# گزارش پروژه چهارم درس شبکههای کامپیوتری دکتر یزدانی بهار ۱۴۰۱ بهار ۱۴۰۱ پرنیان فاضل ۸۱۰۱۹۸۵۱۶ - پریا خوش تاب ۸۱۰۱۹۸۵۱۶

The purpose of this exercise is to get acquainted with the function of **TCP** and to implement the mechanism of sending correct information along with congestion control using UDP sockets in the network.

# **Implementation**

Our project consists of the following \*.c and \*.h files:

 packet.c / packet.h: These files are mainly used for creating tcp packets with header and space for data of a given size.

### Packet.h:

```
enum packet_type {
    DATA,
    ACK,
};

typedef struct {
    int dest;
    int src;
    int seqno;
    int ackno;
    int ctr_flags;
    int data_size;
}tcp_header;

#define MSS_SIZE 1500
#define UDP_HDR_SIZE 8
```

Note that the size of the packets that are sent is 1500 bytes(MSS\_SIZE), so if we subtract the size of the headers from it, the size of the pure data that is sent each time becomes 1448 bytes (DATA\_SIZE).

```
1500 - 20 - 8 - (4*6) = 1448 byte
```

#### packet.c:

```
#include"packet.h"

static tcp_packet zero_packet = {.hdr={0}};

/*
    * create tcp packet with header and space for data of size len
    */
tcp_packet* make_packet(int len)
{
        tcp_packet *pkt;
        pkt = malloc(TCP_HDR_SIZE + len);

        *pkt = zero_packet;
        pkt->hdr.data_size = len;
        return pkt;
}

int get_data_size(tcp_packet *pkt)
{
        return pkt->hdr.data_size;
}
```

• queue.c / queue.h: In these files, we have implemented a FIFO queue data structure from scratch, because the C language does not support the queue library.

## Queue.h:

```
#define QUEUE H
 NODE;
  NODE *tail; // tail of queue
QUEUE *queueCreate(void);
```

```
0 on success, -1 on failure
int enqueue (QUEUE *, void *);
nt dequeue(QUEUE *);
void *peek(QUEUE *);
#endif // QUEUE H
```

#### queue.c:

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#include "queue.h"
```

```
QUEUE *queueCreate()
    // allocate space for a queue struct
   QUEUE *queue = malloc(sizeof(QUEUE));
   if ( queue == NULL )
        fprintf(stderr, "queueCreate(): malloc error\n");
       return NULL;
    }
    // set default values
   queue->head = NULL;
   queue->tail = NULL;
   queue->size = 0;
    return queue;
int enqueue(QUEUE *q, void *item)
   NODE *node;
   // set up new node
   node = malloc(sizeof(NODE));
   if ( node == NULL )
        fprintf(stderr, "enqueue(): malloc error\n");
       return -1;
   else
    {
       node->next = NULL;
       node->item = item;
    }
    // queue is empty
   if ( q->size == 0 )
    {
       q->head = node;
       q->tail = node;
```

```
q->size = 1;
   // queue is non-empty
   else
    {
       q->tail->next = node;
       q->tail = node;
       q->size++;
    }
   return 0;
int dequeue(QUEUE *q)
   // queue is empty
   if (q->size == 0)
        fprintf(stderr, "dequeue(): cannot remove item from empty
queue\n");
       return -1;
    }
    // queue contains 1 item
   if ( q->size == 1 )
    {
       q->head = NULL;
       q->tail = NULL;
       q->size = 0;
   // queue contains multiple items
   else
       q->head = q->head->next;
       q->size--;
   return 0;
void *peek(QUEUE *q)
    // queue is empty
```

```
if ( q->size == 0 )
{
    fprintf(stderr, "peek(): queue is empty\n");
    return NULL;
}
// queue is non-empty
return q->head->item;
}
int size(QUEUE *q)
{
    return q->size;
}
```

• common.c / common.h: In these files, we have implemented a wrapper for perror.

#### common.h:

```
#define NONE     0x0
#define INFO     0x1
#define WARNING 0x10
#define INFO     0x1
#define DEBUG     0x100
#define ALL      0x111

#define VLOG(level, ...) \
    if(level & verbose) { \
        fprintf(stderr, ##__VA_ARGS___); \
        fprintf(stderr, "\n"); \
    }\

void error(char *msg);
#endif
```

#### common.c:

```
#include <stdio.h>
#include"common.h"

int verbose = ALL;

void error(char *msg) {
    perror(msg);
    exit(1);
}
```

• sender.c: For reliability, we implement a sliding window on the sender side. We keep a limit on the max number of packets that can be sent and not ACKed at any given time. If we receive back an ACK with sequence number higher than our current send base, then we slide the window forward and transmit new packets. Of course, we are also keeping a timer on the earliest unACKed packet. There can be a cumulative ACK which will indicate all the earlier packets have been successfully received. If there is a time-out, we retransmit the missing packet(s).

```
#define STDIN_FD 0
#define RETRY 200 //milli second
#define PORT_ROUTER 8080
#define PORT_B 8081

int transmission_round = 0; //keep track of the round
int next_seqno; // next byte to send
int exp_seqno; // expected byte to be acked
int send_base = 0; // first byte in the window
float window_size = 50; // window size at the beginning of slow
start
int portno;
int timer_on = 0;
```

```
FILE *fp;
int sockfd, serverlen;
struct sockaddr in serveraddr, senderAddr;
tcp packet *recvpkt;
sigset t sigmask;
void send packet(char* buffer, int len, int seqno);
void resend_packets(int sig);
void init timer(int delay, void (*sig handler)(int));
void start timer();
void stop timer();
void send packet(char* buffer, int len, int seqno) {
  tcp_packet *sndpkt = make_packet(len);
  memcpy(sndpkt->data, buffer, len);
  sndpkt->hdr.seqno = seqno;
  sndpkt->hdr.src = portno;
  sndpkt->hdr.dest = PORT B;
  printf("Sending packet of sequence number %d of data size %d
to %s\n", seqno, len, inet ntoa(serveraddr.sin addr));
   if (sendto(sockfd, sndpkt, TCP HDR SIZE +
get data size(sndpkt), 0,
                      (const struct sockaddr *)&serveraddr,
serverlen) < 0)
      error("sendto");
  free(sndpkt);
void resend packets(int sig)
  char buffer[DATA SIZE];
  if (sig == SIGALRM)
      VLOG(INFO, "Timeout happened");
```

```
for (int i = send base; i < next seqno; i += DATA SIZE)</pre>
          fseek(fp, i, SEEK SET);
          len = fread(buffer, 1, DATA SIZE, fp);
          send packet(buffer, len, i);
void start timer()
  sigprocmask(SIG UNBLOCK, &sigmask, NULL);
  setitimer(ITIMER REAL, &timer, NULL);
void stop timer()
  sigprocmask(SIG BLOCK, &sigmask, NULL);
void init timer(int delay, void (*sig handler)(int))
  signal(SIGALRM, resend packets);
  timer.it interval.tv sec = delay / 1000; // sets an interval
  timer.it interval.tv usec = (delay % 1000) * 1000;
```

```
timer.it_value.tv_usec = (delay % 1000) * 1000;
  sigemptyset(&sigmask);
  sigaddset(&sigmask, SIGALRM);
int main(int argc, char **argv)
  char *hostname;
  char buffer[DATA SIZE];
  int dup_cnt; // count of continuous duplicate ACKs
  FILE* plot; // file that contains plot of window size,
  struct timeval start, now;
  if (argc != 4)
      fprintf(stderr, "usage: %s <hostname> <port> <FILE>\n",
argv[0]);
      exit(0);
  hostname = argv[1];
  portno = atoi(argv[2]);
  fp = fopen(argv[3], "r");
     error(argv[3]);
  if (sockfd < 0)
      error("ERROR opening socket");
  memset(&senderAddr, 0, sizeof(senderAddr));
  senderAddr.sin family = AF INET; // IPv4
  senderAddr.sin addr.s addr = INADDR ANY;
  senderAddr.sin port = htons(portno);
```

```
if (bind(sockfd, (const struct sockaddr *)&senderAddr,
sizeof(senderAddr)) < 0){</pre>
      perror("bind failed");
      exit(EXIT FAILURE);
  bzero((char *)&serveraddr, sizeof(serveraddr));
  serverlen = sizeof(serveraddr);
      fprintf(stderr, "ERROR, invalid host %s\n", hostname);
      exit(0);
  serveraddr.sin port = htons(PORT ROUTER);
  init_timer(RETRY, resend_packets);
  next_seqno = 0;
  exp seqno = DATA SIZE;
  plot = fopen("CWND.csv", "w+");
  gettimeofday(&start,NULL);
  while (1)
      gettimeofday(&now, NULL);
```

```
while (next seqno < send base + (int) (window size)</pre>
    fseek(fp, next_seqno, SEEK_SET);
    len = fread(buffer, 1, DATA SIZE, fp);
    if (len <= 0)
        send packet(buffer, 0, 0);
    send packet(buffer, len, next seqno);
    next seqno += len;
int n = recvfrom(sockfd, buffer, MSS SIZE, MSG WAITALL,
recvpkt = (tcp packet *)buffer;
assert(get_data_size(recvpkt) <= DATA SIZE);</pre>
if (recvpkt->hdr.ackno % DATA SIZE > 0) {
    printf("***All packets sent successfully\n");
    double time taken = ((double)duration)/CLOCKS PER SEC;
```

• receiver.c: On the receiver side, we have an out-of-order buffer that contains all out-of-order packets (there is a limit for this buffer, if it's full then all other out-of-order packets will be dropped). After receiving an out-of-order packet, the receiver immediately sends back a duplicate ACK to indicate packet loss.

```
FILE *fp;
pthread mutex t m;
int dest;
int src;
void write to file(char* file name, FILE* fp, int pos, char*
data, int len);
void send ACK(int ackno);
void write to file(char* file name, FILE* fp, int pos, char*
data, int len) {
  strcat(file name, ".txt");
  fp = fopen(file name, "a+");
  fseek(fp, pos, SEEK SET);
  pthread mutex lock(&m);
  fwrite(data, 1, len, fp);
  pthread_mutex_unlock(&m);
  fclose(fp);
void send ACK(int ackno) {
  tcp packet *sndpkt = make packet(0);
  sndpkt->hdr.ackno = ackno;
  sndpkt->hdr.dest = dest;
  sndpkt->hdr.src = src;
  sndpkt->hdr.ctr flags = ACK;
  if (sendto(sockfd, sndpkt, TCP HDR SIZE, MSG CONFIRM,
      error("ERROR in sendto");
int main(int argc, char **argv)
  int optval;
  tcp_packet *recvpkt;
  char* file name;
  char buffer[MSS SIZE];
  struct timeval tp;
```

```
int cur_seqno;
int buffer size = 4;
tcp packet* buffer pkts[4]; // buffer to store out-of-order
if (argc != 3)
    fprintf(stderr, "usage: %s <port> FILE RECVD\n", argv[0]);
   exit(1);
portno = atoi(argv[1]);
file name = argv[2];
fp = fopen(file name, "w");
    error(argv[2]);
sockfd = socket(AF INET, SOCK DGRAM, 0);
    error("ERROR opening socket");
optval = 1;
setsockopt(sockfd, SOL SOCKET, SO REUSEADDR,
           (const void *)&optval, sizeof(int));
```

```
bzero((char *)&serveraddr, sizeof(serveraddr));
  serveraddr.sin family = AF INET;
  serveraddr.sin addr.s addr = INADDR ANY;
  serveraddr.sin port = htons(PORT B);
  clientaddr.sin port = htons(PORT ROUTER);
  if (bind(sockfd, (struct sockaddr *)&serveraddr,
            sizeof(serveraddr)) < 0)</pre>
      error("ERROR on binding");
  VLOG(DEBUG, "epoch time, bytes received, sequence number");
  clientlen = sizeof(clientaddr);
  cur seqno = 0;
  while (1)
(struct sockaddr *)&clientaddr, (socklen t *)&clientlen) < 0)</pre>
```

```
recvpkt = (tcp packet *)buffer;
       assert(get_data_size(recvpkt) <= DATA_SIZE);</pre>
       if (recvpkt->hdr.data size == 0)
       dest = recvpkt->hdr.src;
       src = recvpkt->hdr.dest;
       gettimeofday(&tp, NULL);
       VLOG (DEBUG, "%lu, %d, %d", tp.tv sec,
recvpkt->hdr.data size, recvpkt->hdr.seqno);
       if (recvpkt->hdr.seqno - cur seqno <= DATA SIZE)</pre>
           cur seqno = recvpkt->hdr.seqno; // update current
           char name[20];
           sprintf(name, "%d", dest);
           write to file(name, fp, recvpkt->hdr.seqno,
recvpkt->data, recvpkt->hdr.data size);
           ackno = recvpkt->hdr.seqno + recvpkt->hdr.data size;
           int new seqno = cur seqno + recvpkt->hdr.data size;
               if (new seqno == buffer pkts[i]->hdr.seqno) {
                   cur seqno = new seqno;
                   char name[20];
                   sprintf(name, "%d", dest);
                   write to file (name, fp, cur seqno,
buffer pkts[i]->data, buffer pkts[i]->hdr.data size);
```

```
new seqno += get_data_size(buffer_pkts[i]);
        ackno = new seqno;
    for (int i = cnt; i < buffer_size; i++) {</pre>
        buffer pkts[i-cnt] = buffer_pkts[i];
send ACK(ackno);
    buffer pkts[ind] = recvpkt;
```

• router.c: To implement the router, we use two threads, one of which is used to send data and one of which is used to receive data from the hosts. In fact, when receiving data, we push the received packet to the queue, and when sending data, we pop the first element from the queue, and depending on the type of the packet (ACK or DATA), we send the data to the desired host(A or B). Note that in order

to find the host (A or B) we want to send the data to, we have added src and dest fields to the tcp\_header struct.

```
pthread_mutex_t bufferMutex;
#define PORT ROUTER 8080
#define PORT B 8081
#define MAXLINE 1500
#define DROP PERCENTAGE 10
int sockfd, sockfd2;
double maxp = 0.02;
double tempP = 0;
double avg = 0;
double weight = 0.003;
int minThreshold = 5, maxThreshold = 20;
struct sockaddr in servaddr;
QUEUE* routerBuffer;
int random drop(int percentage) {
  int randomP = (rand()%100);
  if (randomP < percentage) {</pre>
int randomEarlyDetection(int s) {
  avg = ((1 - weight) * avg) + (weight * size(routerBuffer));
  if (minThreshold <= avg && avg < maxThreshold) {</pre>
       Count++;
```

```
tempP = ((avg - minThreshold) * maxp)/(maxThreshold -
minThreshold);
       double P = tempP/(1 - (count * tempP));
           P = 1.0;
       double randomP = (rand()%100)/100.00;
       if(randomP <= P) {</pre>
           if(count != 50) {
              res = -1;
           res = 0;
           count = -1;
   } else if(maxThreshold <= avg) {</pre>
      res = -1;
      count = 0;
      res = 0;
      count = -1;
   return res;
void* recieve(void * x) {
      char buffer[MAXLINE];
       len = sizeof(servaddr);
       n = recvfrom(sockfd, buffer, MAXLINE, MSG WAITALL, (
       srand(time(NULL));
           tcp packet *recvpkt = (tcp packet *)buffer;
           pthread mutex lock(&bufferMutex);
           enqueue (routerBuffer, recvpkt);
           pthread mutex unlock(&bufferMutex);
```

```
void *sendHandler(void *x) {
  while(1){
       if(size(routerBuffer) != 0) {
           pthread mutex lock(&bufferMutex);
           tcp packet *recvpkt = peek(routerBuffer);
           dequeue (routerBuffer);
           if(random drop(DROP PERCENTAGE) == 0 /*&&
               int port = recvpkt->hdr.dest;
               memset(&clientaddr, 0, sizeof(clientaddr));
               clientaddr.sin family = AF INET;
               clientaddr.sin port = htons(port);
               int len = sizeof(clientaddr);
               if(recvpkt->hdr.ctr flags == ACK) {
                   sendto(sockfd, recvpkt, TCP HDR SIZE,
MSG CONFIRM, (const struct sockaddr *) &clientaddr, len);
                   sendto(sockfd, recvpkt, TCP HDR SIZE +
get data size(recvpkt), MSG CONFIRM, (const struct sockaddr *)
&clientaddr, len);
           pthread mutex unlock(&bufferMutex);
int main() {
  routerBuffer = queueCreate();
  if ( (sockfd = socket(AF INET, SOCK DGRAM, 0)) < 0 ) {
      perror("socket creation failed");
      exit(EXIT FAILURE);
  if ( (sockfd2 = socket(AF INET, SOCK DGRAM, 0)) < 0 ) {</pre>
      perror("socket2 creation failed");
      exit(EXIT FAILURE);
  memset(&servaddr, 0, sizeof(servaddr));
```

```
servaddr.sin_family = AF_INET; // IPv4
servaddr.sin_addr.s_addr = INADDR_ANY;
servaddr.sin_port = htons(PORT_ROUTER);

// Bind the socket with the server address
if (bind(sockfd, (const struct sockaddr *)&servaddr,
sizeof(servaddr)) < 0) {
    perror("bind failed");
    exit(EXIT_FAILURE);
}

pthread_t thread_receiver_A, thread_receiver_B, thread_sender;
pthread_create(&thread_receiver_A, NULL, sendHandler, NULL);
pthread_create(&thread_receiver_B, NULL, recieve, NULL);
pthread_join(thread_receiver_A, NULL);
pthread_join(thread_receiver_B, NULL);
return 0;
}</pre>
```

# Part 1

1)

# Go Back N Protocol(GBN):

- Go-Back-N ARQ is a specific instance of the automatic repeat request (ARQ)
  protocol, in which the sending process continues to send a number of frames
  specified by a window size even without receiving an acknowledgement (ACK)
  packet from the receiver.
- It uses the principle of protocol pipelining in which multiple frames can be sent before receiving the acknowledgment of the first frame. If we have five frames and the concept is Go-Back-3, which means that the three frames can be sent, i.e.,

- frame no 1, frame no 2, frame no 3 can be sent before expecting the acknowledgment of frame no 1.
- In Go-Back-N ARQ, the frames are numbered sequentially as Go-Back-N ARQ sends the multiple frames at a time that requires the numbering approach to distinguish the frame from another frame, and these numbers are known as the sequential numbers.
- The number of frames that can be sent at a time totally depends on the size of the sender's window. So, we can say that 'N' is the number of frames that can be sent at a time before receiving the acknowledgment from the receiver.
- If the acknowledgment of a frame is not received within an agreed-upon time period, then all the frames available in the current window will be retransmitted. Suppose we have sent the frame no 5, but we didn't receive the acknowledgment of frame no 5, and the current window is holding three frames, then these three frames will be retransmitted.
- The sequence number of the outbound frames depends upon the size of the sender's window.

# Selective Repeat Protocol(SR):

• Selective repeat protocol, also called Selective Repeat ARQ (Automatic Repeat reQuest), is a data link layer protocol that uses a **sliding window** method for reliable delivery of data frames. Here, only the erroneous or lost frames are retransmitted, while the good frames are received and buffered.

- It provides for sending multiple frames depending upon the availability of frames in the sending window, even if it does not receive acknowledgement for any frame in the interim.
- If the receiver receives a corrupt frame, it does not directly discard it. It sends a negative acknowledgement to the sender. The sender sends that frame again as soon as on the receiving negative acknowledgment. There is no waiting for any time-out to send that frame.
- It is mainly used because we need the receiver to be able to accept packets out-of-order using buffer space, for a superior protocol to combine advantages of both Stop-Wait and GBN. Selective Repeat attempts to retransmit only those packets that are actually lost (due to errors).
- The maximum number of frames that can be sent depends upon the size of the sending window.
- Individual acknowledgements are used in Selective Repeat Protocol

When Buffer Space is of more concern than bandwidth then Go Back N Protocol is used, if Bandwidth is of more concern than buffer space then Selective Repeat Protocol is used. In Go-Back-N we have Less complexity, less CPU cycles while for Selective Repeat Protocol More processing power, cpu cycles at receiver. If error rate is low, use Go-back-N else if error rate is high use Selective Repeat Protocol.

# **GBN** vs SR:

GBN	SR
In Go-Back-N Protocol, if the sent frame	In selective Repeat protocol, only those
are find suspected then all the frames are	frames are re-transmitted which are found
re-transmitted from the lost packet to the	suspected.
last packet transmitted.	
Receiver window size of Go-Back-N	Receiver window size of selective Repeat
Protocol is 1.	protocol is N.
In Go-Back-N Protocol, neither sender nor	In selective Repeat protocol, receiver side
at receiver need sorting.	needs sorting to sort the frames.
Go-Back-N Protocol is less complex.	selective Repeat protocol is more
	complex.
In Go-Back-N Protocol, Out-of-Order	In selective Repeat protocol, Out-of-Order
packets are NOT Accepted (discarded) and	packets are Accepted.
the entire window is re-transmitted.	
If error rate is high, it wastes a lot of	Comparatively less bandwidth is wasted in
bandwidth.	retransmitting.
Receiver do not store the frames received	Receiver stores the frames received after
after the damaged frame until the damaged	the damaged frame in the buffer until the
frame is retransmitted.	damaged frame is replaced.

- **2)** In this project, we use GBN protocol for Sliding Window to provide reliable data transfer. The steps are shown below:
  - For the first task, we set a fixed cwnd=10.
  - For the logic, for every packet that the sender sends, it contains sequo which is the number of the first byte of the data from the data stream. The fixed MSS size is
     1.5 KB.
  - On the receiver side, we keep track of the expected sequo number of the packet to check if the next packet we receive is in-order to the data stream. Immediately after receiving the packet, we send back an ACK which tells the sender that we have received the packets. Even if our ACK gets lost, higher ACK seq will still confirm that earlier packets have been received. If we receive an out-of-order packet, we discard it and wait for retransmission.
  - On the router side, We receive the data from the sender (A) and send it to the receiver (B). We also receive the ACK from the receiver (B) and send it to the sender (A).
  - On the sender side, we check the ACK that we receive. If the ACK is larger than our sendbase (meaning former packets received), then we increase our sendbase to the next seque number right after the last packet received. If we timeout before receiving ACK, we restransmit all the packets starting from the sendbase.

3)

• 200KB file:

```
Sending packet of sequence number 431504 of data size 1448 to 127.0.0.1
Sending packet of sequence number 432952 of data size 1448 to 127.0.0.1
Sending packet of sequence number 434400 of data size 1448 to 127.0.0.1 Sending packet of sequence number 435848 of data size 1448 to 127.0.0.1
Sending packet of sequence number 437296 of data size 1448 to 127.0.0.1
Sending packet of sequence number 438744 of data size 1448 to 127.0.0.1
Sending packet of sequence number 440192 of data size 1448 to 127.0.0.1
Sending packet of sequence number 441640 of data size 1448 to 127.0.0.1
Sending packet of sequence number 443088 of data size 1448 to 127.0.0.1
Sending packet of sequence number 444536 of data size 1448 to 127.0.0.1 Sending packet of sequence number 445984 of data size 1448 to 127.0.0.1
Sending packet of sequence number 447432 of data size 1448 to 127.0.0.1
Sending packet of sequence number 448880 of data size 1448 to 127.0.0.1
Sending packet of sequence number 450328 of data size 1448 to 127.0.0.1
Sending packet of sequence number 451776 of data size 1448 to 127.0.0.1
Sending packet of sequence number 453224 of data size 1448 to 127.0.0.1
Sending packet of sequence number 454672 of data size 1448 to 127.0.0.1
Sending packet of sequence number 456120 of data size 1448 to 127.0.0.1
Sending packet of sequence number 457568 of data size 1448 to 127.0.0.1
Sending packet of sequence number 459016 of data size 1448 to 127.0.0.1
Sending packet of sequence number 460464 of data size 1448 to 127.0.0.1
Sending packet of sequence number 461912 of data size 1448 to 127.0.0.1
Sending packet of sequence number 463360 of data size 1448 to 127.0.0.1
Sending packet of sequence number 464808 of data size 1448 to 127.0.0.1
Sending packet of sequence number 466256 of data size 1448 to 127.0.0.1
Sending packet of sequence number 467704 of data size 1448 to 127.0.0.1
Sending packet of sequence number 469152 of data size 1448 to 127.0.0.1
Sending packet of sequence number 470600 of data size 1448 to 127.0.0.1
Sending packet of sequence number 472048 of data size 1448 to 127.0.0.1
Sending packet of sequence number 473496 of data size 1448 to 127.0.0.1
Sending packet of sequence number 474944 of data size 1448 to 127.0.0.1
Sending packet of sequence number 476392 of data size 1448 to 127.0.0.1
Sending packet of sequence number 477840 of data size 1448 to 127.0.0.1
Sending packet of sequence number 479288 of data size 1448 to 127.0.0.1
Sending packet of sequence number 480736 of data size 1448 to 127.0.0.1
Sending packet of sequence number 482184 of data size 1448 to 127.0.0.1
Sending packet of sequence number 483632 of data size 1448 to 127.0.0.1
Sending packet of sequence number 485080 of data size 1448 to 127.0.0.1
Sending packet of sequence number 486528 of data size 1108 to 127.0.0.1
***All packets sent successfully
                                      ----->Time taken: 0.039811 secconds
```

#### • 1MB file:

```
End Of File read
Sending packet of sequence number 0 of data size 0 to 127.0.0.1
***All packets sent successfully
----->Time taken: 0.537036 secconds
```

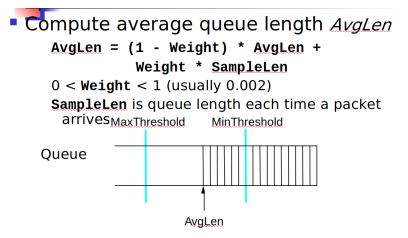
**4)** Size of buffer = 10 & 1MB file

# Part 2

1)

#### Random Early Drop

rather than wait for queue to become full, drop each arriving packet with some drop probability whenever the queue length exceeds some drop level.



- P of a particular flow's packet(s) is roughly proportional to the share of the flow's bandwidth.
- MaxP is typically 0.02, meaning when is halfway between the two thresholds,
   router drops roughly one out of 50 packets.

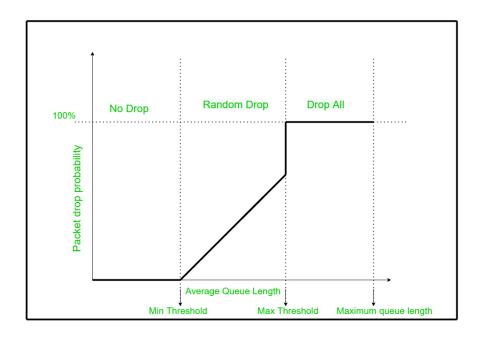
- If traffic is bursty, then MinTh. should be sufficiently large to allow link utilization to be acceptably high.
- Difference between two thresholds should be larger than the typical increase in the calculated average queue length in one RTT; setting MaxThreshold to twice
   MinThreshold is reasonable.

Here is the algorithm of RED:

Two queue length thresholds

if AvgLen <= MinThreshold then
enqueue the packet

if MinThreshold < AvgLen < MaxThreshold
then
calculate probability P
drop arriving packet with probability P
if ManThreshold <= AvgLen then
drop arriving packet



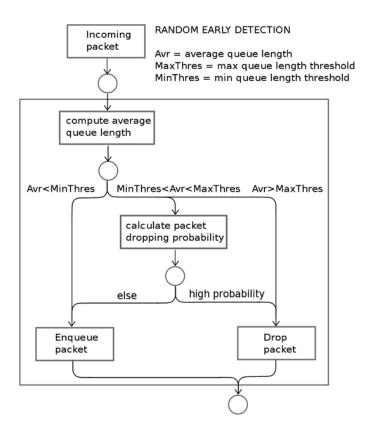
RED monitors the average queue size and drops packets based on statistical probabilities.

If the buffer is almost empty then all incoming packets are accepted. As the queue grows,

the probability for dropping an incoming packet grows too. When the buffer is full, the probability has reached 1 and all incoming packets are dropped.

RED is more fair than tail drop, in the sense that it does not possess a bias against bursty traffic that uses only a small portion of the bandwidth. The more a host transmits, the more likely it is that its packets are dropped as the probability of a host's packet being dropped is proportional to the amount of data it has in a queue. Early detection helps avoid TCP global synchronization.

To sum up:



```
int randomEarlyDetection(int s) {
  int res = 0;
  avg = ((1 - weight) * avg) + (weight * size(routerBuffer));

if(minThreshold <= avg && avg < maxThreshold) {
    count++;</pre>
```

```
tempP = ((avg - minThreshold) * maxp)/(maxThreshold -
minThreshold);
       double P = tempP/(1 - (count * tempP));
       if(count == 50) {
       if(randomP <= P) {</pre>
              res = -1;
           count = 0;
          res = 0;
           count = -1;
   } else if(maxThreshold <= avg) {</pre>
      res = -1;
      count = 0;
      res = 0;
      count = -1;
```

### Assumptions:

```
double maxp = 0.02;
double tempP = 0;
double avg = 0;
int count = -1;
double weight = 0.003;
int minThreshold = 5, maxThreshold = 20;
```

2)

With 3 computers:

----->Time taken: 0.029933 secconds

Total time: 0.507361 s

# References:

https://www.icir.org/floyd/papers/early.twocolumn.pdf

https://techdifferences.com/difference-between-go-back-n-and-selective-repeat-protocol.ht

<u>ml</u>

https://www.geeksforgeeks.org/difference-between-go-back-n-and-selective-repeat-protoco

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