

# stateful computation

continuation

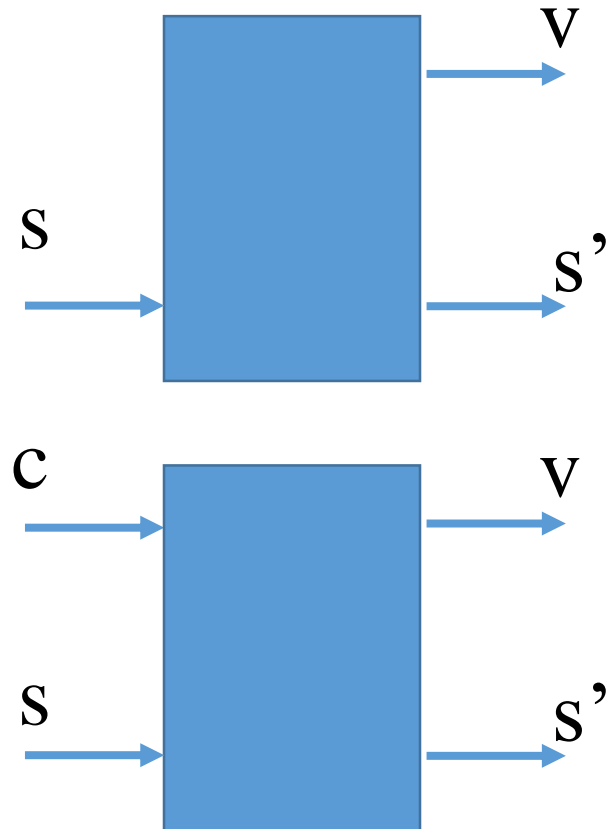
# State Monad

`type State = Int`

`type ST = Int -> Int`

`type ST a = Int -> (a, Int)`

`Char -> ST a`



```
newtype ST a = S (State ->(a, State))
```

```
app :: ST a -> State ->(a, State)
```

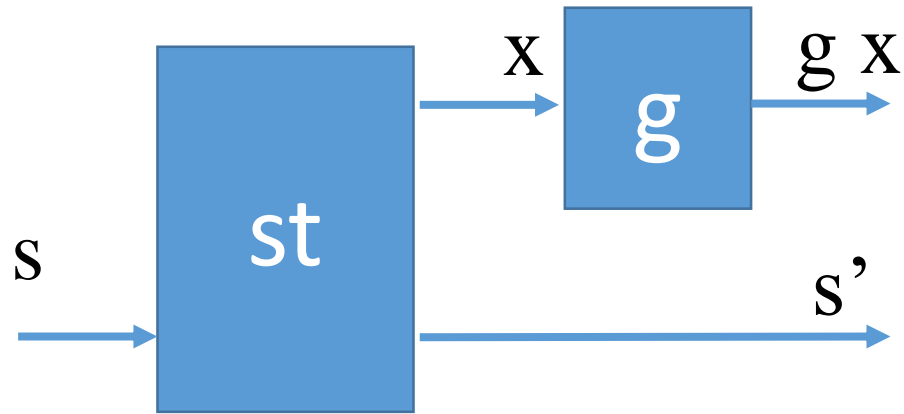
```
app (S st) x = st x
```

ST a is a Functor, an Applicative and also a Monad

instance Functor ST where

--fmap :: (a -> b) -> ST a -> ST b

fmap g st = S(\ s -> let (x,s') = app st s in (g x, s'))



recall that in general `fmap` applies a function `g` to the values contained in a structure

in this case `g` is applied to the output of `st`

instance Applicative ST where

--pure :: a -> ST a

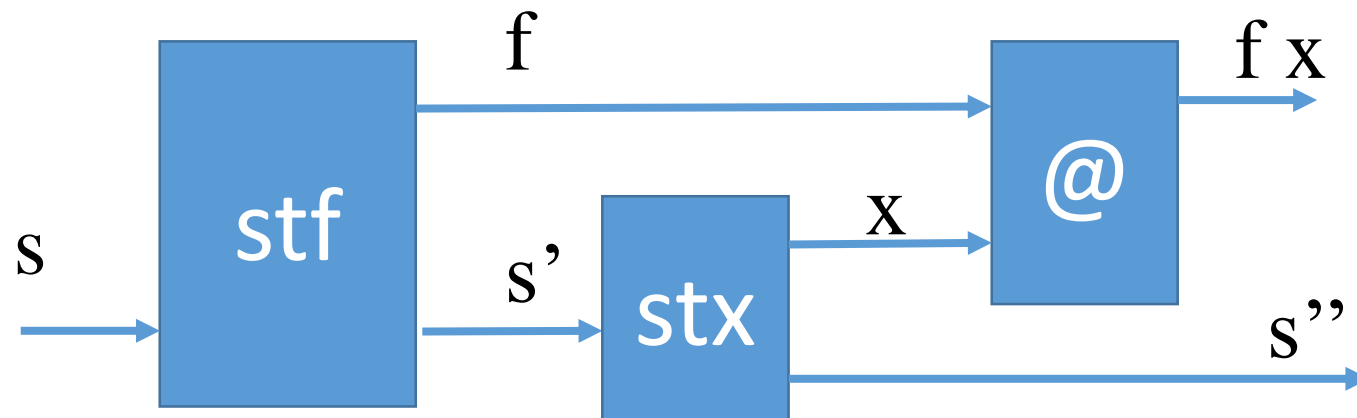
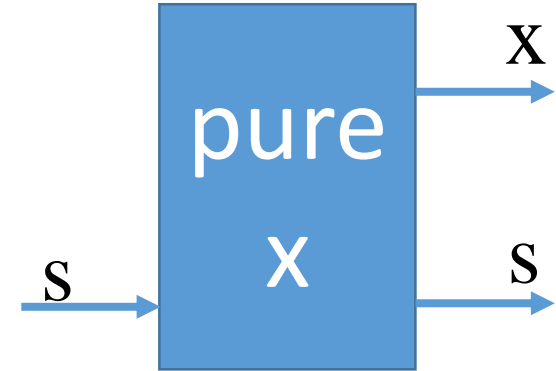
pure x = S (\s -> (x,s))

--(<\*>) :: ST (a->b) -> ST a -> ST b

stf <\*> stx = S (\s ->

let (f, s') = app stf s

(x, s'') = app stx s' in (f x, s''))

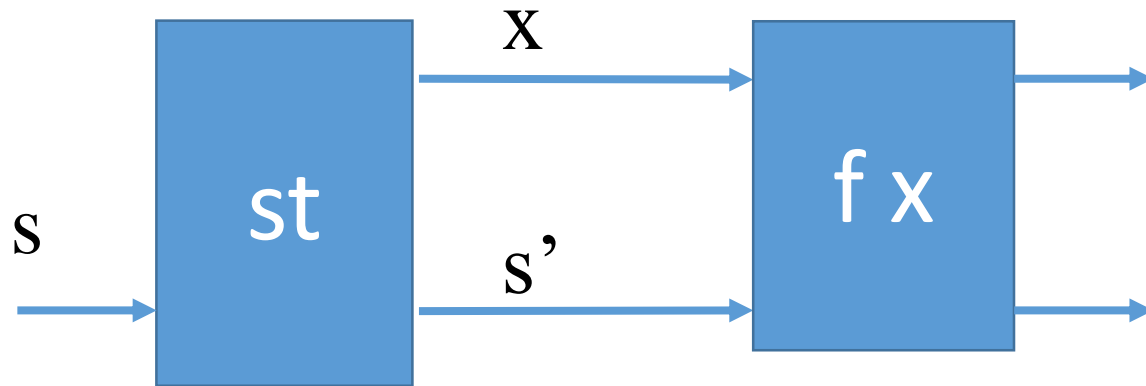


stf and stx are both  
ST and thus both  
produce a change of  
state

instance Monad ST where

--(>>=) :: ST a -> (a -> ST b) -> ST b

st >>= f = S (\s -> let (x,s')=app st s in app (f x) s')



this ST depends on x

use ST a to define a function that takes a binary tree and relabels the leaves with fresh increasing integer according with the infix traversal 0, 1, 2, 3 .... or n, n+1, n+2,....

```
data Tree a = Leaf a | Node (Tree a) (Tree a)
    deriving Show
```

1st solution

```
rlabel :: Tree a -> Int -> (Tree Int, Int)
```

```
rlabel (Leaf _) n = (Leaf n, n + 1)
```

```
rlabel (Node l r) n = (Node l' r', n'')
```

where (l',n') = rlabel l n

(r',n'') = rlabel r n'

this solution is complicated : we need to pass around a growing integer  $n$ . We can do better exploiting the fact that

$$\text{rlabel} :: \text{Tree } a \rightarrow \text{Int} \rightarrow (\text{Tree } \text{Int}, \text{Int}) \rightarrow \text{Tree } a \rightarrow \text{ST } (\text{Tree } \text{Int})$$

the integer that we passed around can be the State

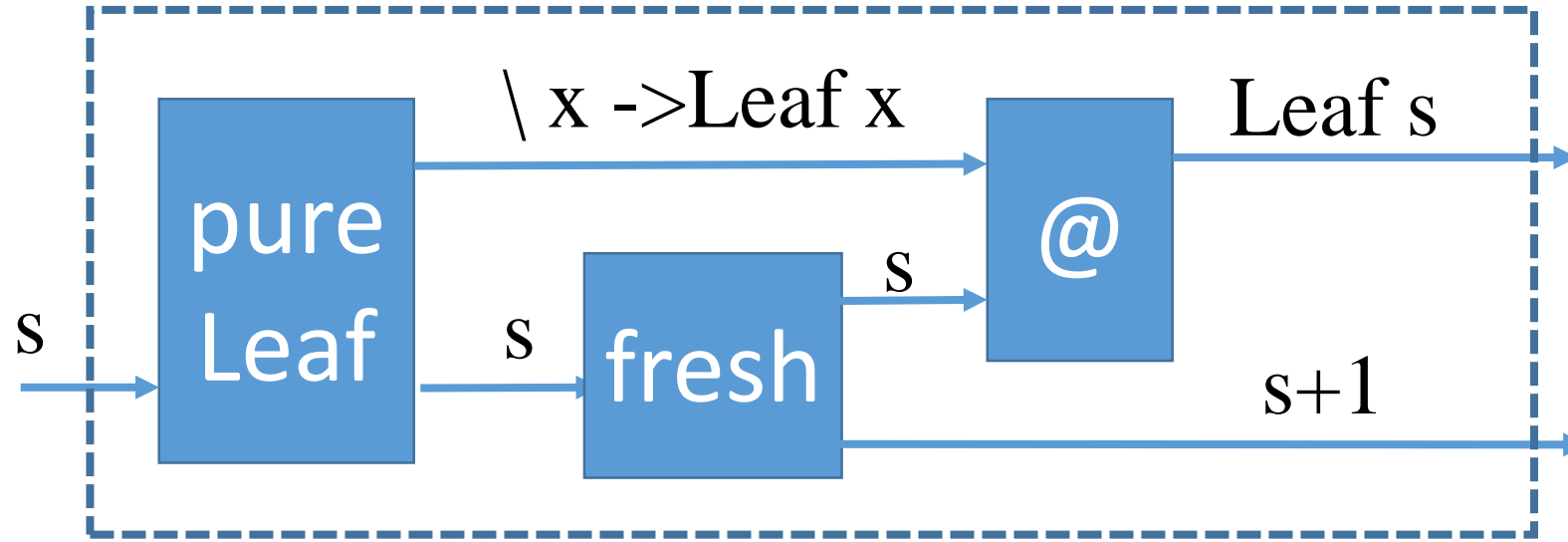
$$\text{fresh} :: \text{ST } \text{Int}$$
$$\text{fresh} = \text{S } (\backslash n \rightarrow (n, n+1))$$

we use the fact that  $\text{ST } a$  is Applicative

$$\text{alabel} :: \text{Tree } a \rightarrow \text{ST } (\text{Tree } \text{Int})$$
$$\text{alabel } (\text{Leaf } \_) = \text{pure Leaf } \langle * \rangle \text{fresh}$$
$$\text{alabel } (\text{Node } l \ r) = \text{pure Node } \langle * \rangle \text{alabel } l \ \langle * \rangle \text{alabel } r$$

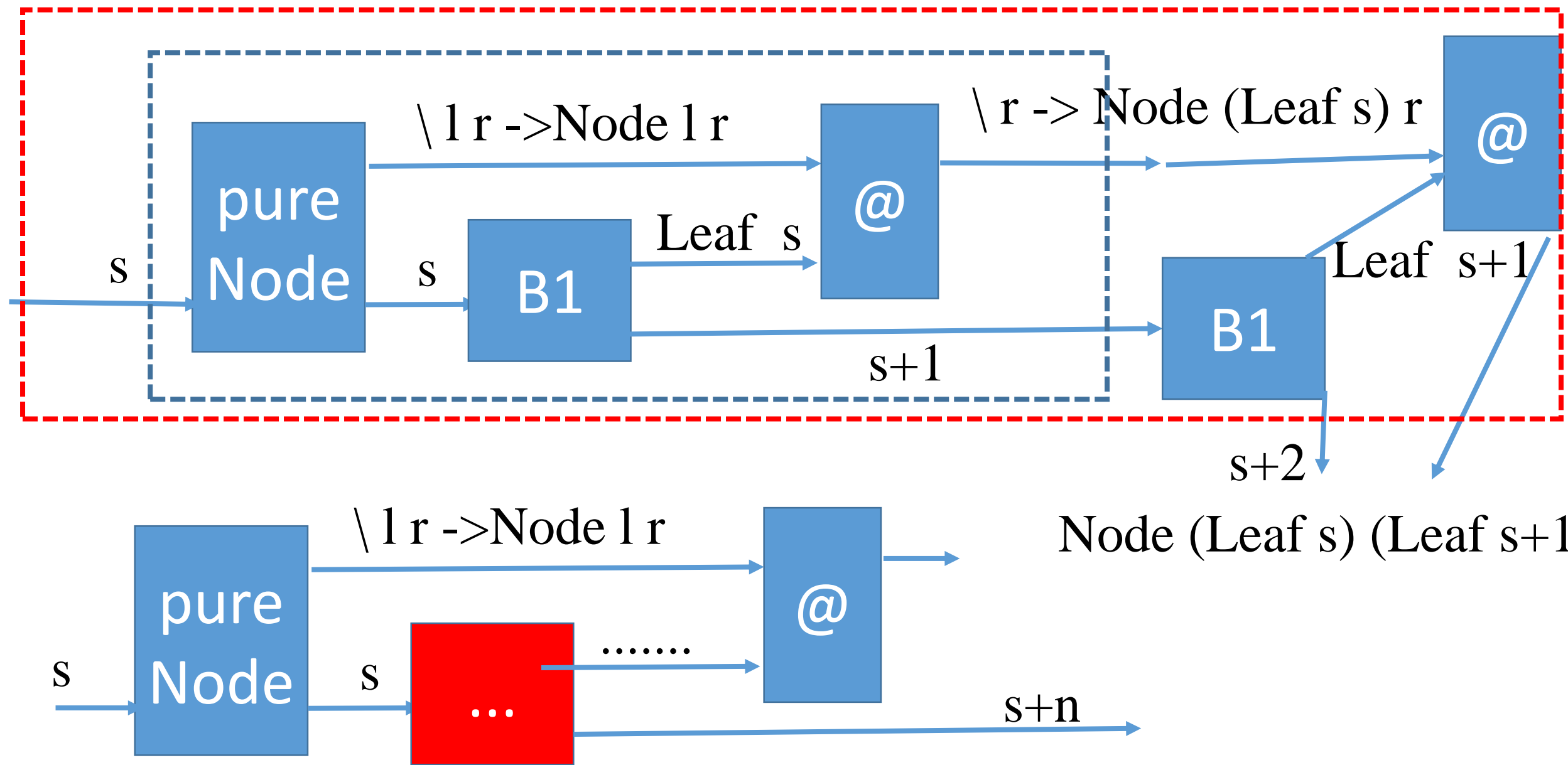


$\text{alabel (Leaf \_)} = \text{pure Leaf} \langle * \rangle \text{fresh} :: \text{ST (Tree Int)}$   
 $= \text{Int} \rightarrow (\text{Tree Int}, \text{Int})$



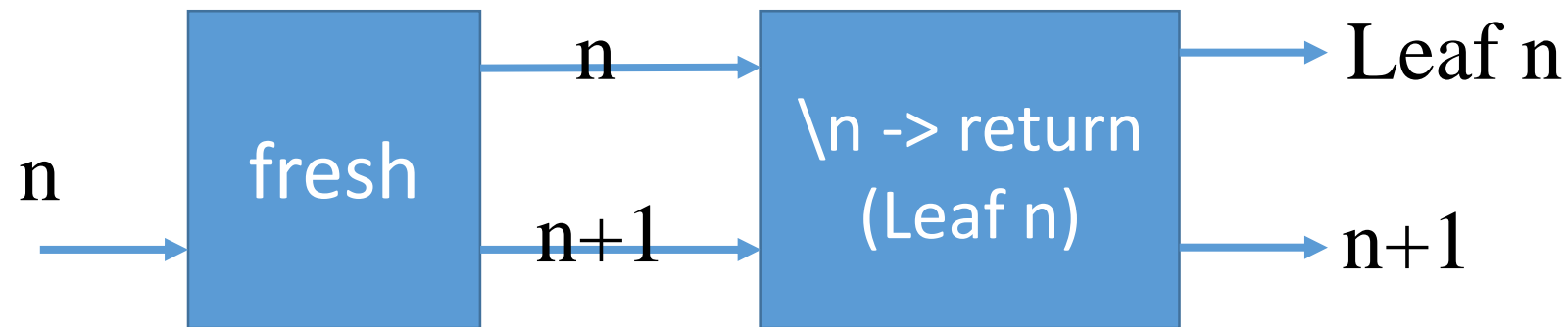
**B1**

$\text{alabel (Node l r)} = \text{pure Node } \langle * \rangle \text{ alabel l } \langle * \rangle \text{ alabel r}$



using the fact that ST is a Monad:

```
mlabel (Leaf _) = fresh >>= \n -> return Leaf n  
mlabel (Node l r) = mlabel l >>=  
    (\l' -> mlabel r >>=  
    (\r' -> return Node l' r'))
```

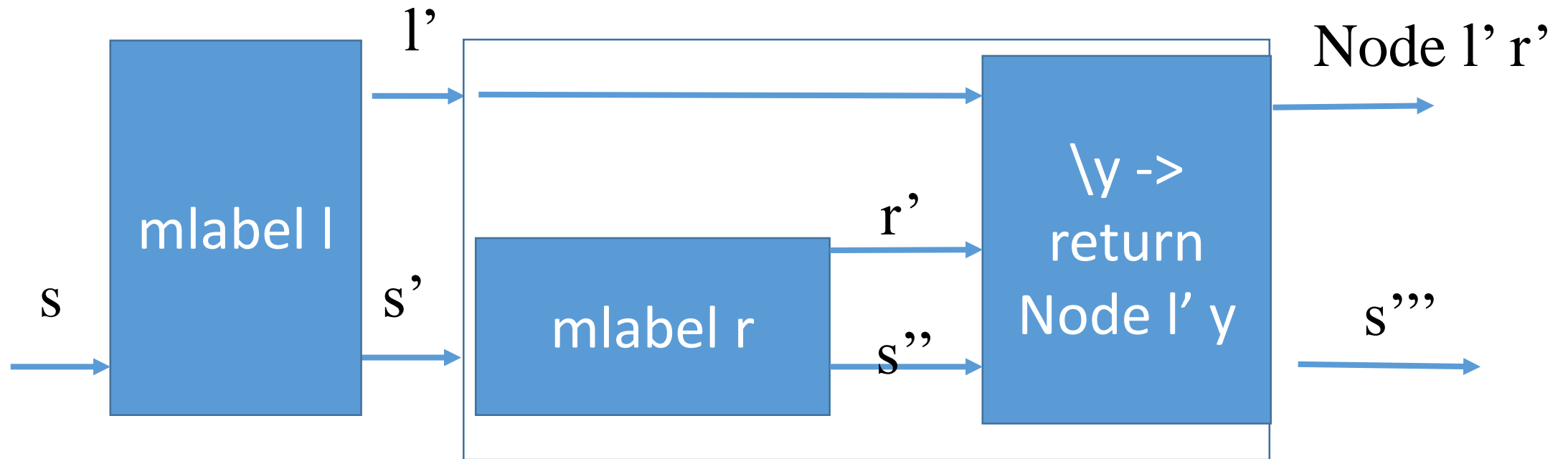


type  
ST Tree Int

```

mlabel (Node l r) = mlabel l >>=
  (\x -> mlabel r >>=
    (\y -> return Node x y))

```



using the do notation of Monads

```
mlabel :: Tree a -> ST (Tree Int)
```

```
mlabel (Leaf _) = do n <- fresh  
                  return (Leaf n)
```

```
mlabel (Node l r) = do l' <- mlabel l  
                       r' <- mlabel r  
                       return (Node l' r')
```

## exercise

Tree a = Leaf a | Node (Tree a) a (Tree a)

mlabel (Leaf \_) = fresh >>= \n -> return (Leaf n)

mlabel (Node l \_ r) = (mlabel l) >>=  
(\l' -> fresh >>=  
(\n -> (mlabel r) >>=  
(\r' -> return (Node l' n r')))))

with do notation:

```
mlabel (Leaf _)    = do n <- fresh  
                    return (Leaf n)
```

```
mlabel (Node l a r) = do l' <- mlabel l  
                          n <- fresh  
                          r' <- mlabel r  
                          return (Node l' n r')
```

## generic functions

`mapM :: Monad m => (a -> m b) -> [a] -> m [b]`

`mapM f [] = return []`

`mapM f (x:xs) = do y <- f x  
                  ys <- mapM f xs  
                  return (y : ys)`

`conv :: Char -> Maybe Int`

`conv c | isDigit c = Just (digitToInt c)  
      | otherwise = Nothing`

`mapM conv «1234»`

`Just [1,2,3,4]`

`mapM conv «123b» → Nothing`



`filterM :: Monad m => (a -> m Bool) -> [a] -> m [a]`

`filterM p [] = return []`

`filterM p (x:xs) = do b <- px`

`ys <- filterM p xs`

`return (if b then x:ys else ys)`

`filterM (\x -> [True,False]) [1,2,3]`

`[[1,2,3],[1,2],[1,3],[1],[2,3],[2],[3],[]]`

```
join :: Monad m => m (m a) -> m a
```

```
join mmx = do mx <- mmx
```

```
    x <- mx
```

```
    return x
```

```
join [[1],[2,3],[],[4]]
```

```
[1,2,3,4]
```

```
join (Just (Just 1))
```

```
Just 1
```

```
join (Just Nothing)
```

```
Nothing
```

## Monad laws

1.  $\text{return } x \gg= f = f \ x$
2.  $m x \gg= \text{return} = m \ x$
3.  $(m x \gg= f) \gg= g = m x \gg= (\backslash x \rightarrow (f \ x \gg= g))$

check the types !!