# Unit Exam: Persistence

Total points 140/215

Take this exam alone. It is closed book, closed notes, but feel free to use a single 8.5" x 11" sheet of paper (both sides) of notes. Take the exam only once.

The respondent's email (parth2@pdx.edu) was recorded on submission of this form.

✓ OS developers generally assume that I/O is slow *	5/5
True	<b>✓</b>
☐ False	

In our simplified model of I/O devices, what types of registers are contained by a canonical I/O device? \*

	Yes, it is part of a canonical I/O device	No, not part of a canonical I/O device	Score	
data register		0	1/1	<b>✓</b>
DMA register	$\bigcirc$		1/1	<b>~</b>
command register	•	0	1/1	<b>✓</b>
floating point register	0	•	1/1	<b>✓</b>
status register		$\bigcirc$	1/1	<b>~</b>
stack pointer	0		1/1	<b>~</b>

<b>✓</b>	Which statement best characterizes "Interrupt-driven I/O"? * 5/5
0	the CPU gets an asynchronous notification after each individual I/O operation completes
0	the I/O device can handle all of the data transfer itself and only notify/interrupt the CPU when the complete sequence of transfers is complete.
	the CPU can do other work and not block while I/O operations complete
0	the CPU synchronously interacts with the I/O device and controls the timing of each I/O operation
<b>✓</b>	Which statement best characterizes DMA (direct memory access) I/O"? * 5/5
•	the I/O device can handle all of the data transfer itself and only notify/interrupt the CPU when the complete sequence of transfers is complete.
0	the CPU synchronously interacts with the I/O device and controls the timing of each I/O operation
0	the CPU can do other work and not block while I/O operations complete
0	the CPU gets an asynchronous notification after each individual I/O operation completes

	Programmed I/O	Interrupt- driven I/O	DMA	Score	
when the I/O device is capable of transferring large amounts of data	0	0	•	1/1	<b>✓</b>
when the I/O device is very slow	0	•	0	1/1	<b>✓</b>
when interrupts are rare	0	0	•	1/1	<b>~</b>
when the excessive I/O interrupts might slow down the system		0	0	1/1	<b>✓</b>
when the I/O device is trusted to make independent memory reads and writes	0	0		1/1	<b>✓</b>
when the I/O device has very fast bursts	•	$\circ$	0	1/1	<b>✓</b>

	reason for OS	not a reason for		
	developer to dislike device drivers	OS developers to dislike device drivers	Score	
there are so many of them	0		0/1	×
Device Drivers make it easy to integrate new I/O devices with existing OSs	0		1/1	<b>✓</b>
Device Drivers are a common source of crashes		0	1/1	<b>✓</b>
device drivers are written by many different developers, many of whom are not OS kernel developers		0	1/1	<b>✓</b>
Device Drivers can handle both character devices (like mouse and keyboard) and block devices (like HDDs and SSDs)	0		1/1	<b>✓</b>
Correct answers	rea	ason for OS develope dislike device drivers	r to developers	eason for OS to dislike device drivers

Match each advantage with each type of storage device *					
	SSDs	HDDs	Score		
faster	•	$\circ$	1/1	<b>~</b>	
longer lasting data integrity	$\circ$		1/1	<b>~</b>	
quieter		$\bigcirc$	1/1	<b>✓</b>	
less expensive (per byte stored)	$\circ$	•	1/1	<b>~</b>	
random reads/writes are just as fast as sequential reads/writes		0	1/1	<b>✓</b>	
no moving parts		0	1/1	<b>~</b>	
larger total capacity	0	•	1/1	<b>✓</b>	

Rank each item from smallest to biggest (as they relate to flash chips) *					
	smallest	bigger	biggest	Score	
plane	0	$\circ$		2/2	<b>✓</b>
block	0	•	$\circ$	2/2	<b>✓</b>
page		0	0	2/2	<b>✓</b>

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✓ Data written to a flash chip will tend to degrade if the chip is left without 5/5 power for 1-2 years. *
True
False
Feedback
the class lecture said that this is True, but the actual research seems to be inconclusive and contradictory. All seem to agree that if flash-based storage devices are stored at low temperatures then they retain data much, much longer.

What must we do before we update (program) a page of a flash chip (assuming the page's contents are currently/previously valid)? *					
	Yes, must be done prior to updating a page of a flash chip	No, not necessarily done prior to updating a page of a flash chip	Score		
sacrifice a chicken at dawn	0		1/1	<b>✓</b>	
copy the block containing the to-be-updated page to a temporary location			1/1	<b>✓</b>	
erase the entire block containing the page			1/1	<b>✓</b>	
copy the page's data to a temporary location			1/1	<b>✓</b>	
erase the page	•	0	0/1	×	
Correct answers	Yes	s, must be done prior	to No, not n	ecessarily done	
		lating a page of a fla chip	sh prior to upo	lating a page of a ash chip	
erase the page		0		•	

X An SSD will typically keep logical blocks located in same physical flash chip blocks to optimize for spatial locality and wear leveling. *	0/5
<ul><li>True</li><li>False</li></ul>	×
Correct answer  False	
Feedback  False. SSDs will tend to move a logical block to varying flash memory locations each ti it is updated. moving logical disk blocks around to varying physical locations achieves wear-leveling.	me
SSDs tend to present the exact same client interface as is used for Hai Disk Drives *	<sup>r</sup> d 5/5
True. this is done for backward compatibility and interoperability	<b>✓</b>
False. SSDs present a flash-aware programming interface that is optimized for performance and reliability of flash-based devices	

~	An SSD's flash controller (specifically the flash translation layer) will 5/5 commonly program pages within an erased block in order, from low page to high page to achieve what goal? *
0	wear leveling
0	write spreading
	minimize disturbance
0	improve performance
0	prevent data corruption

Why does an SSD include volatile memory (in addition to its persistent flash storage chips)? *				
	yes, a reason for including volatile memory on an SSD	no, not a reason	Score	
improve performance	•	0	1/1	<b>✓</b>
reduced cost	0		1/1	<b>✓</b>
provide a temporary storage for page updates			1/1	<b>✓</b>
caching of persistent data		0	1/1	<b>✓</b>
persistent storage of data	0	•	1/1	<b>✓</b>
reduce disturbance	•	0	0/1	×
mapping tables for logical/physical translation	0		0/1	×
Correct answers				
		, a reason for includ Itile memory on an		eason for volatile memory
reduce disturbance		$\circ$		•
mapping tables for logical/physical tra				0

	ading writes across the blo	ocks of the flash chips as evenly as *	5/5
wear	leveling		<b>✓</b>
) flash	dancing		
O repro	gramming		
write	spreading		

Match each OS concept with the physical entity that it virtualizes *							
	CPU	Physical Memory	Persistent Storage Device	Score			
File System	0	0		2/2	<b>✓</b>		
Virtual Memory	0	•	0	2/2	<b>✓</b>		
Process		0	0	2/2	<b>✓</b>		

Match each linux system call with its purpose. *								
	unlink()	open()	lseek()	rename()	stat()	link()	Score	
create a new location for a file without changing old name	0	0	0		0	0	0/2	×
move a file to a different location	0	0	•	0	0	0	0/2	×
jump to a new location within a file	0	0	0	0	0	•	0/2	×
open a file for reading or writing	0	0	0	0	0	•	0/2	×
remove a file	•	0	$\bigcirc$	0	0	0	2/2	<b>✓</b>
get a file's metadata	$\circ$	0	$\circ$	0	•	0	2/2	<b>✓</b>
create a file	0	•	$\bigcirc$	0	0	0	2/2	<b>✓</b>
Correct answers								
	unlin	k()	open()	lseek()	renam	e()	stat()	link()
create a new location for a file	C	)	0	0	0		0	•

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without changing old name							
move a file to a different location	0	0	0	•	0	0	
jump to a new location within a file	0	0	•	0	0	0	
open a file for reading or writing	0	•	0	0	0	0	

	Yes, usually found in a linux file's metadata	No, not a common metadata item for linux files	Score	
date of deletion		$\circ$	0/1	×
access control permissions	•	0	1/1	<b>~</b>
revision history		$\bigcirc$	0/1	×
date of creation		0	2/2	<b>✓</b>
reference to parent	•	$\bigcirc$	1/1	<b>~</b>
name		$\bigcirc$	1/1	<b>✓</b>
owner		0	1/1	<b>✓</b>
favorite Beatles tune	0		1/1	<b>✓</b>
mother's maiden name	0		1/1	<b>✓</b>
time of last update	•	0	1/1	<b>✓</b>
list of update times	•	0	0/1	×
favorite ice cream flavor	0		1/1	<b>✓</b>
Correct answers				
	Yes,	usually found in a linux file's metadata		ommon metadata or linux files

revision history	$\circ$	
list of update times	0	•
✓ In Linux, many items (proce devices, etc.) can be represented.		
<ul><li>True</li><li>False</li></ul>		<b>✓</b>
Linux directories are impler contents of the directory. *		ntaining a list of the 0/5
<ul><li>True</li><li>False</li></ul>		×
Correct answer		
False		
Feedback  False. This was true in the early day directories are implemented in various size and scalability. Typically, each in a unique way while conforming to	ous other ways to optimi different type of file syst	ize for performance, reliability, tem implements its directories

✓ Why does Linux need the fsync() system call? * 5/5	;
To increase performance	
To decrease wear out of SSD flash chips	
To force recently updated file data to be written to persistent storage.	
To implement looser consistency guarantees	
The linux system call allows you to remove a storage device's 0/5 files and directories from the system's directory structure. *	;
omount()	
link()	
unmount()	
unlink()	
Correct answer	
unmount()	

X The linux system call allows you to remove a file. *	0/5
unmount()	
link()	
omount()	
unlink()	
rm()	×
Correct answer	
unlink()	
If we use a free list to keep track of free blocks, then how big is the f list? *	ree 0/5
the free list will have one element (one reference) per storage block	×
it depends on how much of the disk storage is utilized or free	
Correct answer	
it depends on how much of the disk storage is utilized or free	

×	When keeping track of all of the allocated storage blocks for a given file, one rarely used strategy is to use contiguous allocation. Why is this technique rarely used? *	0/5
0	External Fragmentation	
	Internal Fragmentation	×
Corr	ect answer	
	External Fragmentation	
×	Identify one significant disadvantage of using a linked list for file block allocation tracking *	0/5
$\bigcirc$	Simple to implement	
$\bigcirc$	for HDDs, random access into the file is slow	
$\bigcirc$	difficult to grow/shrink the file after it is created	
0	external fragmentation	
	when used with SSDs linkked-list allocation causes bit disturbance across adjacent flash memory pages	×
Corr	ect answer	
•	for HDDs, random access into the file is slow	
Cons	Allocation Table ider the following (small) file allocation table for a FAT filesystem. A "Cluster" is a contiguous grants by the blocks. The special value 0xffff indicates the terminating block of a file.	roup

Cluster	Next
2	0xffff
3	5
4	6
5	0xffff
6	7
7	2
8	3

Details for the file access table (FAT)

Each cluster contains 4 blocks and each block contains 512 bytes

X If file A begins at cluster 4 then how many clusters does it contain? *	0/5
O 1	
O 2	
O 3	
O 4	
O 5	
Other: 16	×
Correct answer	
4	
Feedback	
file A comprises clusters 4, 6, 7 and 2.	

X	if file A begins at cluster 4 and currently uses 50% of the space in its	0/5
	terminating cluster then how large is file A? *	

2048 bytes

X

8192 bytes

1792 bytes

7168 bytes

Other:

### Correct answer

7168 bytes

#### Feedback

file A complete fills 3 clusters (4, 6 and 7). each of those clusters has 4 blocks.

and then it fills one half of the fourth cluster (cluster 2) so 2 more blocks for a total of 3\*4 + 2 = 14 blocks. Each block is 512 bytes, so we have:

 $((3 \times 4) + 2) \times 512 = 14 \times 512 = 7168$  bytes

X	file B ends at cluster 5 and currently uses only 1 of the 4 blocks within	0/5
	cluster 5. How large is file B? *	

	6144 bytes	5	
--	------------	---	--

X

## Correct answer



4608 bytes

## Feedback

if file B ends at cluster 5 then it also comprises clusters 3 and 8. each cluster is 4 blocks in size. the file only uses 1/4 of cluster 5. It therefore uses:

$$(2*4*512) + (1*512) = 4608$$
 bytes

file B terminates at cluster 5 then how many clusters are used (either fully0/5 or partially) by file B? *
O 1
O 2
○ 3
Other:
Correct answer
3
Feedback  File B uses clusters 8, 3 and (at least part of) 5.
In linux it is possible to give a file an extension (such as .c, .o, .py, etc.) 5/5 even though that file does not contain the corresponding type of information. *
True
False

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When memory mapping a file, the OS must reserve virtual memory o/5 space equal to the size of the specified region of the file. *
○ True
False
Correct answer
True
Feedback  True. The OS does not need to reserve the full amount of *physical* memory space, but the *virtual* address space needs to be reserved so that the program can read or write any byte in the mapped file.
✓ If one process updates a file with write() and another process calls read() 5/5 on the same file then can linux automatically make the first process's updates available to the second process's read() call without first writing to persistent storage? *
<ul> <li>Yes. The linux block buffer cache allows updates to be done in memory and subsequent reads to be satisfied from the cache</li> <li>Nope, not gonna happen.</li> </ul>
✓ If an application writes to a file multiple times in succession, does each 5/5 write() call cause the data to be updated in persistent storage? *
<ul> <li>Yes</li> <li>Not always. Many of the updates will be done in the Linux block.buffer cache</li> </ul>

<u>:</u>

~	✓ If a storage block is found to be not assigned to a file and not on the free 5/5 list then what should the file system consistency checker do? *		
0	delete duplicate free list entries		
0	remove the block from the free list		
•	assign the block to the free list	<b>✓</b>	
0	allocate a new block and copy contents to the new block		
<b>~</b>	If a storage block is found to be both assigned to a file and on the free list then what should the file system consistency checker do? *	5/5	
0	allocate a new block and copy contents to the new block		
•	remove the block from the free list	<b>✓</b>	
0	delete duplicate free list entries		
0	assign the block to the free list		
<b>~</b>	If a storage block is found to be on the free list multiple times then what should the file system consistency checker do? *	5/5	
0	remove the block from the free list		
0	allocate a new block and copy contents to the new block		
•	delete duplicate free list entries	<b>✓</b>	
0	assign the block to the free list		

✓ If a storage block is found to be assigned to more than one file then what 5/5 should the file system consistency checker do? *	i
remove the block from the free list	
<ul> <li>allocate a new block and copy contents to the new block</li> </ul>	
delete duplicate free list entries	
assign the block to the free list	
Feedback	
the OS should also assign the newly copied block to one of the files in place of the shared block. it should also notify the owners of the file sbecause this it is likely their data is corrupted and requires manual inspection.	

This form was created inside of Portland State University.

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