## CMO

## Sheet 2 — E0230

## Assignment (Due: 14 December 2020)

## Instructions

- Answer all questions
- For numerical questions, answer without any leading or trailing spaces
- You can submit MS Teams Form only once
- If the value of answer is an **integer**, give answer in box provided as an integer alone **without any decimal point**. For example 0, 10
- Otherwise, if the value of answer is a **decimal number**, give answer in x.yy format **rounded to 2 decimal places**. For example 0.50, 1.50
- Follow this instruction unless otherwise specified in the question
- Late submissions will be penalised

1. Quadratic minimization. Consider the following minimization problem

$$5x^2 + 5y^2 - xy - 11x + 11y + 11$$

- (a) Let  $[\hat{x}, \hat{y}]$  be a point satisfying the first order necessary conditions for a solution.
  - (i) (1 point)  $\hat{x} =$ \_\_\_\_\_
  - (ii) (1 point)  $\hat{y} =$ \_\_\_\_\_
- (b) (1 point) Is this point is a global minimum
- (c) (1 point) What would be the rate of convergence of steepest descent for this problem
- (d) (1 point) Starting at x=y=0, how many steepest descent iterations would it take (at-most) to reduce the function value to  $10^{-11}$  (in integer format)
- 2. Quadratic minimization. Let  $f: \mathbb{R}^4 \longrightarrow \mathbb{R}$  be twice continuously differentiable function. Assume that the hessian of f is positive definite and the largest absolute eigen value of the Hessian matrix of f at all points is bounded above by 25. We are given that  $x_0 = [4,0,-2,1]^T$ ,  $f(x_0) = 6$ , and  $\nabla f(x_0) = [8,4,4,2]^T$ . Using 2nd order Taylor series, find a quadratic function  $g: \mathbb{R}^4 \longrightarrow \mathbb{R}$  such that  $g(x_0) = f(x_0)$  and  $f(x) \leq g(x), \forall x \in \mathbb{R}^4$ . What is the minimum value of g (rounded to the nearest integer) (5 points)?
- 3. Constant step-size. Consider the function

$$f(x) = 3(x_1^2 + x_2^2) + 4x_1x_2 + 5x_1 + 6x_2 + 7$$

where  $x = [x_1, x_2]^T \in \mathbb{R}^2$ . Suppose we use a fixed step-size gradient descent to find minimizer of f.

$$x^{k+1} = x^k - \alpha \nabla f(x^{(k)})$$

Let  $a < \alpha < b$  be the largest range of values of  $\alpha$  for which the algorithm is globally convergent. Find a,b rounded to 2 decimal places in x.yy format.

- (a) (2 points) a =\_\_\_\_\_
- (b) (3 points)  $b = ___$
- 4. Constant step-size. Consider the function  $f: \mathbb{R}^2 \longrightarrow \mathbb{R}$  given by

$$f(x) = \frac{3}{2}(x_1^2 + x_2^2) + (1+a)x_1x_2 - (x_1 + x_2) + b$$

where a and b are some unknown real valued parameters.

(a) (2 points) Find the largest value of a (rounded to the nearest integer) for which unique global minimizer of f exists.

(b) (3 points) Consider the following algorithm

$$x^{(k+1)} = x^k - \frac{2}{5}\nabla f(x^{(k)})$$

Find largest value of a (rounded to the nearest integer) for which the above algorithm converges to the global minimizer of f for any initial point  $x^{(0)}$ .

- 5. Steepest descent. Write a subroutine (in Python) for implementing steepest descent using exact line search.
  - (a) Use the given function f5.pkl, its gradient grad\_f5.pkl and hessian hess\_f5.pkl. The function in f5.pkl takes an argument  $x \in \mathbb{R}^2$  as a python list of size 2 and returns a real number f(x). Similarly, the function in grad\_f5.pkl takes an argument x as a python list and returns  $\nabla f(x)$  as a python list. The function in hess\_f5.pkl does not take an argument and returns hessian matrix of f(x). The functions can be loaded as dill.loads(pickle.load(file\_pointer)), where dill and pickle are python libraries and file\_pointer points to the pickle file to be read. Now, test your subroutine on f(x) using initial condition  $[0, 10]^T$ . For the stopping criterion, use  $||g^{(k)}||_2 \le \epsilon$  where  $\epsilon = 10^{-6}$  and  $||.||_2$  is 2-norm.
    - (i) (2.5 points) Determine number of iterations required to satisfy above stopping condition?
    - (ii) (2.5 points) What is value of objective function f(x) at final point?
  - (b) Now, test the subroutine for following quadratic problem  $f(x) = \frac{1}{2}x^T A x b^T x, x \in \mathbb{R}^4$ . For the stopping criterion, use  $||g^{(k)}||_2 \le \epsilon$  where  $\epsilon = 10^{-6}$  and  $||.||_2$  is 2-norm.

with 
$$A = \begin{pmatrix} 0.78 & -0.02 & -0.12 & -0.14 \\ -0.02 & 0.86 & -0.04 & 0.06 \\ -0.12 & -0.04 & 0.72 & -0.08 \\ -0.14 & 0.06 & -0.08 & 0.74 \end{pmatrix}, b = \begin{pmatrix} 0.76 \\ 0.08 \\ 1.12 \\ 0.68 \end{pmatrix}, x_0 = 0$$

- (i) (2.5 points) Determine number of iterations required to satisfy above stopping condition?
- (ii) (2.5 points) What is value of objective function f(x) at final point?

Give answers for (ii) rounded to 2 decimal places in x.yy format.

6. In-exact line search. Backtracking is a form of inexact line search in which a step size is determined at each step which satisfies the Armio-Goldstein condition. Given constants  $\alpha, \beta \in (0,1)$ , at each step of the algorithm, if the current point is  $x \in \mathbb{R}^d$ , the direction of line search is chosen as  $u = -\nabla f(x)$ , and for determining the step size, an initial step size t = 1 is chosen and is repeatedly updated as  $t \leftarrow \beta t$  until  $f(x+tu) \leq f(x) + \alpha t \nabla f(x)^T u$  and then x is updated as  $x \leftarrow x + tu$ . Once the update distance  $||tu||_2$  for the point x

becomes less than  $\epsilon$  during any epoch, the algorithm is stopped.

Use the given function f6.pkl and its gradient grad\_f6.pkl. The function in f6.pkl takes an argument  $x \in \mathbb{R}^4$  as a python list of size 4 and returns a real number f(x). Similarly, the function in grad\_f6.pkl takes an argument x as a python list and returns  $\nabla f(x)$  as a python list. The functions can be loaded as

dill.loads(pickle.load(file\_pointer)), where dill and pickle are python libraries and file\_pointer points to the pickle file to be read.

Apply backtracking line search algorithm with initial point [10, 100, 100, 10],  $\alpha = 0.5$ ,  $\beta = 0.5$  and  $\epsilon = 10^{-7}$ .

- (a) Let  $[x_1, x_2, x_3, x_4]$  the solution vector with each element rounded to 2 decimal places in x.yy format.
  - (i) (1 point)  $x_1 =$ \_\_\_\_\_
  - (ii) (1 point)  $x_2 =$ \_\_\_\_\_
  - (iii) (1 point)  $x_3 =$ \_\_\_\_\_
  - (iv) (1 point)  $x_4 =$ \_\_\_\_\_
- (b) (3 points) Give the number of iterations it took to obtain the result.
- (c) (i) (1.5 points) Give the least number of function calls to f6 to obtain the result.
  - (ii) (1.5 points) Give the least number of function calls to grad\_f6 used to obtain the result.
- 7. One-dimensional search methods. Consider the one-dimensional minimization problem

$$\min_{x \in [a,b]} \quad f(x) \tag{1}$$

We are given a, b and a subroutine 'foo'. The subroutine 'foo' returns the value function f(x) for any  $x \in [a, b]$ . Let  $x^*$  be the unique minimizer for this problem. Given a tolerance value  $\epsilon > 0$ , we are interested in finding  $\hat{x} \in [a, b]$  such that  $|\hat{x} - x^*| \le \epsilon$ . Here is a pseudo-code to find  $\hat{x}$ .

- Initialize:  $x_l = a, x_u = b$
- In loop:

$$-d = (x_u - x_l) * \rho, 0 < \rho < 1$$

$$-x_{-}=x_{u}-d, x_{+}=x_{l}+d$$

- if  $f(x_{-}) < f(x_{+})$  then  $x_{u} = x_{+}$  otherwise  $x_{l} = x_{-}$
- Output:  $0.5(x_l + x_u)$ , tolerance =  $0.5(x_u x_l)$ , NSC = number of times the subroutine 'foo' was called.

Using the generic pseudo-code described above, we want you to implement two line search techniques for uni-modal functions. First is Golden section search (GS) where  $\rho = \frac{\lambda}{(1+\lambda)}$  where  $\lambda = 0.5(1+\sqrt{5})$ . Now implement Golden section search as a function (in Python). Note that the loop part of GS continues until tolerance  $\leq \epsilon$  is satisfied.

Second is Fibonacci search. In FS,  $k^{th}$  iteration uses  $\rho = \frac{F_{N-k}}{F_{N-k+1}}$ , where  $F_0 = 1$ ,  $F_1 = 1$ ,  $F_n = F_{n-1} + F_{n-2} \ge 2$  is the Fibonacci sequence. Note that in FS, the loop will be executed N-1 times.

Implement the 'foo' subroutine such that 'foo(x)' returns  $e^{-x} - cos(x)$  for  $x \in [0, 1]$ . First call FS subroutine for N = 20, N = 10. Then call GS subroutine with  $\epsilon$  returned by tolerance value of FS. Give answers for  $\hat{x}$  and tolerance rounded to 3 decimal places in the form **x.yyy**.

- (a) Run FS subroutine for N=20 and report the following values
  - (i) (1 point)  $\hat{x} =$ \_\_\_\_\_
  - (ii) (1 point)  $tolerance = \underline{\hspace{1cm}}$
  - (iii) (0.5 points) NSC =
- (b) Run GS subroutine using  $\epsilon$  calculated in (a) and report the following values
  - (i) (1 point)  $\hat{x} =$ \_\_\_\_\_
  - (ii) (1 point)  $tolerance = \underline{\hspace{1cm}}$
  - (iii) (0.5 points) NSC =
- (c) Run FS subroutine for N=10 and report the following values
  - (i) (1 point)  $\hat{x} =$ \_\_\_\_\_
  - (ii) (1 point) *tolerance* = \_\_\_\_\_
  - (iii) (0.5 points) NSC =
- (d) Run GS subroutine using  $\epsilon$  calculated in (c) and report the following values
  - (i) (1 point)  $\hat{x} =$ \_\_\_\_\_
  - (ii) (1 point)  $tolerance = \underline{\hspace{1cm}}$
  - (iii) (0.5 points) NSC =