

Anonymity in the Bitcoin Peer-to-Peer Network

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“Untraceable Bitcoin”

Teenagers using untraceable currency Bitcoin to buy dangerous drugs online

Fears have been raised as children as young as 14 are getting parcels of legal highs delivered to their home

Mirror

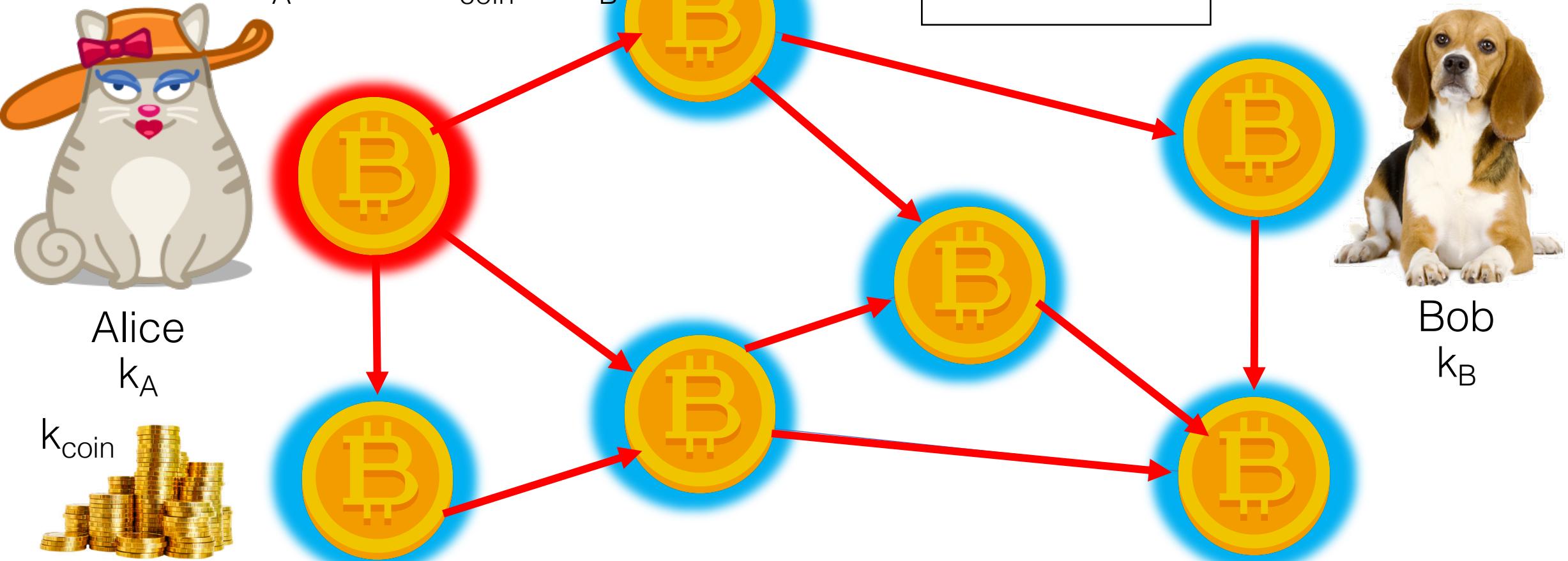


This is false.

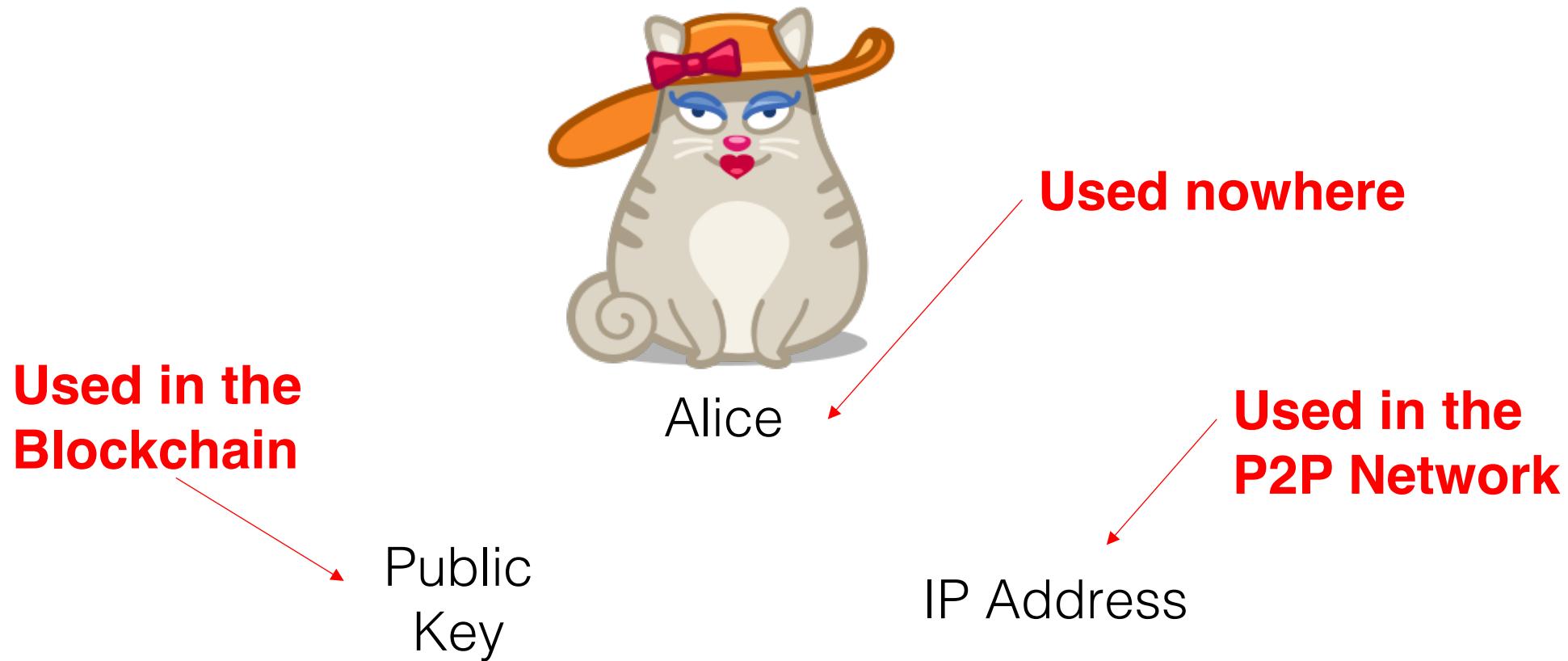
Bitcoin Primer

Transaction

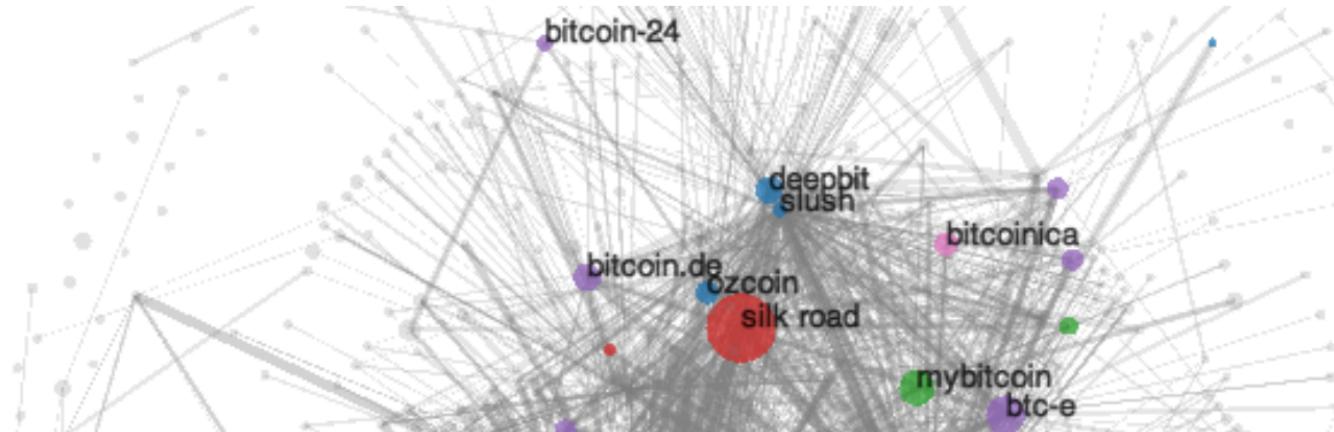
k_A sends k_{coin} to k_B



Multiple Identities



How can users be deanonymized?



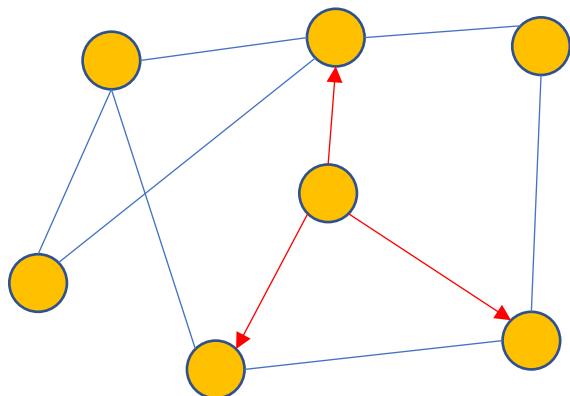
Entire transaction histories
can be compromised.

What about the peer-to-peer
network?

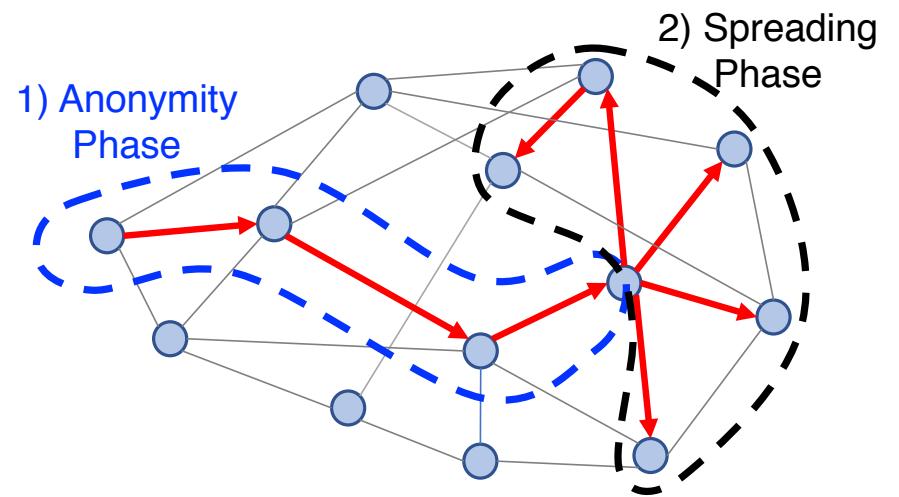
Public Key ← → IP Address

This Talk

How to break privacy



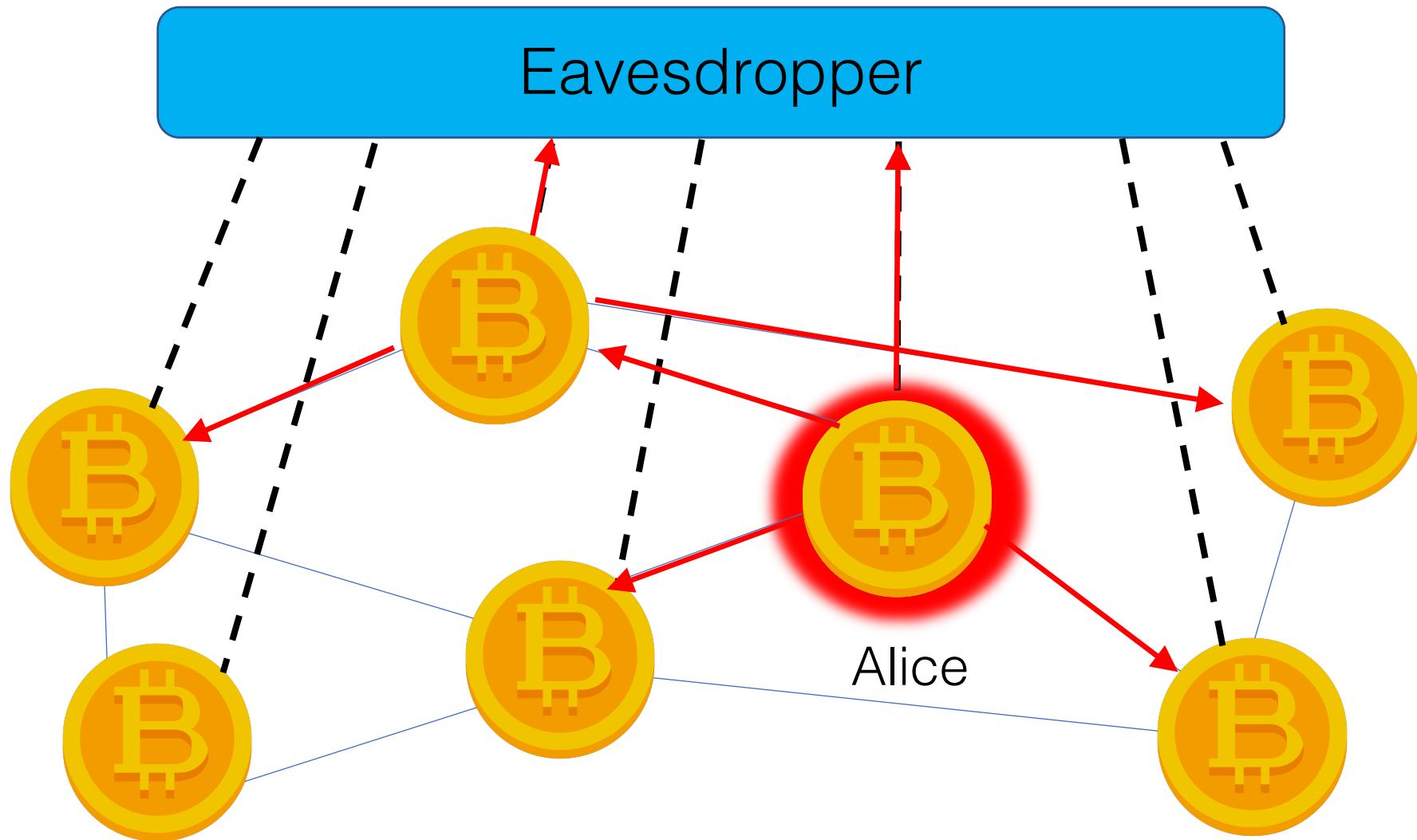
How to fix it



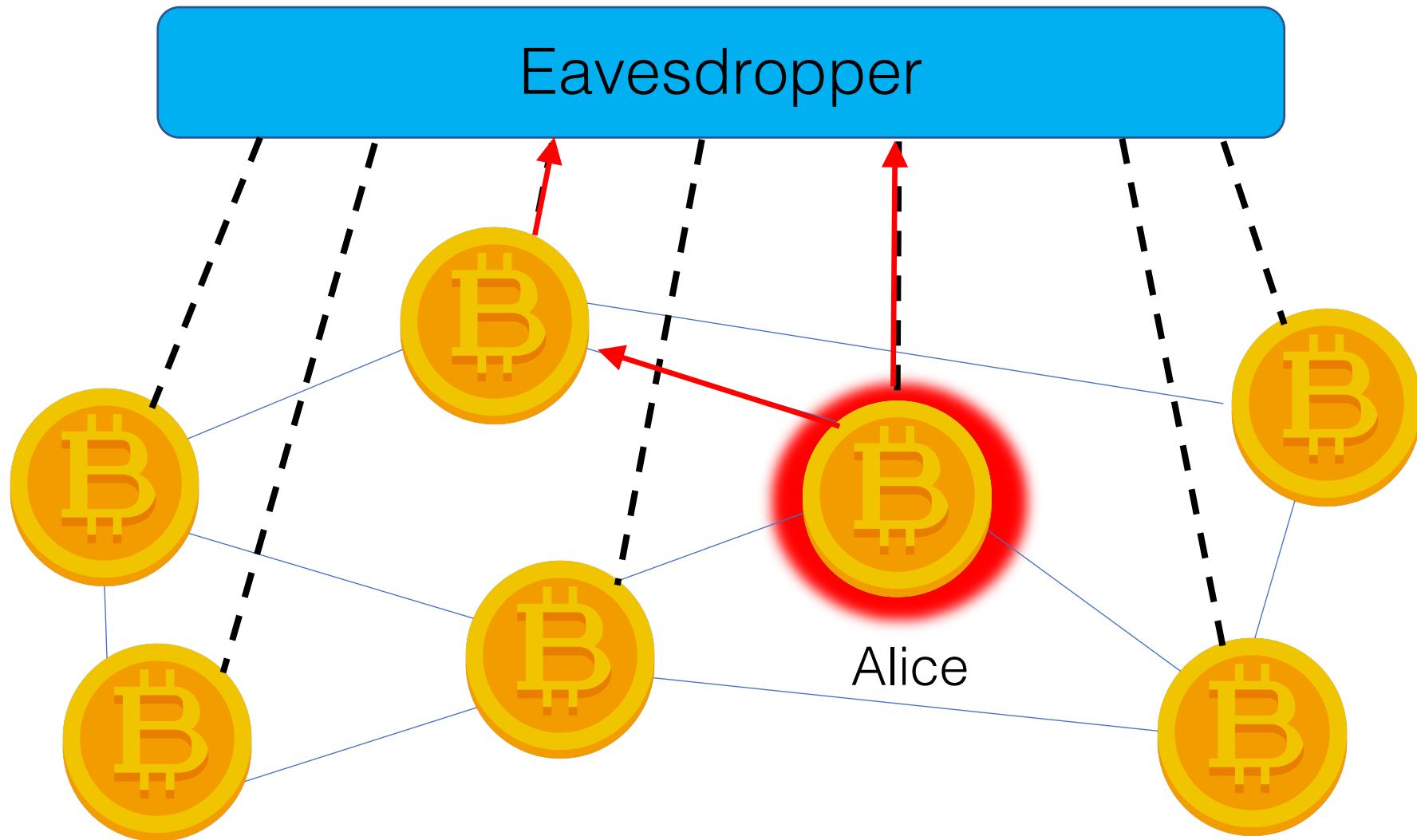
Early attacks

- A. Biryukov, D. Khovratovich, I. Pustagurov, “*Deanonymisation of clients in Bitcoin P2P network*”, CCS 2014
- P. Koshy, D. Koshy, P. McDaniel, “*An analysis of anonymity in Bitcoin using P2P network traffic*”, Financial Crypto 2014

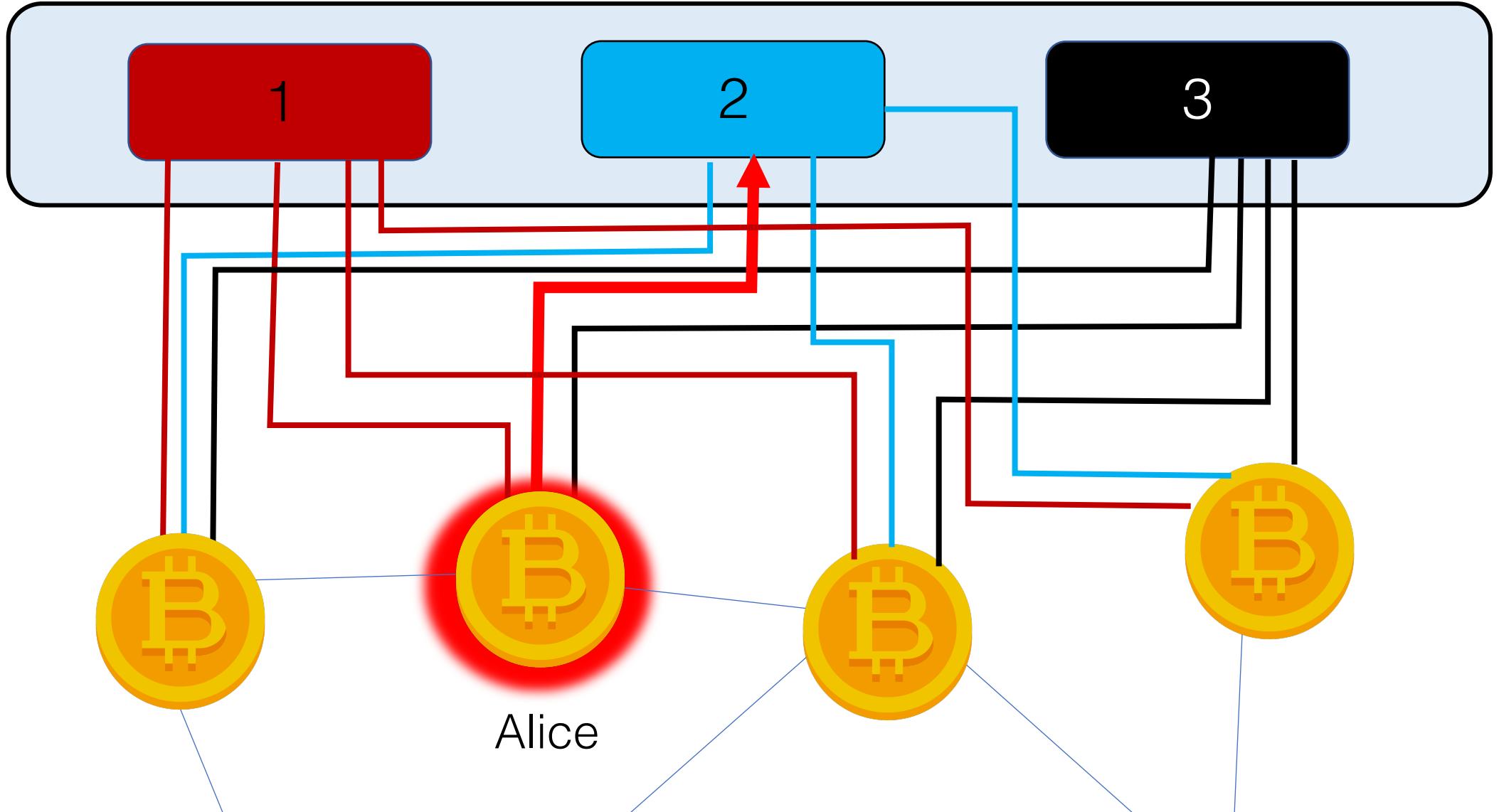
Attacks on the Network Layer



What can go wrong?



What the eavesdropper can do about it

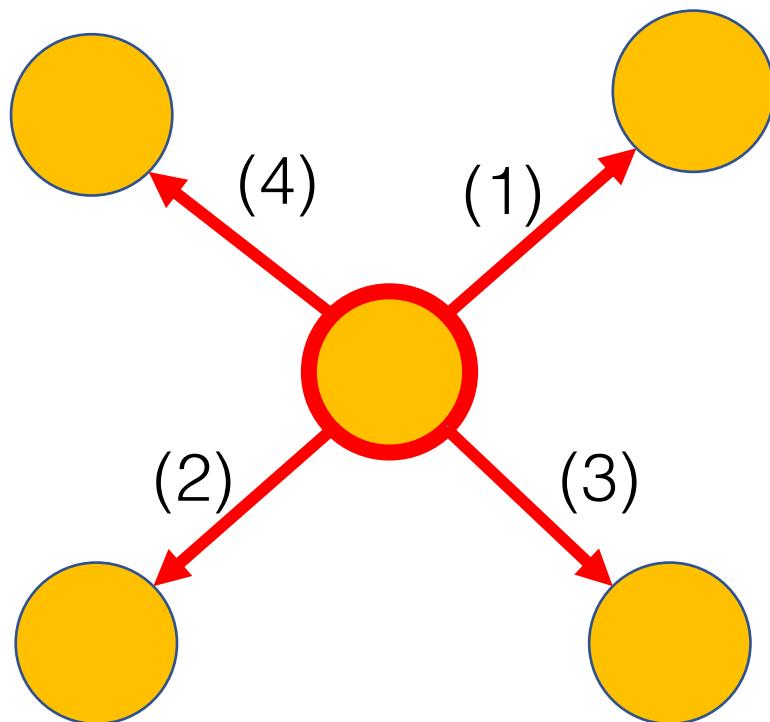


Key Results

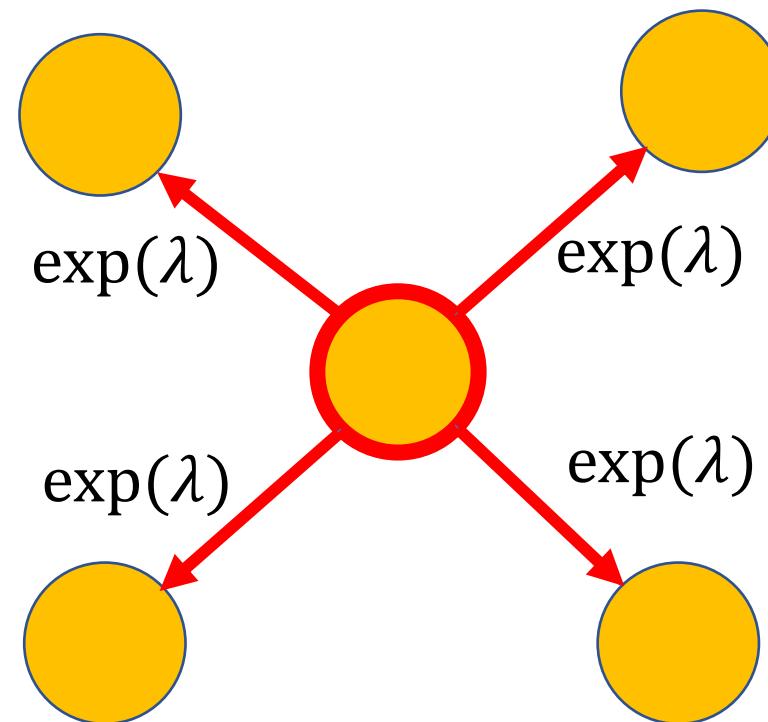
- Make $\theta \approx 50$ connections per node
- Between **11-34% of users deanonymized**, *even behind NAT!*

Bitcoin Core Responds

Trickle (pre-2015)



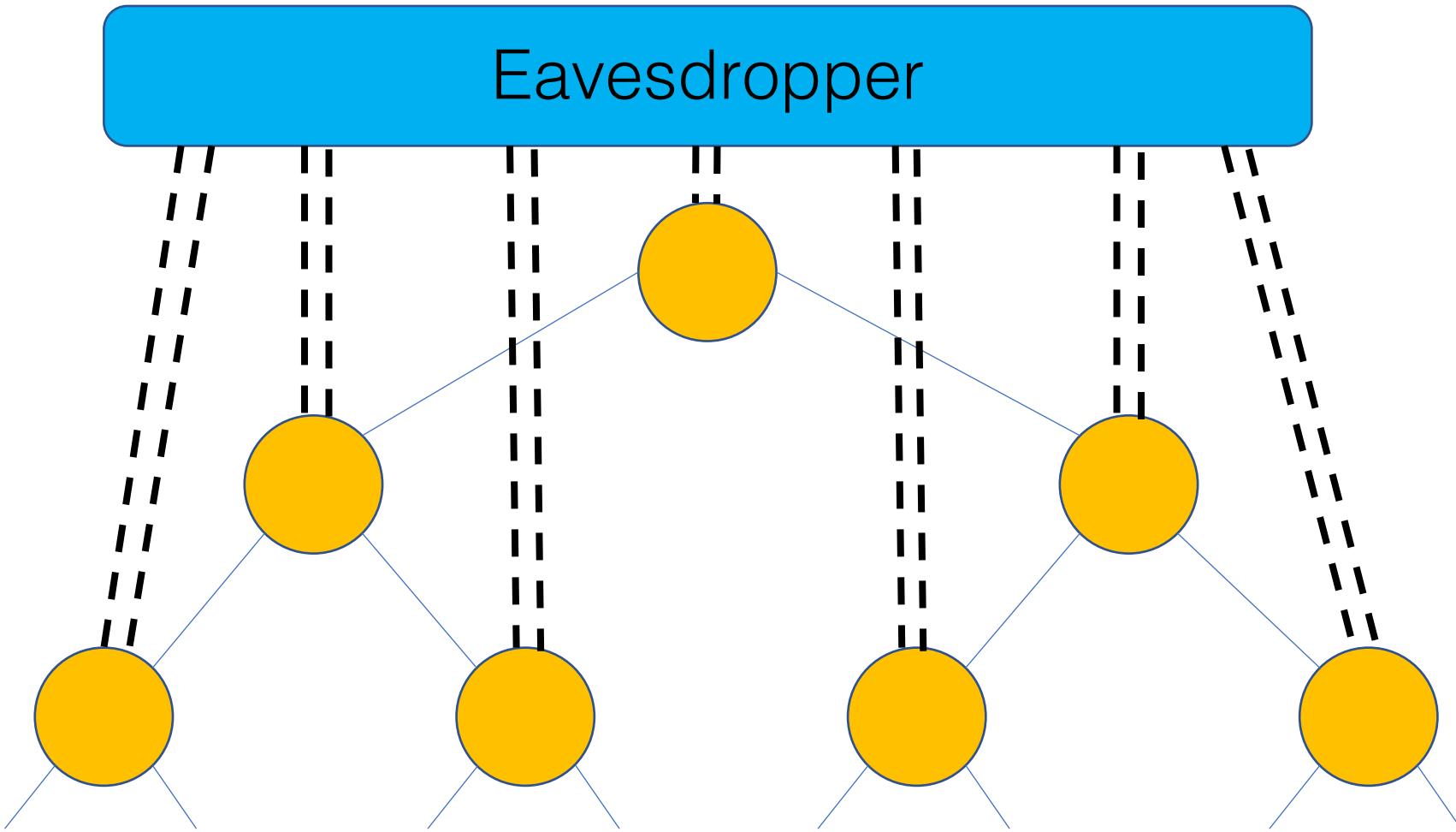
Diffusion (post-2015)



Does diffusion provide stronger
anonymity than trickle spreading?

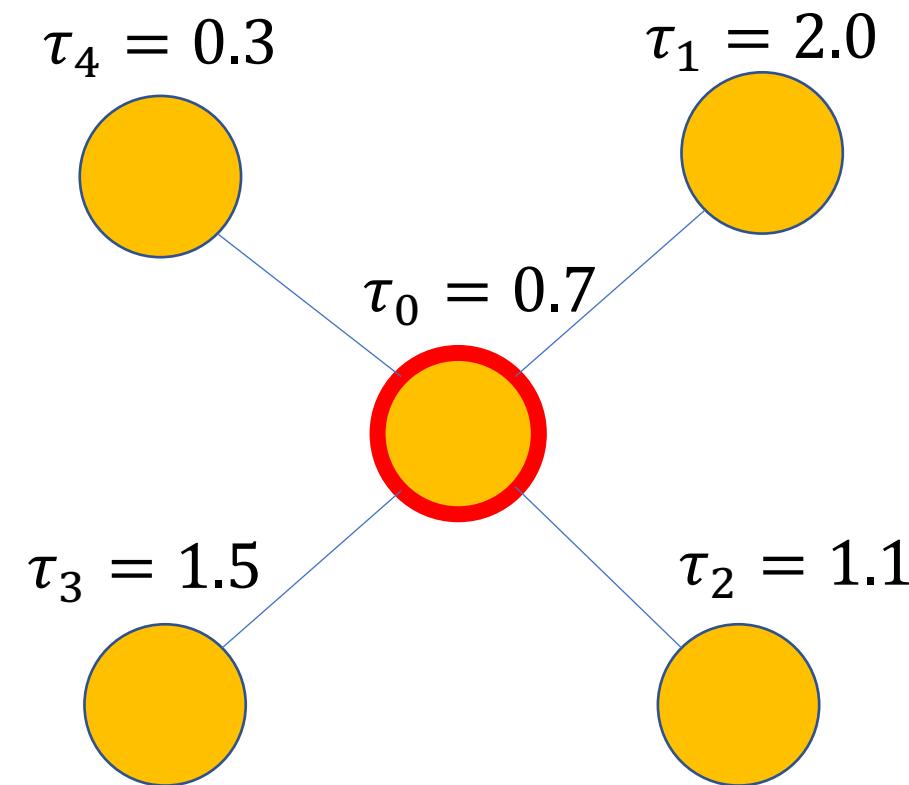
G. F., P. Viswanath, “*Anonymity in the Bitcoin P2P Network*”,
NeurIPS 2017

d -regular trees



Anonymity Metric

$$\boldsymbol{\tau} = \begin{bmatrix} \tau_1 \\ \tau_2 \\ \dots \\ \tau_n \end{bmatrix}$$



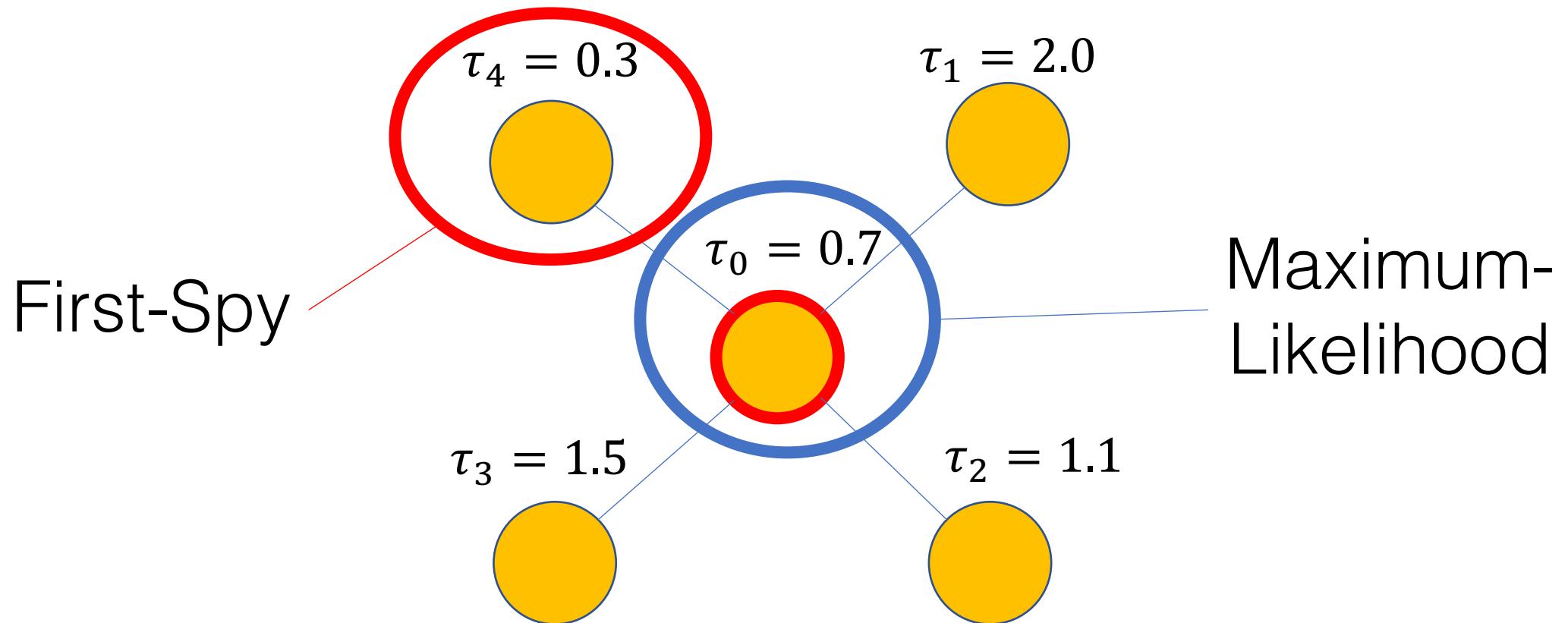
$$P(\text{detection} | \boldsymbol{\tau}, G)$$

↑
timestamps
graph

Estimators

$P(\text{detection} | \tau, G)$

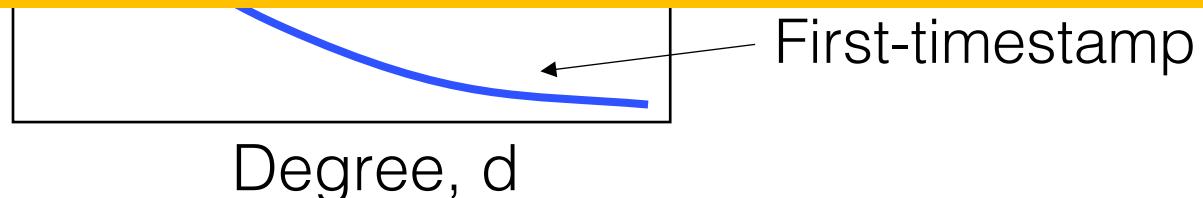
timestamps
graph



Results: d-Regular Trees

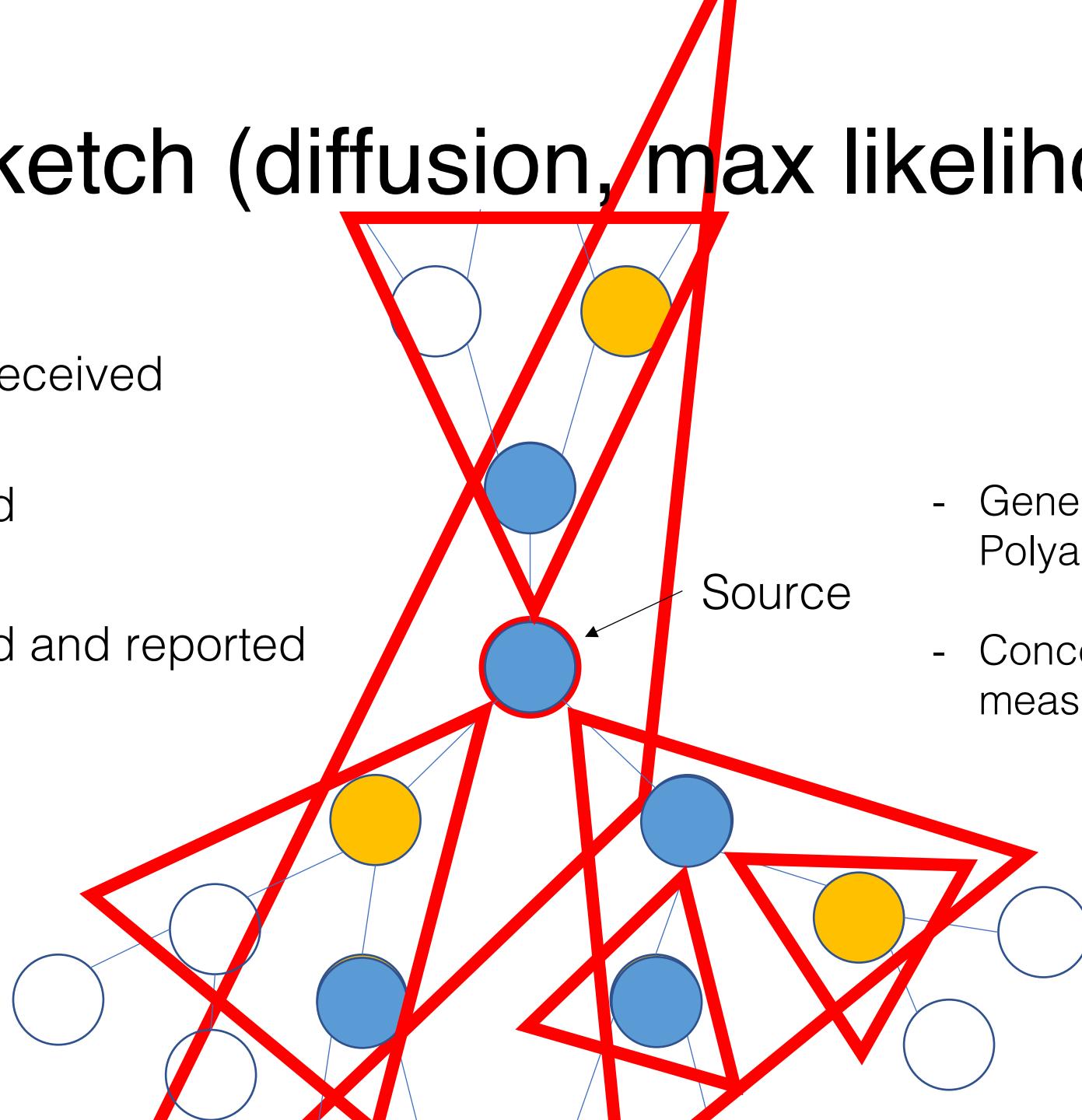
	Trickle	Diffusion
First-Timestamp	$o\left(\frac{\log d}{d}\right)$	$o\left(\frac{\log d}{d}\right)$
Maximum-Likelihood	$\Omega(1)$	$\Omega(1)$

Intuition: Symmetry outweighs local randomness!



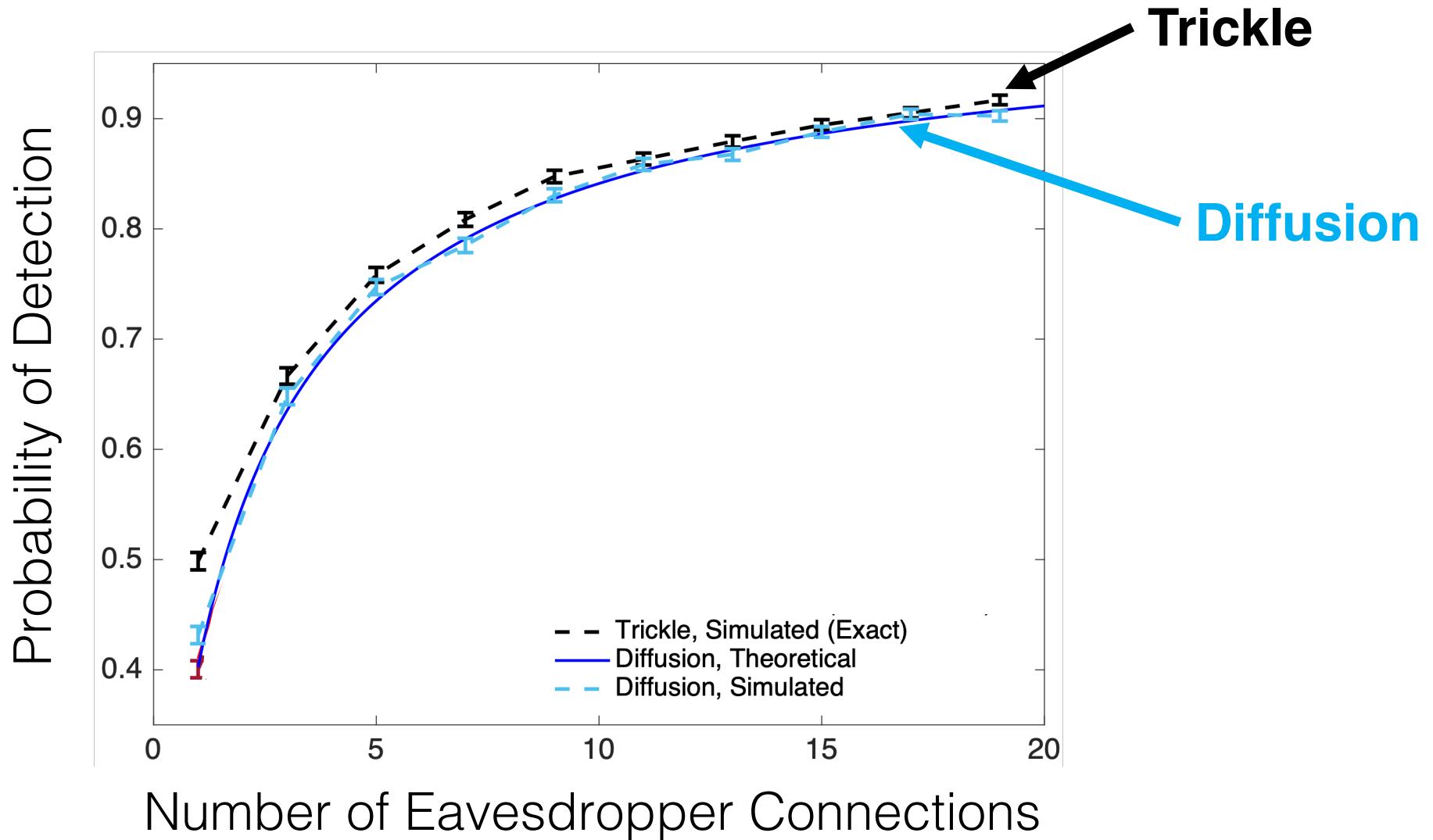
Proof sketch (diffusion, max likelihood)

- Not yet received
- Received
- Received and reported

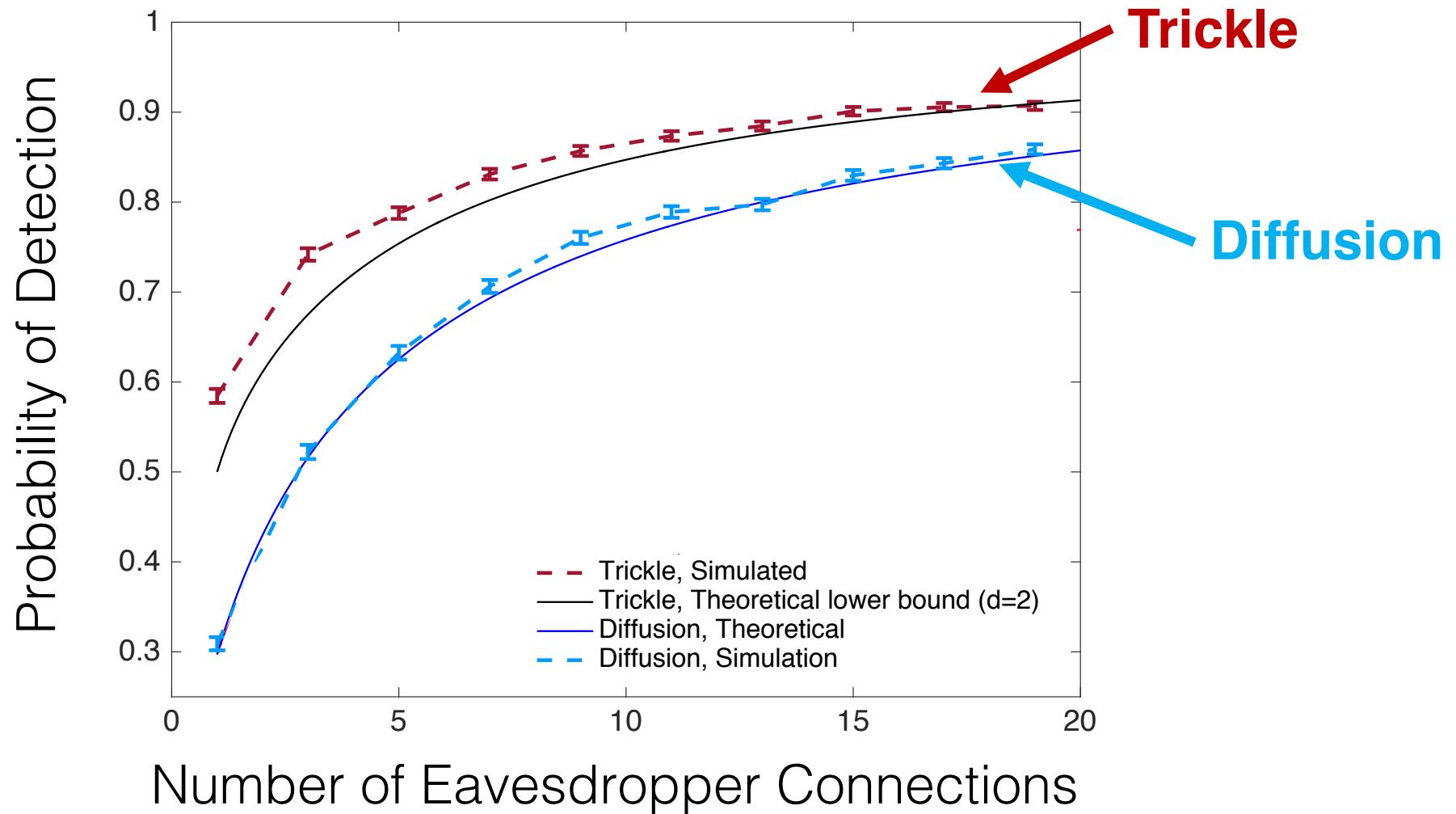


- Generalized Polya Urns
- Concentration of measure

Results: Trees



Results: Bitcoin Graph



Diffusion does not have
(significantly) better anonymity
properties than trickle.

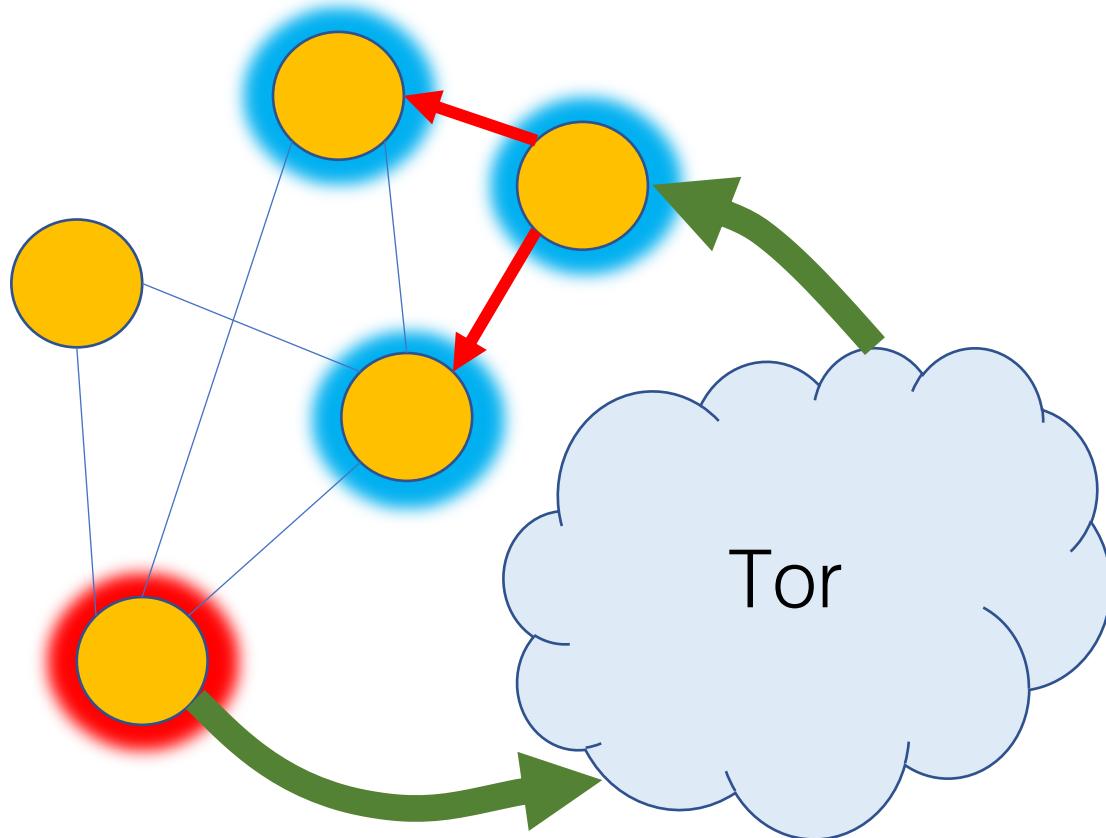
Redesign

Can we fix this problem?

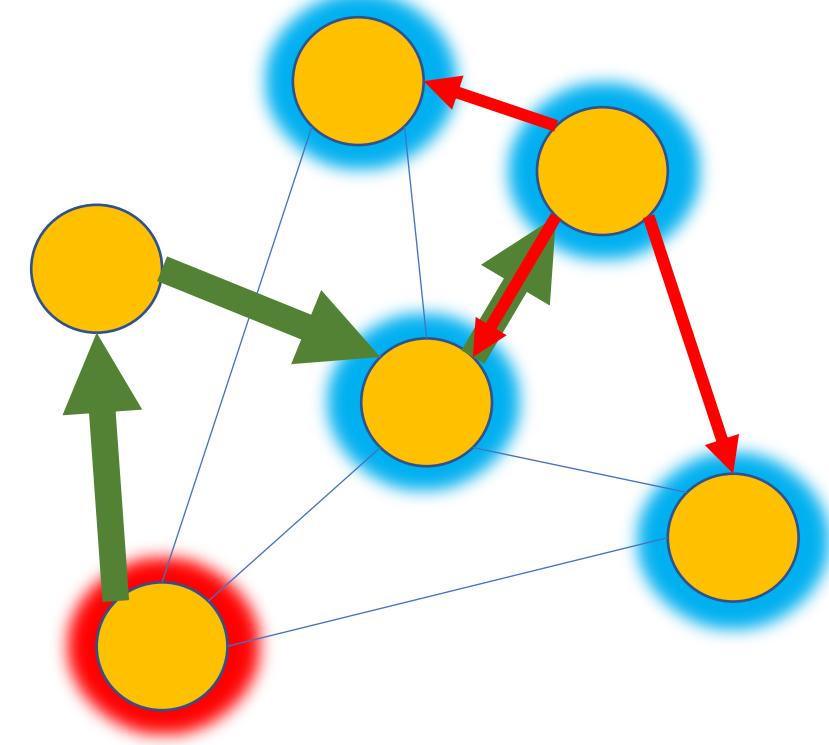


First-order solutions

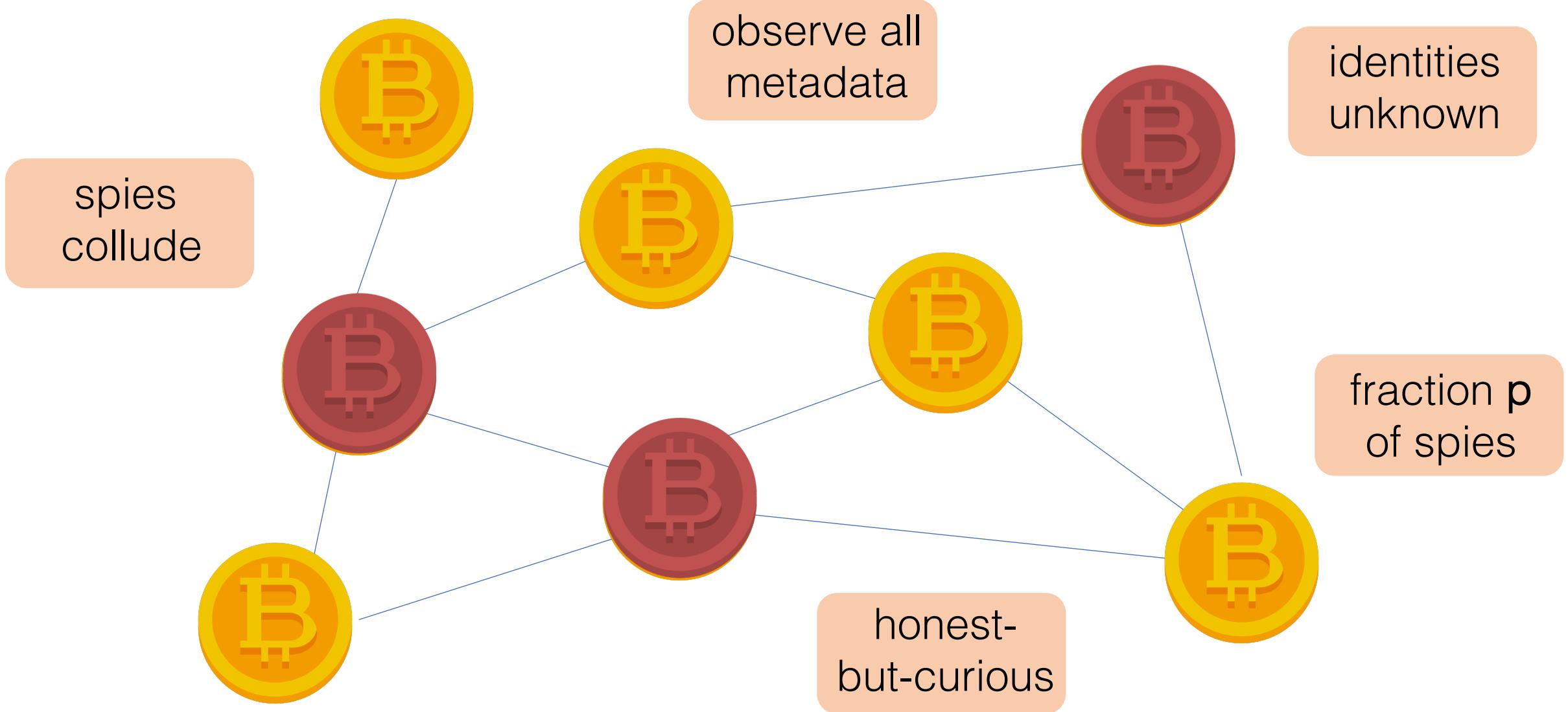
Connect through Tor



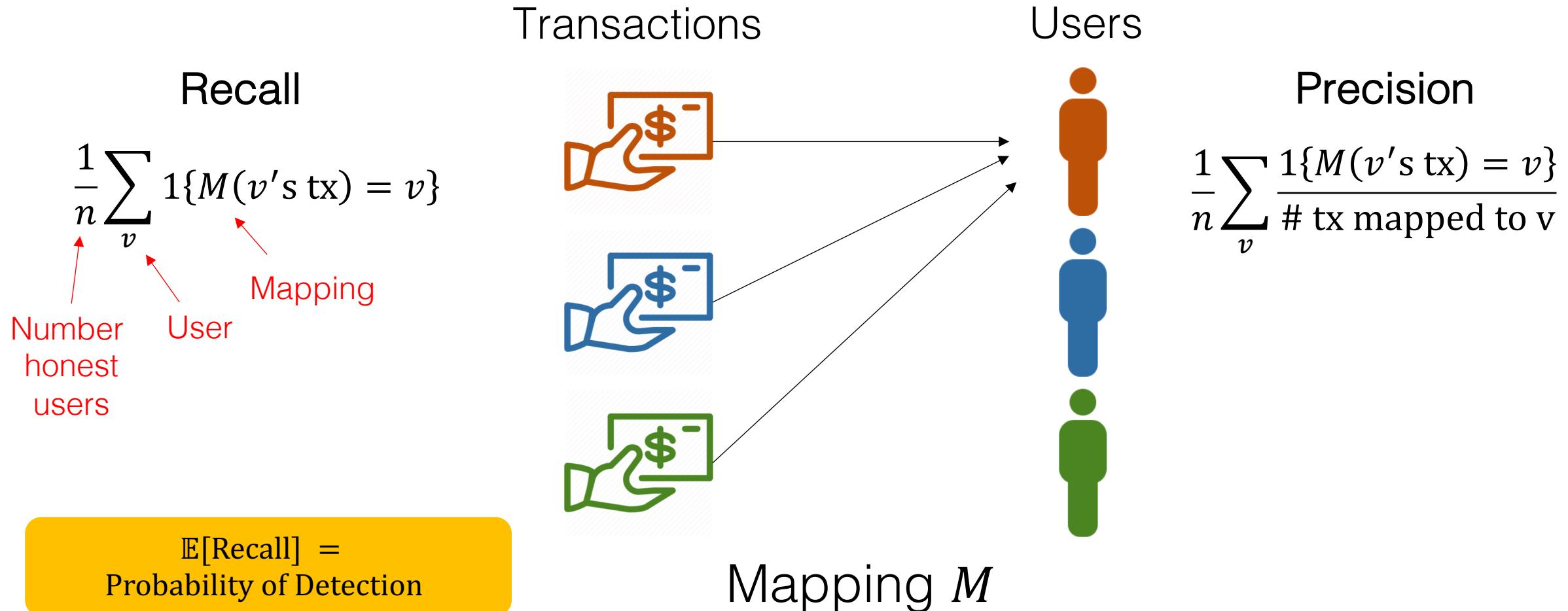
I2P Integration (e.g. Monero)



Botnet adversarial model



Metric for Anonymity

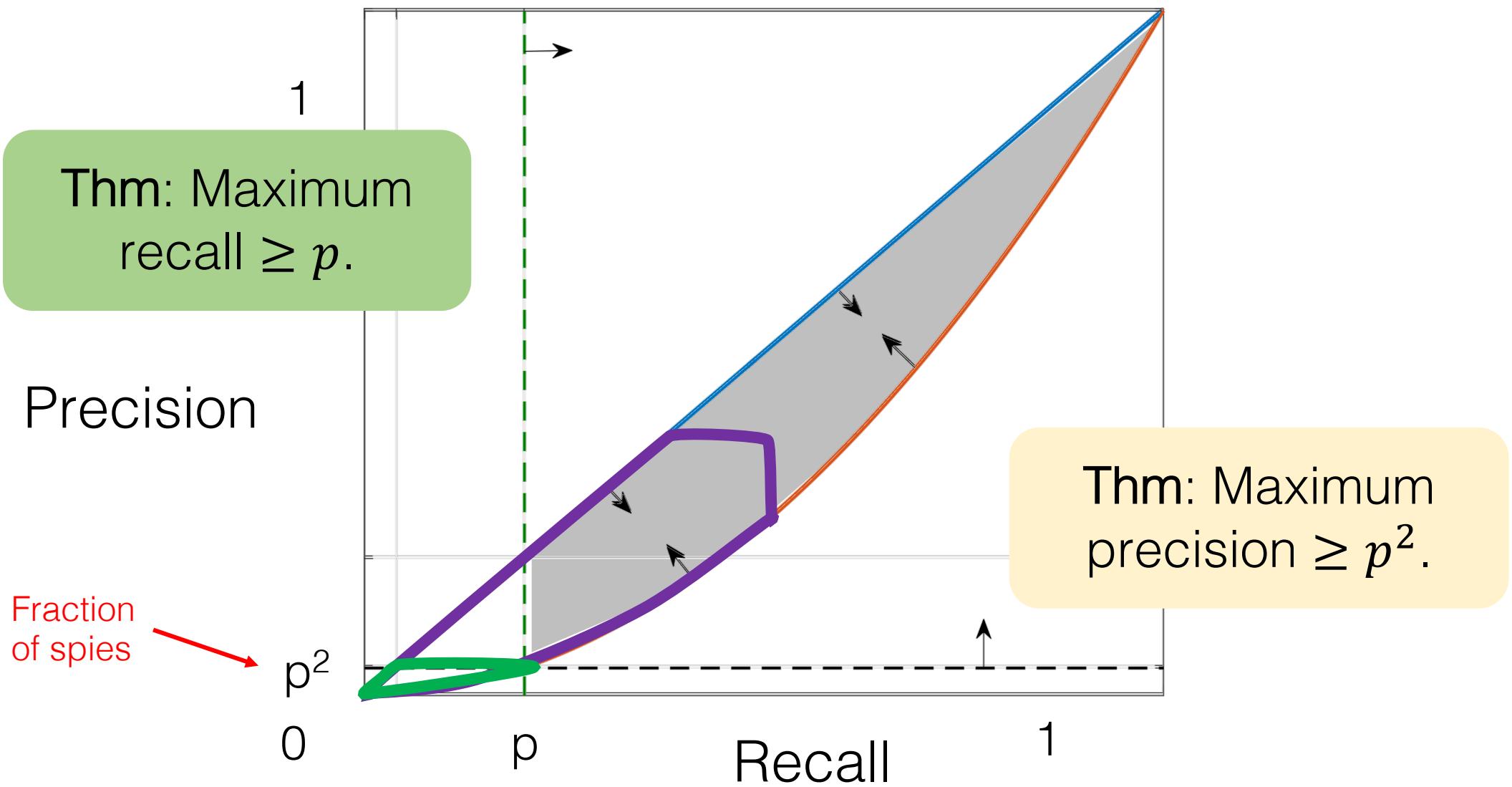


Goal:

Design a distributed flooding protocol that minimizes the maximum **precision** and **recall** achievable by a computationally-unbounded adversary.

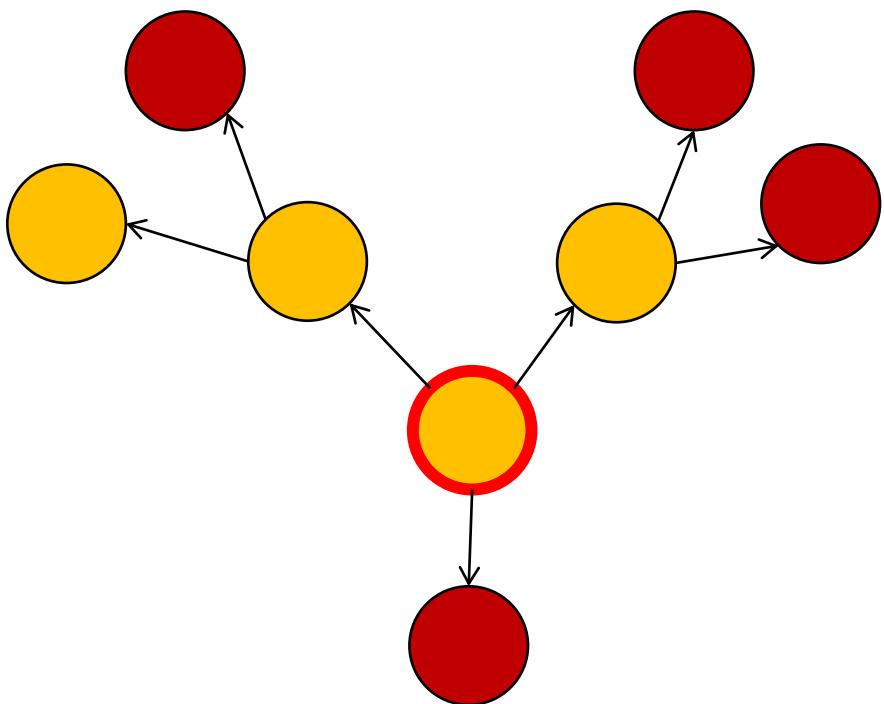
S. B. Venkatakrishnan, G. F., P. Viswanath, “*Dandelion: Redesigning the Bitcoin Network for Anonymity*”, Sigmetrics 2017

Fundamental Limits

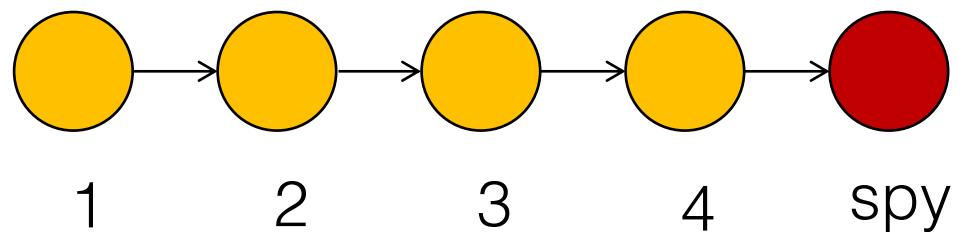


What are we looking for?

Asymmetry

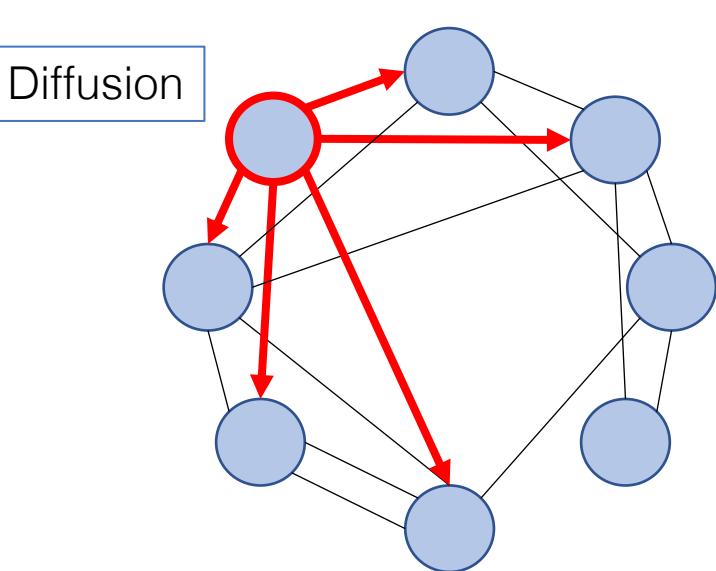


Mixing



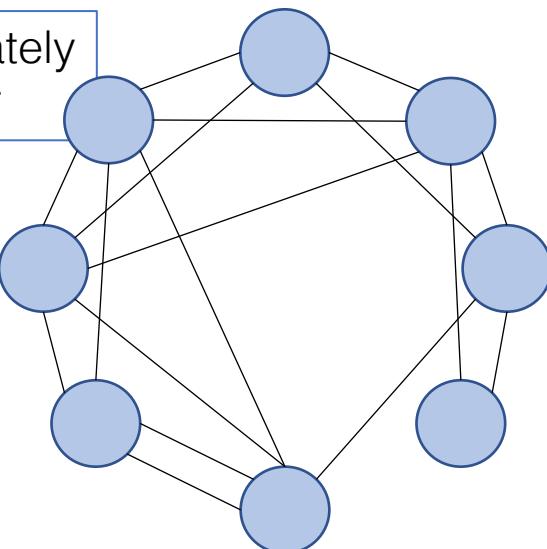
What can we control?

Spreading Protocol



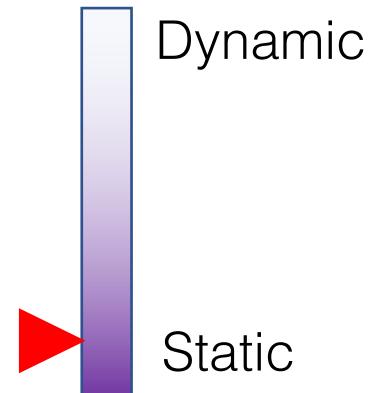
*Given a graph, how
do we spread content?*

Topology



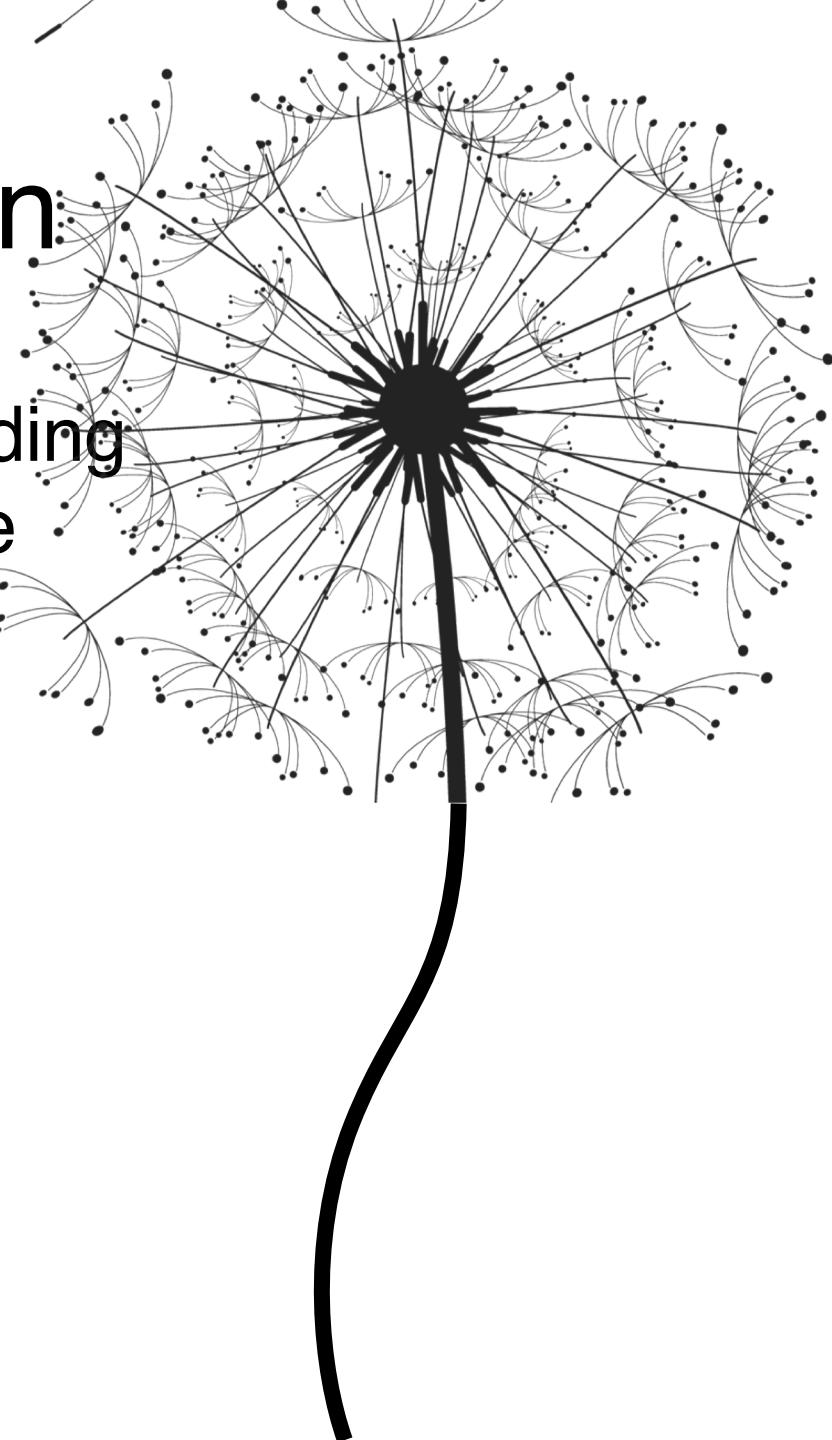
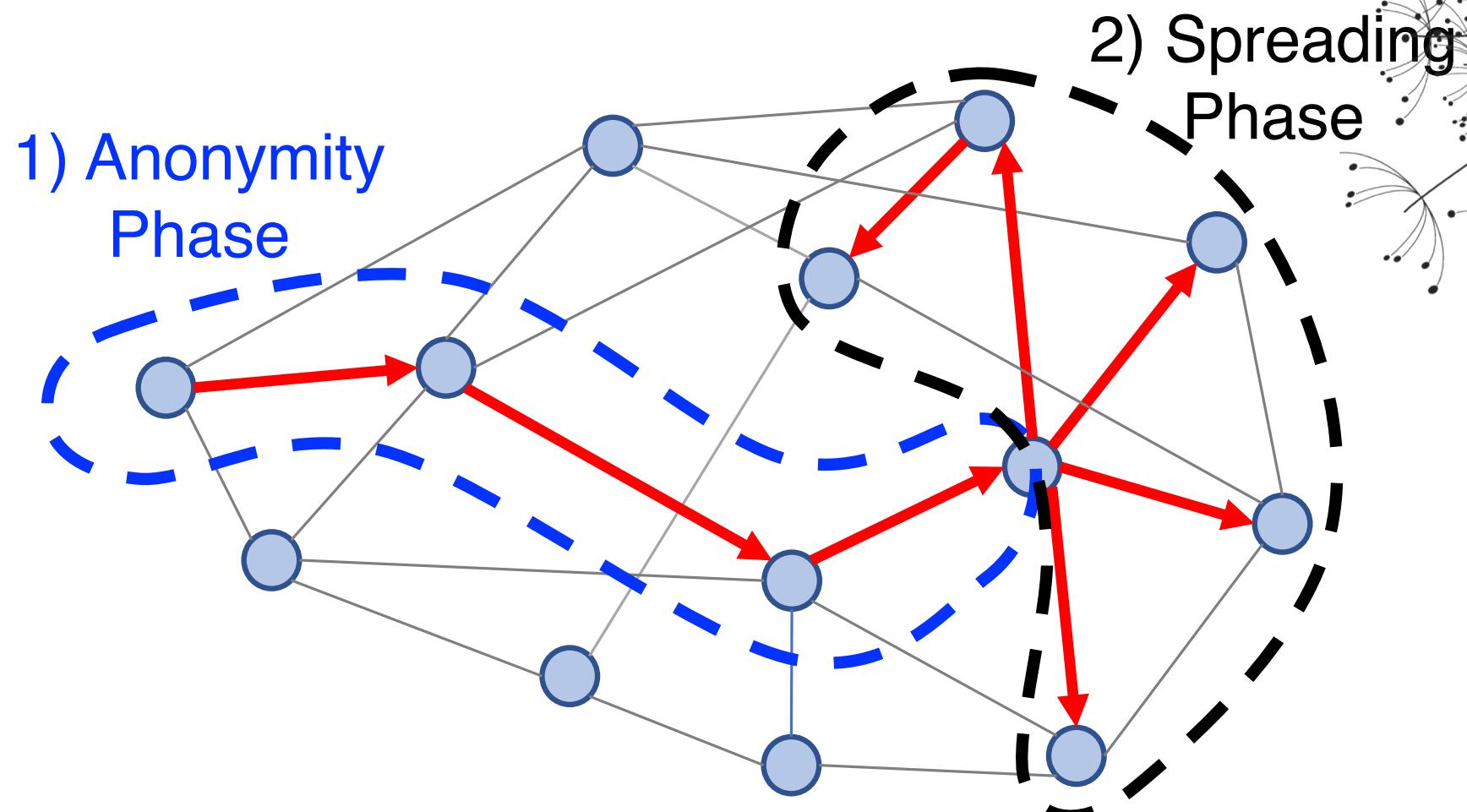
*What is the underlying
graph topology?*

Dynamicity



*How often does the
graph change?*

Spreading Protocol: Dandelion



Why Dandelion spreading?

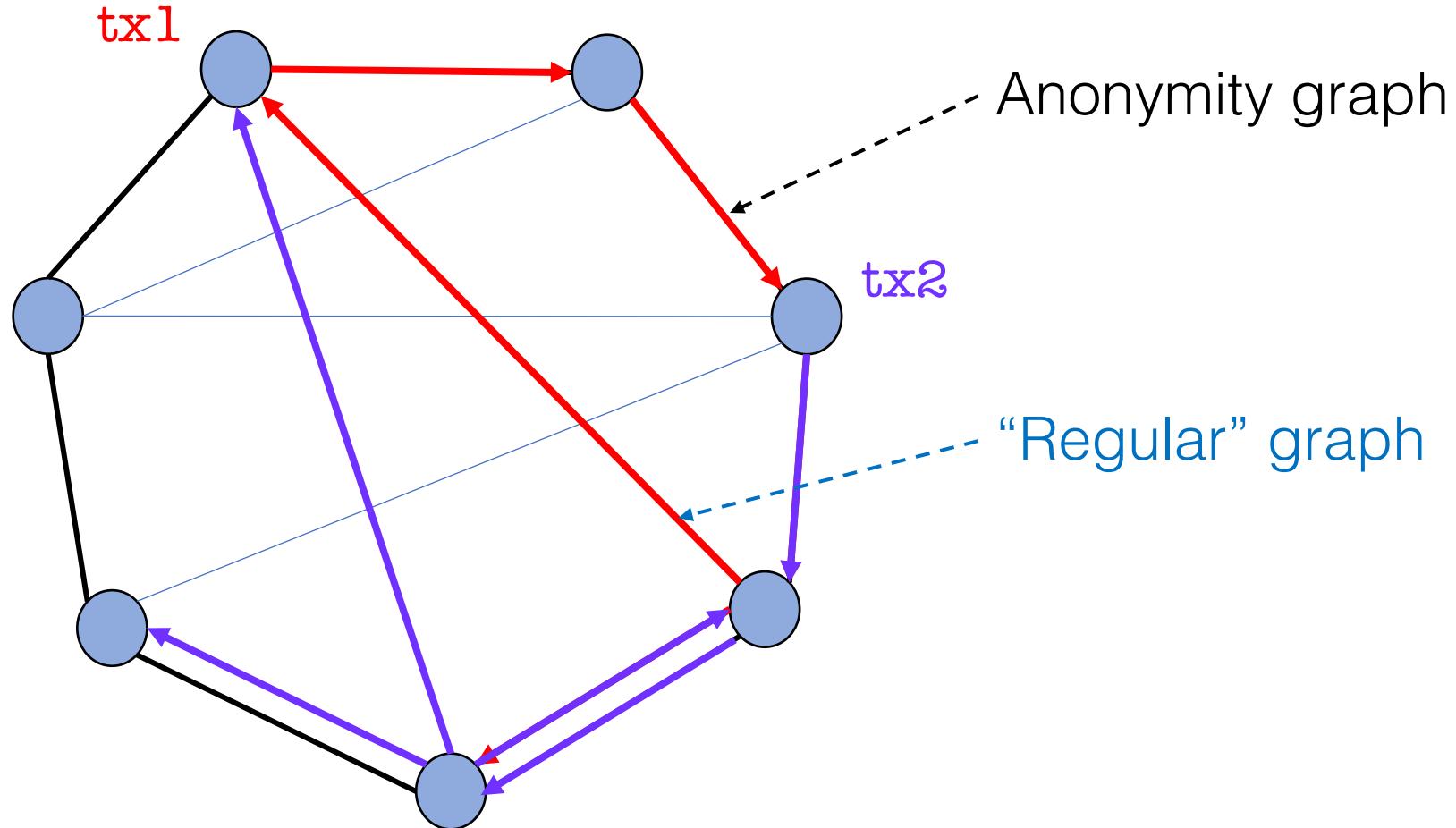
Theorem: Dandelion spreading has an optimally low maximum recall of $p + o\left(\frac{1}{n}\right)$.

Theorem: Fundamental lower bound = p

fraction
of spies

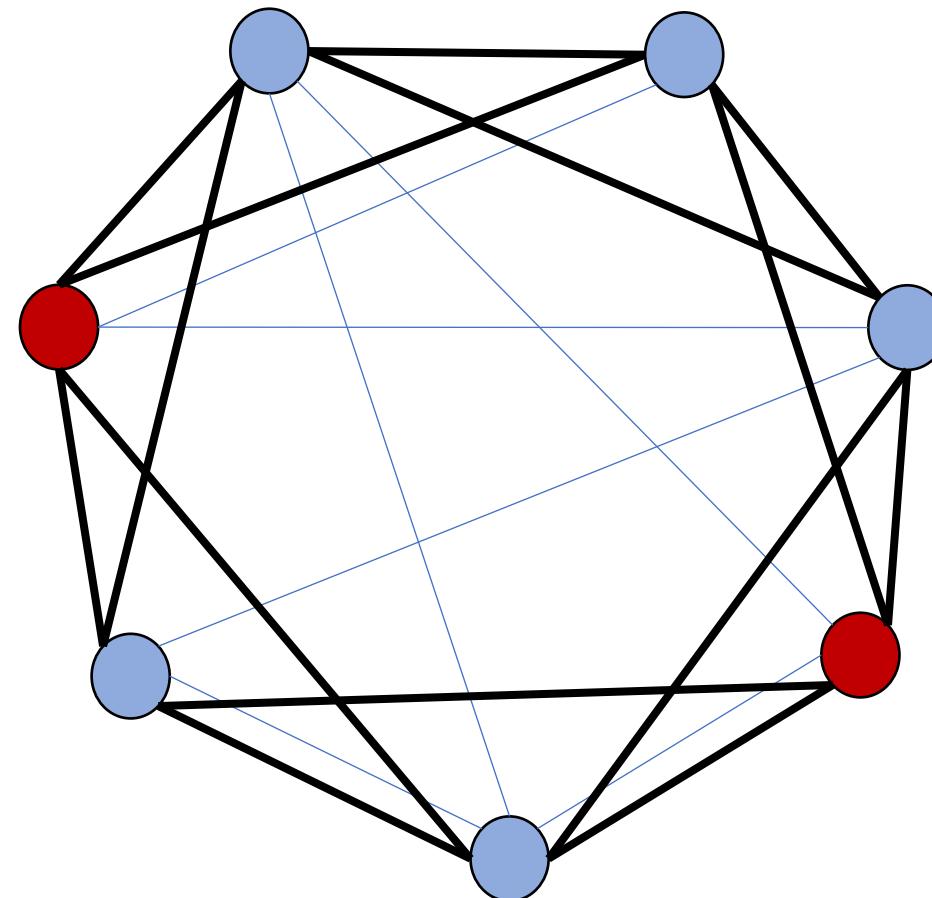
number of
nodes

Graph Topology: Line



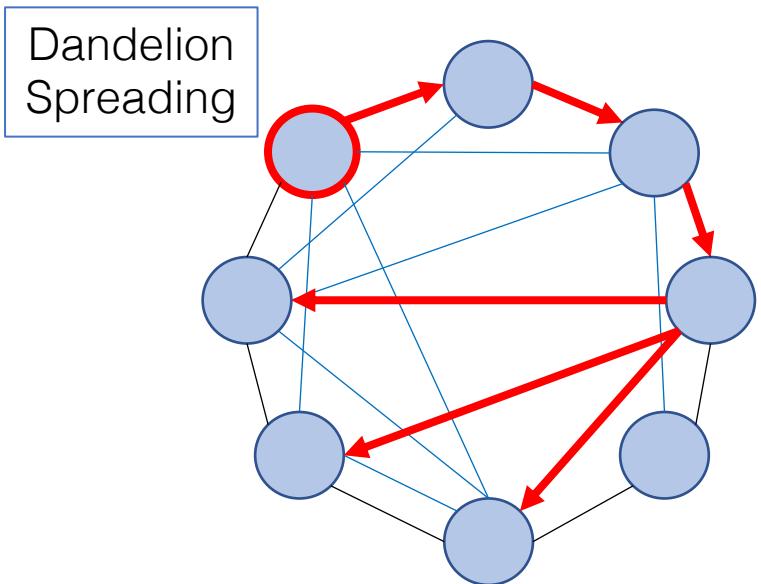
Dynamicity: High

Change the anonymity graph frequently.



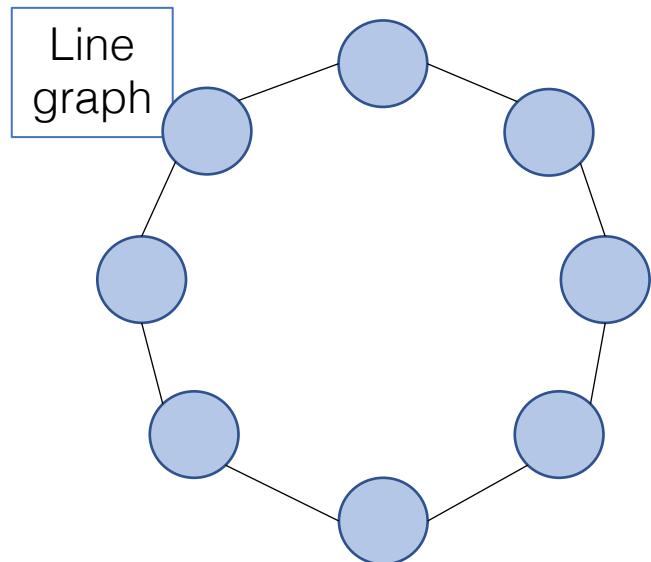
DANDELION Network Policy

Spreading Protocol



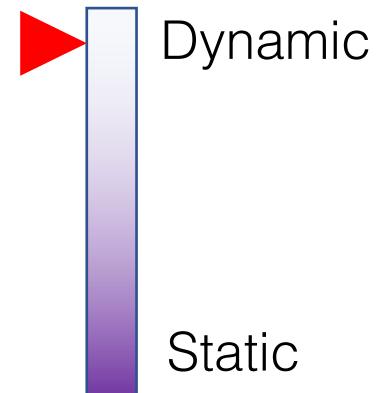
*Given a graph, how
do we spread content?*

Topology



*What is the anonymity
graph topology?*

Dynamicity



*How often does the
graph change?*

Theorem: Fundamental lower bound = p^2

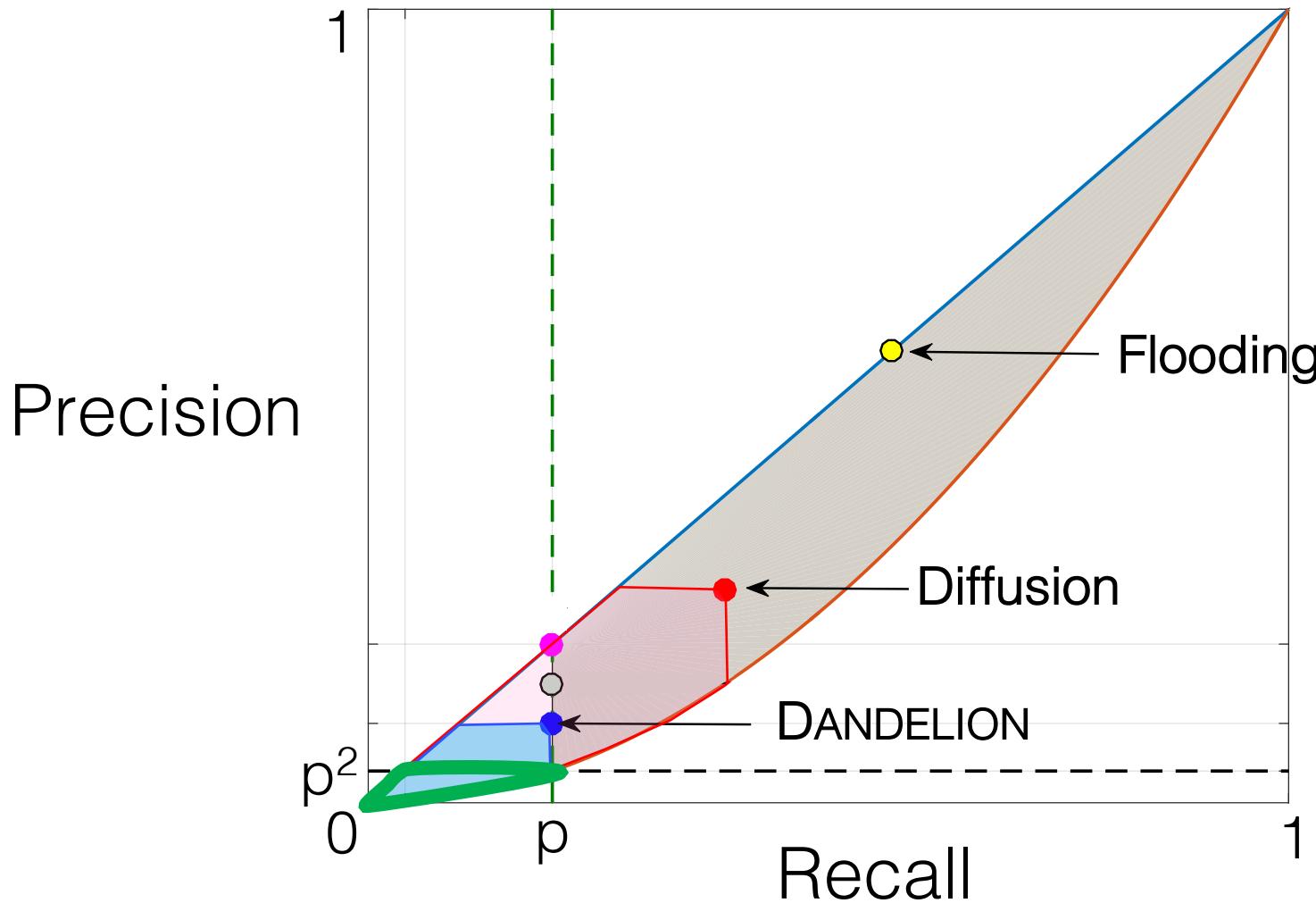
Theorem: DANDELION has a nearly-optimal maximum precision of $\frac{2p^2}{1-p} \log\left(\frac{2}{p}\right) + O\left(\frac{1}{n}\right)$.*

fraction
of spies

number of
nodes

*For $p < \frac{1}{3}$

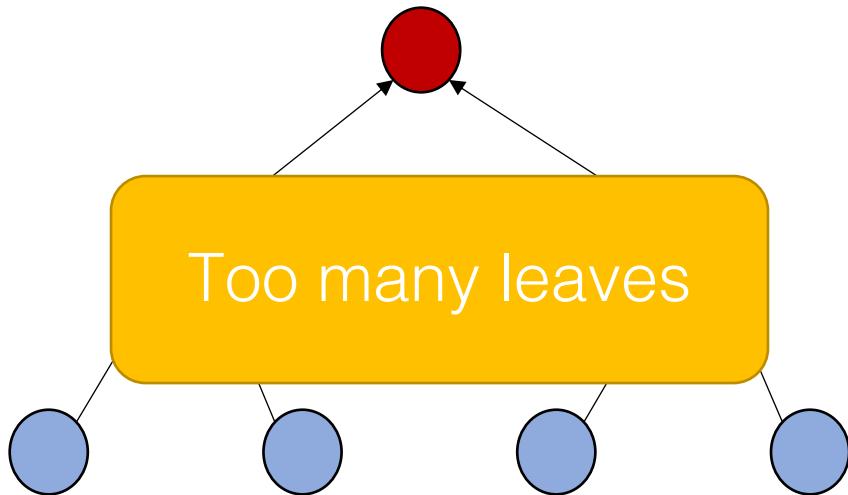
Performance: Achievable Region



Why is DANDELION good?

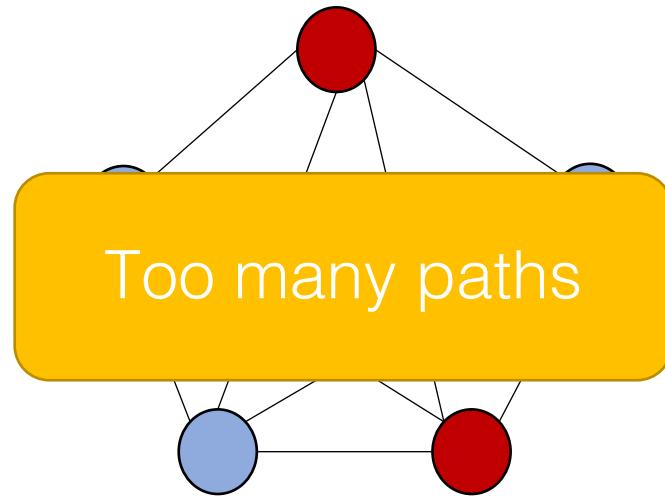
Strong mixing properties.

Tree



Precision: $O(p)$

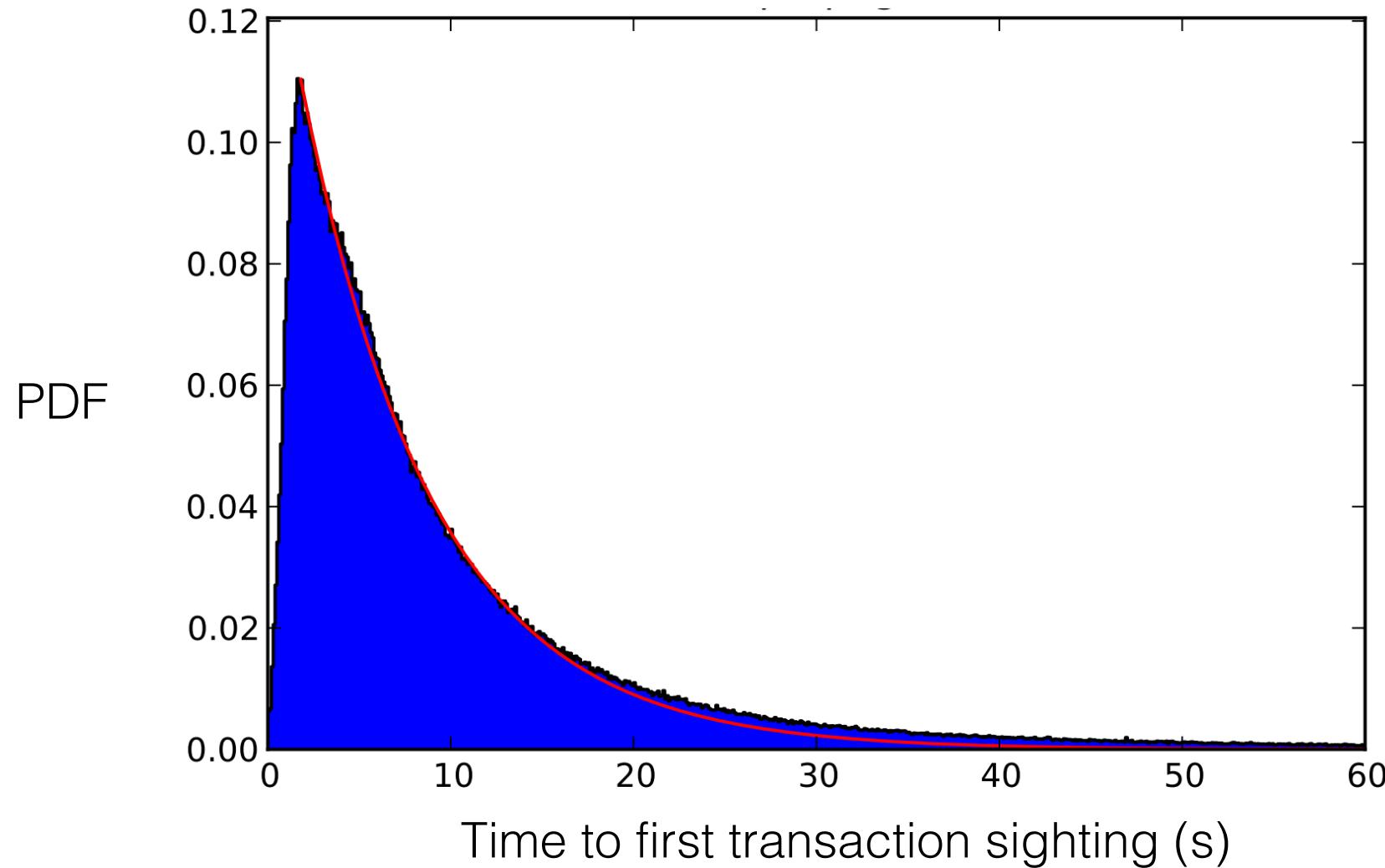
Complete graph



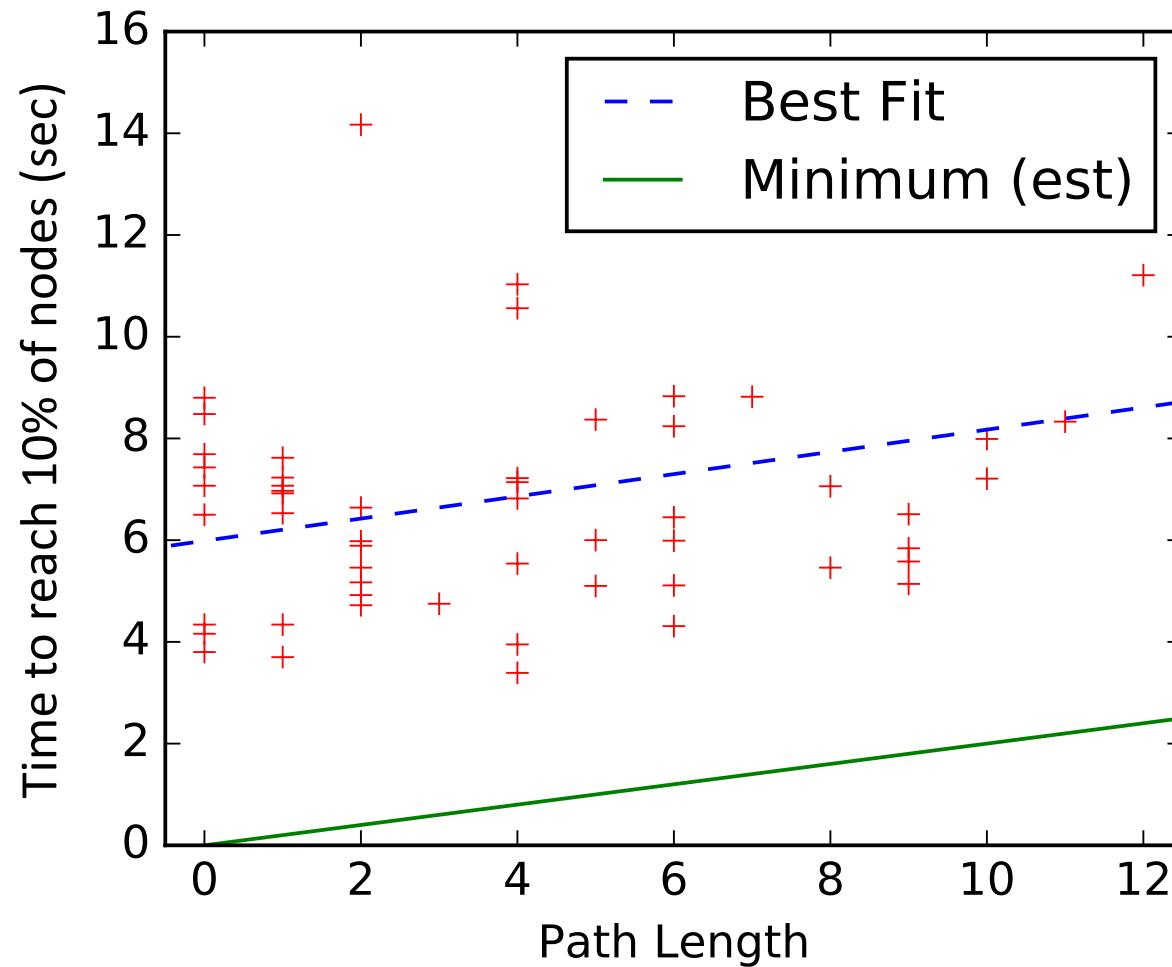
Precision: $\frac{p}{1-p} (1 - e^{p-1})$

How practical is this?

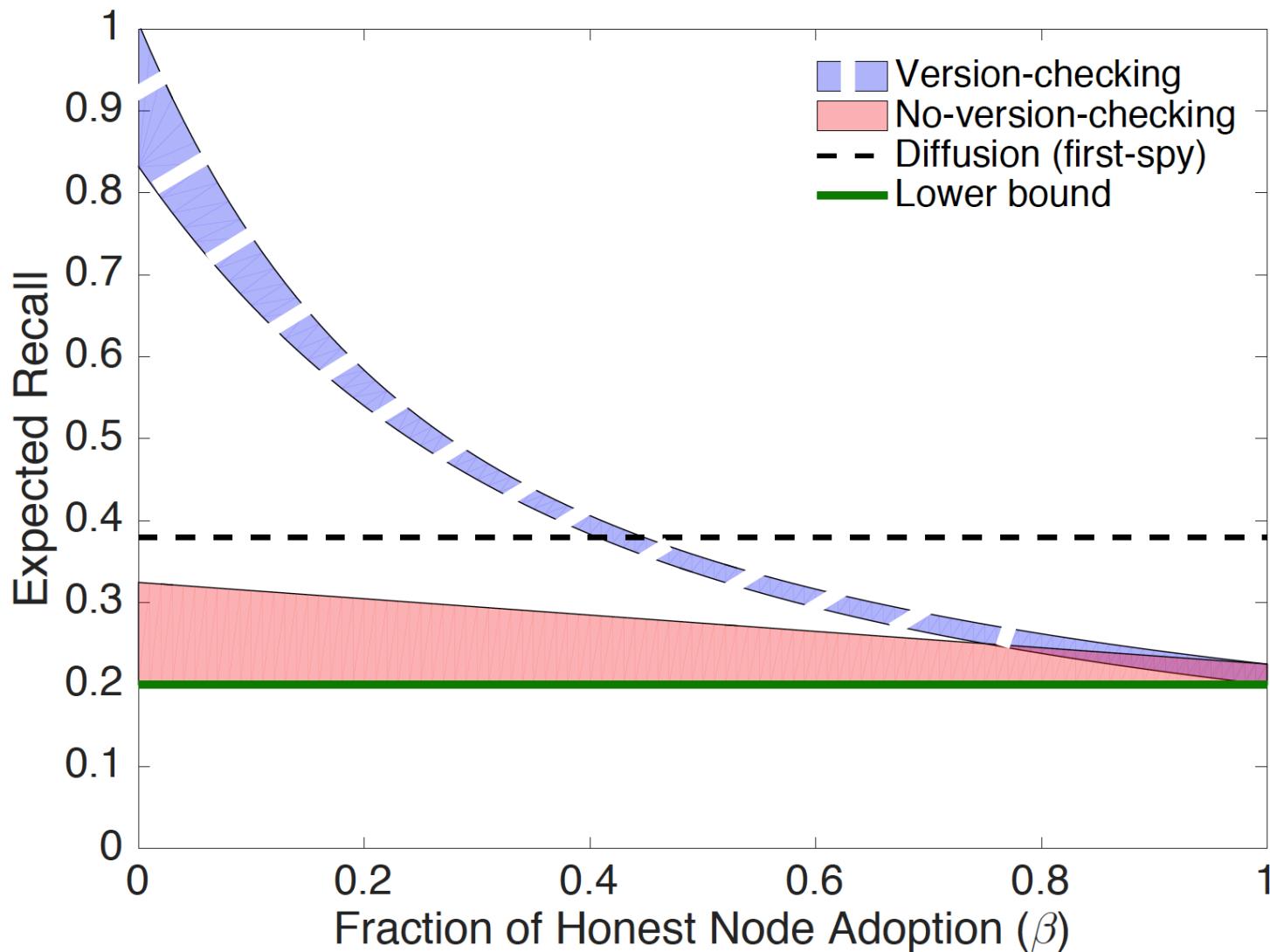
Latency Overhead: Estimate



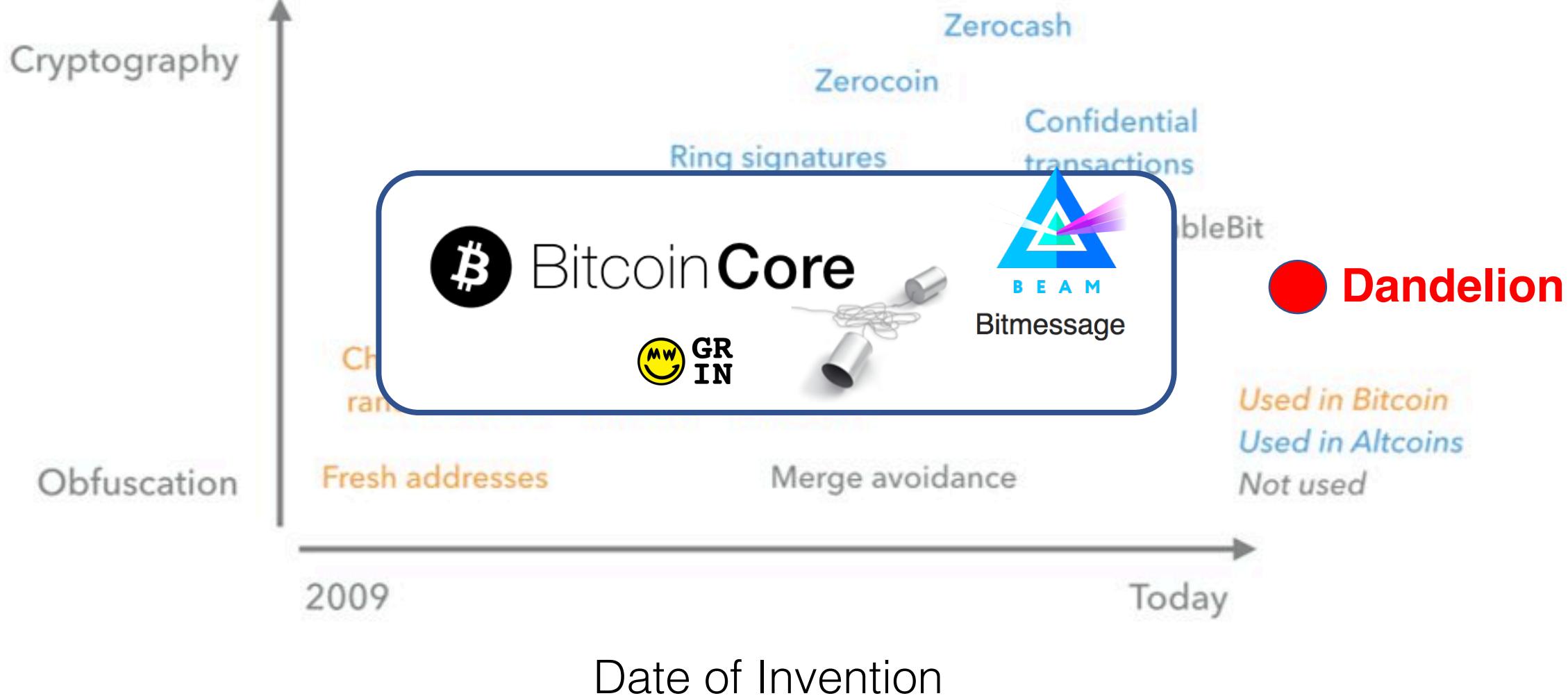
Empirical Delay Distribution



Practical Challenges: Partial deployment



Strength of
Guarantees



Take-Home Messages

- 1) Bitcoin's P2P network has poor anonymity.
- 2) Moving from trickle to diffusion did not help.
- 3) DANDELION may be a lightweight solution for certain classes of adversaries.

<https://github.com/dandelion-org/bitcoin>