

# Report on question 2

## Problem:

Professor Stewart is consulting for the president of a corporation that is planning a company party. The company has a hierarchical structure; that is, the supervisor relation forms a tree rooted at the president. The personnel office has ranked each employee with a conviviality rating, which is a real number. In order to make the party fun for all attendees, the president does not want both an employee and his or her immediate supervisor to attend.

Professor Stewart is given the tree that describes the structure of the corporation, using the left-child, right-sibling representation. Each node of the tree holds, in addition to the pointers, the name of an employee and the employee's conviviality ranking. Describe an algorithm to make up a guest list that maximizes the sum of the conviviality ratings of the guests. Analyze the running time of your algorithm.

## Remarks and assumptions of the solution:

- Remark : Data Structure to store the company employee details and organisational hierarchy is a Tree with left-child, right-sibling representation, where root of the tree is the president of the corporation and parents-child relationship represents the immediate supervisor-employee relationship. Each node will have the name of the employee and the conviviality rating (metric to maximize) of the employee.
- Assumption: an Employee has only one immediate supervisor(a child has only one parent according to our data structure)
- Assumption inviting President of the company is not compulsory

## Approach on the solution:

Our solution uses Dynamic programming to solve this problem,  
Corresponding recursion would be (to find the maximum conviviality ratings) :

```
maximumConviviality(Node):  
    If Node.Children = null  
        Return Node.rating  
    Else  
        maxRatingChildren = 0  
        For C in Node.Children  
            maxRatingChildren += maximumConviviality(C)  
        maxRatingGrandChildren = Node.rating  
        For G in Node.GrandChildren  
            maxRatingGrandChildren += maximumConviviality(G)  
        If maxRatingGrandChildren >= maxRatingChildren  
            Return maxRatingGrandChildren  
        Else  
            Return maxRatingChildren
```

By applying Dynamic Programming to above recursion we calculate the maximum possible Conviviality up to that node and keep track of who are to be invited to maximize Conviviality. When applying dynamic programming we traverse the tree from bottom to top and if “maxRatingGrandChildren >= maxRatingChildren” is True we assign “fitt to be invited” of that node as True. Then to create the invitee list we traverse the tree top to bottom. If “fitt to be invited” is True for a node we add that node to the invitee list and sets “fitt to be invited” of it's children False. If “fitt to be invited” is False for a node we traverse to the next node.

We apply dynamic programming by storing the maximumConviviality up to that node in the Node.rating field (we doesn't use a separate data structure to store maximumConviviality since it's not necessary). “fitt to be invited” parameter tells us that if it's true, to achieve maximumConviviality marked on that node we must invite that node.

## Algorithm(pseudo code) / Time complexity analysis :

Macro function(**getRatingSumOfChildren**) :

Child = no of child node each node has

C = constant

Time	Pseudo Code
<div>C C Child C C C</div>	<pre>def <b>getRatingSumOfChildren</b>(parent):     currentNode = parent.leftChild     ratingSum = 0.0     while not currentNode is None:         ratingSum += currentNode.rating         currentNode = currentNode.rightSibling     return ratingSum</pre>

Time complexity(**getRatingSumOfChildren**) macro =  $3c + (\text{Child} \times 2c) = O(\text{Child})$

Macro function(**getRatingSumOfGrandChildren**) :

Child = no of children each node has

Gchild = no of grandchildren each node has

C = constant

Time	Pseudo Code
<div>C C Child C Gchild C C C C</div>	<pre>def <b>getRatingSumOfGrandChildren</b>(parent):     ratingSum = 0.0     child = parent.leftChild     while not child is None:         grandChild = child.leftChild         while not grandChild is None:             ratingSum += grandChild.rating             grandChild = grandChild.rightSibling         child = child.rightSibling     return ratingSum</pre>

Time complexity(**getRatingSumOfGrandChildren**) macro =  $3c + \text{Child} \times (2c + (\text{Gchild} \times 2c)) = O(\text{Child} \times \text{Gchild})$

Preprocessing :

C = constant

N = no of nodes in the tree

H = no of children a node can have

G = no of grandchildren a node can have

Time	Pseudo Code
C	stack = Stack() #to accesses the tree from bottom to top
C	queue = Queue()
C	queue.enqueue(tree.root)
C	stack.push(tree.root)
C	currentNode = tree.root
N	<b>while not</b> queue.isEmpty():
C	currentNode = queue.dequeue()
C	<b>if not</b> currentNode <b>is None</b> :
C	currentNode = currentNode.leftChild
C	<b>if not</b> currentNode <b>is None</b> :
C	queue.enqueue(currentNode)
C	stack.push(currentNode)
H	<b>while not</b> currentNode.rightSibling <b>is None</b> :
C	currentNode = currentNode.rightSibling
C	queue.enqueue(currentNode)
C	stack.push(currentNode)
N	<b>while(not</b> stack.isEmpty()):#calculate maximum possible rating for each sub tree initiating from Node
C	currentNode = stack.pop()
O(H)	childrenRatingSum = getRatingSumOfChildren(currentNode)
O(HxG)	grandChildrenRatingSum = getRatingSumOfGrandChildren(currentNode)
C	<b>if</b> ((currentNode.rating + grandChildrenRatingSum)>=childrenRatingSum):
C	currentNode.invite = <b>True</b>
C	currentNode.rating += grandChildrenRatingSum
C	<b>else</b> :
C	currentNode.invite = <b>False</b>
C	currentNode.rating = childrenRatingSum

Time complexity(**Preprocessing**) =  $5C + N \times (6C + H \times (3C)) + N \times (7C + O(H) + O(H \times G))$   
 $= O(N \times H + N \times H \times G)$   
 $= O(N \times H \times G)$

### Creating Invitee List :

C = constant

N = no of nodes in the tree

H = no of children a node can have

G = no of grandchildren a node can have

Time	Pseudo Code
C	inviteList = []
C	queue = Queue()
C	queue.enqueue(tree.root)
C	currentNode = tree.root
C	<b>if</b> (currentNode.invite):
C	inviteList.append(currentNode.name)
C	childNodes = currentNode.leftChild
H	<b>while not</b> childNode <b>is</b> None:
C	childNodes.invite = <b>False</b>
C	childNodes = childNodes.rightSibling
N	<b>while not</b> queue.isEmpty():
C	currentNode = queue.dequeue()
C	<b>if not</b> currentNode <b>is</b> None:
C	currentNode = currentNode.leftChild
C	<b>if not</b> currentNode <b>is</b> None:
C	queue.enqueue(currentNode)
C	<b>if</b> (currentNode.invite):
C	inviteList.append(currentNode.name)
C	childNodes = currentNode.leftChild
H	<b>while not</b> childNode <b>is</b> None:
C	childNodes.invite = <b>False</b>
C	childNodes = childNodes.rightSibling
H	<b>while not</b> currentNode.rightSibling <b>is</b> None:
C	currentNode = currentNode.rightSibling
C	queue.enqueue(currentNode)
C	<b>if</b> (currentNode.invite):
C	inviteList.append(currentNode.name)
C	childNodes = currentNode.leftChild
H	<b>while not</b> childNode <b>is</b> None:
C	childNodes.invite = <b>False</b>
C	childNodes = childNodes.rightSibling
C	<b>return</b> inviteList

Time complexity(**Creating Invitee List**) =  $8C + Hx2C + Nx(8C + Hx2C + Hx(5C + Hx2C))$

$$\begin{aligned}\text{Time complexity(Creating Invitee List)} &= O(H + Nx(H + (H \times H))) \\ &= O(NxH \times H)\end{aligned}$$

Time complexity of the algorithm = Time complexity(Creating Invitee List) + Time complexity(Preprocessing)

$$\text{Time complexity of the algorithm} = O(NxH \times H) + O(N \times H \times G)$$

Assume  $G = H \times H$

Then :

$$\begin{aligned}\text{Time complexity of the algorithm} &= O(NxH \times H) + O(N \times H \times H \times H) \\ &= O(N \times H \times H \times H)\end{aligned}$$

Where : N - Total no of nodes in the tree

H - Maximum no of children a node can have