

## Agenda $\Rightarrow$ 2D matrices

- 1) Submatrix sum Queries
- 2) Sum of all submatrices
- 3) Max Submatrix Sum.

$\Theta$  Given a matrix of size  $N \times M$ .  $\neq \Theta$  queries.

$\approx$  Amazon for each query, find the sum of given

Submatrix.  $c_1$        $c_2$

|                   | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
|-------------------|---|---|---|---|---|---|---|
| 0                 | 0 |   |   |   |   |   |   |
| 1                 | 1 |   |   |   |   |   |   |
| $r_1 \rightarrow$ | 2 |   |   |   |   |   |   |
| 3                 |   |   |   |   |   |   |   |
| 4                 |   |   |   |   |   |   |   |
| $r_2 \rightarrow$ | 5 |   |   |   |   |   |   |
|                   |   |   |   |   |   |   |   |

$6 \times 7$

$\nabla$  identify a submatrix.

(Top Left)

(Bottom Right)

$(r_1, c_1)$

$(r_2, c_2)$

|                 | 0  | 1  | 2  | 3  | 4  | 5  | 6 |
|-----------------|----|----|----|----|----|----|---|
| $\Rightarrow 0$ | 7  | 1  | -6 | 3  | 12 | -2 | 4 |
| -1              | 10 | 5  | -2 | 0  | 9  | 4  | 8 |
| 2               | 6  | 4  | -3 | 8  | 11 | 3  | 2 |
| 3               | 13 | -8 | -5 | 12 | 4  | 6  | 7 |
| -4              | 3  | 2  | 1  | 9  | 3  | 9  | 6 |
| 5               | 4  | 3  | -2 | 6  | 8  | 8  | 8 |

$6 \times 7$

$\Rightarrow 0,0 \text{ to } 5,6$

Q

TL  
( $\tau_1$ ,  $c_1$ )

BR  
( $\tau_2$ ,  $c_2$ )

Q1

(1, 2)

(4, 3)  $\Rightarrow$  20

Q2

(1, 1)

(4, 2)  $\Rightarrow$  72

## Sol<sup>n</sup> ① Brute force

sum = 0;

for ( $i = \tau_1$ ;  $i \leq \tau_2$ ;  $i++$ )  $\nmid$  // All rows

for ( $j = c_1$ ;  $j \leq c_2$ ;  $j++$ )  $\nmid$

    sum += M[i][j];

b

T.C of 1 query  $\Rightarrow$   $O(N \times M)$   $\Rightarrow$  O(1)  
" 1 queries  $\Rightarrow$  O(1 \times N \times M)

## 2) Can we optimise ??

1D Array  $\Rightarrow$  prefix sum array.

PF[i]  $\Rightarrow$  sum of all elements from index 0 to i.

$\Rightarrow$  Matrix  $\Rightarrow$  Prefix Sum Matrix

$\underline{PF[i][j]}$   $\Rightarrow$  sum of all elements  
in submatrix from  
 $(0, 0)$  to  $i, j$

|   | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
|---|---|---|---|---|---|---|---|
| 0 | 0 |   |   |   |   |   |   |
| 1 |   |   |   |   |   |   |   |
| 2 |   |   |   |   |   |   |   |
| 3 |   |   |   |   |   |   |   |
| 4 |   |   |   |   |   |   |   |
| 5 |   |   |   |   |   |   |   |

$6 \times 7$

$PF[3, 5]$

Assume we already have a prefix sum matrix.  $C_1$

|   | 0 | 1 | 2 | 3 | 4 | 5 | 6               |
|---|---|---|---|---|---|---|-----------------|
| 0 | 0 | / | / | / | / | / | x $\Rightarrow$ |
| 1 | x | 0 | / | / | / | / | x               |
| 2 | x | x | 0 | / | / | / | x               |
| 3 | x | x | x | 0 | / | / | x               |
| 4 | x | x | x | x | 0 | / | x               |
| 5 | x | x | x | x | x | 0 | x               |
| 6 | x | x | x | x | x | x | x               |

$6 \times 7$

$\Rightarrow M$

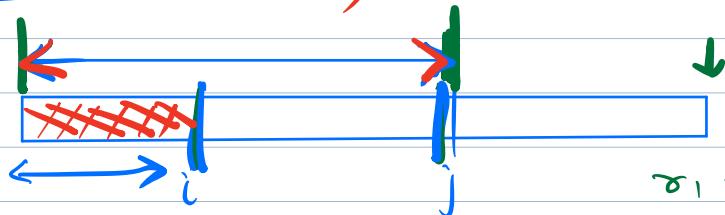
|   | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
|---|---|---|---|---|---|---|---|
| 0 |   |   |   |   |   |   |   |
| 1 |   |   |   |   |   |   |   |
| 2 |   |   |   |   |   |   |   |
| 3 |   |   |   |   |   |   |   |
| 4 |   |   |   |   |   |   |   |

$6 \times 7$

$\Rightarrow PF$

$s$  | | | | | | |

Sum  $(\underline{2}, \underline{2})$  to  $(\underline{4}, \underline{5})$ )



$$\tau_1 = 2$$

$$\tau_2 = 4$$

$$c_1 = 2$$

$$c_2 = 5$$

$$PS[j] - PS[i-1]$$

$$(PF[\underline{4}][\underline{5}] - PF[\underline{1}][\underline{5}] - PF[\underline{4}][\underline{2}] + PF[\underline{1}][\underline{1}])$$

|   | 0 | M | 2 |
|---|---|---|---|
| 0 | 3 | 4 | 1 |
| 1 | 6 | 2 | 9 |
| 2 | ≤ | 3 | 1 |

|   | 0  | P.F | 2  |
|---|----|-----|----|
| 0 | 3  | 7   | 8  |
| 1 | 9  | 15  | 25 |
| 2 | 14 | 23  | 34 |

$$PF[0][0] \rightarrow (0,0) \text{ to } (0,0)$$

$$PF[0][1] \rightarrow (0,0) \text{ to } (0,1)$$

$$PF[0][2] \rightarrow (0,0) \text{ to } (0,2)$$

$$PF[1][0] \rightarrow (0,0) \text{ to } (1,0)$$

$$PF[1][1] \rightarrow (0,0) \text{ to } (1,1)$$

$$PF[1][2] \rightarrow (0,0) \text{ to } (1,2)$$

$$PF[2][0] \rightarrow (0,0) \text{ to } (2,0)$$

$$PF[2][1] \rightarrow (0,0) \text{ to } (2,1)$$

$$PF[2][2] \rightarrow (0,0) \text{ to } (2,2)$$

$\delta_{N \times M}$

Using a prefix matrix, we can calculate

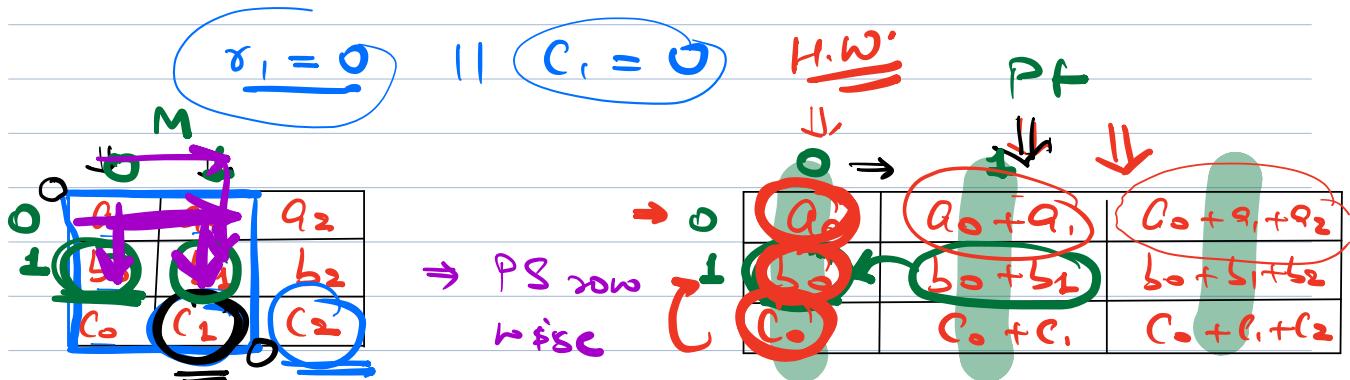
each query in O(1) time

$$T.C. = \text{ } \bigcirc \text{ } \frac{1 \times \delta}{\text{PF matrix}} + \frac{\text{Time to calculate}}{\text{matrix}}$$

Edge Cases  $\Rightarrow$  if ( $x_1 == 0$ ) || ( $c_2 == 0$ )

$$\text{Sum} \left( \begin{matrix} (\tau_1, c_1) & \rightarrow \\ (\tau_2, c_2) & \end{matrix} \right) = \text{PF}[\underline{\tau_2}][\underline{c_2}] - \text{PF}[\underline{\tau_1-1}][\underline{c_2}] - \text{PF}[\underline{\tau_2}][\underline{c_{1-1}}] + \text{PF}[\underline{\tau_1-1}][\underline{c_{1-1}}]$$

$$\delta \Rightarrow O(\epsilon)$$



| $i=0$              | $a_0$             | $a_0 + a_1$  | $a_0 + a_1 + a_2$  |
|--------------------|-------------------|--|--|
| $\rightarrow$      | $a_0 + b_0$       | $a_0 + b_1 + b_0 + b_1$  | $a_0 + a_1 + b_0 + b_1 + b_2$  |
| $\curvearrowright$ | $a_0 + b_0 + c_0$ | $a_0 + \underline{a_1} + \underline{b_0} + \underline{b_1} + \underline{c_0}$<br>$+ c_1$ | $a_0 + \underline{a_1} + \underline{a_2} + \underline{b_0} + \underline{b_2}$<br>$+ b_2 + c_0 + c_1 + c_2$ |

# Code

13

NXM

PF [N] [M]  
row wise

O( $2 \times N \times M$ )

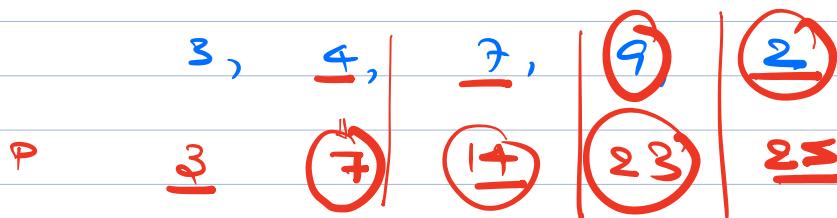
① for ( $i = 0$ ;  $i < N$ ;  $i++$ ) {

$$PF[i][0] = M[i][0]; \text{ // Edge Case}$$

for ( $j = 1$ ;  $j < M$ ;  $j++$ ) {

$$PF[i][j] = PF[i][j-1] + M[i][j];$$

}



Col wise

② for ( $j = 0$ ;  $j < M$ ;  $j++$ ) {

for ( $i = 1$ ;  $i < N$ ;  $i++$ ) {

$$PF[i][j] = PF[i-1][j] + PF[i][j];$$

}

$$T.C. = O(\underline{N} \times \underline{M})$$

$$\text{Overall } T.C. = O\left(\underline{0} + \underline{N} \times \underline{M}\right)$$

PS in array

$$S.C. = O(\underline{N} \times \underline{M})$$

1  
2  
APPS

D-2 Given a matrix of size  $N \times M$

Calculate the sum of all submatrix

Eg

$$\begin{bmatrix} 4 & 9 & 6 \\ 5 & -1 & 2 \end{bmatrix}$$

$$\underline{\underline{[4]}} \Rightarrow 4$$

$$\underline{\underline{[9]}} \Rightarrow 9$$

$$\underline{\underline{[6]}} \Rightarrow 6$$

$$\underline{\underline{[5]}} \Rightarrow 5$$

$$\underline{\underline{[-1]}} \Rightarrow -1$$

$$\underline{\underline{[2]}} \Rightarrow 2$$

$$\underline{\underline{[4, 9]}} \Rightarrow 13$$

$$\underline{\underline{[9, 6]}} \Rightarrow 15$$

$$\underline{\underline{[5, -1]}} \Rightarrow 4$$

$$\underline{\underline{[-1, 2]}} \Rightarrow 1$$

$$\underline{\underline{[4, 9, 6]}} \Rightarrow 19$$

$$\underline{\underline{[5, -1, 2]}} \Rightarrow 6$$

$$\underline{\underline{[4, 5]}} \Rightarrow 9$$

$$\underline{\underline{[9, -1]}} \Rightarrow 8$$

$$\underline{\underline{[5, 2]}} \Rightarrow 8$$

$$\underline{\underline{[4, 9]}} \Rightarrow 17$$

$$\underline{\underline{[9, -1]}} \Rightarrow 16$$

$$\underline{\underline{[5, 2]}} \Rightarrow 16$$

$$\underline{\underline{[4, 9, 6]}} \Rightarrow 23$$

$$\text{Sum} = \underline{\underline{166}} \text{ Ans}$$

$$4 \times 6 + 9 \times 8 + 6 \times 6 + 5 \times 6 + (-1) \times 8$$

$$+ 2 \times 6$$

$$= 24 + 72 + 36 + 30 - 8 + 12$$

$$= \underline{\underline{166}}$$

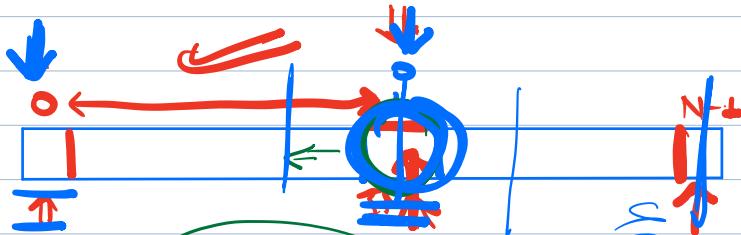
Sum of all submatrices

Sum

$$= \left( \begin{array}{l} \text{No. of submatrices} \\ \text{A}[i][j] \text{ is a part} \\ \text{of} \end{array} \right) \times \underline{\underline{A[i][j]}}$$



Subarray

$$(N-i) - i + 1 = N-i$$

$$(i+1) \times (N-i)$$

|   | 0 | 1 | 2 | 3 | 4 | 5 | 6 |
|---|---|---|---|---|---|---|---|
| 0 | ✓ | ✓ | ✓ | ✓ |   |   |   |
| 1 | ✓ | ✓ | ✓ | ✓ |   |   |   |
| 2 | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ | ✓ |
| 3 |   |   |   | ✓ | ✓ | ✓ | ✓ |

$$6 \times 7$$

$$\underline{\underline{(i, j)}}$$

|   |  |  |  |   |   |   |   |
|---|--|--|--|---|---|---|---|
| + |  |  |  | - | - | - | - |
| S |  |  |  | - | - | - | - |

start  $\Rightarrow$  green

End  $\Rightarrow$  red

|   |                                 |  |
|---|---------------------------------|--|
| $\text{TL}$<br>$(x_1, c_1)$                       | $i - 0 + 1 = i$<br>$\downarrow$ | $\text{BR}$<br>$(x_2, c_2)$<br>$[i, N-1]$                                |
| $x_1 \leq i \Rightarrow [0, i] \Rightarrow (i+1)$ | $c_1 \leq j \Rightarrow (j+1)$  | $x_2 > i \Rightarrow (N-i)$<br>$c_2 > j \Rightarrow M-j$<br>$= [j, M-1]$ |

# submatrix

$(i, j)$  is a part of

$$= \frac{(i+1)}{\downarrow} \frac{(j+1)}{\downarrow} \times \frac{(N-i)}{\downarrow} \frac{(M-j)}{\downarrow}$$

## Code

Sum = 0;

for ( $i = 0$ ;  $i < N$ ;  $i++$ ) {

    for ( $j = 0$ ;  $j < M$ ;  $j++$ ) {

$$\text{Contri} = M \underline{i}[i][j] \times \left( \frac{(i+1) \{ j+1 \} (N-i)}{(M-j)} \right)$$

$\text{sum} = \text{sum} + \text{contri};$

{

T.C. =  $O(\underline{N} \times \underline{M})$

S.C. =  $O(1)$

$\Theta^3$  Given a matrix of size NxM.

google find the max submatrix sum.

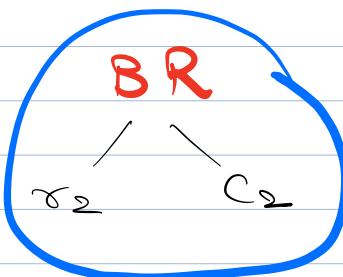
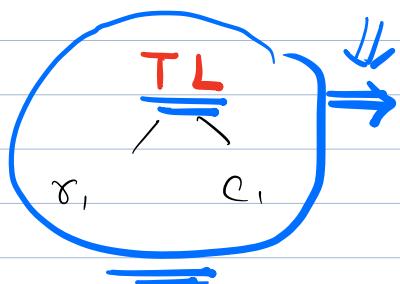
Duo's  
Facebook

|    |   |    |
|----|---|----|
| -2 | 3 | 1  |
| 3  | 6 | -3 |

$\Rightarrow \underline{\underline{10}}$

Soln ) Brute force

→ Create all submatrices  
for each submatrix, calc sum &  
take max



r for ( $r_1 = 0$ ;  $r_1 < N$ ;  $r_1 ++$ ) d // N

TL { } - for ( $c_1 = 0$ ;  $c_1 < M$ ;  $c_1 ++$ )  $\{ \text{ // } M$

BR { for ( $r_2 = r_1$ ;  $r_2 < N$ ;  $r_2 ++$ )  $\{ \text{ // } N$   
for ( $c_2 = c_1$ ;  $c_2 < M$ ;  $c_2 ++$ )  $\{ \text{ // } M$

//  $r_1, c_1, r_2, c_2$

↓  
// Prefix Sum Matrix

{      }

↓

$N^4$  X

T.C. =  $O(N^2 M^2)$

Q

How to optimise  $\Rightarrow$  2D Kadane's

1) find the max submatrix sum

starting at row 0 & ending at

row  $N-1$

2)

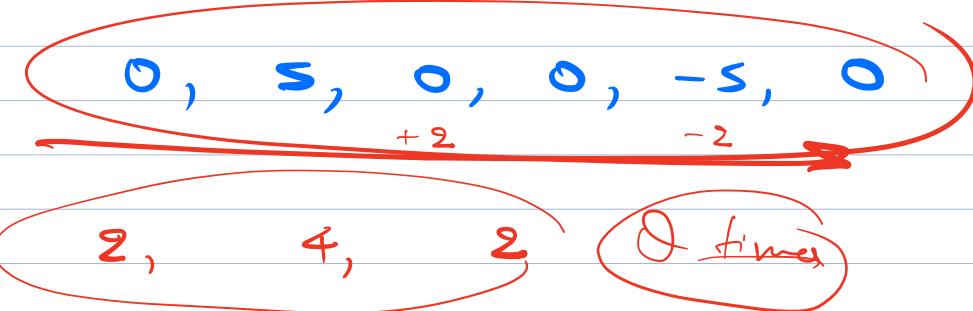
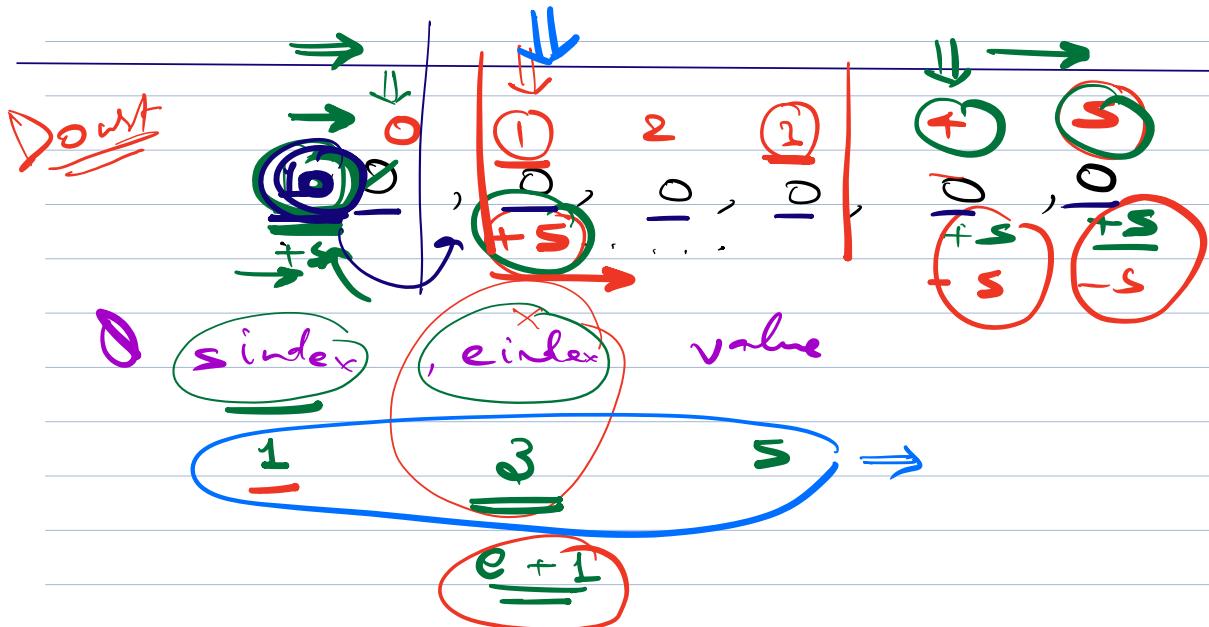
" " "

starting at row 0 & ending at  
any row.

3)

" " "

starting at any row & ending at any row.



LLD  $\Rightarrow$

Normalisation

$$\begin{matrix} 1+1 \\ (+-) \end{matrix}$$

1) Resummet gathering

✓ 7 ♂ 5