

Homework 1.

- The file name of your homework (in PDF) should be in the format: “學號-作業編號.pdf”. For example: 00957999-hw1.pdf
- Please submit your homework to Tronclass **before 23:59, October 19, 2022.**
(可以用 word 檔寫完後轉成 pdf 檔上傳，或是手寫後拍照後存成 pdf 檔上傳)

1. (10%) Order the following function by growth rate in increasing order (由小到大排序):

$5n^2$, 2^n , $n \log(\log n)$, $n \log(n^2)$, n^n , $n!$

	1	2	3	4	5	6	7	8	9
$5n^2$	5	20	45	80	125	180	245	320	405
2^n	2	4	8	16	32	64	128	256	512
$n \log(\log n)$	$\log(\log 1)$	$2\log(\log 2)$	$9\log(\log 9)$
$n \log(n^2)$	$\log 1$	$2\log 4$	$3\log 9$	$4\log 16$	$5\log 25$	$9\log 81$
n^n	1	4	27	64	3125	46656	823543	16777216	387420489
$n!$	1	2	6	24	120	720	5040	40320	362880

Ans: $n \log(\log n) < n \log(n^2) < 5n^2 < 2^n < n! < n^n$

2. (20%) Show that the following statements are correct:

(a) $n! = O(3n^n)$

(b) $5n^2 - 20n = \Theta(n^2)$

(a): 存在 $C, n_0 > 0$, 當 $n \geq n_0$ 時, $n! \leq C3n^n$

$$n! \leq 3n! \leq 3n^n \leq C3n^n \quad \text{當 } C = 1, n_0 = 1 \text{ 不等式成立}$$

(b): 證明 $5n^2 - 20n = O(n^2)$ & 證明 $5n^2 - 20n = \Omega(n^2)$

$$\text{存在 } C, n_0 > 0, \text{ 當 } n \geq n_0 \text{ 時, } 5n^2 - 20n \leq Cn$$

$$5n^2 - 20n \leq 5n^2 \leq Cn^2 \quad \text{當 } C = 5, n = 1 \text{ 時不等式成立}$$

$$\text{存在 } C, n_0 > 0, \text{ 當 } n \geq n_0 \text{ 時, } 5n^2 - 20n \geq Cn^2$$

$$4n^2 + n^2 - 20n \geq 4n^2 \text{ (當 } n \geq 20) \geq Cn^2 \quad \text{當 } C = 2, n_0 = 20 \text{ 得證}$$

3. (20%) Show that the following statements are incorrect:

(a) $10n^2 + 100n + 9 = O(n)$

(b) $3^n = O(2^n)$

(a): 不存在 $C, n_0 > 0$, 當 $n \geq n_0$ 時, $10n^2 + 100n + 9 \leq Cn$

反證:

$$\text{假設 } C, n_0 > 0 \text{ 存在 當 } n \geq n_0, 10n^2 + 100n + 9 \leq Cn$$

$$Cn \leq n^2 \leq 10n^2 \quad 10n^2 + 100n + 9 \leq Cn$$

(當 $n \geq C$ 的時, 左式成立) 則 $n > \max(n_0, C)$ 不等式成立, 矛盾

(b): 不存在 $C, n_0 > 0$, 當 $n \geq n_0$ 時, $3^n \leq C2^n$

反證:

$$\text{假設 } C, n_0 > 0 \text{ 存在 當 } n \geq n_0, 3^n \leq C2^n$$

$$\text{取 } \log \rightarrow n \log 3 \leq (\log C + n \log 2)$$

$$\log C + n \log 2 < n \log 3 \leq (\log C + n \log 2)$$

當 $C < n(\log 3 - \log 2)$ 時, 左式成立 矛盾

4. (20%) Fill the “accurate” step count table (yellow regions) for the following program:

Let $MAX_SIZE = M$

Statement	s/e	Frequency	Total Steps
void mult(int a[][MAX_SIZE], int b[][MAX_SIZE], int c[][MAX_SIZE])	0	0	0
{	0	0	0
int i, j, k;	0	0	0
for (i = 0; i < MAX_SIZE; i++)	1	$M + 1$	$M + 1$
for (j = 0; j < MAX_SIZE; j++) {	1	$M(M + 1)$	$M^2 + M$
c [i][j] = 0;	1	M^2	M^2
for (k = 0; k < MAX_SIZE; k++)	1	$M^2(M + 1)$	$M^3 + M^2$
c [i][j] += a [i][k] * b [k][j];	1	M^3	M^3
}	0	0	0
}	0	0	0
Total			$2M^3 + 3M^2 + 2M + 1$

5. (10%) Given an four-dimensional array $A[100][200][300][400]$, if the address of $A[4][18][29][36]$ is $\alpha + X$, what is the value of X ?

Note: (1) α is the address of $A[0][0]$

(2) Only consider row major here.

$$4 * 200 * 300 * 400 + 18 * 300 * 400 + 29 * 400 + 36 = 98171636$$

Ans: 98171636

6. (10%) 請說明如何使用課本 Program 2.6 的 `padd` 函式，計算多項式乘法 $A(x) * B(x)$ ，其中

$$A(x) = 3x^2 + 2x + 1$$

$$B(x) = 5x^3 + 4x - 1$$

```
int solve(int n, int k){
    int idx = 0;
    for(int i = 0; i < n; i++){
        for(int j = 0; j < k; j++){
            C[idx].coef = A[i].coef * B[j].coef;
            C[idx++].expon = A[i].expon + B[j].expon;
        }
    }
    int now = 0;
    for(int i = 0; i < idx; i++){
        if(C[i].vis) continue;
        ans[now].expon = C[i].expon;
        ans[now].coef = C[i].coef;
        for(int j = i + 1; j < idx; j++){
            if(C[i].expon == C[j].expon){
                ans[now].coef += C[j].coef;
                C[j].vis = 1;
            }
        }
        now++;
    }
    return now;
}
```

將 A 和 B 的所有項相乘

乘完以後同項合併

其中 `now - 1` 代表最高項次

7. (10%) Ackerman's function $A(m, n)$ is defined as:

$$A(m, n) = \begin{cases} n + 1, & \text{if } m = 0 \\ A(m - 1, 1), & \text{if } n = 0 \\ A(m - 1, A(m, n - 1)), & \text{otherwise} \end{cases}$$

What is the value of $A(3, 4)$?

$$A(m, n) = n + 1 \text{ when } m = 0$$

$$n + 2 \text{ when } m = 1$$

$$A(3, 4) = A(2, A(3, 3))$$

$$A(3, 3) = A(2, A(3, 2))$$

$$A(3, 2) = A(2, A(3, 1))$$

$$A(3, 1) = A(2, A(3, 0))$$

$$A(3, 0) = A(2, 1)$$

$$A(2, 1) = A(1, A(2, 0)) = A(1, 3) = 5$$

$$A(2, 0) = A(1, 1) = 3$$

$$A(2, 5) = A(1, A(2, 4)) = 13$$

$$A(2, 4) = A(1, A(2, 3)) = 11$$

$$A(2, 3) = A(1, A(2, 2)) = 9$$

$$A(2, 2) = A(1, A(2, 1)) = A(1, 5) = 7$$

由這幾項得出

$$\rightarrow A(2, n) = 2(n + 3) - 3 \text{ 此規律後再反推}$$

$$A(3, 4) = A(2, A(3, 3))$$

$$A(3, 3) = A(2, A(3, 2))$$

$$A(3, 2) = A(2, A(3, 1))$$

$$A(3, 0) = A(2, 1) = 5$$

$$A(3, 1) = A(2, A(3, 0)) = A(2, 5) = 13$$

$$A(3, 2) = A(2, A(3, 1)) = A(2, 13) = 29$$

$$A(3, 3) = A(2, A(3, 2)) = A(2, 29) = 61$$

$$A(3, 4) = A(2, A(3, 3)) = A(2, 61) = 125$$

Ans: 125