

Direct¹
Code Nation
Gogh.
JP Morgan

Q

Given an array of size N.

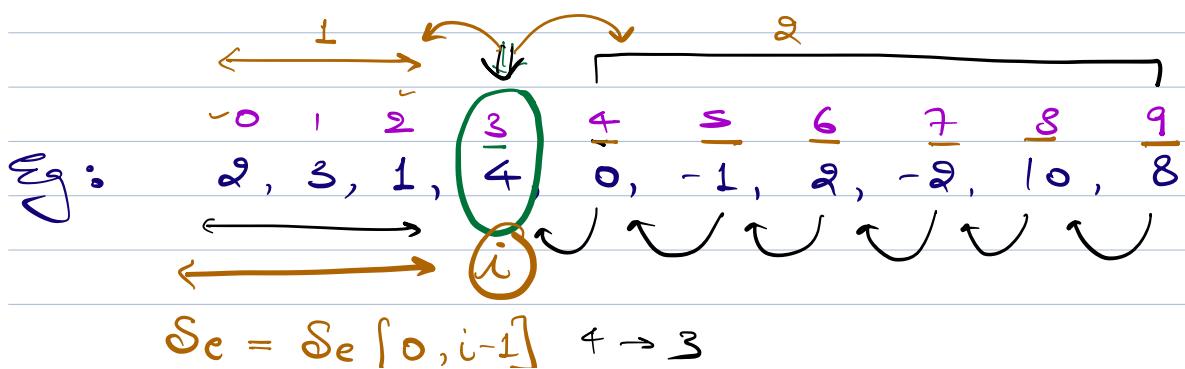
Count the no. of special indexes.

Special index : An index after removing which.

Sum of all odd = sum of all even ..

in the resulting array.

	0 1 2 3 4 5
A :	4, 3, 2, 7, 6, -2
i = 0	3, 2, 7, 6, -2 → $S_e = 8, S_o = 8$ ✓
i = 1	4, 2, 7, 6, -2 $S_e = 9, S_o = 8$ ✗
i = 2	4, 3, 7, 6, -2 $S_e = 9, S_o = 9$ ✓



$$S \rightarrow 4$$

$$G \rightarrow S$$

$$7 \rightarrow 6$$

Index i \Rightarrow after i, all odd indexes becomes even
even becomes odd.

After removal of index i

$$S_e = \underbrace{S_e [0, i-1]}_{\text{Before } i} + \underbrace{S_o [i+1, N-1]}_{\text{After } i}$$

$$S_o = \underbrace{S_o [0, i-1]}_{\text{Before } i} + \underbrace{S_e [i+1, N-1]}_{\text{After } i}$$

$$\underline{\underline{PSe[N-1]}} - \underline{\underline{PSe[i]}}$$

Given an array of size N.

Build an array leftMax. where

leftMax[i] \Rightarrow the maximum value in the array from index 0 to i

A:

0	1	2	3	4	5	6	7	8	9
-3	6	2	4	5	2	8	-9	3	1

LM: -3, 6, 6, 6, 6, 6, 8

$$\underline{M}[0] = A[0];$$

$$\begin{aligned} \underline{LM[1]} &\rightarrow \text{Max}[0, 1] \Rightarrow \text{Max}(\underline{LM[0]}, A[1]) \\ LM[2] &\Rightarrow \text{Max}[0, 2] \Rightarrow \text{Max}(\underline{LM[1]}, A[2]) \end{aligned}$$

$$\underline{LM[s]} \Rightarrow \text{Max}[0, s]$$

$$\underline{LM[6]} \Rightarrow \text{Max}[0, 6] \Rightarrow \text{Max}(\underline{LM[s]}, A[6])$$

LM_{i-1}

$$LM[i] \Rightarrow \max(LM[\underline{i-1}], A[i])$$

1

Code

L_M[o] = A[o]; // Necessary ??

for ($i = 1$; $i < N$; $i++$) {

$$LMT[i] = \max(LMT[i-1], A[i]);$$

1

$$T.C. = O(N)$$

$$S.C. = O(N)$$

Right Max [i] \Rightarrow Max of all elements from

$$A: -3, \underline{6}, \underline{2}, \underline{4}, \underline{5}, \underline{2}, \underline{\frac{6}{8}}, \underline{-9}, \underline{\frac{3}{3}}, \underline{1}$$

RM : 8 8 8 8 8 8 8 , 3 , 3 , 1 ↙

$$RM[N-1] = \underline{A[N-1]}$$

$$RM[N-2] \\ (RM[8]) = \max[N-2, N-1] \Rightarrow \max(RM[N-1], A[N-2])$$

$$RM[i] \Rightarrow \max[i, N-1]$$

$$\Rightarrow \max(\underline{RM[i+1]}, A[i])$$

Code

i ??

N-1
i+1 = N

$$RM[N-1] = \underline{A[N-1]}$$

for (i = N-2; i >= 0; i--) {

$$RM[i] = \max(RM[i+1], A[i]);$$

}

Google

Given a string of lowercase alphabet

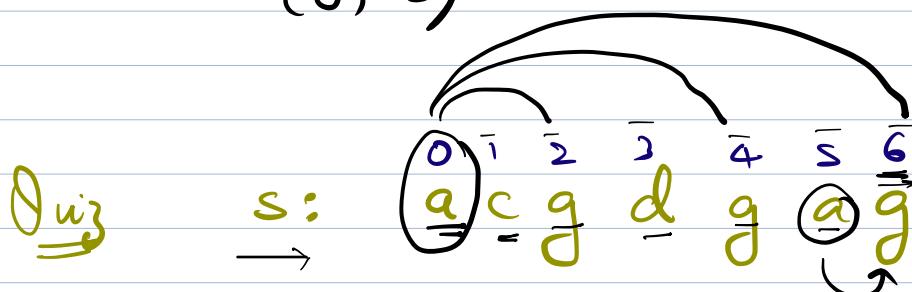
Return the count of pairs (i, j). s.t.

i < j

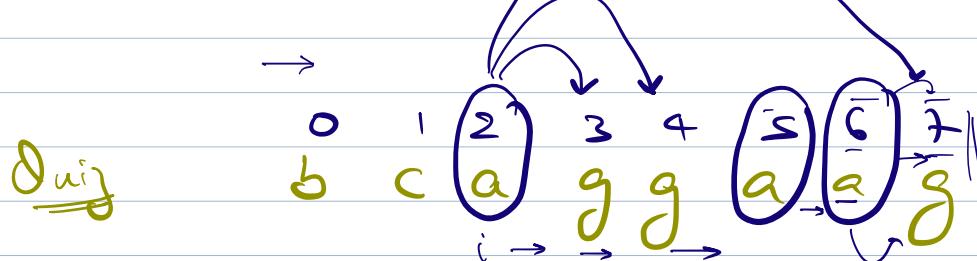
$$s[i] = 'a' \\ s[j] = 'g'$$

$s:$ $\begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 \end{matrix}$

$$(0, 3) \\ (0, 5) \Rightarrow 3$$



$$(0, 2) \\ (0, 4) \\ (0, 6) \Rightarrow 4$$



$$(2, 3) \\ (2, 4) \\ (2, 7) \Rightarrow 5$$

2) Brute force

$a b c d e \}^{N-2} \binom{N-1}{1, 1, 1, 1, 1}$

Code

$i=1$

$i=N-2$

$ans = 0;$

$\text{for } (i=0; i < N; i++) \{$

$j=N$

$\Rightarrow \text{if } (s[i] == 'a') \{$

$\quad \quad \quad \text{for } (j=i+1; j < N; j++) \{$

$\quad \quad \quad \Rightarrow \text{if } (s[j] == 'g') \{$

$\quad \quad \quad \quad \quad \quad ans++;$

$\}$

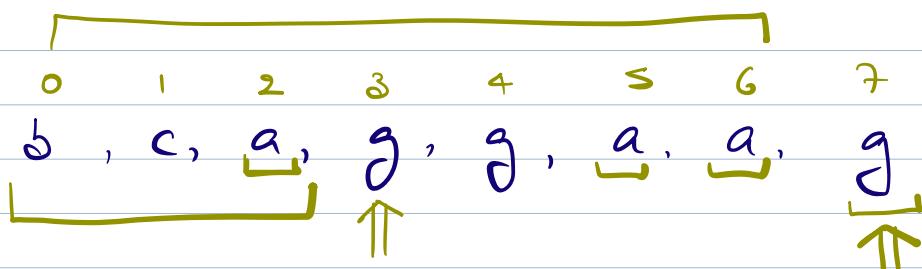
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T.C. = $O(N^2)$

S.C. = $O(1)$

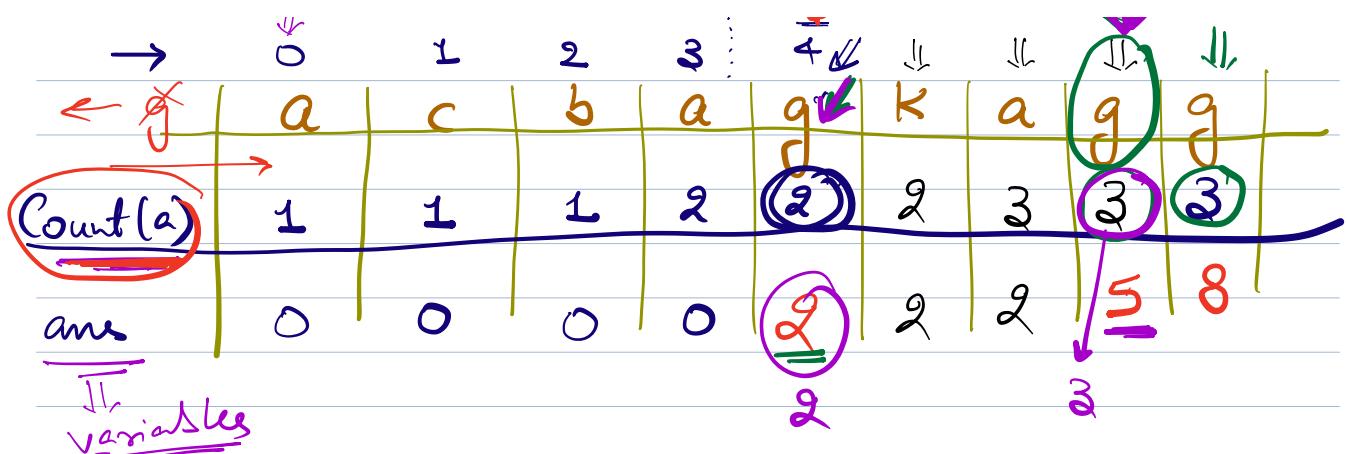


\Rightarrow Every 'g' will make a valid pair with all the 'a' on the left of it



11





T.C. = $O(N)$

S.C. = ~~$O(\underline{\underline{N}})$~~ $O(\underline{\underline{1}})$

Code

$ans = 0;$

$Counta = 0;$

for ($i=0; i < N; i++$) {

 if ($s[i] == 'a'$) {
 $Counta++$;

 } else if ($s[i] == 'g'$) {

$ans = ans + Counta$;

}

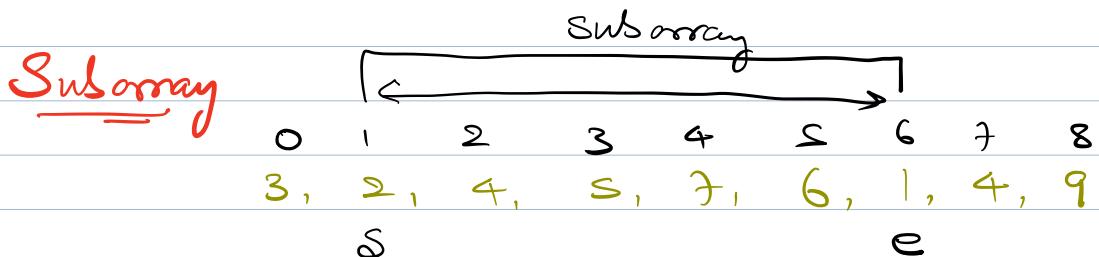
return ans ;

Amazon

Q

Given an array of size N .
Return the length of the smallest
subarray which contains both the min

& the max of the array.

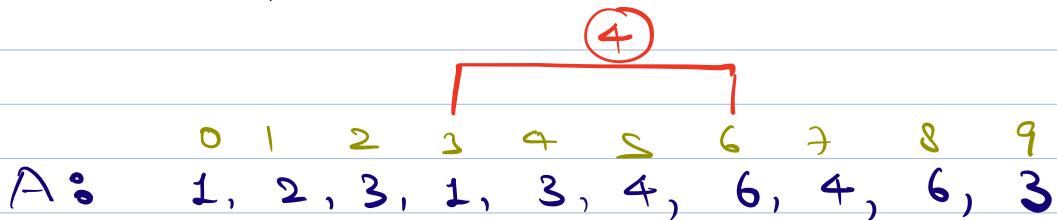


s
 1

e
 6

\Rightarrow 2 variables

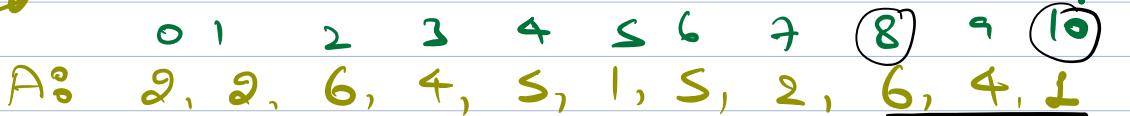
$$\Rightarrow \begin{bmatrix} s & e \\ l & r \end{bmatrix} \Rightarrow e - s + 1 \Rightarrow \text{total no. of elements}$$



$$\text{Max} = 6$$

$$\text{Min} = 1$$

Ques



$$\text{Max} = 6$$

$$\text{Min} = 1$$

Q

A :

 $A = [1, 6, 4, 2, 7, 7, 5, 1, 3, 1, 1, 5]$

$$\begin{aligned} \text{Max} &= 7 & \Rightarrow (3) \\ \text{Min} &= 1 \end{aligned}$$

$A = [1, 8, 8, 8, 8, 8, 8]$

$$\text{Max} = 8 \quad \checkmark$$

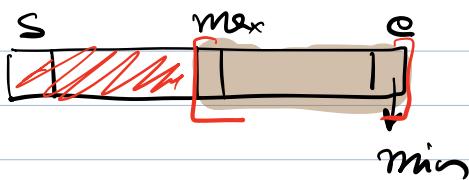
$$\text{Min} = 8 \quad \checkmark$$

$s == e$

$[49, 49]$ \Rightarrow 1 element

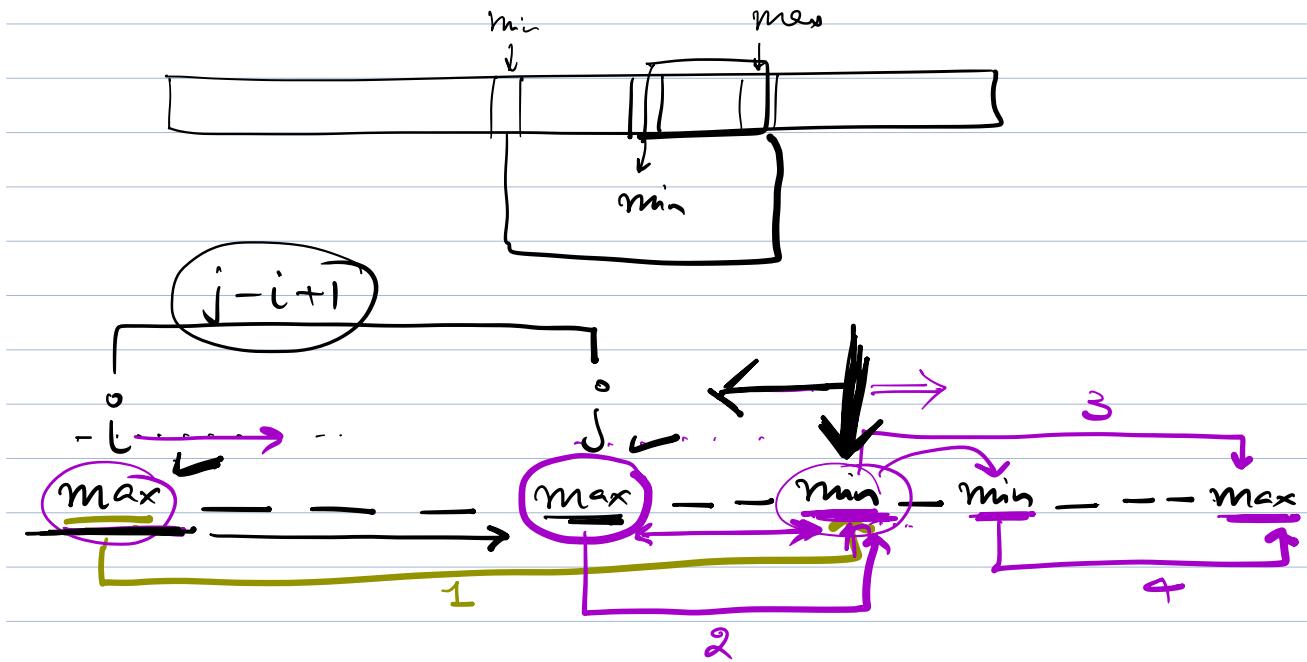
Observations

1) An subarray will always have max / min in the corner.



2) There will only be one min & one max

present in the 0 ans subarray.



① Brute force

$O(n)$

1) find max & min

2) iterate from L to R

$\forall \text{max}$, find nearest min in right

$\forall \text{min}$, find nearest max in right

Code

ans = N

// for every min \rightarrow find max in the right

// for every max \rightarrow find nearest min ...

\Rightarrow for ($i = 0$; $i < N$; $i++$) {

if ($A[i] == A_{\min}$) {

\Rightarrow for ($j = i$; $j < N$; $j++$) {

if ($A[j] == A_{\max}$) {

$ans = \min (ans, j - i + 1);$
break;

}

}

else if ($A[i] == A_{\max}$) {

for ($j = i$; $j < N$; $j++$) {

if ($A[j] == A_{\min}$) {

$ans = \min (ans, j - i + 1);$
break;

}

}

}

$$T.C = O(N^2)$$

A: 9, 9, 9, 9, 9, 9, 9, 9, 9, 9, 9, 1

Doubts

$$A \rightarrow \mathbb{Z}$$

$$B \rightarrow M \times 2$$



Sum of \leftarrow rays

$1 \rightarrow$ indeed

\mathbb{Z}

$\mathbb{Z} + 1$

\mathbb{O}

$A^{\circ} : 1, 2, 3, 4, \dots$

$\mathbb{Z} - 1$

$\mathbb{Z} \Rightarrow \mathbb{Z} + 1$

$A^{\circ} : \cancel{0} \quad 0$

\mathbb{Z}

\mathbb{L}

B

\longleftrightarrow

\longleftrightarrow

$A^{\circ} : 1, 4, 7, 11, 21, 35, 66 \dots \mathbb{Z}$

$B = 3$

$\mathbb{Z} \rightarrow \mathbb{Z} + 1$

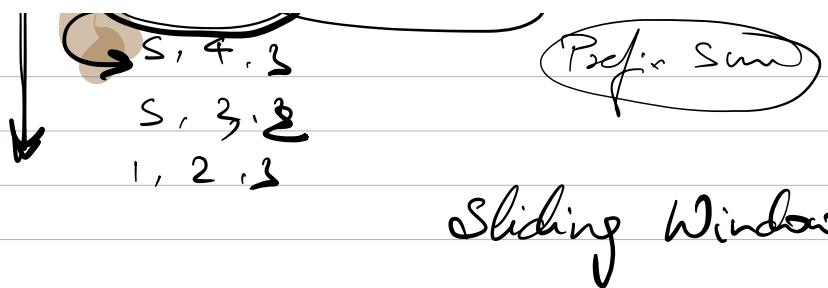
Rw

$R \sim S$

$S, 4, 2 - 2 + 3$

$S, 4, 2$

2



1, 2, 3, 4, 5, 6

1 3 6 10 15 21

21 20 18 15 11 6

