

Solution: The diameter of a tree T is the largest of the following quantities:

- The diameter of T's left subtree.
- The diameter of T's right subtree.
- The longest path between leaves that goes through the root of T (this can be computed from the heights of the subtrees of T).

The longest path between leaves that goes through a particular node say,

nd

can be calculated as:

$$1 + \text{height of left subtree of } nd + \text{height of right subtree of } nd$$

Therefore, final **Diameter** of a node can be calculated as:

$$\text{Diameter} = \text{maximum}(\text{lDiameter}, \text{rDiameter}, 1 + \text{lHeight} + \text{rHeight})$$

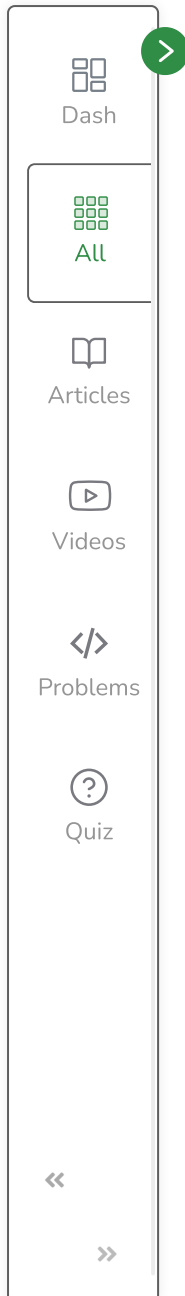
Where,

lDiameter = Diameter of left subtree

rDiameter = Diameter of right subtree

lHeight = Height of left subtree

rHeight = Height of right subtree



Implementation :

C++

Java

```
// Recursive optimized Java program to find the diameter of a
// Binary Tree

/* Class containing left and right child of current
node and key value*/
class Node
{
    int data;
    Node left, right;

    public Node(int item)
    {
        data = item;
        left = right = null;
    }
}

/* Class to print the Diameter */
class BinaryTree
{
    Node root;

    /* Method to calculate the diameter and return it to main */
    int diameter(Node root)
    {
        /* base case if tree is empty */
        if (root == null)
            return 0;

        /* get the height of left and right sub trees */
        int lheight = height(root.left);
        int rheight = height(root.right);
```



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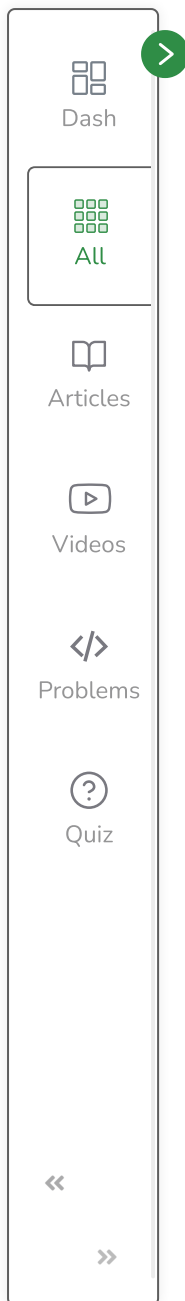


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```

/* get the diameter of left and right subtrees */
int ldiameter = diameter(root.left);
int rdiameter = diameter(root.right);

/* Return max of following three
  1) Diameter of left subtree
  2) Diameter of right subtree
  3) Height of left subtree + height of right subtree + 1 */
return Math.max(lheight + rheight + 1,
                Math.max(ldiameter, rdiameter));
}

/* A wrapper over diameter(Node root) */
int diameter()
{
    return diameter(root);
}

/*The function Compute the "height" of a tree. Height is the
  number of nodes along the longest path from the root node
  down to the farthest leaf node.*/
static int height(Node node)
{
    /* base case tree is empty */
    if (node == null)
        return 0;

    /* If tree is not empty then height = 1 + max of left
       height and right heights */
    return (1 + Math.max(height(node.left), height(node.right)));
}

public static void main(String args[])
{
    /* creating a binary tree and entering the nodes */
    BinaryTree tree = new BinaryTree();
    tree.root = new Node(1);
    tree.root.left = new Node(2);
    tree.root.right = new Node(3);
    tree.root.left.left = new Node(4);

```



```

tree.root.left.right = new Node(5);

System.out.println("The diameter of given binary tree is : "
    + tree.diameter());
    }
}

```

Output:

Diameter of the given binary tree is 4

Time Complexity: $O(N^2)$, where N is the number of nodes in the binary tree.

Auxiliary Space: $O(N)$ for call stack

Efficient Approach: To solve the problem follow the below idea:

The above implementation can be optimized by calculating the height in the same recursion rather than calling a height() separately.

Below is the implementation of the above approach:

C++ Java Python3 C# Javascript

```

// Recursive Java program to find the diameter of a
// Binary Tree

// Class containing left and right child of current

```



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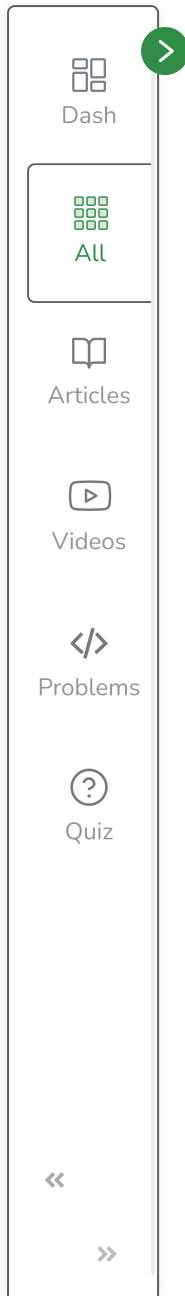


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```
// node and key value
class Node {
    int data;
    Node left, right;

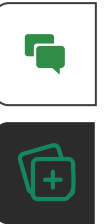
    public Node(int item)
    {
        data = item;
        left = right = null;
    }
}

// A utility class to pass height object
class Height {
    int h;
}

// Class to print the Diameter
class BinaryTree {
    Node root;

    // define height =0 globally and call
    // diameterOpt(root,height) from main
    int diameterOpt(Node root, Height height)
    {
        // lh --> Height of left subtree
        // rh --> Height of right subtree
        Height lh = new Height(), rh = new Height();

        if (root == null) {
            height.h = 0;
            return 0; // diameter is also 0
        }
        /*
            ldiameter --> diameter of left subtree
            rdiameter --> Diameter of right subtree
            Get the heights of left and right subtrees
            in lh and rh. And store the returned values in
            ldiameter and ldiameter*/
        int ldiameter = diameterOpt(root.left, lh);
        int rdiameter = diameterOpt(root.right, rh);
    }
}
```





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```
// Height of current node is max of heights of left
// and right subtrees plus 1
height.h = Math.max(lh.h, rh.h) + 1;
```

```
return Math.max(lh.h + rh.h + 1,
                Math.max(ldiameter, rdiameter));
```

}

```
// A wrapper over diameter(Node root)
int diameter()
```

{

```
Height height = new Height();
```

```
return diameterOpt(root, height);
```

}

```
// The function Compute the "height" of a tree. Height
// is
// the number of nodes along the longest path from the
// root node down to the farthest leaf node.
```

```
static int height(Node node)
```

{

```
// base case tree is empty
```

```
if (node == null)
```

```
return 0;
```

```
// If tree is not empty then height = 1 + max of
```

```
// left height and right heights
```

```
return (1
```

,

```
// Driver Code
```

```
public static void main(String args[])
```

{

```
// creating a binary tree and entering the nodes
```

```
BinaryTree tree = new BinaryTree();
```

```
tree.root = new Node(1);
```

```
tree.root.left = new Node(2);
```

```
tree.root.right = new Node(3);
```



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```
tree.root.left.left = new Node(4);
tree.root.left.right = new Node(5);

// Function Call
System.out.println(
    "The diameter of given binary tree is : "
    + tree.diameter());
}
```

Output

Diameter of the given binary tree is 4

Time Complexity: $O(N)$

Auxiliary Space: $O(N)$ due to recursive calls.

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