

Signals

System Programming

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Introduction

- A long and thorough look at Unix signals
- The earlier implementations of signals
- POSIX.1 reliable-signal concept and all the related functions

☐ Signal

- Software interrupts: a way of handling asynchronous events, e.g. Ctrl-C
- 15 signals in Version 7, 31 signals in SVR4/4.4BSD, FreeBSD 5.2.1, Mac
 OS X 10.3, and Linux 2.4.22, and 38 signals for Solaris 9
- <signal.h>

☐ Conditions to generate a signal

- Terminal-generated signals, e.g. DELETE key (SIGINT)
- Hardware exceptions such as divide by 0, invalid memory reference (SIGSEGV), and the like
- kill(2) and kill(1)
- Software conditions, e.g. when out-of-band data arrives over a network connection (SIGURG), when a process writes to a pipe after the reader has terminated (SIGPIPE), and when an alarm clock expires (SIGALRM).

☐ Disposition (or action) of the signal

- Ignore the signal
 - SIGKILL and SIGSTOP can never be ignored.
 - Ignoring some signals, e.g. SIGFPE and SIGSEGV, results in undefined program behaviors.
- Catch the signal
- Default action
 - For most signals, it is to terminate the process

☐ Figure 10.1 Unix System signals

 "terminate+core" means that a memory image of the process is left in the file named core.

□ SIGABRT

Generated by abort function

□ SIGALRM

When alarm or setitimer function expires

□ SIGBUS

An implementation-defined hardware fault

□ SIGCHLD

When a child terminates or stops

□ SIGCONT

Sent to a stopped process when it is continued

□ SIGEMT

An implementation-defined hardware fault

□ SIGFPE

 An arithmetic exception, such as divide-by-0, floating point overflow, and so on

□ SIGHUP

- Sent to the controlling process if a disconnect is detected by the terminal interface
- Sent to each process in the foreground process group if the session leader terminates

□ SIGILL

When an illegal hardware instruction is execu

□ SIGINFO

 Sent to all processes in the foreground process group when we type the status key (often Ctrl-T)

□ SIGINT

 Sent to all processes in the foreground process group in case of the interrupt key (often DELETE or Ctrl-C)

□ SIGIO

To indicate an asynchronous I/O event

□ SIGIOT

To indicate implementation-defined hardware fault

□ SIGKILL

- Can't be caught or ignored. A sure way to kill any process.

□ SIGPIPE

 Generated when we write to a pipeline (a socket) when the reader (the other end) has terminated

□ SIGPOLL

When a specific event occurs on a pollable device

☐ SIGPROF

When a profiling interval timer (set by the setitimer) expires

☐ SIGPWR

- On a system with a UPS, to instruct the init process to shutdown everything
- System V's powerfail and powerwait in inittab file

□ SIGQUIT

 Sent to all processes in the foreground process group in case of the terminal quit key (often Ctrl-backslash)

□ SIGSEGV

To indicate an invalid memory reference

□ SIGSTOP

To stop a process, can't be caught or ignored

□ SIGSYS

To signal an invalid system call

□ SIGTERM

By the kill(1) command (by default)

□ SIGTRAP

An implementation-defined hardware fault

□ SIGTSTP

 Sent to all processes in the foreground process group in case of the terminal suspend key (often Ctrl-Z)

□ SIGTTIN

When a background process tries to read from its controlling terminal

□ SIGTTOU

When a background process tries to write to its controlling terminal

□ SIGURG

 To notify that an urgent condition has occurred, or in case of out-of-band data on a network connection

☐ SIGUSR1/SIGUSR2

A user-defined signal for use in application programs

□ SIGVTALRM

When a virtual interval timer (set by setitimer) expires

☐ SIGWINCH

When a window size (associated with (pseudo) terminal) is changed

□ SIGXCPU/SIGXFSZ

If soft CPU time limit / soft file size limit is exceeded

signal Function

☐ Figure 10.2

```
$ ./a.out & start process in background
[1] 4720 job-control shell prints job number and process ID
$ kill -USR1 4720 send it SIGUSR1

$ kill -USR2 4720 send it SIGUSR2

received SIGUSR2

$ kill 4720 now send it SIGTERM
[1] + Terminated ./a.out &
```

signal Function

Figure 10.2

```
static void
                  sig_usr(int); /* one handler for both signals */
int
main(void)
{
         if (signal(SIGUSR1, sig_usr) == SIG_ERR)
                  err sys("can't catch SIGUSR1");
         if (signal(SIGUSR2, sig_usr) == SIG_ERR)
                  err_sys("can't catch SIGUSR2");
         for (;;)
                  pause();
static void
sig_usr(int signo)
                          /* argument is signal number */
         if (signo == SIGUSR1)
                  printf("received SIGUSR1\n");
         else if (signo == SIGUSR2)
                  printf("received SIGUSR2\n");
         else
                  err_dump("received signal %d\n", signo);
```

Unreliable Signals

☐ Unreliable signals in earlier versions of the Unix System

- Signals could get lost
- The action for a signal was reset to its default action each time the signal occurred.
- Unable to turn a signal off when it is not wanted (i.e. no signal blocking)

```
int sig int flag;
main()
  int sig int();
  signal(SIGINT, sig int);
  while (sig int flag == 0)
     pause();
                                                  What if another signal occurs
                                                         here?
sig int()
  signal(SIGINT, sig int);
  sig int flag = 1;
```

Interrupted System Calls

☐ With earlier Unix systems, if a process caught a signal while being blocked in a "slow" system call, the system call was interrupted. It returned an error with error set to EINTR.

☐ Slow system calls

- reads from and write to certain file types (pipes, terminal devices, and network devices)
- opens of files that block until some condition occurs
- pause and wait
- certain ioctl operations
- some of the IPC functions (Chapter 15)



Interrupted System Calls

☐ We now have to handle the error return explicitly.

again:

```
if ( (n = read(fd, buf, BUFFSIZE)) < 0) {
  if (errno == EINTR)
    goto again; /* just an interrupted system call */
  /* handle other errors */
}</pre>
```

☐ Automatic restarting of certain interrupted system calls under 4.2BSD

- ioctl, read, readv, write, writev, wait, and waitpid

☐ A signal is *generated*, *delivered*, or *pending*. ☐ If a signal is *blocked*, and if its action is either SIG_DFL or to catch the signal, then the signal remains *pending* until the process unblocks the signal or change the action to SIG_IGN. ☐ What if a blocked signal is generated more than once before the signal is unblocked? Most Unix systems do not queue signals (i.e. deliver the signal once.) ☐ No order in which different signals are delivered to a process. ☐ Signal mask that defines the set of signals blocked.

Reliable Signal Terminology and Semantics

kill and raise Functions

```
#include <signal.h>
int kill(pid_t pid, int signo)
int raise(int signo);
```

- □ kill sends a signal to a process or a group of processes
 - pid > 0
 - Sent to the process whose process ID is *pid*.
 - pid == 0
 - Sent to all processes whose pgid equals the pgid of the sender.
 - pid < 0
 - Sent to all processes whose pgid equals the absolute value of pid.
 - pid == -1
 - Sent to all process for which the sender has permission to send a signal
 - Permission to send a signal
 - The real or effective UID of the sender has to equal the real or effective UID of the receiver. (If _POSIX_SAVED_IDS is supported, then the receiver's saved set-user-ID is checked instead of its effective UID.)
- ☐ raise sends a signal to itself.



alarm and pause Functions

```
#include <unistd.h>
unsigned int alarm(unsigned int seconds);

□ When the timer expires, SIGALRM is generated.
□ It returns 0 or number of seconds until previously set alarm
```

 Only one alarm clock per process. If there is a not-yet-expired clock for the process, the remaining seconds is returned.

```
#include <unistd.h>
int pause(void);
```

☐ pause suspends the calling process until a signal is caught. (it returns -1 with errno set to EINTR).

Thank you for your attention!!

Q and A