

CHAR ARRAYS

char ch[5];

'a', 'b', 'c', ... 'z'
'A', 'B', 'C', ... 'Z'
'0', '1', '2', ... '9'

0	1	2	3	4	5

1 char → 1 byte

char name[100];

cin >> name;  parth

p	a	r	t	h	\0			
0	1	2	3	4	5	6	7	8...99

→ cout << int(name[s]);

→ 0 is ascii value of null

cin
→ space ' '
→ tab \t
→ newline \n
enter

solve → cin.getline(ch, length)
20

→ strlen(ch)

→ Reverse string

P	A	R	T	H
0	1	2	3	4

j = n-1 index

H	A	R	T	P
0	1	2	3	4

swap

H	T	R	A	P
0	1	2	3	4

swap

i > j
Break

→ palindrome → racecar

① input → racecar
reverse input → racecar

$$TC = O(n+n) = O(2n) = On$$

$$SC = O(n)$$

↳ bcz for reverse string
↓

→ uppercase lowercase convert Uppercase

$$\text{lowercase char} - 'a' + 'A' = \text{uppercase char}$$

$$\rightarrow \text{lowercase} \quad \text{lowercase char} = \text{uppercase char} + 'a' - 'A'$$

→ if input is like Parth → PARTH , here we can put conditions like

$$\text{IF } (\text{ch}[i] \geq 'a' \& \& \text{ch}[i] \leq 'z')$$



$$TC = O(n)$$
$$SC = O(1)$$

while ($i \leq j$)

↳ if ($\text{arr}[i] == \text{arr}[j]$)
 $i++$
 $j--$

Else

↳ not palindrome

$$\rightarrow 'c' - 'a' + 'A'$$

$$\rightarrow 99 - 97 + 65$$

$$\rightarrow 2 + 65$$

$$\rightarrow 67 \rightarrow \text{ascii of 'C'}$$

* String

```
#include <string>
```

```
string str; ~~~> cin >> str;
```

```
[P a r t h \o]
```

not accessible

```
getline (cin, str)
```

→ str.length()

.empty() → isEmpty ? T/F

.push_back('A')

.pop_back()

.substr(0, 6) (index, length)

Parth Patel ~~~> Parth
0 1 2 3 4 5 6 s e t s g
1 2 3 4 5 6

→ str1.compare(str2) == 0 → exactly same

a zzyy
a ≠ b

b zzyy
a → b

g7 < g8 → return (-1)

for b → a
g8 > g7
return (+1)

i
[l o v e]

→ in starting compare
length len1 ≠ len2
so not same

[l o v e r]
j

→ comparing i == j

→ sentence.find(target) ~~~> return index from where target starts

string sentence = ' hello Jee kaise ho ';

6 1 2 3 4 5 6 7 8 9 10

target

string target = 'Kaise' ;

```
string sentence = "hello Jee kaise ho saare";  
string target = "Everyone";
```

```
cout << sentence.find(target);  
if(sentence.find(target) == std::string::npos) {  
    cout << "Not Found" << endl;  
}
```

no position

→ replace (startIndex , howmanycharsremove , replacingString)

```
string str = "This is my First Message";  
string word = "Babbar"; "  
5 remove
```

```
str.replace(11, 5 , "Second"); → This is my Second Message  
cout << str << endl;
```

→ str.erase (startIndex , howmany char erase)

A	B	C	D	E	F	G	H	I	J
0	1	2	3	4	5	6	7	8	9
starting index		3							

~~~~~ A B C I J

5 char remove

## Remove All Adjacent Duplicates in string

untill s

→ ans = ""

ans = "a"

a b b a c a

ans = "ab"

here both are same so  
pop last char from ans

ans = "a"

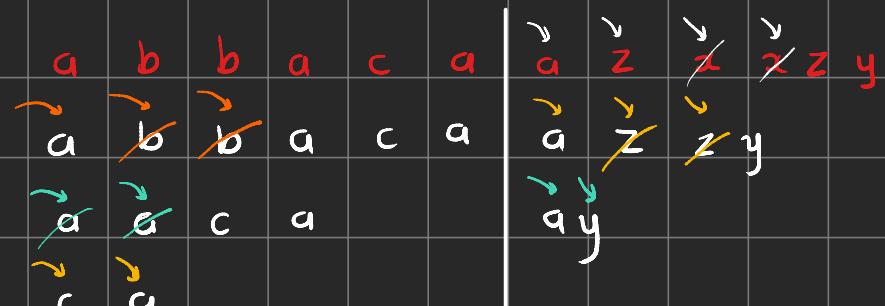
a b b a c a

ans = "

a b b a c a

ans = "c"

ans = "ca"



## Remove All occurrences of a substring

main →

d a a b c b a a b c b c

part →

a b c

d a a b c b a a b c b c ~> d a b a b c  
x x  
dab

## # Valid Palindrome II

a b c a

remove one char and check that palindrome or not

a b c a       $i=j \rightarrow$  no need to remove  
 $i$              $j$

a b c a       $i \neq j \rightarrow$  remove one  
 $i$              $j$

remove  $i$  to make palind

or  
remove  $j$  to make palind

~~a~~ b c d  $\rightarrow$  bcd  
~~i~~ itl         $j$   
check it is palindrome or not

a b c ~~d~~  $\rightarrow$  abc  
 $i$          $j-1$  ~~j~~  
check it is palindrome or not

## \* Minimum time difference

[ "23:59", "00:00" ]  $\rightsquigarrow$  [ 1439 + 0 ] mins

23 : 59  
01 234

convert string to int } stoi("23")

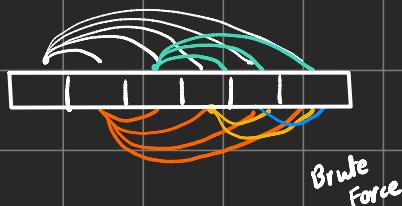
$$= 23 * 60 + 59 \\ = 1439$$

12:10 | 10:15 | 13:15 | 17:20 | 18:00 | 19:47 | 23:59

730 | 615 | 795 | 1040 | 1080 | 1187 | 1439

converting to mins

sort      Just to reduce TC



Brute Force

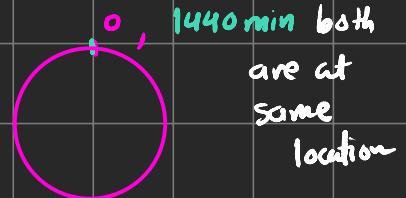
IMP test case

$$\text{lastdiff} = (\min[0] + 1440) - \min[n-1]$$

sorted

because in clock

circular  
difference



1440 min = 24 hr

|     |  |     |  |     |  |      |  |      |  |      |  |      |
|-----|--|-----|--|-----|--|------|--|------|--|------|--|------|
| 615 |  | 730 |  | 795 |  | 1040 |  | 1080 |  | 1187 |  | 1439 |
|-----|--|-----|--|-----|--|------|--|------|--|------|--|------|

diff<sub>1</sub> diff<sub>2</sub> d<sub>3</sub> d<sub>4</sub> d<sub>5</sub> d<sub>6</sub> → find min from diff all

end → minimum of (min, lastdiff)

## \* Palindromic Substrings (MIMP)

→ contiguous sub strings

|   |  |   |  |   |
|---|--|---|--|---|
| a |  | b |  | c |
|---|--|---|--|---|

→ substrings find

|   |   |   |
|---|---|---|
| a | b | c |
| i | j |   |

|     |  |    |  |   |
|-----|--|----|--|---|
| a   |  | b  |  | c |
| ab  |  | bc |  |   |
| abc |  |    |  |   |

find palindrom  
from all  
substrings → a, b, c

TC.  $O(n^3)$

## ② odd length substring

$\rightarrow \boxed{n | o | o | n}$   $i=j \Rightarrow n \checkmark$

$n | o | o | n$   $i$  is out of scope  
 $i$   $j$  stop

$\rightarrow \boxed{n | o | o | n}$   $\rightsquigarrow i=j \rightarrow o \checkmark$   
 $i$   $j$  so palindrome

$n | o | o | n$   $\rightsquigarrow i \neq j \rightarrow noo \times$   
 $i$   $j$  stop

$\rightarrow \boxed{n | o | o | n}$   $\rightsquigarrow i=j \rightarrow o \checkmark$

$n | o | o | n$   $\rightsquigarrow i \neq j \rightsquigarrow oon \times$   
 $i$   $j$  stop

## even length substring

$\rightarrow \boxed{n | o | o | n}$   $\rightsquigarrow i \neq j \rightsquigarrow no \times$   
 $i$   $j$  stop

$n | o | o | n$   $i=j \rightsquigarrow oo \checkmark$   
 $i$   $j$  so palindrome

$n | o | o | n$   $\rightsquigarrow i=j \rightsquigarrow noon \checkmark$   
 $i$   $j$

$n | o | o | n$   $\rightsquigarrow$  out of scope  
 $i$   $j$  stop

$\rightarrow \boxed{n | o | o | n}$   $\rightsquigarrow i \neq j \rightsquigarrow on \times$   
 $i$   $j$  stop

$\rightarrow \boxed{n | o | o | n}$   $\rightsquigarrow$  out of scope  
 $i$   $j$  stop

$\rightarrow$ 

|   |   |   |   |
|---|---|---|---|
| n | o | o | n |
| i | j |   |   |

 $\rightsquigarrow i=j \rightarrow n \checkmark$

$$TC = O(n^2)$$

n o o n  $\rightsquigarrow j$  is out of scope  
 i j  $\underline{\text{stop}}$

## \* How to pass custom comparator into sort function

string s = "babbar"

sort(s.begin(), s.end())  $\xrightarrow{\text{according to}} \text{lexicography, ascii values}$

babbar  
 $\xrightarrow{\text{ascending order}}$

what if I need into  
 decreasing order  $\xrightarrow{\text{rbbra}}$

```
bool cmp(char first, char second) {
    return first > second;
}
```

descending

first character is first and all  
 if condition is becoming true then

can be used with INT value  
 like bool (int first, int sec)  
 return first > sec

```
int main() {
    string s = "babbar";
    sort(s.begin(), s.end(), cmp);
```

$\xrightarrow{\text{rbbaa}}$

$\rightarrow$  return first < second  $\rightarrow$  increasing order

→

|                                                     |                                                                   |         |           |
|-----------------------------------------------------|-------------------------------------------------------------------|---------|-----------|
| "love"                                              | "babbar"                                                          | "rahul" | "sandeep" |
| $= l + o + v + e$<br>$= 12 + 15 + 22 + 5$<br>$= 54$ | $= b + a + b + b + a + r$<br>$= 2 + 1 + 2 + 2 + 1 + 18$<br>$= 26$ |         |           |

~~~~~ we can make custom comparator according to the total sum of all characters

→ using auto

```
vector<string> ans{"love", "babbar", "rahul", "sandeep"};

for( auto x: ans ){
    cout << x << " ";
}
cout << endl;
```

NOTE: while finding maximum number, int maxi = INT_MIN
— — minimum number, int mini = INT_MAX

❖ Valid Anagram

s = "anagram" t = "nagaram" → TRUE

s = "cat" t = "rat" → False

anagram of a string is another string that contains the same characters, only the order of characters can be different. For example, "act" and "tac" are anagrams of each other.

① sorting

- sorting according to ascii value

s → aaagmnrv } same
t → aaagmnrv }

TC. O(n log n)

| cook@pop-os: ~ \$ ascii -d | | | | | | | | | | | | |
|----------------------------|--------|-------|------|------|------|-------|---------|--|--|--|--|--|
| 0 NUL | 16 DLE | 32 | 48 0 | 64 ☐ | 80 P | 96 ` | 112 p | | | | | |
| 1 SOH | 17 DC1 | 33 ! | 49 1 | 65 A | 81 Q | 97 a | 113 q | | | | | |
| 2 STX | 18 DC2 | 34 " | 50 2 | 66 B | 82 R | 98 b | 114 r | | | | | |
| 3 ETX | 19 DC3 | 35 # | 51 3 | 67 C | 83 S | 99 c | 115 s | | | | | |
| 4 EOT | 20 DC4 | 36 \$ | 52 4 | 68 D | 84 T | 100 d | 116 t | | | | | |
| 5 ENQ | 21 NAK | 37 % | 53 5 | 69 E | 85 U | 101 e | 117 u | | | | | |
| 6 ACK | 22 SYN | 38 & | 54 6 | 70 F | 86 V | 102 f | 118 v | | | | | |
| 7 BEL | 23 ETB | 39 ' | 55 7 | 71 G | 87 W | 103 g | 119 w | | | | | |
| 8 BS | 24 CAN | 40 (| 56 8 | 72 H | 88 X | 104 h | 120 x | | | | | |
| 9 HT | 25 EM | 41) | 57 9 | 73 I | 89 Y | 105 i | 121 y | | | | | |
| 10 LF | 26 SUB | 42 * | 58 : | 74 J | 90 Z | 106 j | 122 z | | | | | |
| 11 VT | 27 ESC | 43 + | 59 ; | 75 K | 91 [| 107 k | 123 { | | | | | |
| 12 FF | 28 FS | 44 , | 60 < | 76 L | 92 \ | 108 l | 124 | | | | | |
| 13 CR | 29 GS | 45 - | 61 = | 77 M | 93] | 109 m | 125 } | | | | | |
| 14 SO | 30 RS | 46 . | 62 > | 78 N | 94 ^ | 110 n | 126 ~ | | | | | |
| 15 SI | 31 US | 47 / | 63 ? | 79 O | 95 _ | 111 o | 127 DEL | | | | | |

② counting

$s = \text{"anagram"}$

↓
frequency
table

$a \rightarrow 3$
 $g \rightarrow 1$
 $m \rightarrow 1$
 $n \rightarrow 1$
 $r \rightarrow 1$

$t = \text{"nagaram"}$

$a \rightarrow 3$
 $g \rightarrow 1$
 $m \rightarrow 1$
 $n \rightarrow 1$
 $r \rightarrow 1$

if both freq. table are
same then

→ we can go with hashmaps

→ Here



0 1 2 3 . . .

255 → total 256 ASCII values

→ for first string's chars

do arr[a] ≈ arr [int(a)] ++

→ second string's char

do (--)

} in end if not
0 then not ...

* Reverse only letters

$s = \text{"ab-cd"}$ ↽ "dc-ba"

① two pointer Approach with some rules

T e s t i n g - L e e t = c o d e - Q !

L

H

① if (s[l] → alphabet & s[h] → alphabet) \rightsquigarrow swap (s[l], s[h]), l++, h++
elseif (s[l] → not alphabet) \rightsquigarrow l++
else \rightsquigarrow h--

* isAlphabet \rightarrow a-z \rightarrow [97 - 122] and A-Z \rightarrow [65 - 90] ascii ch
 \rightarrow isalpha(ch) \rightarrow T/F C++ func STL

T e s t i n g - L e e t = c o d e - Q !

K L

Q e d o

H H K

S e - T !

* longest common prefix

["codehelp", "codenothelp"]

P P

| | | | | | |
|---|---|---|---|---|---|
| f | l | o | w | e | r |
| x | x | i | | | |

if $i = j = k \rightarrow$ go above

| | | | |
|---|---|---|---|
| f | l | o | w |
| j | j | j | |

| | | | | | |
|---|---|----------|---|---|---|
| f | l | i | g | h | t |
| k | k | k | | | |
| | | ↓ | | | |
| | | not same | | | |

* Reverse vowels of string

a,e,i,o,u

s = "IceCreAm" \rightarrow "Ae_CreIm"

| | | | | | | | |
|---|---|---|---|----|---|---|---|
| I | c | e | C | r | e | A | m |
| l | x | l | h | k- | k | | |
| A | e | | e | I | | | |

If $i = h \rightarrow$ vowel such
 $i \rightarrow$ not vowel $\rightarrow i++$

$h \rightarrow$ not vowel $\rightarrow h--$

#

* Isomorphic strings

$s = "egg"$ $t = "add"$ \rightsquigarrow True

$e \leftrightarrow a$
 $g \leftrightarrow d$

} unique mapping

$s = "foo"$ $t = "bar"$ \rightsquigarrow False

$f \leftrightarrow b$

$O \leftrightarrow a$

$o \leftrightarrow r$ ✗ unique mapping not possible

Here O is already mapped with a
and we are trying to map O to r again not possible

→ Sol. hash {256}



0 1 2 ...

255 → total 256

hash [s[i]] = t[i]

hash [p] = t

hash [112] = 116

$p \rightarrow t$

$p \rightarrow t$

} mapping with same char
means okay

$a \rightarrow i$

$p \rightarrow t$

$e \rightarrow l$

$r \rightarrow c$

$s = P a p e r$

i

$t = t i t l e$

i



0 1 2 ...

255 → total 256

isMapped (t[i]) = true

But if there is something
like $p \rightarrow x$ then False
bcz p is already mapped
with t that's why

egg
add

foo
bar

paper
title

budc
ki kp

e → a
g → d

agg
add

f → b
o → a

boo
baa

p → t
a → i

taper
tiper

b → k
a → i

kadc
kidc

True

False

do
bcz o → a

p → t
e → l
r → e

titer
titlr
title

d → k
False

↳ you cannot do it
bcz K is mapped with
b already

* Reorganize String

so that same chars
are not side by side
adjacent

s = "aab" → "aba"

s = "aaab" → "

aba x
aba x

s = aaabc → abaca

s = acabb → ababa

① Using priority queue ($n \cdot \log n$)

② greedy

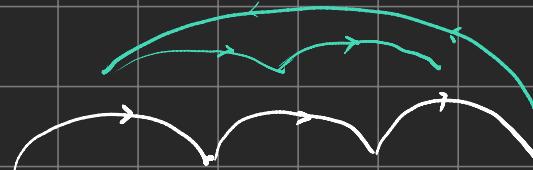
① Most occurrent character and fit it non-adjacently in one go (one iteration)

② fill the rest characters

→

a a a b b e f

$$\begin{array}{l} a \rightarrow 3 \\ b \rightarrow 2 \\ e \rightarrow 1 \\ f \rightarrow 1 \end{array}$$



a b a e a f b
0 1 2 3 4 5 6

→ a a a b b b e f g g

$$\begin{array}{l} a \rightarrow 3 \\ b \rightarrow 3 \\ e \rightarrow 1 \\ f \rightarrow 1 \\ g \rightarrow 2 \end{array}$$

a b a e a f b g b g
0 1 2 3 4 5 6 7 8 9

→ a a a b

$$\begin{array}{l} a \rightarrow 3 \\ b \rightarrow 1 \end{array}$$

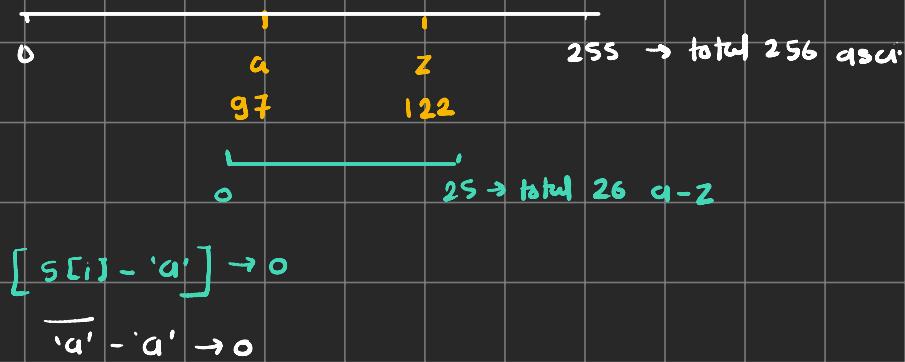
a a a b
0 1 2 3

When I am placing most occurrent character it should be placed in one go like

a a a but here
0 1 2 3

it is not possible so FALSE

- ① count hash
- ② Try to place most occurrent in one go \rightsquigarrow If (not possible) \rightarrow return false
- ③ Place other chars with one index gap



Group Anagrams

$strs = ["eat", "tea", "tan", "ate", "nat", "bat"] \rightsquigarrow [["bat"], ["nat", "tan"], ["ate", "eat", "tea"]]$

① eat, tea, tan, ate, nat, bat

$\xrightarrow{\text{sorted}}$ aet, aet, ant, aet, ant, abt



| Map \rightarrow | first
key | second
value |
|-------------------|---------------|-----------------|
| ① aet | eat, tea, ate | |
| ② ant | tan, nat | |
| ③ abt | bat | |

```

//#include <map>
map<string, vector<string>> mp;

for (auto str : strs) {
    string s = str;
    sort(s.begin(), s.end());
    mp[s].push_back(str);
}

vector<vector<string>> ans;
for (auto it = mp.begin(); it != mp.end(); it++) {
    cout << "key: " << it->first << " | value: ";
    for (const auto& str : it->second) {
        cout << str << " ";
    }
    cout << endl;
}

ans.push_back(it->second);
}

```

② without sorting

hash[256] = {}

string A string B
 ↓ ↓
 $\text{hashmap A} == \text{hashmap B}$ } anagram

- ① 'eat' creating hashmap above
- ② for 'tea' same hash key is coming so store value of tea

TC \Rightarrow str whole string $\rightarrow N$
 sorting $N \log N \rightarrow N \log K$
 map \rightarrow ignore now
 $O(N \cdot K \log K)$

SC $\Rightarrow O(N \cdot K)$ \rightarrow bcz storing ans
 length of strs so $N = 6$
 length of longest ele $K = 3$

- length = 6
- longest ele's length under cons = 3

| key | value |
|-----------|----------------|
| hash[256] | vector<string> |
| ① [] | eat, tea, ate |
| e - 1 | |
| a - 1 | |
| t - 1 | |
| ② [] | tan |
| t - 1 | |
| n - 1 | |
| a - 1 | |

```

std::array<int, 256> hashFunc(string s) {
    std::array<int, 256> hash = {0};
    for (int i = 0; i < s.size(); i++) {
        hash[s[i]]++;
    }
    return hash;
}

vector<vector<string>> groupAnagrams(vector<string>& strs) {
    map<std::array<int, 256>, vector<string>> mp;
    for (auto str : strs) {
        mp[hashFunc(str)].push_back(str);
    }

    vector<vector<string>> ans;
    for (auto it = mp.begin(); it != mp.end(); it++) {
        ans.push_back(it->second);
    }

    return ans;
}

```

new way to create Array

`std::array < datatype, size > currName ;`

`= {0}`

$$TC = O(N \cdot K)$$

$$SC = O(N \cdot K) + \text{Map size}$$

* Longest Palindromic substring

$s = "babad"$  "bab" or "aba"

① Sub strings of string

② extract Palindromic ones

↳ find longest from them

```

bool isPalindrome(string& sub, int start, int end) {
    while (start < end) {
        if (sub[start] != sub[end]) {
            return false;
        }
        start++;
        end--;
    }
    return true;
}

string longestPalindrome(string s) {
    string ans = "";
    for (int i = 0; i < s.size(); i++) {
        for (int j = i; j < s.size(); j++) {
            if (isPalindrome(s, i, j)) {
                string t = s.substr(i, j - i + 1);
                ans = t.size() > ans.size() ? t : ans;
            }
        }
    }
    return ans;
}

```

① substrings



• b a b a b → a
i j
ab

• b a b a b → b
i j
b

* Find index of first occurrence

s = "sadbutsad"

t = "sad"

s a d b u t s a d
0 1 2
↓
0 th index

→ L e t s a d u
0 1 2 3 4 5 6

s a d
s a d
s a →
s a d
v

{ sliding window }

→ if initial character of t matched then

match other chars of t

* String to Int (atoi) / run length encoding

-- lakshay 4139 --

```
#include <iostream>
#include <cstdlib> // Required for atoi
#include <string>

int main() {
    std::string str1 = "12345"; // C++ std::string
    const char* str2 = "-6789"; // C-style string (const char*)

    // Using stoi with std::string (str1)
    int num1_stoi = std::stoi(str1); // Converts "12345" to 12345
    std::cout << "Using stoi, str1: " << num1_stoi << std::endl; // Output: 12345

    // Using atoi with C-style string (str2)
    int num2_atoi = atoi(str2); // Converts "-6789" to -6789
    std::cout << "Using atoi, str2: " << num2_atoi << std::endl; // Output: -6789

    // Using atoi with str1 converted to C-style string (str1.c_str())
    int num3_atoi = atoi(str1.c_str()); // Converts "12345" to 12345
    std::cout << "Using atoi, str1.c_str(): " << num3_atoi << std::endl; // Output: 12345
```

- ① ignore leading white space
- ② determine sign
- ③ if digit found after finding sign - integer prepare
- ④ till next nondigit char found
- ⑤ special handling of if num is out of range so return INT-MIN OR INT-MAX

| | | | | | | | | |
|---|------|---|---|---|---|---|---|---|
| → | a | a | b | b | c | d | e | e |
| | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 |
| | ↓ | ↓ | | | | | | |
| | prev | i | | | | | | |

s → ^{index}
aa bb cdee
 a 2

cont = 1

prev == i → cont++

prev ≠ i → s[index] = prev

cont > 1 then s[index + 1] = cont

cont = 1

a a b b c d e e
↓ i
prev

s → ~~a a~~ bb c d e e
index
a 2

INT to char

digit + '0'

a a b b c d e e
prev i

cont = 2

a a b b c d e e
p i

s → a 2 ~~b b~~ c d e e
index
b 2

a a b b c d e e
p i

s → a 2 b 2 ~~c~~ d e e
index
c
cont = 1 so do not store it
on (index + 1)

a a b b c d e e
p i

s → a 2 b 2 c ~~d~~ e e
index
d

a a b b c d e e
p i

s → a 2 b 2 c d ~~e~~ e e
index
e 2

* Int to Roman

27 → X X V I I

$$10 + 10 + 5 + 1 + 1 = 27$$

58 → num >= RomanValue

$$\begin{aligned} 58 &>= 50 \longrightarrow 50 + 8 \\ &\quad L + 8 \end{aligned}$$

$$\begin{aligned} \text{num} &= 58 - 50 \\ &= 8 \end{aligned}$$

8 >= RomanValue

$$\begin{aligned} 8 &>= 5 \longrightarrow 5 + 3 \\ &\quad V + 3 \end{aligned}$$

$$\begin{aligned} \text{num} &= 8 - 5 \\ &= 3 \end{aligned}$$

3 >= RomanValue

$$3 >= 1 \longrightarrow \text{while } (\text{num} >= 1) \rightsquigarrow \text{odd symbol I}$$

$$\text{num} = \text{num} - \text{value}$$

| Symbol | Value |
|--------|-------|
| I | 1 |
| V | 5 |
| X | 10 |
| L | 50 |
| C | 100 |
| D | 500 |
| M | 1000 |

largest
to smaller



$$3333 = 1000 + 1000 + 1000 + 333 \quad \rightsquigarrow \text{num} = 3333 - 3000$$

M M M + 333

$$= 333$$

$$333 = 100 + 100 + 100 + 33 \quad \rightsquigarrow \text{num} = 333 - 300$$

C C C + 33

$$= 33$$

$$33 = 10 + 10 + 10 + 3 \quad \rightsquigarrow \text{num} = 33 - 30$$

X X X + 3

$$= 3$$

$$3 = 1 + 1 + 1 \quad \rightsquigarrow \text{num} = 3 - 3$$

I I I

$$= 0$$

$$1994 = 1000 + 994 \quad \rightsquigarrow \text{num} = 1994 - 1000$$

M

$$= 994$$

$$994 = 100 + \dots \text{9 times} + 94$$

BUT ans is MCMXCIV

$$\overline{900} \rightarrow CM = |100 - 1000| = 900 \quad \left. \right\} \text{so we need}$$

$$\text{like } IV = |1 - 5| = 4 \quad \left. \right\} \text{new List with specials}$$

① Map Iterate (larger to small)

```
while (num >= value[i])  
{  
    roman = roman + symbol  
    num = num - value[i]  
}
```

```
string romanSymbols[] = {"M", "CM", "D", "CD", "C", "XC", "L", "XL", "X", "IX", "V", "IV", "I"};  
int values[] = {1000, 900, 500, 400, 100, 90, 50, 40, 10, 9, 5, 4, 1};
```

```
string ans = "";  
  
for(int i=0; i<13; i++){  
    while(num>=values[i]){  
        ans = ans + romanSymbols[i];  
        num = num - values[i];  
    }  
}
```

M = 1000
CM = 900
D = 500
CD = 400
C = 100
XC = 90
L = 50
XL = 40
X = 10
IX = 9
V = 5
IV = 4
I = 1

* Zig Zag conversion

s = "PAYPALISHIRING" rows =



take 3 vector string
and store values like
this

0 → PAHN
1 → APLSIIIG
2 → YIR

* Largest Number

1 | 2 | 3 | 4 | 5

sort
decrease

5 4 3 2 1
largest num

3 30 34 5 9

→

9 5 3 4 3 0 3

9 5 3 4 3 3 0 → Ans

lexicography sort

{
bool mycom (string a, string b)
return a > b

this say that
 $30 > 3$

but we need
Something that
say $3 > 30$

solve

$a = 30, b = 3$

string t1 = 303 = a+b
t2 = 330 = b+a

means t1 સી હોલે ← TRUE
અની

return t1 > t2

$303 > 330$ F

t2 સી હોલે કહો

