#### 1. The Cherlin-Hrushovski rank

# Definable means "definable with parameters".

**Definition 1.1.** Let  $D \subseteq \mathcal{M}$  be a definable set. We define *Cherlin-Hrushovski rank* of D, denoted  $\mathsf{CH-rk}(D)$ . First, we inductively define  $\mathsf{CH-rk}(D) \geq n$  for  $n \in \mathbb{N}$ :

- (1)  $\mathsf{CH}\text{-rk}(D) \geq 0$  if D is non-empty.
- (2)  $\mathsf{CH-rk}(D) > 0$  if D is infinite. [In particular, if  $\mathsf{CH-rk}(D) = 0$  then D is finite.]
- (3) For  $n \in \mathbb{N}$ , CH-rk $(D) \geq n+1$  if there exist definable sets  $D_1, D_2$  and definable functions  $\pi: D_1 \to D$  and  $f: D_1 \to D_2$  such that:
  - (a) For all  $d \in D$  we have CH-rk  $(\pi^{-1}(d)) = 0$ . [In particular,  $\pi$  is surjective and finite-to-one.]
  - (b)  $\mathsf{CH-rk}(D_2) > 0$ . [In particular,  $D_2$  is infinite.]
  - (c) For all  $d \in D_2$  we have  $\mathsf{CH-rk}(f^{-1}(d)) \geq n$ . [In particular, if  $n \geq 0$ , then f is also surjective.]

As is usual,  $\mathsf{CH-rk}(D) = n$  if  $\mathsf{CH-rk}(D) \ge n$  and  $\mathsf{CH-rk}(D) \not \ge n+1$ . If  $\mathsf{CH-rk}(D) \ge n$  for all  $n \in \mathbb{N}$  then we write  $\mathsf{CH-rk}(D) = \infty$ .

# Remarks 1.2.

- (1) If  $\mathcal{M}$  is not  $\aleph_0$ -categorical (in particular  $\aleph_0$ -saturated), then CH-rk should be defined in a saturated model. This will not be an issue here, as we will only work in the  $\aleph_0$ -categorical context.
- (2) The definition above can be applied either to  $\mathcal{M}$  or to  $\mathcal{M}^{eq}$ . When the distinction becomes significant, the following terminology will be used:
  - CH-rk calculated in  $\mathcal{M}$  will be referred to as *pre-rank*;
  - CH-rk calculated in  $\mathcal{M}^{eq}$  will be referred to simply as rank (somehow indicating that  $\mathcal{M}^{eq}$  is the right context to carry out these calculations).

To begin to understand this definition, we can imagine the simplifying scenario where we enforce that  $D_1 = D$  and  $\pi = \mathrm{id}$ . With this additional assumption, the definition simply says that  $\mathsf{CH-rk}(D) \geq n$  if there is a definable set  $D_2$  and a definable function  $f: D \to D_2$  whose fibres all have rank at least n. Equivalently, (again, assuming that  $D_1 = D$  and  $\pi = \mathrm{id}$ )  $\mathsf{CH-rk}(D) \geq n+1$  means that there is a uniformly definable partition[A partition of of the form  $\{\phi(x,t): t \in E\}$ , where E is a definable set.] of D into infinitely many parts of rank at least n.

The definition weakens this a little bit, since  $\mathsf{CH-rk}(D) \geq n$  if there is a definable finite cover D' of D [A definable surjective map  $\pi: D' \to D$  such that  $\pi^{-1}(d)$  is finite for all  $d \in D$ .] with  $\mathsf{CH-rk}(D') \geq n$ .

The following example, from [Sim22], illustrates why this may be necessary:

**Example 1.3.** Let  $\mathcal{M}$  be an infinite set with no extra structure (so we are working in the language of pure equality). And let  $D \subseteq \mathcal{M}^{eq}$  be the set of subsets of M of size 2 (naturally a definable set in  $\mathcal{M}^{eq}$ , as it is the quotient of the definable set  $\{(a,b): a \neq b\}$  by the definable equivalence relation  $(a,b) \sim (c,d) \iff \{a,b\} = \{c,d\}$ ). Then:

# Claim 1.3.1. CH-rk(D) $\geq 2$ .

Proof of Claim. First, we show that  $\mathsf{CH-rk}(D) \geq 2$ . Let  $D_1 = \{(a,b) \in M^2 : a \neq b\}$ ,  $D_2 = M$ ,  $f: D_1 \to D_2$  the projection onto the first coordinate and  $\pi_{\sim} : D_1 \to D$  the "selector function" for classes of the equivalence relation  $\sim$  (from above). Since  $D_2$  is infinite and so is  $f^{-1}(d)$  for all  $d \in D_2$ , this inequality follows.  $\blacksquare$  [In fact,  $\mathsf{CH-rk}(D) = 2$ , but one inequality suffices to make the point I want to make, and this note is already getting out of hand.]

**Claim 1.3.2.** Let D' be an infinite definable subset of  $\mathcal{M}^{eq}$  and  $f: D \to D'$  a definable map. Then f has finite fibres.

*Proof of Claim.* [The intuition is that we cannot definably choose an element from each member of D.]

Fix an infinite definable subset D' of  $\mathcal{M}^{eq}$  and suppose toward a contradiction that there is a definable map  $f:D\to D'$  with infinite fibres. Observe that if for some  $d\neq d'\in D'$  there are  $\{a,b\}\in f^{-1}(d)$  and  $\{a',b'\}\in f^{-1}(d')$  such that  $|\{a,b\}\cap \{a',b'\}|=1$ , then, as all elements in the same fibre are conjugates, this means that for all  $\{a,b\}\in f^{-1}(d)$  and all  $\{a',b'\}\in f^{-1}(d')$  we have that  $|\{a,b\}\cap \{a',b'\}|=1$ . This is easily seen to be impossible.

So any two distinct fibres of f must consist of pairs of elements of M with trivial intersection. But, of course, this means that f can have only one fibre, which is impossible.

[The point of the second claim is that if we enforce in the definition that  $D_1 = D$  and  $\pi = \mathrm{id}$ , then the rank of D would be 1.]

We now give two more examples with computations of CH-rk. These are lifted from [Wol20].

**Example 1.4.** Let  $\mathcal{L} = \{E_1, E_2\}$  be a language with two binary relation symbols and  $\mathcal{M}$  an  $\mathcal{L}$ -structure in which  $E_1$ ,  $E_2$  are both equivalence relations with infinitely many infinite classes, so that each  $E_1$ -class is refined by infinitely many  $E_2$ -classes.

- Let D be an  $E_2$ -class. By definition  $\mathsf{CH-rk}(D) \geq 0$ , since D is infinite. Of course, there is no partition of D into infinitely many infinite sets
- The rank of an  $E_1$  class in  $\mathcal{M}$  is 2.
- The rank of  $\mathcal{M}$  is 3.

In each case, it is easy to see why the rank is at least the given value. The other direction requires a bit more care.

The next example is also a continuation from the end of David's talk:

**Example 1.5** ([CH02, Example 2.1.11]). Let  $\mathcal{L} = \{0, +\}$  and p a fixed prime. Let  $A = \bigoplus_{n \in \omega} \mathbb{Z}/p^2\mathbb{Z}$ , i.e.

 $\mathsf{dom}(\mathcal{M}) = \{(a_i)_{i \in \omega} : a_i \in \mathbb{Z}/p^2\mathbb{Z}, \text{ and } a_i = 0 \text{ for all but finitely many } i \in \omega\}.$ 

This is, of course, a countable  $\aleph_0$ -categorical structure.

Recall that, for  $a \in A$  we write pa for  $\underbrace{a + \cdots + a}_{p \text{ times}}$ . We write A[p] for the subgroup

of A consisting of p-th powers, or, equivalently of all the elements of A of order p:

$$A[p] := \{ a \in A : pa = 0 \}.$$

This is an  $\mathbb{F}_p$ -vector space of infinite dimension, and we shall denote its projectivisation by  $J_0$ , so

$$J_0 = A[p] \setminus \{0\} / \sim$$

where  $a \sim b$  if, and only if a = rb for some  $r \in \mathbb{F}_p$ . For  $c \in A$  we will write  $A_c$  for the set  $\{a \in A : pa = c\}$ . We Lie coordinatised this structures, resulting in a tree of height 4. The calculations in [Wol20] give us:

$$\mathsf{CH}\text{-rk}(A) = 2.$$

(More precisely,  $\mathsf{CH-rk}(A_c) = 1$ , for all  $c \in A$  and  $\mathsf{CH-rk}(J_0) = 1$ .)

The examples above are special cases of Corollary 1.27

1.1. **The very basics.** We start by listing some of the most basic properties of CH-rk.

**Lemma 1.6** (Lemma 2.2.2(1)). Let  $D \subseteq \mathcal{M}$  be a definable set. Then  $\mathsf{CH}\text{-rk}(D) = 0$  if, and only if D is finite.

*Proof.* It suffices to show that if D is finite, then  $\mathsf{CH-rk}(D) \not\geq 1$ . Suppose not. Then, by definition, there is a definable set  $D_1$ , an infinite definable set  $D_2$ , a surjective map  $f: D_1 \to D_2$  (since  $\mathsf{CH-rk}(f^{-1}(d)) \geq 0$ , for all  $d \in D_2$ ) and a surjective finite-to-one map  $\pi: D_1 \to D$ . But, it is clear that  $D_1$  is also infinite, and hence D must be infinite, a contradiction.

**Lemma 1.7** (Lemma 2.2.2(2)). Let A, B be definable sets. If  $A \subseteq B$  then  $\mathsf{CH-rk}(A) \leq \mathsf{CH-rk}(B)$ .

*Proof.* Formally we prove, by induction on  $n \in \mathbb{N}$ , that, for definable sets A, B, if  $A \subseteq B$  and  $\operatorname{CH-rk}(A) \ge n$  then  $\operatorname{CH-rk}(B) \ge n$ . The base case (n=0) is trivial. For the inductive step, suppose that the result holds for definable sets of rank at least n and  $\operatorname{CH-rk}(A) \ge n+1$ . By definition, there exist definable sets  $D_1, D_2$  and definable functions  $\pi: D_1 \to A$ , and  $f: D_1 \to D_2$ , such that:  $D_2$  is infinite;  $\operatorname{CH-rk}(\pi^{-1}(d)) = 0$  for all  $d \in D_1$ ; and  $\operatorname{CH-rk}(f^{-1}(d)) \ge n$ , for all  $d \in D_2$ .

Without loss of generality, we may assume that  $D_1$  and B are disjoint (e.g. by appending a new fixed coordinate to all elements of  $D_1$ ). Let  $D'_1 = D_1 \sqcup (B \setminus A)$ , pick a point  $d_* \in D_2$  and define:

$$g: D_1' \to D_2$$
 
$$d \mapsto \begin{cases} f(d) & \text{if } d \in D_1 \\ d_{\star} & \text{otherwise.} \end{cases}$$

and

$$\pi': D_1' \to B$$
 
$$d \mapsto \begin{cases} \pi(d) & \text{if } d \in D_1 \\ d & \text{otherwise.} \end{cases}$$

Clearly, for each  $d \in D_2$  we have  $g^{-1}(d) \subseteq f^{-1}(d)$ . Since  $\mathsf{CH-rk}\left(f^{-1}(d)\right) \geq n$ , by inductive hypothesis,  $\mathsf{CH-rk}\left(g^{-1}(d)\right) \geq n$ . Thus, by definition, we have that  $\mathsf{CH-rk}(B) \geq n + 1$ , as required.

**Lemma 1.8** (Remark in the proof of Lemma 2.2.2). Let A and B be definable sets. Then:

$$\mathsf{CH-rk}(A \cup B) = \max\{\mathsf{CH-rk}(A), \mathsf{CH-rk}(B)\}.$$

*Proof.* It is clear form Lemma 1.7 that  $\max\{\mathsf{CH-rk}(A),\mathsf{CH-rk}(B)\} \leq \mathsf{CH-rk}(A \cup B)$ . For the inequality  $\mathsf{CH-rk}(A \cup B) \leq \max\{\mathsf{CH-rk}(A),\mathsf{CH-rk}(B)\}$ , we start by observing that it suffices to prove it when A and B are disjoint. Indeed, once we have proved it for two disjoint sets, we can immediately generalise it by induction to unions of n disjoint sets. Then, clearly:

$$\begin{aligned} \mathsf{CH-rk}(A \cup B) &\leq \max\{\mathsf{CH-rk}(A \setminus (A \cap B)), \mathsf{CH-rk}(B \setminus (A \cap B)), \mathsf{CH-rk}(A \cap B)\} \\ &\leq \max\{\mathsf{CH-rk}(A), \mathsf{CH-rk}(B)\}, \end{aligned}$$

where the last inequality follows immediately from Lemma 1.7.

Let  $n \in \mathbb{N}$  be arbitrary and suppose that A and B are disjoint sets and  $\mathsf{CH-rk}(A \cup B) \geq n+1$  (the case n=0 is trivial). The result will follow almost immediately from the definition and the pigeonhole principle. Explicitly, if  $\mathsf{CH-rk}(A \cup B) \geq n+1$ , then, there are definable sets  $D_1, D_2$  and surjective definable maps  $\pi: D_1 \to D$  and  $f: D_1 \to D_2$  such that  $D_2$  is infinite,  $\pi^{-1}(d)$  is finite, for all  $d \in D$  and

 $\mathsf{CH-rk}(f^{-1}(d)) \geq n$ , for all  $d \in D_2$ . We may partition  $D_2$  into two disjoint definable sets, as follows:

$$D_2 = \{ d \in D_2 : \pi(f^{-1}(d)) \in A \} \sqcup \{ d \in D_2 : \pi(f^{-1}(d)) \in B \}.$$

At most one of these sets is finite, by the pigeonhole principle, and thus at least one of CH-rk(A) or CH-rk(B) must be greater than n+1.

## 1.2. Ranks of elements.

**Definition 1.9.** Let  $a \in \mathcal{M}$  and  $B \subseteq \mathcal{M}$ . The Cherlin-Hrushovski rank of a over B, denoted CH-rk(a/B), is:

$$\mathsf{CH-rk}(a/B) := \min\{\mathsf{CH-rk}(D) : D \in \mathsf{tp}(a/B)\}.$$

Remark 1.10. Let  $\mathcal{M}$  be an  $\aleph_0$ -categorical structure  $a \in \mathcal{M}$  and  $B \subseteq \mathcal{M}$  a finite subset. Then, there is a smallest B-definable subset of  $\mathcal{M}$  containing a. [By  $\aleph_0$ -categoricity of  $\mathcal{M}$  and finiteness of B there are only finitely many B-definable subsets, up to equivalence. Finite lattices have minimal elements.]

Thus, the following is well-defined:

**Definition 1.11.** Let  $\mathcal{M}$  be an  $\aleph_0$ -categorical structure  $a \in \mathcal{M}$  and  $B \subseteq \mathcal{M}$  a finite subset. The *locus of a over B* is the smallest *B*-definable subset containing a.

In particular:

Remark 1.12. Let  $\mathcal{M}$  be  $\aleph_0$ -categorical,  $a \in \mathcal{M}$  and  $B \subseteq \mathcal{M}$  a finite set. Then  $\mathsf{CH-rk}(a/B)$  is precisely the rank of the locus of a over B.

We may now translate the results of the previous subsection, in the context of ranks of elements:

**Lemma 1.13** (Lemma 2.2.2(1') and (2")). Let  $a \in \mathcal{M}$  and  $B, B_1, B_2 \subseteq \mathcal{M}$ . Then:

- (1)  $\mathsf{CH}\text{-rk}(a/B) = 0$  if, and only if  $a \in \mathsf{acl}(B)$ .
- (2) If  $B_1 \subseteq B_2$  then  $\mathsf{CH-rk}(a/B_2) \le \mathsf{CH-rk}(a/B_1)$ . [In the book, they say that this amounts to  $D_1 \subseteq D_2 \implies \mathsf{CH-rk}(D_1) \le \mathsf{CH-rk}(D_2)$ . I think it's just from the definition. Say  $\mathsf{CH-rk}(a/B_1) = \mathsf{CH-rk}(D)$  for some  $B_1$ -definable set D. Then D is also  $B_2$ -definable, so  $\mathsf{CH-rk}(a/B_2) \le \mathsf{CH-rk}(D) = \mathsf{CH-rk}(a/B_1)$ .]

**Lemma 1.14** (Extension Property, Lemma 2.2.2(2')). Let D be a non-empty B-definable set. Then, there is a complete type over B containing D and having the same rank.

*Proof.* The lemma follows easily from the fact that:

$$\mathsf{CH-rk}(A \cup B) = \max\{\mathsf{CH-rk}(A), \mathsf{CH-rk}(B)\}.$$

Let's actually go through the details:

Let D be B-definable. To fix notation, say D is given by  $\phi(x, b_0)$  for some  $\mathcal{L}$ -formula  $\phi(x, y)$  and some  $b_0 \in B$ . Now, let  $\pi(x)$  be the following partial type:

$$\pi(x) = \{ \psi(x, b) \in \mathcal{L}(B) : \mathsf{CH-rk}(\phi(x, b_0) \land \neg \psi(x, b)) < \mathsf{CH-rk}(\phi(x, b_0)) \}.$$

Clearly  $\phi(x, b_0) \in \pi(x)$ . We claim that  $\pi(x)$  is finitely consistent. Indeed, suppose that  $\psi_1(x,b), \psi_2(x,b) \in \pi(x)$ , and assume toward a contradiction that  $\psi_1(x,b) \wedge \psi_2(x,b)$  is empty. By definition, then,  $\mathsf{CH-rk}(\psi_1(x,b) \wedge \psi_2(x,b)) = 0$ , but:

$$\mathsf{CH-rk}(\phi(x,b_0)) = \max \left\{ \begin{aligned} \mathsf{CH-rk}(\phi(x,b_0) \wedge \neg \psi_1(x,b)), \\ \mathsf{CH-rk}(\phi(x,b_0) \wedge \neg \psi_2(x,b)), \\ \mathsf{CH-rk}(\phi(x,b_0) \wedge \psi_1(x,b) \wedge \psi_2(x,b)) \end{aligned} \right\},$$

and by assumption, the RHS above is less than CH-rk( $\phi(x, b_0)$ ), a contradiction. So, let p(x) extend  $\pi(x)$  to a complete type over B.

To finish the proof, suppose that for some formula  $\psi(x,b) \in p(x)$  we have that  $\mathsf{CH-rk}(\psi(x,b)) < \mathsf{CH-rk}(\phi(x,b_0))$ . Then  $\mathsf{CH-rk}(\psi(x,b) \land \phi(x,b_0)) < \mathsf{CH-rk}(\phi(x,b_0))$ , so  $\neg \psi(x,b) \in p(x)$ , a contradiction.

1.3. **The not so very basics.** The next lemma is a useful tool for computing ranks. We will repeatedly use it in the proof of 'additivity' (Proposition 1.17):

**Lemma 1.15** (Lemma 2.2.3). Let  $\mathcal{M}$  be  $\aleph_0$ -categorical. Then, the following are equivalent:

- (1) CH-rk $(a/b) \ge n + 1$ .
- (2) There are a', c with  $a' \in \operatorname{acl}(abc) \setminus \operatorname{acl}(bc)$ , and  $\operatorname{CH-rk}(a/a'bc) \geq n$

*Remark* 1.16. Observe, for instance, that by the lemma above, the following are equivalent:

- (1) CH-rk(a) = 1.
- (2) There are a', c with  $a' \in \operatorname{acl}(ac) \setminus \operatorname{acl}(c)$  such that  $a \in D$ , for some a'c-definable D, and for all a', c with  $a' \in \operatorname{acl}(ac) \setminus \operatorname{acl}(c)$  we have  $a \in \operatorname{acl}(a'c)$ .

*Proof.* For the entirety of the proof, let D be the locus of a over b. [The point being that if  $a, \alpha \in D$  then  $a \equiv_b \alpha$ . Indeed, suppose toward a contradiction that  $\alpha \in D$  and there is some b-definable set D' such that  $\alpha \in D'$ , but  $a \notin D'$ . Then, the locus of a should have been  $D \cap (\neg D')$ . Thus by the locus of a over b we essentially mean the formula isolating  $\operatorname{tp}(a/b)$ .]

(1)  $\Rightarrow$  (2) . By definition, CH-rk(a/B) = CH-rk(D). Let  $D_1, D_2, \pi$  and f witness that CH-rk(D)  $\geq n+1$ . Let  $c \in \mathcal{M}$  be a finite tuple such that  $D_1, D_2, \pi$  and f are all c-definable. By  $\aleph_0$ -categoricity, it follows that  $D_2 \setminus \operatorname{acl}(bc) \neq \emptyset$ , since  $D_2$  is infinite.

[The algebraic closure of finite sets in  $\aleph_0$ -categorical structures is finte.]

So, we may pick some  $a' \in D_2 \setminus \operatorname{acl}(bc)$ . We know that  $\operatorname{CH-rk}(f^{-1}(a')) \geq n$ , by assumption, and since  $f^{-1}(a')$  is a'bc-definable, by the Extension Property (Lemma 1.14) and  $\omega$ -saturation, we can find some  $a_1 \in f^{-1}(a')$  such that  $\operatorname{CH-rk}(a_1/a'bc) = \operatorname{CH-rk}(f^{-1}(a'))$ . Now, to finish the proof, let  $a_0 = \pi(a_1)$ . It is easy to see that  $a' \in \operatorname{acl}(a_0bc)$ , since the fibre of  $\pi$  above  $a_0$  is an algebraic set.

Claim 1.16.1. CH-rk $(a_0/a'bc) \ge n$ .

Proof of Claim. We know that  $\mathsf{CH-rk}(a_1/a'bc) \geq n$ . Let  $A_1$  be the locus of  $a_1$  over a'bc and  $A_0$  the locus of  $a_2$  over a'bc. Clearly  $\pi$  is a finite-to-one map from  $A_1$  to  $A_0$ , so the claim follows.

Now, to finish the proof, since  $a_0 \in D$ , we have that  $\operatorname{tp}(a_0/b) = \operatorname{tp}(a/b)$ . By strong  $\omega$ -homogeneity there is an automorphism  $\sigma$  taking  $a_0$  to a. But then we are done, after replacing a' by  $\sigma(a')$  and c by  $\sigma(c)$ .

 $(2) \Rightarrow (1)$  Suppose that there are a' and c as in (2). Let  $D_1$  be the set:

$$\{(x,y): \operatorname{tp}(xy/bc) = \operatorname{tp}(aa'/bc)\}$$

[This is definable, because  $\mathcal{M}$  is  $\aleph_0$ -categorical, and types over finite sets are isolated.]

Let f be the projection of  $D_1$  onto the second coordinate and take  $D_2$  to be  $f(D_1)$ .

Claim 1.16.2. Let  $\pi$  be the projection of  $D_1$  onto the first coordinate. Then  $\pi(D_1) = D$ .

Proof of Claim. Recall that D is the locus of a over b. Suppose first that  $\alpha \in D$ . In particular, this means that  $\operatorname{tp}(a/b) = \operatorname{tp}(\alpha/b)$ . By strong  $\omega$ -homogeneity, there is an automorphism  $\sigma \in \operatorname{Aut}(\mathcal{M}/bc)$  taking a to  $\alpha$ . Taking  $\alpha' = \sigma(a')$  gives us  $\operatorname{tp}(\alpha\alpha') = \operatorname{tp}(aa'/bc)$ , so  $\alpha \in \pi(D_1)$ . Conversely, suppose that  $(\alpha, \alpha') \in D_1$ , then, we claim that  $\alpha \in D$ . Since  $\alpha \in D_1$ , we have  $a \equiv_b \alpha$ , and thus  $\alpha \in D_1$ .

[The point of the assumption  $a' \in \operatorname{acl}(abc) \setminus \operatorname{acl}(bc)$  is that the set  $D_2$  is infinite and that the fibre in  $D_1$  above any  $d \in D$  is finite.]

To finish the proof, it remains to show that for all  $d \in D_2$  we have  $\mathsf{CH-rk}\left(f^{-1}(d)\right) \geq n$ .

Claim 1.16.3. CH-rk 
$$(f^{-1}(a')) \ge n$$

*Proof of Claim.* By assumption, we have  $\mathsf{CH-rk}(a/a'bc) \geq n$ . So, by definition, the locus of a over a'bc has rank at least n.

Finally, observe that all fibres of f are conjugates under automorphisms, so we are done.

**Proposition 1.17** (Lemma 2.2.4). Let  $\mathcal{M}$  be  $\aleph_0$ -categorical. If CH-rk(a/bc) and CH-rk(b/c) are finite then:

$$\mathsf{CH}\text{-rk}(ab/c) = \mathsf{CH}\text{-rk}(a/bc) + \mathsf{CH}\text{-rk}(b/c).$$

*Proof.* The proof is by induction on the  $n = \mathsf{CH-rk}(a/bc) + \mathsf{CH-rk}(b/c)$ . The base case is trivial. Indeed, if  $\mathsf{CH-rk}(a/bc) + \mathsf{CH-rk}(b/c) = 0$ , then  $\mathsf{CH-rk}(a/bc) = 0$ , so  $a \in \mathsf{acl}(bc)$  and  $\mathsf{CH-rk}(b/c) = 0$ , so  $b \in \mathsf{acl}(c)$ . Combining the two conclusions, we obtain that  $a \in \mathsf{acl}(b)$ , so  $a, b \in \mathsf{acl}(c)$  and thus  $\mathsf{CH-rk}(ab/c) = 0$ .

Now, for the inductive step. First, we shall show the inequality  $\mathsf{CH-rk}(ab/c) \leq n$ . To this end, pick some d such that  $\mathsf{acl}(abcd) \setminus \mathsf{acl}(abc) \neq \emptyset$  and let  $e \in \mathsf{acl}(abcd) \setminus \mathsf{acl}(cd)$ . By Lemma 1.15, it suffices to show that  $\mathsf{CH-rk}(ab/cde) < n$ . Explicitly, our inductive hypothesis will give us the following:

(IH) If 
$$\mathsf{CH}\text{-rk}(a/bcde) + \mathsf{CH}\text{-rk}(b/cde) < n$$
 then:

$$\mathsf{CH-rk}(ab/cde) = \mathsf{CH-rk}(a/bcde) + \mathsf{CH-rk}(b/cde).$$

At this point, there are two cases to consider:

• <u>Case 1</u>: If  $e \notin \operatorname{acl}(bcd)$ , then  $e \in \operatorname{acl}(abcd) \setminus \operatorname{acl}(bcd)$ . In this case, we claim that CH-rk(a/bcde) < CH-rk(a/bc) ≤ n. Indeed, assume toward a contradiction that CH-rk(a/bcde) = CH-rk(a/bc). Now, applying Lemma 1.15, we see that we have found d, e with  $e \in \operatorname{acl}(abcd) \setminus \operatorname{acl}(bcd)$  and CH-rk(a/bcde) ≥ n. Thus CH-rk(a/bc) ≥ CH-rk(a/bc) + 1, which is, of course, a contradiction. Thus CH-rk(a/bcde) < CH-rk(a/bc). We always have that CH-rk(b/cde) ≤ CH-rk(b/cde), so

$$\mathsf{CH-rk}(a/bcde) + \mathsf{CH-rk}(b/cde) < \mathsf{CH-rk}(a/bc) + \mathsf{CH-rk}(b/c) = n,$$

and, by (IH) we obtain:

$$\mathsf{CH-rk}(ab/cde) = \mathsf{CH-rk}(a/bcde) + \mathsf{CH-rk}(b/cde) < n,$$

as required.

• <u>Case 2</u>: If  $e \in \operatorname{acl}(bcd)$ , then  $e \in \operatorname{acl}(bcd) \setminus \operatorname{acl}(cd)$ . In this case, we claim that  $\operatorname{CH-rk}(b/cde) < \operatorname{CH-rk}(b/cd)$ . The argument is similar. If not, then we have that  $\operatorname{CH-rk}(b/cde) = \operatorname{CH-rk}(b/cd)$ , and again, we have found d, e such that  $e \in \operatorname{acl}(bcd) \setminus \operatorname{acl}(cd)$ . So again by Lemma 1.15 we can conclude that  $\operatorname{CH-rk}(b/cd) \ge \operatorname{CH-rk}(b/cd) + 1$ , which is a contradiction. We again conclude by (IH) exactly as in the previous case.

It remains to show the inequality  $\mathsf{CH-rk}(ab/c) \geq n$ . Observe, first, that if  $\mathsf{CH-rk}(b/c) = 0$  then:

$$\mathsf{CH-rk}(ab/c) \ge \mathsf{CH-rk}(a/c) \ge \mathsf{CH-rk}(a/bc) = n,$$

and we are done.

[The first inequality holds because the locus of ab over c certainly contains the locus of a over c.]

Thus, we may assume that  $\mathsf{CH}\text{-rk}(b/c) > 0$ . By Lemma 1.15 it suffices to find some b', d such that  $binacl(bcd) \setminus \mathsf{acl}(cd)$  and  $\mathsf{CH}\text{-rk}(ab/b'cd) \ge n-1$ , for then we have  $\mathsf{CH}\text{-rk}(ab/c) \ge n$ .

To this end, pick b',d such that  $b'\in \mathsf{acl}(bcd)\setminus \mathsf{acl}(cd)$  and  $\mathsf{CH-rk}(b/b'cd)=\mathsf{CH-rk}(b/c)-1$ .

[Such elements exist precisely by Lemma 1.15.]

To finish the proof, by Lemma 1.15 it suffices to show that  $\mathsf{CH-rk}(ab/b'cd) \geq n-1$ . By the Extension Property (Lemma 1.14) we may assume that  $\mathsf{CH-rk}(a/bb'cd) = \mathsf{CH-rk}(a/bc)$ .

[The point is that the locus D of a over bc is, of course, definable over  $\{b,b',c,d\}$ , thus, by  $\omega$ -saturation we can find an element  $a_{\star} \in D$ , such that  $\mathsf{CH-rk}(a_{\star}/bb'cd) = \mathsf{CH-rk}(D) = \mathsf{CH-rk}(a/bc)$ . By strong  $\omega$ -homogeneity, we can find an automorphism  $\sigma$  which fixes b and c sending  $a_{\star}$  to a. Replacing d and b', by their images under  $\sigma$  the claimed equality holds. Of course, since  $\sigma$  fixes b and c, the previous conditions on b' and d are still true.]

Thus:  $\mathsf{CH-rk}(a/bb'cd) + \mathsf{CH-rk}(b/b'cd) = n-1,$ 

so, in particular, by (IH) we have that:

$$\mathsf{CH-rk}(ab/b'cd) = \mathsf{CH-rk}(a/bb'cd) + \mathsf{CH-rk}(b/b'cd) < n,$$

and the result follows.

**Corollary 1.18** (Corollary 2.2.5). If CH-rk(D) = 1, then acl defines a pregeometry on D.

*Proof.* For notational convenience, throughout this proof we will write  $\operatorname{acl}(-)$  to mean  $\operatorname{acl}(-) \cap D$ . We must show that for all singletons  $a, b \in D$  and finite tuples c from D we have that:

$$a \in \operatorname{acl}(bc) \setminus \operatorname{acl}(c) \implies b \in \operatorname{acl}(ac).$$

Let a, b and c satisfy the antecedent of the implication above. We know that

$$\mathsf{CH-rk}(ab/c) = \mathsf{CH-rk}(b/ac) + \mathsf{CH-rk}(a/c)$$

and we wish to show that  $\mathsf{CH-rk}(b/ac) = 0$ . Thus, it suffices to show that  $\mathsf{CH-rk}(ab/c) = \mathsf{CH-rk}(a/c)$ . Since  $a \not\in \mathsf{acl}(c)$  we must have that  $\mathsf{CH-rk}(a/c) \ge 1$ , and since  $\mathsf{CH-rk}(a) = 1$ , we must have that  $\mathsf{CH-rk}(a/c) = 1$ . Similarly,  $\mathsf{CH-rk}(ab/c) = 1$ , and we are done.

**Definition 1.19.** We say that a and b are independent over C, written  $a \downarrow_C^{\mathsf{ch}} b$  if:

$$\mathsf{CH}\text{-rk}(a/bC) = \mathsf{CH}\text{-rk}(a/C) + \mathsf{CH}\text{-rk}(b/C).$$

Remark 1.20. Equivalently,  $a \downarrow_C^{\mathsf{ch}} b$  if, and only if,  $\mathsf{CH-rk}(a/bC) = \mathsf{CH-rk}(a/C)$ .

We collect here some of the properties of  $\bigcup^{ch}$ :

**Lemma 1.21** (Lemma 2.2.7). The independence relation  $\bigcup$ <sup>ch</sup> satisfies the following properties:

(1) Symmetry: 
$$a \downarrow_E^{\mathsf{ch}} b \iff b \downarrow_E^{\mathsf{ch}} a$$
.

- (2) Monotonicity:  $a \downarrow_E^{\mathsf{ch}} bc \implies a \downarrow_E^{\mathsf{ch}} b$ .
- (3) Base Monotonicity: For  $E \subseteq F \subseteq G$  we have  $a \downarrow_E^{\mathsf{ch}} F \implies a \downarrow_B^{\mathsf{ch}} F$ .
- (4) Transitivity:
- (5) If  $a \in \operatorname{acl}(bC)$  then  $a \downarrow_C^{\operatorname{ch}} b \iff a \in \operatorname{acl}(C)$ .

Remark 1.22. The conjunction of properties (2) - (4) above are equivalent to the following equivalence:

$$a \stackrel{\mathsf{ch}}{\underset{E}{\bigcup}} bc \iff a \stackrel{\mathsf{ch}}{\underset{Ec}{\bigcup}} b \text{ and } a \stackrel{\mathsf{ch}}{\underset{E}{\bigcup}} c$$

Proof.

(1) Symmetry: Suppose that  $a \downarrow_E^{\mathsf{ch}} b$ . Then, by definition, we have that:

$$\mathsf{CH} ext{-rk}(a/bC) = \mathsf{CH} ext{-rk}(a/C)$$

By Proposition 1.17 we know that:

$$\mathsf{CH-rk}(a/bC) + \mathsf{CH-rk}(b/C) = \mathsf{CH-rk}(b/aC) + \mathsf{CH-rk}(a/C).$$

and thus  $\mathsf{CH}\text{-rk}(b/C) = \mathsf{CH}\text{-rk}(b/aC)$ .

(2) Monotonicity, Base Monotonicity, and Transitivity: We show the equivalent equivalence. Suppose first that  $a \, \, \bigcup_E^{\mathsf{ch}} bc$ . By definition, we have that:

$$\mathsf{CH}\text{-rk}(a/bcE) = \mathsf{CH}\text{-rk}(a/E)$$

We must show that  $a \downarrow_E^{\mathsf{ch}} cb$ , i.e. that  $\mathsf{CH}\text{-rk}(a/bcE) = \mathsf{CH}\text{-rk}(a/cE)$  and  $a \downarrow_E^{\mathsf{ch}} c$ , i.e. that  $\mathsf{CH}\text{-rk}(a/cE) = \mathsf{CH}\text{-rk}(a/E)$ . Both of these follow immediately, since:

$$\mathsf{CH}\text{-}\mathsf{rk}(a/bcE) \leq \mathsf{CH}\text{-}\mathsf{rk}(a/cE) \leq \mathsf{CH}\text{-}\mathsf{rk}(a/E),$$

and

$$\mathsf{CH}\text{-}\mathsf{rk}(a/E) \geq \mathsf{CH}\text{-}\mathsf{rk}(a/cE) \geq \mathsf{CH}\text{-}\mathsf{rk}(a/bcE).$$

Conversely, by symmetry we have that  $\mathsf{CH-rk}(b/acE) = \mathsf{CH-rk}(b/cE)$  and  $\mathsf{CH-rk}(c/aE) = \mathsf{CH-rk}(c/E)$ . Then:

$$\begin{split} \mathsf{CH-rk}(abc/E) &= \mathsf{CH-rk}(b/acE) + \mathsf{CH-rk}(c/aE) + \mathsf{CH-rk}(a/E) \\ &= \mathsf{CH-rk}(a/E) + \mathsf{CH-rk}(b/E) + \mathsf{CH-rk}(c/E), \end{split}$$

and we are done.

(3) Suppose that  $a\in \operatorname{acl}(bC)$ , so  $\operatorname{CH-rk}(a/bC)=0$ . First, we show that if  $a\downarrow^{\operatorname{ch}}_C b$  then  $a\in\operatorname{acl}(C)$ . Indeed, if  $\operatorname{CH-rk}(a/bC)=\operatorname{CH-rk}(a/C)$ , it follows that  $\operatorname{CH-rk}(a/C)=0$ , so  $a\in\operatorname{acl}(C)$ . Conversely, suppose that  $a\in\operatorname{acl}(C)$ . Then  $\operatorname{CH-rk}(a/C)=0$ , so  $\operatorname{CH-rk}(a/C)=0=\operatorname{CH-rk}(a/bC)$  and thus  $a\downarrow^{\operatorname{ch}}_C b$ .

# 1.4. ... And geometries. Brief reminder (from Paolo's talk):

**Definition 1.23** ((Weak) Linear Geometries). A weak linear geometry is one of the following six (types) of structures.

- (1) A Degenerate Space
- (2) A Pure Vector Space
- (3) A Polar Space
- (4) An Inner Product Space
- (5) An Orthogonal Space
- (6) A Quadratic Space

A linear geometry is a weak linear geometry  $\mathcal{M}$  expanded by a set of algebraic constants in  $\mathcal{M}^{eq}$ , i.e. expanded by adding to the language a subset of  $\operatorname{acl}^{eq}(\emptyset)$ .

These are the building blocks of the following kinds of geometries:

#### Definition 1.24.

- (1) An unoriented weak linear geometry a weak linear geometry of type (1)-(5) or a reduct of a QUADRATIC SPACE in which we have forgotten the Witt defect function.
- (2) A basic linear geometry is a linear geometry with the elements of K and L named by constants, and in the case of a Polar Space, the two vector spaces V and W named by unary predicates.
- (3) A projective geometry is a structure obtained by a linear geometry by factoring out the equivalence relation  $\operatorname{acl}(x) = \operatorname{acl}(y)$ .
- (4) An affine geometry is a pair (J, A), consisting of a linear geometry J with underlying vector space V (one of the two vector spaces in the Polar Space case), and a definable subset A on which V acts definably and regularly, where J carries its given structure and A.

**Lemma 1.25** (Lemma 2.2.10). The linear, affine, and projective geometries are all of pre-rank 1.

*Proof.* "Do them one-by-one, using QE, if you don't trust their proof."

**Definition 1.26.** A structure  $\mathcal{M}$  is *Lie coordinatised* if it admits a tree structure < of finite height (where < is an invariant partial order) with a unique 0-definable root such that:

- (1) Coordinatisation: For all  $a \in M$  above the root, one of the following holds: (A) a is algebraic over its <-predecessor. OR
  - (B) There is b < a and a b-definable projective geometry  $J_b$ , fully embedded in  $\mathcal{M}$  such that either:
    - (i)  $a \in J_b$ , or
    - (ii) There is some  $c \in M$  such that b < c < a and a c-definable affine or quadratic geometry  $(J_c, A_c)$  such that  $a \in A_c$  and the projectivisation of  $J_c$  is  $J_b$ .
- (2) Orientation: If  $a, b \in M$  have the same type over  $\emptyset$  and are associated with coorinatising quadratic geometries  $J_a$ ,  $J_b$ , then any definable map between them which preserves everything other than w also preserves w.

We say that  $\mathcal{M}$  is  $Lie\ coordinatisable$  if it is interpretable with a structure  $\mathcal{N}$  with finitely many 1-types over  $\emptyset$  which is Lie coordinatised.

Corollary 1.27 (Corollary 2.2.11). If  $\mathcal{M}$  is Lie coordinatisable then  $\mathcal{M}$  has finite rank, at most the height of the coordination tree.

*Proof.* [Idea:Every time we meet a geometry the rank goes up by 1. The algebraic steps don't increase the rank.]  $\Box$ 

**Corollary 1.28.** Let J be a linear, projective, or affine geometry and a, b are finite tuples from J. If  $acl(a) \cap acl(b) = C$  then  $a \cup_C^{ch} b$ . So J is one-based.

2. More Cherlin-Hrushovski ranks

## Definition 2.1. Induction

- $\mathsf{CH}\text{-rk}_0(D) = \begin{cases} 0 & \text{if } D \text{ is finite} \\ \infty & \text{otherwise} \end{cases}$
- Let  $\alpha \in \mathbf{Ord}$ , and assume that  $\mathsf{CH-rk}_{\beta}(D)$  has been defined for all  $\beta < \alpha$ . Then:
  - (1)  $\mathsf{CH}\text{-rk}_{\alpha}(D) \geq 0$  if D is non-empty.

- (2)  $\mathsf{CH}\text{-rk}_{\alpha}(D) > 0$  if  $\mathsf{CH}\text{-rk}_{\beta}(D) = \infty$  for all  $\beta < \alpha$ .
- (3) For  $n \in \mathbb{N}$ ,  $\mathsf{CH-rk}_{\alpha}(D) \geq n+1$  if there exist definable sets  $D_1, D_2$  and definable functions  $\pi: D_1 \to D$  and  $f: D_1 \to D_2$  such that:
  - (a) For all  $d \in D$  we have  $\mathsf{CH}\text{-rk}_{\alpha}\left(\pi^{-1}(d)\right) = 0$ .
  - (b) CH-rk<sub>\alpha</sub>(D<sub>2</sub>) > 0.
  - (c) For all  $d \in D_2$  we have  $\mathsf{CH}\text{-rk}_{\alpha}(f^{-1}(d)) \geq n$ .

TODO: Look at their pseudofinite example...

# 3. An alternative definition

In this section, we follow [Sim22, Section 2.4] rather closely. Diversions from that source will be indicated in red.

By a uniformly definable family  $(X_t:t\in E)$  we mean that E is a definable set and there is a formula  $\phi(x,y)$  such that, for each  $t\in E$ , the set  $X_t$  is precisely  $\phi(x,t)$ . Such a family is weakly k-inconsistent, for  $k\in \mathbb{N}$ , if whenever  $X_{t_1},\ldots,X_{t_k}$  are pairwise distinct members of the family then  $\bigcap_{i\leq k}X_{t_i}=\emptyset$ .

[A uniformly definable family  $(X_t : t \in E)$  is called k-inconsistent, for  $k \in \mathbb{N}$  if whenever  $t_1, \ldots, t_k$  are pairwise distinct then  $\bigcap_{i \leq k} X_{t_i} = \emptyset$ .]

**Definition 3.1** ([Sim22, Definition 2.2]). Let D be a definable subset of  $\mathcal{M}$  (not  $\mathcal{M}^{eq}$ ). We define p-rk(D)  $\geq n$ , by induction on  $n \in \mathbb{N}$ :

- $b\text{-rk}(D) \ge 0$  if D is consistent.
- b-rk(D) > 0 if D is infinite.
- $p\text{-rk}(D) \ge n+1$  if there is a uniformly definable weakly k-inconsistent family  $(X_t: t \in E)$  of subsets of D, containing infinitely many pairwise distinct sets, such that  $p\text{-rk}(X_t) \ge n$  for all  $t \in E$ .

**Theorem 3.2.** For any structure  $\mathcal{M}$  and any definable set  $D \subseteq M^k$  not  $\mathcal{M}^{eq}$  we have that:

$$p$$
-rk $(D) = CH$ -rk $(D)$ .

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