Permutation polynomials over finite fields and their applications to Cryptography

Francesco Pappalardi

Beirut, March 9th, 2002

Finite Fields

Arr Let $\mathbb{F}_p = \mathbb{Z}/p\mathbb{Z} = \{0, 1, \dots, p-1\}$

(field if p prime);

$$|\mathbb{F}_p[x]/(f) = \{a_0 + a_1t + \dots + a_{m-1}t^{m-1} \mid a_i \in \mathbb{F}_p, \}$$

 $rightharpoonup \mathbb{F}_p[x]/(f)$ is a field

$$(g_1 \star g_2 \in \mathbb{F}_p[x]/(f) \text{ is } g_1g_2 \mod f)$$

 $\mathbb{F}_p[x]/(f)$ does not depend on f(i.e. if $h \in \mathbb{F}_p[x]$ irreducible, $\partial f = \partial h \Longrightarrow \mathbb{F}_p[x]/(f) \cong \mathbb{F}_p[x]/(h)$)

$$\mathbb{F}_{p^m} = \mathbb{F}_p[x]/(f)$$

any choice of f with $m = \partial f$ is the same

$$|\mathbb{F}_{p^m}| = p^m$$



$lacksquare{\mathbb{F}_q}$

Set
$$q = p^m$$

- Produce $\mathbb{F}_q \iff \text{find } f \in I_m(q);$ $(I_m(q) = \{ f \in \mathbb{F}_p[x], f \text{ irreducible}, \partial f = m \});$
- $\sum_{d|m} dI_d(q) = q^m;$
- $I_m(q) = \frac{q^m q}{m}$ (if m is prime) $I_m(q) \sim \frac{q^m}{m}$;
- $rightharpoonup If <math>m \nmid p-1 \& m \text{ is prime} \Longrightarrow \frac{x^m-1}{x-1} \in I_{m-1}(q);$
- Some fields: $\mathbb{F}_{2503} = \mathbb{F}_2[x]/(x^{503} + x^3 + 1), \mathbb{F}_{5323^{20}} = \mathbb{F}_{5323}[x]/(f)$ $f = x^{20} + 1451x^{18} + 5202x^{17} + 752x^{16} + 3778x^{15} + 4598x^{14} + 2563x^{13} + 5275x^{12} + 4260x^{11} + 862x^{10} + 4659x^9 + 3484x^8 + 1510x^7 + 4556x^6 + 2317x^5 + 2171x^4 + 3100x^3 + 4100x^2 + 682x + 5110$
- rightharpoonup Good to find <math>f sparse.





Interpolation on \mathbb{F}_q

Given $h: \mathbb{F}_q \to \mathbb{F}_q$ a function.

h can always be interpolated with a polynomial in $\mathbb{F}_q[x]$!

$$f_h(x) = \sum_{c \in \mathbb{F}_q} h(c) \prod_{\substack{d \in \mathbb{F}_q \\ d \neq c}} \frac{x - d}{c - d} \in \mathbb{F}_q[x]$$

FINITE FIELDS INTERPOLATION.

$$f_h(x) = \sum_{c \in \mathbb{F}_q} h(c) \left(1 - (x - c)^{q-1} \right) \in \mathbb{F}_q[x]$$

 \mathbb{F}_q^* is a (ciclic) group under multiplication

$$\Longrightarrow d^{q-1} = \begin{cases} 1 & d \neq 0 \\ 0 & d = 0. \end{cases}$$





More on interpolation in \mathbb{F}_q

- If $f_1, f_2 \in \mathbb{F}_q[x]$ with $f_1(c) = f_2(c) \forall c \in \mathbb{F}_q$, $\implies x^q x \mid f_1(x) f_2(x) \mid ;$
- The interpolant polynomial is unique mod $x^q x$ \implies unique with degree $\leq q 1$;
- If $c_h = \#\{c \in \mathbb{F}_q \mid h(c) \neq c\},\$ $q c_h \leq \partial f_h \leq q 2;$
- Problem. Find functions with sparse interpolation polynomial.





Permutation polynomials

$$\mathcal{S}(\mathbb{F}_q) = \{ \sigma : \mathbb{F}_q \to \mathbb{F}_q \mid \sigma(1:1) \}$$

permutations of \mathbb{F}_q .

- $f \in \mathbb{F}_q[x]$ is called permutation polynomial (PP) if "f (as a funtion) is a permutation"; (i.e. $\exists \sigma \in \mathcal{S}(\mathbb{F}_q), \sigma(c) = f(c) \ \forall c \in \mathbb{F}_q$)
- If $f_{\sigma}(x) = \sum_{c \in \mathbb{F}_q} \sigma(c) \left(1 (x c)^{q 1}\right) \in \mathbb{F}_q[x] \Longrightarrow$ $f \in \mathbb{F}_q[x] \text{ is PP} \iff \exists \sigma \in \mathcal{S}(\mathbb{F}_q), f \equiv f_{\sigma} \bmod x^q x.$
- **Examples:**
 - $ax + b, \qquad a, b \in \mathbb{F}_q, a \neq 0;$





More examples of PP

ightharpoonup Composition. $f \circ g$ is PP

if f, g are PP;

 $x^{(q+m-1)/m} + ax$ is a PP

if m|q-1;

ightharpoonup Linearized Polynomials. Let $q=p^m$,

$$L(x) = \sum_{s=0}^{r-1} \alpha_s x^{q^s} \qquad (\alpha_s \in \mathbb{F}_{p^m})$$

- $L(c_1 + c_2) = L(c_1) + L(c_2);$
- $L \in \mathrm{GL}_m(\mathbb{F}_p) \subset \mathcal{S}(\mathbb{F}_{p^m}) \iff \det(\alpha_{i-j}^{q^j}) \neq 0.$ $\iff L(x) = 0 \text{ has } 1 \text{ solution.}$





One more example of PP

ightharpoonup Dickson Polynomials. If $a \in \mathbb{F}_q$, $k \in \mathbb{N}$

$$D_k(x,a) = \sum_{j=0}^{\lfloor k/2 \rfloor} \frac{k}{k-j} {\binom{k-j}{j}} (-a)^j x^{k-2j}$$

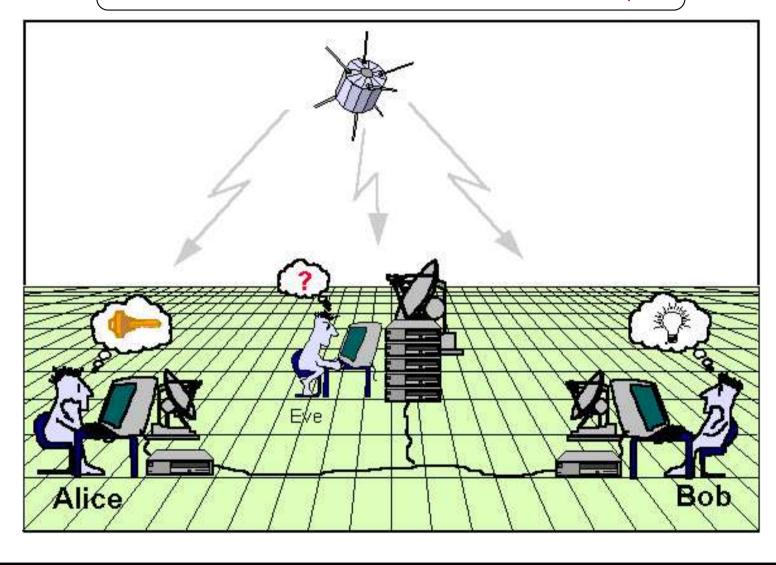
- if $a \neq 0$, $D_k(x, a)$ is a PP \iff $(k, q^2 1) = 1$;
- $D_k(x,0) = x^k \text{ is a PP} \iff (k,q-1) = 1.$
- Note: if $(mn, q^2 1) = 1$,

$$D_m(D_n(x,\pm 1),\pm 1) = D_{mn}(x,\pm 1)$$
.





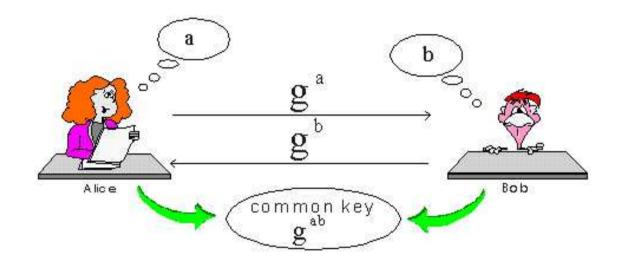
Diffie-Hellmann key exchange 1/2







Diffie-Hellmann key exchange 2/2



- **1** Alice and Bob agree on a finite field \mathbb{F}_q , and a generator $g \in \mathbb{F}_q$;
- **2** Alice picks a secret $a \in [0, q-1]$, Bob picks a secret $b \in [0, q-1]$;
- **3** They compute and publish g^a (Alice) and g^b (Bob);
- 4 The common secret key is g^{ab} .





Dickson analogue of DH Key-exchange

- ① Alice and Bob agree on a finite field \mathbb{F}_q , and a $\gamma \in \mathbb{F}_q$ $(\gamma \ not \ necessarily \ a \ generator);$
- ② Alice picks a secret $a \in [0, q^2 1]$, Bob picks a secret $b \in [0, q^2 1]$;
- 3 They compute and publish $D_a(\gamma, 1)$ (Alice) and $D_b(\gamma, 1)$ (Bob);
- 4 The common secret key is

$$(D_{ab}(\gamma, 1) = D_a(D_b(\gamma, 1, 1)) = D_b(D_a(\gamma, 1, 1))$$

NOTE. There is a fast algorithm to compute the value of a Dickson polynomial at an element of \mathbb{F}_q .

Problem. Find new classes of PP.





The problem of enumeration of PP by degree

$$N_d(q) = \{ \sigma \in \mathcal{S}(\mathbb{F}_q) \mid \partial(f_\sigma) = d \}$$

Problem. Compute $N_d(q)$

$$\sum_{d \le q-2} N_d(q) = q! \qquad (\partial f_\sigma \le q-2);$$

$$N_1(q) = q(q-1);$$

$$N_d(q) = 0 \text{ if } d|q-1$$
 (Hermite criterion);

- $\sim N_d(q)$ is known for $d \leq 6$;
- Almost all permutation polynomials have degree q-2.

(S. Konyagin, FP)
$$M_q = \{ \sigma \in \mathcal{S}(\mathbb{F}_q) \mid \partial f_\sigma < q - 2 \}$$

$$|\# M_q - (q-1)!| \leq \sqrt{2e/\pi} q^{q/2}$$





Other ways of counting

If
$$\sigma \in \mathcal{S}(\mathbb{F}_q)$$
,

$$c_{\sigma} = \#\{a \in \mathbb{F}_q \mid \sigma(a) \neq a\}$$

$$\sigma \neq id \Longrightarrow q - c_{\sigma} \leq \partial f_{\sigma} \leq q - 2$$

(since $f_{\sigma}(x) - x$ has at least $q - c_{\sigma}$ roots)

Consequences.

- \sim 2-cycles have degree q-2;
- 3-cycles have degree q-2 or q-3;
- \gg k-cycles have degree in [q-k, q-2].

(Wells)
$$\#\{\sigma \in 3\text{-cyle}, \ \partial(f_{\sigma}) = q - 3\} = \begin{cases} \frac{2}{3}q(q - 1) & q \equiv 1 \bmod 3\\ 0 & q \equiv 0 \bmod 3\\ \frac{1}{3}q(q - 1) & q \equiv 0 \bmod 3 \end{cases}$$





Beirut, March 9th 2002

More enumeration functions

- $\sigma_1, \sigma_2 \text{ conjugated} \Longrightarrow c_{\sigma_1} = c_{\sigma_2};$
- \sim C conjugation class of permutations;
- $c_{\mathcal{C}} = \#\{ \text{ elements } \in \mathbb{F}_q \text{ moved by any } \sigma \in \mathcal{C} \};$ (i.e. $c_{\mathcal{C}} = c_{\sigma} \text{ for any } \sigma \in \mathcal{C} \qquad q - c_{\mathcal{C}} \leq f_{\sigma})$
- $\mathcal{C} = [k] = k \text{cycles} \implies c_{[k]} = k.$
- Natural enumeration functions:

 $(minimal\ degree);$

 \mathcal{K} $M_{\mathcal{C}}(q) = \#\{\sigma \in \mathcal{C}, \partial f_{\sigma} < q - 2\}$

 $(non-maximal\ degree).$





Permutation Classes with non maximal degree

Let $C = (m_1, \ldots, m_t)$ be the class of permutations with m_1 1-cycles, ..., m_t t-cycles. The number c_C of elements in \mathbb{F}_q moved by any element of C is

$$c_{\mathcal{C}} = 2m_2 + 3m_3 + \dots + tm_t.$$

$$M_{\mathcal{C}}(q) = \#\{\sigma \in \mathcal{C}, \partial f_{\sigma} < q - 2\}$$

THEOREM 1 (C. Malvenuto, FP). $\exists N = N_{\mathcal{C}} \in \mathbb{N}, f_1, \dots, f_N \in \mathbb{Z}[x],$ $f_i \text{ monic}, \partial f_i = c_{\mathcal{C}} - 3 \text{ such that if } q \equiv a \text{ mod } N, \text{ then}$

$$M_{\mathcal{C}}(q) = \frac{q(q-1)}{m_2! 2^{m_2} \cdots m_t! t^{m_t}} f_a(q).$$





Consequences of Theorem 1

ightharpoonup If \mathcal{C} is fixed,

$$\operatorname{Prob}(\partial f_{\sigma} < q - 2 \mid \sigma \in \mathcal{C}) \sim \frac{1}{q};$$

ightharpoonup If $q=2^r$, \mathcal{C}_r is the conjugation class of r transposition,

$$M_{\mathcal{C}_r}(q) = \frac{q!}{r!2^r(q-2r+1)!} - \frac{q-2(r-1)(2r-1)}{2r} M_{\mathcal{C}_{r-1}}(q);$$

ightharpoonup One can compute $M_{\mathcal{C}}(q)$ for $c_{\mathcal{C}} \leq 6$.





Table 1. $\#c_{\mathcal{C}} \leq 6$, (q odd)

$$M_{[4]}(q) = \frac{1}{4}q(q-1)(q-5-2\eta(-1)-4\eta(-3))$$

$$M_{[5]}(q) = \frac{1}{5}q(q-1)\left(q^2 - (9-\eta(5)-5\eta(-1)+5\eta(-9))q + 26+5\eta(-7)+15\eta(-3)+15\eta(-1)\right)$$

$$M_{[2\ 3]}(q) = \frac{1}{6}q(q-1)\left(q^2 - (9+\eta(-3)+3\eta(-1))q + (24+6\eta(-3)+18\eta(-1)+6\eta(-7))\right) + \eta(-1)(1-\eta(9))q(q-5).$$





Table 2. $\#c_{\mathcal{C}} \leq 6$, (q even)

$$M_{[4]}(2^n) = \frac{1}{4}2^n(2^n-1)(2^n-4)(1+(-1)^n)$$

$$M_{[2\ 2]}(2^n) = \frac{1}{8}2^n(2^n-1)(2^n-2)$$

$$M_{[5]}(2^n) = \frac{1}{5}2^n(2^n-1)(2^n-3-(-1)^n)(2^n-6-3(-1)^n)$$

$$M_{[2\ 3]}(2^n) = \frac{1}{6}2^n(2^n-1)(2^n-3-(-1)^n)(2^n-6).$$





Table 3. $\#c_{\mathcal{C}} = 6$, $(q \text{ odd}, 3 \nmid q)$

$$\begin{split} M_{[6]}(q) &= \frac{q(q-1)}{6} \{q^3 - 14\,q^2 + [68 - 6\,\eta(5) - 6\,\eta(50)]q - \\ &= [154 + 66\,\eta(-3) + 93\,\eta(-1) + 12\,\eta(-2) + 54\,\eta(-7)]\} \\ M_{[4\ 2]}(q) &= \frac{q(q-1)}{8} (q^3 - [14 - \eta(2)]q^2 + \\ &= [71 + 12\,\eta(-1) + \eta(-2) + 4\,\eta(-3) - 8\,\eta(50)]q \\ &= -[148 + 100\,\eta(-1) + 24\,\eta(-2) + 44\,\eta(-3) + 40\,\eta(-7)]) \\ M_{[3\ 3]}(q) &= \frac{q(q-1)}{18} (q^3 - 13\,q^2 + [62 + 9\,\eta(-1) + 4\,\eta(-3)]q \\ &= -[150 + 99\,\eta(-1) + 42\,\eta(-3) + 72\,\eta(-7)]) \\ M_{[2\ 2\ 2]}(q) &= \frac{q(q-1)}{48} (q^3 - [14 + 3\,\eta(-1)]q^2 + [70 + 36\,\eta(-1) + 6\,\eta(-2)]q \\ &= -[136 + 120\,\eta(-1) + 48\,\eta(-2) + 8\,\eta(-3)]) \end{split}$$





Table 4. $\#c_{\mathcal{C}} = 6$

$$M_{[6]}(3^n) = \frac{3^n (3^n - 1)}{6} \{3^{3n} - [14 + 2(-1)^n] 3^{2n} + [71 + 39(-1)^n] 3^n - [162 + 147(-1)^n] \}$$

$$M_{[4\ 2]}(3^n) = \frac{3^n(3^n-1)}{8} \{3^{3n} - [14+3(-1)^n]3^{2n} + [72+40(-1)^n]3^n - [164+140(-1)^n]\}$$

$$M_{[3\ 3]}(3^n) = \frac{3^n(3^n-1)}{18} \{ (1+(-1)^n) 3^{3n} - [14+15(-1)^n] 3^{2n} + [71+81(-1)^n] 3^n - [150+171(-1)^n] \}$$

$$M_{[2\ 2\ 2]}(3^n) = \frac{3^n(3^n-1)}{48} \{3^{3n} - [14+3(-1)^n]3^{2n} + [76+36(-1)^n]3^n - [168+120(-1)^n]\}$$





Table 5. $\#c_{\mathcal{C}} = 6$)

$$M_{[6]}(2^n) = \frac{2^n (2^n - 1)}{6} \left\{ (2^n - 3 - (-1)^n)(2^{2n} - (11 - (-1)^n)2^n + (41 + 7(-1)^n)) \right\}$$

$$M_{[4\ 2]}(2^n) = \frac{2^n(2^n-1)}{8} \{(2^n-3-(-1)^n)(2^{2n}-11\cdot 2^n+37+(-1)^n)\}$$

$$M_{[3\ 3]}(2^n) = \frac{2^n(2^n-1)}{18} \{(2^n-3-(-1)^n)(2^{2n}-(10-(-1)^n)2^n+45-3(-1)^n))\}$$

$$M_{[2\ 2\ 2]}(2^n) = \frac{2^n(2^n-1)}{48} \{(2^n-2)(2^n-4)(2^n-8)\}.$$





k-cycles with minimal degree

$$m_{[k]}(q) = \#\{\sigma \text{ k-cycle}, \partial f_{\sigma} = q - k\}$$

THEOREM 2 (C. Malvenuto, FP).

 \blacktriangleleft If $q \equiv 1 \mod k \implies$

$$m_{[k]}(q) \ge \frac{\varphi(k)}{k} q(q-1).$$

 $rightharpoonup If <math>q = p^f, p \ge 2 \cdot 3^{[k/3]-1} \implies$

$$m_{[k]}(q) \le \frac{(k-1)!}{k} q(q-1).$$





Sketch of the Proof of Theorem 2. (1/3)

STEP 1. Translate the problem into one on counting points of an algebraic varieties;

$$m_k(q) = \frac{q(q-1)}{k} n_k(q)$$

where $n_k(q) = \{ \sigma \in [k] \mid \partial f_{\sigma} = q - k, \sigma(0) = 1 \}.$

Need to show $|n_k(q)| \leq (k-1)!$. Now

$$f_{\sigma}(x) = \sum_{c \in \mathbb{F}_q} \sigma(c) \left(1 - (x - c)^{q-1} \right) = A_1 x^{q-2} + A_2 x^{q-3} + \dots + A_{q-1}.$$

with
$$A_j = \sum_{c \in \mathbb{F}_q} \sigma(c) c^j = \sum_{c \in \mathbb{F}_q} \sigma(c) \left(c^j - c^{j-1} \right) = \sum_{\substack{c \in \mathbb{F}_q \\ \sigma(c) \neq c}} (\sigma(c) - c) c^j.$$





Sketch of the Proof of Theorem 2. (2/3)

If
$$\sigma = (0, 1, x_1, x_2, \dots, x_{k-2}) \in \mathcal{S}(\mathbb{F}_q),$$

$$A_j(\sigma) = (1 - x_1) + (x_1 - x_2)x_1^j + \dots + (x_{k-2} - x_{k-2})x_{k-3}^j + x_{k-2}^{j+1}.$$

Def. (Affine k-th Silvia set)

$$\mathcal{A}_{k} : \begin{cases}
(1 - x_{1}) + x_{1}(x_{1} - x_{2}) + \dots + x_{k-3}(x_{k-3} - x_{k-2}) + x_{k-2}^{2} &= 0 \\
(1 - x_{1}) + x_{1}^{2}(x_{1} - x_{2}) + \dots + x_{k-3}^{2}(x_{k-3} - x_{k-2}) + x_{k-2}^{3} &= 0 \\
\vdots & \vdots & \vdots \\
(1 - x_{1}) + x_{1}^{k-2}(x_{1} - x_{2}) + \dots + x_{k-3}^{k-2}(x_{k-3} - x_{k-2}) + x_{k-2}^{k-1} &= 0
\end{cases}$$

$$n_k(q) = \#\{\underline{x} = (x_1, \dots, x_{k-2}) \in \mathbb{F}_q^{k-2} \mid \underline{x} \in \mathcal{A}_k(\mathbb{F}_q), x_i \neq x_j\} \le \#\mathcal{A}_k(\mathbb{F}_q)$$

$$\dim_{\overline{\mathbb{F}}_q} \mathcal{A}_k = 0 \quad \stackrel{\text{Bezout Thm.}}{\Longrightarrow} \quad \# \mathcal{A}(\mathbb{F}_q) \le (k-1)!$$





Sketch of the Proof of Theorem 2. (3/3)

STEP 2.

Theorem. If K is an algebrically closed field,

char(**K**) =
$$\begin{cases} 0 & \text{or} \\ > 2 \cdot 3^{[k/3]-1}. \end{cases}$$

Then

$$\dim_{\mathbf{K}} \mathcal{A}_k = 0.$$

NOTE.

- \triangle Proof is based on finding projective hyperplanes disjoint from A_k ;
- \triangleq There are examples of small values of q with $\dim_{\mathbf{K}} A_k > 0$;





Numerical Examples (4–cycles)

$$m_{[4]}(\mathbb{F}_q) = \frac{1}{4}q(q-1) \cdot \begin{cases} 6 & \text{if } q \equiv 1 \pmod{20} \\ 4 & \text{if } q \equiv 11 \pmod{20} \\ 2 & \text{if } q \equiv 9,13,17 \pmod{20} \\ 0 & \text{if } q \equiv 3,7,19 \pmod{20}, \end{cases}$$

$$m_{[4]}(\mathbb{F}_{5^n}) = \frac{1}{2}5^n(5^n-1), \quad m_{[4]}(\mathbb{F}_{2^n}) = \begin{cases} 2^n(2^n-1) & \text{if } 4|n \\ 0 & \text{otherwise.} \end{cases}$$





Numerical Examples (5-cycles)

If
$$q \notin \{2, 13, 61, 3719, 3100067\} \Rightarrow m_{[5]}(\mathbb{F}_q) = \frac{q(q-1)}{5}(r_q + t_q + u_q),$$

$$t_q = \begin{cases} 4 & \text{if } q \equiv 1 \pmod{5} \\ 1 & \text{if } q \equiv 0 \pmod{5} \end{cases} \quad u_q = \begin{cases} -1 & \text{if } p = 11, 41 \\ 0 & \text{otherwise,} \end{cases} \quad r_q = \#\{ \begin{cases} \mathbb{F}_q - \text{roots} \\ \text{of } g_2 \end{cases} \}$$

$$g_2(x) = 2x^{20} - 29x^{19} + 229x^{18} - 1249x^{17} + 5187x^{16} - 17222x^{15} + 47040x^{14} - 107505x^{13} + 207622x^{12} - 340496x^{11} + 474638x^{10} - 560999x^9 + 559052x^8 - 465487x^7 + 319628x^6 - 177653x^5 + 77807x^4 - 25797x^3 + 6074x^2 - 904x + 64.$$





$$g_2(\alpha) = 0, \sigma_{\alpha} = (0, 1, \alpha, y(\alpha), z(\alpha)) \Rightarrow \partial f_{\sigma_{\alpha}} = q - 5 \text{ (minimal)}$$

```
y(x) = \frac{1}{(2)^3(13)(61)(3719)(3100067)} (6245340990732510 - 74275247020348477 x + 425897367479627411 x^2 - 1556772755104088477 x^3 + 4068122356423765520 x^4 - 8092377944341897339 x^5 + 12739155747072503154 x^6 - 16281608694400072277 x^7 + 17191467892889878476 x^8 - 15176855331347725064 x^9 + 11289210111615920188 x^{10} - 7103742513094855073 x^{11} + 3782081407301444460 x^{12} - 1696979431552752820 x^{13} + 635807089991226023 x^{14} - 195705738631474759 x^{15} + 48121368022605621 x^{16} - 9009616966592957 x^{17} + 1165803130533438 x^{18} - 82558295396232 x^{19}
```

```
 z(x) = \frac{1}{(2)^3(13)(61)(3719)(3100067)} -292290150269490 \, x^{19} + 3950333490943181 \, x^{18} \\ -29484664428617801 \, x^{17} + 152268243151302965 \, x^{16} - 599002775464475543 \, x^{15} \\ +1880438345917167218 \, x^{14} - 4841135989461751552 \, x^{13} + 10378374551469856881 \, x^{12} \\ -18679878403151115130 \, x^{11} + 28303942873286020848 \, x^{10} - 36041151267474587782 \, x^{9} \\ +38336702176933085823 \, x^{8} - 33711958096174593304 \, x^{7} + 24129466512539278343 \, x^{6} \\ -13742359416000756136 \, x^{5} + 6020424561116746133 \, x^{4} - 1925677501494324283 \, x^{3} \\ +413273185040891961 \, x^{2} - 51203861193252214 \, x + 2593061963570136)
```





Numerical Examples (6-cycles)

If
$$p \gg 1 \implies m_{[6]}(\mathbb{F}_p) = \frac{p(p-1)}{6}(s_1 + s_2 + s_3 + s_4)$$
 where
$$s_i = \# \left\{ \frac{\mathbb{F}_q - \text{roots}}{\text{of } f_i} \right\},$$

$$f_1(x) = x^2 - 3x + 3;$$

 $f_2(x) = x^4 - 3x^3 + 9x^2 - 9x + 3;$
 $f_3(x) = x^6 - 4x^5 + 12x^4 - 22x^3 + 25x^2 - 14x + 3;$
 $f_4(x) =$ Devil's Hat.





Galois Structure of the Silvia set

$$\operatorname{Gal}(\mathbb{Q}(f_1)/\mathbb{Q})) \cong \mathbb{Z}/2\mathbb{Z} \text{ (cyclotomic permutations)}$$

$$\operatorname{Gal}(\mathbb{Q}(f_2)/\mathbb{Q})) \cong D_4$$

$$\operatorname{Gal}(\mathbb{Q}(f_3)/\mathbb{Q})) \cong (\mathbb{Z}/3\mathbb{Z})^2 \rtimes S_2$$

 $Gal(\mathbb{Q}(Devil's\ Hat)) \cong ???$

(exponent probably = $(2)^5(3)^3(5)(7)(11)(13)(17)$)

Later discovered that

$$\operatorname{Gal}(\mathbb{Q}(\operatorname{Devil's} \operatorname{Hat})) \leq (\mathbb{Z}/6\mathbb{Z})^{18} \rtimes S_{18}$$



