

Differential 2-rigs

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Motivations

Galois theory asserts that to each polynomial equation one can attach a finite group, the **Galois group** of the polynomial, in such a way that the roots of a polynomial can be found via algebraic operations and root extraction if and only if its Galois group is **solvable**.

A group G is solvable if there exists a chain of subgroups $1 \leq G_1 \leq \cdots \leq G_n \leq G$ with the property that each G_i is normal in G_{i+1} , and each quotient G_{i+1}/G_i is abelian.

Soon after Galois' tragic death people started wondering if a similar theorem could be attained for **differential** equations. The work of Lie and Liouville goes precisely in the direction of establishing such conditions.

Turns out that this is possible, paying a price:

- the group attached to a differential equation is not finite any more;
- it has a non-trivial topology, and the “correct” subgroups to consider are the closed ones;
- sometimes, such groups are algebraic manifolds of **infinite dimension**.

What we now call **Lie groups** was nothing but the attempt by S. Lie to classify the objects arising as Galois groups of differential equations.

The theory of **reductive** algebraic groups arose as a way to understand the operation of adding a solution to $y' = y$ to the ring of polynomials. (=exponential elements; they live in rings of power series).

A **differential algebra** over k is a k -algebra R endowed with an endomorphism $d : R \rightarrow R$ that is k -linear, and satisfies the **Leibniz rule**:

$$d(a.b) = da.b + a.db$$

A **differential equation** in R is an equation of the form $F(y, y^{(1)}, y^{(2)}, \dots) = 0$ where F is a polynomial with coefficients in R , and $y^{(1)} := dy, y^{(n)} = d(y^{(n-1)})$.

A **differential extension** of R is a bigger differential k -algebra $F \supset R$ obtained from R by adding solutions to differential equations;

(Liouville) Solvability of the **Galois group** of a differential extension $F \supset R$ allows the possibility to solve differential equations in R , finding their solutions in F , by means of ‘elementary operations’:

- Ring operations
- addition of integrals (=solutions to $y' = a$, $a \in R$)
- addition of exponentials (=solutions to $y' = by$, $b \in R$).

Slogan

Abstracting the theory of differential rings is a secure path to unexpectedly profound mathematics.

Guided by this, the plan is to transform rings into (rig) categories and endomorphisms $d : R \rightarrow R$ into functors $\partial : \mathcal{R} \rightarrow \mathcal{R}$, and the Leibinz condition

$$\partial(A \otimes B) \cong \partial A \otimes B + A \otimes \partial B$$

into an isomorphism natural in A, B .

Rings

Definition

A **2-rig** is a category \mathcal{C} such that

- \mathcal{C} admits finite coproducts, denoted $A \cup B$;
- \mathcal{C} admits a monoidal structure $\otimes : \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$ that is **bilinear**, i.e. the functors $A \otimes -$ and $- \otimes B$ commute with coproducts.

In simple terms, we want to capture a notion that categorifies **rigs** (rings without additive inverses).

A 2-rig is a particular instance of a more general notion, first introduced by Laplaza: a category \mathcal{C} with two monoidal structures \otimes, \oplus , such that ' \otimes distributes over \oplus '.

As natural as his axiomatics may seem, the precise formalisation of a coherence theorem for a distributive category requires a lot of effort and numerous diagrams: see Laplaza for the precise definition.

Many of our examples will also satisfy additional assumptions:

- A **commutative** 2-rig is a category \mathcal{C} , where \otimes is also symmetric;
- A **monoidally cocomplete** 2-rig is a category \mathcal{C} , which moreover admits **all** small colimits, and such that $A \otimes -$ and $- \otimes B$ distribute over all colimits;
- A **closed** 2-rig is a category \mathcal{C} such that each $A \otimes -$ and $- \otimes B$ have right adjoints; in this case, of course, they preserve all colimits that exist in \mathcal{C} .

Combinations are possible: it is clear what a **commutative closed 2-rig** is.

The following are examples of 2-rigs:

- The category $(\text{Set}, \times, 1)$ of **sets and functions** is a commutative closed 2-rig; more generally, all cartesian closed categories with coproducts $(\mathcal{A}, \times, 1)$ are bicartesian categories.
- The category $(\text{Mod}_R, \otimes, R)$ of **modules over a ring** R is a commutative closed 2-rig.
- The category of (real or complex) topological **vector bundles** over a topological space X , equipped with the tensor product of vector bundles is a 2-rig (where \cup is the direct sum of vector bundles taking the bundle associated with fiberwise Vect-coproduct).

- Given a monoidal category (\mathcal{A}, \oplus, j) the category $([\mathcal{A}^{\text{op}}, \text{Set}], *, yj = \mathcal{A}(j, -))$ of presheaves over \mathcal{A} endowed with the **Day convolution** monoidal structure

$$F * G := \int^{U, V \in \mathcal{A}} FU \times GV \times \mathcal{A}(U \oplus V, -)$$

is a closed 2-rig.

Note that $[\mathcal{A}^{\text{op}}, \text{Set}]$ is closed no matter what \oplus is, and the internal hom can be computed as

$$\{G, H\} : A \mapsto \int_X \text{Set}(GX, H(A \oplus X))$$

Derivations

Definition (Derivation on a 2-rig)

A **derivation** on a 2-rig is a functor $\partial : \mathcal{C} \rightarrow \mathcal{C}$ having the following properties:

- $\partial(A \cup B) \cong \partial A \cup \partial B$, and naturally so; this means that ∂ is a strong monoidal functor with respect to the \cup monoidal structure.
- $\partial(A \otimes B) \cong \partial A \otimes B \cup A \otimes \partial B$ and naturally so.

$\partial(A \otimes B) \cong \partial(A) \otimes B \cup A \otimes \partial B$ and naturally so

This second condition deserves to be spelled out completely: it means that ∂ is equipped with a 2-cell \mathfrak{l} filling the diagram

$$\begin{array}{ccccc}
 \mathcal{C} \times \mathcal{C} & \xrightarrow{\Delta} & \mathcal{C} \times \mathcal{C} \times \mathcal{C} \times \mathcal{C} & \xrightarrow{(\partial \otimes \mathcal{C}, \mathcal{C} \otimes \partial)} & \mathcal{C} \times \mathcal{C} \\
 \downarrow \otimes & & \uparrow \mathfrak{l} & & \downarrow \cup \\
 \mathcal{C} & \xrightarrow{\partial} & & & \mathcal{C}
 \end{array}$$

where $\Delta_{\mathcal{C} \times \mathcal{C}}$ is the diagonal functor $(A, B) \mapsto (A, B, A, B)$ and $(\partial \otimes \mathcal{C}, \mathcal{C} \otimes \partial)$ does the obvious thing.

The cell \mathfrak{l} is called the **leibnizator**, has components

$\mathfrak{l}_{AB} : \partial(A \otimes B) \Rightarrow \partial A \otimes B \cup A \otimes \partial B$, and it is subject to the following coherence conditions:

Compatibility with the right distributor:

$$\begin{array}{ccc}
 \partial((Y \cup Z) \otimes X) & \xrightarrow{\delta^R} & \partial(Y \otimes X \cup Z \otimes X) \\
 \downarrow \mathfrak{l} & & \downarrow \wr \\
 (Y \cup Z) \otimes \partial X \cup \partial(Y \cup Z) \otimes X & & \partial(Y \otimes X) \cup \partial(Z \otimes X) \\
 \downarrow \delta^R \cup \delta^R & & \downarrow \mathfrak{l} \cup \mathfrak{l} \\
 Y \otimes \partial X \cup Z \otimes \partial X \cup \partial Y \otimes X \cup \partial Z \otimes X & \xrightarrow{\sim} & \partial Y \otimes X \cup Y \otimes \partial X \cup \partial Z \otimes X \cup Z \otimes \partial X
 \end{array}$$

where the unnamed isomorphisms are symmetries of \cup or arising from the strong \cup -monoidality of ∂ ;

Compatibility with the left distributor:

$$\begin{array}{ccc}
 \partial(X \otimes (Y \cup Z)) & \xrightarrow{\delta^L} & \partial(X \otimes Z \cup Y \otimes Z) \\
 \downarrow \wr & & \downarrow \\
 \partial X \otimes (Y \cup Z) \cup X \otimes \partial(Y \cup Z) & & \partial(X \otimes Y) \cup \partial(X \otimes Z) \\
 \downarrow \delta^L \cup \delta^L & & \downarrow \wr \cup \wr \\
 \partial X \otimes Y \cup \partial X \otimes Z \cup X \otimes \partial Y \cup X \otimes \partial Z & \longrightarrow & \partial X \otimes Y \cup X \otimes \partial Y \cup \partial X \otimes Z \cup X \otimes \partial Z
 \end{array}$$

where the unnamed isomorphisms are symmetries of \cup or arising from the strong \cup -monoidality of ∂ ;

Compatibility with the right and left annihilator:

$$\begin{array}{ccc}
 \partial(X \otimes 0) & \xrightarrow{\sim} & \partial 0 \\
 \downarrow \iota_{X0} & & \downarrow \wr \\
 \partial X \otimes 0 \cup X \otimes \partial 0 & \twoheadrightarrow 0 + 0 \twoheadrightarrow & 0
 \end{array}$$

$$\begin{array}{ccc}
 \partial(0 \otimes X) & \xrightarrow{\sim} & \partial 0 \\
 \downarrow \iota_{0X} & & \downarrow \wr \\
 \partial 0 \otimes X \cup 0 \otimes \partial X & \twoheadrightarrow 0 + 0 \twoheadrightarrow & 0
 \end{array}$$

where the unnamed isomorphisms come from the fact that $A \otimes -$ and $- \otimes B$ preserve the initial object for all $A, B \in \mathcal{C}$;

Compatibility with the right and left \otimes -unitor:

$$\begin{array}{ccc}
 \partial(X \otimes I) & \xrightarrow{\partial\rho_X} & \partial X \\
 \downarrow \iota_{XI} & & \uparrow \iota_{\partial X} \\
 \partial X \otimes I \cup X \otimes \partial I & & \\
 \downarrow \rho_{\partial X} \cup X \otimes \partial I & & \\
 \partial X \cup X \otimes \partial I & &
 \end{array}
 \qquad
 \begin{array}{ccc}
 \partial(I \otimes X) & \xrightarrow{\partial\lambda_X} & \partial X \\
 \downarrow \iota_{IX} & & \uparrow \iota_{\partial X} \\
 \partial I \otimes X \cup I \otimes \partial X & & \\
 \downarrow \partial I \otimes X \cup \lambda_{\partial X} & & \\
 \partial I \otimes X \cup \partial X & &
 \end{array}$$

(the properties of ∂I , where I is the \otimes -monoidal unit, are quite a subtle business; we will come back to this later; note in particular that **no axiom entails that $\partial I = 0$**).

Compatibility with the associator:

$$\begin{array}{ccc}
 \partial((A \otimes B) \otimes C) & \xrightarrow{\partial\alpha} & \partial(A \otimes (B \otimes C)) \\
 \downarrow \wr_{A \otimes B, C} & & \downarrow \wr_{A, B \otimes C} \\
 \partial(A \otimes B) \otimes C \cup (A \otimes B) \otimes \partial C & & \partial A \otimes (B \otimes C) \cup A \otimes \partial(B \otimes C) \\
 \searrow \wr_{A, B \otimes C \cup (A \otimes B) \otimes \partial C} & & \swarrow \\
 \partial A \otimes B \otimes C \cup A \otimes \partial B \otimes C \cup A \otimes B \otimes \partial C & &
 \end{array}$$

First remarks

Directly from these definitions we can easily see that

- The derivative of the **initial object** \emptyset must be the initial object (because ∂ preserves the empty coproduct, or because of the Leibniz rule);
- by induction on n ,

$$\partial(A^{\otimes n}) \cong n \cdot A^{\otimes(n-1)} \otimes \partial A$$

where $n \cdot -$ sends an object $X \in \mathcal{C}$ to the n -fold coproduct $X \cup \dots \cup X$.

- again by induction,

$$\partial^n(X \otimes Y) = \coprod_{k=0}^n \binom{n}{k} \cdot \partial^{n-k} X \otimes \partial^k Y$$

where $\binom{n}{k}$ is the set of k -elements subsets of $\{1, \dots, n\}$.

Examples

Edge cases

Any 2-rig \mathcal{C} , endowed with the **trivial derivation** $\mathcal{C} \rightarrow \mathcal{C}$ that is the constant functor at the initial object (regarded as empty coproduct).

Let P be a **distributive lattice**; the identity functor $P \rightarrow P$ is, trivially, a derivation (because every element of P is \vee -idempotent).

Idempotents

let \mathcal{C} be a 2-rig; an object $E \in \mathcal{C}$ is **\cup -idempotent** if $E \cup E \cong E$. Given an idempotent object, the endofunctor $\partial_E : A \mapsto E \otimes A$ is a derivation, since

$$\begin{aligned}\partial_E(A \otimes B) &= E \otimes A \otimes B \cong (E \cup E) \otimes A \otimes B \\ &= (E \otimes A) \otimes B \cup A \otimes (E \otimes B).\end{aligned}$$

More generally, given a derivation ∂ on a 2-rig \mathcal{C} , the endofunctor $A \mapsto \partial A \cup E \otimes A$ is a derivation: the former is a particular case of this construction when $\partial = 0$ is the trivial derivation.

Polynomials

Let \mathcal{C} be a 2-rig; let $Y : \mathcal{C} \rightarrow \mathcal{C}$ a strong monoidal endofunctor, we define a category of **\mathcal{C} -valued polynomials**:

- objects are ‘polynomials’ $\sum_{i=0}^d A_i \otimes Y^i$, regarded as endofunctors $\mathcal{C} \rightarrow \mathcal{C}$, with the convention that $Y^0 = 1_{\mathcal{C}}$ and the action on an object X is given by

$$X \mapsto \sum_{i=0}^d A_i \otimes Y^i(X) \in \mathcal{C};$$

- morphisms are natural transformations of functors.

The category $\mathcal{C}[Y]$ so obtained is a 2-rig where the sum is ‘component-wise’, and the \otimes -product is a similar ‘Cauchy product’ of polynomials.

Polynomials

let \mathcal{C} be a 2-rig; $\mathcal{C}[Y]$ becomes a differential 2-rig if we endow \mathcal{C} with the trivial derivation, and we put $\partial Y = I$, suitably extended on a generic expression $\sum_{i=0}^d A_i \otimes Y^i$ by linearity and Leibniz rule:

$$\partial \left(\sum_{i=0}^d A_i \otimes Y^i \right) \cong \sum_{i=1}^d i \cdot A_i \otimes Y^{i-1}$$

let \mathcal{C} be a differential 2-rig, with derivation denoted $a \mapsto \partial a$. One can define the **2-rig of differential polynomials with coefficients in \mathcal{C}** introducing an infinite set of ‘variables’

$\mathcal{Y} := \{Y, Y^{(1)}, Y^{(2)} \dots, Y^{(n)}, \dots\}$ and defining

$$\mathcal{C}[\mathcal{Y}] := \varinjlim \left(\mathcal{C} \rightarrow \mathcal{C}[Y] \rightarrow \mathcal{C}[Y, Y^{(1)}] \rightarrow \mathcal{C}[Y, Y^{(1)}, Y^{(2)}] \rightarrow \dots \right)$$

where we define inductively $\mathcal{C}[Y, Z] := \mathcal{C}[Y][Z]$, and the derivation as $\partial : Y^{(i)} \mapsto Y^{(i+1)}$.

A **co-Heyting algebra** is a bounded distributive lattice K such that $x \vee - : K \rightarrow K$ has a left adjoint $-\backslash x$ for all $x \in K$:

$$y \backslash x \leq z \iff y \leq x \vee z$$

Define $\partial x := x \wedge \bot$ it's easy to see that $\partial : K \rightarrow K$ is a derivation when K is regarded as a distributive 2-rig. Leibniz rule takes the form

$$\partial(a \wedge b) = (\partial a \wedge b) \vee (a \wedge \partial b).$$

Example: lattices of closed subsets of topological spaces.

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Example: lattice of subtoposes of a given topos \mathcal{E} ; define the **boundary** $\partial\mathcal{A}$ of a subtopos $\mathcal{A} \subseteq \mathcal{E}$ in this lattice and then in turn the boundary ∂T of the **geometric theory** T that \mathcal{A} classifies (see Caramello, 2009).

Results

Theorem

Let \mathcal{C} be a 2-rig, and M a internal semigroup with multiplication $m : M \otimes M \rightarrow M$; then the map $\partial m : \partial M \otimes M \cup M \otimes \partial M \rightarrow \partial M$ splits as a pair of maps

$$\begin{cases} i_R : \partial M \otimes M \rightarrow \partial M \\ i_L : M \otimes \partial M \rightarrow \partial M \end{cases}$$

Then, i_R (resp., i_L) is a right (resp., left) action of M over ∂M .

- Let \mathcal{C} be a 2-rig with a natural number object \mathbf{N} ; then, \mathbf{N} is a monoid in a canonical way, with respect to the morphism $\mathbf{N} \times \mathbf{N} \rightarrow \mathbf{N} : \lambda p q. s^p q$.

If \mathcal{C} is a differential bicartesian category, then $\partial\mathbf{N}$ is a **Lawvere dynamical system**.

- Let \mathcal{C} be an elementary topos, regarded as a bicartesian category; if \mathcal{C} has a differential structure, the derivative $\partial\Omega$ of the subobject classifier is a module for the monoid $(\Omega, \wedge, \text{true})$.

Theorem

Let \mathcal{C} be a category that satisfies the following assumptions:

- additive category with byproducts noted \oplus ;
- it has the structure of a differential 2-rig with multiplicative structure \odot , and a derivation ∂ that is \oplus -linear and \odot -Leibniz.

Then there exists a canonical extension $\bar{\partial}$ of ∂ to the additive presheaf category $\hat{\mathcal{C}} = [\mathcal{C}, \text{Ab}]$ (coproduct-linear and convolution-Leibniz), that hence becomes a differential 2-rig.

$$\begin{array}{ccc} \mathcal{C}^{\text{op}} & \xrightarrow{\partial} & \mathcal{C}^{\text{op}} \\ y \downarrow & \swarrow & \downarrow y \\ \check{\mathcal{C}} & \xrightarrow{\bar{\partial}} & \check{\mathcal{C}} \end{array}$$

Recall the equivalence

$$[\mathbf{Fin}, \mathbf{Set}] \cong [\mathbf{Set}, \mathbf{Set}]_\omega$$

given by **left Kan extension** along the embedding $J : \mathbf{Fin} \rightarrow \mathbf{Set}$, where at the right hand side we put *finitary* endofunctors of \mathbf{Set} .

This equivalence can be promoted to a **monoidal** equivalence, if $[\mathbf{Set}, \mathbf{Set}]_\omega$ is considered a monoidal category with respect to composition.

Given $F : \mathbf{Fin} \rightarrow \mathbf{Set}$, $T_F = \mathbf{Lan}_J F$ is the associated finitary functor and there exists a unique monoidal structure \diamond such that

$$\mathbf{Lan}_J(F \diamond G) \cong \mathbf{Lan}_J F \circ \mathbf{Lan}_J G$$

Theorem (The chain rule)

Let $F, G : \mathbf{Fin} \rightarrow \mathbf{Set}$ be non- Σ -species, and let ∂ be a derivation with respect to the cartesian monoidal structure; then,

$$\begin{cases} \partial T_G(T_F A) \times \partial T_F A \cong \partial(T_G \circ T_F) \\ \partial(G \diamond F) = \mathsf{Lan}_J \partial G(Fn) \times \partial Fn. \end{cases}$$

Questions

Vices and virtues of being $\partial 1$

In a differential ring, using the Leibniz rule: $d1 = d1 + d1$, which entails $d1 = 0$.

In a rig things are way more complicated: one has to **postulate** that $d1 = 0$, or derivation isn't even well-defined:

$$da = d(a \cdot 1) = da + a.d1 = da + \sum_{k=1}^{\infty} a.d1$$

Something similar happens in a differential 2-rig:

$$\partial I \cong \partial I \cup \partial I$$

from which we get the idea that ∂I is 'either empty or big' (e.g., in Set, ∂I is empty or -at least- countable.)

Vices and virtues of being $\partial 1$

Theorem

There is no nontrivial finite colimit-preserving derivation on the 2-rig $(\mathbf{Fin}, \times, 1)$ of finite sets and functions: such $\partial : \mathbf{Fin} \rightarrow \mathbf{Fin}$ is completely determined by its action on the point, so that

$$\partial A \cong \partial(A \cdot 1) \cong A \cdot \partial 1.$$

Same in the category of finite dimensional vector spaces, where $d = \dim V = 2d$ has 0 (thus the zero space) as unique solution. Same in every category with a choice of **dimension** $\mathcal{C} \rightarrow \mathbb{N}$ for objects.

Vices and virtues of being $\partial 1$

Sometimes $\partial 1 = \partial 1 + \partial 1$ is forced to have just trivial solutions due to naturality;

Let ∂ be a derivation in a category of functors $\mathcal{A} \rightarrow \mathbf{Set}$; then, $\partial 1$ is a functor $\mathcal{A} \rightarrow \mathbf{Set}$ such that $F \cong F + F$, and naturally so; such functors must be constant on connected components of \mathcal{A} .

The unbearable largeness of $\partial 1$

Use Yoneda lemma.

Consider the hom-set $\text{hom}(\partial 1, Z)$ for a generic object Z ;

$\partial 1 \cong \partial 1 + \partial 1$ yields

$$\text{hom}(\partial 1, Z) \cong \text{hom}(\partial 1 + \partial 1, Z) \cong \text{hom}(\partial 1, Z) \times \text{hom}(\partial 1, Z).$$

So, $\text{hom}(\partial 1, Z)$ can only be empty, a singleton or infinite.

Thus if a category is finite $\text{hom}(\partial 1, Z)$ is either empty or a singleton, and in particular it must be a singleton when $Z = \partial 1$.

Coalgebras

If $\partial 1 \cong \partial 1 + \partial 1$, this means that $\partial 1$ is naturally a coalgebra for the "leave it or double it" functor $S : A \mapsto A + A$, in such a way that there is a unique map

$$\begin{array}{ccc} \partial 1 & \cong & \partial 1 + \partial 1 \\ \downarrow & & \downarrow \\ C & \cong & C + C \end{array}$$

between $\partial 1$ and the terminal coalgebra of S ; but wait, in the category of topological space C is the Cantor set! **What just happened here?**

Napier objects

As an endofunctor, ∂ might have interesting fixed points, and there is a standard procedure to build its initial algebra and terminal coalgebra.

Initial algebras are trivial, in that $\partial 0 = 0$ by using the Leibniz rule. On the other hand, the triviality of terminal coalgebras is governed by the shape of $\partial 1$:

$$1 \leftarrow \partial 1 \leftarrow \partial \partial 1 \leftarrow \partial \partial \partial 1 \leftarrow \dots$$

and the first ordinal λ for which the transition morphism $v : \partial^\lambda 1 \leftarrow \partial^{\lambda+1} 1$ is invertible realises the terminal coalgebra.

Derivatives of models

Let (\mathcal{C}, J) be a *Grothendieck monoidal site* i.e. a Grothendieck topology such that the category of sheaves is monoidal; let $\partial : \hat{\mathcal{C}} \rightarrow \hat{\mathcal{C}}$ be a convolution-derivation; if ∂ is such that

a big black hole of ignorance

then, ∂ restricts to a derivation on the category of J -sheaves (that is itself monoidal).

Derivatives of models

Let \mathcal{T} be an algebraic theory of some sort, with the property that $\text{Mod}(\mathcal{T})$ is monoidal; let $\partial : \hat{\mathcal{T}} \rightarrow \hat{\mathcal{T}}$ be a convolution-derivation; if ∂ is such that

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then ∂ sends a \mathcal{T} -model to a \mathcal{T} -model.

