

# CHAPTER 2

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## Boolean Algebra

***This chapter in the book includes:***

Objectives

Study Guide

2.1 Introduction

2.2 Basic Operation

2.3 Boolean Expression and Truth Table

2.4 Basic Theorem

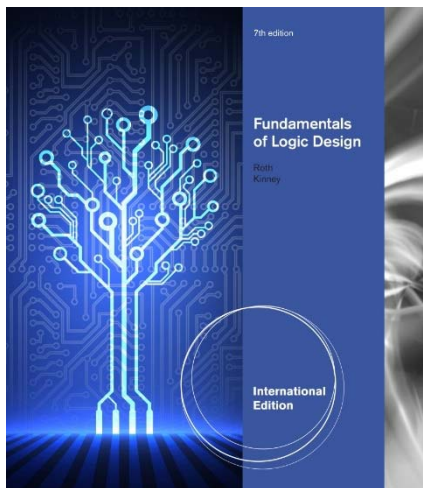
2.5 Commutative, Associative and Distributive Laws

2.6 Simplification Theorem

2.7 Multiplying Out and Factoring

2.8 DeMorgan's Laws

Problems



# Objectives

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Topics introduced in this chapter:

- Understand the basic operations and laws of Boolean algebra
- Relate these operations and laws to AND, OR, NOT gates and switches
- Prove these laws using a truth table
- Manipulation of algebraic expression using
  - Multiplying out
  - Factoring
  - Simplifying
  - Finding the complement of an expression

## 2.1 Introduction

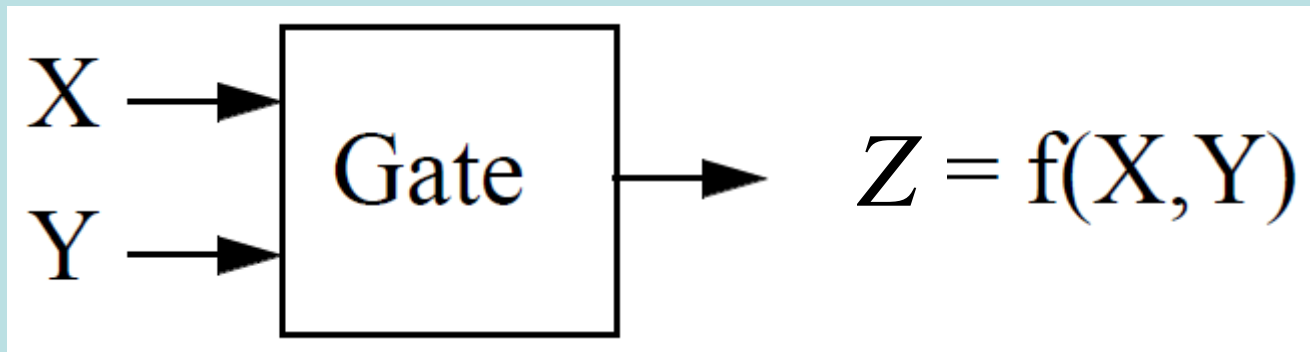
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- Basic mathematics for logic design: Boolean algebra
- Restrict to switching circuits( Two state values 0, 1) – Switching algebra
- Boolean Variable : X, Y, ... can only have two state values (0, 1)
  - representing True(1) False (0)

# Gate

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- Gate : A simple electronic circuit (a system) that realizes a logical operation.



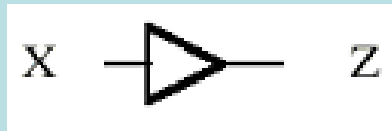
# Truth Table

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- **Truth Table** : listing all of its possible input configurations and the corresponding output signal
  - » The use of the symbols L and H usually correlates with the **high** and **low** voltages.
  - » The use of 0 (F) and 1 (T) must be associated with the voltages. It does not matter which way it is done.
    - \* If 1 is assigned to H and 0 to L  $\implies$  *positive logic*
    - \* If 0 is assigned to H and 1 to L  $\implies$  *negative logic*
- We will use the positive logic convention unless explicitly indicated otherwise.

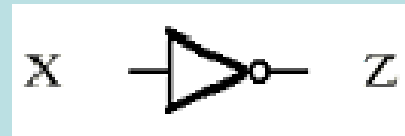
# Standard Gates & Symbols

- Buffer



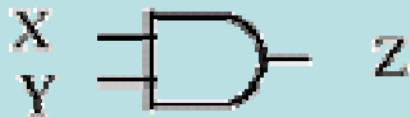
X	Z
0	0
1	1

- Not (Inverter or Complement)



X	Z
0	1
1	0

- AND



X	Y	Z
0	0	0
0	1	0
1	0	0
1	1	1

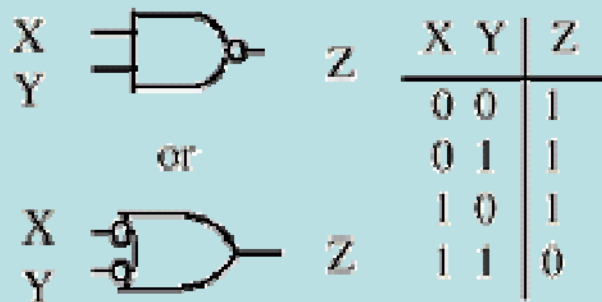
- OR



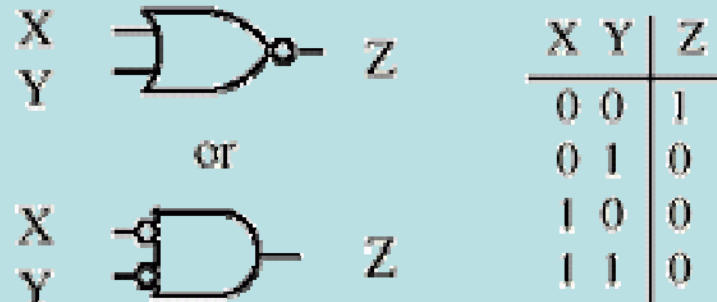
X	Y	Z
0	0	0
0	1	1
1	0	1
1	1	1

# Standard Gates & Symbols

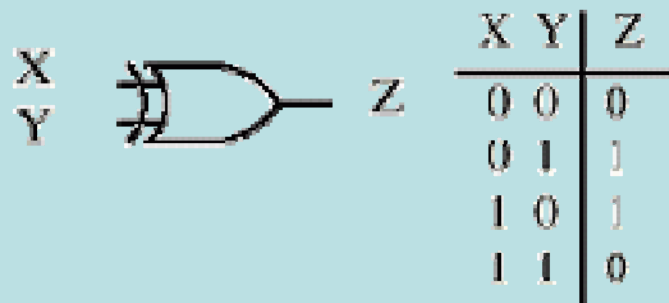
- NAND



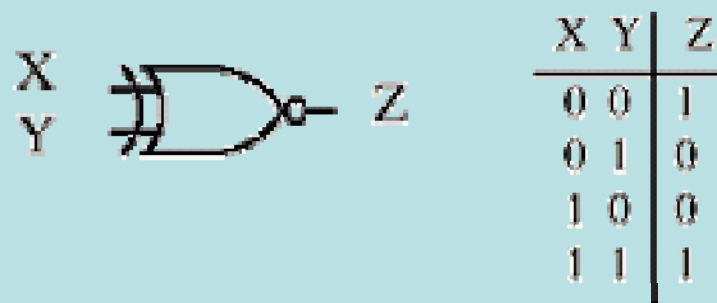
- NOR



- XOR (exclusive OR)



- Equivalence



## 2.2 Basic Operations

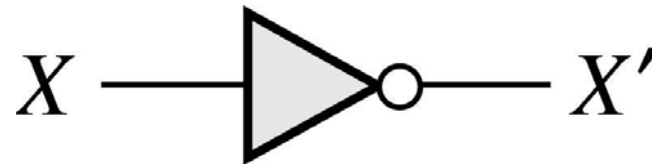
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### NOT(Inverter)

$$0' = 1 \quad \text{and} \quad 1' = 0$$

$$X' = 1 \quad \text{if} \quad X = 0 \quad \text{and} \quad X' = 0 \quad \text{if} \quad X = 1$$

### Gate Symbol





## 2.2 Basic Operations

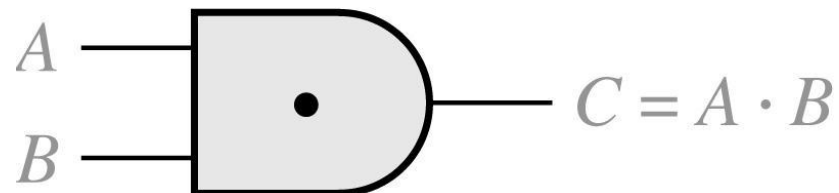
### AND

$$0 \cdot 0 = 0 \quad 0 \cdot 1 = 0 \quad 1 \cdot 0 = 0 \quad 1 \cdot 1 = 1$$

### Truth Table

A	B	$C = A \cdot B$
0	0	0
0	1	0
1	0	0
1	1	1

### Gate Symbol



## 2.2 Basic Operations

**OR**

$$0 + 0 = 0$$

$$0 + 1 = 1$$

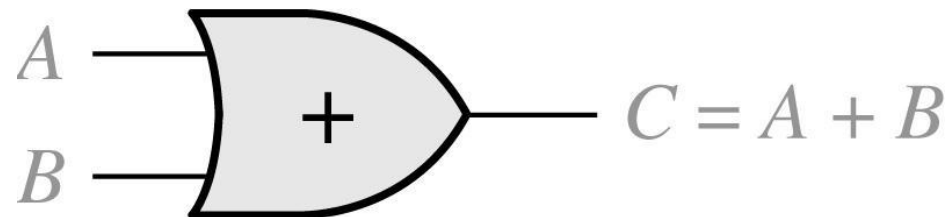
$$1 + 0 = 1$$

$$1 + 1 = 1$$

**Truth Table**

A B	$C = A + B$
0 0	0
0 1	1
1 0	1
1 1	1

**Gate Symbol**



## 2.2 Basic Operations

### Apply to Switch



$X = 0 \rightarrow$  switch open

$X = 1 \rightarrow$  switch closed

### AND

$$T = A \cdot B$$

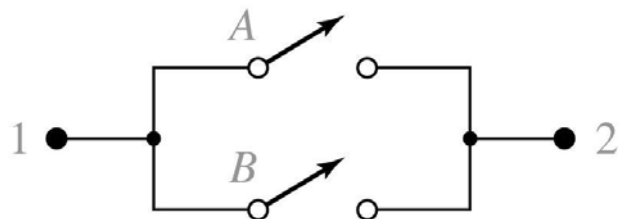


$T = 0 \rightarrow$  open circuit between terminals 1 and 2

$T = 1 \rightarrow$  closed circuit between terminals 1 and 2

### OR

$$T = A + B$$



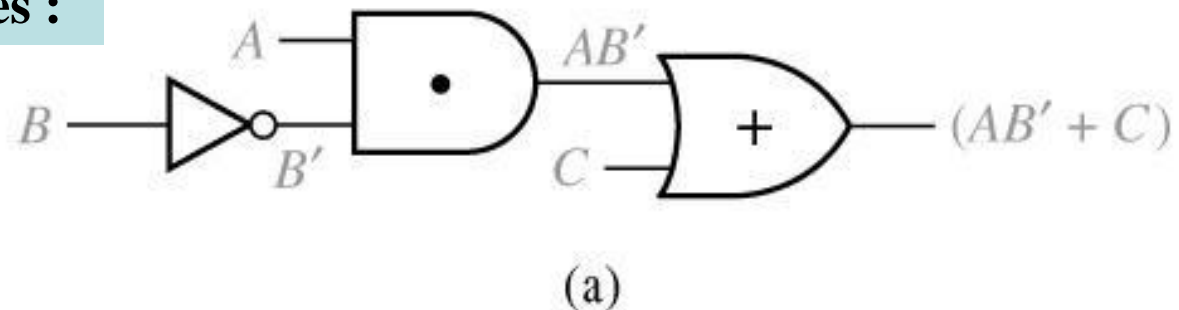
## 2.3 Boolean Expressions and Truth Tables

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**Logic Expression :**

$$AB' + C$$

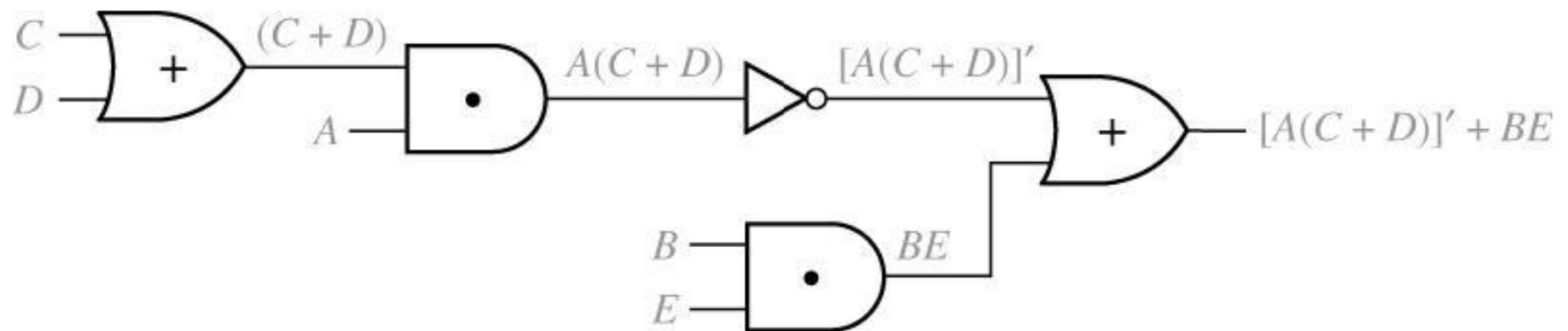
**Circuit of logic gates :**



## 2.3 Boolean Expressions and Truth Tables

**Logic Expression :**  $[A(C + D)]' + BE$

**Circuit of logic gates :**



**Logic Evaluation : A=B=C=1, D=E=0**

$$[A(C + D)]' + BE = [1(1 + 0)]' + 1 \cdot 0 = [1(1)]' + 0 = 0 + 0 = 0$$

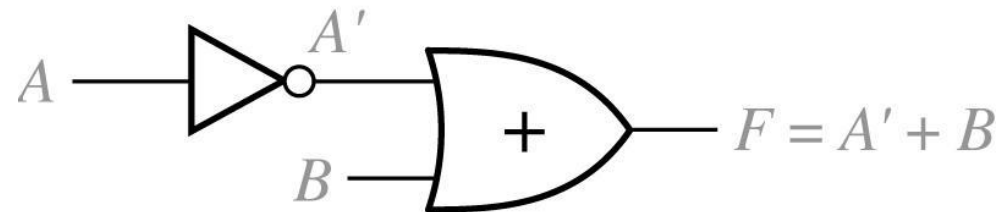
**Literal : a variable or its complement in a logic expression**

$$ab'c + a'b + a'bc' + b'c'$$

10 literals

## 2.3 Boolean Expressions and Truth Tables

### 2-Input Circuit and Truth Table



(a)

A	B	$A'$	$F = A' + B$
0	0	1	1
0	1	1	1
1	0	0	0
1	1	0	1

## 2.3 Boolean Expressions and Truth Tables

Proof using Truth Table

$$AB' + C = (A + C)(B' + C)$$

$n$  variable needs

$$\underbrace{2 \times 2 \times 2 \times \dots}_{n \text{ times}} = 2^n \text{ rows}$$

$n$  times

TABLE 2.1

A B C	B'	AB'	AB' + C	A + C	B' + C	(A + C)(B' + C)
0 0 0	1	0	0	0	1	0
0 0 1	1	0	1	1	1	1
0 1 0	0	0	0	0	0	0
0 1 1	0	0	1	1	1	1
1 0 0	1	1	1	1	1	1
1 0 1	1	1	1	1	1	1
1 1 0	0	0	0	1	0	0
1 1 1	0	0	1	1	1	1

## 2.3 Boolean Expressions and Truth Tables

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- So there are three ways of defining a switching (Boolean) function:

- (1) Logical expression

- (2) Truth table

- (3) Logic Network (circuit)

==> Three representations all describe the same function

- Precedence(우선순위) in algebraic expressions:

NOT AND OR except for brackets



## 2.4 Basic Theorems

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Operations with 0, 1

$$X + 0 = X$$

$$X \cdot 1 = X$$

$$X + 1 = 1$$

$$X \cdot 0 = 0$$

Idempotent Laws

$$X + X = X$$

$$X \cdot X = X$$

Involution Laws

$$(X')' = X$$

Complementary Laws

$$X + X' = 1$$

$$X \cdot X' = 0$$

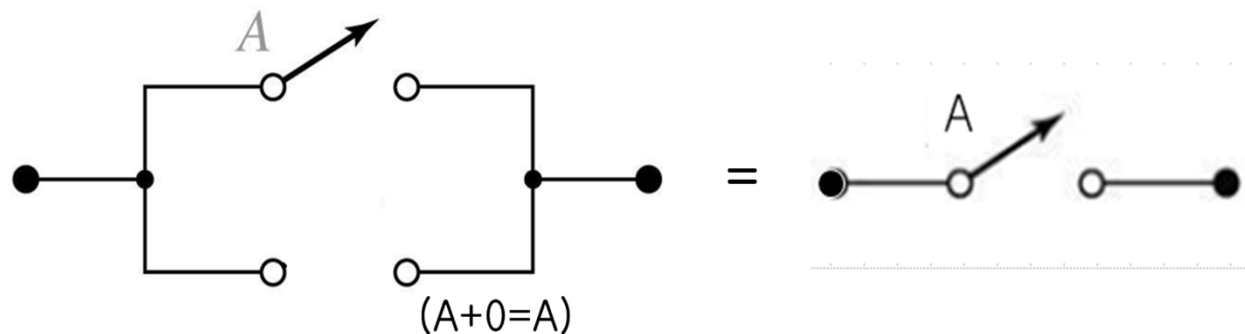
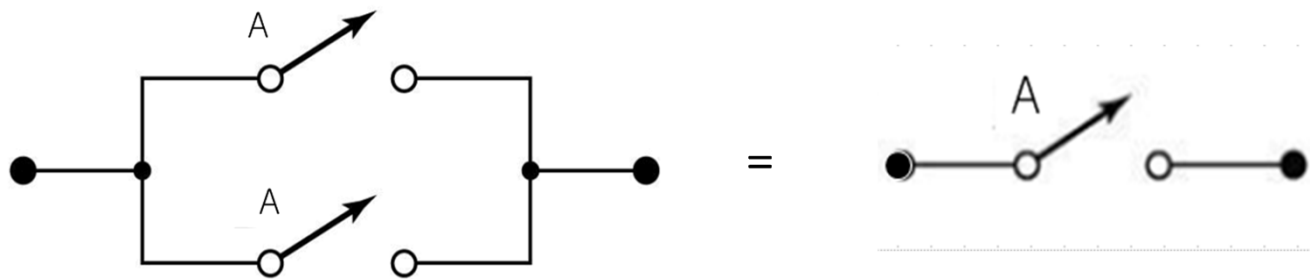
Proof  $X = 0$ ,  $0 + 0' = 0 + 1 = 1$ , and if  $X = 1$ ,  $1 + 1' = 1 + 0 = 1$

Example

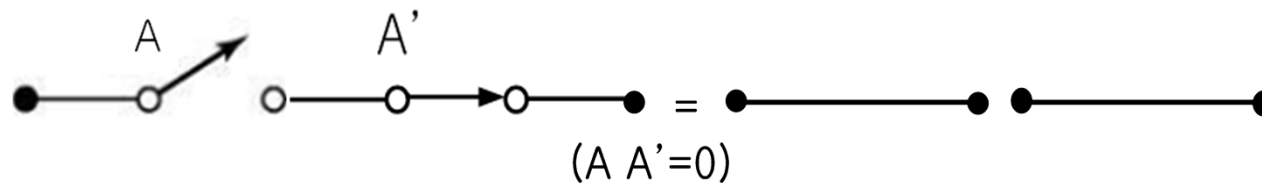
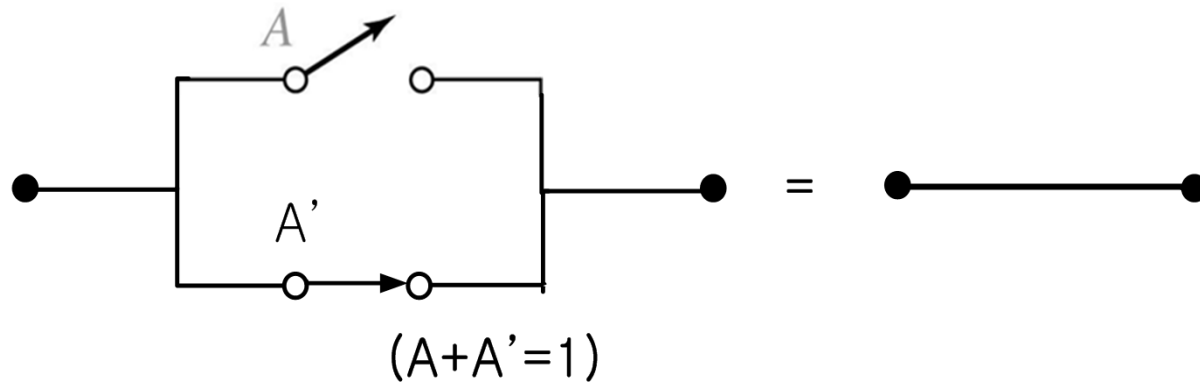
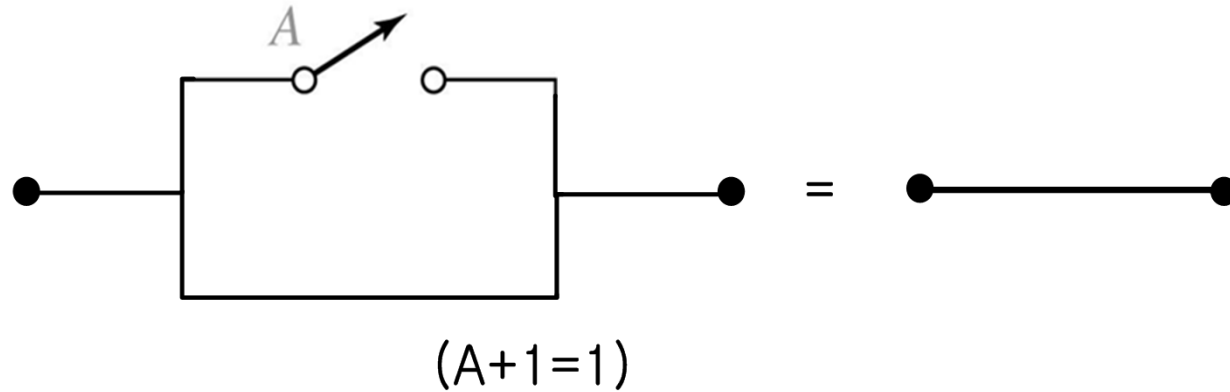
$$(AB' + D)E + 1 = 1$$

$$(AB' + D)(AB' + D)' = 0$$

## 2.4 Basic Theorems with Switch Circuits



## 2.4 Basic Theorems with Switch Circuits



## 2.5 Commutative, Associative, and Distributive Laws

Commutative Laws:

$$XY = YX$$

$$X + Y = Y + X$$

Associative Laws:

$$(XY)Z = X(YZ) = XYZ$$

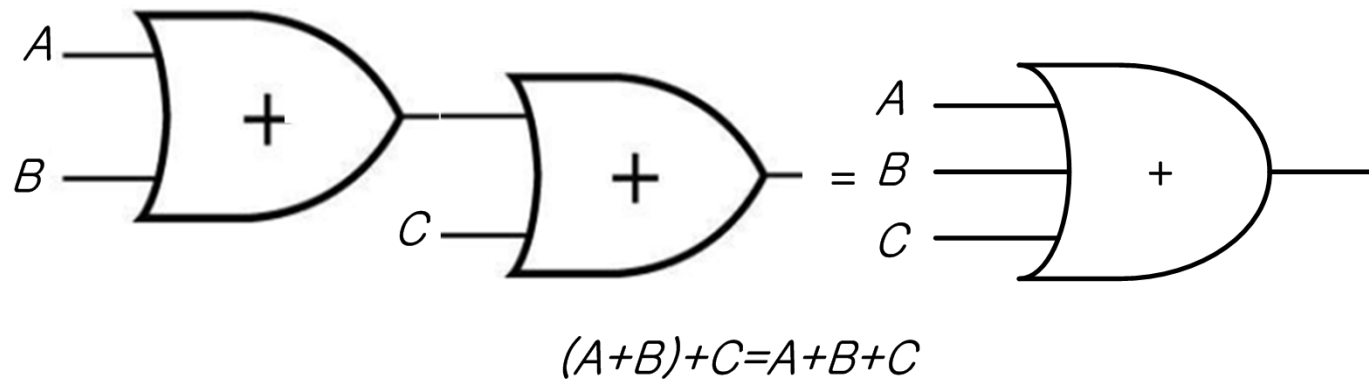
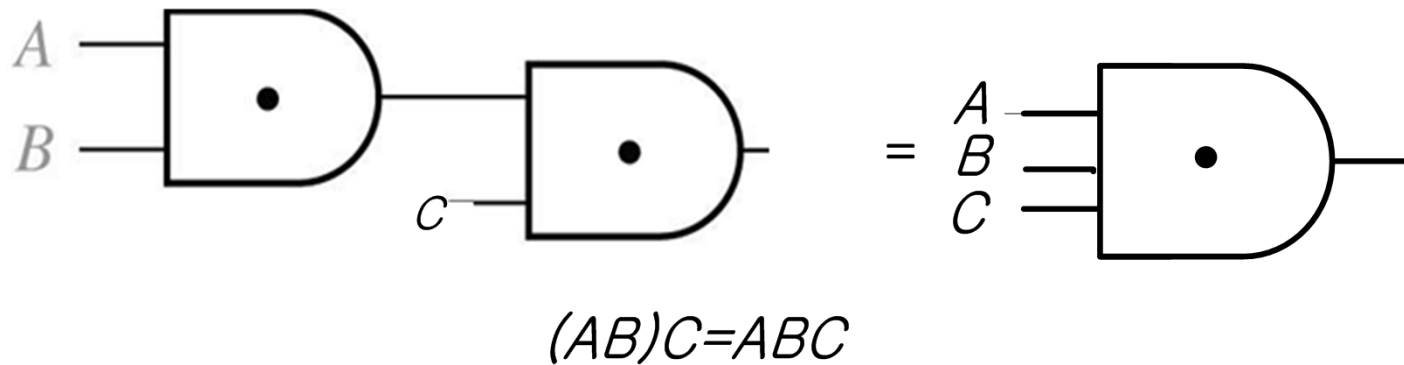
$$(X + Y) + Z = X + (Y + Z) = X + Y + Z$$

Proof of Associate Law for AND

X	Y	Z	XY	YZ	(XY)Z	X(YZ)
0	0	0	0	0	0	0
0	0	1	0	0	0	0
0	1	0	0	0	0	0
0	1	1	0	1	0	0
1	0	0	0	0	0	0
1	0	1	0	0	0	0
1	1	0	1	0	0	0
1	1	1	1	1	1	1

# Associative Laws for AND and OR

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## 2.5 Commutative, Associative, and Distributive Laws

**AND**  $XYZ = 1$  iff  $X = Y = Z = 1$

**OR**  $X + Y + Z = 0$  iff  $X = Y = Z = 0$

Distributive Laws:

$$X(Y + Z) = XY + XZ$$

$$X + YZ = (X + Y)(X + Z)$$

*Valid only Boolean algebra not for ordinary algebra*

자주 활용됨

**Proof**

$$\begin{aligned}(X + Y)(X + Z) &= X(X + Z) + Y(X + Z) = XX + XZ + YX + YZ \\ &= X + XZ + XY + YZ = X \cdot 1 + XZ + XY + YZ \\ &= X(1 + Z + Y) + YZ = X \cdot 1 + YZ = X + YZ\end{aligned}$$

## 2.6 Simplification Theorems

### Useful Theorems for Simplification

$$XY + XY' = X$$

$$(X + Y)(X + Y') = X$$

$$X + XY = X$$

$$X(X + Y) = X$$

$$(X + Y')Y = XY$$

$$XY' + Y = X + Y$$

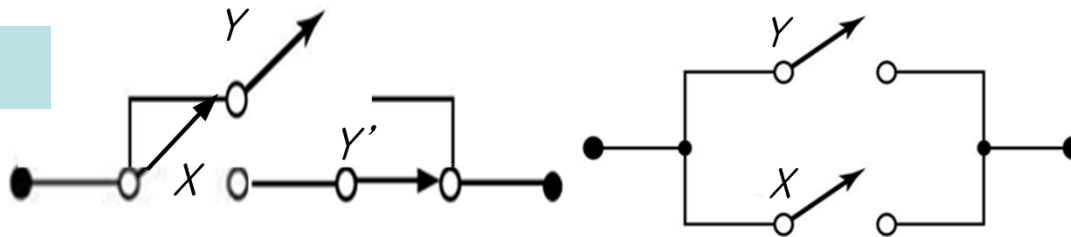
### Proof

$$X + XY = X \cdot 1 + XY = X(1 + Y) = X \cdot 1 = X$$

$$X(X + Y) = XX + XY = X + XY = X$$

$$Y + XY' = (Y + X)(Y + Y') = (Y + X)1 = Y + X$$

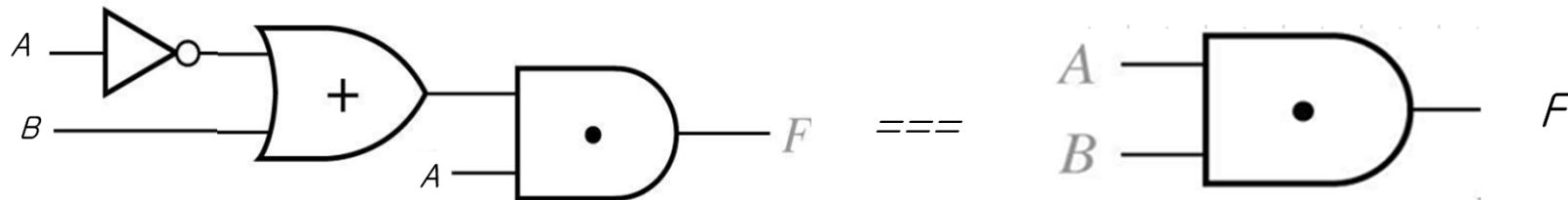
### Proof with Switch



## 2.6 Simplification Theorems

### Equivalent Gate Circuits

$$F = A(A' + B) = AB$$





## 2.7 Multiplying Out and Factoring

To obtain a sum-of-product form → Multiplying out using distributive laws

Sum of product form:  $AB' + CD'E + AC'E$

Still considered to be in sum of product form:  $ABC' + DEFG + H$   
 $A + B' + C + D'E$

Not in Sum of product form:  $(A + B)CD + EF$

Multiplying out and eliminating redundant terms

$$\begin{aligned}(A + BC)(A + D + E) &= A + AD + AE + ABC + BCD + BCE \\ &= A(1 + D + E + BC) + BCD + BCE \\ &= A + BCD + BCE\end{aligned}$$

## 2.7 Multiplying Out and Factoring

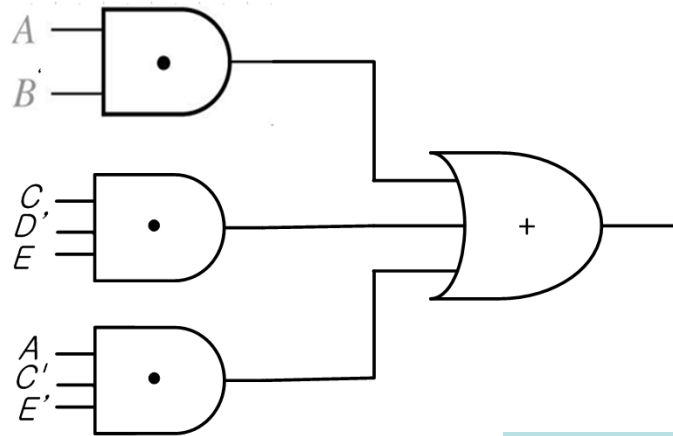
To obtain a product of sum form  $\rightarrow$  all sums are the sum of single variable

Product of sum form:  $(A + B')(C + D' + E)(A + C' + E')$

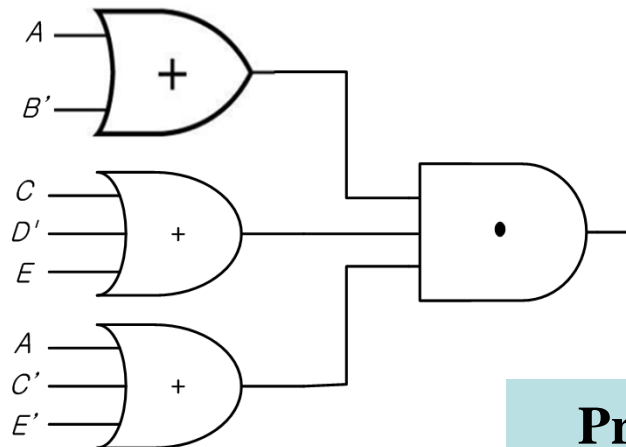
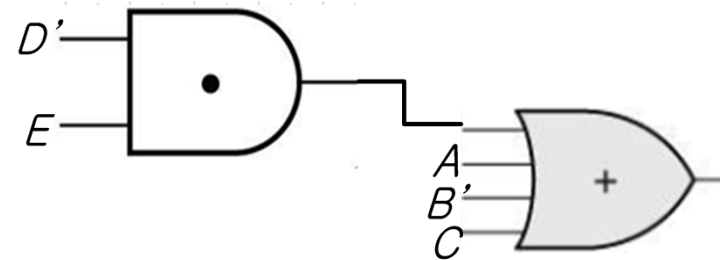
Still considered to  
be in product of  
sum form:

$$\begin{array}{l} (A + B)(C + D + E)F \\ AB'C(D' + E) \end{array}$$

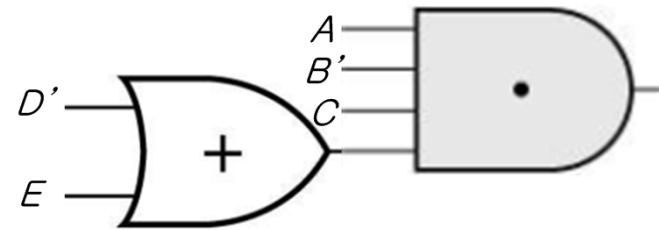
# Circuits for SOP and POS form



**Sum of product form:**



**Product of sum form:**



## 2.8 DeMorgan's Laws

### DeMorgan's Laws

$$(X + Y)' = X'Y'$$

$$(XY)' = X' + Y'$$

### Proof

X	Y	X'	Y'	X + Y	(X + Y)'	X'	Y'	XY	(XY)'	X' + Y'
0	0	1	1	0	1	1	1	0	1	1
0	1	1	0	1	0	0	0	0	1	1
1	0	0	1	1	0	0	0	0	1	1
1	1	0	0	1	0	0	0	1	0	0

### DeMorgan's Laws for $n$ variables

$$(X_1 + X_2 + X_3 + \dots + X_n)' = X_1' X_2' X_3' \dots X_n'$$

$$(X_1 X_2 X_3 \dots X_n)' = X_1' + X_2' + X_3' + \dots + X_n'$$

### Example

$$(X_1 + X_2 + X_3)' = (X_1 + X_2)' X_3' = X_1' X_2' X_3'$$

## 2.8 DeMorgan's Laws

**Inverse of**  
 $F = A'B + AB'$

$$F' = (A'B + AB')' = (A'B)'(AB')' = (A + B')(A' + B) \\ = AA' + AB + B'A' + BB' = A'B' + AB$$

A B	A' B	A B'	F = A'B + AB'	A' B'	A B	F' = A'B' + AB
0 0	0	0	0	1	0	1
0 1	1	0	1	0	0	0
1 0	0	1	1	0	0	0
1 1	0	0	0	0	1	1

**Dual:** 'dual' is formed by replacing AND with OR, OR with AND, 0 with 1, 1 with 0

$$(XYZ...) ^D = X + Y + Z + ... \quad (X + Y + Z + ...) ^D = XYZ...$$

The dual of an expression may be found by complementing the entire expression and then complementing each individual variable

$$(AB' + C)' = (AB')' C' = (A' + B) C', \quad \text{so} \quad (AB' + C)^D = (A + B') C \quad 30/30$$