

Prefix Sum, Carry forward technique

Subarrays.

⇒ Contiguous part of an array (l , r)

$$\underline{r > l}$$

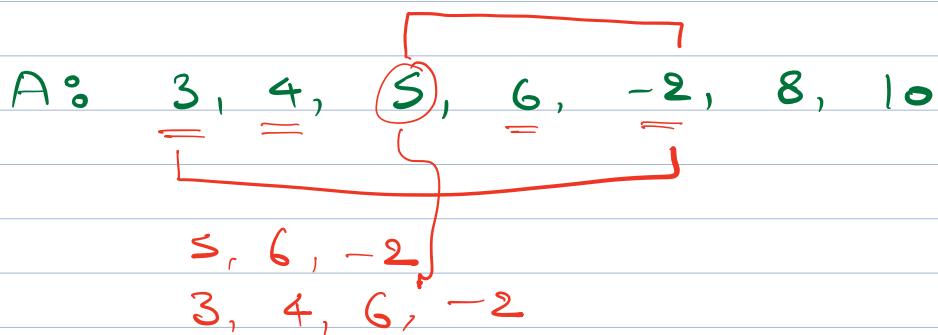
⇒ Largest subarray $l = 0$
 $r = N - 1$

Complete array is also a subarray of itself

⇒ Single element can also be a subarray.

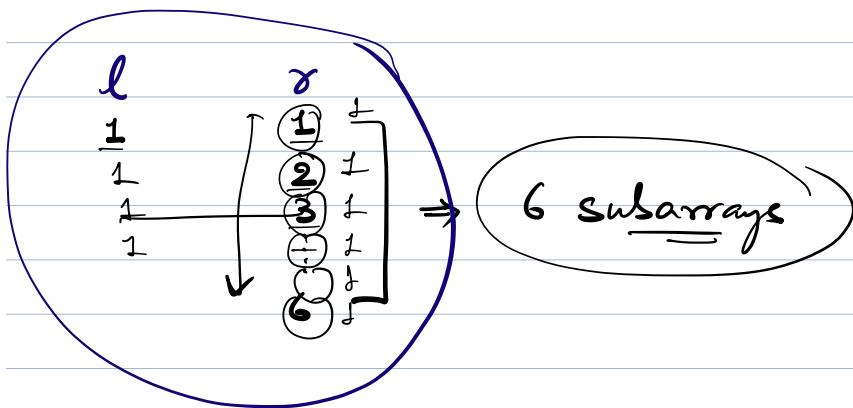
⇒ Empty array is also a subarray

⇒ But we will only talk about non empty subarray.



Ques

A : 0, 1, 2, 3, 4, 5, 6
A : 4, 2, 10, 3, 12, -2, 15



No. of subarrays in the given array.

$A = [4, 2, 10, 3, 12, -2, 15]$ \Rightarrow Tel.

Starts at Ends at

$\Rightarrow 0$

$[0, 6] \Rightarrow 7$

$4, [4, 2], [4, 2, 10], [4, 2, 10, 3]$

1	$[1, 6]$	$\Rightarrow 6$
2	$[2, 6]$	$\Rightarrow 5$
3	$[3, 6]$	$\Rightarrow 4$
4	$[4, 6]$	$\Rightarrow 3$
5	$[5, 6]$	$\Rightarrow 2$
6	$[6, 6]$	$\Rightarrow 1$
\sum		
28		

$$\frac{7 \times 8}{2} = \frac{56}{2} = 28$$

Θ = No. of non Empty subarrays in an array
of size N

$$\# \text{ Subarrays} = \# \text{ sa. starting at } 0 = Z + \dots$$


 " " " 1 = Z + 1
 " " " 2 = Z + 2
 " " " 3 = Z + 3
 ...
 " " " n = Z + n

$\# \text{ subarrays starting at } i = 1$

$$1 + 2 + 3 + \dots + N = \frac{N(N+1)}{2}$$

$$\# \text{ No. of subarrays} = \frac{N \times (N+1)}{2}$$

Q Print all values of a subarray.
(s, c)

Code

✓ void printSubarray (A[], s, e) {

P --- > V . . .

→ for ($i = \underline{s}$; $i \leq \underline{e}$; $i++$) {

 print (A[i]); } ↗

L ↘

Q Find the sum of a given subarray.

Q Print all subarrays of a given array of size N.

A : $\begin{matrix} 0 & 1 & 2 \\ 2, & 8, & 9 \end{matrix}$

s	e	
0	0	[2] ↗
0	1	[2, 8] ↗
0	2	[2, 8, 9] ↗
1	1	[8] ↗
1	2	[8, 9] ↗
2	2	[9] ↗

Code

sc → [(for ($i = \underline{0}$; $i < \underline{N}$; $i++$) { }] $\Rightarrow O(n^2)$

ei → [for ($i = i$; $i < N$; $i++$) { }] ↗

'D' \cup - 'U' \cup
// Subarray $i \underline{\rightarrow} j$

\Rightarrow [for ($K = i$; $K \leq j$; $K++$) {
 Print A[K];
}

\$

$$T.C. = O(n^3)$$

$$\frac{O(n^2)}{\text{# of subarray}} \times \frac{O(n)}{\text{Print}}$$

$$\begin{matrix} i=0 \\ j=0 \end{matrix} \quad j=1 \quad j=2 \quad j=3 \Rightarrow \Sigma$$

$$\text{Print} \quad \text{Print} \quad \text{Print} \quad \text{Print} \Rightarrow O(n)$$

$$\begin{matrix} i=1 \\ j \Rightarrow \underline{\Sigma} \end{matrix}$$

$$\begin{matrix} i=2 \\ j \Rightarrow \underline{\Sigma-2} \end{matrix}$$

1

O Print the sum of every single
subarray.

A: 3, 2, -1, 4
0 1 2 3

s	e	
0	0	[3] \Rightarrow 3
0	1	[3, 2] \Rightarrow 5
0	2	[3, 2, -1] \Rightarrow 4
0	3	[3, 2, -1, 4] \Rightarrow 8

Code

O(n²) [for (i = 0; i < N; i++) { // fix i
for (j = i; j < N; j++) {

Sum = 0;

= for (k = i; k <= j; k++) {

Sum += A[k];

O(1)

prefix
sum

}
Print (sum);

5

$$\text{T.C.} = O(n^3) \xrightarrow{\text{P.S.array}} O(n^2) + O(n)$$

S.C. = $O(n)$

$$\# \text{subarray} = \frac{(n)(n+1)}{2} = \frac{n^2+n}{2}$$

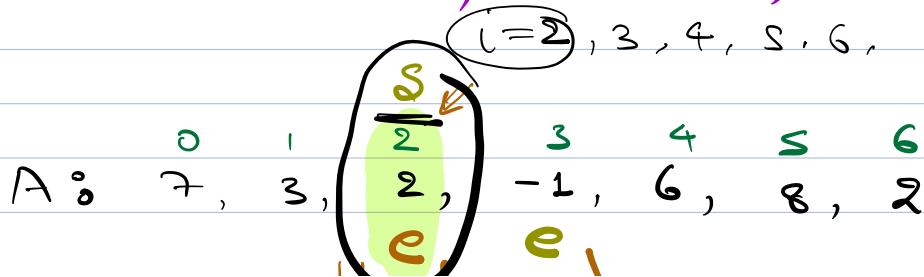
$$O\left(\frac{n^2+n}{2}\right) = \underline{\underline{O(n^2)}}$$

$$1, 1, 1, 1, 1, 1, \dots, \frac{n^2}{2} \Rightarrow O(n^2)$$

Θ Print sum of all subarrays starting from the index 3 (No extra space allowed)

$$A : 0, 1, 2, 3, 4, 5, 6$$

$$\text{T.C.} = O(n^2) \Rightarrow O(n)$$



	s	e	
2	2	<u>a[2]</u>	
2	3	<u>a[2] + a[3]</u>	
2	4	<u>a[2] + a[3] + a[4]</u>	
2	5	<u>a[2] + a[3] + a[4] + a[5]</u>	

Code

$\Rightarrow s = 2; \quad \leftarrow \leftarrow \rightarrow$

```

Sum = 0;
for (j = s; j < N; j++) {
    Sum = Sum + arr[j];
    print (sum);
}

```

Code - 2

$[i, i] \Rightarrow$ single element sum.

\Rightarrow for (i = 0; i < N; i++) {

```

Sum = 0;
for (j = i; j < N; j++) {
    Sum = Sum + A[j];
    print (sum);
}

```



Google
FB.
Amazon

Q

Given an array of size N .
Find the sum of all subarray sums.

A : 1, 2, 3

s	c	Sum
0	0	$\frac{1}{1} \rightarrow a[0]$
0	1	$\frac{1}{3} \rightarrow a[0] + a[1]$
0	2	$\frac{1}{6} \rightarrow a[0] + a[1] + a[2]$
1	1	$\frac{1}{2} \rightarrow a[1]$
1	2	$\frac{1}{5} \rightarrow a[1] + a[2]$
2	2	$\frac{1}{3} \rightarrow a[2]$
$\sum = 20$		$3 \times a[0] + 4 \times a[1] + 3 \times a[2]$

① Brute force

total Sum = 0;

for (i = 0; i < N; i++) {

Sum = 0;
for (j = i; j < N; j++) {
Sum = Sum + A[j]; } } \Rightarrow
totalSum = totalSum + Sum;

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print (totalSum);

T.C. = $O(N^2)$ rejected.

$$\frac{O(N)}{\text{↓}} \Rightarrow O(L)$$

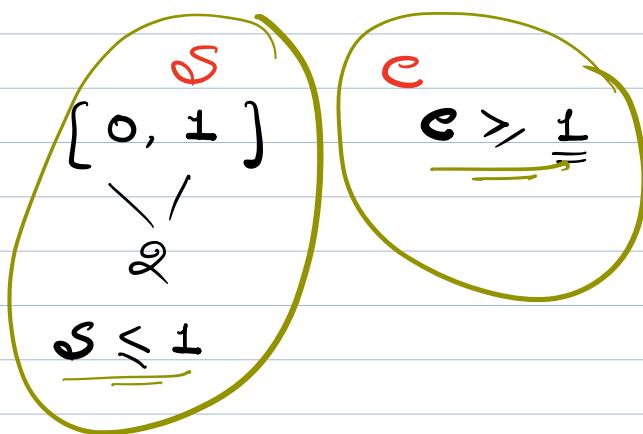
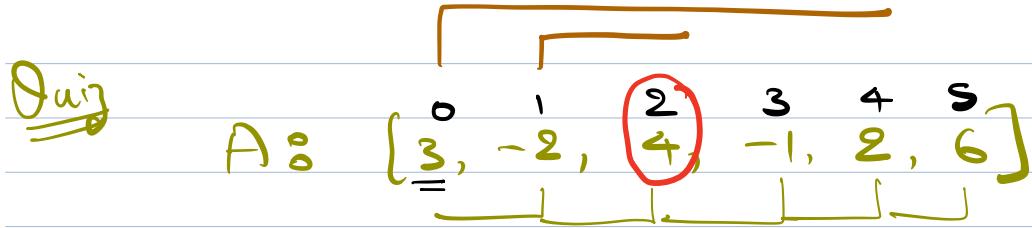
no. of elements present
in the array.

Contribution Technique

→ If we are able to calculate the contribution of each element towards the sum.

Contribution = # times the element is occurring
×
the value of the element.

Total no. of subarrays, a given element
is a part of



\Rightarrow # subarrays of which an element or index i will be a part of

~~total~~ $(s \leq i) \underline{88} (e \geq i)$

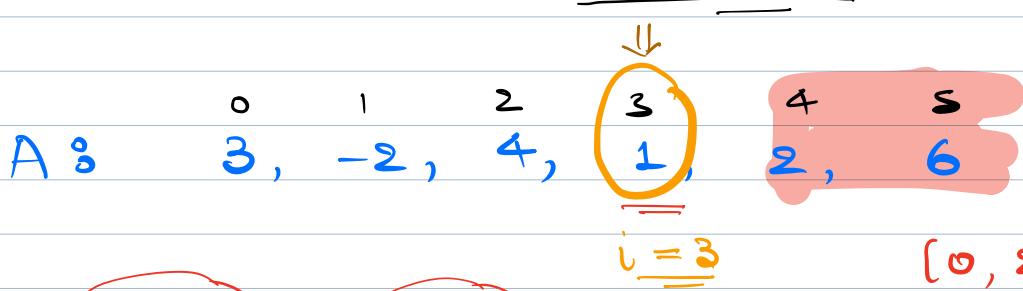
$$[0-i] \quad [i^l, N^{-1}] \quad \Rightarrow (r-l+1)$$

$$(i+1) \quad \left(\frac{N-1}{l} - \frac{i+1}{i} \right) = \underline{(N-i)}$$

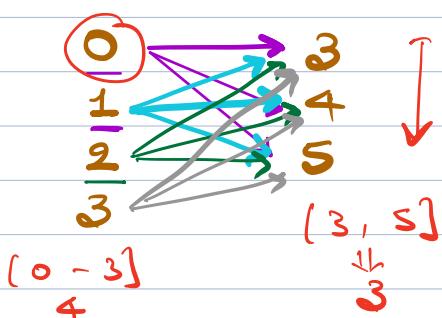
$$\underline{(i+1)} \quad \underline{(N-i)}$$

$$\text{Total no. of subarrays} = |s| \times |e|$$

$$= \underline{(i+1)} \times \underline{(N-i)} \Rightarrow$$



$s \leq i$ $e > i$



$s=0$	$s=1$
$[0, 3]$	$[1, 3]$
$[0, 4]$	$[1, 4]$
$[0, 5]$	$[1, 5]$
$[2, 3]$	$[2, 5]$
$[2, 4]$	$[3, 4]$
$[2, 5]$	$[3, 5]$

Code

$\text{sum} = 0;$

$\text{for } (i=0; i < N; i++) \{$

$s = (i+1);$
 $e = (N-i);$

$\text{count} = s \times e;$

$\text{contribution} = \text{count} \times A[i];$

Sum = Sum + contribution;

5

$$T.C. = O(N)$$

$$S.C. = O(1)$$

Doubts

A: $\begin{matrix} 0 & 1 & 2 \\ 1, 2, 3 \end{matrix}$

$\begin{array}{lll} \downarrow \\ - [1] & [2] & [3] \\ - [1, 2] & [2, 3] \\ - [1, 2, 3] \end{array}$

