

# Optimization Modeling — Part 1/4

First glances

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Areas of research:

- Multi-agent optimization: Bilevel programs, Game theory
- Optimization modeling: mainly focused on energy and environmental applications

A coordinated course with Metropolitan Electricity Authority (MEA) on  
STOCHASTIC OPTIMIZATION

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17 Nov – 18 Dec 2025



# Overview

The main objective of this lecture is to get you to know the **framework and workflow of optimization and modeling**.

The lecture consists of 4 parts:

Part 1: First glances

Part 2: Deterministic models

Part 3: Recalls on probability theory

Part 4: Basic stochastic models

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## Section 1

General ideas about optimization

## Subsection 1

### An optimization dictionary

# What is optimization ?

## **Optimization**

(n.) A process of making something as good as it can be.

## **Improvement**

(n.) The act of making something better.

## **Mathematical optimization**

(n.) Using and/or creating mathematical tools to do optimization.

# Where is optimization in the map of data analytics?

The four types of data analytics:

- **Descriptive:** What happened in the past ?  
Classical statistical tools, Data visualization, Clustering, etc.
- **Diagnostic:** Why something happened in the past ?  
Data discovery, Data mining, Root cause analysis, etc.
- **Predictive:** What is likely to happen in the future ?  
Regression, Classification, etc.
- **Prescriptive:** How to make something happen ?  
Optimization, Operations research, Reinforcement learning (repeated decision), etc.

- !! One should notice that Optimization remains the main tool when it comes to prescriptive analytics.
- !! Machine learning tools are largely predictive and descriptive in nature.
- !! ML tools are also useful to incorporate with Optimization.

# Ingredient dictionary of an optimization problem

| Ingredient                | Intuition  | Expression  |
|---------------------------|--|---|
| <b>Decision variable</b>  | What you control.  | $u \in \mathcal{U}$                                       |
| <b>Optimization space</b> | The scope of what you control.                             | $\mathcal{U}$   |
| <b>Objective function</b> | The quantitative evaluation (outcome) of the decision.     | $J : \mathcal{U} \rightarrow \mathbb{R} \cup \{+\infty\}$ |
| <b>Constraints</b>        | Limitations, represented as a set of admissible decisions. | $\mathcal{U}^{\text{ad}}$                                 |

## Decision variable

It is typical that  $\mathcal{U} = \mathbb{R}^n$  ( $n$  real variables), which means  $u$  is actually a vector

$$u = (u_1, \dots, u_n).$$

# Optimization sense

An objective function  $J$  could represent different quantities, e.g.

- **badness:** cost, disutility, etc. (The lower, the better.)
- **goodness:** profit, utility, etc. (The higher, the better.)

If  $J$  displays the badness, then we are looking at a **minimization problem**

$$\min_{u \in U^{\text{ad}}} J(u).$$

If  $J$  displays the goodness, then we are looking at a **maximization problem**

$$\max_{u \in U^{\text{ad}}} J(u).$$

## Useful notations

- $\arg \min_{u \in U^{\text{ad}}} J(u) = \{\text{Admissible decisions that minimizes } J(u)\}$
- $\arg \max_{u \in U^{\text{ad}}} J(u) = \{\text{Admissible decisions that maximizes } J(u)\}$

# Optimization sense

Here and forward, it is conventional that the min, max, arg min and arg max are taken over  $U^{\text{ad}}$ , unless stated otherwise.

## Observations

- $\min J(u) = -\max [-J(u)]$
- $\max J(u) = -\min [-J(u)]$
- $\arg \min J(u) = \arg \max [-J(u)]$  (arg min is obtained from arg max of negative objective.)
- $\arg \max J(u) = \arg \min [-J(u)]$  (arg max is obtained from arg min of negative objective.)

## Useful notes

For  $\lambda > 0$  and  $\beta \in \mathbb{R}$ , we have the following.

- $\min[\lambda J(u)] = \lambda \min J(u)$
- $\min[J(u) + \beta] = [\min J(u)] + \beta$
- $\arg \min[\lambda J(u)] = \arg \min J(u)$  (The positive scaling does not affect arg min.)
- $\arg \min[J(u) + \beta] = \arg \min J(u)$  (Shifting up and down do not affect arg min.)

That being said, we shall focus mainly on the theory of minimization.

## Constraints

Usually the constraint set  $U^{\text{ad}}$  is **explicit**, which means it is defined using equations and inequalities

$$U^{\text{ad}} = \left\{ u \mid \begin{array}{ll} g_i(u) \leq 0 & i = 1, \dots, r \\ h_j(u) = 0 & j = 1, \dots, \ell. \end{array} \right\}$$

= {Decisions satisfying certain equations and inequalities.}.

Additionally, we may require some slices of  $u$  to be integers. In this case, we have  $\mathcal{U} = \mathbb{R}^n$  and

$$U^{\text{ad}} = \left\{ u \mid \begin{array}{ll} g_i(u) \leq 0 & i = 1, \dots, r \\ h_j(u) = 0 & j = 1, \dots, \ell \\ u_k \in \mathbb{Z} & k \in K_0 \end{array} \right\}.$$

Here,  $K_0 \subset \{1, \dots, n\}$  represents the set of indices of  $u_i$ 's that are required to be integers.

# General framework

An **optimization problem** (or a **mathematical program**, or a **mathematical programming problem**) takes the following form.

$$\left\{ \begin{array}{ll} \text{optimization sense} & \text{objective function} \\ \min / \max & J(u) \\ \text{s.t.} & \begin{array}{ll} g_i(u) \leq 0 & i = 1, \dots, r \\ h_j(u) = 0 & j = 1, \dots, \ell \\ u \in C & \end{array} \\ & \text{possibly further limitations} \end{array} \right\} \text{constraints}$$

## Subsection 2

A tale of a hungry wizard — a toy modeling example

# A tale of a hungry wizard

Once upon a time in a fruit market...

- A hungry wizard has a magic sack that carries 1 ton of anything.
- He loves apple and guava, so he wants to fill his magic sack full with these fruits.
- So he goes to a fruit merchant and asks for a combination of apple and guava.
- A ton of apple is sold for 3 gold bars, and a ton of guava for 2 gold bar.
- The merchant has 0.8 ton of each fruit.

Since the merchant knows optimization, he formulates his decision model.

# A mathematical tale of a hungry wizard

## Decision variable

In the merchant's model, he sets  $\mathcal{U} = \mathbb{R}^2$  and

$$u = \left( \underbrace{u_1}_{\text{sales amount of apple}}, \underbrace{u_2}_{\text{sales amount of guava}} \right).$$

## Objective function

The objective function is an evaluation of his sales performance (income), which is

$$J(u) = J(u_1, u_2) = \underbrace{3u_1}_{\text{apple sales}} + \underbrace{2u_2}_{\text{guava sales}}.$$

## Constraints

There are a few constraints.

- The two fruits fill the sack full:  $u_1 + u_2 = 1$
- Stock limitations:  $u_1 \leq 0.8, u_2 \leq 0.8$
- Non-negativity:  $u_1 \geq 0, u_2 \geq 0$

## Subsection 3

What could it be ?

## What could it be ? — Manufacturing

$$\left\{ \begin{array}{l} \text{min} \quad \text{Manufacturing cost} \\ \text{s.t.} \quad \text{Demand constraint} \\ \qquad \quad \text{Manufacturing limit} \\ \qquad \quad \text{Material limit} \end{array} \right.$$

## What could it be ? — Resources

$$\left\{ \begin{array}{ll} \min & \text{Resources used} \\ \text{s.t.} & \text{Order quantity} \\ & \text{Product quality} \\ & \text{Production formula} \end{array} \right.$$

## What could it be? — Delivery

$$\left\{ \begin{array}{l} \text{min} \quad \text{Distance travelled} \\ \text{s.t.} \quad \text{Visits all customers} \\ \qquad \text{No subtour} \end{array} \right.$$

$$\left\{ \begin{array}{l} \text{min} \quad \text{Travel time} \\ \text{s.t.} \quad \text{Visits all customers} \\ \qquad \text{No subtour} \end{array} \right.$$

## What could it be ? — Load shifting

$$\left\{ \begin{array}{l} \text{min} \quad \text{Bill payment} + \text{discomfort} \\ \text{s.t.} \quad \text{Demand dynamics} \\ \qquad \text{Shift limits} \\ \qquad \text{Total energy preserved} \end{array} \right.$$

## What could it be ? — Electricity production

$$\left\{ \begin{array}{l} \text{min} \quad \text{Production cost} + \text{Startup cost} + \text{Shutdown cost} \\ \text{s.t.} \quad \text{Generation limits} \\ \qquad \text{Transmission limits} \\ \qquad \text{Demand balance} \\ \qquad \text{Reserve margin constraint} \\ \qquad \text{Emission limit} \\ \qquad \text{Generator ramping constraints} \end{array} \right.$$

## What could it be ? — Transmission expansion planning

$$\left\{ \begin{array}{l} \text{min} \quad \text{Investment cost + Operation cost} \\ \text{s.t.} \quad \text{Power flow balance} \\ \qquad \text{Material constraint} \\ \qquad \text{Budget constraint} \\ \qquad \text{Reliability constraint} \end{array} \right.$$

## What could it be? — Electricity market clearing

$$\left\{ \begin{array}{l} \text{max } \text{Social welfare} = \text{Consumer benefit} - \text{Generation cost} \\ \text{s.t. } \text{Market balance} \\ \quad \text{Network limits} \\ \quad \text{Nodal pricing constraints} \end{array} \right.$$

## What could it be? — Storage management

$$\left\{ \begin{array}{l} \text{min} \quad \text{Generation cost} + \text{Battery operation cost} \\ \text{s.t.} \quad \text{Power balance: generation} + \text{discharge} = \text{demand} + \text{charge} \\ \qquad \qquad \text{Storage limits} \\ \qquad \qquad \text{Storage behaviors} \end{array} \right.$$

## What could it be ? — Microgrid optimization

$$\left\{ \begin{array}{ll} \min & \text{Cost of local generation + Import/export cost} \\ \text{s.t.} & \text{Power balances} \\ & \text{Renewable constraints} \end{array} \right.$$



## Section 2

### Some theory

## Subsection 1

Linearity, quadraticity and convexity

# Linear programs (LP)

A **linear program**: LP is the problem where all the functions are defined by linear (affine) equations, equalities and inequalities.

A linear program has the form

$$\begin{cases} \min & e_1 u_1 + e_2 u_2 + \cdots + e_n u_n (+e_0) \\ \text{s.t.} & a_{i1} u_1 + a_{i2} u_2 + \cdots + a_{in} u_n \leq b_i \quad i = 1, \dots, r \\ & c_{j1} u_1 + c_{j2} u_2 + \cdots + c_{jn} u_n = d_j \quad j = 1, \dots, \ell. \end{cases}$$

In vector-matrix form, we write

$$\begin{cases} \min & e^\top u (+e_0) \\ \text{s.t.} & Au \leq b \\ & Cu = d. \end{cases}$$

## Quadratic programs (QP)

A **quadratic program**: **QP** is an optimization problem taking the form

$$\begin{cases} \min & u^\top Eu + e^\top u + e_0 = \sum_{i,j} e_{ij} u_i u_j + \sum_i e_i u_i + e_0 \\ \text{s.t.} & Au \leq b \\ & Cu = d. \end{cases}$$

# Convexity

A set  $C \subset \mathbb{R}^n$  is **convex** if we have

$$(1 - t)u + tv \in C \quad \text{for any } t \in [0, 1] \text{ and any } u, v \in C.$$

## Examples

- A **line segment** is convex.
- A **hyperplane** (a set defined by linear equalities) is convex.
- A **half space** (a set defined by linear inequalities) is convex.
- An **intersection of convex sets** is convex.

# Convexity

A function  $J : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{+\infty\}$  is

- **convex** if

for any  $u, v \in \mathbb{R}^n$  and any  $t \in [0, 1]$ ,

$$J((1-t)u + tv) \leq (1-t)J(u) + tJ(v),$$

- **strictly convex** if

for any  $u, v \in \mathbb{R}^n$ ,  $u \neq v$  and any  $t \in (0, 1)$ ,

$$J((1-t)u + tv) < (1-t)J(u) + tJ(v),$$

- **strongly convex** (with modulus  $a > 0$ ) if

for any  $u, v \in \mathbb{R}^n$  and any  $t \in [0, 1]$ ,

$$J((1-t)u + tv) \leq (1-t)J(u) + tJ(v) - \frac{a}{2}t(1-t)\|u - v\|^2.$$

## Fact

A strongly convex function is strictly convex, and a strictly convex function is convex.

# Convexity

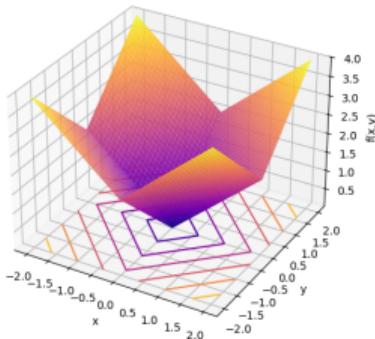
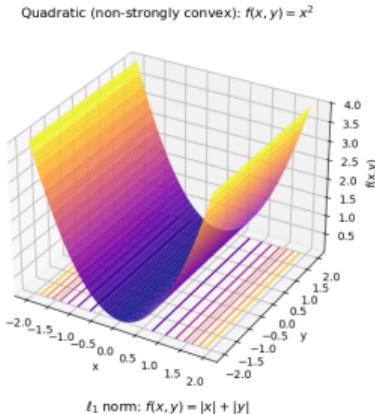
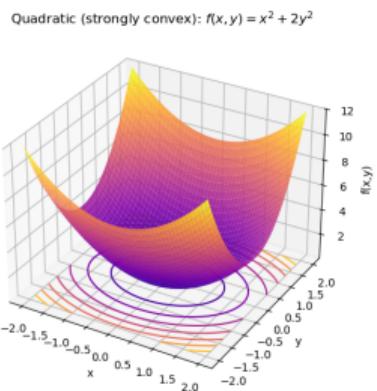
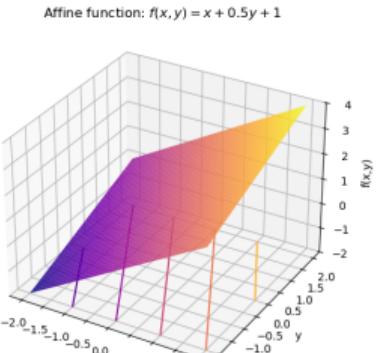


Figure: TL: affine    TR: convex quadratic    /    BL: Strongly convex quadratic    BR: Convex

# Convexity

## Examples (and facts)

- An **affine function** (a function defined by a linear equation) is convex.
- A **quadratic function** is convex if  $E$  has no negative eigenvalues.
- A **quadratic function** is strongly convex if  $E$  has all positive eigenvalues. In this case  $a = \min\{\text{eigenvalues of } E\}$ .
- A **positive scaling of a convex function** is convex.
- A **sum of convex functions** is convex.
- A **sum of convex functions with a strongly convex function** is strongly convex.
- A **pointwise supremum (maximum) of convex functions** is convex.
- The set  $\{u \mid g_i(u) \leq 0, \quad i = 1, \dots, r\}$ , where  $g_i$ 's are convex, is convex.

# Convexity with calculus

If  $J$  is continuously twice differentiable, then its **Hessian matrix** is defined by

$$H_J(u) = \left[ \frac{\partial^2 J}{\partial u_i \partial u_j}(u) \right]_{i,j}.$$

In this case, we have

$J$  is convex  $\iff H_J(u)$  has nonnegative eigenvalues at every  $u$

$J$  is strictly convex  $\iff H_J(u)$  has positive eigenvalues at every  $u$

$J$  is strongly convex (with modulus  $a > 0$ )  $\iff H_J(u)$  has eigenvalues  $\geq a$  at every  $u$

## Examples

- A quadratic function  $J(u) = u^\top E u + e^\top u + e_0$  has a Hessian  $H_J(u) = 2E$  at all  $u$ .
- An affine function  $J(u) = e^\top u + e_0$  has a Hessian  $H_J(u) = 0$  at all  $u$ .

# Convex programs

A **convex program** is an optimization problem where  $J$  is a convex function and  $U^{\text{ad}}$  is a convex set.

Usually,  $U^{\text{ad}}$  is given by linear equalities and convex inequalities

$$U^{\text{ad}} = \left\{ u \mid \begin{array}{l} g_i(u) \leq 0 \quad i = 1, \dots, r \\ Au = b \end{array} \right\},$$

where  $g_i$ 's are convex,  $A \in \mathbb{R}^{m \times n}$  and  $b \in \mathbb{R}^m$ .

## Mixed-integer programs

A **mixed-integer program** is an optimization problem where some slices of  $u$  are required to be integers.

This is often combined with LPs or QPs, which respectively result in MILPs or MIQPs.

## Subsection 2

Local vs global optima

# Local vs global optima

## Definition

A decision  $\bar{u}$  is

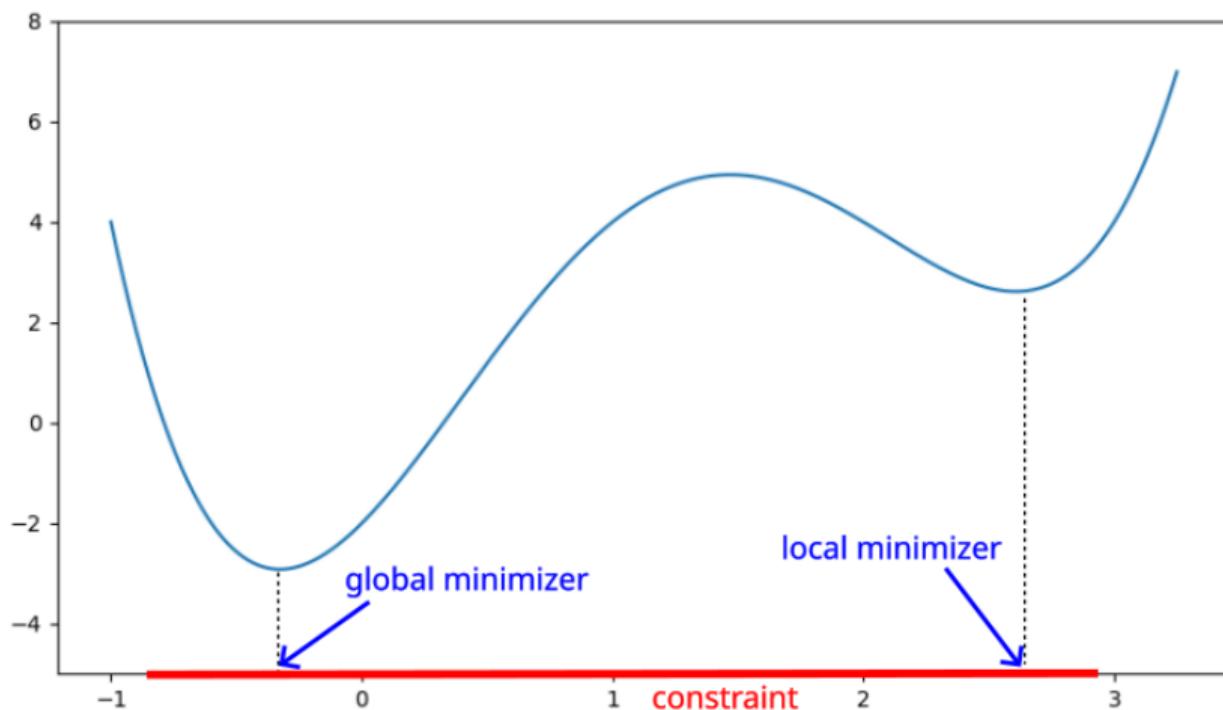
- a **global minimizer** of  $J$  if

$$J(\bar{u}) \leq J(u) \quad \text{for any admissible decision } u \in U^{\text{ad}}.$$

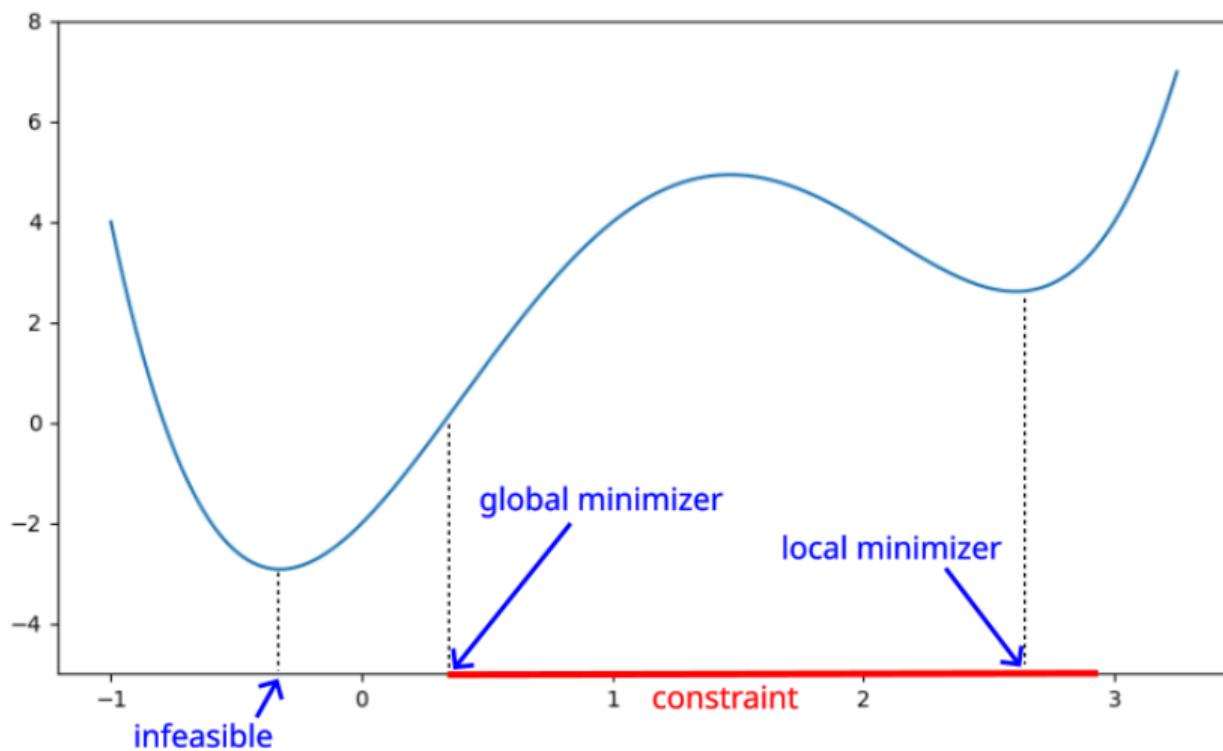
- a **local minimizer** of  $J$  if

$$J(\bar{u}) \leq J(u) \quad \text{for any admissible decision } u \in U^{\text{ad}} \text{ near } \bar{u}.$$

# Local vs global optima



# Local vs global optima



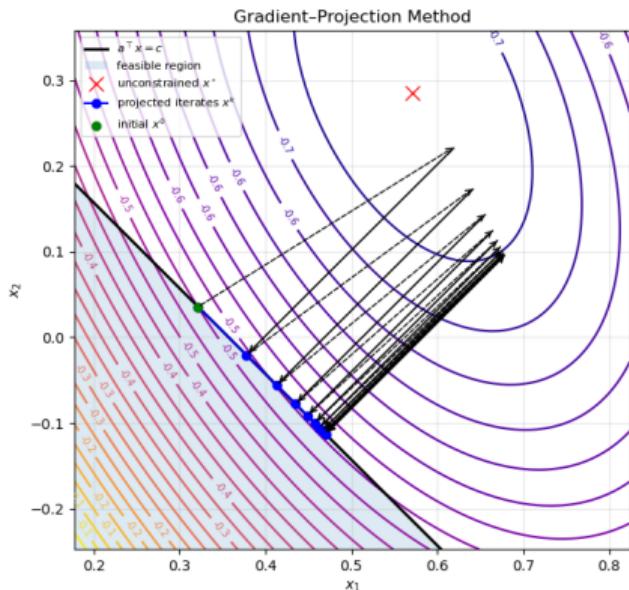
# Detecting local optima

## Fact

- In a convex program, a local minimizer is a global minimizer. (Local becomes global)
- An explicit convex program has a global minimizer if  $U^{\text{ad}}$  is bounded. (Existence)
- A strongly convex function has exactly one global minimizer. (Existence and uniqueness)

# Follow the slopes — Gradient-projection algorithm

A useful way to find a minimizer is to follow the slopes, using the gradient-projection algorithm.



## Gradient-projection algorithm.

**Input.** Initial guess  $u^{(0)}$ , step length  $\rho > 0$ , tolerance  $\varepsilon > 0$ .

**Output.** Optimal decision  $\bar{u}$ .

**Repeat**

$$| \quad u^{(k+1)} \leftarrow \text{proj}_{U^{\text{ad}}} [u^{(k)} - \rho \nabla J(u^{(k)})]$$

**until**  $\|\nabla J(u^{(k)})\| < \varepsilon$ .

### Subsection 3

Optimization with only equality constraints

# Optimization with only equality constraints

We consider here the following optimization problem

$$\begin{cases} \min & J(u) \\ \text{s.t.} & h_j(u) = 0 \quad j = 1, \dots, \ell \end{cases}$$

with  $J$  and  $h_j$ 's being differentiable.

Thus we may take advantage of the gradients

$$\nabla J(u) = \left( \frac{\partial J}{\partial u_i}(u) \right)_i \quad \text{and} \quad \nabla h_j(u) = \left( \frac{\partial h_j}{\partial u_i}(u) \right)_i$$

# Lagrange's approach

If we define a Lagrange function  $L : \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R} \cup \{+\infty\}$  by

$$L(u, \lambda) = \underbrace{J(u)}_{\text{original cost}} + \underbrace{\sum_{j=1}^{\ell} \lambda_j h_j(u)}_{\text{weighted penalty}},$$

then we have the gradient

$$\nabla L(u, \lambda) = (\nabla_u L(u, \lambda), \nabla_\lambda L(u, \lambda))$$

where

$$\begin{aligned}\nabla_u L(u, \lambda) &= \nabla J(u) + \sum_{j=1}^{\ell} \lambda_j \nabla h_j(u) \\ \nabla_\lambda L(u, \lambda) &= (h_j(u))_j.\end{aligned}$$

# Necessary condition (Sieve)

## Theorem

Let  $\bar{u} \in \mathbb{R}^n$ . We suppose that the constraints are **qualified\*** at  $\bar{u}$ .

Then a **necessary condition** for  $\bar{u}$  to be a **local minimizer** of  $J$  over  $U^{\text{ad}}$  is:

there exist  $\bar{\lambda}_j$ 's (called **Lagrange multipliers**) in which

$$\nabla L(\bar{u}, \bar{\lambda}) = 0. \quad (1)$$

Note that (1) is exactly the same as

$$\nabla J(\bar{u}) + \sum_{j=1}^{\ell} \bar{\lambda}_j \nabla h_j(\bar{u}) = 0, \quad h_j(\bar{u}) = 0 \quad \forall j = 1, \dots, \ell.$$

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\*For example, either (1) all  $h_j$ 's are linear, or (2) the gradients of  $h_j(\bar{u})$ 's are linearly independent, or (3) there is only one constraint.

# Sufficient condition

## Theorem

Let  $\bar{u} \in \mathbb{R}^n$ . We suppose that the objective function is convex and the constraints are linear.

Then a **sufficient condition** for  $\bar{u}$  to be a **global minimizer** of  $J$  over  $U^{\text{ad}}$  is:

there exist  $\bar{\lambda}_j$ 's (called **Lagrange multipliers**) in which

$$\nabla L(\bar{u}, \bar{\lambda}) = 0.$$

## Section 3

### Solvers and packages

# Some off-the-shelf solvers

| Solver            | Comments   |
|-------------------|--|
| CPLEX             | Commercial. Free for small problems. Fully free for academics. |
| Gurobi            | Commercial. Free for small problems. Fully free for academics. |
| GAMS              | Commercial. Free trial. Fully free for academics.              |
| SCIP <sup>†</sup> | Free and open-source. Also works with nonlinear problems.      |
| HiGHS             | Free and open-source. Only for LPs.                            |
| CBC               | Free and open-source.  |

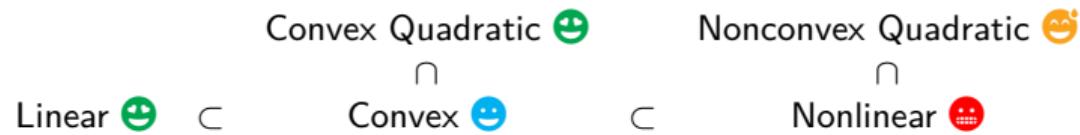
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<sup>†</sup>We shall use SCIP in our demo problems.

## Some off-the-shelf packages

| <b>Package</b> | <b>Language</b> | <b>Description</b>                               |
|----------------|-----------------|--|
| PySCIPOpt      | Python          | Interface to SCIP (MIP, MINLP)                   |
| Pyomo          | Python          | General modeling language; supports many solvers |
| PuLP           | Python          | Lightweight LP/MIP modeling                      |
| CVXPY          | Python          | Convex optimization modeling                     |
| JuMP           | Julia           | High-performance modeling for LP/QP/MIP/NLP      |

# The classes of optimization problems



Large mixed-integer problems are known to be slow to solve exactly. Hence we might have to rely on heuristic methods to *improve* the outcome and hope it is close enough to the minimizer.

# Getting started with SCIP

Let's see a demo of how to use SCIP through PySCIPOpt.

-» Continue to **Part 2**.