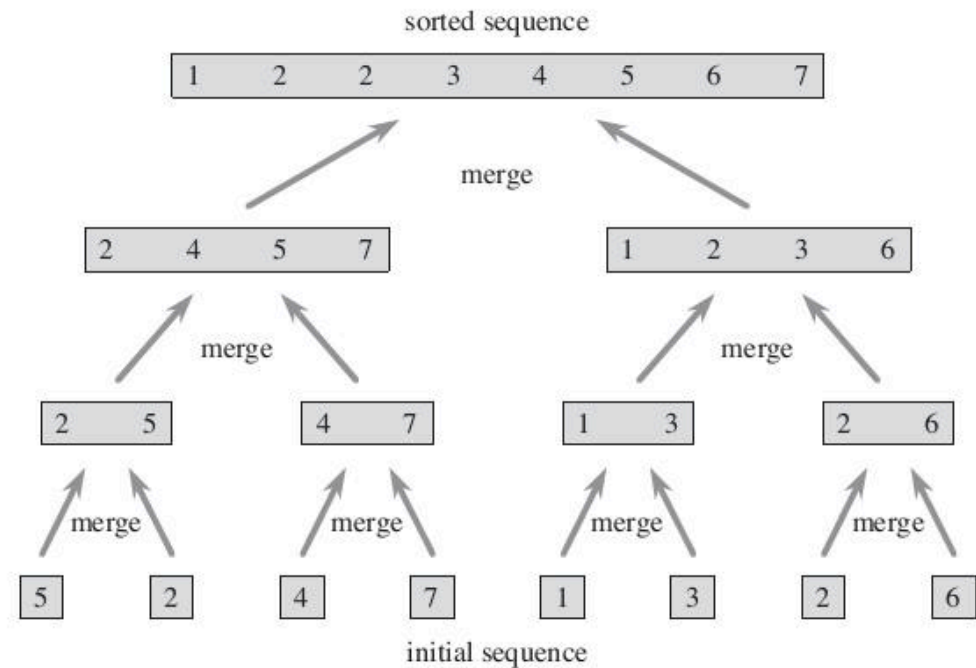
if $p < r$

$q = (p + r) / 2$

 MERGE_SORT(A, p, q)

 MERGE_SORT(A, q + 1, r)

 MERGE(A, p, q, r)



Quicksort $O(N \log N)$

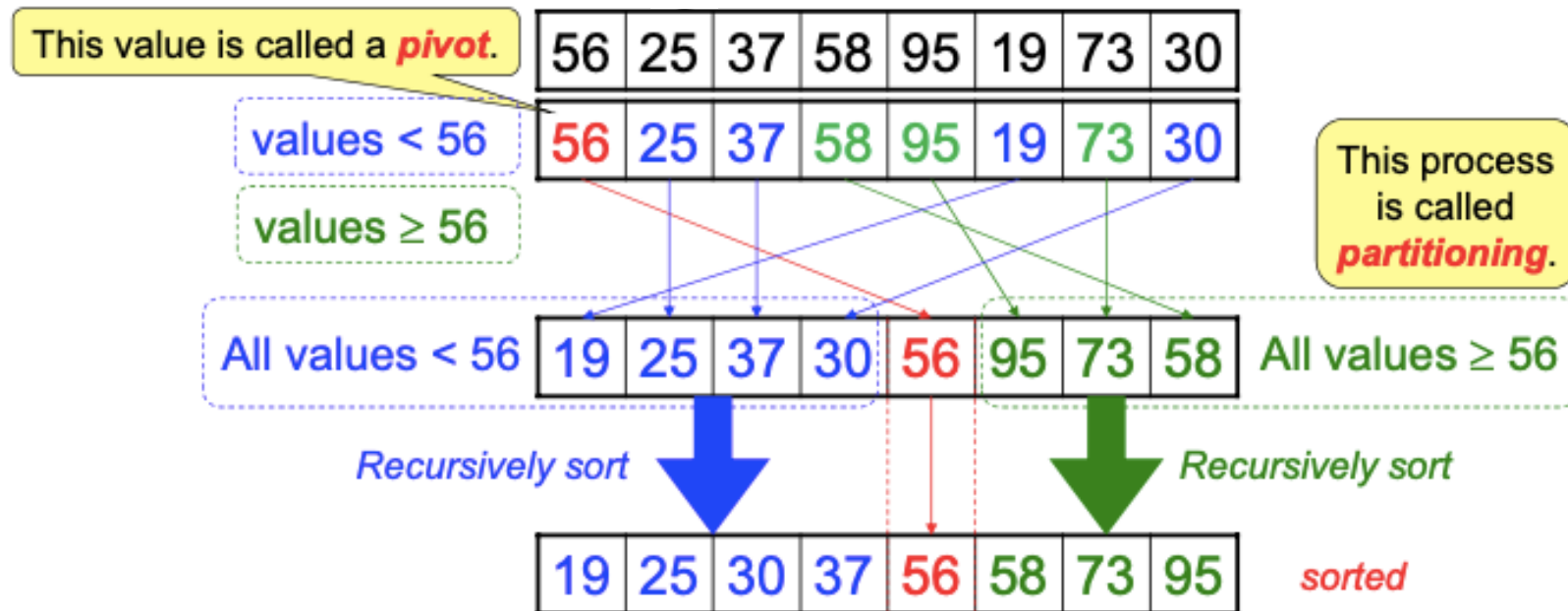
Like Merge Sort, QuickSort is a Divide and Conquer algorithm. It picks an element as pivot and partitions the given array around the picked pivot.

Any array element can be chosen as the pivot.

```
/* low --> Starting index, high --> Ending index */
quickSort(arr[], low, high)
{
    if (low < high)
    {
        /* pi is partitioning index, arr[pi] is now
           at right place */
        pi = partition(arr, low, high);

        quickSort(arr, low, pi - 1); // Before pi
        quickSort(arr, pi + 1, high); // After pi
    }
}
```

Quicksort $O(N \log N)$



Counting Sort $O(N + K)$

Counting sort is an algorithm for sorting a collection of objects according to keys that are small integers.

Pseudocode:

```
count = array of k+1 zeros
for x in input do
    count[key(x)] += 1

total = 0
for i in 0, 1, ... k do
    count[i], total = total, count[i] + total

output = array of the same length as input
for x in input do
    output[count[key(x)]] = x
    count[key(x)] += 1

return output
```


Radix Sort $O(D(N + K))$

Given N D-digits number in which each digit can take on up to K possible values, radix sort correctly sorts these numbers in $O(D(N + K))$.

- There are D counting sorts which each takes $O(N + K)$.

Original	Digit 0	Digit 1	Digit 2
329	72 <u>0</u>	7 <u>2</u> 0	<u>3</u> 29
457	35 <u>5</u>	3 <u>2</u> 9	<u>3</u> 55
657	43 <u>6</u>	4 <u>3</u> 6	<u>4</u> 36
839	45 <u>7</u>	8 <u>3</u> 9	<u>4</u> 57
436	65 <u>7</u>	3 <u>5</u> 5	<u>6</u> 57
720	32 <u>9</u>	4 <u>5</u> 7	<u>7</u> 20
355	83 <u>9</u>	6 <u>5</u> 7	<u>8</u> 39

Exercise 1

Which of the following sorting algorithms in its typical implementation gives best performance when applied on an array which is sorted or almost sorted (maximum 1 or two elements are misplaced).


- A. Quicksort
- B. Merge Sort
-  C. Insertion Sort

Insertion sort takes linear time when input array is sorted or almost sorted (maximum 1 or 2 elements are misplaced). All other sorting algorithms mentioned above will take more than linear time in their typical implementation.

Exercise 2

Suppose we are sorting an array of eight integers using quicksort, and we have just finished the first partitioning with the array looking like this:


2	5	1	7	9	12	11	10
---	---	---	---	---	----	----	----

-  A. The pivot could be either the 7 or the 9.
- B. The pivot could be the 7, but it is not the 9.
- C. The pivot is not the 7, but it could be the 9.
- D. Neither the 7 nor the 9 is the pivot.

7 and 9 both are at their correct positions (as in a sorted array). Also, all elements on left of 7 and 9 are smaller than 7 and 9 respectively and on right are greater than 7 and 9 respectively.


Exercise 3

Given an array where numbers are in range from 1 to n^6 , which sorting algorithm can be used to sort these number in linear time?

- A. Not possible to sort in linear time.
-  B. Radix Sort.
- C. Counting Sort.
- D. Quick Sort.

Exercise 4

Randomized quicksort is an extension of quicksort where the pivot is chosen randomly. What is the worst case complexity of sorting n numbers using randomized quicksort?

- A. $O(n)$
- B. $O(n \log n)$
-  C. $O(n^2)$
- D. $O(n!)$

Randomized quicksort has expected time complexity as $O(n \log n)$, but worst case time complexity remains same. In worst case the randomized function can pick the index of corner element every time.


Exercise 5

Assume that a mergesort algorithm in the worst case takes 30 seconds for an input of size 64. Which of the following most closely approximates the maximum input size of a problem that can be solved in 6 minutes?

- A. 256
- B. 512
- C. 1024
- D. 2048

Exercise 5

Assume that a mergesort algorithm has a running time of $c \cdot n \log n$ for an input of size n . Which of the following is the input size of a problem that can be solved in 6 minutes?

- A. 256
-  B. 512
- C. 1024
- D. 2048

Time complexity of merge sort is $\theta(n \log n)$

$c \cdot 64 \log 64$ is 30

$c \cdot 64 \cdot 6$ is 30

c is $5/64$

For time 6 minutes


$$5/64 \cdot n \log n = 6 \cdot 60$$

$$n \log n = 72 \cdot 64 = 512 \cdot 9$$

$$n = 512.$$

Exercise 6


Which is the correct order of the following algorithms with respect to their time Complexity in the best case ?

- A. Merge sort > Quick sort > Insertion sort > selection sort
-  B. Insertion sort < Quick sort < Merge sort < selection sort
- C. Merge sort > selection sort > quick sort > insertion sort
- D. Merge sort > Quick sort > selection sort > insertion sort

In best case Quick sort: $O(n \log n)$
Merge sort: $O(n \log n)$
Insertion sort: $O(n)$
Selection sort: $O(n^2)$

Exercise 7

Which one of the following in place sorting algorithms needs the minimum number of swaps?

- A. Quick sort
- B. Insertion sort
-  C. Selection sort

Selection sort takes minimum number of swaps to sort an array. It takes maximum of $O(n)$ comparisons to sort an array with n elements.