

Q1(a) We have two direct pointers and one indirect pointer.

Each direct pointer points to a data block, so it has 4KB, totally, we have 8KB.

Each indirect pointer points to a data block with all of its data storing direct pointer, i.e. $\frac{4096 \text{ Byte}}{4 \text{ Byte for each direct pointer}} = 1024 = 1K$

$$1K \times 4KB = 4MB$$

Finally, we conclude the biggest size of a file with SFS is $4MB + 8KB$.

Q1(bi) `read_t(inum, 7000, buf1, 10000)` is

First, we calculate a and b :

$$a = \lceil 7000 / 4096 \rceil = 1$$

$$b = 7000 \bmod 4096 = 2904$$

Because $a < 2$, `read_t()` starts at `direct_blk[1] = 9`

Then we calculate a' and b' :

$$a' = \lceil (2904 + 10000 - 1) / 4096 \rceil = 3$$

$$b' = (2904 + 10000 - 1) \bmod 4096 = 615$$

Because $a' \geq 2$, `read_t()` ends at indirect block. The block number

$$\text{is } \text{Cell}[3-4] = \text{Cell}[1] = 19$$

Thus, `read_t()` will read data blocks 9, 17, 19

Q1(bii) read-t(inum, 12000, buf1, 10000);

$$a = \lfloor 12000 / 4096 \rfloor = 2$$

$$b = 12000 \bmod 4096 = 3808$$

read-t() starts at indirect block. The block number is $\text{cell}[(2-2)]$
 $= \text{cell}[0] = 17$

$$a' = \lfloor (3808 + 26000 - 1) / 4096 \rfloor = 7$$

$$b' = (3808 + 26000 - 1) \bmod 4096 = 1135$$

read-t ends at indirect block. The block number is $\text{cell}[(7-2)]$
 $= \text{cell}[5] = 27$

Thus, read-t() will read data block 17, 19, 20, 21, 25, 27

Q1(biii) read-t(inum, 10000, buf1, 36000);

$$a = \lfloor 10000 / 4096 \rfloor = 2$$

$$b = 10000 \bmod 4096 = 1808$$

read-t() starts at indirect block. The block number is $\text{cell}[(2-2)]$
 $= \text{cell}[0] = 17$

$$a' = \lfloor (1808 + 36000 - 1) / 4096 \rfloor = 9$$

$$b' = (1808 + 36000 - 1) \bmod 4096 = 943$$

read-t ends at indirect block. The block number is $\text{cell}[(9-2)]$
 $= \text{cell}[7] = 29$

Thus read_t() will read data block 17, 19, 20, 21, 25, 27, 28, 29

Q1(biv) read_t(inum, 1000, buf1, 31000);

$$a = \lfloor 1000 / 4096 \rfloor = 0$$

$$b = 1000 \bmod 4096 = 1000$$

read_t() starts at direct_blk[0] = 3

$$a' = \lfloor (1000 + 31000 - 1) / 4096 \rfloor = 7$$

$$b' = (1000 + 31000 - 1) \bmod 4096 = 3327$$

read_t ends at indirect block. The block number is cell[(7-2)]
= cell[5] = 27

Thus read_t() will read data block 3, 9, 17, 19, 20, 21, 25, 27

| | read_t (inum, offset, buf1, count) | The data block numbers in sequence that will be read from (only list the data blocks that contain file data) |
|-----------|------------------------------------|--|
| Example 1 | read_t(inum, 133, buf1, 400); | 3 |
| Example 2 | read_t(inum, 133, buf1, 6000); | 3, 9 |
| (i) | read_t(inum, 7000, buf1, 10000); | 9, 17, 19 |
| (ii) | read_t(inum, 12000, buf1, 26000); | 17, 19, 20, 21, 25, 27 |
| (iii) | read_t(inum, 10000, buf1, 36000); | 17, 19, 20, 21, 25, 27, 28, 29 |
| (iv) | read_t(inum, 1000, buf1, 31000); | 3, 9, 17, 19, 20, 21, 25, 27 |

Q2

Data block 0: "/"

| | |
|-------|----|
| . | 0 |
| .. | 0 |
| dir7 | 7 |
| dir6 | 6 |
| dir11 | 11 |
| file5 | 18 |

Data block 7: "dir7"

| | |
|------|---|
| . | 7 |
| .. | 0 |
| dir3 | 3 |
| dir1 | 1 |

Data block 8: "dir8"

| | |
|-------|----|
| . | 8 |
| .. | 11 |
| dir10 | 10 |
| dir5 | 5 |

Data block 4: "dir4"

| | |
|-----|---|
| . | 4 |
| .. | 3 |
| dir | 9 |

Data block 10: "dir10"

| | |
|-------|----|
| . | 0 |
| .. | 8 |
| file2 | 15 |
| file4 | 17 |

Q2(b)

inode 0 \rightarrow data block 0 \rightarrow inode 11 \rightarrow data block 11

\rightarrow inode 8 \rightarrow data block 8 \rightarrow inode 10 \rightarrow data block 10

\rightarrow inode 15