

CSC 225

addie@uvic.ca

the time

Announcements

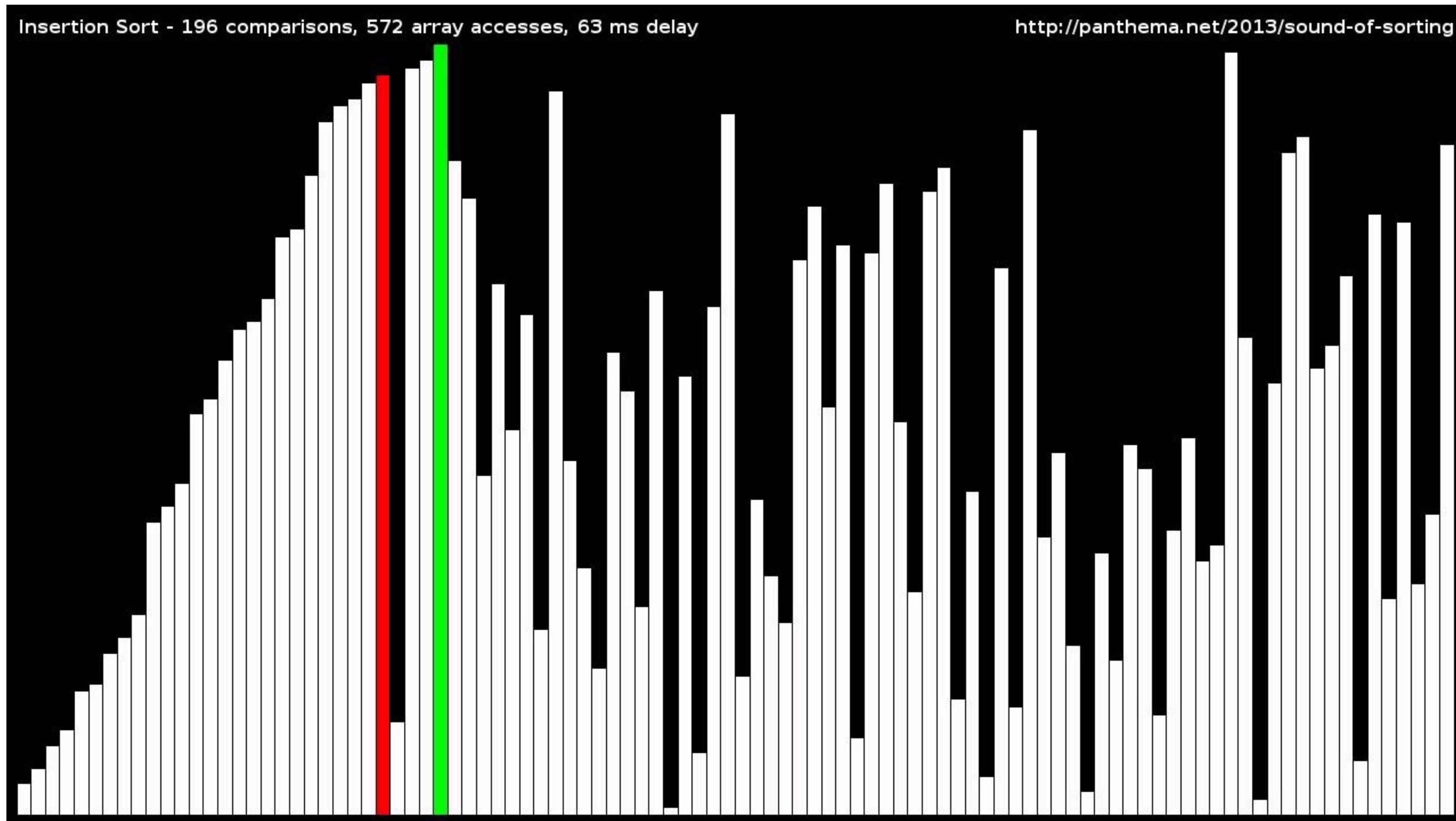
- program 1 due on Saturday
- please be sure to do your own testing!
- quiz 6 out today, due tomorrow



Poll

[menti.com](https://www.menti.com)

Insertion Sort



Insertion Sort

Idea: like adding new elements to an already sorted list. slot in the element.

Show what the list
looks like after each
pass

{10, 3, 26, 9, 4, 8, 12}

Insertion Sort

Algorithm InsertionSort(A, n) :

```
  for  $k \leftarrow 1$  to  $n-1$  do
     $val \leftarrow A[k]$ 
     $j \leftarrow k-1$ 
    while  $j \geq 0$  and  $A[j] > val$  do
       $A[j+1] \leftarrow A[j]$ 
       $j \leftarrow j-1$ 
    end
     $A[j+1] \leftarrow val$ 
  end
end
```

What is the worst-case run-time?

Insertion Sort

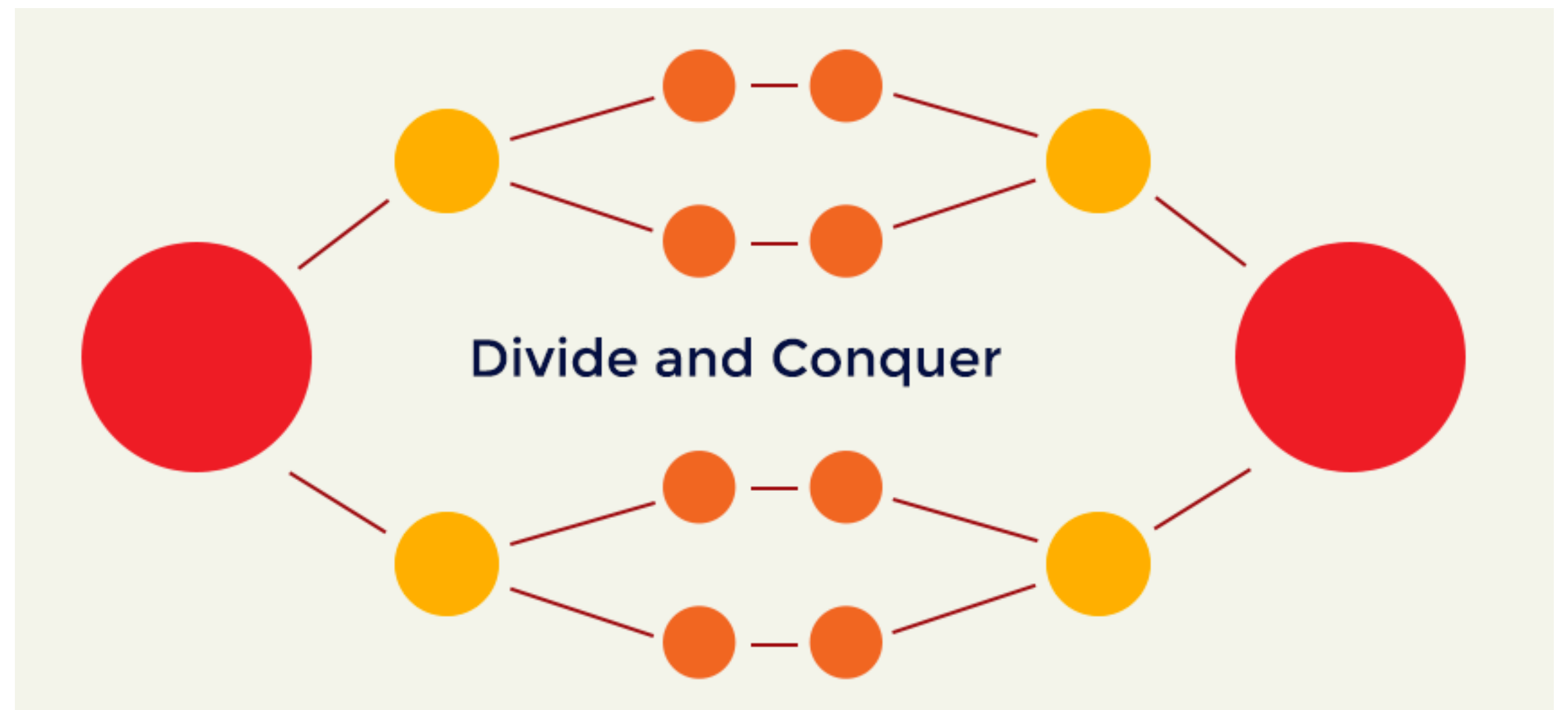
Algorithm InsertionSort(A, n) :

```
  for k ← 1 to n-1 do
    val ← A[k]
    j ← k-1
    while j ≥ 0 and A[j] > val do
      A[j+1] ← A[j]
      j ← j-1
    end
    A[j+1] ← val
  end
end
```

Best case input?
Worst case input?

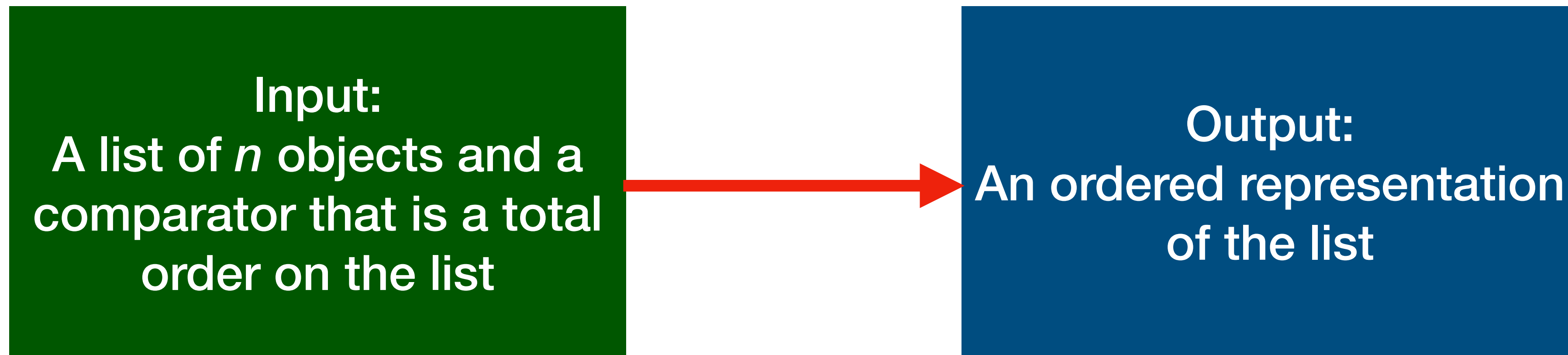
Divide and Conquer

- an efficient and easy-to-understand algorithm technique
 - Binary Search
 - Merge sort
 - Quick sort
- repeatedly divide a problem into smaller pieces
 - solve smaller pieces
 - sometimes may need to combine smaller solutions



Merge Sort

Definition + requirements



Merge Sort

Intuition

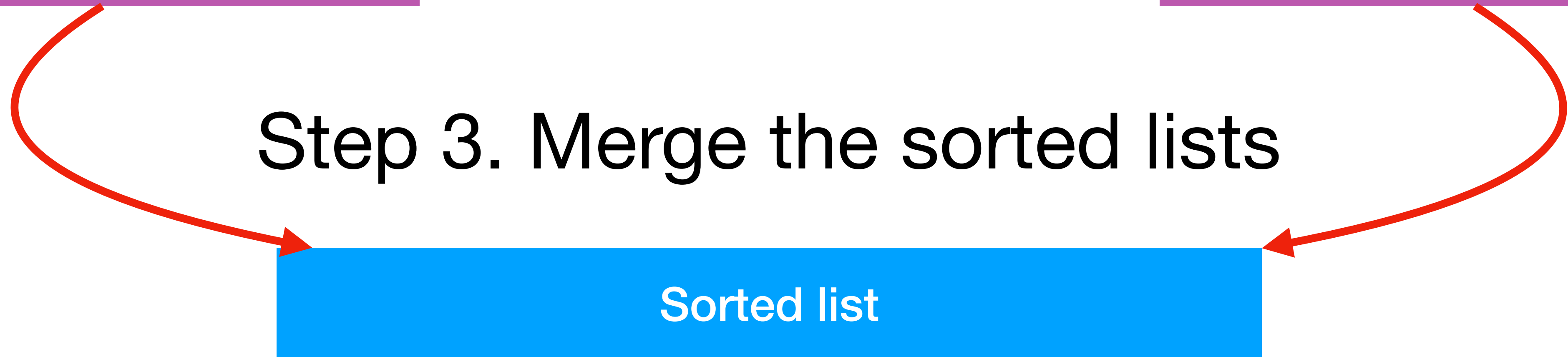
Step 1. Split the list in half



Step 2. Sort the sublists



Step 3. Merge the sorted lists



Merge sort: 3 steps

Step 1. Split the list in half

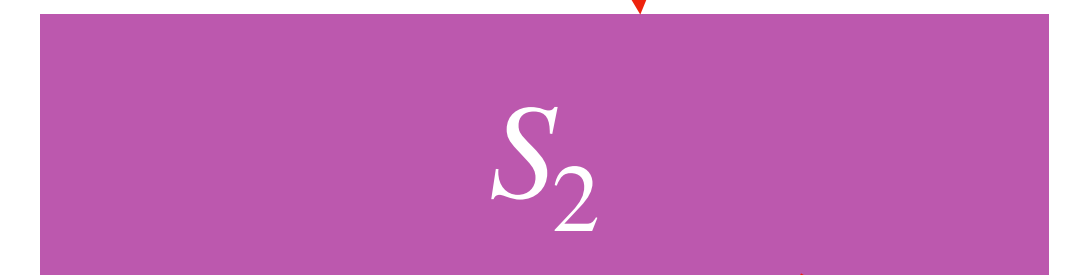


- Divide

Step 2. Sort the sublists

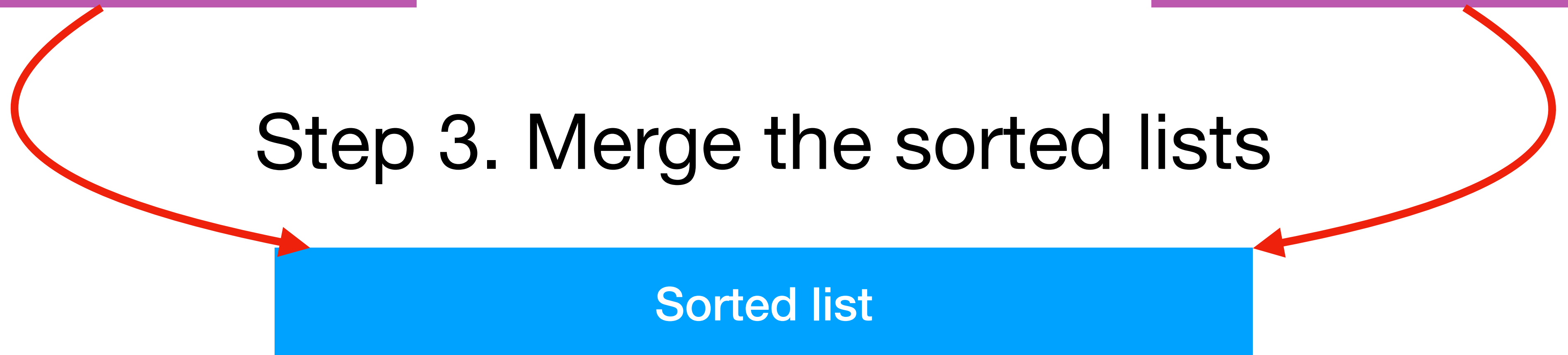


- Recur

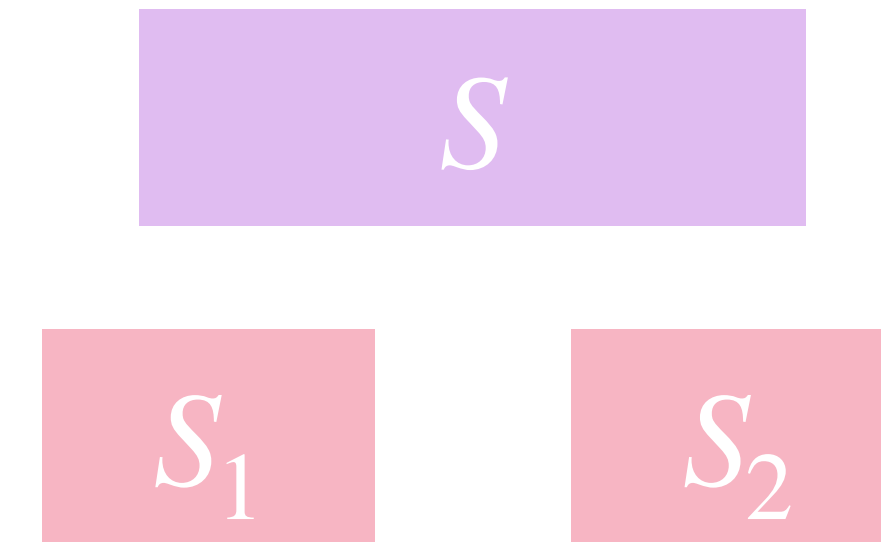


- Merge

Step 3. Merge the sorted lists



Step 1: Divide

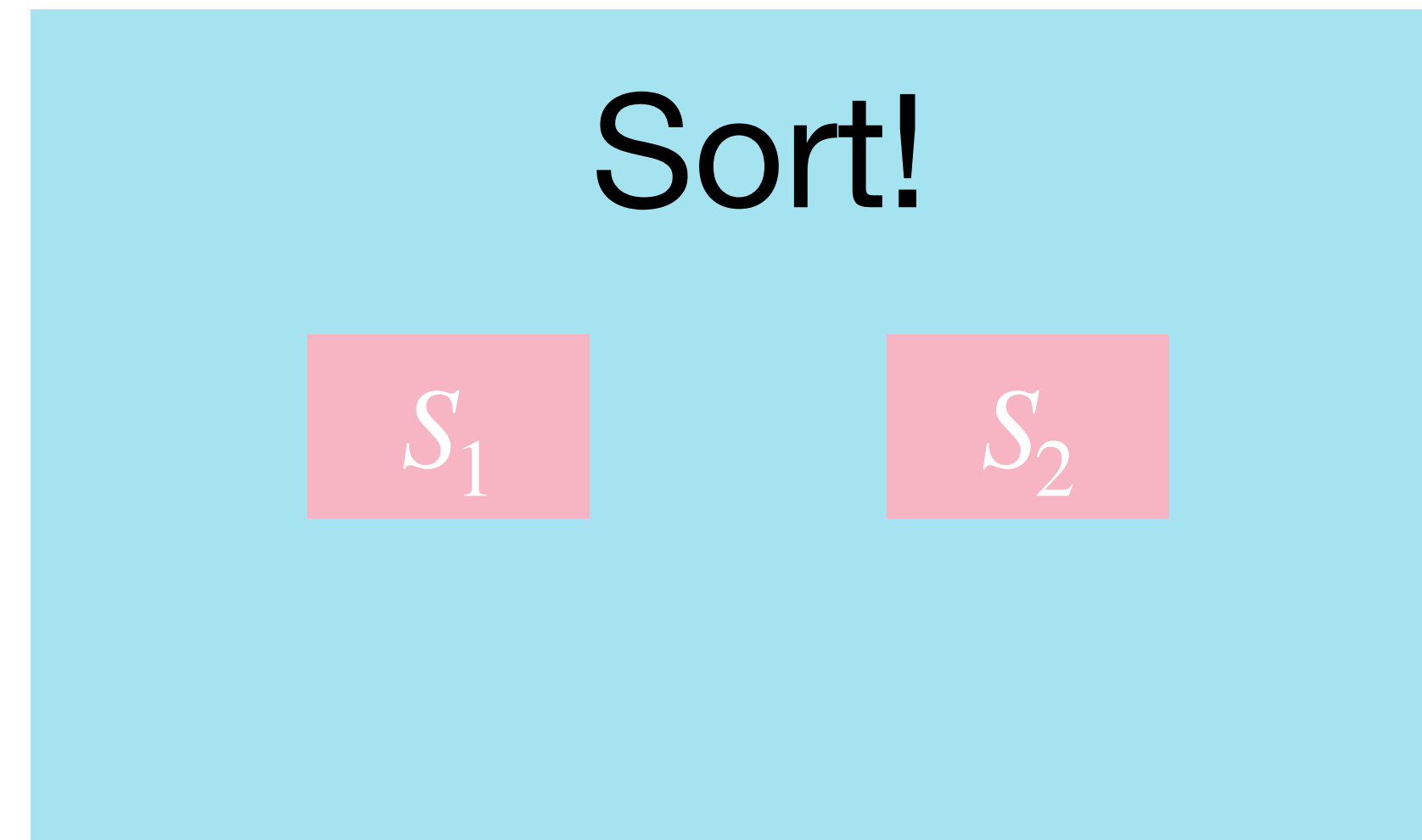


Let S be a sequence of n elements.

- If $|S| = 0$ or $|S| = 1$, return S
- Otherwise, break S into two parts S_1 and S_2 such that they each contain about half of the elements in S

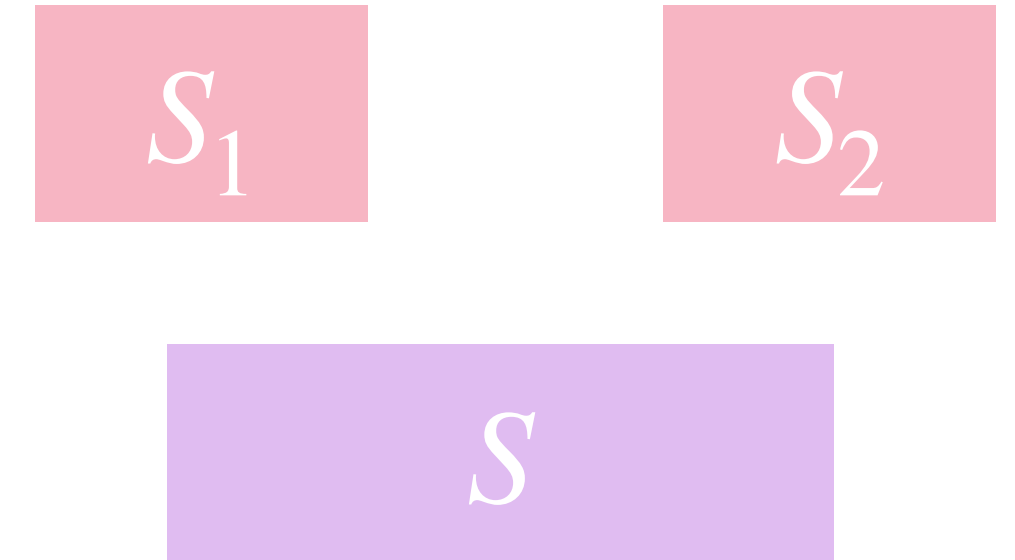
Step 2: Recur

- Recursively sort S_1 and S_2



Step 3: Merge

- Merge S_1 and S_2 to obtain a sorted S



Merging two sorted lists

- Look at the smallest elements in either list and compare
- Add the smaller one to the final sorted list
- Repeat until both sublists are empty

Merging two sorted lists

S_1

1	3	5	7
---	---	---	---

S_2

2	4	9	10
---	---	---	----

Example

Let $S = [8, 1, 11, 4, 12, 3, 7, 5]$. Draw the Merge sort tree.

Merge sort algorithm

Algorithm MergeSort(S) :

if $S.size() < 2$ **then**
 return S

 divide(S_1, S_2, S)
 $S_1 \leftarrow \text{MergeSort}(S_1)$
 $S_2 \leftarrow \text{MergeSort}(S_2)$
 merge(S_1, S_2, S)

return S

Merge sort algorithm

Algorithm MergeSort(S):

if $S.size() < 2$ **then**
 return S

$divide(S_1, S_2, S)$
 $S_1 \leftarrow MergeSort(S_1)$
 $S_2 \leftarrow MergeSort(S_2)$
 $merge(S_1, S_2, S)$

return S

Algorithm $divide(S_1, S_2, S)$:

for $i \leftarrow 0$ **to** $\lceil \frac{n}{2} \rceil - 1$ **do**
 $S_1[i] \leftarrow S[i]$
end

for $i \leftarrow \lceil \frac{n}{2} \rceil$ **to** $n-1$ **do**
 $S_2[i - \lceil \frac{n}{2} \rceil] \leftarrow S[i]$
end

end

Merging two sorted lists

- Look at the smallest elements in either list and compare
- Add the smaller one to the final sorted list
- Repeat until both sublists are empty

Merge sort algorithm

Algorithm MergeSort(S):

if $S.size() < 2$ **then**
 return S

 divide(S_1, S_2, S)
 $S_1 \leftarrow \text{MergeSort}(S_1)$
 $S_2 \leftarrow \text{MergeSort}(S_2)$
 merge(S_1, S_2, S)

return S

Algorithm merge(S_1, S_2, S):

$i \leftarrow 1$

$j \leftarrow 1$

while $i \leq n_1$ **and** $j \leq n_2$ **do**

if $S_1[i] \leq S_2[j]$ **then**

$S[i+j-1] \leftarrow S_1[i]$

$i \leftarrow i + 1$

else

$S[i+j-1] \leftarrow S_2[j]$

$j \leftarrow j + 1$

while $i \leq n_1$ **do**

$S[i+j-1] \leftarrow S_1[i]$

$i \leftarrow i + 1$

while $j \leq n_2$ **do**

$S[i+j-1] \leftarrow S_2[j]$

$j \leftarrow j + 1$

Merge sort algorithm

Algorithm MergeSort(S) :

```
if  $S.size() < 2$  then  
    return  $S$ 
```

```
divide( $S_1, S_2, S$ )  
 $S_1 \leftarrow$  MergeSort( $S_1$ )  
 $S_2 \leftarrow$  MergeSort( $S_2$ )  
merge( $S_1, S_2, S$ )
```

```
return  $S$ 
```

Worst case running time?

Merge sort alg

Worst case running time?

Algorithm MergeSort(S):

if $S.size() < 2$ **then**
 return S

$divide(S_1, S_2, S)$
 $S_1 \leftarrow MergeSort(S_1)$
 $S_2 \leftarrow MergeSort(S_2)$
 $merge(S_1, S_2, S)$

return S

Algorithm $divide(S_1, S_2, S)$:

for $i \leftarrow 0$ **to** $\lceil \frac{n}{2} \rceil - 1$ **do**
 $S_1[i] \leftarrow S[i]$
end

for $i \leftarrow \lceil \frac{n}{2} \rceil$ **to** $n-1$ **do**
 $S_2[i - \lceil \frac{n}{2} \rceil] \leftarrow S[i]$
end

end

Merge sort

Worst case running time?

Algorithm MergeSort(S):

if $S.size() < 2$ **then**
 return S

 divide(S_1, S_2, S)
 $S_1 \leftarrow \text{MergeSort}(S_1)$
 $S_2 \leftarrow \text{MergeSort}(S_2)$
 merge(S_1, S_2, S)

return S

Algorithm merge(S_1, S_2, S):

$i \leftarrow 1$

$j \leftarrow 1$

while $i \leq n_1$ **and** $j \leq n_2$ **do**

if $S_1[i] \leq S_2[j]$ **then**

$S[i+j-1] \leftarrow S_1[i]$

$i \leftarrow i + 1$

else

$S[i+j-1] \leftarrow S_2[j]$

$j \leftarrow j + 1$

while $i \leq n_1$ **do**

$S[i+j-1] \leftarrow S_1[i]$

$i \leftarrow i + 1$

while $j \leq n_2$ **do**

$S[i+j-1] \leftarrow S_2[j]$

$j \leftarrow j + 1$

Merge sort algorithm

Algorithm MergeSort(S) :

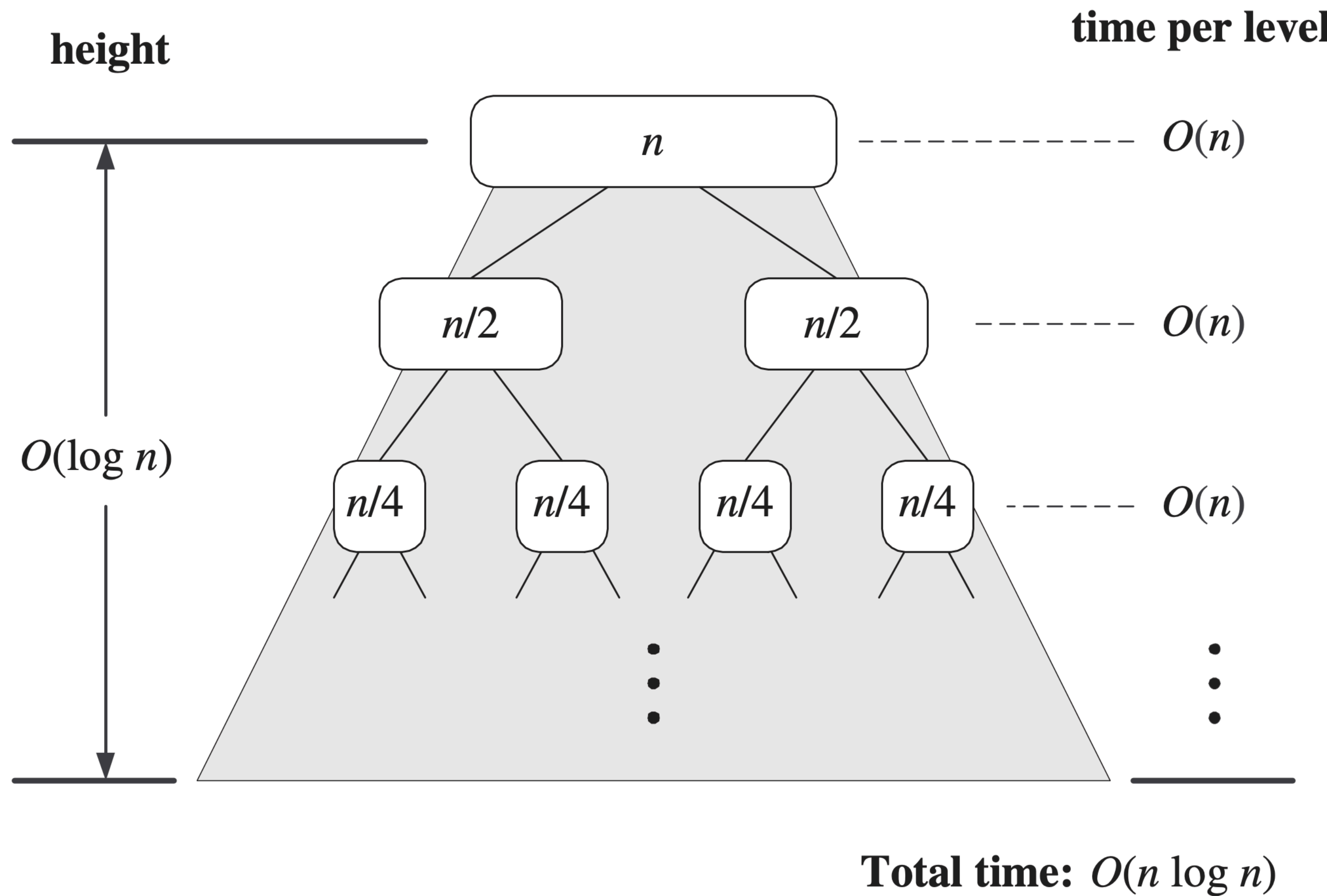
```
if  $S.size() < 2$  then  
    return  $S$ 
```

```
    divide( $S_1, S_2, S$ )  
     $S_1 \leftarrow \text{MergeSort}(S_1)$   
     $S_2 \leftarrow \text{MergeSort}(S_2)$   
    merge( $S_1, S_2, S$ )
```

```
return  $S$ 
```

Worst case running time?

Depth recursion of Merge sort



Depth recursion of Merge sort

