

Application
without service
time

→ FCFS & RR can work;
SPN cannot work.

→ SPN is non pre
emptive

Throughput

→ Number of process
completed per time
unit → High for shorter
processes

Shortest
Remaining
Time

→ Preemptive version
of SPN

→ Running process can
be kept in ready queue

Example

A	0	3
B	2	6
C	4	4
D	6	9
E	8	2

→ $T=0$ A

$T=2$ B A (A:1, B:6)

$T=3$ B

$T=4$ B C (B:5, C:4)

FCFC
Tie Breaker
used

$T=6$ D B C (C:2, B:5, D:5)

$T=8$ E D B (D:5, B:5, E:2)

$T=10$ D B

$T=15$ D

$T=20$

Highest
Response Ratio
Next

→ We are looking to minimize Normalized TAT.

$$\rightarrow R = \frac{\overset{\text{Waiting Time}}{\downarrow} (W+S) \overset{\text{Service Time}}{\nwarrow}}{S}$$

→ Non-preemptive policy

→ Shorter process favored (small denom)

→ Long / IO / old
favoured (large denom)

→ Balances favour of
short processes along
with starvation prevention

Is HRRN
the best?

→ For measurable
performance ; SRT is
better

→ For non-measurable
(fair treatment) ;
HRRN is better

Real World
Strategies

→ Completely fair
scheduling is used
in Linux systems

Feed back
Scheduling

→ If we want to
avoid short process
penalty; penalise
longer processes.

→ If you can't estimate
time remaining; use

What is the main idea?

time execution.

→ Penalise longer running jobs

→ Internally using Round Robin

→ Priority based queues after some fixed quantum of execution time.

Steps

① RQO

② Preempted to RQ.

② Every subsequent preemption has lower priority

→ All queue follows FCFS ; Last queue follows Round Robin (as no lower queue to demote)

→ Imagine a moving treadmill that is descending after every turn

Analysis for
FB

↓
Or multi
level feedback

→ Turnaround is stretched alarmingly for longer processes

→ $q = 2^i$ reduces harshness on longer processes

Other possible
strategies

→ Priority scheduling (Preemptive / NonPreemptive)

→ Longest Job First (Non Preemptive)

→ Longest Remaining Time First (Preemptive)

→ None is used in general purpose OS in pure form; hybrid version is used.

HW-1

3 processes
 $Cid = 0, 1, 2$
 $(S = 2, 4, 8)$
Arrival = 0 for
all. Consider
LRTF; ties are
broken by high
priority to lowest
process id. What
is the average
turn around time?

hw-2

3 proc at $t=0$
arrival; total
 $S = 10, 20, 30$.
20-1. is I/O,
70-1. is computation
& 10-1. is I/O.
Shortest rem
compute time
first. Non
preemptive. For
what % of time
does CPU remain

idle?

Summary: SPN - SRT - Throughput
- HRRN - Perfect balance
achieved - FB - Modified FB with
 $q = 2^i$ - Other possible scheduling
algorithms - None used in real world