

## ▼ Q3

A)

we define a function called m: length of the first string

n: length of the second string.

we consider words A and B as arrays of characters.  $A = a_0 a_1 \dots a_m$ ,  $B = b_0 b_1 \dots b_n$

We first initialize two-dimensional array  $DP_{(m+1) \times (n+1)}$  as 0 to store all the distance between substring of the words.  $DP_{ij}$  denotes the edit distance of  $a_0 a_1 \dots a_i$  and  $b_0 \dots b_j$ .

we use a bottom up manner.

$DP_{ij} =$

1. If  $i=0$ : it means the first string is empty, only option is to insert all characters of  $b_0 \dots b_j$  into  $a_0$  so  $DP_i = i$
2. If  $j=0$ : it means the second string is empty, only option is to remove all characters of first string so  $DP_j = j$
1. otherwise if  $a_i = b_j$ : it means no change is needed so the edit distance of  $DP_{ij}$  would be as the same edit distance  $DP_{i-1 j-1}$ . so  $DP_{ij} = DP_{i-1 j-1}$
2. o. w: when  $a_i$  and  $b_j$  are not the same, we need to whether insert, remove, or replace the letter so  $DP_i = \min(DP_{ij-1} + 1, DP_{i-1j} + 1, DP_{i-1j-1} + 2)$ :
  - insert:  $DP_{ij} = DP_{ij-1} + 1$  we insert  $b_j$  into  $a_0 \dots a_i$
  - remove:  $DP_{ij} = DP_{i-1j} + 1$  we delete  $a_i$  from  $a_0 \dots a_i$
  - replace:  $DP_{ij} = DP_{i-1j-1} + 2$  we replace  $a_i$  with  $b_j$

```
def edit_cost(a, b, m, n):
    dp = [[0 for x in range(n+1)] for x in range(m+1)]
    for i in range(m + 1):
        for j in range(n + 1):
            if i == 0:
                dp[i][j] = j
            elif j == 0:
                dp[i][j] = i
            elif a[i-1] == b[j-1]:
                dp[i][j] = dp[i-1][j-1]
            else:
                dp[i][j] = min(dp[i][j-1]+1, dp[i-1][j]+1, dp[i-1][j-1]+2)
    return dp[m][n], dp
```

in order to find the actions, we start from  $DP_{mn}$  and go up until we reach  $DP_{00}$  when we are at state  $DP_{ij}$ :

1. if  $i,j=0$ : it means we have finished converting the strings.
  2. if  $j=0$ : it means we are in the first column, and we want to reach  $b_0$ . so we delete  $a_j$  and go to state  $DP_{i-1j}$ . so the action is to delete  $a_j$  from the  $j$  position
  3. if  $i=0$ : we are in the first row, and our first string is  $a_0$  so we insert  $b_j$  to the string and go to  $DP_{ij-1}$ . so the action is to insert  $b_i$  from the  $i$  position
  4. o.w:
- if  $a_i = b_j$  are the same, we do not do any actions and go to state  $DP_{i-1j-1}$
  - o.w: we first calculate  $\min(DP_{ij-1} + 1, DP_{i-1j} + 1, DP_{i-1j-1} + 2)$ :
    - 1. if  $DP_{i-1j-1} + 2$  is the minimum, we replace  $a_i$  with  $b_j$  and go to  $DP_{i-1j-1}$  state.
    - 2. if  $DP_{ij-1} + 1$  is the minimum, we insert  $b_j$  and go to  $DP_{ij-1}$  state.
    - 3. if  $DP_{i-1j} + 1$  is the minimum, we delete  $a_i$  and go to  $DP_{i-1j}$  state.

when we go to the previous state, we execute the algorithm until we reach  $DP_{00}$

```
def find_action(m,n,dp,a,b):
    if m==0 and n==0:print("Done")
    elif m>0 and n==0:print("Delete",a[m-1]);find_action(m-1,n,dp,a,b);
    elif m==0 and n>0:print("Insert",b[n-1]);find_action(m,n-1,dp,a,b);
    elif a[m-1]==b[n-1]:find_action(m-1,n-1,dp,a,b)
    elif (dp[m-1][n-1]+2)<=(dp[m-1][n]+1) and (dp[m-1][n-1]+2)<=(dp[m][n-1]+1):
        print("Replace",a[m-1],"with",b[n-1])
        find_action(m-1,n-1,dp,a,b)
    elif (dp[m][n-1]+1)<(dp[m-1][n-1]+2) and (dp[m][n-1]+1)<=(dp[m-1][n]+1):
        print("Insert",b[n-1])
        find_action(m,n-1,dp,a,b)
    else:
        print("Delete",a[m-1])
        find_action(m-1,n,dp,a,b)
```

```
a="index"
b="inside"
min_cost,dp=edit_cost(a, b,len(a),len(b))
print("minimum cost to turn index into inside : "+str(min_cost))
find_action(len(a),len(b),dp,a,b)
```

```
minimum cost to turn index into inside : 3
Delete x
Insert i
Insert s
Done
```

```
a = "sunday"
```

```

b = "saturday"
min_cost,dp=edit_cost(a, b,len(a),len(b))
print("minimum cost to turn sunday to saturday : "+str(min_cost))
find_action(len(a),len(b),dp,a,b)

```

```

    minimum cost to turn sunday to saturday : 4
    Replace n with r
    Insert t
    Insert a
    Done

```

## B)

we consider words A and B as arrays of characters.  $A = a_0 a_1 \dots a_n$ ,  $B = b_0 b_1 \dots b_m$  where  $a_0 = b_0 = \emptyset$ .

we first create a two dimensional array  $DP_{(n+1) \times (m+1)}$  to store the longest common string of the substrings.

$DP_{ij}$  denotes the length of the longest common string of  $a_0 a_1 \dots a_i$  and  $b_0 \dots b_j$ .

for each  $0 \leq i \leq n$ , we loop through  $0 \leq j \leq m$ , since we iterate i and j in an increasing order, when we want to calculate  $DP_{ij}$ , if  $r + t < i + j$ , we already have the value of  $DP_{rt}$

$DP_{ij} =$

1. if  $i=0$ : it means we want to find  $LCS(\emptyset, b_0 \dots b_j)$  which is 0 so  $DP_{0j} = 0$
2. if  $j=0$ : it means we want to find  $LCS(a_0 \dots a_i, \emptyset)$  which is 0 so  $DP_{i0} = 0$
3. o.w:
  - if  $a_i = b_j = x$ : it means we want to find  $LCS(a_0 \dots a_{i-1}x, b_0 \dots b_{j-1}x)$ . by removing  $x$  from both strings, we just need to find  $LCS(a_0 \dots a_{i-1}, b_0 \dots b_{j-1})$  and increment it by 1 (because the last character  $x$  is common) so  $DP_{ij} = DP_{i-1,j-1} + 1$
  - if  $a_i \neq b_j$ : since  $a_i$  and  $b_j$  are not equal, first we remove  $a_i$  and calculate  $l_1 = LCS(a_0 \dots a_{i-1}, b_0 \dots b_j)$  and then remove  $b_j$  and calculate  $l_2 = LCS(a_0 \dots a_i, b_0 \dots b_{j-1})$ . since we want the length of the longest common string, we choose the maximum of  $(l_1, l_2)$ . so  $DP_{i,j} = \max(DP_{i-1,j}, DP_{i,j-1})$

so  $DP_{nm}$  would be the length of longest common string of A and B

since we first initialized  $DP$  as zeros, and the first column and value are all 0, we do not iterate the first row and column in the code.

```

def lcs(a , b):
    n = len(a)
    m = len(b)
    dp = [[0]*(m+1) for i in range(n+1)]
    for i in range(1,n+1):
        for j in range(1,m+1):

```

```

    for j in range(1,m+1):
        if i == 0 or j == 0 :
            dp[i][j] = 0
        elif a[i-1] == b[j-1]:
            dp[i][j] = dp[i-1][j-1]+1
        else:
            dp[i][j] = max(dp[i-1][j] , dp[i][j-1])
    return dp[n][m]

```

```

s1 = "abdacbab"
s2 = "acebfca"
print ("Length of the longest common string of abdacbab and acebfca: ", lcs(s1, s2) )

```

Length of the longest common string of abdacbab and acebfca: 4

```

s1 = "hello"
s2 = "what"
print ("Length of the longest common string of hello and what: ", lcs(s1, s2) )

```

☞ Length of the longest common string of hello and what: 1