

SEGSQRSS - Sum of Squares with Segment Tree

#tree

Segment trees are extremely useful. In particular "Lazy Propagation" (i.e. [see here, for example](#)) allows one to compute sums over a range in $O(\lg(n))$, and update ranges in $O(\lg(n))$ as well. In this problem you will compute something much harder:

The sum of squares over a range with range updates of 2 types:

1) increment in a range

2) set all numbers the same in a range.

Input

There will be **T** ($T \leq 25$) test cases in the input file. First line of the input contains two positive integers, **N** ($N \leq 100,000$) and **Q** ($Q \leq 100,000$). The next line contains **N** integers, each at most 1000. Each of the next **Q** lines starts with a number, which indicates the type of operation:

2 st nd -- return the sum of the squares of the numbers with indices in **[st, nd]** {i.e., from **st** to **nd** inclusive} ($1 \leq st \leq nd \leq N$).

1 st nd x -- add "**x**" to all numbers with indices in **[st, nd]** ($1 \leq st \leq nd \leq N$, and $-1,000 \leq x \leq 1,000$).

0 st nd x -- set all numbers with indices in **[st, nd]** to "**x**" ($1 \leq st \leq nd \leq N$, and $-1,000 \leq x \leq 1,000$).

Output

For each test case output the “Case <caseno>:” in the first line and from the second line output the sum of squares for each operation of type 2. Intermediate overflow will not occur with proper use of 64-bit signed integer.

Example

```
Input:
2
4 5
1 2 3 4
2 1 4
0 3 4 1
2 1 4
1 3 4 1
2 1 4
1 1
1
2 1 1
```

```
Output:
Case 1:
30
7
13
Case 2:
1
```

Solution Idea:

This problem is a nice example of segment tree with lazy propagation. All you need to solve this kind of problem is to do some experiment with the formula and convert them into a suitable form. From which we can drive the segment tree solution.

In this problem given an integer array $A=[a_1, a_2, a_3, a_4 \dots]$. You need to perform 3 kinds of operation on it.

type 0: Set all the element in the interval l to r to x .

type 1: Increase all the element in the interval l to r by x .

type 2: Print the value of square sum of the interval l to r . For example let $l=1$ and $r=3$. then you need to print $a_1^2 + a_2^2 + a_3^2$.

At first think we have a pre-calculated square sum array. Then if we perform type 1 operation which is increase it by x then what scenario will happen?

Our sum now become $(a_1+x)^2 + (a_2+x)^2 + (a_3+x)^2$. Let expand this equation-

$$\begin{aligned} & (a_1+x)^2 + (a_2+x)^2 + (a_3+x)^2 \\ &= a_1^2 + 2*a_1*x + x^2 + a_2^2 + 2*a_2*x + x^2 + a_3^2 + 2*a_3*x + x^2 \\ &= (a_1^2 + a_2^2 + a_3^2) + 3*x^2 + 2*x*(a_1 + a_2 + a_3) \end{aligned}$$

for a range l to r the generalized equation is –

$$(a_{l^2} + a_{l+1^2} + \dots + a_{r^2}) + (r-l+1)*x^2 + 2*x*(a_l + a_{l+1} + \dots + a_r)$$

We can store the information about sum, square_sum, type_0 lazy and type_1 lazy of an interval in each segment tree node. And perform type 1 operation then we can calculate then sum value for a node over a range by multiplication and square sum value by above equation.

For every type 0 operation we can simply put the value of sum and square sum by using some basic calculation and arithmetic multiplication operation.

And for every type 2 operation we just need to get the sum of square sum value over a range l to r .

So we can perform each of 3 operation using a segment tree with lazy propagation.

```

1 //https://www.spoj.com/problems/SEGSQRSS/
2 #include<bits/stdc++.h>
3 using namespace std;
4 #define ll long long
5 #define mx 100005
6 #define sqr(x) (x)*(x)
7 struct data{
8     ll sum,sqrsum,lazy,upd;
9 };
10 data tree[3*mx];
11 ll ara[mx];
12
13 void push_down(int n, int b, int e){
14     int l=n*2,r=l+1,mid=(b+e)/2;
15
16     if(tree[n].upd){
17         tree[l].lazy=0;
18         tree[r].lazy=0;
19         tree[l].sum=(mid-b+1)*tree[n].upd;
20         tree[l].sqrsum=(mid-b+1)*sqr(tree[n].upd);
21         tree[r].sum=(e-mid)*tree[n].upd;
22         tree[r].sqrsum=(e-mid)*sqr(tree[n].upd);
23         tree[l].upd=tree[n].upd;
24         tree[r].upd=tree[n].upd;
25         tree[n].upd=0;
26     }
27     if(tree[n].lazy){
28         tree[l].sqrsum+=(tree[l].sum*(2*tree[n].lazy))+(mid-b+1)*sqr(tree[n].lazy);
29         tree[r].sqrsum+=(tree[r].sum*(2*tree[n].lazy))+(e-mid)*sqr(tree[n].lazy);
30         tree[l].sum+=(mid-b+1)*tree[n].lazy;
31         tree[r].sum+=(e-mid)*tree[n].lazy;
32         tree[l].lazy+=tree[n].lazy;
33         tree[r].lazy+=tree[n].lazy;
34         tree[n].lazy=0;
35     }
36 }
37
38 void init(int n, int b, int e){
39     if(b==e){
40         tree[n].sum=ara[b];
41         tree[n].sqrsum=sqr(ara[b]);
42         tree[n].lazy=0;
43         tree[n].upd=0;
44         return;
45     }
46
47     int l=n*2,r=l+1,mid=(b+e)/2;
48
49     init(l,b,mid);
50     init(r,mid+1,e);
51
52     tree[n].sum=tree[l].sum+tree[r].sum;
53     tree[n].sqrsum=tree[l].sqrsum+tree[r].sqrsum;
54     tree[n].lazy=0;
55     tree[n].upd=0;
56 }
57
58 void update(int n, int b, int e, int x, int y, ll val, int type){
59     if(b>y || e<x) return;
60     if(b>=x && e<=y){
61         if(type==0){
62             tree[n].sum=(e-b+1)*val;
63             tree[n].sqrsum=(e-b+1)*sqr(val);
64             tree[n].lazy=0;
65             tree[n].upd=val;
66         }else{
67             tree[n].sqrsum+=(tree[n].sum*2*val)+((e-b+1)*sqr(val));
68             tree[n].sum+=(e-b+1)*val;
69             tree[n].lazy+=val;
70         }
71     }
72 }

```

```
73     push_down(n,b,e);
74
75     int l=n*2,r=l+1,mid=(b+e)/2;
76
77     update(l,b,mid,x,y,val,type);
78     update(r,mid+1,e,x,y,val,type);
79
80     tree[n].sum=tree[l].sum+tree[r].sum;
81     tree[n].sqrsum=tree[l].sqrsum+tree[r].sqrsum;
82
83 }
84
85 ll query(int n, int b, int e, int x, int y){
86     if(b>y || e<x) return 0;
87     if(b>=x && e<=y) return tree[n].sqrsum;
88
89     push_down(n,b,e);
90
91     int l=n*2,r=l+1,mid=(b+e)/2;
92
93     ll p=query(l,b,mid,x,y);
94     ll q=query(r,mid+1,e,x,y);
95     return p+q;
96 }
97 int main(){
98     int tt;scanf("%d",&tt);
99
100    for(int ks=1; ks<=tt; ks++){
101
102        int n,q; scanf("%d%d",&n,&q);
103        for(int i=1; i<=n; i++) scanf("%lld",&ara[i]);
104
105        init(1,1,n);
106
107        printf("Case %d:\n",ks);
108        while(q--){
109            int typ,x,y; scanf("%d%d%d",&typ,&x,&y);
110
111            if(typ==2){
112                ll ans=query(1,1,n,x,y);
113                printf("%lld\n",ans);
114            }else{
115                int v; scanf("%d",&v);
116
117                if(typ==0) update(1,1,n,x,y,v,0);
118                else update(1,1,n,x,y,v,1);
119            }
120        }
121    }
122
123    return 0;
124 }
```