

# Shortest Path

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## Dijkstras:

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Implementation with PQ

- $O(|E| + |V| \log |V|)$
- Make source current node with its distance 0
- Repeat until no elements in PQ
  - If current node is not in distance map or has a greater value than computed value place the current node in the distance map (mapped to distance)
  - Place its neighbors in the PQ with their distances to current node + current node distance
  - Dequeue min element from PQ and make current node

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## Bellman Ford:

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psuedo java code

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```
for i=1 to size(vertices)-1
  for each edge (u,v)
    if distance[u]+w < distance[v]
      distance[v]=distance[u]+w
      predecessor[v]=u
```

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c++ code

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```
function BellmanFord(list vertices, list edges, vertex source)
  ::distance[],predecessor[]

  // This implementation takes in a graph, represented as
  // lists of vertices and edges, and fills two arrays
  // (distance and predecessor) with shortest-path
  // (less cost/distance/metric) information

  // Step 1: initialize graph
  for each vertex v in vertices:
    if v is source then distance[v] := 0
    else distance[v] := inf
    predecessor[v] := null

  // Step 2: relax edges repeatedly
  for i from 1 to size(vertices)-1:
    for each edge (u, v) in Graph with weight w in edges:
      if distance[u] + w < distance[v]:
        distance[v] := distance[u] + w
        predecessor[v] := u
```

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```
// Step 3: check for negative-weight cycles
for each edge (u, v) in Graph with weight w in edges:
    if distance[u] + w < distance[v]:
        error "Graph contains a negative-weight cycle"
return distance[], predecessor[]
```

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### Notes

- $O(|V| |E|)$
- Allows negative weights