Lesson 5

Declaring types and classes

3 type definitions:

1.**type** → Transparent types: a new name for already existing types

2. data

Opaque types, user defined types

3. **newtype** → Especially simple opaque types: more efficiently implemented

1. Type

```
type String = [Char]
```

type Pos = (Int , Int)

type Trans = Pos -> Pos

type Tree = (Int, [Tree]) --not allowed, recursive!!

polymorphic : type Pair a = (a,a)

type Assoc k v = [(k,v)]

find :: Eq k => k -> Assoc k v -> v find k t = head [v | (k',v) <- t, k==k']

2. Data

from Prelude, data Bool = False | True

False and True are constructors

data A = B | C would also work

data Move = North | South | East | West deriving Show

move :: Move -> Pos -> Pos

```
move North (x,y) = (x,y+1)
move South (x,y) = (x, y-1)
move East (x,y) = (x+1,y)
move West (x,y) = (x-1,y)
```

```
moves :: [Move] -> Pos -> Pos
moves [] p = p
moves (m:ms) p = moves ms (move m p)
```

esercizio invmoves ??

a more advanced example

data Shape = Circle Float | Rect Float Float

Circle and Rect are costructor functions:

> :t Rect

Rect:: Float -> Float -> Shape

>Rect 2.3 1.4

Rect 2.3 1.4

insomma il risultato di applicare Rect 2.3 1.4 è l'espressione stessa. The values of Shape are Rect 2.3 1.4 and Circle 3.3

square :: Float -> Shape square x = Rect x x

area :: Shape -> Float area (Circle r) = pi * r^2 area (Rect x y) = x*y

data Maybe a = Nothing | Just a

is used to represent a success/failure behaviour of functions

Nothing stands for failure

Just x stands for success with result x

divisione safe

```
safediv : Int -> Int -> Maybe Int
safediv _ 0 = Nothing
safediv m n = Just (m 'd' n)
```

safe head :: [a] -> Maybe a
safehead [] = Nothing
safehead xs = Just (head xs)

3. newtype

when a data definition has one costructor only with one parameter only then it is possible to define it like this,

newtype Nat = N Int

how does it compare with

type Nat = Int data Nat = N Int ?? Recursive types types defined using data and newtype can be rcursive

Nat = Zero | Succ Nat

Zero, Succ Zero, Succ (Succ Zero),

nat2int :: Nat -> Int
nat2int Zero =0
nat2int (Succ n) = 1 + nat2int n

int2nat :: Int -> Nat int2nat 0 = Zero int2nat n = Succ (int2nat (n-1)) add :: Nat -> Nat -> Nat add m n = int2nat (nat2int n + nat2int m)

ma anche add Zero n = n add (Succ m) n = Succ (add m n)

binary trees

data Tree a = Leaf a | Node (Tree a) a (Tree a)

```
flatten :: Tree a -> [a]
flatten (Leaf x) = [x]
flatten (Node I x r) = flatten I ++ [x] ++ flatten r
```

data Tree a = Node a [Tree a] data Tree a b = Leaf a | Node (Tree a b) b (Tree a b)

Class and instance declarations

a new class can be defined with the class mechanism

for a type a to be in Eq it must support == and /=, and is enough to define ==

instance Eq Bool where
False == False = True
True == True = True
== False

only types defined by data and new types can me made into instances of classes

default definitions of methods (such as /=) can be overridden.

Classes can be extended to form new classes

min x y | x
$$\leq$$
 y = x | otherwise = y

$$\max x y \mid x \le y = y$$

| otherwise = x

if we have a type in Eq, to make it also in Ord one needs to define 4 methods

instance Ord Bool where

Derived instances

when new types are defined it is often convenient to make them into instances of built-in classes

data Bool = False | True deriving (Eq, Ord, Show, Read) in this way all metods that belong to the classes can be used for Bool:

➤ False == False

True

> False < True

True

the order False < True is determined by the declaration data Bool = False | True

in case of costructors with arguments, the types of those arguments must be in the classes

data Shape = Circle Float | Rect Float Float

data Maybe a = Nothing | Just a

if me make Shape a deriving Class then Float in Class

lexigographical order

- > Rect 1.0 4.0 < Rect 2.0 3.0
- True
- > Rect 1.0 4.0 < Rect 1.0 3.0
- False
- > Circle 1.0 < Circle 1.1
- True