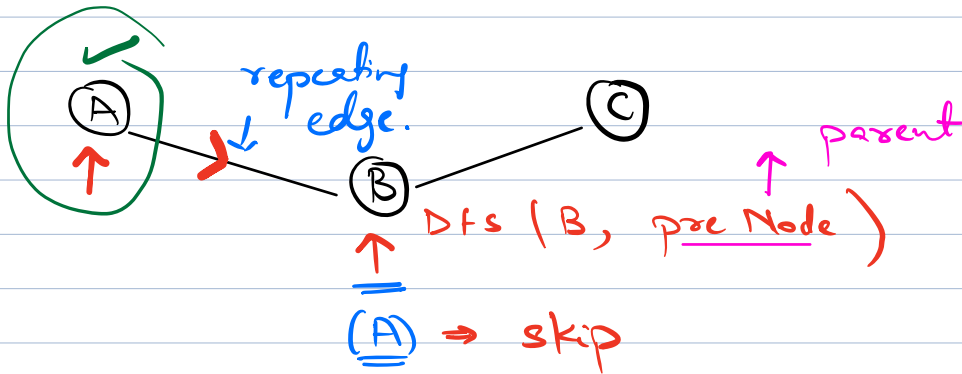
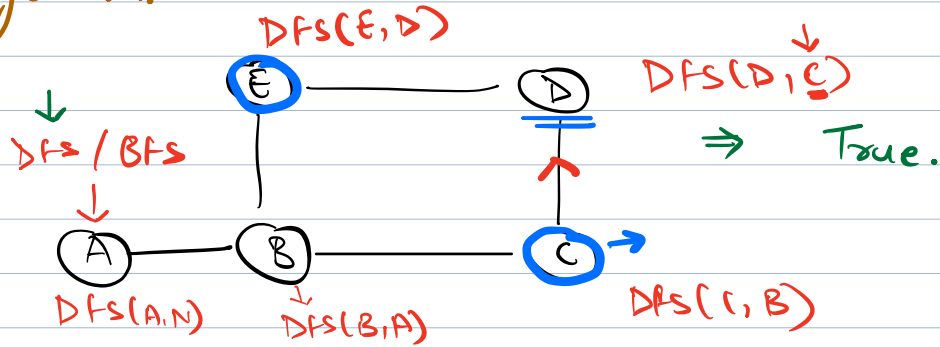




Given an undirected graph.

Return True if the graph contains a cycle...

eg



Code

```

bool isCyclic (u, parentNode) {
    DFS
    Source
    NULL
    visited[u] = true;
    for (all v connected to u) {
        if (visited[v] == True &&
            v != parentNode) {
            return True;
        }
    }
}

```

```

    }
    else if (visited[v] == false) {
        if (isCyclic(v, u)) {
            return True;
        }
    }
    return false;

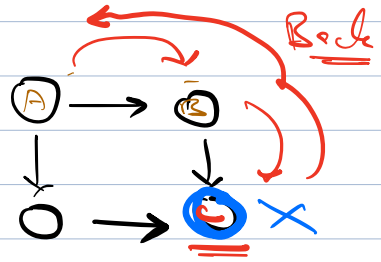
```

}

$$T.C. = O(V+E)$$

H.W.

Directed graph.



Q

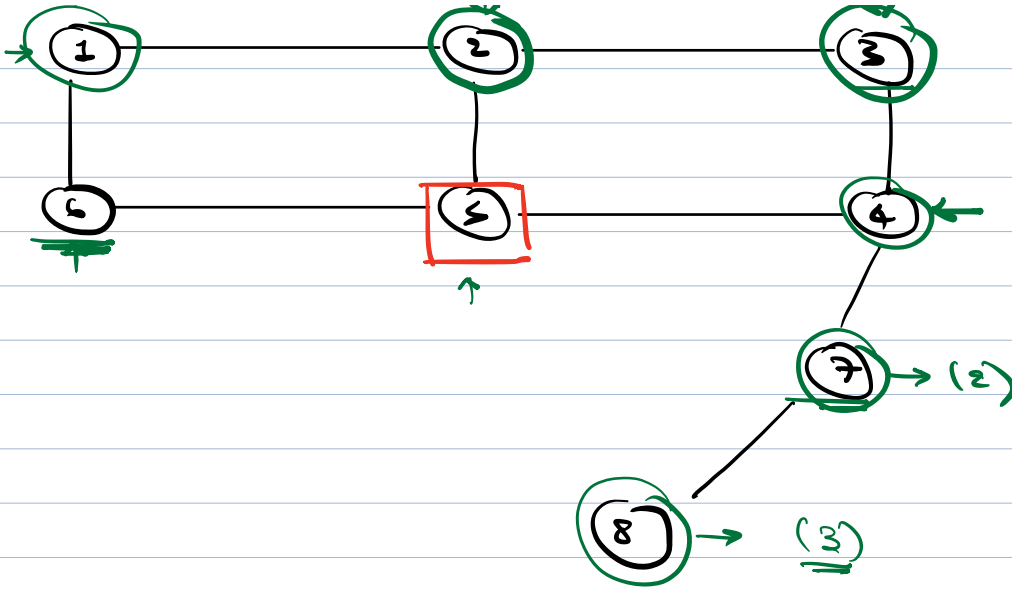
Given an undirected / unweighted graph.

Find min distance from a source to all nodes.

↓

↓

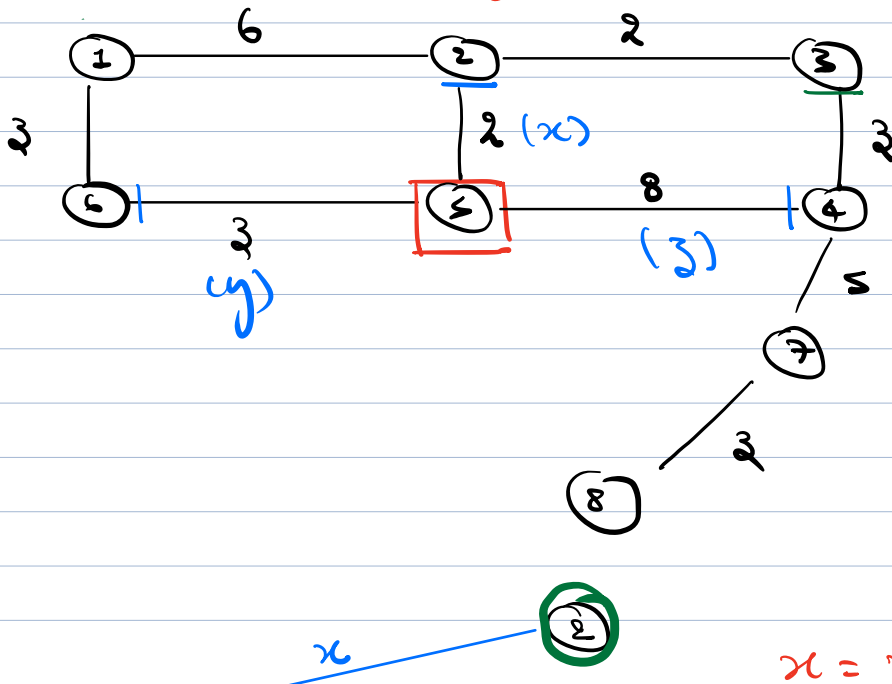
↓



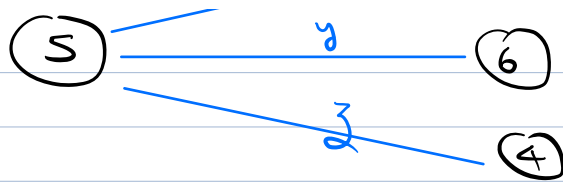
O/P

	2	1	2	1	0	1	2	3
0	1	2	3	4	5	6	7	8

Undirected / Weighted. (No -ve weights)  
(Dijkstra's Algo)



$$x = \min(1, 1, 2)$$



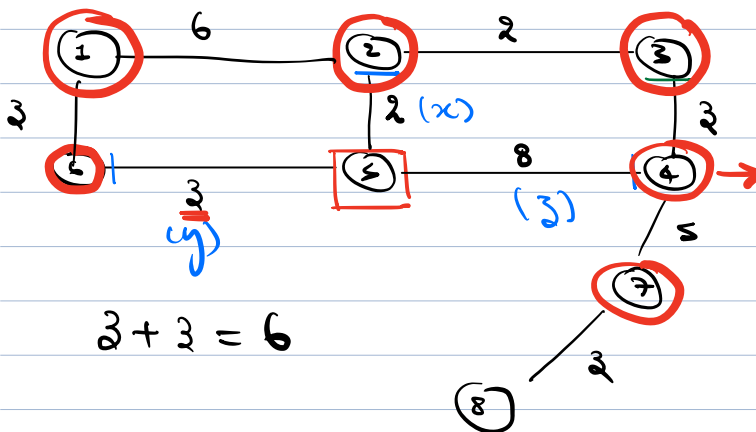
(1, 2, 3)

# Ways to reach 2 is  $\rightarrow$  Direct (x)

2) Through 6 (y+)

3) Through 4 (z+)

$$x = \min(x, y + \underline{w}, z + \underline{w})$$



1 : ~~2~~ ~~8~~ 6  
 2 : ~~6~~ 2  
 3 : ~~3~~ 4  
 4 : ~~6~~ ~~8~~ 7  
 5 : 0  
 6 : ~~6~~ 3  
 7 : ~~3~~ 12  
 8 : ~~2~~ 15

Min Heap  $\Rightarrow$  Pair  $\langle$  Length, Node  $\rangle$

~~$\langle 2, 2 \rangle$~~   
 ~~$\langle 3, 6 \rangle$~~   
 ~~$\langle 8, 4 \rangle$~~   
 ~~$\langle 4, 3 \rangle$~~   
 ~~$\langle 8, 1 \rangle$~~   
 ~~$\langle 6, 1 \rangle$~~

$\Rightarrow$  Edges

$O(E) \rightarrow$  space

~~$\langle 7, 4 \rangle$~~   
 ~~$\langle 12, 7 \rangle$~~   
 $\langle 15, 8 \rangle$

T.C. =  $E \log(E)$

## Code

Min Heap  $\rightarrow h$

Min Dist  $[N+1] = \infty$

Min Dist [source] = 0;

void dijkstra ( source ) {

for ( all  $v$  connected to source ) {

$h.add( Pair \langle W_{s-v}, v \rangle );$

Min Distance  $[v] = W_{s-v};$

}

while ( !  $h.is\ Empty()$  ) {

$Pair \langle l, u \rangle = h.getMin();$

if (  $l \leq MinDistance[u]$  ) {

for ( all  $x$  connected to  $u$  ) {

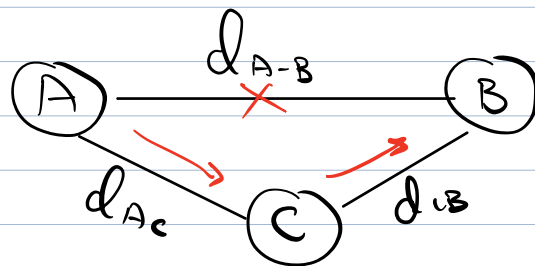
$D_{s-x} = \underset{\downarrow}{l} + W_{u-x};$

$D_{s-u}$

if (  $D_{s-x} < MinDist[x]$  ) {

$\text{Min Dist}[x] = D_{s-x}$   
 h. insert (Pair  $\langle D_{s-x}, x \rangle$ );

$$\begin{aligned}
 \text{T.C.} &= O(V + \cancel{E} + E \log E) \\
 &= O(V + E \log E)
 \end{aligned}$$



$\Rightarrow$  Relaxing an edge

$$d_{Ac} + d_{CB} < d_{A-B}$$

★ Floyd Warshall Algo.

(All possible shortest path algo)

Q find min distance to reach any node from every node.

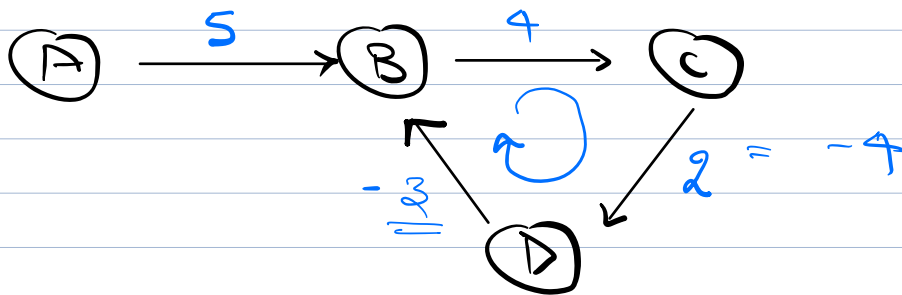
O/P  $\rightarrow$  Matrix of size  $V \times V$

Undirected  $\rightarrow$  -ve weight edge. X

A  $\xrightarrow{2}$  B  $\xrightarrow{-4}$  C

$$2 - 4 - 4 = \underline{\underline{-6}} \quad -4 - 4 = \underline{\underline{-8}}$$

Directed (No -ve weight cycle)



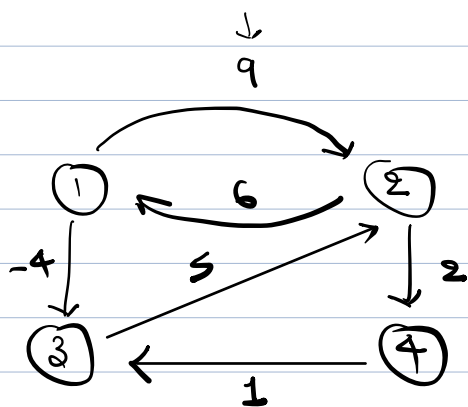
O/P  $\Rightarrow$   $(V+1) \times (V+1)$  matrix

$D[i][j] \rightarrow$  Min distance from  $i$  to  $j$

Iterate over all nodes

↳ Treat this node as int. node

↳ Try to relax every edge.

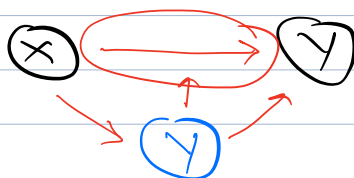


Adjacency matrix

	1	2	3	4
1	0	9	-4	$\infty$
2	6	0	$\infty$	2
3	$\infty$	5	0	$\infty$
4	$\infty$	$\infty$	1	0

$\Delta_0$

1) Node 1 as inter node.



2, 4  $\rightarrow$  2, 1 + 1, 4  
(2) (6)  $\infty$

	1	2	3	4
1	0	9	-4	$\infty$
2	6	0	$\infty$	2
3	$\infty$	5	0	$\infty$
4	$\infty$	$\infty$	1	0

$\Delta_0$

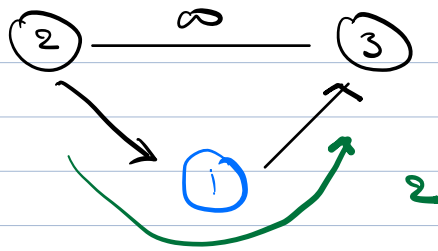
$\rightarrow$

	1	2	3	4
1	0	9	-4	$\infty$
2	6	0	2	2
3	$\infty$	5	0	$\infty$
4	$\infty$	$\infty$	1	0

$\Delta_1$

$$M[2][3], \quad M[2][1] + M[1][3] = 2$$





$$M[3][2], \quad (M[3][1] + M[1][2])$$

$\downarrow$   
 $\infty$

$$M[3][4], \quad (M[3][1] + M[1][4])$$

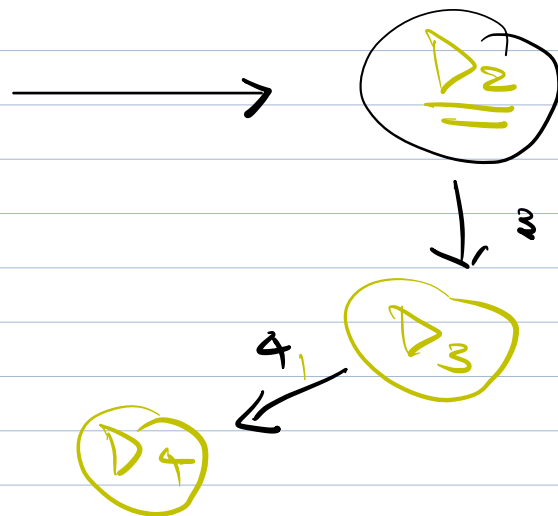
$\downarrow$   
 $\infty$

$$M[4][2], \quad (M[4][1] + M[1][2])$$

2) Node 2 as Int Node

	1	2	3	4
1	0	9	-4	$\infty$
2	6	0	2	2
3	$\infty$	5	0	$\infty$
4	$\infty$	$\infty$	1	0

$\Delta_1$



Code

$$T.C. = O(V^3)$$

Graph  $\rightarrow$  i) Topology set  
ii) Bipartite graph (Graph coloring)

$S \rightarrow$  Prims, Kruskal (MST)