

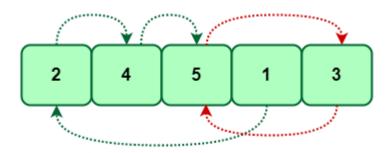
## Cycle Sort

Cycle sort is an in-place sorting Algorithm, <u>unstable sorting algorithm</u>, and a comparison sort that is theoretically optimal in terms of the total number of writes to the original array.

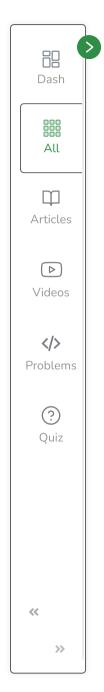


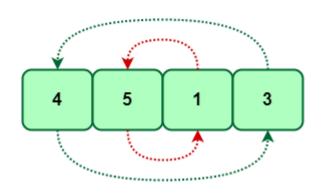
- It is optimal in terms of the number of memory writes. It <u>minimizes the number of memory writes</u> to sort (Each value is either written zero times if it's already in its correct position or written one time to its correct position.)
- It is based on the idea that the array to be sorted can be divided into cycles. Cycles can be visualized as a graph. We have n nodes and an edge directed from node i to node j if the element at i-th index must be present at j-th index in the sorted array.

Cycle in 
$$arr[] = \{2, 4, 5, 1, 3\}$$



Cycle in  $arr[] = \{4, 3, 2, 1\}$ 







## **Implementation**

```
C++ Java

// Java program to implement cycle sort

import java.util.*;
import java.lang.*;

class GFG {
    // Function sort the array using Cycle sort
    public static void cycleSort(int arr[], int n)
    {
        // count number of memory writes
        int writes = 0;
}
```

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Problems
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  Quiz
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```

```
// traverse array elements and put it to on
// the right place
for (int cycle start = 0; cycle start <= n - 2; cycle start++) {</pre>
    // initialize item as starting point
    int item = arr[cycle start];
    // Find position where we put the item. We basically
    // count all smaller elements on right side of item.
    int pos = cycle start;
    for (int i = cycle start + 1; i < n; i++)</pre>
        if (arr[i] < item)</pre>
            pos++;
    // If item is already in correct position
    if (pos == cycle_start)
        continue;
    // ignore all duplicate elements
    while (item == arr[pos])
        pos += 1;
    // put the item to it's right position
    if (pos != cycle start) {
        int temp = item;
        item = arr[pos];
        arr[pos] = temp;
        writes++;
```



```
// Rotate rest of the cycle
                              while (pos != cycle start) {
                                  pos = cycle start;
  // Find position where we put the element
 Dash
                                  for (int i = cycle start + 1; i < n; i++)</pre>
                                      if (arr[i] < item)</pre>
  pos += 1;
  Αll
                                  // ignore all duplicate elements
  \Box
                                  while (item == arr[pos])
Articles
                                      pos += 1;
 // put the item to it's right position
Videos
                                  if (item != arr[pos]) {
                                      int temp = item;
  </>>
                                      item = arr[pos];
Problems
                                      arr[pos] = temp;
              Tutorials
 Courses
                           Jobs
                                      Practice
                                                  Contests
                      // Driver program to test above function
                      public static void main(String[] args)
<<
                          int arr[] = \{ 1, 8, 3, 9, 10, 10, 2, 4 \};
                          int n = arr.length;
   >>
```







## **Time Complexity Analysis:**

- Worst Case: O(n<sup>2</sup>)
- Average Case: O(n<sup>2</sup>)
- Best Case: O(n<sup>2</sup>)

## **Auxiliary Space:** O(1)

• The space complexity is constant cause this algorithm is in place so it does not use any extra memory to sort.

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