

Process Environment

System Programming

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I. Process Environment

- □ How a program is executed
- □ How command-line arguments are passed to the new program
- □ What the typical memory layout looks like
 Different ways to terminate
- □ Environment variable, longjmp/setjmp functions, and process resource limits

I. Program Execution

- ☐ int main(int *args*, char **argv*[])
- □ When executed by the kernel *exec*, a special start-up routine sets up things, including the command-line arguments, so that the main function can be called.
- ☐ Five ways for a process to terminate
 - Normal termination: return from main, calling exit, and calling _exit or _Exit
 - Abnormal termination: calling abort and terminated by a signal

I. Process Termination

```
#include <stdlib.h>
void exit(int status);

void _Exit(int status);

#include <unistd.h>

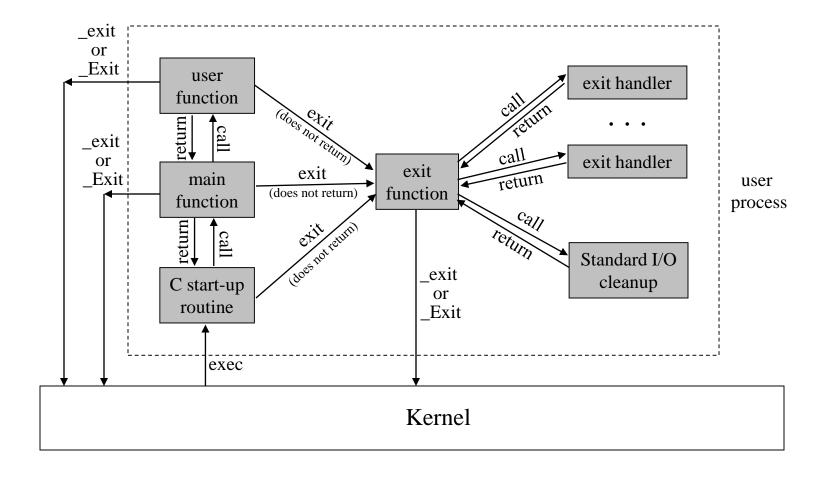
void _exit(int status);

□ exit performs certain cleanup processing and then returns to the kernel, while _exit/ Exit returns to the kernel
```

- In the case of exit, fclose is called for all open streams.
- ☐ The *exit status* is undefined if
 - Any of these functions is called without an exit status,
 - main does a return without a return value, or
 - main is not declared to return an integer.

immediately.

I. How a C Program is Started and Terminated



I. Process Termination

□ ISO C exit handlers: up to 32 functions that are automatically called by exit

```
#include <stdlib.h>
int atexit(void (*func)(void));
```

- □ Called in reverse order of their registration
- □ Program 7.3

I. Process Termination

```
#include "apue.h"
static void my_exit1(void);
static void my_exit2(void);
int main(void)
  if (atexit(my exit2) != 0)
    err_sys("can't register my_exit2");
  if (atexit(my exit1) != 0)
    err_sys("can't register my_exit1");
  if (atexit(my_exit1) != 0)
     err_sys("can't register my_exit1");
  printf("main is done\n");
  return(0);
static void my_exit1(void) {
  printf("first exit handler\n");
static void my exit2(void) {
  printf("second exit handler\n");
```

\$./a.out
main is done
first exit handler
first exit handler
second exit handler

Program 7.3

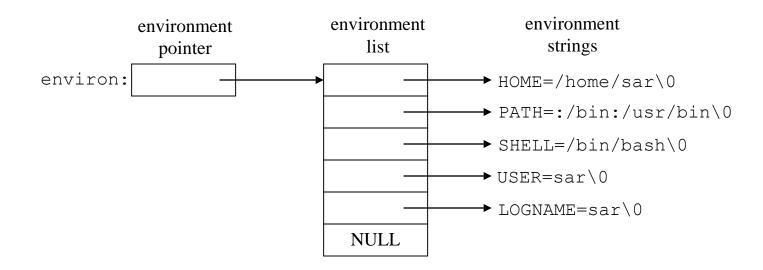


I. Command-Line Arguments

```
#include "apue.h"
int
main(int argc, char *argv[])
  int i;
  for (i = 0; i < argc; i++)
   printf("argv[%d]: %s\n", i, argv[i]);
  exit(0);
$ ./echoarg arg1 TEST foo
argv[0]: ./echoarg
arqv[1]: arq1
argv[2]: TEST
argv[3]: foo
```

I. Environment List

- □ Each program is passed an environment list, an array of char pointers.
- ☐ A global variable: extern char **environ;
- ☐ getenv/putenv functions



I. Memory Layout of a C Program

- □ Text segment: machine instructions
 - Sharable (a single copy for frequently executed programs such as text editors, C compiler, shells, etc.)
 - Often read-only
- □ (Initialized) data segment: variables specifically initialized in the program

int maxcount = 99;

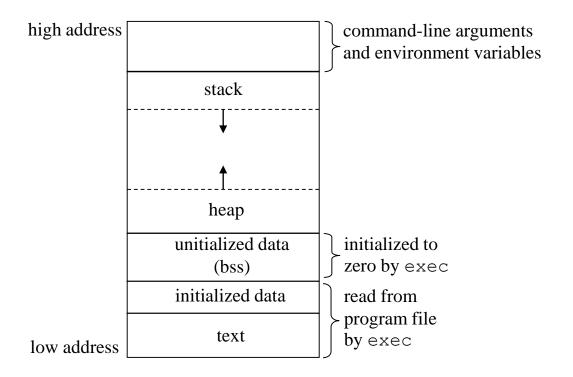
 Uninitialized data segment (bss segment): initialized by the kernel to arithmetic 0 or null pointers

long sum[1000];

- Stack: return addr & certain info about the caller's environ (such as some of the machine registers), auto variables
- ☐ Heap: dynamic memory allocation

I. Memory Layout of a C Program

```
$ size /bin/cc /bin/sh
text
       data
                bss
                        dec
                                hex
81920
       16384
                664
                                18298
                                        /bin/cc
                        98968
                                        /bin/sh
90112
       16384
                        106496 la000
                0
```



I. Shared Libraries

- □ Static library (*.a) vs. dynamic library (*.so)
- □ A single copy of the library routine somewhere in memory
- □ Reduces executable size at the cost of run-time overhead (the first time shared libraries are called)
- □ Can be replaced with new version without re-linking

```
$ ldd tcsh
libsocket.so.1 => /usr/lib/libsocket.so.1
libnsl.so.1 => /usr/lib/libnsl.so.1
libc.so.1 => /usr/lib/libc.so.1
libdl.so.1 => /usr/lib/libdl.so.1
libmp.so.2 => /usr/lib/libmp.so.2
```

□ LD_LIBRARY_PATH

I. Memory Allocation

```
#include <stdlib.h>
void *malloc(size_t size);
void *calloc(size_t nobj, size_t size);
void *realloc(void *ptr, size_t newsize);
void free(void *ptr);
```

☐ malloc

- allocates a specified number of bytes of memory
- initial value is indeterminate.

☐ calloc

- allocates space for a specified number of objects of a specified size.
- initialized to all 0 bits.

☐ realloc

- changes the size of a previously allocated area.
- The initial value of increased space is indeterminate.

I. Memory Allocation

□ Additional space for record keeping

- Blk size, a ptr to the next blk, and the like. Overwriting this record-keeping information often leads to a catastrophe.
- Freeing an already freed block or a block not obtained from the allocation functions – can also be fatal.

☐ Memory leakage

- ☐ void *alloca(size_t size)
 - It allocates size bytes of space in the stack frame.
 - No need to free the space.
 - automatically freed when the function from which alloca() is called returns.



I. Environment Variables

- □ Environment variables
 - name=value
- □ Defined environment variables
 - HOME, LANG, LC_ALL, LC_COLLATE, LC_CTYPE,
 LC_MONETARY, LC_NUMERIC, LC_TIME, LOGNAME, NLSPATH,
 SHELL, PATH, TERM, TZ
- □ csh built-in commands: env, setenv, unsetenv

```
#include <stdlib.h>
char *getenv(const char *name); /* ISO C */
```

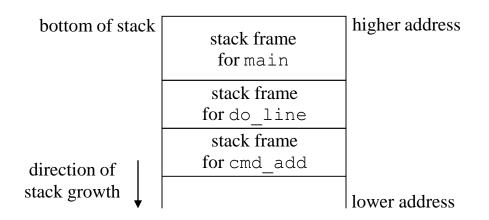
I. Environment Variables

```
#include <stdlib.h>
int putenv(const char *str);
int setenv(const char *name, const char *value, int rewrite);
void unsetenv(const char *name);

putenv takes a string of the form name=value. If name already exists, its old definition is first removed.
setenv sets name to value.
```

- If name already exists, then
 - If rewrite is nonzero, the existing definition is first removed.
 - Otherwise, no effect.
- ☐ unsetenv removes any definition of name.

- □ Useful for handling error conditions in a deeply nested function call.
- □ Program 7.9



□ Nonlocal goto to handle nonfatal errors

- If the cmd_add encounters an error, it might want to ignore the rest of the line, and return to the main to read the next input line.
- A chain of return values?



```
#include "apue.h"
#define TOK ADD 5
void do_line(char *);
void cmd add(void);
int get_token(void);
int main(void) {
  char line[MAXLINE];
  while (fgets(line, MAXLINE, stdin) != NULL)
     do_line(line);
  exit(0);
char *tok ptr; /* global pointer for get token() */
void do_line(char *ptr) /* process one line of input */
  int cmd:
  tok_ptr = ptr;
  while ((cmd = get_token()) > 0) {
     switch (cmd) { /* one case for each command */
       case TOK ADD:
          cmd_add();
          break;
```

```
void cmd_add(void)
{
  int token;
  token = get_token();
  /* rest of processing for this command */
}
int get_token(void)
{
  /* fetch next token from line pointed to by tok_ptr */
}
```

Program 7.9

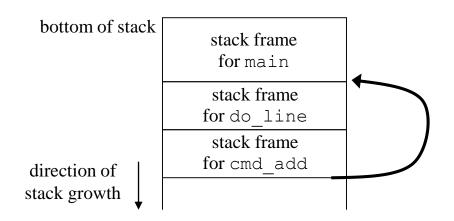


□ Nonlocal goto (setjmp/longjmp) as opposed to C goto statement within a function.

```
#include <setjmp.h>
int setjmp(jmp_buf env);
void longjmp(jmp_buf env, int val);
```

- □ setjmp returns 0 if called directly, nonzero if returning from a call to longjmp.
- □ jmp_buf is some form of array holding all the information required to restore the status of the stack.
- □ env is a global variable.

- □ longjmp(env, val)
 - The same env used in a call to setjmp.
 - A nonzero value val becomes the return value from setjmp.
- □ A longjmp call in cmd_add



```
#include "apue.h"
#include <setjmp.h>
jmp_buf jmpbuffer;
int main(void) {
          char line[MAXLINE];
          if (setjmp(jmpbuffer) != 0)
                     printf("error");
          while (fgets(line, MAXLINE, stdin) != NULL)
                    do_line(line);
          exit(0);
void cmd_add(void) {
          int token;
          token = get_token();
          if (token < 0)
                                          /* an error has occurred */
                     longjmp(jmpbuffer, 1);
          /* rest of processing for this command */
```

- □ What are the states of the automatic vars and register vars in the main function?
 - Rolled back or left alone? The answer is "indeterminate."
 - If you don't want your automatic var rolled back, define it with the volatile attribute. If you want it "left alone", declare it global or static.
- □ Program 7.13 on a system with vars in registers rolled back and vars in memory left alone.

```
$ cc testjmp.c; ./a.out
in f1():
globval = 95, autoval = 96, regival = 97, volaval = 98, statval = 99
after longjmp:
globval = 95, autoval = 96, regival = 97, volaval = 98, statval = 99
$ cc -O testjmp.c; ./a.out
in f1():
globval = 95, autoval = 96, regival = 97, volaval = 98, statval = 99
after longjmp:
globval = 95, autoval = 2, regival = 3, volaval = 98, statval = 99
```

```
#include "apue.h"
#include <setjmp.h>
static void f1(int, int, int, int);
static void f2(void);
static imp_buf impbuffer;
static int globval;
int main(void)
  int autoval:
  register int regival;
  volatile int volaval:
  static int statval;
  globval = 1; autoval = 2; regival = 3; volaval = 4;
  statval = 5;
  if (setimp(impbuffer) != 0) {
     printf("after longimp:\n");
     printf("globval = %d, autoval = %d,
             regival = %d, volaval = %d,
             statval = %d\n", globval, autoval,
             regival, volaval, statval);
     exit(0);
```

```
/* Change variables after setimp, but before longimp.*/
  globval = 95; autoval = 96; regival = 97;
  volaval = 98; statval = 99;
  /* never returns */
  f1(autoval, regival, volaval, statval);
  exit(0);
static void f1(int i, int j, int k, int l)
  printf("in f1():\n");
  printf("globval = %d, autoval = %d, regival = %d,
         volaval = %d, statval = %d\n",
         globval, i, j, k, l);
  f2();
static void f2(void)
  longimp(impbuffer, 1);
```

I.getrlimit and setrlimit

```
#include <sys/time.h>
#include <sys/resource.h>
int getrlimit(int resource, struct rlimit *rlptr);
int setrlimit(int resource, const struct rlimit *rlptr);
struct rlimit {
 rlim t rlim cur; /* soft limit: current limit */
 rlim t rlim max; /* hard limit: max value for rlim cur */
□ resource: RLIMIT AS, RLIMIT CORE, RLIMIT CPU,
  RLIMIT DATA, RLIMIT FSIZE, RLIMIT LOCKS,
  RLIMIT_MEMLOCK, RLIMIT_NOFILE, RLIMIT_NPROC,
  RLIMIT_RSS, RLIMIT_SBSIZE, RLIMIT_STACK, RLIMIT_VMEM
☐ The resource limits for a process are established by process 0
```

on boot and then inherited by each successive process.

Thank you for your attention!!

Q and A