

Profunctorial Semantics I

Fosco Loregian 

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Algebraic structures

A **group** is a set equipped with operations

- $m : G \times G \rightarrow G$
- $i : G \rightarrow G$
- $e : 1 \rightarrow G$

...

you know the drill

Algebraic structures

Theorem (Higman-Neumann 1953)

A **group** is a set equipped with a single binary operation $/ : G \times G \rightarrow G$ subject to the single equation

$$x / (((x/x)/y)/z) / (((x/x)/x)/z) = y$$

for every $x, y, z \in X$.

Well.

This is awkward.

The theory of equationally definable classes of algebras, initiated by Birkhoff in the early thirties, is [...] hampered in its usefulness by two defects. [...] the second is the awkwardness inherent in the presentation of an equationally definable class in terms of operations and equations.

Quite recently, Lawvere, by introducing the notion - closely akin to the clones P. Hall - of an algebraic theory, rectified the second defect.

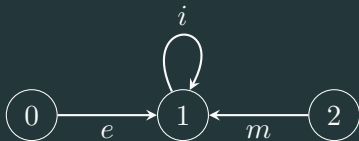
Definition

An **operator domain** is a sequence $\underline{\Omega} = (\Omega_n \mid n \in \mathbb{N})$; the elements of Ω_n are called **operations** of **arity** n .

Definition

An **interpretation** \underline{E} of an operator domain $\underline{\Omega}$ consists of a pair $(E, (f_\omega \mid \omega \in \Omega_n, n \in \mathbb{N}))$ where $f_\omega : E^n \rightarrow E$ is an n -ary operation on the set E called the *carrier* of \underline{E} .

An operator domain can be represented as a (rooted) graph: for example, for groups



Way better to use functors.

A **Lawvere theory** is an identity-on-objects functor $p : \mathbf{Fin}^{\text{op}} \rightarrow \mathcal{L}$ that commutes with finite products.

Unwinding the definition:

- \mathcal{L} is a category with the same objects as \mathbf{Fin} , the category of finite sets and functions;
- p is a functor that acts trivially on objects
- The only thing that can change between \mathbf{Fin} and \mathcal{L} is the number of morphisms $[n] \rightarrow [m]$.

Equivalently: p is a **promonad** on the opposite of \mathbf{Fin} , regarded as an object of the bicategory of profunctors, that preserves the monoidal structure. \mathcal{L} is the Kleisli object of p .

$$\left\{ \begin{array}{l} \text{identity on obj} \\ \text{left adjoints} \\ p: [\mathcal{L}, \mathbf{Set}] \rightarrow [\mathbf{Fin}^{\text{op}}, \mathbf{Set}] \end{array} \right\} \rightleftharpoons \left\{ \begin{array}{l} \text{monads in Prof} \\ p: \mathbf{Fin}^{\text{op}} \rightsquigarrow \mathbf{Fin}^{\text{op}} \end{array} \right\}$$

- The trivial theory is the identity functor $1_{\text{Fin}} : \text{Fin}^{\text{op}} \rightarrow \text{Fin}^{\text{op}}$
- Since p preserves products, it is **uniquely determined by its value on $[1]$** . This means that if $p : \text{Fin}^{\text{op}} \rightarrow \mathcal{L}$ is a Lawvere theory, then every object of \mathcal{L} is X^n if $p[1] = X$.
- The only difference between Fin and \mathcal{L} is thus the set of morphisms $[n] \rightarrow [m]$, added on top of those in Fin .

$$\mathcal{L}_{\text{Grp}} =$$

$$\begin{array}{ccccc}
 & & i & & \\
 & & \text{⤿} & & \\
 [0] & \xrightarrow{e} & [1] & \xleftarrow{m} & [2]
 \end{array}$$

A **model** for a Lawvere theory p is a product-preserving functor $\ell : \mathcal{D} \rightarrow \mathbf{Set}$.

The category $\mathbf{Mod}(p)$ for a Lawvere theory is a full, reflective subcategory of the category $[\mathcal{L}, \mathbf{Set}]$ of all functors $\mathcal{D} \rightarrow \mathbf{Set}$.

Theorem

The following conditions are equivalent:

- *ℓ is a model for a Lawvere theory \mathcal{L} ;*
- *The composition $\ell \circ p$ preserves finite products;*
- *The composition $\ell \circ p$ is representable (with respect to the inclusion $J : \mathbf{Fin} \rightarrow \mathbf{Set}$), i.e.*

$$\ell(X[n]) \cong \mathbf{Set}(J[n], A)$$

for some $A \in \mathbf{Set}$.

As a consequence, the square

$$\begin{array}{ccc}
 \text{Mod}(p) & \xrightarrow{r} & [\mathcal{D}, \text{Set}] \\
 u \downarrow & & \downarrow - \circ X \\
 \text{Set} & \xrightarrow{[J, 1]} & [\text{Fin}^{\text{op}}, \text{Set}]
 \end{array}$$

is a pullback.

- $\text{Mod}(p)$ is a **reflective** subcategory of $[\mathcal{D}, \text{Set}]$. We write $r_! \dashv r$ for the resulting adjunction.
- The functor u is **monadic**, with left adjoint f .
- This sets up a functor

$$\mathfrak{M} : \text{Th}_L(\text{Fin}) \rightarrow \text{Mnd}_{<\omega}(\text{Set})$$

because the monad uf above is finitary.

There is an equivalence of categories between $\text{Th}_L(\text{Fin})$ and $\text{Mnd}_{<\omega}(\text{Set})$.

We have to construct a functor in the opposite direction,

$\exists : \text{Mnd}_{<\omega}(\text{Set}) \rightarrow \text{Th}_L(\text{Fin})$; given T , we consider the composition $\text{Fin} \hookrightarrow \text{Set} \xrightarrow{F^T} \text{Set}^T$ and its bo-ff factorization, in a square

$$\begin{array}{ccc} \mathcal{L}^{\text{op}} & \xrightarrow{ff} & \text{Set}^T \\ \text{bo} \uparrow & & \uparrow F^T \\ \text{Fin} & \xrightarrow{J} & \text{Set} \end{array}$$

- the left vertical arrow is a Lawvere theory almost by definition.
- Set^T has the universal property of the category of \mathcal{L} -models.

Theories as promonads

There is a 2-monad $\tilde{S} : \mathbf{Prof} \rightarrow \mathbf{Prof}$ whose algebras are exactly promonoidal categories.

Given a profunctor $p : \mathcal{A} \rightsquigarrow \mathcal{B}$ between promonoidal categories $(\mathcal{A}, \mathfrak{P}, J_A), (\mathcal{B}, \mathfrak{Q}, J_B)$:

- p is a **pseudo- \tilde{S} -algebra** morphism;
- The cocontinuous left adjoint \hat{p} associated to p is **strong monoidal** with respect to the convolution monoidal product on presheaf categories;

Assume the promonoidal structures $\mathfrak{P}, \mathfrak{Q}$ on \mathcal{A}, \mathcal{B} are representable; then, the conditions above are in turn equivalent to

- Both mates $p^{\triangleleft} : \mathcal{A} \rightarrow P\mathcal{B}$ and $p^{\triangleleft} : \mathcal{B} \rightarrow P^*\mathcal{A}$ are **strong monoidal** wrt convolution on their codomains.

Theories as $[\mathbf{Fin}, \mathbf{Set}]$ -categories

Theorem

$$[\mathbf{Fin}, \mathbf{Set}] \cong \mathbf{End}_{<\omega}(\mathbf{Set})$$

If the LHS is endowed with the monoidal structure induced by composition of endofunctors; this is called the **substitution** monoidal product of functors $F, G : \mathbf{Fin} \rightarrow \mathbf{Set}$:

$$F * G : m \mapsto \int^n Fn \times (Gm)^n$$

The substitution monoidal product is a highly non-symmetric, right closed monoidal structure (not left closed).

The category $[\mathbf{Fin}, \mathbf{Set}]$ works as base of enrichment.

Theories as $[\mathbf{Fin}, \mathbf{Set}]$ -categories

From [Garner]

From now on we blur the distinction between the categories $[\mathbf{Fin}, \mathbf{Set}] \cong \mathbf{End}_{<\omega}(\mathbf{Set})$:

- A **finitary monad** is a monoid in $\mathbf{End}_{<\omega}(\mathbf{Set})$, i.e. a $\mathbf{End}_{<\omega}(\mathbf{Set})$ -category with a single object, i.e. a $[\mathbf{Fin}, \mathbf{Set}]$ -category with a single object.
- A Lawvere theory is a $[\mathbf{Fin}, \mathbf{Set}]$ -category that is **absolute** (Cauchy-, Karoubi-)**complete** as an enriched category and generated by a single object.
- Lawvere theories form a reflective subcategory in finitary monads; reflection is the enriched **Cauchy completion** functor.

Theories as $[\mathbf{Fin}, \mathbf{Set}]$ -categories

Equivalently,

- A Lawvere $[\mathbf{Fin}, \mathbf{Set}]$ -category is an enriched category where every object A is the tensor $y[n] \odot X$ for a distinguished object $X \cong y[1] \odot X$. All such categories are enriched-Cauchy complete.
- A $[\mathbf{Fin}, \mathbf{Set}]$ -category is a special kind of **cartesian multicategory**: one where a multimorphism $f : X_1 \dots X_n \rightarrow Y$ is such that $X_1 = X_2 = \dots = X_n$.

Generalisations/extensions:

- let \mathbb{N} be the **discrete** category over natural numbers;
- let \mathbf{P} be the **groupoid** of natural numbers;

The categories $[\mathbb{N}, \mathbf{Set}]$ and $[\mathbf{P}, \mathbf{Set}]$ become monoidal with respect to substitution products $*_N, *_P$:

$$F *_N G = \coprod_{k \in \mathbb{N}} G_k \times \coprod_{\vec{n} | \sum n_i = n} X_{n_1} \times \cdots \times X_{n_k}$$

$$F *_P G = \int^{k, \vec{n}} Y_k \times X_{n_1} \times \cdots \times X_{n_k} \times \mathbf{P}(\sum n_i, n)$$

PRO(P)S

$*_N$ and $*_P$ -monoids are respectively non-symmetric and symmetric operads.

- A **PRO** is an identity-on-objects strong monoidal functor $p : \mathbb{N} \rightarrow \mathcal{P}$. \mathcal{P} is possibly non-cartesian.
- A **PROP** is an identity-on-objects strong monoidal functor $p : \mathbb{N} \rightarrow \mathcal{P}$. \mathcal{P} is symmetric monoidal.

These are, of course, other examples of promonoidal promonads.

PRO(P)s and operads

Every PRO $p : \mathbf{Fin}^{\text{op}} \rightarrow \mathcal{T}$ gives rise to the operad $O(\mathcal{T}) = (\mathcal{T}(n, 1) \mid n \in \mathbb{N})$. 2. Conversely, any operad $(\mathcal{O}(n) \mid n \in \mathbb{N})$ gives rise to a pro $T(\mathcal{O})$, where

$$T(\mathcal{O})(n, m) = \coprod_{k_1 + \dots + k_m = n} \mathcal{O}(k_1) \times \dots \times \mathcal{O}(k_m).$$

(It would be helpful to imagine a picture of m trees stacked vertically.)

If we begin with an operad \mathcal{O} , we have $\mathcal{O} = O(T(\mathcal{O}))$. (This is because $T(\mathcal{O})(n, 1) = \mathcal{O}(n)$, according to the above formula.)

On the other hand, if we start with a PRO \mathcal{T} , then there exists a canonical map of PROs $T(O(\mathcal{T})) \rightarrow \mathcal{T}$, given by, for each n and m , a canonical function

$$\coprod_{k_1 + \dots + k_m = n} T(k_1, 1) \times \dots \times T(k_m, 1) \rightarrow T(n, m) \quad (\star)$$

induced from the monoidal product on T .

This sets up an adjunction

$$T : \mathbf{Opd}[\mathbf{S}] \rightleftarrows \mathbf{PRO}[\mathbf{P}] : O$$

with fully faithful left adjoint, so that [symmetric] operads can be regarded as a $\mathbf{PRO}[\mathbf{P}]$ s \mathcal{T} such that each function (\star) is bijective.

Re-enact [Garner] away from \mathbf{Set} .

Let \mathcal{V} be a locally presentable base of enrichment; let $\mathfrak{F}(\mathcal{V})$ be the subcategory of finitely presentable objects:

- $\mathfrak{F}(\mathcal{V})$ is the free **finite weighted** cocompletion of the point;
- There is a strong monoidal equivalence of categories

$$[\mathfrak{F}(\mathcal{V}), \mathcal{V}] \cong [\mathcal{V}, \mathcal{V}]_{<\omega}$$

between functors $\mathfrak{F}(\mathcal{V}) \rightarrow \mathcal{V}$ and finitary endo- \mathcal{V} -functors;

- The \mathcal{V} -substitution product on LHS is

$$F * G = A \mapsto \int^B FB \otimes_{\mathcal{V}} (GA)^B$$

- There is an equivalence of categories between **finitary \mathcal{V} -monads** and **enriched-Cauchy-complete categories** generated by a single object under iterated finite powers.
- **Models** for a Lawvere theory correspond to **algebras** for the associated finitary monad; free models are free algebras are representables in

$$\begin{aligned} \text{Alg}(T, \mathcal{C}) &= [\mathfrak{F}(\mathcal{V}), \mathcal{V}]\text{-Cat}(T, \mathcal{C}) \\ &\cong [\mathfrak{F}(\mathcal{V}), \mathcal{V}]\text{-Cat}(\hat{T}, \mathcal{C}) \\ &= \text{Mod}(\hat{T}, \mathcal{C}) \end{aligned}$$

class of lims	finite \times	\mathbb{D} -lims	finite powers	weighted \mathbb{D} -limits	bicat \times
theory	\mathbf{Fin}^{op}	completion of $\{*\}$	completion of $\{*\}$	completion of $\{*\}$	completion of $\{*\}$
semantics	\mathbf{Set}	\mathbf{Set}	\mathcal{V}	\mathcal{V}	\mathbf{Prof}
eq. with	finitary	\mathbb{D} -accessible	$[\mathfrak{F}(V), V]$ -monoids	$[?, V]$ -monoids	???

Profunctorial semantics

- Characterise the free carbicat $\mathbb{CB}(\ast)$ on a singleton: see [link here](#));
- Check if the univ property of \mathbf{Fin} remains true for $\mathbb{CB}(\ast)$;
- Take $\mathbb{CB}(\ast) = F$, and consider its **free cocompletion** in the bicolimit sense
- Prove that

$$\begin{aligned}[PF, PF] &\cong [\mathbb{CB}(\ast), PF] \\ &\cong PF\end{aligned}$$

monoidally; \odot -**monoids** := monoids in PF wrt composition in $[PF, PF]$.

Profunctorial semantics

- Prove that there is a **syntax-VS-semantics** adjunction here: theories are promonoidal promonads T on (a 1-skeleton of) $\mathbb{CB}(\ast)$, and models are carbicat homomorphisms $\mathbf{Kl}(T) \rightarrow \mathbf{Prof}$. There is an equivalence

$$\{\text{theories}\} \cong \{\text{??? monads}\}$$

- Let PROs come into play: analogue of the adjunction between PROs and operads.

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