

//_

Question 1

$$A^0 = \begin{bmatrix} 0 & 2 & \infty & 1 & 8 \\ 6 & 0 & 3 & 2 & \infty \\ \infty & \infty & 0 & 4 & \infty \\ \infty & \infty & 2 & 0 & 3 \\ 3 & \infty & \infty & \infty & 0 \end{bmatrix}$$

Using formula:

$$A^k(i, j) = \min(A^{k-1}(i, j), A^{k-1}(i, k) + A^{k-1}(k, j))$$

$$A^1 = \begin{bmatrix} 0 & 2 & \infty & 1 & 8 \\ 6 & 0 & 3 & 2 & 14 \\ \infty & \infty & 0 & 4 & \infty \\ \infty & \infty & 2 & 0 & 3 \\ 3 & 5 & \infty & 4 & 0 \end{bmatrix}$$

$$A^1(2, 3) = \min(3, 6 + \infty) \\ = 3$$

$$A^1(3, 5) = \min(\infty, \infty + 8) \\ = \infty$$

$$A^1(2, 4) = \min(2, 6 + 1) \\ = 2$$

$$A^1(4, 2) = \infty$$

$$A^1(4, 3) = 2$$

$$A^1(4, 5) = 3$$

$$A^1(2, 5) = \min(\infty, 6 + 8) \\ = 14$$

$$A^1(5, 2) = 5$$

$$A^1(5, 3) = \infty$$

$$A^1(3, 2) = \min(\infty, \infty + \infty) \\ = \infty$$

$$A^1(5, 4) = 4$$

$$A^1(3, 4) = \min(4, \infty + 1) \\ = 4$$

$$A^2 = \begin{bmatrix} 0 & 2 & 5 & 1 & 8 \\ 6 & 0 & 3 & 2 & 14 \\ \infty & \infty & 0 & 4 & \infty \\ \infty & \infty & 2 & 0 & 3 \\ 3 & 5 & 19 & 4 & 0 \end{bmatrix}$$

$$A^2(1,3) = 5 \quad A^2(3,1) = \infty \quad A^2(4,1) = \infty$$

$$A^2(1,4) = 1 \quad A^2(3,4) = 4 \quad A^2(4,3) = 2$$

$$A^2(1,5) = 8 \quad A^2(3,5) = \infty \quad A^2(4,5) = 3$$

$$A^2(5,1) = 3 \quad A^2(5,3) = 19 \quad A^2(5,4) = 4$$

$$A^3 = \begin{bmatrix} 0 & 2 & 5 & 1 & 8 \\ 6 & 0 & 3 & 2 & 14 \\ \infty & \infty & 0 & 4 & \infty \\ \infty & \infty & 2 & 0 & 3 \\ 3 & 5 & 19 & 4 & 0 \end{bmatrix}$$

$$A^3(1,2) = 2 \quad A^3(2,1) = 6 \quad A^3(4,1) = \infty$$

$$A^3(1,4) = 1 \quad A^3(2,4) = 2 \quad A^3(4,2) = \infty$$

$$A^3(1,5) = 8 \quad A^3(2,5) = 14 \quad A^3(4,5) = 3$$

$$A^3(5,1) = 3 \quad A^3(5,2) = 5 \quad A^3(5,4) = 4$$

$$A^4 = \begin{bmatrix} 0 & 2 & 3 & 1 & 4 \\ 6 & 0 & 3 & 2 & 5 \\ \infty & \infty & 0 & 4 & 7 \\ \infty & \infty & 2 & 0 & 3 \\ 3 & 5 & 6 & 4 & 0 \end{bmatrix} \quad \begin{array}{l} A^4(1,2) = 2 \\ A^4(1,3) = 3 \\ A^4(1,5) = 4 \end{array}$$

$$A^4(2,1) = 6 \quad A^4(3,1) = \infty \quad A^4(5,1) = 3$$

$$A^4(2,3) = 3 \quad A^4(3,2) = \infty \quad A^4(5,2) = 5$$

$$A^4(2,5) = 5 \quad A^4(3,5) = 7 \quad A^4(5,3) = 6$$

$$A^5 = \begin{bmatrix} 0 & 2 & 3 & 1 & 4 \\ 6 & 0 & 3 & 2 & 5 \\ 10 & 12 & 0 & 4 & 7 \\ 6 & 8 & 2 & 0 & 3 \\ 3 & 5 & 6 & 4 & 0 \end{bmatrix}$$

Algorithm:

Given a matrix A^0 , let $n=0$

1. Create a matrix with A^0 's size called A^{n+1}
2. Copy the n^{th} row and column of A^n to A^{n+1}
3. Apply the formula for $A^{n+1}(i,j)$

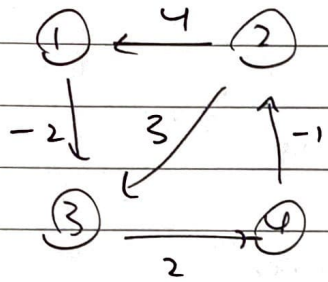
$$A^{n+1}(i,j) = \min [A^n(i,j), A^n(i,n+1) + A^n(n+1,j)]$$

4. Once all the entries are filled, increment n
5. Repeat the steps from 1 until $n = \text{size of matrix}$

Time complexity

$O(V^3)$ where $V = \text{no. of vertices}$

Question 2



List of edges :

$(1,3), (2,1), (2,3), (3,4), (4,2)$

$n = 4$

Phase 1

$$d(1) = 0$$

$$d(3) = d(1) - 2 \\ = -2 < \infty$$

$$d(1) = \infty > 0$$

$$d(3) = \infty > -2$$

$$d(4) = d(3) + 2 \\ = 0 < \infty$$

$$d(2) = d(4) - 1 \\ = -1 < \infty \\ = -1$$

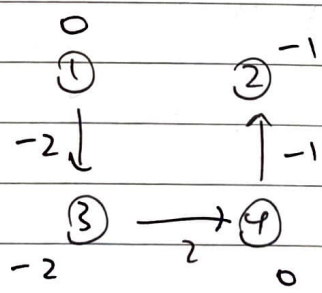
Phase 2

Initial distances:

$$d(1) = 0 \quad d(3) = -2 \\ d(2) = -1 \quad d(4) = 0$$

No changes on relaxation

\therefore The shortest paths are



Algorithm:

1. loop from 0 to $n-1$
2. Go through the list of edges
3. If $d(u) + \text{cost}(u, v) < d(v)$, then
update $d(v)$ at $d(u)$
4. Repeat until the looping is done
or no change occurs

Time Complexity

$O(v \times e)$ where v = vertices count
 e = edges count

Question 3

1. Loop from the start of the string till the end
2. Check if the i^{th} to $i+l^{\text{th}}$ index match the pattern of length l
3. If the pattern matches, display i
4. If no pattern matches, display as such

KMP Algorithm:

1. Create a prefix table for the string pattern as pi
2. Keep track of letters in the string as i
3. Set j as 0
4. If $i = \text{pattern}[j+1]$, increment j
5. If the letters don't match, then set j as $\text{pattern}[j]$'s pi and set i as the next letter in the string
6. If $j = \text{length of pattern}$, sequence is found

Rabin Karp Algorithm:

1. Hash the string and pattern and find the expected remainder
2. Loop through the numbers of same length as the pattern at a time
3. If the $\text{hit}(\text{mod calculation})$ is spurious, then check if the numbers match
4. If match, sequence is detected

Examples:

String: ABC CDD A EFG

Pattern: CDD

Naive method:

1. ABC \neq CDD

2. BCC \neq CDD

3. CCD \neq CDD

4. CDD = CDD, sequence found at pos. 4

KMP method:

Pi table:

index	0	1	2
char	C	D	D
pi	0	0	1

1. ~~ABC~~ A \neq C

2. B \neq C

3. C = C, inc j

4. C \neq D, go to 0

5. C = C, inc j

6. D = D, inc j

7. D = D, inc j, j = len of pattern, sequence found

//_

Rabin Karp method:

$$m = 10, n = 3$$

{ A, B, C, D, E, F, G, H, I, J }

1 2 3 4 5 6 7 8 9 10

choosing prime = 13

$$p = 13$$

$$\text{hash}(p) = \sum (v * m^x) \cdot 13 \text{ where } v = \text{character}$$

$$m = 10$$

$x = \text{position}$

$$\begin{aligned} \therefore \text{hash}(p) &= (3 \times 10^2) + (4 \times 10^1) + (4 \times 10^0) \mod 13 \\ &= 344 \cdot 13 \\ &= \underline{6} \end{aligned}$$

hash value for each pair of 3 (len of pattern)

A B C C D D A E F G

$$\Rightarrow \quad 6 \quad 12 \quad 9 \quad 6 \quad 6^{12} \quad 12 \quad 0 \quad 8$$

↑
spurious
hit

↑
spurious
hit
↑
exact
match

\therefore Sequence found at position 4