

Announcements

Pset 2 out

Recitation 1 grades up?

Chomsky Normal Form

$$S \rightarrow \epsilon$$

$$A \rightarrow a$$

$$A \rightarrow BC \quad (BC \neq S)$$

$$S, A, B, C \in V$$

$$S \rightarrow (S) \mid SS \mid \epsilon \leftarrow$$

Step 1: (of 5) Make a new start variable S_0
with rule $S_0 \rightarrow S$

$$\begin{array}{l} S_0 \rightarrow S \\ S \rightarrow (S) \mid SS \mid \epsilon \end{array}$$

\leftarrow

$$S_0 \Rightarrow S \Rightarrow \epsilon$$

Step 2: "Eliminate" ϵ -rules.

Call a variable A nullable if

- $A \rightarrow \epsilon$ is a rule

- $A \rightarrow X_1 X_2 \dots X_n$

where each X_i is a nullable variable

$$\begin{pmatrix} S \rightarrow TA \\ T \rightarrow \epsilon \\ A \rightarrow \epsilon \end{pmatrix}$$

Ex: $B \rightarrow a \underline{C} \underline{D} \underline{A} \underline{D}$ (C, D, A are nullable)

add:
rules: $B \rightarrow a D A D$

$B \rightarrow a D$;

$B \rightarrow a C A$;

nullable vars:
 S, S_0

$S_0 \rightarrow S$
 $S \rightarrow (S) | SS | \epsilon$

$$S_0 \rightarrow \underline{S} \mid \epsilon$$

$$S \rightarrow \underline{(S)} \mid \check{SS} \mid \underline{()}/\underline{S} \leftarrow$$

Step 3: "Eliminate" unit rules $S \rightarrow T$
 $(A \rightarrow B)$ $T \rightarrow A$

Form a directed graph

$$A \rightarrow B$$

$$B \rightarrow a$$

where the vertices are the variables,
 and add a directed edge from A to B
 if $A \rightarrow B$ is a unit rule.

if $A \rightarrow A$ is a unit rule, delete it.

for all rules $A \rightarrow uBy$, then if
 B has a path in the graph to C
 then add rule $A \rightarrow uCy$.

$$S_0 \rightarrow S / \epsilon$$

$$S \rightarrow (S) / SS / ()$$



$$S_0 \rightarrow \epsilon / (S) / SS / ()$$

$$S \rightarrow (S) / SS / ()$$

Office Hours / Problem Solving

Pset 1:

$$\begin{pmatrix} 0 \\ 0 \\ 1 \end{pmatrix} \begin{matrix} 1 \\ 1 \\ 1 \end{matrix} \begin{matrix} 0 \\ 1 \\ 1 \end{matrix}$$

Proof of reversal!

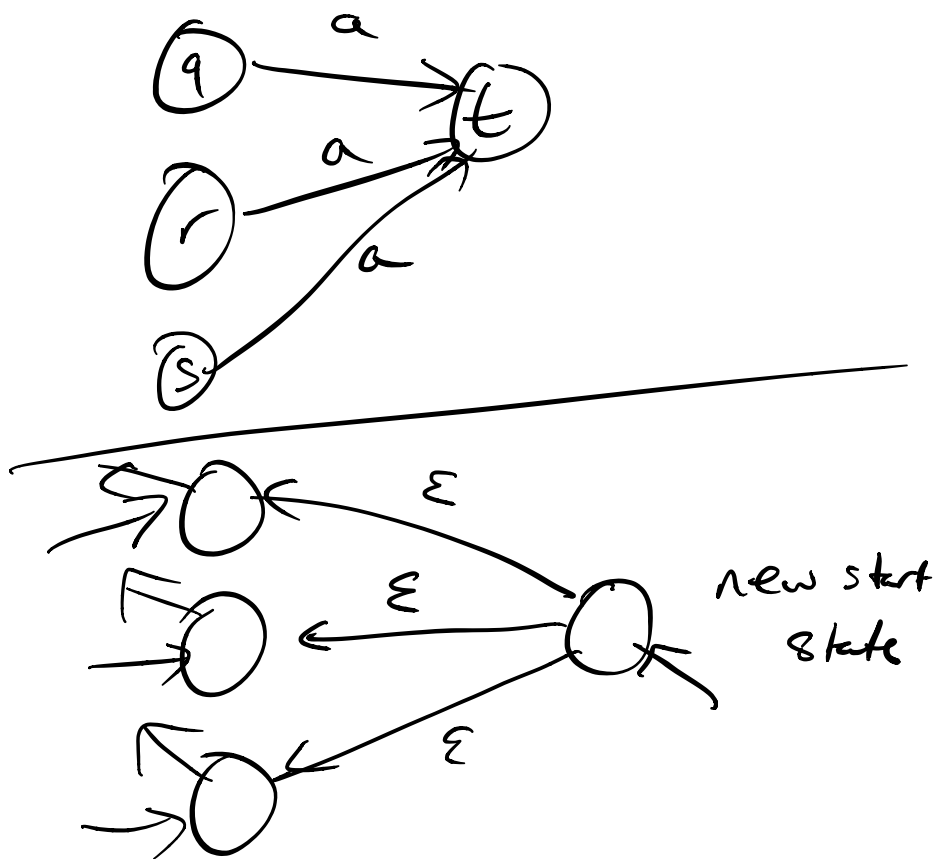
DFA D for L .

Want DFA/NFA/... for L^R .

if $w = w_1 w_2 \dots w_n$

then $w^R = w_n w_{n-1} \dots w_2 w_1$

$$L^R = \{w^R : w \in L\}$$



now make original start state final

Regexes if R is a regex for L :

1. if $R = a$ then $\text{rev}(R) = a$
2. if $R = \epsilon$ $\text{rev}(R) = \epsilon$
3. if $R = \emptyset$ $\text{rev}(R) = \emptyset$
4. if $R = R_1 \cup R_2$ $\text{rev}(R) = \text{rev}(R_1) \cup \text{rev}(R_2)$
5. if $R = R_1 R_2$ $\text{rev}(R) = \text{rev}(R_2) \text{rev}(R_1)$
6. if $R = R_1^*$ $\text{rev}(R) = (\text{rev}(R_1))^*$

$w_1 \quad w_2 \quad \dots \quad w_n$

$w_n^{\text{rev}} \quad w_{n-1}^{\text{rev}} \quad \dots \quad w_1^{\text{rev}}$

(Note: extra scratch paper)

$$L = \{a^m b^n : \underbrace{m < \sqrt{n}} \text{ or } \underbrace{n < \sqrt{m}}\}$$

Claim: L is not reg.

Proof: Assume it is

$\Rightarrow \exists$ a p for L .

Choose $w = a^p b^{p^2+1}$

$$x = a^\alpha$$

$$y = a^\beta$$

$$z = a^{p-\alpha-\beta} b^{p^2+1}$$

$$z = a^{p-\alpha-\beta} b^{p^2+1}$$

$$\text{Pump up: } xy^2z = a^{p+\beta} b^{p^2+1}$$

$$(m < \sqrt{n})$$

$$(n < \sqrt{m})$$

$$p+\beta < \sqrt{p^2+1}$$

$$(p+\beta)^2 < p^2+1$$

$$2p\beta + \beta^2 < 1$$

\Rightarrow contradiction!

$$p^2+1 < \sqrt{p+\beta}$$

$$\leq \sqrt{2p}$$

\Rightarrow contradiction!

$$\delta: Q \times \Sigma \rightarrow Q \times \underset{\uparrow}{\Sigma L, R}$$