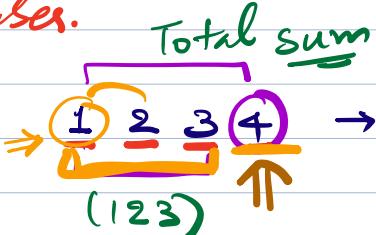


Agenda

- 1) More problems on Recursion
- 2) TC / SC analysis.

Q Given a number. Write a recursive code to calculate the sum of digits of given numbers.

Solⁿ 

$$1 + 2 + 3 + 4 = 10$$

$$\underline{N \% 10} = \text{last digit}$$
$$\underline{\text{Sum}(N / 10)} = \text{no. formed by remaining digits}$$


$$55 + 2 = SS$$

Code

```
int calculateSum (N) {
```

// Ass: calculateSum(N) returns the correct digit sum of the no. N.

// Base Case

```
if (N < 10) { return N; }
```

return $(N \cdot 10) + \underline{\text{calculateSum}}(\underline{N/10})$;

5

calculate $(\boxed{1234})$ {

return $\underline{4} + \underline{\text{calculate}}(\underline{123})$

5

N digit $\Rightarrow (N-1)$ digits $\Rightarrow (N-2)$

⋮
⋮
⋮

1 digit no. $\begin{array}{c} \text{1} \\ \text{2} \\ \text{3} \end{array} \Rightarrow \underline{\underline{N}}$

Q-2 Implement pow function

$\text{pow}(a, n) \Rightarrow a^n$

$\text{pow}(2, 3) = 8$

$a^n = a \times a \times a \times a \dots \dots a$ (N times)

$$a \times a^{n-1}$$

$$2^5 = 2 \times 2^4$$

$$\underline{n=0}$$

int pow (a, n) {

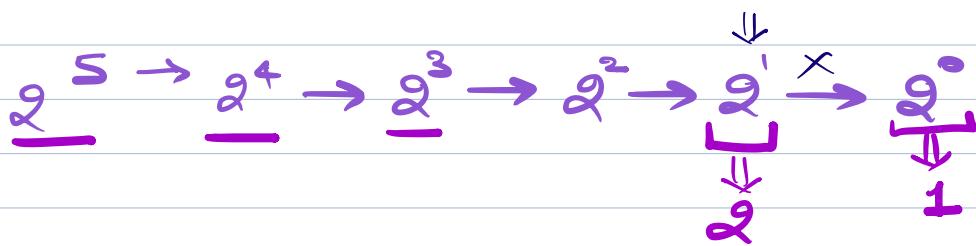
// Base Case

if (n==0) { return 1; }

if (n==1) { return a; }

return a \times pow (a, n-1);

}



$$a^{\cancel{5}} \Rightarrow (n-1)$$

$$a \times a^{\cancel{4}}$$

$$a \times a^{\cancel{3}}$$

$$a \times a^{\cancel{2}}$$

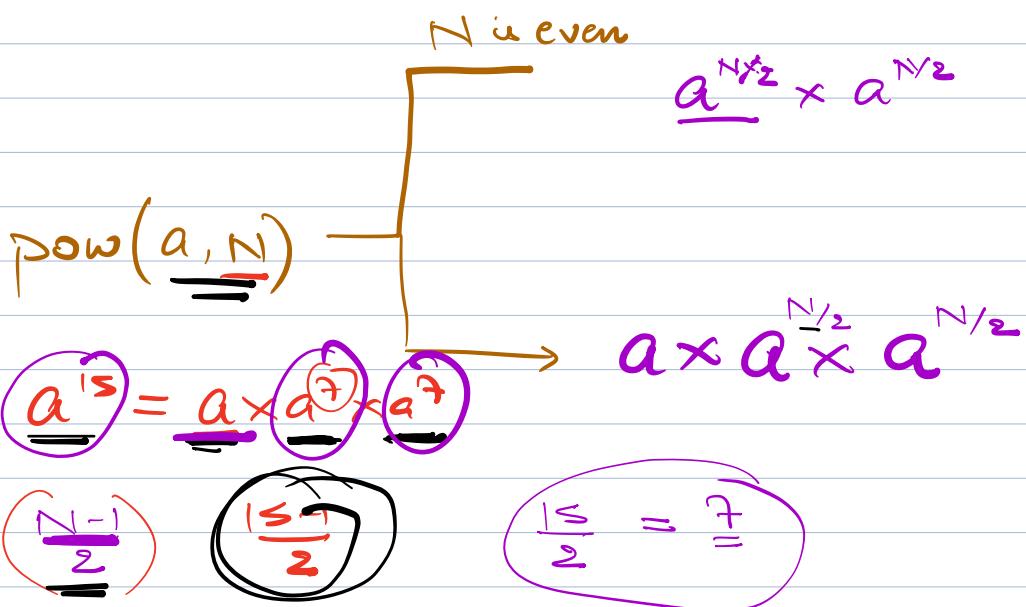
$$a \times (\underline{a})^{\cancel{1}}$$

$$a^{\cancel{n}} \in \underline{O(N)}$$

$$\approx T.C = \underline{O(N)}$$

$$\text{Eq} \Rightarrow (\underline{a}^{10}) = \underline{a}^{\frac{s}{2}} \times (\underline{a}^9) \\ (\underline{a}^s)^2 \times \underline{a}^s \quad (\perp \text{RC})$$

$$\begin{aligned} \underline{a}^{15} &= a \times a^{\frac{14}{2}} \\ &= a \times a^7 \times a^7 \quad (\perp \text{RC}) \end{aligned}$$



Code

```
int pow(a, n) {
    if (n == 0) return 1;
    if (n % 2) {
        x = pow(a, (n-1)/2);
    }
```

fast power method

```
int pow(a, N) {
    if (n == 0) return 1;
    x = pow(a, n/2);
    xsquare = x*x;
    if (n % 2) {
        return a * xsquare;
    }
}
```

return axxn;

5
else {

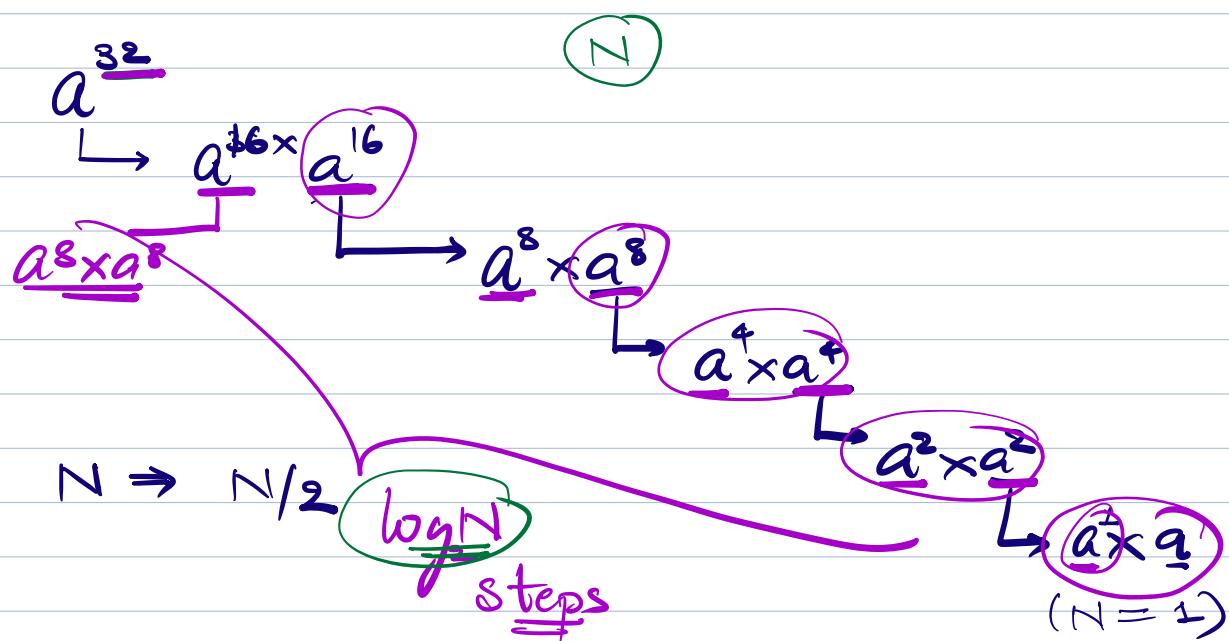
x = pow (a, n/2);

return xx;

5

- -
else if
 return square;

5



$$T.C. = O(\log_2 N)$$

$a^n \uparrow$

232-1 (Unsigned int)

240

(100)250

$\xrightarrow{\text{mod}}$

Google



pow(a, n, m) \Rightarrow $\underline{a}^n \cdot \underline{1 \cdot m}$

Eg \rightarrow

$a = 5$
 $n = 3$
 $m = 20$

5^3

$$(125) \cdot 1 \cdot 20 = \underline{\underline{125}} \cdot \underline{\underline{20}}$$

Code

$T(N)$

$10^9 + 7$

```

int pow (a, n, m) {
    if (n == 0) { return 1; } // O(1)
    x = pow(a, n/2, m); // O(1)
    // O(1) xsquare = ((x * m) * (x * m)) * 1 · m;
    O(1) if (n & 1) {
        return ((a * m) * (xsquare)) * 1 · m;
    }
    else
        O(1) return xsquare;
}

```

b

$$a \cdot m = [0, \underbrace{10^9 + 7 - 1}]$$

$$\underline{10^s} \times 10^s = \textcircled{10^s} \cdot \textcircled{10^{9+7}}$$

pow(s, 3, 20) ↴

$$x = \cancel{\text{pow}(s, 1, 20)}; \quad \textcircled{s} \in [0, m-1]$$

$$\Rightarrow \underline{x \text{ square}} = ((\underline{x \cdot 20}) \times (\underline{x \times 20})) \cdot \underline{20}$$
$$(\underline{s \times s}) \cdot \underline{20}$$
$$(2s) \cdot \underline{20}$$

$$\therefore \text{return } (\underline{s \times s}) \cdot \underline{20} \quad (\underline{s})$$

pow (s, 1, 20) ↴

$$x = \text{pow}(s, \cancel{0}, \cancel{20}),$$

$$x \text{ square} = ((1 \cdot 20) \times (1 \cdot 20)) \cdot 1 \cdot 20$$

$$\text{return } (\underline{s \times 1}) \cdot \underline{20} \quad (\underline{s})$$

b

Calculate x (very sig)

$$[0 - \underline{M-1}]$$

↓

$$\cancel{\underline{2M+4}} \cdot \underline{1 \cdot M} = \cancel{4}$$

$$\cancel{3M+4} \cdot \underline{1 \cdot M} = \cancel{4}$$

$$\cancel{2M+8} \cdot \underline{1 \cdot M} = \cancel{8}$$

TC of Recursion

1) $\underline{\text{Sum}(N)} = N + \underline{\text{sum}(N-1)}$ (Main Logic)

Assume: $\underline{T(N)}$ $\underline{T(N-1)}$

$$\underline{T(N)} = \underline{T(N-1)} + 1 \quad \text{--- (1)}$$

$T(N) \Rightarrow \text{TC of } \underline{\text{sum}(N)}$

$T(N-1) \Rightarrow \text{TC of } \underline{\text{sum}(N-1)}$

In eqn (1) dat $N = N-1$

$T(N) =$

$$\underline{T(N-1)} = \underline{(T(N-2) + 1)}$$

put this in eq. ①

$$\Rightarrow T(N) = \underline{(T(N-2) + 1)} + 1$$

$$T(N) = \underline{T(N-2)} + \underline{2} \quad - \textcircled{2}$$

$$\underline{T(N-2)} = \underline{T(N-3)} + 1$$

$$T(N) = \underline{(T(N-3) + 1)} + 2$$

$$\underline{T(N)} = \underline{\underline{T(N-3)}} + \underline{\underline{3}} \quad - \textcircled{3}$$

⋮
: K times (

$$T(N) = T(\underline{\underline{N-K}}) + \underline{\underline{K}} \quad -$$

Base Case

$$N - K = 0 \quad (K = N)$$

After N steps

$$T(N) = T(0) + N$$



$$T(N) = 1 + T(N)$$

$$T.C. = O(N)$$

TC

$$\Sigma_x T(N) = 2T(N-1) + 1$$

\Rightarrow

$\begin{aligned} N &= N-2 \\ &\quad \downarrow \\ N &= N-1 \end{aligned}$

$T(N-1) = 2T(N-2) + 1$

$\boxed{T(0) = 1}$

$$T(N) = 2(2T(N-2) + 1) + 1$$

$$T(N) = 4[T(N-2)] + 2 + 1$$

$$T(N) = 4[2T(N-3) + 1] + 2 + 1$$

$$T(N) = \frac{8}{2^3} [T(N-3)] + \frac{4+2+1}{2^3 - 1}$$

$$T(N) = 8[2T(N-4) + 1] + 4+2+1$$

$$T(N) = \frac{16}{2^4} T(N-4) + \frac{8+4+2+1}{(2^4 - 1)}$$

1 - - - -

$$T(N) = 2^k (T(N-k)) + 2^k - 1$$

⋮
⋮

κ steps

$$T(N) = 2^k (T(N-k)) + 2^k - 1$$

$\text{if } \kappa = N$

$$T(N) = 2^N (T(0)) + 2^N - 1$$

$$T(N) = 2^N + 2^N - 1$$

$$\underline{T(N) = 2^{N+1} - 1}$$

$$T.C. = O(2^N)$$

H.W.

$$T(N) = T(N-1) + T(N-2) + 1$$

$$\begin{aligned} T(0) &= 1 \\ T(1) &= 1 \end{aligned}$$

$T(N) = T(\underline{N/2}) + 1$ ($= 1$)
 power form
 $N = N/2$
 $T(N/2) = T(\underline{N/4}) + 1$

$$T(N) = T(\underline{\underline{N/4}}) + 2$$

$$T(N) = \left(T(\underline{\underline{N/8}}) + 1 \right) + 2$$

$$= T(\underline{\underline{N/8}}) + 3$$

$$T(N) = T(\underline{\underline{\underline{N/2^3}}}) + 3$$

⋮

~~IT~~ \downarrow

$$T(N) = T\left(\underline{\underline{\underline{\frac{N}{2^k}}}}\right) + k$$

$\frac{N}{2^k} = 1$
 $N = 2^k$

$$\log_2 N = \log_2 2^k$$

$$(K = \log_2 N)$$

$$T(N) = 1 + \log_2 N$$

$$T.C. = O(\log N)$$

~~POW~~ \rightarrow return $a \times \underline{\text{pow}(a, n/2)} \times \underline{\text{pow}(a, n/2)}$

$$\underline{T(N)} = 2 \times \underline{T(N/2)} + 1$$

$$N = N/2$$

$$T(N/2) = 2 \times T(N/4) + 1$$

$$\begin{aligned} T(N) &= 2 \left(2 T\left(\frac{N}{4}\right) + 1 \right) + 1 \\ &= 4 T\left(\frac{N}{4}\right) + 2 + 1 \end{aligned}$$

$$= 8 T\left(\frac{N}{8}\right) + 4 + 2 + 1$$

$$= 16 T(N) + 8 + 2 + 1$$

$$(1/16)$$

$$= 2^4 T\left(\frac{N}{2^4}\right) + 2^4 - 1$$

⋮
⋮

K steps

$$\Rightarrow T(N) = 2^K T\left(\frac{N}{2^K}\right) + 2^K - 1$$

$$\frac{N}{2^K} = 1 \Rightarrow K = \log_2 N$$

$$T(N) = 2^{\log_2 N} T(1) + 2^{\log_2 N} - 1$$

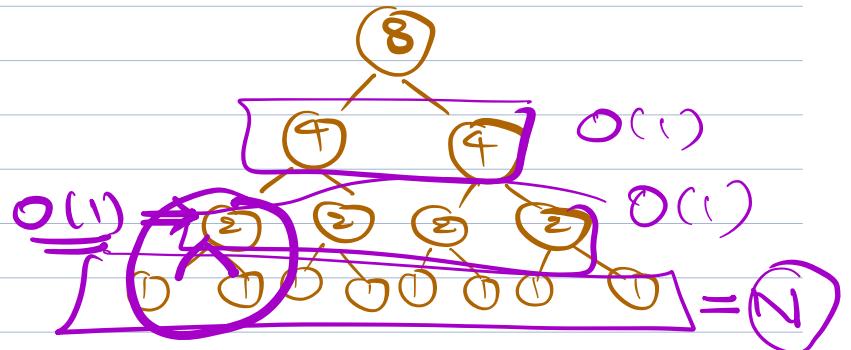
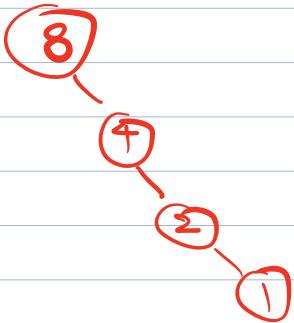
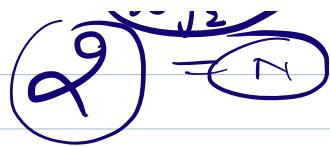
$$\underline{T(N) = N + N - 1}$$

$$T.C. = \underline{\underline{O(N)}}$$

$$\log_2 8 = 3$$

$\log_2 N \times$

$$2^x = N$$



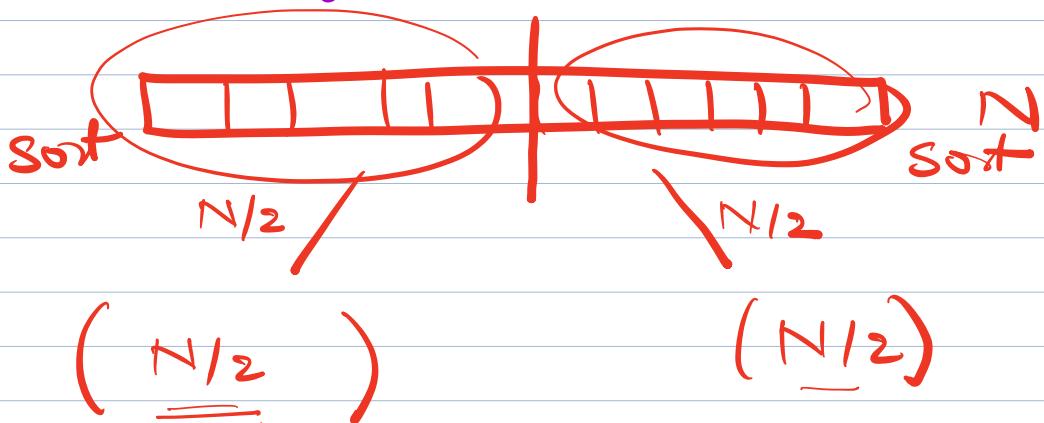
α^2 , $\underline{\alpha}$ $\underline{\alpha}$

H.W

$$T(N) = 2T(N/2) + O(N) \Rightarrow N \log N$$

$T(1) = 1$
 $T(0) = 1$

\Rightarrow Merge Sort (Recursion)



\sqrt{N} $\log N$



Don't

$$x = s;$$

$$y = \underline{x} + \underline{1};$$

$$\text{print}(y) \quad \cancel{s} \checkmark$$

postfix

$$y = \underline{++x};$$

$$\text{print}(y) \rightarrow 6 \checkmark$$

prefix

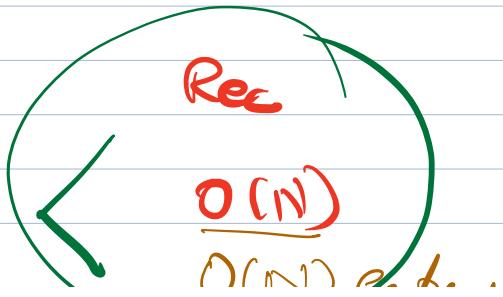
$$x+1$$

I swap
naman

$$\cancel{\text{ayush}} \\ \text{buuya}$$

sum \Rightarrow if

$$\frac{O(N)}{\checkmark}$$



~

UML Y/curve spans

88

if ($B \& z$) == 1)

if (B & z == 1)

~~if ($B \& z$)~~