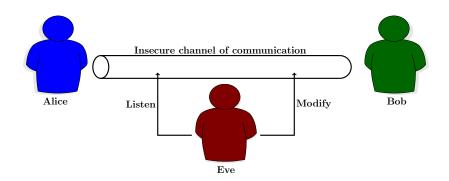
## Symmetric Key Cryptography

University of Birmingham

#### Outline of This Lecture

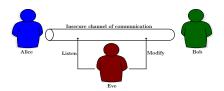
- The framework and the model
- Classical Cryptography and One Time Pad
- Stream Ciphers and Block Ciphers
- Modes of Operations

## Setup



Alice and Bob needs to communicate "securely"

#### Model



- ► Alice, Bob, and Eve are Algorithms.
- Questions: What is the computation power of Eve?
- Question: What kind of tampering could Eve do?

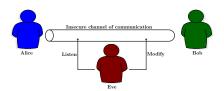
## Modelling Eve

- Assume computation power: Million Teraflops  $\approx 2^{60}$  computations per second.
- ▶ 1 common year=  $3600 \times 24 \times 365$  seconds = 3153600 seconds  $\approx 2^{22}$  seconds.
- ▶ Total one year of computation  $\approx 2^{82}$  computations

## Modelling Eve

- Assume computation power: Million Teraflops  $\approx 2^{60}$  computations per second.
- ▶ 1 common year=  $3600 \times 24 \times 365$  seconds = 3153600 seconds  $\approx 2^{22}$  seconds.
- ▶ Total one year of computation  $\approx 2^{82}$  computations
- ▶ Ballpark estimate of Eve's power  $\approx 2^{100}$  computations: accepted standard for non-classified data.

#### Model



- Alice, Bob, and Eve are Algorithms.
- Questions: What is the computation power of Eve? 2<sup>100</sup> computations.
- Question: What kind of tampering could Eve do?

#### Modelling Eve: Channel Modification

Could Eve erase everything? If yes, no communication could be done.

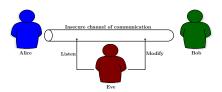
### Modelling Eve: Channel Modification

- Could Eve erase everything? If yes, no communication could be done.
- Could Eve modify small fraction of data (say 1/4)?
  - We use error correcting codes to correct errors. (beyond the scope)

### Modelling Eve: Channel Modification

- Could Eve erase everything? If yes, no communication could be done.
- Could Eve modify small fraction of data (say 1/4)?
  - We use error correcting codes to correct errors. (beyond the scope)
- Cryptographic modeling: Eve could modify any fraction, we care about error detection.

#### Model



- Alice, Bob, and Eve are Algorithms.
- Computation power of Eve? 2<sup>100</sup> for non-classified data
- Eve could modify any part: Alice and Bob need error detection.

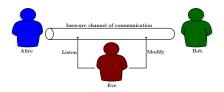
#### Question

What does Eve know?

## Kerckhoffs's principle: Second Rule

System should not require secrecy. Algorithms of Alice and Bob are public information.

#### Model

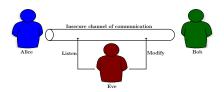


- Alice, Bob, and Eve are Algorithms.
- Computation power of Eve? 2<sup>100</sup> for non-classified data
- Eve could modify any part: Alice and Bob need error detection.

#### Question

What does Eve know? Alice and Bob's algorithms.

#### Model



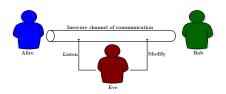
- Alice, Bob, and Eve are Algorithms.
- Computation power of Eve? 2<sup>100</sup> for non-classified data
- Eve could modify any part: Alice and Bob need error detection.

#### Question

What does Eve know? Alice and Bob's algorithms.

No Secrecy Yet: Eve could run Bob's algorithm on the communication!

#### Model: the KEY to secure communication



- Alice, Bob, and Eve are Algorithms.
- Computation power of Eve? 2<sup>100</sup> for non-classified data
- Eve could modify any part: Alice and Bob need error detection.

#### Question

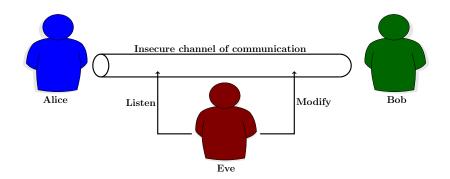
What does Eve know? Alice and Bob's algorithms.

#### secret key

A secret information known to Bob; unknown to Eve.



## Symmetric Key Cryptography

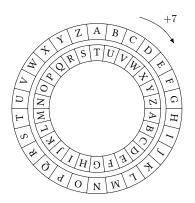


#### Setup in symmetric key cryptography

Alice and Bob <u>both</u> know the secret key. Eve does not know the secret key.

## Symmetric Key Cryptography: Historic account

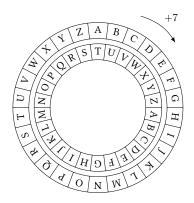




► SECURITY → ZLJBYPAF

## Symmetric Key Cryptography: Historic account

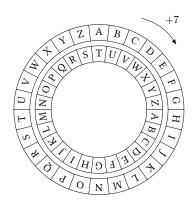




- SECURITY → ZLJBYPAF
- Decryption requires going back 7 characters

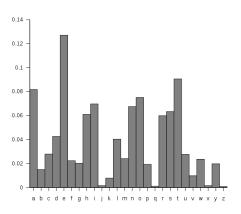
## Symmetric Key Cryptography: Historic account





- SECURITY → ZLJBYPAF
- Decryption requires going back 7 characters
- ► Could be generalised to any number between 1 to 26. The number is going to be the key.

## Ceaser Cipher: Cryptanalysis



- Broken using frequency analysis: 'e' is the most frequent character, followed by 't' then 'a'
- for sufficiently long ciphertext, shift by fixed length maintains the relative frequency.
  - For +7 shift, 'I' is most frequent, followed by 'a' and 'h'



#### One Time Pad

- For each character a random shift is chosen.
- the sequence of shift (number) is the key

Plain text: THIS IS SECRET OTP-Key: XVHE UW NOPGDZ

Ciphertext: Q C P W C O F S R X H S In groups: QCPWC OFSRX HS

Perfect Secrecy when the key is random.

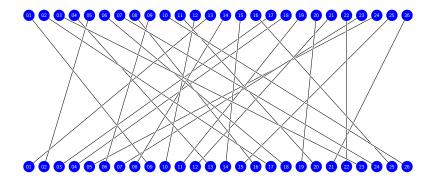
#### Issues with one time pad

- Key has to be as long as the text.
- Can not repeat key.

## **Encrypting with Smaller Keys**

- ▶ **Block Cipher** Encrypt *n*-bit block via a randomly chosen permutation.
- ► Stream Cipher Generate a random looking bit-stream from a smaller key and xor the message with the stream

## **Encryptions are Permutations**



#### **Block Ciphers**

▶ Message Space  $\mathcal{M}$ , typically  $\{0,1\}^{128}$ 

Number of possible Permutations factorial  $\left(2^{128}\right)$ 

#### **Block Ciphers**

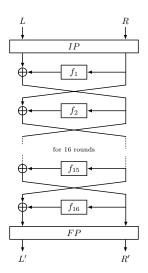
- ▶ Message Space  $\mathcal{M}$ , typically  $\{0,1\}^{128}$
- ▶ Key Space  $\mathcal{K}$ , say  $\{0,1\}^{128}$

#### Number of possible Permutations

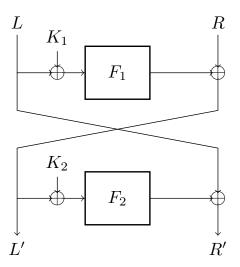
factorial  $(2^{128})$ 

A block cipher over the keyspace  $\{0,1\}^{128}$  is a family of  $2^{128}$  many permutations

## **Block Cipher Designs:DES**



## Design Principle: Feistel Network



## Block Cipher Designs:AES

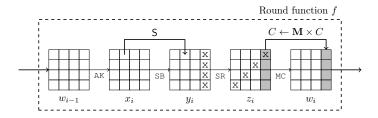
 $\mathcal{M} = \{0,1\}^{128}$ . Each domain element is 16-bytes long.

| 0 | 4 | 8  | 12 |
|---|---|----|----|
| 1 | 5 | 9  | 13 |
| 2 | 6 | 10 | 14 |
| 3 | 7 | 11 | 15 |

## Block Cipher Designs:AES

 $\mathcal{M} = \{0,1\}^{128}$ . Each domain element is 16-bytes long.

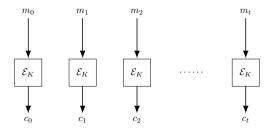
| 0 | 4 | 8  | 12 |
|---|---|----|----|
| 1 | 5 | 9  | 13 |
| 2 | 6 | 10 | 14 |
| 3 | 7 | 11 | 15 |



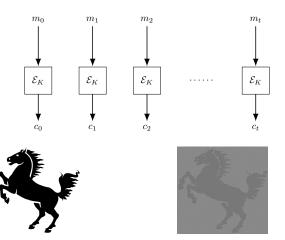
### **Modes of Operations**

- Block ciphers encrypt fixed length strings: AES encrypts 128 bits.
- ► How to encrypt large messages? Modes of operations

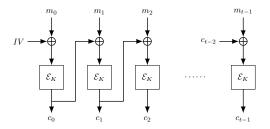
# Electronic Code Book: Parallel applications of block cipher



# Electronic Code Book: Parallel applications of block cipher

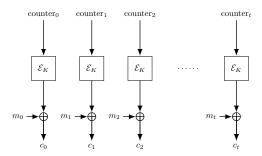


#### **CBC** mode



## Caution IV needs to be random!

#### Counter mode



#### Caution

Ciphertexts are malleable.