

Week 4 Thursday Lecture Notes

Prep:

- Read DVP **Chapter 6.3**
- **Implementation 1** Due tonight @Midnight.
- *There is an update to the specifications*: No Longer need to output every single pair that is tied closest together.
- **You only need to output the distance of the shortest and one of the pairs of points.
- Might be useful to make input file changeable for Implementation 1.

Dynamic Programming: Edit Distance:

- Minimum number of edit operations needed to turn s into t. (arbitrary strings t and s)
- Editing Operations:
 - o I = Insertion
 - o D = Deletion
 - o S = Substitution

Example

s:	I	N	T	E	*	N	T	I	O	N
t:	*	E	X	E	C	U	T	I	O	N
	<u>d</u>	<u>s</u>	<u>s</u>				<u>i</u>	<u>s</u>		

• Distance: 5 (assuming unit cost for each operation)

- **d = deletion || s = substitution || I = insertion.**

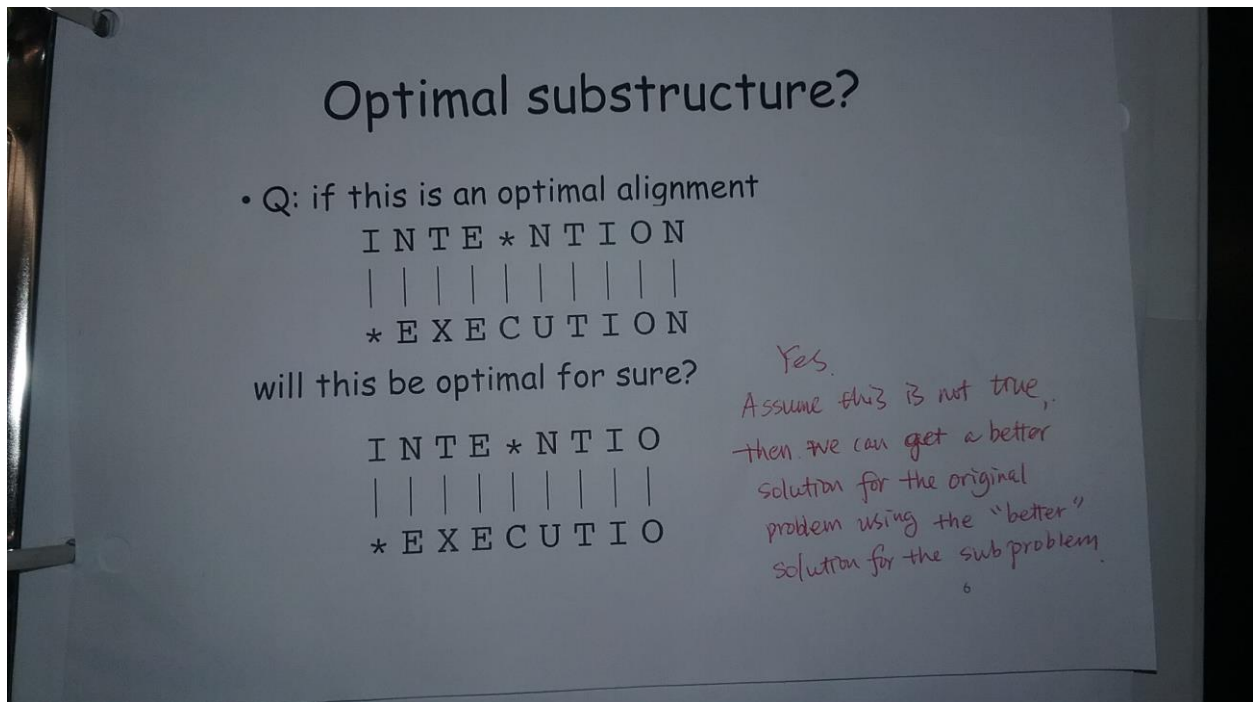
Finding the number of operations to change. String s to String t

Trying to find the most optimal edit distance.

Real Life Application:

- Computational Biology
- Google Translation and speech recognition. (Very Often will not be a 1-1 Language Conversion)

Optimal substructure. If substructures have optimal alignment will them combined be the overall optimal structure? **YES**



Picture Shows that you can Insert, Delete, Or Align the Strings

$D(i,j)$ = the edit distance between $s_1, s_2 \dots s_i$ and $t_1, t_2 \dots t_j$

$$s = s_1 s_2 \dots s_m, t = t_1 t_2 \dots t_n$$

- We can create sub-problems by considering the prefixes of s and t

$D(i, j)$ = the edit distance between $s_1 s_2 \dots s_i$ and $t_1 t_2 \dots t_j$

Q: To figure out $D(i, j)$, what are the possible choices we can make regarding the last positions i and j ?

A: Three possibilities:

Align i with j

$s_1 s_2 \dots s_{i-1}$	s_i
$t_1 t_2 \dots t_{j-1}$	t_j

$$D(i-1, j-1) + \begin{matrix} 0 \\ 1 \end{matrix}$$

Align i with *

$s_1 s_2 \dots s_{i-1}$	s_i
$t_1 t_2 \dots t_{j-1}$	*

$$D(i-1, j) + 1$$

Align j with *

$s_1 s_2 \dots s_{i-1}$	s_i	*
$t_1 t_2 \dots t_{j-1}$		t_j

$$D(i, j-1) + 1$$

- $D(i, i)$: the minimum of the three possible choices:

Picture Shows 3 Possibilities: Align, Insert, and Delete.

Align i with j : $D(i-1, j-1) + 0$ or 1 (0 for same, 1 for substitution)

Align i with *: $D(i-1, j) + 1$

Align j with *: $D(i, j-1) + 1$

Recurrence Relation for $D(i,j)$:

For $i, j \geq 1$

$$D(i, j) = \min \begin{cases} D(i-1, j) + 1 & \text{deletion} \\ D(i, j-1) + 1 & \text{insertion} \\ \underline{D(i-1, j-1) + \text{diff}(s_i, t_j)} & \text{align } i \text{ with } j \end{cases}$$
$$\text{diff}(a, b) = \begin{cases} 0 & \text{if } a = b \\ 1 & \text{if } a \neq b \end{cases}$$

Base case?
any others?

$D(0,0) = 0$ $D(0,1) = 1$
 $D(1,0) = 1$ $D(0,2) = 2$
 $D(2,0) = 2$ $D(0,3) = 3$
 \vdots \vdots
 $D(m,0) = m$ $D(0,n) = n$

Edit Distance

Picture Shows examples of why the base cases are true

$$D(0,n) = n$$

$$D(0,0) = 0$$

$$D(n,0) = m$$

Edit Distance: Pseudocode. Runtime speed of $O(m*n)$ roughly a quadratic function runtime (very expensive)

Edit Distance Table: How expensive it is for each position:

The Edit Distance Table

9	N	9	8	8															
8	O	8	7	7															
7	I	7	6	6															
6	T	6	5	5															
5	N	5	4	4															
4	E	4	3	4															
3	T	3	3	3															
2	N	2	2	2															
1	I	1	1	2															
0	#	0	1	2															
j=	#	E	X																
i=	0	1	2																

$\min(\swarrow, \leftarrow, \downarrow) = 1$
1, 1+1, 1+1

$\min(\swarrow, \leftarrow, \downarrow) = 2$
1+1, 2+1, 1+1

For This Table example we start at point $[2,9] = 8$.

8,7,6,5,4,3,3,2,1,0.

Computing Alignment

- ## Preformance:

Space: $O(nm)$

Backtrace: $O(n+m)$

Ending Notes:

- **Implementation Assignment 1:** Due Midnight 2/2/2017
- (**You can submit up to 24 hours late without penalty**)
- **Contact TA/Instructor** with questions about the implementation.
- **Set 3 Quiz Problems** will probably be available over the weekend to study.
- **Quiz 3** will be next Thursday
- (Recitation for Quiz Review at normal location **Wednesday 5:45-6:45pm**)

Finished Dynamic Programming Edit Distance Lecture (W4D2)

End of Week 4 Thursday Lecture Notes

~Information composed by Notetaker Scott Russell for CS 325 **DAS** student