

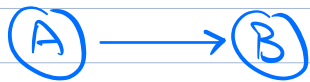


Given N dependent jobs. find the order in which these jobs can be performed.



(1, 2, 3, 4, 5, 6)

or (1, 2, 4, 3, 5, 6)

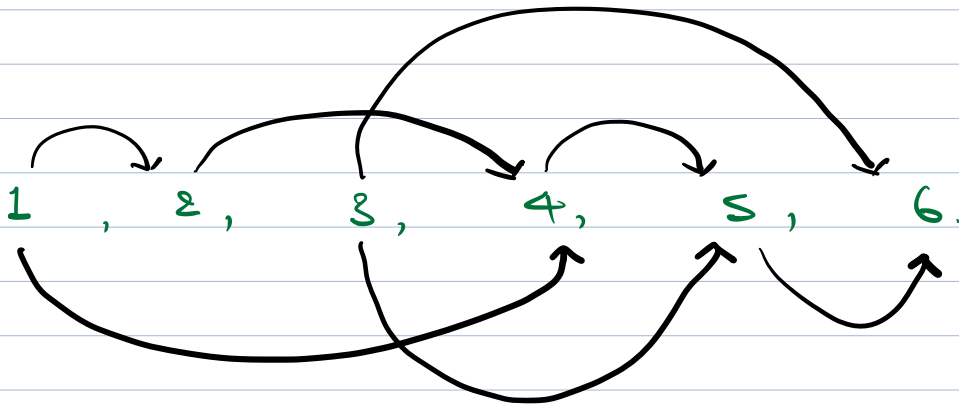


(B is dependent on A)

perform A before you can perform B.

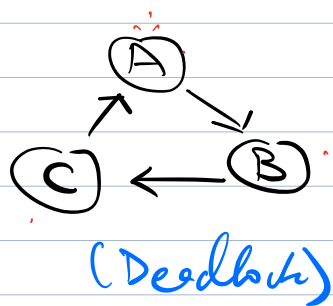
Topological Sort

⇒ Linear ordering of nodes. s.t. if there is a path from the node i to the node j then i is on the left of j



~~xxx~~

Directed. Acyclic Graph

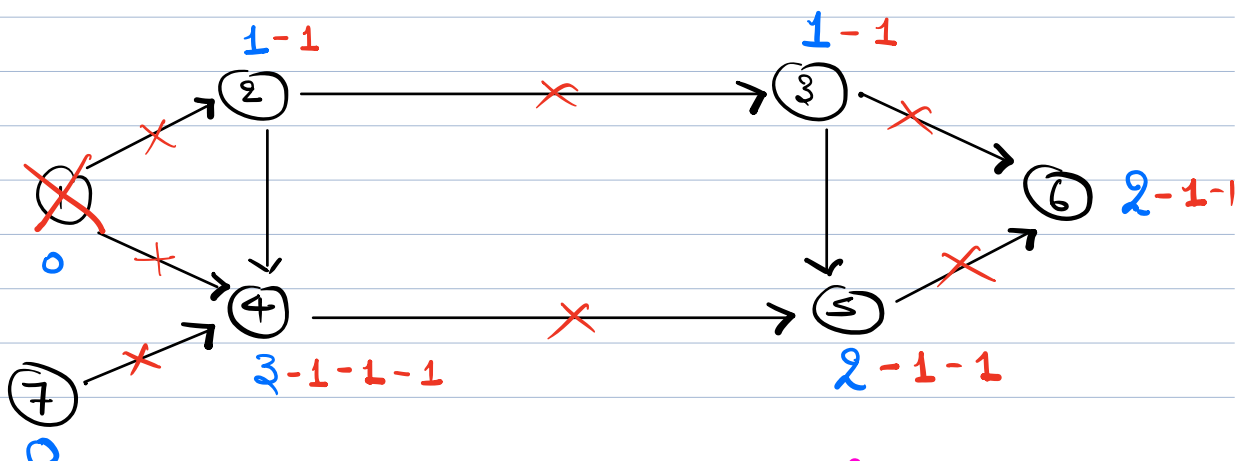


$P1 \rightarrow P2$

* Ways to Perform Topological Sort

1) Using Indegree of A Node

No. of incoming edges.



✓

set (Degree 0)

1 0
 2 ~~1~~ 0
 3 ~~1~~ 0
 4 ~~3~~ ~~2~~ ~~1~~ 0
 5 ~~2~~ ~~1~~ 0
 6 ~~2~~ ~~1~~ 0
 7 0

~~1~~, ~~7~~, ~~2~~
~~3~~, ~~4~~, ~~5~~, ~~6~~

1, 2, 3, 7, 4, 5, 6

Code

↓ Degree (N+1);
 1 → 0
 2 → ~~0~~ 1
 3 → ~~0~~ 1
 4 → ~~0~~ ~~1~~ ~~2~~ 3
 5 → ~~0~~ ~~1~~ 2
 6 → ~~0~~ ~~1~~ 2
 7 → 0

→ 1 → 2, 4
 → 2 → 4, 3
 → 3 → 5, 6
 → 4 → 5
 → 5 → 6
 → 6 → ~~x~~
 → 7 → 4

list < int > order;
 Queue < int > degreeZero;

for (i = 1; i <= N; i++) {

if (Degree[i] == 0) {

degreeZero.enqueue(i);

}

}

while (!degreeZero.isEmpty()) {

```

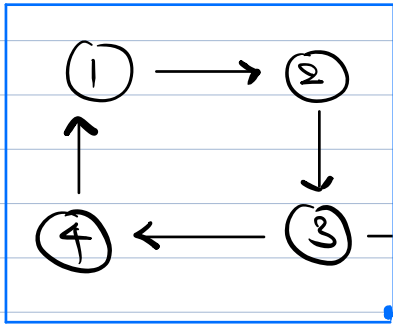
int x = degreeZero.enqueue();
Order.add(x);
for (all u connected to x) {
    Degree[u]--;
    if (Degree[u] == 0) {
        degreeZero.enqueue(u);
    }
}
}

```

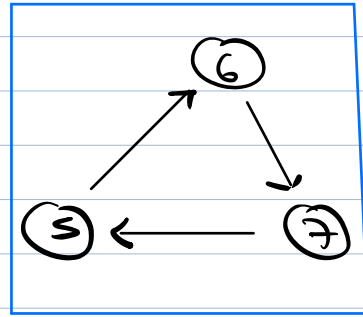
$$\begin{aligned}
 T.C. &= O(V+E) + O(V+E) \\
 &= \underline{\underline{O(V+E)}}
 \end{aligned}$$

Strongly Connected Component

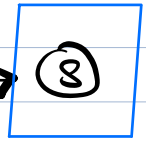
\Rightarrow ^{Path exists} Path from every node to every other node.



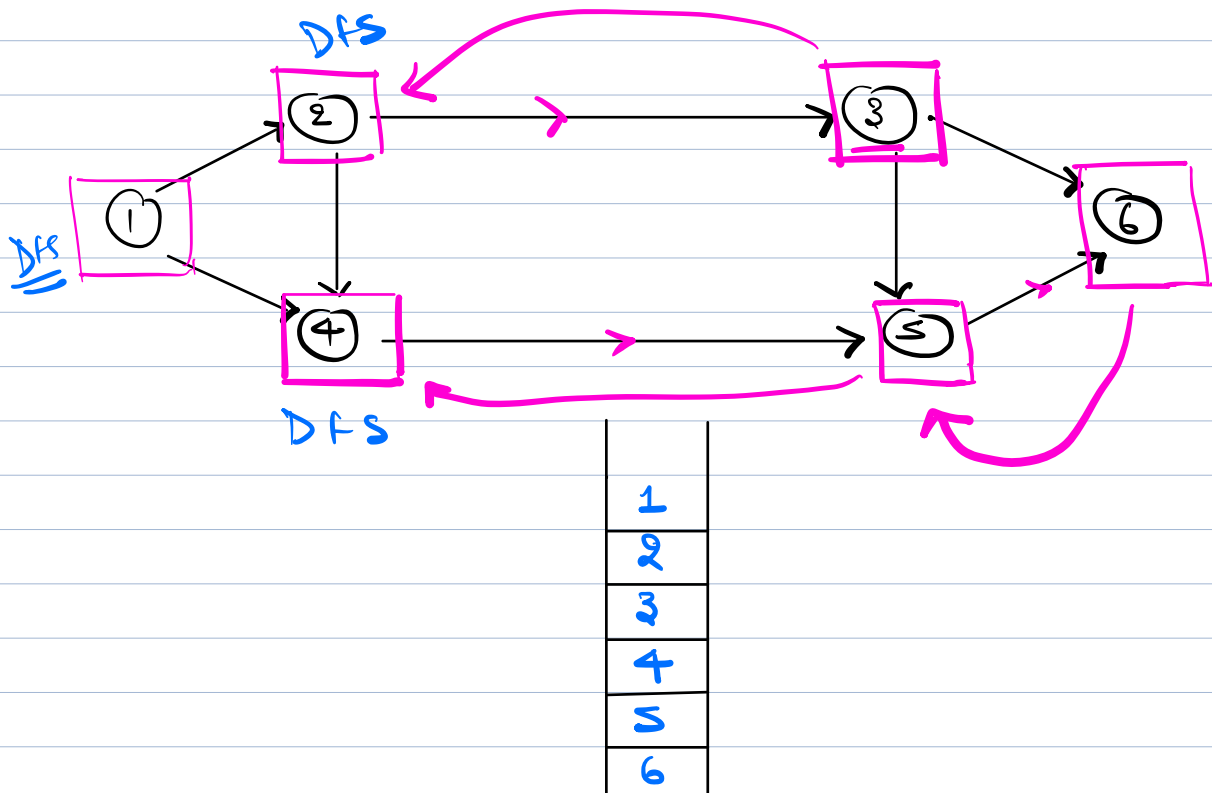
A



B

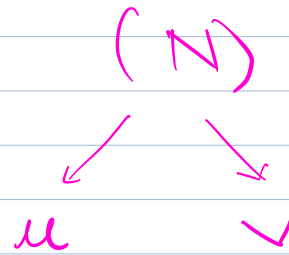
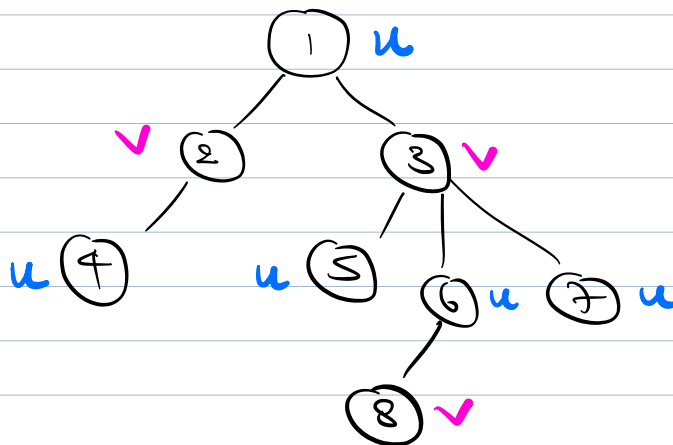


2) Using Stacks (Outdegree)

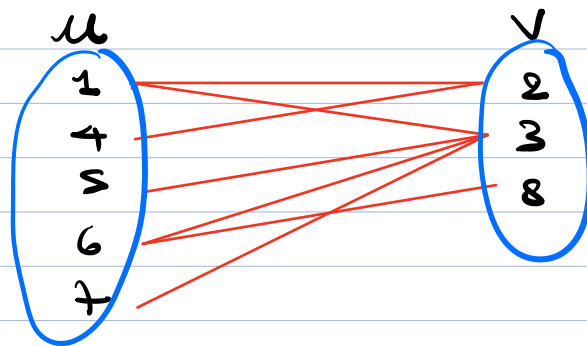


$$T.C. = O(\underline{\underline{V+E}})$$

★ Bipartite Graph



⇒ Bipartite graph



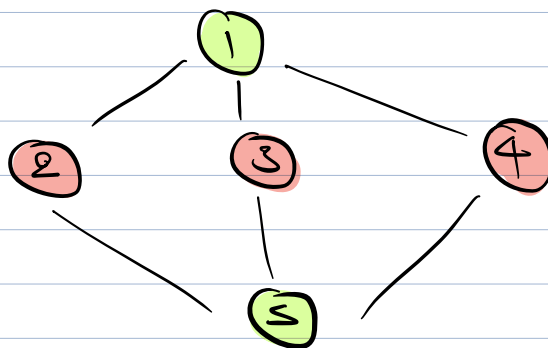
⇒ A graph whose vertices can be divided into 2 disjoint & independent sets (u & v) s.t. every edge connects a vertex in u to a vertex in v .



Given a graph, identify if it is bipartite ??

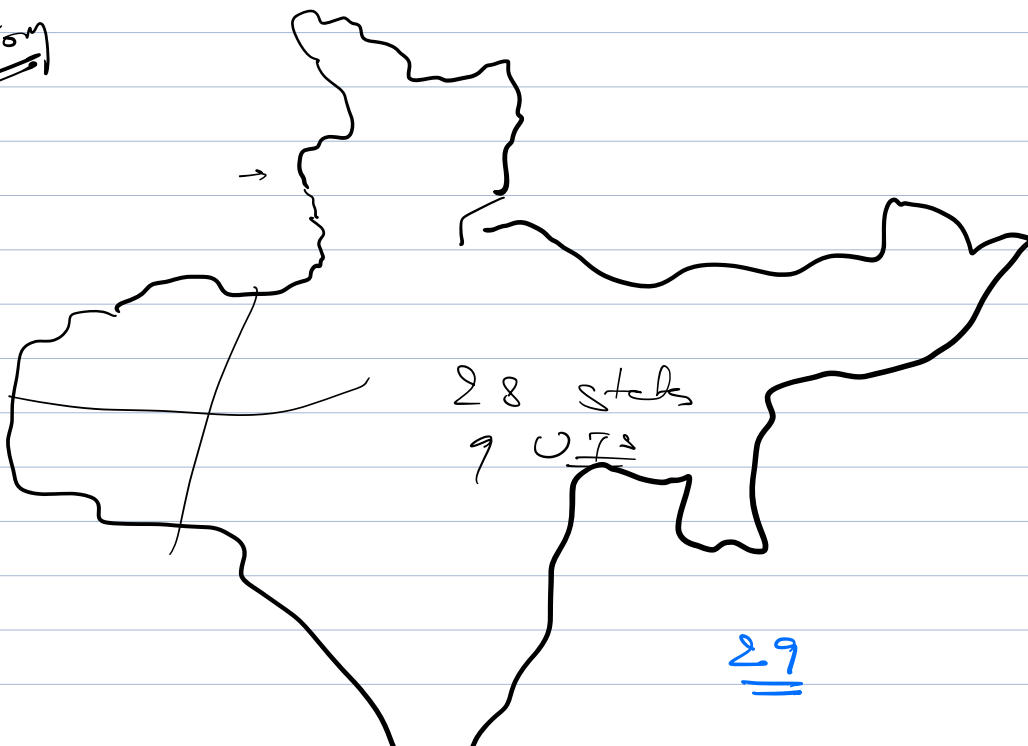
★ Graph Coloring

⇒ Color the nodes of the graph s.t.
no. to adjacent nodes have the same
color. using min no. of colors.



(2)
⇓
Bipartite

History



29

✓

1852 → European Responder

⇒ 4 colors to color a map.

Augustus De Morgan

↳ Set Theory
↳ Mathematical induction.

✓ X ✓ X ✓ X

124 years. ⇒ first proof supported by
a computer using a
simulator.

1990 →

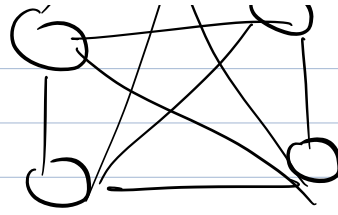
~~Chromatic~~ Chromatic No.

↳ Min no. of colors req. to color a
graph.

Fully connected N nodes

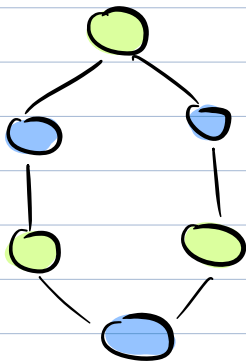


(12)

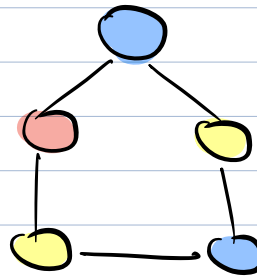


(BackTracking)

Q Chromatic no. of a cycle ??



(2)



(2)

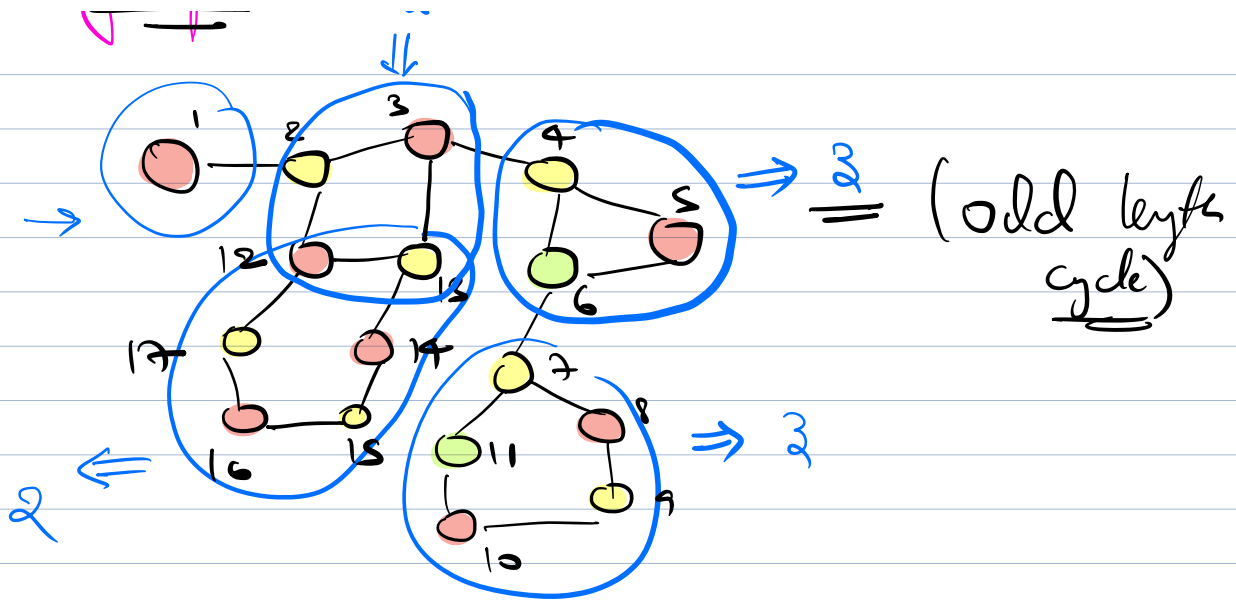
nodes (Even) = 2

nodes (Odd) = 3

Q Chromatic no. of a tree ??

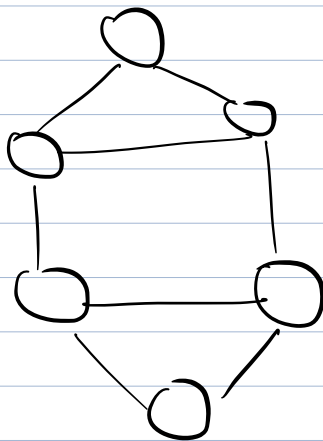
= 2

Graph



No Odd length
Cycle

\Rightarrow Bipartite
graph



H.W.

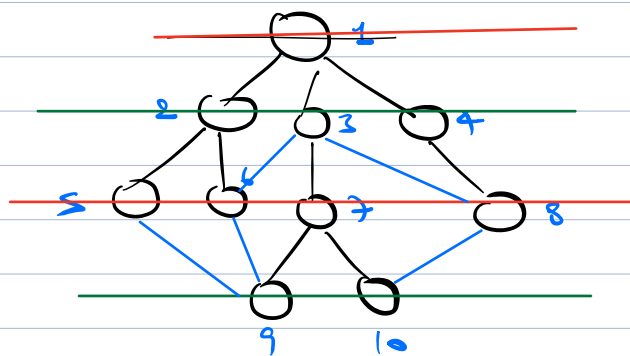
Calculate the length of every cycle.

Hint: Maintain all nodes in the
path (DFS)

Q

Given a tree with N nodes.

find the max no. of edges that can be added to this tree so that it remains bipartite.



$u(\text{odd})$

1
5
6
7
8

X

$v(\text{even})$

2
3
4
9
10

Y

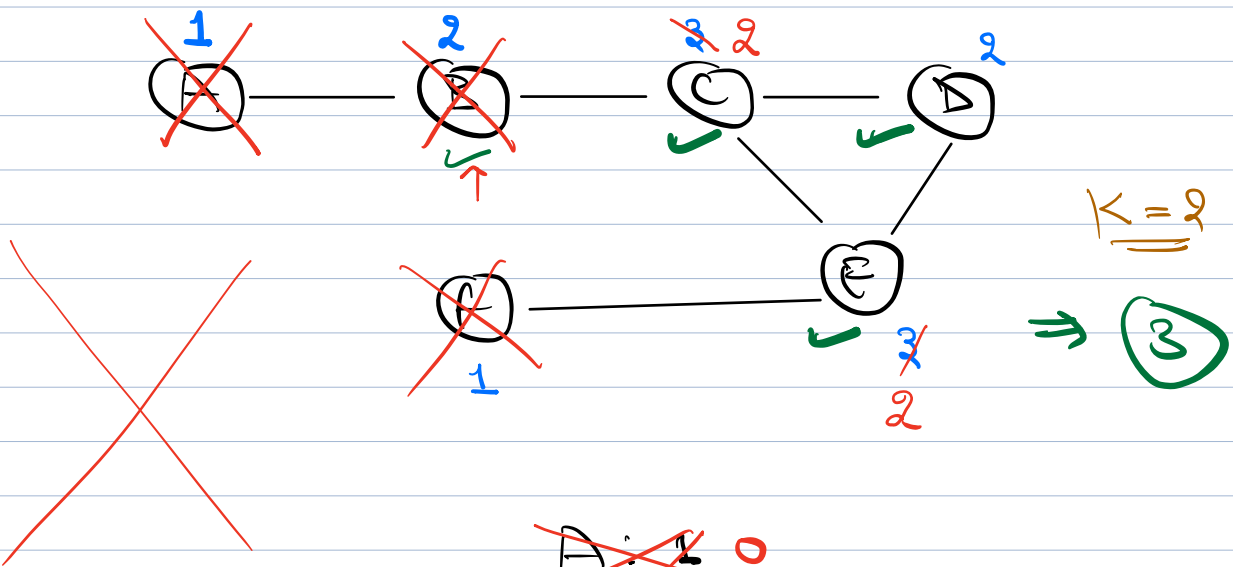
$$(x \times y) - (N-1)$$

Q N people who want to attend a party.

You can only attend a party if K of your friends are already attending the party.

Max no. of people that can attend the party

Given friendship graph (Undirected)



~~A: 1 0~~
~~B: 2 1 < K~~
C: 2 2
D: 2
E: 3 2
~~F: 1~~

~~A, K, S~~

PS Session (Last Session)



Sunday (9:00 AM)



11:00 - 11:30



Gautam