anorder - LNR (BST)  $T \cdot C = O(N)$ S.C. = Stack space = min RST 6, 6, 30,50 morder predecessor morder LST , 10 , RST

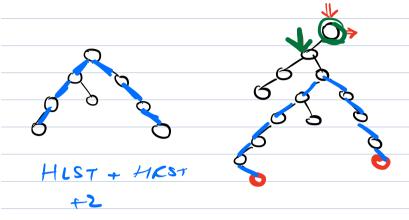
> trodes, we try to find the inorder predecessor. & update it's right to the node. 2, 3, 5, 6, 10, 20, so while (cur /= null) & if (curr. left == mull) & print carr. value, curr = curr. right; else d

```
pred = find Predecessor (cur);
           (pred. right = = null) «
              pred. right = null;
               print (curr. value);
               carr = carr. night;
         z
2, 3, 5, 6, 10, 30, 50
       T.C = O(3N) = O(N)
        S_{i}(t) = O(t)
```

birdPredecessor (root) &

curr = root;

curr = curr. left; while (Curr. night /= null) &s (arr.right 1 = 200t))7 carr = carr right; octorn curi Find the diameter of the longest path b/w any 2 nodes.



Code

int dia (root) x

y (root = = null) d ret -1;

int lh = height (root. left); int vh = height (root. right);

ans = lh+~h+2

int ld = dia (root. left); int -d = dia (root. right);

retorn mex (ld, rd, ans);

ኔ

## $T.C. = O(N^2)$

Tree onfo class int height; int dia; constractor () .... b Tree Inp dia (root) K if (root = = NULL) d return new Treedif (-1, -1); Tree onfo l = dia ( root. left); True dup v = dia (root. right); return new Treeding (max (l. height),
max (l. diameter, B. diameter) ( l. height + v. height) + 2)) 1

Height (Hught Balanced Tree) = 
$$\log N$$

| let (LST) - let (RST) |  $l = 1$ 

| wodes of the tree.

| height is Balanced.

| Tree chifo is Balanced (rost) &

| H.W. (Code)