

# WEI Notes

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Many of the details of the proofs here come from some old notes by David Evans.

## 1 Projective and semiprojective geometries

For  $\Gamma \subset \text{Aut}(M)$ , write  $x^\Gamma := \{x^\gamma : \gamma \in \Gamma\}$ . Recall the following definitions:

- Definition 1.1.** (i) A **projective geometry**  $J$  is obtained from a linear geometry by factoring out the equivalence relation  $\text{acl}(x) = \text{acl}(y)$ .
- (ii) A **semiprojective geometry**  $J$  is obtained from a basic linear geometry by factoring out the equivalence relation  $x^Z = y^Z$ , where  $Z$  is the centre of the automorphism group.

**Remark 1.2.** Note that, since we start with a basic linear geometry, the elements of the underlying field  $K$  are fixed. Recall that, for any  $K$ -vector space  $V$ , we have  $\text{Aut}(V/K) \cong K^\times$ .<sup>1</sup> Thus, we get that  $x^Z$  is the set of  $\alpha x$  for  $\alpha \in K^\times$  respecting all of the additional structure of the basic linear geometry.

**Example 1.3.** Let  $M = (V, K, L, \beta)$ , where  $V$  is a  $K$ -vector space,  $L$  is a  $K$ -line, and  $\beta$  is a symplectic form with respect to some  $\sigma$ . Take  $f \in \text{Aut}(M)$ . Since in particular  $f$  is an automorphism of  $V$  fixing  $K$ , it sends  $v \mapsto \alpha v$  for some fixed  $\alpha \in K^\times$ . Moreover,  $f(\beta(x, y)) = \beta(f(x), f(y))$ . But  $\beta(f(x), f(y)) = \beta(\alpha x, \alpha y) = \alpha^2 \beta(x, y)$ , and  $f(\beta(x, y)) = \beta(x, y)$ , since  $f$  fixes  $K$ . Therefore,  $\alpha \in \{\pm 1\}$ . In this case, for  $x \in V$ , we get  $x^Z = \{x, -x\}$ .

**Lemma 1.4** (2.3.6). *Let  $J$  be a basic semiprojective geometry. Then  $J$  has weak elimination of imaginaries.*

*Proof.* Let  $a \in J^{\text{eq}}$ , write  $J = V/\sim$  where  $V$  is a basic linear geometry, and let  $A := \text{acl}(a) \cap V$ . Note that by Remark 1.2, we have  $J \subseteq V^{\text{eq}}$ , and hence,  $a \in V^{\text{eq}}$ . So, by Lemma 2.3.5, it follows that  $a \in \text{dcl}^{\text{eq}}(A)$ . Write  $B := \text{acl}(a) \cap J$ . For  $g \in \text{Aut}(J/B)$ , we can find some  $h \in \text{Aut}(V/A)$  inducing  $g$ , and so, in particular, if  $\bar{c}$  is a tuple in  $J$  such that  $a$  is  $\bar{c}$ -definable, the orbit of  $\bar{c}$  in  $J$  over  $B$  is the same as that over  $A$ . Therefore,  $a \in \text{dcl}^{\text{eq}}(B)$ .  $\square$

<sup>1</sup>Indeed, if  $f \in \text{Aut}(V/K)$  is such that  $v \mapsto \alpha v + \sum_i \alpha_i v_i$  (a linearly independent expression), then  $f(0_V) = 0_V + \sum_i \alpha_i v_i = 0_V$ , which implies that  $\alpha_i = 0_V$  for all  $i$ . So any automorphism is given by multiplication by a non-zero scalar.

**Remark 1.5** (cf. 2.3.7). A projective geometry  $J$  does not have weak elimination of imaginaries. As an example, suppose  $J$  is an inner product space with a symplectic form  $\beta$ . First, let  $J_s$  be the semiprojective geometry associated to  $V$ , the basic linear geometry covering  $J$ . Then  $J_s \subseteq J^{\text{eq}}$ . Let  $a \in J_s$ , so that  $a = \{\pm v\}$  for some  $v \in V$ . We claim that  $a \notin \text{dcl}^{\text{eq}}(\text{acl}(a) \cap J)$ : indeed, if  $|K| \geq 4$ , there is an automorphism sending  $v \mapsto \alpha v$  for some  $\alpha \notin \{0, \pm 1\}$ , which clearly fixes  $\langle v \rangle = \text{acl}(a) \cap J$ .

## 2 Affine geometries

**Definition 2.1.** Let  $V$  be a definable  $K$ -vector space and let  $A$  be a definable set on which  $V$  acts regularly. A  **$K$ -affine map** is a map  $\lambda: A \rightarrow K$  such that

$$\lambda\left(\sum_i \alpha_i a_i\right) = \sum_i \alpha_i \lambda(a_i),$$

where  $\alpha_i \in K$  such that  $\sum_i \alpha_i = 1$ .

**Remark 2.2.** Fix a base point  $a \in A$ . Since  $V$  acts regularly on  $A$ , we can write  $a_i = a + v_i$  for some  $v_i \in V$ . Thus,  $\sum_i \alpha_i a_i = a + (\sum_i \alpha_i v_i)$  and  $\sum_i \alpha_i v_i \in V$ , but we need to check that this is well-defined, i.e., that it does not depend on our choice of base point. So choose some other  $a' \in A$  and write  $a_i = a' + v'_i$ . Once again, since  $V$  acts regularly on  $A$ , we can write  $a' = a + v'$  for some  $v' \in V$ , and hence  $v_i = v'_i + v'$ . Therefore, the left-hand side is well-defined.

**Definition 2.3.** Let  $V^*$  be the **definable dual** of  $V$  (i.e., the set of definable linear functionals), and let  $A^*$  be the set of  $M$ -definable  $K$ -affine maps on  $A$ .

**Lemma 2.4** (2.3.9). *There is an exact sequence*

$$(0) \longrightarrow K \xrightarrow{f} A^* \xrightarrow{g} V^* \longrightarrow (0),$$

i.e.,  $f$  is injective,  $g$  is surjective, and  $\text{im } f = \ker g$ .

*Proof.* We first define  $f$  by  $k \mapsto \lambda_k$ , where  $\lambda_k: A \rightarrow K$  is the constant map in  $k$ . This is clearly injective.

Now, define  $g$  by  $\lambda \mapsto \lambda'$ , where we define, for  $v \in V$ ,  $\lambda'(v) := \lambda(a+v) - \lambda(a)$  for some  $a \in A$ . We need to check that  $\lambda'$  is well-defined and linear, so that  $\lambda' \in V^*$ . We first show linearity with respect to a choice of base point  $a$ :

$$\begin{aligned} \lambda'(v + v') &= \lambda(a + v + v') - \lambda(a) \\ &= \lambda((a + v) + (a + v') - a) - \lambda(a) \\ &= \lambda(a + v) + \lambda(a + v') - 2\lambda(a) \\ &= \lambda'(v) + \lambda'(v'). \end{aligned}$$

Having this, we can now show that  $\lambda'$  is well-defined: let  $a' \in A$  be some other base point, and write  $a' = a + v'$ . Then:

$$\begin{aligned}\lambda(a' + v) - \lambda(a') &= \lambda(a + v' + v) - \lambda(a + v') \\ &= \lambda'(v' + v) + \lambda(a) - \lambda'(v) - \lambda(a) \\ &= \lambda'(v').\end{aligned}$$

To show that  $g$  is surjective, take some  $\lambda^* \in V^*$ . Choose a base point  $a \in A$ , set  $\lambda(a) := 0$ , and, for any other  $a' \in A$ , define  $\lambda(a') := \lambda^*(v')$ , where  $a' = a + v'$ .

Finally, we need to show that  $\text{im } f = \ker g$ , i.e., that  $g(\lambda_k)$  is the constant function at 0. But this is clear:  $\lambda'(v) = \lambda(a + v) - \lambda(a) = k - k = 0$ .  $\square$

**Remark 2.5** (2.3.10). (i) Recall the following corollary of quantifier elimination:

**Corollary 2.6** (2.2.9). *The definable linear functions on the vector space  $V$  in a linear geometry are those afforded either by the inner product (if one is given, or comes from a quadratic form), or by the dual in the polar case.*

In particular, in the case of a pure vector space, any linear form in  $V^*$  must be definable using only names for elements in  $K$ , i.e., must be constant. But, by the argument above, the only constant linear form in  $V^*$  is the 0 function. So, in this case,  $A^* = K$  and  $V^* = (0)$ . Clearly this is also the case in the degenerate space.

- (ii) We can code  $A^*$  in  $(V, V^*, A)^{\text{eq}}$ : given some  $\lambda' \in V^*$ ,  $a \in A$ , and  $k \in K$ , we can associate a unique  $\lambda \in A^*$  to  $(\lambda, a, k)$  by declaring  $\lambda(a + v) := \lambda'(v) + k$  for all  $v \in V$ . Thus, quotienting out by the  $\emptyset$ -definable relation  $(\lambda_1, a_1, k_1) \sim (\lambda_2, a_2, k_2)$  iff the corresponding  $K$ -affine maps are equal, we see that  $A^*$  is identified with a sort in  $(V, V^*, A)^{\text{eq}}$ . In particular, the algebraic closure of  $\lambda \in A^*$  in  $(V, V^*, A)$  is the line  $\langle \lambda' \rangle$  in  $V^*$  for  $\lambda' \in V^*$  corresponding to it, and hence,  $(V, V^*, A)$  does not have weak elimination of imaginaries.

Moreover, if we have a non-degenerate bilinear form, we can write each  $\lambda \in V^*$  as  $\lambda(x) = \beta(v, x)$  for some fixed  $v \in V$ , and so we do not have to mention  $V^*$  explicitly.

- (iii) We do have weak elimination of imaginaries in  $(V, V^*, A^*)$ , but this is not stably embedded in  $(V, V^*, A, A^*)$ . Indeed, note that, in  $(V, V^*, A, A^*)$ , after choosing a base point for  $A$ , we can definably identify  $A^*$  with  $K \oplus V^*$ . If  $(V, V^*, A^*)$  is stably embedded in  $(V, V^*, A, A^*)$ , then we must be able to define the isomorphism without invoking any elements from  $A$ . But this is impossible.
- (iv) The surjective map from the previous lemma shows that  $V^*$  is definable in  $(V, A, A^*)$ , so we do not need to include it in the geometry whenever  $A^*$  is already included (even in the polar case!).

**Lemma 2.7** (2.3.11). *Let  $V$  be a vector space, let  $J$  be a basic, linear, nonquadratic geometry covered by  $V$  (or with  $V$  one of the two vector spaces if  $J$  is polar), and let  $A$  be a definable subset on which  $V$  acts regularly. Then  $(J, A, A^*)$  admits quantifier elimination in its natural language.*

*Proof.* We first need to describe the language. This will include:

- The original language for  $J$ .
- Predicates for  $A$  and  $A^*$ .
- Maps  $+$ :  $V \times A \rightarrow A$  to be interpreted as  $(v, a) \mapsto a + v$  and  $-$ :  $A \times A \rightarrow V$  as  $(a_1, a_2) \mapsto v$  where  $a_2 = a_1 + v$ .
- A map  $\text{ev}$ :  $A \times A^* \rightarrow K$  to be interpreted as  $(a, \lambda) \mapsto \lambda(a)$ .
- A  $K$ -vector structure on  $A^*$ .
- Names for the constant functions in  $A^*$ .
- If  $V^*$  is explicitly included in  $J$ , a map  $A^* \rightarrow V^*$  to be interpreted as  $\lambda$  to  $\lambda'$  as in the previous lemma; and if  $V^*$  is just definable (cf. (iv)) but not explicitly in the language, then a map  $A^* \times V \rightarrow K$  to be interpreted as  $(\lambda, v) \mapsto \lambda'(v)$ .

The idea is to run again the same proof as we saw in David's talk for quantifier elimination. Let  $\bar{b}$  and  $\bar{c}$  be tuples from  $J \cup A \cup A^*$  with the same quantifier-free type and  $d \in J \cup A \cup A^*$ . We want to find some  $e \in J \cup A \cup A^*$  such that  $\text{qftp}(\bar{b} \frown d) = \text{qftp}(\bar{c} \frown e)$ . We can split the proof into several cases:

Case 1:  $\bar{b} \subseteq J$ .

Then, by the proof of quantifier elimination for basic linear geometries, there is some  $g \in \text{Aut}(J \cup A)$  with  $g(\bar{b}) = \bar{c}$ , and so we can take  $e := g(d)$ .

Case 2: There is some  $a_0 \in \bar{b} \cap A$ .

Let us write  $\bar{b} = (a_0, \bar{b}_1)$ . Applying a translation, write  $\bar{c} = (a_0, \bar{c}_1)$ . Then, using  $a_0$  as a base point for  $A$ , any points from  $A$  and  $A^*$  are quantifier-free interdefinable with points in  $J$  (cf. Remark 2.5(ii)). So this reduces to the first case.

Case 3:  $\bar{b} \subseteq J \cup A^*$ .

Without loss of generality, we may assume  $\bar{b}$  is quantifier-free definably closed. Let  $\lambda_1, \dots, \lambda_r$  be the elements of  $A^*$  that appear in  $\bar{b}$ . Then, for  $a \in A$ ,  $\text{tp}(a/\bar{b})$  is determined by  $\lambda_1(a), \dots, \lambda_r(a)$ , since, for any  $v \in \bar{b} \cap J$ , we have  $\lambda_i(v + a) = \lambda'_i(v) + \lambda_i(a)$ . Let  $\lambda'_1, \dots, \lambda'_r$  be the corresponding elements from  $A^*$  in  $\bar{c}$ . By an automorphism, we may assume that  $\lambda'_i = (\lambda_i^*)'$ , so that  $\lambda_i - \lambda_i^* = k_i$  with  $k_i$  constant for all  $i$ . Remove any constant forms if necessary, and let  $q \leq r$  be maximal such that  $\lambda'_{i_1}, \dots, \lambda'_{i_q}$  are linearly independent. Then choose  $a' = v' + a$  such that  $\lambda_{i_j}^*(a') = \lambda_{i_j}(a)$ , so that  $\lambda'_{i_j}(v') = k_{i_j}$  for all  $j \leq q$ . Thus, we see that  $\text{qftp}(a/\bar{b}) = \text{qftp}(a'/\bar{c})$ . Hence, this reduces to the second case.  $\square$

**Lemma 2.8** (2.3.12). *Let  $J$  be a basic, nonquadratic, linear geometry and let  $(J, A)$  be a corresponding basic affine geometry. Then  $(J, A, A^*)$  has weak elimination of imaginaries.*

*Proof.* Since we had Lemma 2.3.3 for the affine case and  $A^*$  is algebraic over  $V^*$ , we can just mimic Nick's proof from last week. So it suffices to show that, if  $B \subseteq (J, A, A^*)^{\text{eq}}$  is algebraically closed and  $f: (J, A, A^*) \rightarrow (J, A, A^*)^{\text{eq}}$  is  $B$ -definable, then  $f$  is constant on each 1-type  $D$  over  $B$ .

As we did last time, let  $I := \{(x, y) \in D^2 : \langle xB \rangle \cap \langle yB \rangle = B\}$ , where the span is the algebraic closure in  $(J, A, A^*)$ . We claim that, if  $(x, y) \in I$ , then  $f(x) = f(y)$ . Note that, if  $B \cap A \neq \emptyset$ , then we can apply the same identification of  $A$  with  $V$  and  $A^*$  with  $K \oplus V^*$  as before and return to the linear case. So assume  $B \cap A = \emptyset$ .

Case 1:  $D \subseteq J$ .

Let  $d$  be a realisation of  $D$ . Then  $D = \text{tp}(d/B) = \text{tp}(d/B \cap J)$  (by quantifier-elimination), and we have  $\langle xB \rangle \cap \langle yB \rangle = B$  iff  $\langle xB \cap J \rangle_J \cap \langle yB \cap J \rangle_J = B \cap J$ . Moreover, since  $f$  is  $B$ -definable,  $f(d)$  is  $\langle B \cap J, d \rangle_J$ -definable, and so we may take  $\text{im } f \subseteq J^{\text{eq}}$ . So this reduces to the linear case.

Case 2:  $D \subseteq A$ .

Now  $D$  is determined by the values that affine maps  $\lambda \in B \cap A^*$  take on  $d \in D$ . Since linear maps in  $B$  are covered by affine maps, the relevant part of  $B$  for this information is  $B \cap V$ . Moreover, if  $(x, y) \in D^2$ , then, in particular,  $x \equiv_B y$ , and thus,  $\beta(x - y, v) = \beta(x, v) - \beta(y, v) = 0$  for  $v \in B \cap V$ . Thus,  $x - y$  is orthogonal to  $B \cap V$ . The only remaining information left concerns the values of  $Q(x - y)$  for a nondegenerate quadratic form (assuming there is one). For this, we need to find  $v, w \in B^\perp$  such that  $Q(v)$ ,  $Q(w)$  and  $Q(v + w)$  take on arbitrary values. But this we can do by using similar arguments as in the previous talks.

Case 3:  $D \subseteq A^*$ .

For  $(x, y) \in I$ ,  $\text{tp}(x, y/B)$  is determined by the type of the image in  $V^*$ . So once again we can run the same proof as in the linear case.  $\square$

The upshot of these results about imaginaries is the following:

**Lemma 2.9** (2.3.19). *A Lie coordinatisable structure is  $\omega$ -categorical.*